CS & IT ENGINEERING

Theory of Computation

Regular Languages



Lecture No.- 09

Recap of Previous Lecture







Regular Language Vs Regular Expression

Topics to be Covered









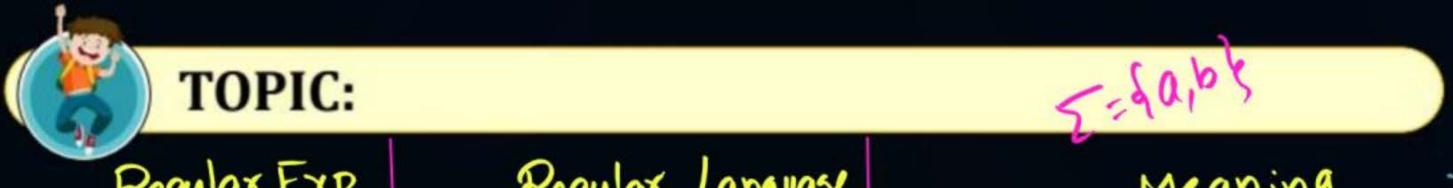


Topic

Regular Language Vs Regular Expression

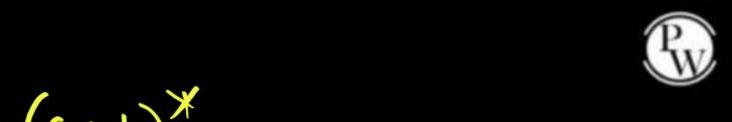


Universal Set





F	Regular Exp.	Regular Language	Meaning	iset
(5)	(a+b)*	€a,b}*	Set of (all) strings over = Set of (all) non zero lingth	I= 40,64 (2) LOUGU
(b)	(a+b) ^t	fa,6}+=fw wE(a+b)*,1w1>0	= Set of (all) non zeno brylk	strings
(3)	a.(a+6)*	= a E* = faw we E*}	= Set of (all) strings where	every string begins will
(8)	b (a+b)*	Starting wilk b		
(19)	(a+6) a	ends wilk a		
20	(atb)*b	11 11 6		
_				



a followed by any bs

Set of (a11) strings starting with a' over E=da, b?

ab __ a\tau



O $ab^* = \{a, ab, abb, ab\}, \dots$

(2) a(a+b)* = fa, aa, ab, aaa, aab,

$$(1)$$
 (2)





14	Regular Exp.	Regular Language	Meaning
(2)	(otb)+	= (a+b) b = Set of	all strings ending will b
	(b*a)+		t of all Strings ending wilk a
23	(a6*)+	= a E*	
ঞ্চ	(bå*)*	= b \(\Sigma^* \)	
क्ड	(ab)* a	$=(a+b)^*=\Sigma^*$	
26)	(ba * b*	$= (a+b)^{*} = \overline{\Sigma}^{*}$	Universal Set over 2
27)	å (bå*)*	= (a+b)* = I*	
Slide	5		

R

 $(a^*b)^+ = (a+b)^*b$

$$\alpha \Sigma^* = (\alpha b^*)^{\dagger}$$

$$b \Sigma^* = (ba^*)^+$$

$$\Sigma^* \alpha = (b^* \alpha)^{\dagger}$$





$$(a^{*}b)^{+} = (a+b)^{*}b$$
 $(a^{*}b)^{-}=b^{*}$
 $(a^{*}b)^{-}=b^{*}$
 $(a^{*}b)^{-}=b^{*}$
 $(a^{*}b)^{-}=b^{*}$
 $(a^{*}b)^{-}=b^{*}$
 $(a^{*}b)^{-}=b^{*}$





$$b^*(ab^*) = (a+b)^*$$
 $b^*(ab^*) = 0$
 $b^*(ab^*) = 0$
 $b^*(ab^*) = 0$



TOPIC: Worte Regular Exp



=
$$\int aa\epsilon, aaa, aaaa, ...$$
}
= $\int a^2, a^3, a^4, ...$ }
= $\int a^n |n=2|$

$$= \frac{x}{a} = \frac{$$





$$= (a+b)^* aa (a+b)^*$$

Set of all strings containing aa as substring.

 $(a+b)^{*}$ $\alpha a (a+b)^{*}$ $(a*b^{*})^{*}$ $\alpha a (b*a^{*})^{*}$



(a+b)=(xb) =(xb) =(xb) =(b*x) =(b*x) =





$$= (a+b)^2 = (aa+ab+ba+bb) = \sum_{i=1}^{2}$$

(32)
$$\{\omega \mid \omega \in \{a,b\}^*, |\omega| \geq 2\}$$

= $(a+b)^2 \cdot (a+b)^* = \sum_{z=1}^{2} \sum_{z=1}^{2} = \sum_{z=1}^{2} \sum_{z=1$

$$(E+a+b)^2$$







$$\begin{cases}
\omega \mid \omega \in \{a,b\}^*, \quad |\omega| = \text{divisible by } \\
= \{\epsilon, aa, ab, ba, bb, aaaa, aaab, \dots \}
\end{cases} = (\sum \sum)^*$$

$$= (aa+ab+ba+bb)^* = (a+b)^2 = (\Sigma^2)^*$$

$$(\Sigma^2)^*$$
, $\Sigma = \Sigma.(\Sigma^2)^*$

$$= ((a+b)^{2})^{2}(a+b) = (a+b).((a+b)^{2})^{2}$$

$$=(atb).((atb))$$



$$\left(\left(a+b\right)^{2}\right)^{*}=\left(aa+ab+ba+bb\right)^{*}$$

$$\left(\left(\mathcal{E} + \mathbf{a} + \mathbf{b} \right)^{2} \right)^{*} = \left(\mathcal{E} + \mathbf{a} + \mathbf{b} + \mathbf{a} \mathbf{a} + \mathbf{a} \mathbf{b} + \mathbf{b} \mathbf{a} + \mathbf{b} \mathbf{b} \right)^{*}$$
$$= \left(\mathbf{a} + \mathbf{b} \right)^{*}$$
$$= \left(\mathbf{a} + \mathbf{b} \right)^{*}$$



$$Ma(w) = 2$$

 $mo.oga's in w is 2$
 $Ha(w) = 2$





$$\{\omega \mid \omega \in \{a,b\}^*, n_a(\omega) = 2\}$$

$$= b^*ab^*ab^*$$

$$37$$
 $\{\omega | \omega \in \{a,b\}^*, n_a(\omega) \geq 2\}$ $baba(a+b)^* = (a+b)^* \alpha (a+b)^* \alpha (a+b)^* = \sum_{i=1}^{n} \alpha \sum_{$



$$N_{a}(\omega) = 0$$
 \Rightarrow b
 $N_{a}(\omega) = 1$ \Rightarrow b
 $N_{a}(\omega) = 1$ \Rightarrow b
 $N_{a}(\omega) = 2$ \Rightarrow b
 $N_{a}(\omega) = 2$ \Rightarrow b
 $A_{a}(\omega) = 2$
 A_{a



Home work





2 mins Summary



Topic

Operators

Topic

Properties

Topic

Simplification



THANK - YOU