Computer Science

Theory of Computation

Regular Languages and Non-regular Languages



Lecture No.- 2

Recap of Previous Lecture







Topics to be Covered







Topic Regular Languages

Non-regular Languages



Language (Set)

Regular Set

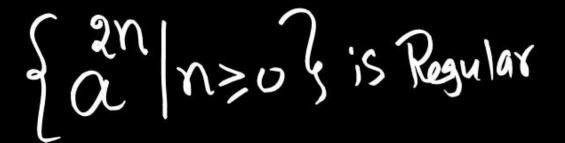
(not regular)



Language Finite Language It is always regular More Kan 1 Symbol L over 1 symbol No Infinite Dependency IP L notforms A.P. - forms A.P than Lis regular Regulat then Lis not regular

 $\begin{cases} x & x \\ x & b \end{cases} = \begin{cases} m & n \\ a & b \end{cases}$ No Insinite dependency Regular

of ab = ge, ab, 2b, 3b, ... Infinite Dependency Not regular



E, a, a, a, ...

0,2,4,6,. A.P. Sixies



Ja no is not regular a, a, a, a, 1,2,4,8,... not A.P. series



Note:

- I) Every finite language is regular.
- I) Regular language is eilker finite or infinite
- III) Infinite language is cilker regular or not regular



$$\int \left\{ \frac{a}{a} \right\} \left\{ \frac{m}{m} \right\} = \frac{x}{a}$$

$$\times$$
 (2) $\{a_b^m\}_{m=n}^n\} = \{a_b^m\}_{m=n}^n$



$$\sqrt{6}$$
 $\frac{d^{2}m^{2}}{d^{2}m^{2}}$ $m=n=10$ $\frac{1}{2}$ $\frac{d^{2}m^{2}}{d^{2}m^{2}}$

$$X = \{a^n b^n m > n > 5\}$$

$$\frac{1}{9} \left\{ \frac{1}{a^m b^n} \right\}_{m=1, n=even} = \frac{1}{a^n b^n} =$$

(10)
$$damba | m=even, n=odd = a bd = (aa) b(bb)$$



$$\varphi$$

$$\chi(19) \left\{ a^n b^{2n} \right\} = \left\{ a^n b^{2n} \middle| n > 0 \right\}$$



m+n=even

I) m=even, n=even

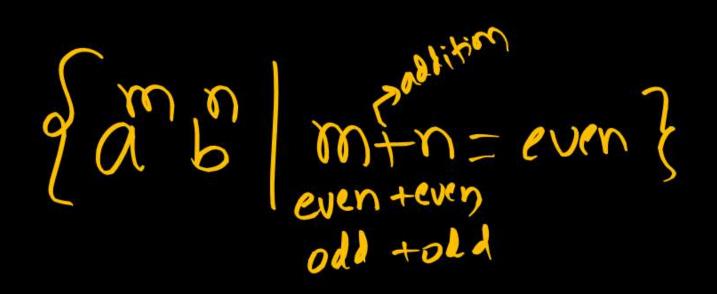
II) modl, noodd

M+N=odd



m=n=even

m=even, n=even



= aven even + add bdd

= (aa)*(bb)* + a(aa)* b(bb)*



even+even=even

even+odd=odd

odd+even=odd

odd+even=odd

odd+odd=even



(16)
$$\{a^{m}b^{n}|_{LCM(m,n)=1}^{n}=\{a^{l}b^{l}\}$$

$$\left| \sum_{n=1}^{\infty} \left(\sum_{n=1}^{\infty} \left(\sum_{n=1}^{\infty} \left(\sum_{n=1}^{\infty} \sum_{n=1}^{\infty} \left(\sum_{n=1}^{\infty} \sum_{n=1}^{\infty} \sum_{n=1}^{\infty} \sum_{n=1}^{\infty} \left(\sum_{n=1}^{\infty} \sum_{n=1}^{$$





$$\{a^mb^k\}=\{a^mb^k|m\geqslant 0,n\geqslant 0,k\geqslant 0\}=a^kb^k\}$$





K. S. Sce

M 7 0,1,2. . .

21 792,4,6,.

311+2 => 2,58,11.

$$\times (28) \{ a^{5} \} = \{a, a, a^{32}, \dots \}$$

$$(30)$$
 (30)

$$(32)_{x}$$
 $\{\alpha^{n}\}_{s} = \{\alpha, \alpha, \alpha, \alpha^{n}, \dots\}$



$$-\left\{\tilde{a}\right\}n+1\right\}=\varepsilon+aaa^*$$



$$gan | n=2 \text{ or } 3$$

$$\sqrt{35}$$
 $\frac{1}{9}$ $\frac{1}{1}$ $\frac{1}{1$

$$\sqrt{39} \quad \{a^{2n} \mid n \ge 0\}^* = \{\epsilon, a^*, a^*, a^*, a^*, a^*, a^*\}^* = \{aa\}^* - \{a^*\}^*$$

$$\sqrt{33} \left\{ \frac{n^2}{n^2} \right\}_{n \geq 0} = \left\{ \frac{1}{a} \right\}^*$$



$$38 \left\{ \frac{2^n}{\alpha^2} \mid n \ge 0 \right\}^* = \left\{ \frac{2^n}{\alpha^2} \right\}^* \cup \dots = \alpha$$

$$(39) \quad \{a^n\}^* = a^*$$





$$\{a^{n+2}b^{n+j}|n,i,j\geq 0\}$$





$$\begin{array}{c}
n=0 \text{ i)} \quad a'b' = a'b' \\
n=1 \text{ i)} \quad a'b' = a'b' \\
n=2 \text{ i)} \quad$$

H.W.:



(45) $\{\omega \omega | \omega \in \{a,b\}^*\}$

(46) $\{\omega\omega^{R} | \omega \in \{a,b\}^{*}\}$

(47) g w#w | w ∈ {a, b}* }

(48) $\{\omega \# \omega^{R}\}$

(49) (49) (49) (49)

 W^{R} : Reverse of WIf W = abb.

then $W^{R} = bba$



2 mins Summary



Topic

Regular Languages

Topic

Non-regular Languages



THANK - YOU