## CS & IT ENGINEERING

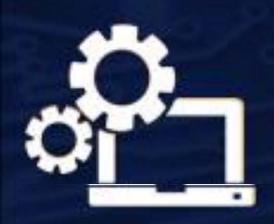


Regular Languages & Non Regular Languages

**DPP 04** Discussion Notes



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TOPICS TO BE COVERED

01 Question

02 Discussion

Q.1

Consider alphabet  $\Sigma = \{a, b\}$ , the empty string  $\in$  and the set of strings S,  $\bigcup$  P, Q and R generated by the corresponding non-terminals of a regular grammar. S, P, Q and R related as follows (S is a start symbol): [MCQ]

ap 
$$S \rightarrow \widehat{aP} | \underline{bQ} | \in$$
 $AbR \qquad P \rightarrow \widehat{bR} | aS$ 
 $AbaO \qquad Q \rightarrow aR | \widehat{bS}$ 
 $AbaO \qquad R \rightarrow \underline{aQ} | \underline{bP}$ 

- A.  $L = \{w: n_a(w) \text{ and } n_b(w) \text{ both are even} \}$
- B.  $L = \{w: n_a(w) \text{ and } n_b(w) \text{ both are odd}\}.$
- C.  $L = \{w: n_a(w) \text{ or } n_b(w) \text{ are even}\}.$
- D. None of these.

Consider the following language L on alphabet

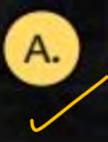


$$\Sigma = \{a, b\}$$

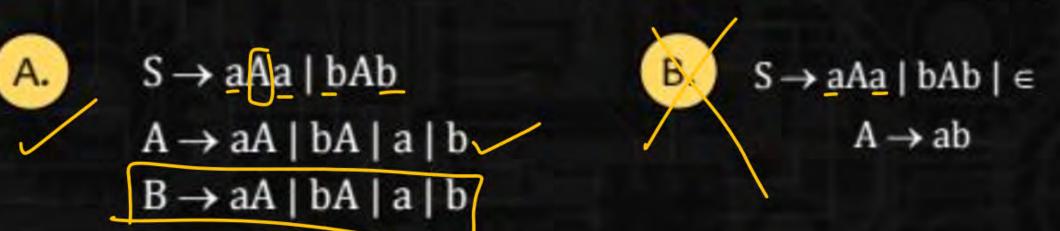
$$\sum = \{a, b\}$$

$$L = \{wxw^{R} \mid w, x \in \{a, b\}^{+}\} = \alpha x \alpha + b x b = \alpha (\alpha + b) \alpha + L(\alpha + b) b$$

The correct regular grammar of above language is/are possible?



$$S \rightarrow \underline{aAa} \mid \underline{bAb}$$
  
 $A \rightarrow aA \mid \underline{bA} \mid \underline{a} \mid \underline{b}$   
 $B \rightarrow aA \mid \underline{bA} \mid \underline{a} \mid \underline{b}$ 

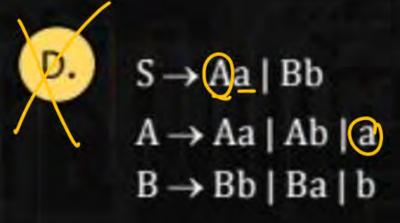


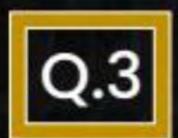


$$S \rightarrow \underline{a}A \mid bB$$

$$A \rightarrow aA \mid bA \mid \underline{a}$$

$$B \rightarrow bB \mid aB \mid b$$





Consider the following statements:



 $S_1$ : If language is regular then, grammar must be regular

[MCQ]

 $S_2$ : If grammar is regular then, language can't be regular.

Which of the following is correct?

