

Computer Science

Theory of Computation

Regular Languages and Non-regular Languages

Lecture No.- 9

A man with a beard and mustache, wearing a black polo shirt, stands with his arms crossed in front of a blurred bookshelf. He is wearing a black watch on his left wrist.

Malleham Devasane Sir

Recap of Previous Lecture



Topic

Closure Properties for Finite Languages

Topic

Closure Properties for Infinite Languages

Topic

Closure Properties for Regular Languages

Topics to be Covered



Topic

Closure Properties for Regular Languages



Closure Properties for Regular Languages



- ① $L_1 \cup L_2$
- ② $L_1 \cap L_2$
- ③ \overline{L}
- ④ $L_1 - L_2$
- ⑤ $L_1 \cdot L_2$
- ⑥ L^{Rev}
- ⑦ L^*
- ⑧ L^+

⑨ $\text{Subset}(L)$

⑩ $\text{prefix}(L)$

⑪ $\text{suffix}(L)$

⑫ $\text{Substring}(L)$

⑬ $f(L) = \text{Substitution}$

⑭ $h(L) = \text{Homomorphism}$

⑮ $h^{-1}(L)$

⑯ L_1 / L_2
Quotient

⑰ $L_1 \oplus L_2$
Symmetric Difference

⑱ $\text{Half}(L) = \frac{1}{2}(L)$

⑲ $\text{Second Half}(L)$

⑳ $\text{one third}(L)$

㉑ $\text{Middle } \frac{1}{3}(L)$

㉒ $\text{Last } \frac{1}{3}(L)$

㉓ Finite Union

㉔ " \cap

㉕ " Difference

㉖ " Concatenation

㉗ " Subset

㉘ " Substitution

㉙ Inf \cup

㉚ Inf \cap

㉛ Inf Concatenation

㉜ Inf Subset

㉝ Inf Substitution

(16) Quotient $L_1 / L_2 = \{ u \mid uv \in L_1, v \in L_2 \}$

$$L_1 = \{ \epsilon, abb \}$$

$$L_2 = \{ a, bb \}$$

$$L_1 / L_2 = \left\{ \underbrace{\epsilon/a}_\times, \underbrace{\epsilon/bb}_\times, \underbrace{abb/a}_\times, \underbrace{abb/bb}_a \right\}$$

$$= \{ a \}$$

$$\boxed{abb} / \cancel{bb} = a$$

$$\boxed{uv / v = u}$$

$$abb / \boxed{a} = \emptyset$$

$$abb / \epsilon = abb$$

$$L/\varepsilon = L$$

$$\varepsilon/\varepsilon = \varepsilon$$

$$a/\varepsilon = a$$

$$\varepsilon/a = \phi$$

$$\textcircled{1} L_1 = a^*$$

$$L_2 = \{a\}$$

$$\begin{aligned} \text{i) } L_1/L_2 &= a^*/a = \{\varepsilon, a, aa, aaa, \dots\}/a \\ &= \{\overset{\times}{\varepsilon/a}, \overset{=\varepsilon}{a/a}, \overset{=a}{aa/a}, \overset{=aa}{aaa/a}, \dots\} \\ &= \{\varepsilon, a, aa, \dots\} = a^* \end{aligned}$$

$$\begin{aligned} \text{ii) } L_2/L_1 &= a/a^* \\ &= \{a/\varepsilon, a/a, a/\underline{aa}, a/\underline{aaa}, \dots\} \\ &= \{a, \varepsilon\} \end{aligned}$$

