Computer Science

Theory of Computation

Turing Machine



Lecture No.- 2

Recap of Previous Lecture









Topic

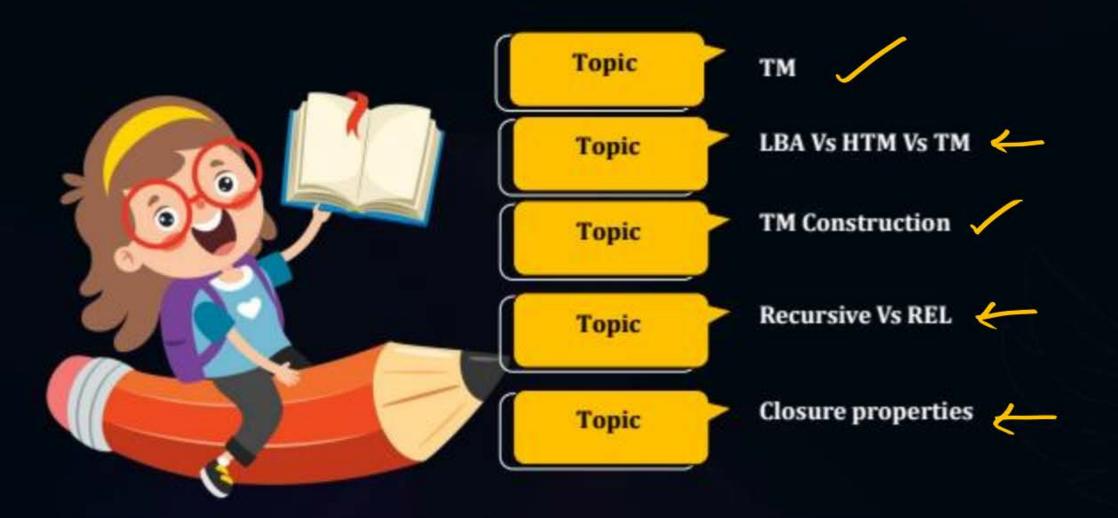
Turing Machine

Topics to be Covered











Recursive Language

RE Language

ATM exist

(TM that always halts)

TM exist

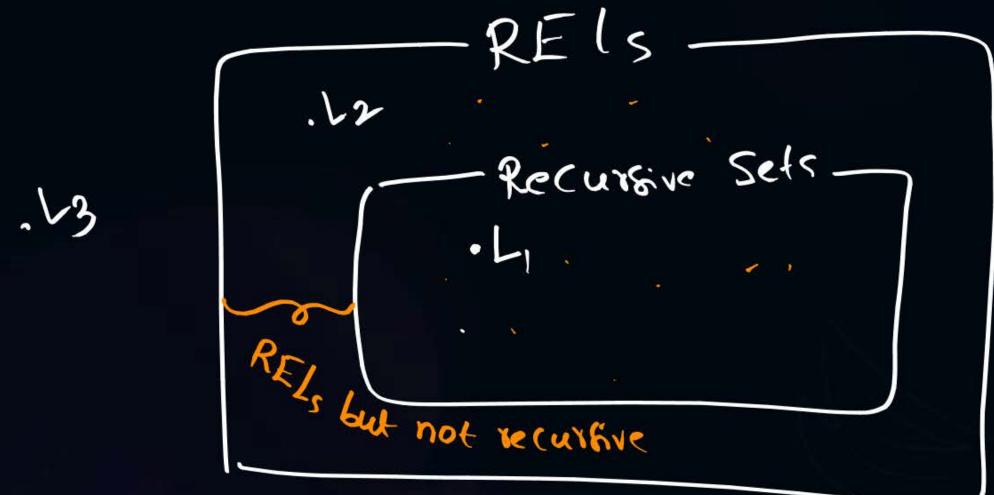
Valid => Halts at sind in Invalid => Halts at my

