# CS & IT ENGINEERING

Theory of Computation

**Regular Languages** 



Lecture No.- 11

### **Recap of Previous Lecture**







Topic

Regular Language Vs Regular Expression

## **Topics to be Covered**











Topic

Regular Language Vs Regular Expression







$$\eta_{\alpha}(\omega) = 3k+2$$
 $k > 0$ 
 $= 2, 5, 8, 11, 14, ...$ 





(47) {w| wefa, b}\*, w starts wik aa or bb}

$$= aa \Sigma^* + bb \Sigma^* = (aa + bb) (a + b)^*$$

(48) fw/wefa,bj\*, w ends wilk aa or bbz

(49) for loss da, by\*, w contains aa or bb as substring?

$$= \sum^{*} (aa+bb) \sum^{*} = \sum^{*} aa \sum^{*} + \sum^{*} bb \sum^{*}$$





&w|wefa,b}x, w starts wilk aa and w starts with tholy There is no string

(5) éw | we la, by\*, wends wilk aa, wends wilk bby

fw | we ta, byt, w contains aa and w contains bby

= 5 aa 5 bb 2 + 2 bb 2 aa 5

= Ex(aas\*bb+bbs\*aa) E\*

aabb bbaav ··· Aa...bb...

... bb...aa...





= 
$$\Sigma^*(aabb + bbaa) \Sigma^*$$
  
=  $\Sigma^*(aabb \Sigma^* + \Sigma^*bbaa \Sigma^*)$ 







d w | w ∈ fa, b}\*, 2<sup>nd</sup> symbol of w is a'?

$$= (a+b) a (a+b)$$

$$= \sum \alpha \sum^{*}$$





 $z^*a(a+b) = z^*aa + z^*ab$ 

5 ta 5 = (a+b)\* a (a+b) and symbol from end is a

any sequence and last any symbol





$$(57) \{(a^mb)|m,n\geq 0\} \Rightarrow a^*b^*$$

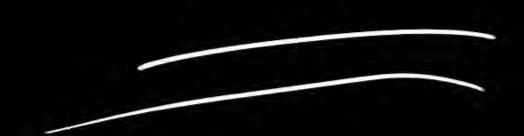
$$\Rightarrow a^+b^+$$

$$\{b^ma^n \mid m,n\geq 2\}$$

$$=bbaa=bbaa$$

69 
$$\{(m,n) \in \mathbb{Z}\} \Rightarrow bbb aaa = bbaa = bbaa$$

60  $\{(m,n) \times | m| = 2\} = a(\epsilon+b) \cdot cc+$ 
 $\{(m,n) \times | m| = 2\} = a(\epsilon+b) \cdot cc+$ 







fambam,n>of

= de, a, b, ab, aa, bb, ...}

= a\* b\*





```
of as n>0 /= fE, ab, aabb, aaabb, ....}
Note:
            Equal no. of as 2 6's
        Can you write Regular Expression
            There is no regular expressions
       >It is not regular language.
```





Language
Regular Language
Not Regular Language

No Regular Exp.





4 symbol from end is a



#### 2 mins Summary



Topic

Regular Languages

Topic

**Regular Expressions** 



# THANK - YOU