- A.  $S_1$  is true.
- $S_2$  is true.
- Both  $S_1$  and  $S_2$  are true.
- D. None of these

## Consider the following grammar G:



G:

$$S \rightarrow ABC$$

$$A \rightarrow aA \mid a \qquad \uparrow \qquad L = abc \stackrel{*}{\sim} = abc \stackrel{*}{\sim}$$

The language generated by above grammar is?



For language  $= \{b \ a*b\}$  the minimum pumping length will be \_\_\_. [NAT]

-0 b - 0 b - 0



Consider some regular expression:



r<sub>1</sub>: a\*bb\*c\*(ab)\*

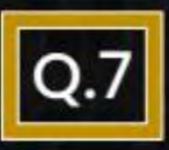
[NAT]



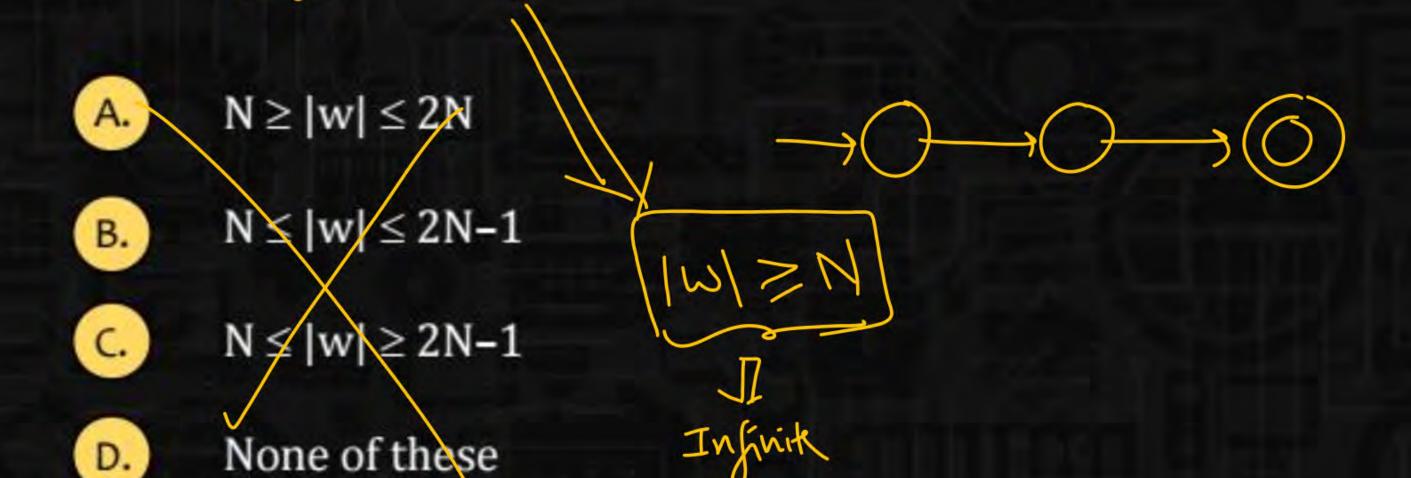
 $\mathbf{r}_2$ :  $\mathbf{a}^*\mathbf{b}^*\mathbf{a}\mathbf{b} \cup (\mathbf{b}\mathbf{b})^*$ 

If minimum pumping length of  $r_1$  is  $P_1$  and minimum pumping of  $r_2$  is  $P_2$  then the value of  $P_1 * P_2$  will be  $\frac{6}{2}$ .

3



Suppose, a language L has finite automata M with N states. The W language generated by FA is L(M) is an infinite if and only if  $\exists_w \in L$  such that



Consider the following grammars G<sub>1</sub> and G<sub>2</sub>:



G<sub>1</sub>: 
$$S \rightarrow aS | S | A$$
  
 $A \rightarrow aA | abA | \in$ 

Which of the following grammar is/are regular?

- A.  $G_1$  only
- $G_2$  only
- C. Both G<sub>1</sub> only G<sub>2</sub>
- D. None of these