$\begin{matrix} =a & =\varepsilon & \times & \times \end{matrix}$

$$2) \quad L_1 = ab^*$$

$$L_2 = a^*b$$

$$i) L_1/L_2 = ab^*/a^*b = \{a, ab, abb, \dots\} / \{b, ab, aab, \dots\}$$

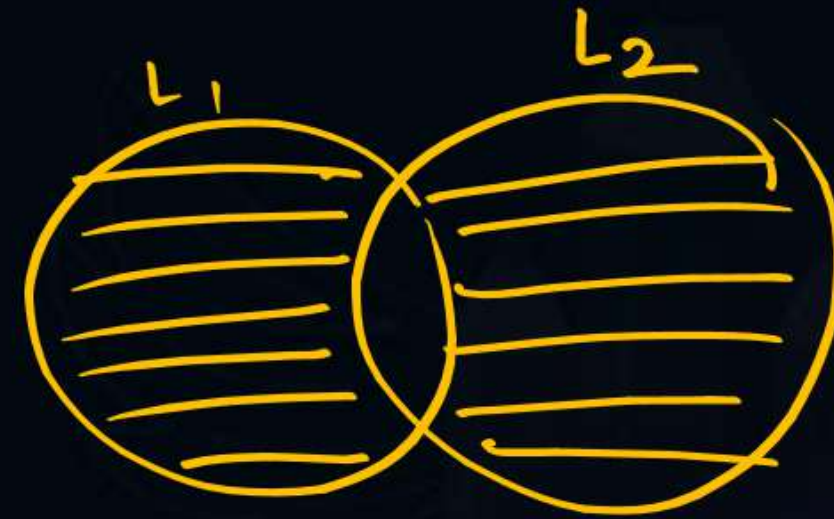
$$= \{ \overset{\times}{a}/b, \overset{\times}{a}/ab, \overset{\times}{a}/aab, \dots, \\ \overset{=a}{ab}/b, \overset{=\epsilon}{ab}/ab, \overset{\times}{ab}/aab, \dots, \\ \overset{=ab}{abb}/b, \overset{\times}{abb}/ab, \overset{\times}{abb}/aab, \dots \} = \underline{\epsilon + ab^*}$$

H.W. ii) $L_2/L_1 =$

⑪ Symmetric Difference

$$L_1 \oplus L_2 = L_1 \Delta L_2 = (L_1 \cup L_2) - (L_1 \cap L_2)$$

$$= (L_1 - L_2) \cup (L_2 - L_1)$$



$$\textcircled{1} \quad L_1 = a^*$$

$$L_2 = b^*$$

$$L_1 \Delta L_2 = (L_1 \cup L_2) - (L_1 \cap L_2)$$

$$= (a^* \cup b^*) - (\{\epsilon\})$$

$$= a^+ + b^+$$

$$\textcircled{2} \quad L_1 = a^+$$

$$L_2 = b^+$$

$$L_1 \Delta L_2 = (L_1 \cup L_2) - (L_1 \cap L_2)$$

$$= (a^+ \cup b^+) - \phi$$

$$= a^+ + b^+$$

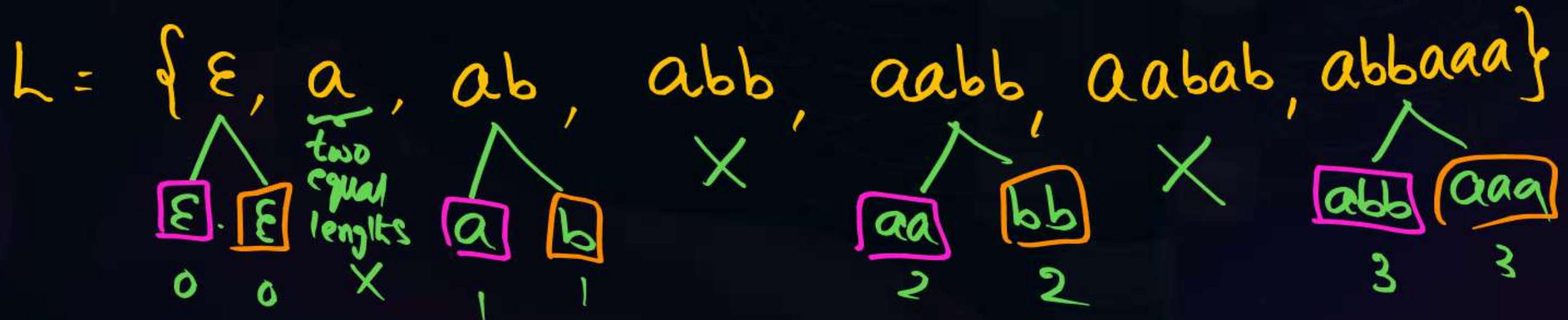
$$(18) \text{Half}(L) = \frac{1}{2}(L) = \text{First Half}(L)$$