Valid strings => logic exist

Invalid strings => logic mayor

maynor exist

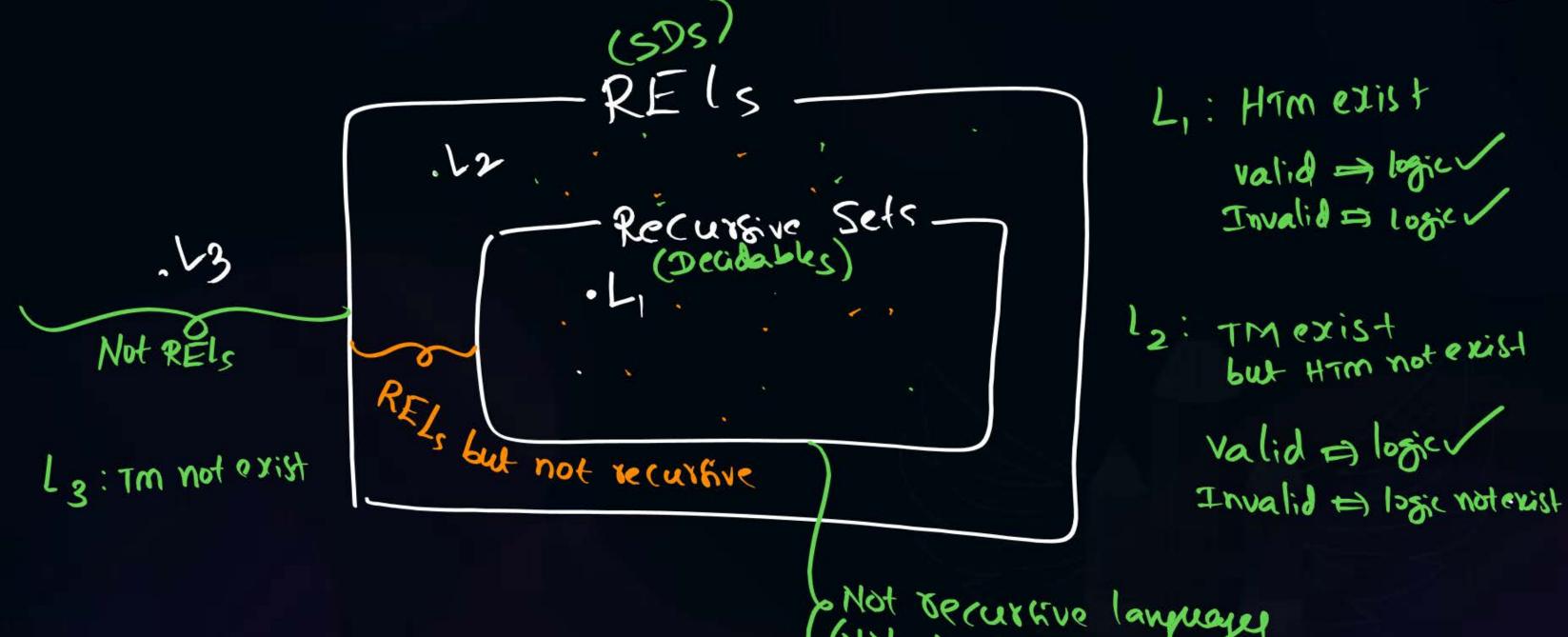




L, is Recursive L, is REL L2 is REL La is REL but rec L2 is not recurrive

L3 is not recurrive





(Not decidable)

(Undecidable)



Linear Bound Automater linearly bounded tape & Recursive Set

HTM exist LITM that always halts

TM exist

Logic exist for valid Logic may a may not exist

Closure properties for Recursive language



- (i) Union
- (2) Intersection
- 3 Complement
- (4) Difference
- (5) Concatenation
- (6) Reversal
- 1) Kleene Stoy
- (8) Kleene plus

- @ Subset
- (10) brefix
- 1 Suffix
- (12) Substring
- (B) f(L)
- (L)
- (15) E- fice h(L)
- (16) h'(L)

- (17) Finite U
- (18) 11
- (19) 11 -
- (20)
- (2) 11 =
- (24) 11 f
- 23) to (28): Infinite (U, n, -, :)



For Recursive language -> Remember Not closed -> Subset · Infinite (U, N, -1., 5, 8) Substitution

Closure properties for REIS



- (i) Union
- (2) Intersection
- (3) Complement
- (I) Difference
- (5) Concatenation
- (6) Reversal
- (7) Kleene Stoy
- (8) Kleene plus

- (Subset
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- (16) h'(L)

- (17) Finite U
- (18)
- 11 -
- (20)
- (2) 11 =
- (22) 11 f
- (23) to (28): Infinite (U, n, -, :,) (5, f)

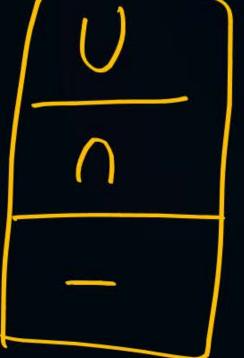


For RELs: L) Remember not Closed Complement D'ifference Finite Difference Subset 7 Infinite (U,n,-,., s,f)

Pw

Moximus)

Recursive Set



Regular

Recurrive Recurrive (may or may not) be regular)

