

# CS & IT ENGINEERING

## Theory of Computation

Regular Languages

Lecture No.- 09

A man with a beard and mustache, wearing a black polo shirt, stands with his arms crossed in front of a blurred bookshelf. He is wearing a black smartwatch on his left wrist.

Mallesham Devasane Sir

# Recap of Previous Lecture



**Topic**

**Regular Language Vs Regular Expression**





# Topics to be Covered



**Topic**

**Regular Language Vs Regular Expression**



Universal Set



# TOPIC:



$\Sigma = \{a, b\}$

Regular Exp.	Regular Language	Meaning
(15) $(a+b)^*$	$\{a, b\}^*$	Set of <u>all</u> strings over $\Sigma = \{a, b\} \Rightarrow$ Universal Language
(16) $(a+b)^+$	$\{a, b\}^+ = \{w \mid w \in (a+b)^*,  w  > 0\}$	Set of <u>all</u> non zero length strings
(17) $\underline{a}(a+b)^*$	$= a \Sigma^* = \{aw \mid w \in \Sigma^*\}$	Set of <u>all</u> strings where every string begins with a.
(18) $b(a+b)^*$	Starting with b	
(19) $(a+b)^*a$	ends with a	
(20) $(a+b)^*b$	" " b	



$$\textcircled{1} \underline{ab^*} \neq \textcircled{2} a(a+b)^*$$

$\downarrow$   
a followed by any b's

Set of (a) strings starting with 'a' over  $\Sigma = \{a, b\}$

$$ab^* \subset a\Sigma^*$$

$$\textcircled{1} \quad ab^* = \{a, ab, abb, abb^2, \dots\}$$

$$\textcircled{2} \quad a(a+b)^* = \{a, aa, ab, aab, aab^2, \dots\}$$

$$\textcircled{1} \subset \textcircled{2}$$





# TOPIC:

Regular Exp.	Regular Language	Meaning
21) $(a^*b)^+$	$= (a+b)^*b$	Set of all strings ending with b
22) $(b^*a)^+$	$= (a+b)^*a = \Sigma^*a$	Set of all strings ending with a
23) $(ab^*)^+$	$= a\Sigma^*$	
24) $(ba^*)^+$	$= b\Sigma^*$	
25) $(a^*b)^*a^*$	$= (a+b)^* = \Sigma^*$	Universal Set over $\Sigma$
26) $(b^*a)^*b^*$	$= (a+b)^* = \Sigma^*$	
27) $a^*(ba^*)^*$	$= (a+b)^* = \Sigma^*$	



$$(a^*b)^+ = (a+b)^*b$$

b ✓

ab ✓

bb ✓

,

,

|

b ✓

ab ✓

bb ✓

,

,

|

,

$$\Sigma = \{a, b\}$$

$$a \Sigma^* = (ab^*)^+$$

$$b \Sigma^* = (ba^*)^+$$

$$\Sigma^* a = (b^* a)^+$$

$$\Sigma^* b = (a^* b)^+$$



$$(a^*b)^+ = (a+b)^*b$$

$$(a^0b)' = b \checkmark$$

$b \checkmark$

$ab \checkmark$

$bb \checkmark$

$$\left(\frac{a^0b}{b}\right)^2 = \left(\frac{a^0b}{b}\right)' \left(\frac{a^0b}{b}\right)' = ab \checkmark$$

$bb \checkmark$



## TOPIC:



$$(28) \quad b^* (ab^*)^* = (a+b)^*$$

$$\downarrow$$
$$b^0 (ab^*)^0 = \varepsilon \checkmark$$

$$b^0 (ab^0)^1 = a \checkmark$$

$$b^1 ( )^0 = b \checkmark$$

$$\downarrow$$
$$\varepsilon \checkmark$$
$$a \checkmark$$
$$b \checkmark$$





## TOPIC:

Write Regular Exp



(29)

$$L = \{ \underbrace{aa}_\text{form} w \mid w \in a^* \}$$

$\begin{array}{c} \swarrow \downarrow \searrow \\ \epsilon \quad a \quad a^2 \end{array}$

$$= \{ aa\epsilon, aa a, aa aa, \dots \}$$

$$= \{ a^2, a^3, a^4, \dots \}$$

$$= \{ a^n \mid n \geq 2 \}$$

$$= aa a^*$$

$$= a^* aa$$

$$= aa^* a$$

$$= aa^+$$

$$= a^+ a$$



## TOPIC:



$$(30) \quad L = \{ w_1 a a w_2 \mid w_1, w_2 \in (a+b)^* \}$$

$$= (a+b)^* a a (a+b)^*$$

Set of all strings containing 'aa' as substring.