$$= \{u \mid uv \in L, |u| = |v|\}$$

$$(19) \text{Second Half}(L) = \{v \mid uv \in L, |u| = |v|\}$$

$$\text{First Half}(L) = \{\epsilon, a, aa, abb\}$$

$$\text{Second Half}(L) = \{\epsilon, b, bb, aaa\}$$



Closure Properties for Regular Languages



(20) $\text{one-third}(L) = \frac{1}{3}(L) = \{ x \mid xyz \in L, \underbrace{|x|=|y|=|z|}_{3 \text{ equal parts}} \}$

(21) $\text{Middle } \frac{1}{3}(L) = \{ y \mid \quad \quad \quad \}$

(22) $\text{Last } \frac{1}{3}(L) = \{ z \mid \quad \quad \quad \}$



$\frac{1}{3}(L) = \{ \epsilon, a, ba \}$, $\text{Middle } \frac{1}{3}(L) = \{ \epsilon, a, bb \}$, $\text{Last } \frac{1}{3}(L) = \{ \epsilon, a, aa \}$

(23) Finite Union

$L_i \rightarrow \text{Regular}$

Union

$$L_1 \cup L_2 = \{ w \mid w \in L_1, \text{ or } w \in L_2 \}$$

$$L_1 \cup L_2 \cup L_3 \cup \dots \cup L_k \Rightarrow \text{Always Regular}$$

k is constant

Each L_i Regular
 k Regulars

(24) Finite Intersection: $L_1 \cap L_2 \cap L_3 \cap \dots \cap L_k \Rightarrow \text{Always Regular}$

(25) Finite Difference: $L_1 - L_2 - L_3 - \dots - L_k \Rightarrow \text{Always Regular}$

(26) Finite Concatenation: $L_1 \cdot L_2 \cdot L_3 \cdot \dots \cdot L_k \Rightarrow \text{Always Regular}$