工)

RELIN

Regular

Always REL (may or may not be veg)



- D Rec, U Rec => Recursive
- 2) Rec, n Rec2 => Recurrive
- 3) Rec

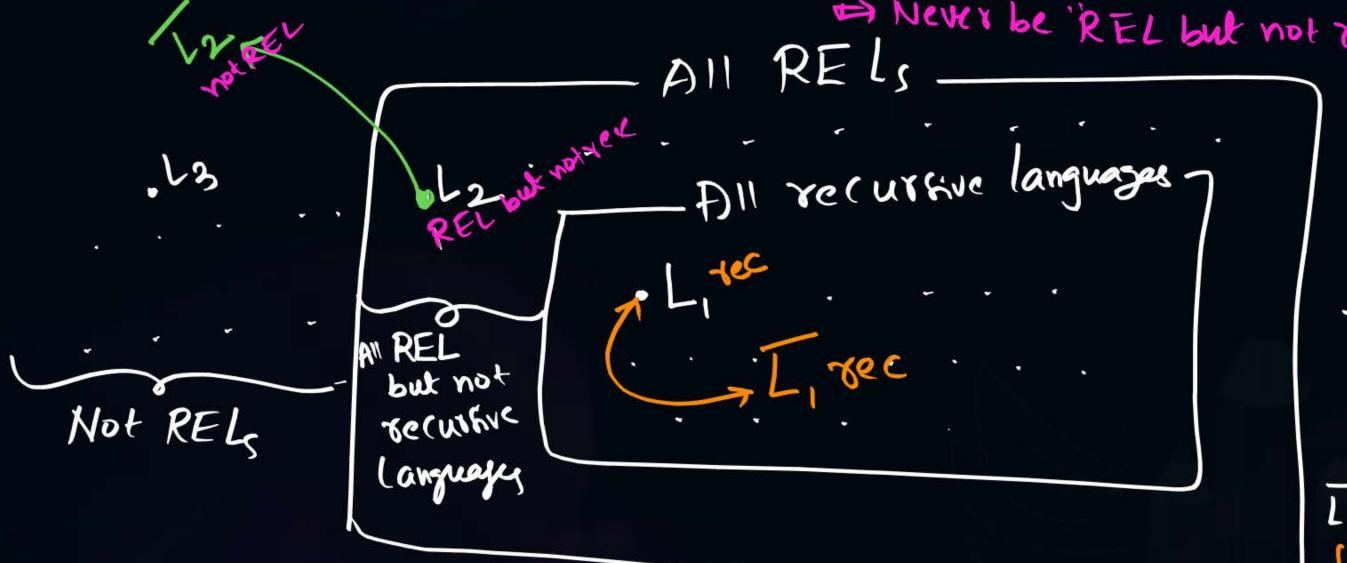
Recursive

- 4) RE, UREL => REL
- 5) RE, ORE, DREL
- (6) RE -> May or may not be
 Peilter returning
 Or not REL
 - > Nevex be "RE but not rec"

