

# CS & IT ENGINEERING

## Algorithms

Analysis of Algorithms

**DPP- 01**

Discussion Notes



By- Aditya sir



## [MCQ]



#Q. Sort the functions in ascending order of asymptotic (big-O) complexity. *increasing*

$f_1(n) = n$ ,  $f_2(n) = 80$ ,  $f_3(n) = n^{\log n}$ ,  $f_4(n) = \log \log^2 n$ ,  $f_5(n) = (\log n)^{\log n}$

- A**  $f_2(n), f_4(n), f_1(n), f_5(n), f_3(n)$  ✓
- B**  $f_2(n), f_1(n), f_4(n), f_5(n), f_3(n)$  ✗
- C**  $f_2(n), f_1(n), f_4(n), f_3(n), f_5(n)$  ✗
- D**  $f_1(n), f_1(n), f_4(n), f_3(n), f_2(n)$  ✗

**A** ✓

Soln:-

Rate of growth:

decr < Const < Poly < Expo

$T_1(n) = n$  ————— Poly

$T_2(n) = 80$  ————— Const

$T_3(n) = n^{(\log n)}$  ————— Expo —

$T_4(n) = \log((\log n)^2)$  ————— Poly

$T_5(n) = \log n^{(\log n)}$  ————— Expo —

$T_2 < T_4 < T_1 < T_5 < T_3$

$n$        $\log((\log n)^2)$

$n > 2 \log(\log n)$



$$n^{\log n} > \log n^{\log n}$$

$$\cancel{\log n} * \log n$$

$$\cancel{\log n} * \log(\log n)$$

$$\text{Let } \log n \rightarrow x$$

$$\log n$$

$$x$$

$$\log(\log n)$$

$$\log(x)$$

[MSQ]

multiple Soln can be correct



#Q. Consider two function  $f(n) = 10n + 2\log n$  and  $g(n) = 5n + 2(\log n)^2$ , then which of the following is correct option?

- ☒ **A**  $f(n) = \theta(g(n))$
- ☒ **B**  $f(n) = O(g(n))$  Big Oh
- ☒ **C**  $f(n) = \omega(g(n^2))$  small omega
- ☒ **D** None of the above
- $f \leq g$
- $f(n) > g(n^2) * c$
- A, B ✓



Soln:  $f(n) = 10 \times n + 2 \times \log(n)$ ,  $g(n) = \underline{5n} + 2(\log n)^2$

$$f(n) = \Theta(n)$$

O    Ω

$$g(n) = \Theta(n)$$

O    Ω

$$\left. \begin{array}{l} f = \Theta(n) \\ g = \Theta(n) \end{array} \right\} \left. \begin{array}{l} f = \Theta(g(n)) \\ \underline{g = \Theta(f(n))} \end{array} \right\}$$

$$g(n) = 5n + 2(\log n)^2$$

$$g(n^2) = 5n^2 + 2(\log(n^2))^2$$

$$= 5n^2 + 2(2\log(n))^2$$

$$= \underline{5n^2} + 8(\log n)^2 \longrightarrow \underline{O(n^2)}$$

$$f(n) = O(n), \quad g(n^2) = O(n^2)$$

$$f(n) \subset \underline{\underline{g(n^2)}}$$



#Q. Consider two function  $f(n) = \sqrt{n}$  and  $g(n) = n \log n + n$  then  $f(n) / g(n)$  is equivalent to how many of the following given below? \_\_\_\_\_.

**A**  $O(n^{-1/2})$  ✓ *Small oh*

✓ **B**  $O(n^{-1/2})$  *Big oh*

**C**  $\Omega(1/\log n)$  ✗

✗ **D**  $\theta(n^{-1/2})$

$$\frac{f(n)}{g(n)} < \frac{1}{\sqrt{n}}$$

Ans: (2)



Soln:-

$$f(n) = \sqrt{n}$$

$$g(n) = n \log n + n$$

$$\boxed{h(n)} = \frac{f(n)}{g(n)} = \frac{\sqrt{n}}{n \log n + n} = \frac{\sqrt{n}}{n(\log n + 1)} = \boxed{\frac{1}{\sqrt{n}(\log n + 1)}}$$

Check option A:  $h(n) = \frac{1}{\sqrt{n}(\log n + 1)}$  is  $O(n^{-1/2})$ ?