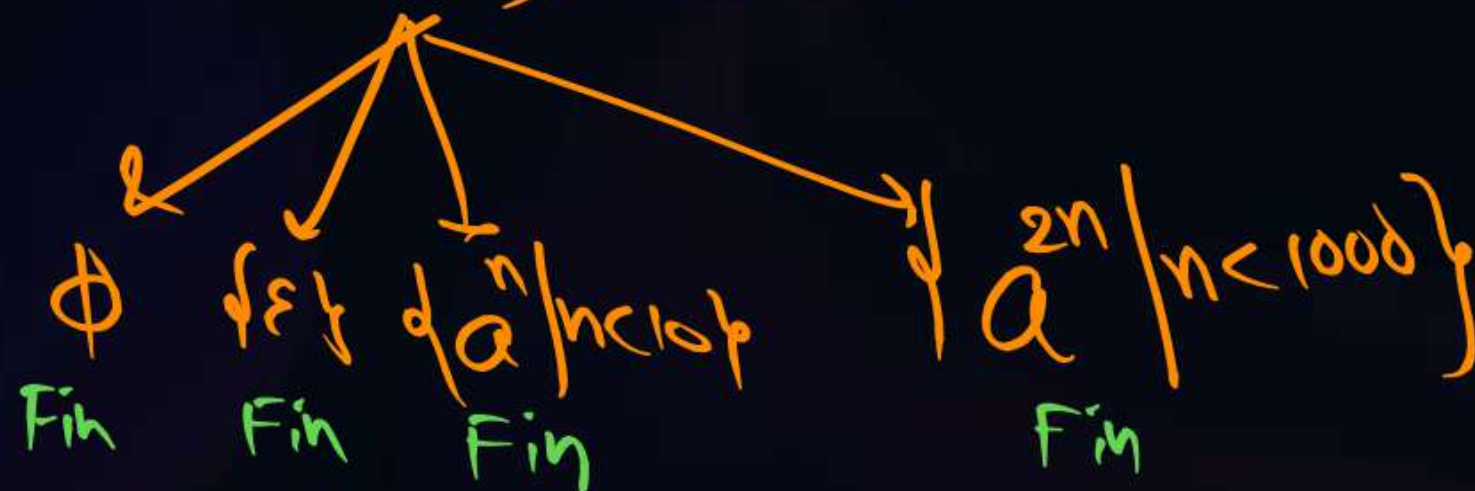
*** (27) Finite Subset

→ Closed for Regulars

Finite subset of Regular is always finite lang
(regular)

$$L = a^*$$

Finite subsets of a^*



Subset

→ Not closed for Regulars

Subset of Regular is either regular or not regular

$$L = a^*$$

Subsets of a^*



(28) Finite Substitution

→ closed for regulars

$L = a^*b$
 L is Regular over $\Sigma = \{a, b\}$



$f(a) = \text{Finite lang}$
 $= \{\epsilon\}$

$f(b) = \text{Finite lang}$
 $= \{0, 1\}$

$f(L) = \epsilon^*(0+1) \Rightarrow \text{Regular}$
 $f(L)$ is Regular

*** (29) Infinite Union : $L_1 \cup L_2 \cup L_3 \cup L_4 \cup \dots \Rightarrow$ either Regular or not reg
 Infinite no. of regulars

Case 1: $\sum_{\text{reg}}^* \cup L_2 \cup L_3 \cup \dots \Rightarrow \sum_{\text{reg}}^*$

Case 2: $\{\epsilon\} \cup \{ab\} \cup \{a^2b^2\} \cup \{a^3b^3\} \cup \dots \Rightarrow a^n b^n$
 Reg reg reg reg Not reg

- (30) Infinite \cap
- (31) Infinite $-$
- (32) Infinite concatenation
- (33) Infinite subset
- (34) Infinite substitution

Not closed for Regulars

Shortcut:

For Regular Languages:

→ Subset
→ Infinite($\cup, \cap, -, \cdot, \subseteq, f$)

Not closed

$$1) L_1 = (aa)^* = \{\epsilon, a^2, a^4, a^6, \dots\}$$

$$L_1 \cup L_2 = (aa)^* + a^* = a^*$$

$$L_2 = a^* = \{\epsilon, a, a^2, a^3, \dots\}$$

$$ii) L_1 \cap L_2 = (aa)^* \cap a^* = (aa)^*$$

$$iii) \overline{L_1} = \overline{(aa)^*} = \sum_{All}^* - (aa)^*_{even} = a(aa)^*_{odd}$$

$$iv) \overline{L_2} = \overline{a^*} = \phi$$

$$v) L_1 - L_2 = \phi$$

$$vi) L_2 - L_1 = a(aa)^*$$

$$vii) L_1 / L_2 = a^*$$

$$L_1 / L_2 = (aa)^* / a^*$$

$$= \{\epsilon, aa, aaaa, aaaaaa, \dots\} / a^*$$

✓	ϵ / ϵ	$aa / \epsilon = aa$
x	ϵ / a	$aa / a = a$
x	ϵ / aa	$aa / aa = \epsilon$
	\vdots	$aa / aaaa = x$
	\vdots	\vdots

$= \epsilon$

2) Which of the following are closed for Regular languages?

i) Union

ii) Subset

iii) Inverse Homomorphism

iv) Finite Intersection

3) Which of the following is TRUE?

✓ A) $\text{prefix}(L) = \{x \mid xy \in L, y \in \Sigma^*\}$

✓ B) $\text{Half}(L) = \{x \mid xy \in L, |x| = |y|\}$

✓ C) $L_1 \cap L_2 = \{x \mid x \in L_1, x \in L_2\}$

✗ D) $L_1 - L_2 = \{x \mid x \in L_2, x \notin L_1\}$

4) $L_1 \rightarrow$ finite language

$L_2 \rightarrow$ Regular language

$L_3 \rightarrow$ Infinite language

$L_4 \rightarrow$ Non-Regular language

TRUE ?