REL > Need not be REL

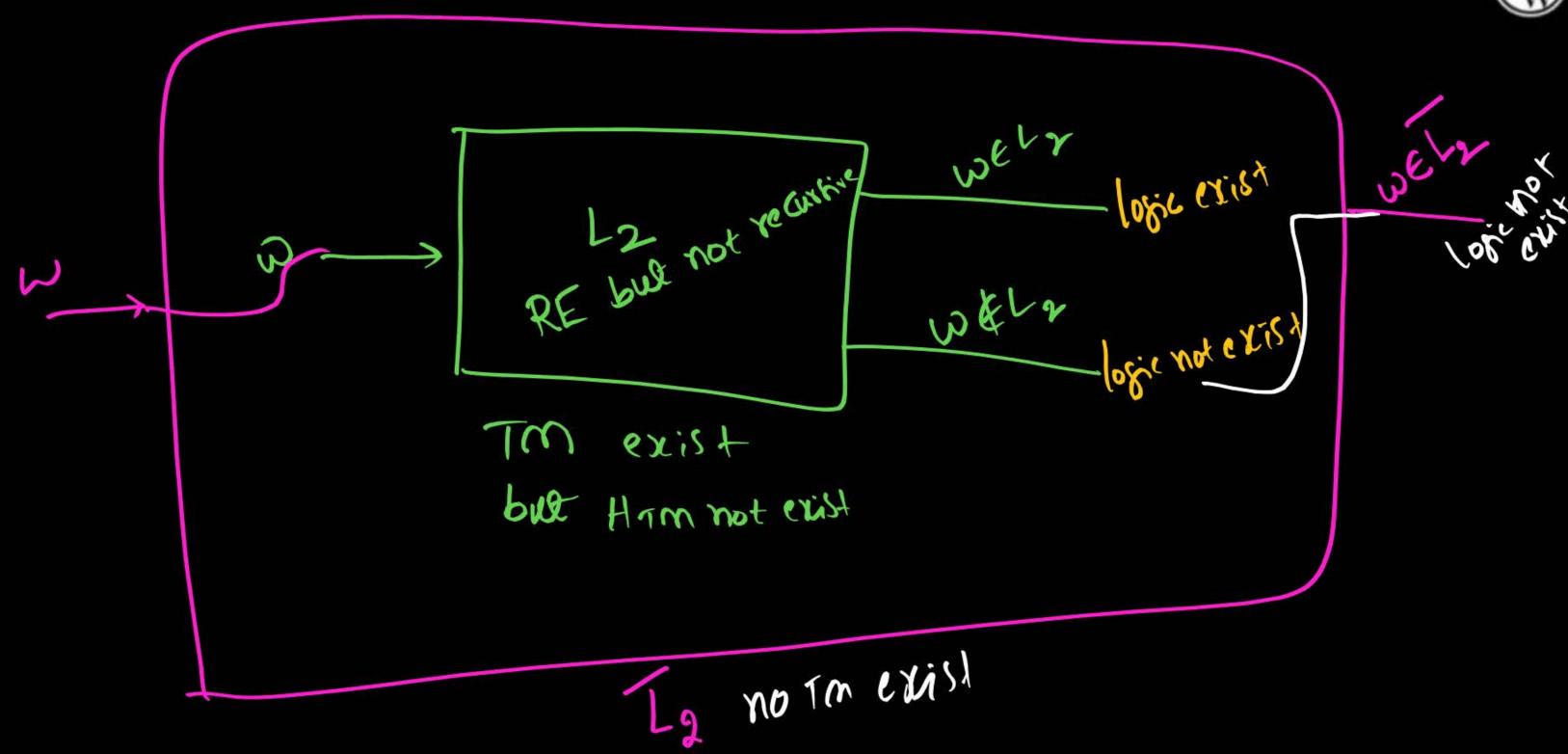


Mever be 'REL but not recursive"



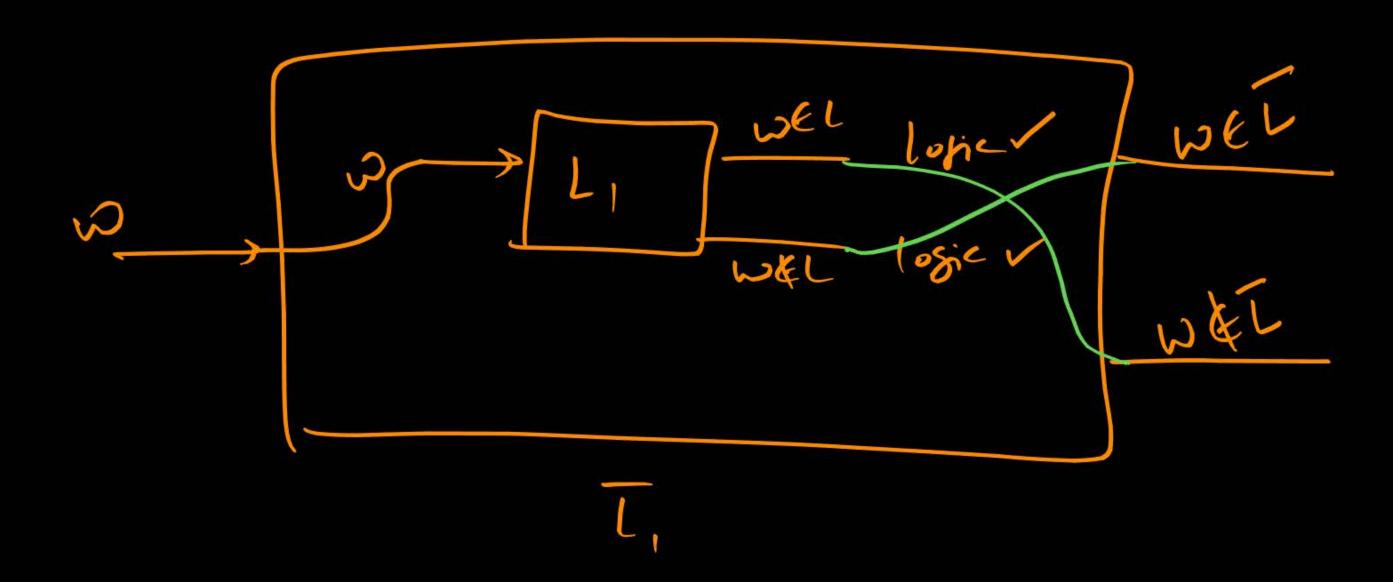
Not REL





not REL





Rec = always Recurrive

Recursive language (Decidable)

