My experience

Monday, December 24, 2012

Segment Trees and lazy propagation

In this topic i will explain a very interesting data structure that can be used to solve a specific set of problems. I will start by explaining its definition and the proceeding with an example problem to solve with it.

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What is segment trees?

Segment Trees is a Tree data structure for storing intervals, or segments, It allows querying which of the stored segments contain a given point. It is, in principle, a static structure; that is, its content cannot be modified once the structure is built. It only uses O(N Ig(N)) storage.

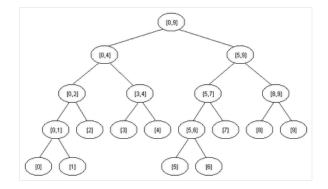
A segment trees has only three operations: build_tree, update_tree, query_tree.

Building tree: To init the tree segments or intervals values Update tree: To update value of an interval or segment Query tree: To retrieve the value of an interval or segment

Example Segment Tree:

- The first node will hold the information for the interval [i, j]
- If i<j the left and right son will hold the information for the intervals [i, (i+j)/2] and [(i+j)/2+1, j]

Notice that the height of a segment tree for an interval with **N** elements is **[logN] + 1**. Here is how a segment tree for the interval **[0, 9]** would look like:



Order of growth of segment trees operations

build_tree: O(N lg(N))update_tree: O(lg(N + k))query_tree: O(lg(N + k))

K = Number of retrieved intervals or segments

Show me your code

- 1 /**
- 2 * In this code we have a very large array called arr, and very large set of operations
- * Operation #1: Increment the elements within range [i, j] with value val

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About Me

Hussein El-Sayed

Cairo, Egypt

Software engineer at ITWorx

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```
4 * Operation #2: Get max element within range [i, j]
5 * Build tree: build tree(1, 0, N-1)
6 * Update tree: update tree(1, 0, N-1, i, j, value)
7 * Query tree: query_tree(1, 0, N-1, i, j)
8 */
10#include<iostream>
11#include<algorithm>
12using namespace std;
14#include<string.h>
15#include<math.h>
16
17#define N 20
18#define MAX (1+(1<<6)) // Why? :D
19#define inf 0x7fffffff
21int arr[N];
22int tree[MAX];
23
24/**
25 * Build and init tree
26 */
27void build_tree(int node, int a, int b) {
28
    if(a > b) return; // Out of range
29
          if(a == b) { // Leaf node
30
                 tree[node] = arr[a]; // Init value
31
32
                  return;
33
34
          build tree(node*2, a, (a+b)/2); // Init left child
35
          build tree (node*2+1, 1+(a+b)/2, b); // Init right child
36
37
38
          tree[node] = max(tree[node*2], tree[node*2+1]); // Init root value
39}
40
41/**
42 * Increment elements within range [i, j] with value value
43 */
44void update tree(int node, int a, int b, int i, int j, int value) {
45
46
          if(a > b \mid \mid a > j \mid \mid b < i) // Current segment is not within range [i, j]
47
                  return;
48
          if(a == b) { // Leaf node
49
50
                  tree[node] += value;
51
                  return;
52
53
          update tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
54
          update_tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right child
55
56
          tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root with max value
57
58}
60/**
61 * Query tree to get max element value within range [i, j]
62 */
63int query_tree(int node, int a, int b, int i, int j) {
64
65
          if(a > b \mid \mid a > j \mid \mid b < i) return -inf; // Out of range
66
67
          if (a >= i && b <= j) // Current segment is totally within range [i, j]
68
                  return tree[node];
69
70
          int q1 = query tree(node*2, a, (a+b)/2, i, j); // Query left child
71
          int q2 = query_tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right child
72
73
          int res = max(q1, q2); // Return final result
74
75
          return res;
76}
77
78int main() {
          for(int i = 0; i < N; i++) arr[i] = 1;
```

```
80
81     build_tree(1, 0, N-1);
82
83     update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
84     update_tree(1, 0, N-1, 7, 10, 12); // Increment range [7, 10] by 12
85     update_tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100
86
87     cout << query_tree(1, 0, N-1, 0, N-1) << endl; // Get max element in range [0, N-1]
88}
```

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segment tree.cpp view raw

Lazy Propagation

Sometimes a segment tree operation wouldn't survive if the problem constraints is too large, here it come lazy propagation along with the segment tree.

In the current version when we update a range, we branch its childs even if the segment is covered within range. In the lazy version we only mark its child that it needs to be updated and update it when needed.

```
/**
   \star In this code we have a very large array called arr, and very large set of operations
1 2 * Operation #1: Increment the elements within range [i, j] with value val
3 4 * Operation #2: Get max element within range [i, j]
5 6 * Build tree: build_tree(1, 0, N-1)
7 8 * Update tree: update tree(1, 0, N-1, i, j, value)
   * Query tree: query_tree(1, 0, N-1, i, j)
10 */
11
12 #include<iostream>
13 #include<algorithm>
14 using namespace std;
16 #include<string.h>
17 #include<math.h>
19 #define N 20
20 #define MAX (1+(1<<6)) // Why? :D
21 #define inf 0x7fffffff
22
23 int arr[N];
24 int tree[MAX];
25 int lazy[MAX];
26
27 /**
28 * Build and init tree
29 */
30 void build tree(int node, int a, int b) {
          if(a > b) return; // Out of range
31
32
          if(a == b) { // Leaf node
33
                  tree[node] = arr[a]; // Init value
34
                  return;
35
36
37
          build tree(node*2, a, (a+b)/2); // Init left child
38
          build tree (node*2+1, 1+(a+b)/2, b); // Init right child
39
40
          tree[node] = max(tree[node*2], tree[node*2+1]); // Init root value
41
42 }
43
44 /**
45 * Increment elements within range [i, j] with value value
46 */
47 void update_tree(int node, int a, int b, int i, int j, int value) {
48
          if(lazy[node] != 0) { // This node needs to be updated
49
                   tree[node] += lazy[node]; // Update it
50
51
52
                   if(a != b) {
                           lazy[node*2] += lazy[node]; // Mark child as lazy
53
                           lazy[node*2+1] += lazy[node]; // Mark child as lazy
54
                   }
55
```