$$= \Sigma^* a a \Sigma^*$$



$$(a+b)^* aa (a+b)^*$$

$$(a^*b^*)^* aa (b^*a^*)^+$$

$$\begin{aligned} (a+b)^* &= (a^*b^*)^* \\ &= (a^*b^*)^+ \\ &= (b^*a^*)^* \\ &= (b^*a^*)^+ \\ &= \vdots \end{aligned}$$



## TOPIC:



$$(31) \{w \mid w \in \{a, b\}^*, |w| = 2\}$$

$$= (a+b)^2 = (aa+ab+ba+bb) = \Sigma^2$$

$$(32) \{w \mid w \in \{a, b\}^*, |w| \geq 2\}$$

$$= (a+b)^2 \cdot (a+b)^* = \Sigma^2 \Sigma^* = \Sigma \Sigma^+ = \Sigma^+ \Sigma$$

$$(33) \{w \mid w \in \{a, b\}^*, |w| \leq 2\}$$

$$= \epsilon + a + b + aa + ab + ba + bb = \epsilon + a + b + (a+b)^2 = (\epsilon + a + b)^2$$



$$(\varepsilon + a + b)^2$$

len

$$\varepsilon \rightarrow 0$$

$$a + b \rightarrow 1$$

$$\varepsilon + a + b \rightarrow \leq 1$$

$$(\varepsilon + a + b)^2 \rightarrow \leq 2$$

$$(\varepsilon + a + b) \cdot (\varepsilon + a + b)$$

$$\varepsilon + a + b + aa + ab + ba + bb$$





## TOPIC:



(34)

$\{w \mid w \in \{a, b\}^*, |w| \equiv \text{even} \}$  *divisible by 2*

$= \{\epsilon, aa, ab, ba, bb, aaaa, aaab, \dots\}$

$= (aa+ab+ba+bb)^* = [(a+b)^2]^* = (\Sigma^2)^*$

(35)

$\{w \mid w \in \{a, b\}^*, |w| \equiv \text{odd} \}$

$(\Sigma^2)^* \cdot \Sigma = \Sigma \cdot (\Sigma^2)^*$

$= ((a+b)^2)^* (a+b)$

$= (a+b) \cdot ((a+b)^2)^*$

$$\left( (a+b)^2 \right)^* = (aa+ab+ba+bb)^*$$


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$$\begin{aligned} \left( (\epsilon+a+b)^2 \right)^* &= (\epsilon + \underline{a} + \underline{b} + aa + ab + ba + bb)^* \\ &= (a+b)^* \end{aligned}$$



$$n_a(w) = 2$$

no. of a's in w is 2

$$\#_a(w) = 2$$



## TOPIC:



$$(36) \quad \{w \mid w \in \{a, b\}^*, n_a(w) = 2\} \\ = b^* a b^* a b^*$$

$$(37) \quad \{\boxed{w} \mid w \in \{a, b\}^*, n_a(w) \geq 2\} \\ b^* a b^* a (a+b)^* = (a+b)^* a (a+b)^* a (a+b)^* = \Sigma^* a \Sigma^* a \Sigma^*$$

$$(38) \quad \{w \mid w \in \{a, b\}^*, n_a(w) \leq 2\} \\ b^* (a+\varepsilon) b^* (a+\varepsilon) b^*$$

$$n_a(w) = 0 \Rightarrow b^*$$

OR

$$n_a(w) = 1 \Rightarrow b^* a b^*$$

OR

$$n_a(w) = 2 \Rightarrow b^* a b^* a b^*$$

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$$n_a(w) \leq 2 \Rightarrow \underbrace{b^*}_{\text{b}} + \underbrace{b^* a b^*}_{\text{bab}} + \underbrace{b^* a b^* a b^*}_{\text{babab}}$$

$$= b^* (a + \epsilon) \quad b^* (a + \epsilon) b^*$$





TOPIC:

Home work



$$(39) \quad \{w \mid w \in \{a, b\}^*, n_a(w) = \text{even}\}$$

$$(40) \quad \{w \mid w \in \{a, b\}^*, n_a(w) = \text{odd}\}$$



## 2 mins Summary



Topic

Operators

Topic

Properties

Topic

Simplification

**THANK - YOU**