~~A)~~ $L_1 \cup L_2$ is Regular

~~B)~~ $L_1 \cup L_3$ is Infinite

~~C)~~ $L_1 \cap L_4$ is Regular

~~D)~~ $L_2 \cap L_3$ is Infinite

I) Non-reg \cap Finite \Rightarrow always Finite set

Note: Finite set \cap Any \Rightarrow Finite set

II) Non-regular \cap Infinite \Rightarrow either fin or Inf
either reg or nonreg

i) $\underbrace{a^n b^n}_{\text{nonreg}} \cap \underbrace{a^*}_{\text{Inf}} \Rightarrow \{\epsilon\}$
reg, fin

ii) $\underbrace{a^n b^n}_{\text{nonreg}} \cap \underbrace{a^n b^n}_{\text{Inf}} \Rightarrow a^n b^n$
nonreg, Inf

Revision :

2 page short notes

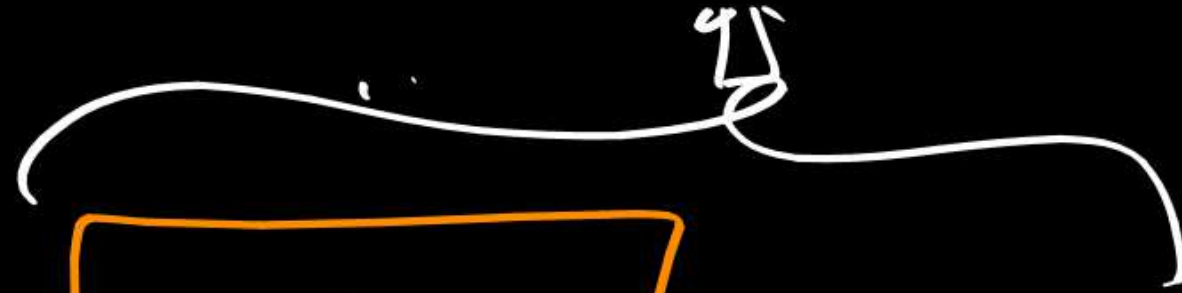
video



regular
Notes
50 page

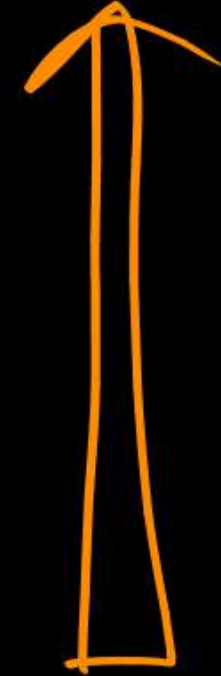


revision
&
practice



$$S \rightarrow bA$$

$$A \rightarrow bS \mid c$$



$$S \rightarrow bA$$

$$S \Rightarrow b \cdot (bb)^* c$$

$$L = (bb)^* bc$$

$$= \underline{\underline{b(bb)^*c}}$$

method 1:

$$A = c + bS$$

$$S \rightarrow b \textcircled{A}$$

$$S \rightarrow bc \mid b \textcircled{b} S$$

$$\boxed{L = (bb)^* bc}$$

method 2:

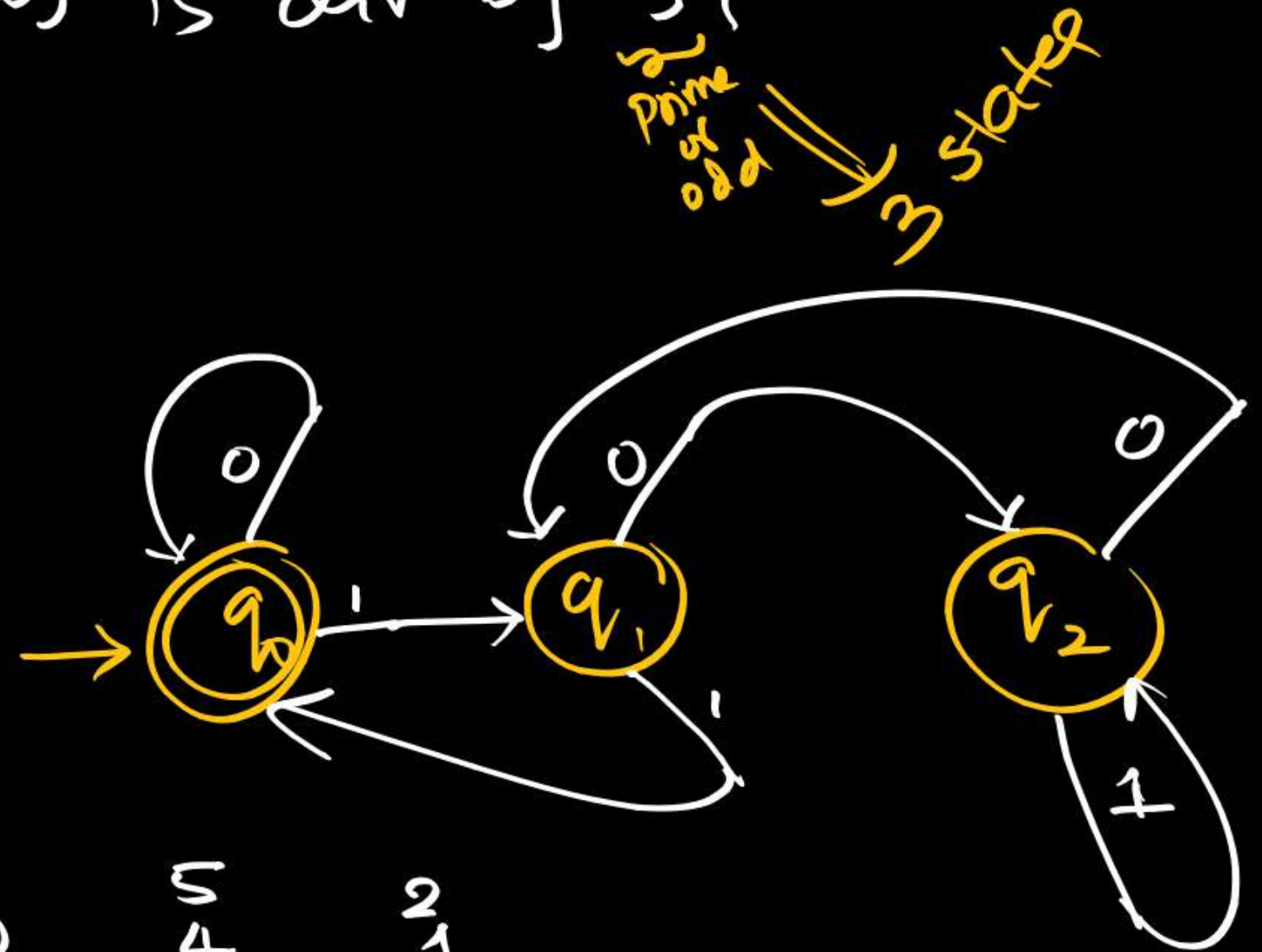
$$A \rightarrow c \mid bS$$

$$A \rightarrow c \mid bbA$$

$$A = (bb)^* c \rightarrow \textcircled{1}$$

$L = \{w \mid w \in \{0,1\}^*, \text{Dec}(w) \text{ is div by } 3\}$

	0	1
$\rightarrow q_0$	q_0	q_1
q_1	q_2	q_0
q_2	q_1	q_2



2 prime or odd \rightarrow 3 states

101
100
Bin

5
4
Dec

2
1
Remainder
0 \rightarrow q_0
1 \rightarrow q_1
2 \rightarrow q_2
0



2 mins Summary



Topic

closure properties for regular languages

Next: Pumping lemma
&

FA with o/p

THANK - YOU