HTM exist

TM can enumerate Lin lexicographic orderm (effective)

Im Decidable language

Halting program exist

Algorithm exist

RE language (Semiduidhu)

Tm exist

Im can enumerate L

Enumerable Set (Relognitable)

program exist



Turing Machine Single Tape TM Multi tape Tm Multi head TM multi head & multi tape TM Multi Dimensional tape TM

= 2 Stack PDA = n stack PDA
(n>2) = Universal TM

\$
HTM X
eilker SDUD Not RE



UD = not decidable = not recurrive

HTM not exist

Whid X X

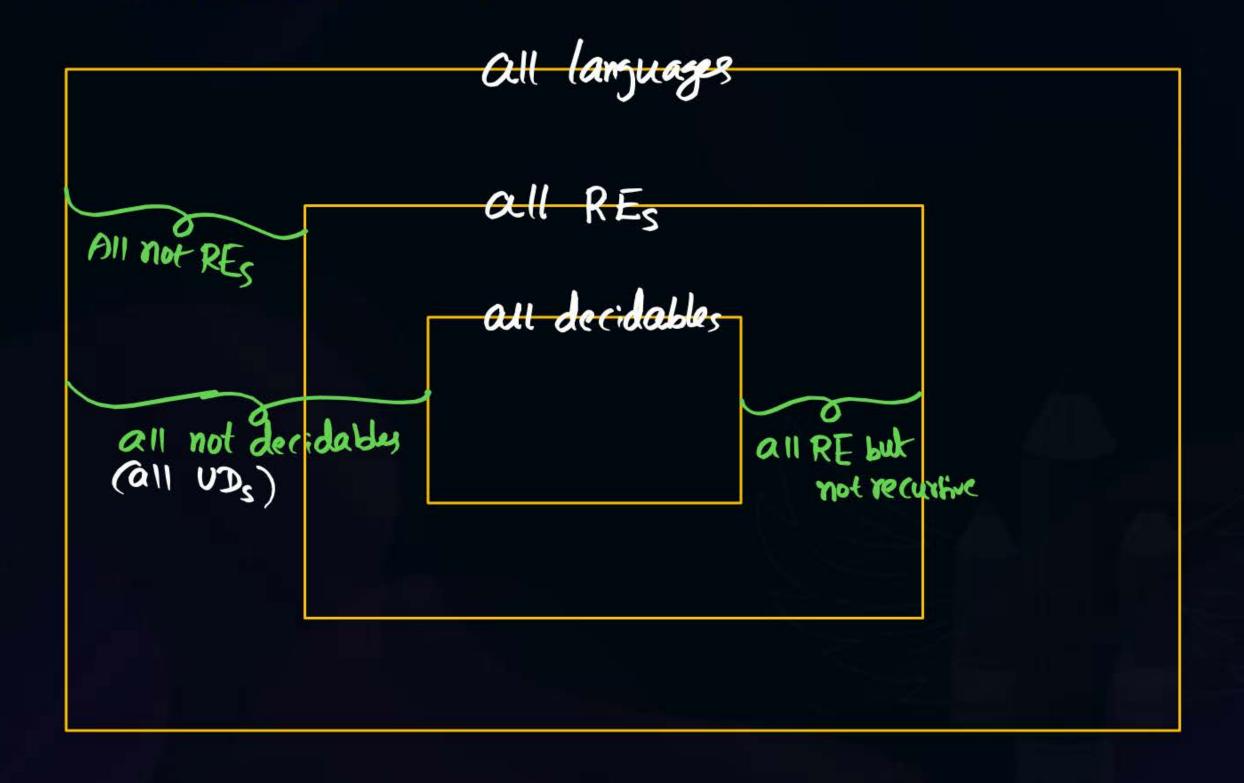
Invalid X X

	lope			ale
	9	5900	1 3	8
Vatid _		1	X	X
Invalid	/	X	1	X
	P		6	7
		- (D	



RE but not recursive Undecidable SD and UD SDUD RE = SD Not rec = Not decidable = undecidable







2 mins Summary



Topic

Closure propulsu

Timport definition

LBA, Him, Tim



THANK - YOU