```
57
                   lazy[node] = 0; // Reset it
58
59
60
           if (a > b \mid | a > j \mid | b < i) // Current segment is not within range [i, j]
                  return;
61
62
          if(a >= i && b <= j) { // Segment is fully within range
63
                   tree[node] += value;
64
65
                   if(a != b) { // Not leaf node
66
                           lazy[node*2] += value;
67
68
                           lazy[node*2+1] += value;
69
70
71
                   return;
72
73
           update_tree(node*2, a, (a+b)/2, i, j, value); // Updating left child
74
75
          update_tree(1+node*2, 1+(a+b)/2, b, i, j, value); // Updating right child
76
           tree[node] = max(tree[node*2], tree[node*2+1]); // Updating root with max value
77
78 }
79
80 /**
81 * Query tree to get max element value within range [i, j]
82 */
83 int query_tree(int node, int a, int b, int i, int j) {
84
85
          if(a > b || a > j || b < i) return -\inf; // Out of range
86
87
           if(lazy[node] != 0) { // This node needs to be updated
                   tree[node] += lazy[node]; // Update it
88
89
                   if(a != b) {
90
                           lazy[node*2] += lazy[node]; // Mark child as lazy
91
                           lazy[node*2+1] += lazy[node]; // Mark child as lazy
92
93
94
                  lazy[node] = 0; // Reset it
95
96
           }
97
           if (a >= i && b <= j) // Current segment is totally within range [i, j]
98
99
                   return tree[nodel:
100
101
           int q1 = query_tree(node*2, a, (a+b)/2, i, j); // Query left child
          int q2 = query_tree(1+node*2, 1+(a+b)/2, b, i, j); // Query right child
102
103
104
           int res = max(q1, q2); // Return final result
105
106
          return res;
107}
108
109int main() {
110
          for (int i = 0; i < N; i++) arr[i] = 1;
111
112
          build_tree(1, 0, N-1);
113
114
          memset(lazy, 0, sizeof lazy);
115
          update_tree(1, 0, N-1, 0, 6, 5); // Increment range [0, 6] by 5
116
117
          update_tree(1, 0, N-1, 7, 10, 12); // Incremenet range [7, 10] by 12
118
          update tree(1, 0, N-1, 10, N-1, 100); // Increment range [10, N-1] by 100
119
          cout << query_tree(1, 0, N-1, 0, N-1) << endl; // Get max element in range [0, N-1]
```

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Sample Problems to try

- Quadrant Queries
- D-Query

References

- Wiki
- Topcoder tutorials

Posted by Hussein El-Sayed at 4:58 AM

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11 comments:



donnghi December 24, 2012 at 12:25 PM

In lazy's version, i thinks it's better if you replace update tree[2*node] and tree[2*node+1] in 49th, 50th, 80th and 81th line by lazy[2*node] and lazy[2*node+1].

Its reason is your query is not really come down to higher level, so lazy[] should be updated Reply



Hussein El-Sayed December 24, 2012 at 10:44 PM

I can't understand you :)

Reply



donnghi December 25, 2012 at 3:03 AM

So, in line 80th:

tree[node*2] += lazy[node]; // Mark child as lazy tree[node*2+1] += lazy[node]; // Mark child as lazy

=> replaced by: lazy[node*2] += lazy[node]; lazy[node*2+1] += lazy[node];

It's correct?

Reply



Hussein El-Sayed December 25, 2012 at 3:29 AM

Yes you are totally right:).. thanks for correcting me;)...

Reply



Hussein El-Sayed December 25, 2012 at 3:30 AM

The same at line 49 and 50, updated check it now and tell me :)

Reply



Sandipan Manna December 31, 2012 at 12:13 PM

but your program gives output as 117 !!!

Reply



Hussein El-Sayed January 1, 2013 at 3:06 AM

No it should be 113, however the size of the array needs to be (1+(1<<6)) as it should be $2^{(1+\lg N)}$.

Also there was some checks needed to be added in the lazy version.. please check it and get back to me.

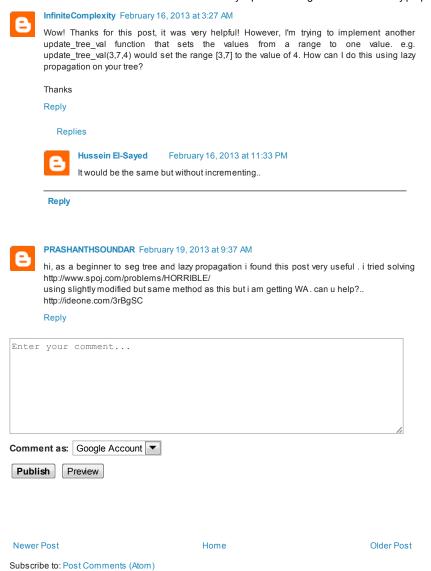
Reply



Sandipan Manna January 11, 2013 at 11:19 PM

Yes your segment tree size should be int x = (int)(ceil(log2(N)))+1; size = (1 << x); This one!

Reply



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