$$\frac{1}{\log n + 1} < 1$$

$h(n)$  is not  
 $\Omega(n^{-1/2})$

↓  
Hence not  
 $O(n^{-1/2})$

$$\frac{1}{\sqrt{n}(\log n + 1)}$$

~~$$\frac{1}{\sqrt{n}(\log n + 1)}$$~~

$$\frac{1}{(\log n + 1)}$$

$$n^{-1/2} \rightarrow \frac{1}{\sqrt{n}}$$

~~$$\frac{1}{\sqrt{n}}$$~~

Hence A & B  
both are True

$$< 1$$



check (  $\frac{1}{\sqrt{n}(\log n + 1)}$  ) is  $\Omega\left(\frac{1}{\log n}\right)$  ?

$$\frac{1}{\sqrt{n} \log n + \sqrt{n}} < \frac{1}{\log n}$$

$$h(n) = \Omega\left(\frac{1}{\log n}\right)$$

X

$$\sqrt{n} \log n + \sqrt{n} > \log n$$

#Q. Consider the following C-code

```
void foo (int n)
```

```
{
```

```
    int a = 1;
```

```
    if (n == 1)
```

```
        return;  $\rightarrow$  exit
```

```
for (; a <= n; a++) for(a=1, a <= n; a++)
```

```
{
```

```
    printf("GATEWALLAH");
```

```
    break;
```

```
}
```

```
}
```

Ans :- (A)

**A**

$O(1)$

**B**

$O(n)$

**C**

$O(\log n)$

**D**

$O(\sqrt{n})$

What is the worst time complexity of above program?



#Q. Consider the following asymptotic functions :

$$f_1 = 2^n$$

$$f_2 = 1.001^n$$

$$f_3 = e^n$$

$$f_4 = n!$$

Ans  $\rightarrow$  (D)

Which of the following is correct increasing order of above functions?

**A**  $f_3, f_4, f_1, f_2$  ✗

**B**  $f_2, f_4, f_1, f_3$  ✗

**C**  $f_3, f_2, f_1, f_4$  ✗

**D**  $f_2, f_1, f_3, f_4$  ✓

Soln :-

$$f_1 = 2^n \longrightarrow \text{Expo}$$

$$f_2 = (1.001)^n \longrightarrow \text{Expo}$$

$$f_3 = e^n \longrightarrow \text{Expo}$$

$$* f_4 = n! \longrightarrow \text{Expo } (n^n)$$

$$L_0 \rightarrow n(n-1) \times (n-2) \dots 1$$

$$\approx n^n$$

$$e \rightarrow \approx 2.71$$

$$1.001 < 2 < 2.71(e)$$

$$(1.001)^n < 2^n < e^n < n^n$$

$$f_2 < f_1 < f_3 < f_4$$



MSC → Multiple can be correct  
options

#Q. Consider ~~two function~~ the following functions:

$$f_1(n) = 4^{2^n}$$

$$f_2(n) = n!$$

$$f_3(n) = 4^{e^n}$$

$$f_4(n) = n^{n^n}$$

Which of the following is/are correct?

Ans:

B, C, D

$$\underline{f_2} < \underline{f_1} < \underline{f_3} < \underline{f_4}$$

$$f_1 = O(f_2)$$

$$f_1 \leq f_2 \quad \times$$

**A**  $f_1(n) = O(f_2(n)) \quad \times$

**B**  $f_1(n) = O(f_4(n))$

$$f_1(n) = O(f_4(n))$$

$$f_1 = O(f_4)$$

$$\downarrow$$
  

$$f_1 \leq f_4$$

**C**  $f_1(n) = O(f_3(n))$

$$f_1 \leq f_3$$

**D**  $f_2(n) = O(f_3(n))$

$$f_2(n) = O(f_3(n))$$

$$f_2 \leq f_3$$

Soln :-

$$\checkmark f_1(n) = (4)^{2^n} \longrightarrow \text{expo}$$

$$\underline{f_2(n) = n!} \longrightarrow (n)^n \longrightarrow \text{expo}$$

$$\checkmark f_3(n) = (4)^{e^n} \longrightarrow \text{expo}$$

$$f_4(n) = (n)^{n^n} \longrightarrow \text{expo}$$

$$f_2 < f_1 < f_3 < f_4$$

$$(4)^{2^n} < (4)^{e^n}$$

$$2^n < e^n$$

$$e \approx 2.71$$

$$f_1 < f_3$$

$$(n)^n < (n)^{n^n}$$

$$f_2 < f_4$$



$$(n)^n$$

<

$$(4)^{e^n}$$

$$\log(n^n)$$

$$\log((4)^{e^n})$$

$$n \times \log n$$

$$e^n \times \log(4)$$

$$n \log n$$

poly

<

$$e^n$$

expo

$$(n)^n$$

<

$$(4)^{2^n}$$

$$\log(n^n)$$

$$\log((4)^{2^n})$$

$$n \neq \log n$$

$$2^n \neq \log(4)$$

$n \log n$

← Poly

<

$2^n$

→ Expo



Ans: C

#Q. Consider two function  $f_1(n) = n^{2^n}$  and  $f_2(n) = n^{n^2}$  then which of the following is true.

**A**  $f_1(n) = O(f_2(n))$  ~~✗~~  $f_1 \leq f_2$

~~✗~~ **B**  $f_1(n) = \theta(f_2(n))$

**C**  $f_1(n) = \omega(f_2(n))$   $\rightarrow$  small omega

~~✗~~ **D** None of these

$$\underline{\underline{f_1 > f_2}}$$

$$f_1 \approx f_2$$

(rate of growth)

$$f_1 > f_2$$

$$\underline{2^n \neq O(n^2)}$$

Soln:-

$$f_1(n) = n^{2^n} > f_2(n) = n^{n^2}$$

log on both sides

$$\log(n^{2^n})$$

$$2^n \times \cancel{\log n}$$

$$\text{expo} \rightarrow 2^n$$

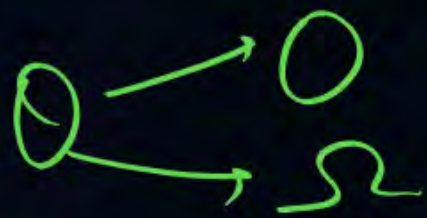
$$\log(n^{n^2})$$

$$n^2 \times \cancel{\log n}$$

$$n^2 \rightarrow \text{poly}$$



[MCQ]



#Q.  $f(n) = \sum_{i=1}^n i^3 = x$ , choices for x

$$x = \frac{1}{4} (n^4 + 2n^3 + n^2)$$

$$x \rightarrow \left. \begin{matrix} O(n^4) \\ \Omega(n^4) \end{matrix} \right\} O(n^4)$$

$$\begin{matrix} x \leq n^4 \rightarrow \text{Tight} \\ x \leq n^5 \\ x \leq n^6 \end{matrix} \left. \vphantom{\begin{matrix} x \leq n^4 \\ x \leq n^5 \\ x \leq n^6 \end{matrix}} \right\} \text{loose}$$

$$x \rightarrow O(n^3) \checkmark$$

$$x \gg n^5 \times$$

$$\begin{aligned} x &= \Omega(n^3) \\ \checkmark x &\gg C * n^3 \end{aligned}$$

I.  $\theta(n^4)$   
III.  $O(n^5)$

II.  $\theta(n^5) \times$   
IV.  $\Omega(n^3)$

**A** I, II, III  $\times$

**D**

**B** II, III, IV

**C** I, II, III, IV

**D** I, III, IV

Soln:

$$x = \sum_{i=1}^n i^3 = 1^3 + 2^3 + 3^3 \dots n^3$$

$$(a+b)^2 = a^2 + 2ab + b^2$$

$$\rightarrow = \left( \frac{n(n+1)}{2} \right)^2 = \left( \frac{n^2+n}{2} \right)^2$$

$$(n^2)^2 = n^{2 \times 2} = \underline{\underline{n^4}}$$

$$= \frac{1}{4} (n^2+n)^2 = \frac{1}{4} \left[ (n^2)^2 + 2 \times n^2 \times n + n^2 \right]$$

$$= \frac{1}{4} \left[ n^4 + 2n^3 + n^2 \right]$$

Dominating  $\downarrow$

$$\underline{\underline{O(n^4)}}$$





**THANK - YOU**

Practice all the  
Concepts  
taught in  
Class.