Project Report-- ##Transaction Manager## (Assignment2)

Division of Labor:

Team 31

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· Neha Thokala(nxt0631@mavs.uta.edu, 1002030631) (implementation of commit, abort)

Overall Status

- Initially, I followed the steps provided to compile and build the code locally. To gain a better understanding of C and C++ programming language we referenced the FAQ section provided by Stroustrop (https://www.stroustrup.com/bs_faq2.html) and other sources online.
- After gaining a better understanding of how the main function runs, we implemented readtx(),writetx(), CommitTx(), and AbortTx() in TxMgr and TxtTx code files.
- Test files were given by sir, and we added new test scenarios to test edge cases.
- After testing locally, copied source files to the UTA omega server.
- The code was compiled and built successfully.
- Test scenarios generated the expected results as log files

Login at omega and create a directory.
-mkdir tmproject2

After making tmproject2 folder in omega, copy files by running this command in a local terminal

- pscp -r /Users/sakshisrivastava/Desktop/uta/Sem3Spring/DBMS/project2/4331-5331_Proj2Spring23_team_31/ sx3702@omega.uta.edu:/home/s/sx/sx3702/tmproject2/

SSH to omega and verify.
-cd /home/s/sx/sx3702/tmproject2/src
- vi zgt_semaphore.C // uncomment union semun ZGT_arg
- chmod +x zgt_test
- make clean && make # compiles the code
- ./zgt_test ../more-test-files/Scenario1.txt # run test cases.
- vi Scenario5.log
- vi ../more-test-files/Scenario5.txt

Where you encountered difficulty

- When building the code locally we faced an error due to the redefinition of semaphore - zgt_semaphore.C:25:7: error: redefinition of 'semun'. The variable was already defined in the <sys/sem.h> header file which caused an error.
- It was difficult to jump into a low-level programming language like C and most of the time was spent on understanding the syntax and semantics. It would help if students were given some references for C to get started with the language. But it was a great learning exercise for us personally.

Overview and file descriptions:

In this project, we were asked to implement one of the lower-level components of DBMS – the Transaction Manager which will be called by higher layers to handle concurrency control. For this purpose, we have implemented Strict Two-Phase Locking (S2PL) protocol where an Exclusive lock is assigned for Write operations and a Shared lock is assigned for Read operations. Hence the Transaction Manager handles Locking and Releasing objects as and when necessary.

The Transaction Manager has been implemented using zgt_tx.c and zgt_tm.c programs.

Transaction manager (zgt_tm.c) manages transactions (zgt_tx) in a multi-threaded environment using the Pthreads library and ensures proper synchronization using mutexes and condition variables.
The code is written in C++ and includes several functions for creating and processing transactions, performing read and write operations, and handling committing or aborting transactions as needed.

Data Structures:

struct param: This structure contains the information required to execute a transaction operation, such as the transaction ID (tid), object number (obno), transaction type (Txtype), and sequence number (count).

<u>Transaction Manager (zgt_tm.c) Functions:</u>

- TxRead: Creates a new thread to perform a read operation in a transaction by initializing the param structure and calling the readtx function within the newly created thread.

- TxWrite: Creates a new thread to perform a write operation in a transaction by initializing the param structure and calling the writetx function within the newly created thread.

- CommitTx: Creates a new thread to commit a transaction by initializing the param structure and calling the committx function within the newly created thread.

- AbortTx: Creates a new thread to abort a transaction by initializing the param structure and calling the aborttx function within the newly created thread.

Transaction (zgt_tx.c) Functions:

- readtx: Receives a struct param as an argument, containing the transaction ID, object number, and count. Calls the process_read_write() function with the transaction ID, object number, count, and a Shared lock mode ('S').

- writetx: Like readtx(), receive a struct param as an argument. Calls the process_read_write() function with the transaction ID, object number, count, and an Exclusive lock mode ('X').

- process_read_write():
    - Starts the operation with the start_operation() function.
    - Sets the lock on the object as an exclusive lock considering the necessary conditions by calling the set_lock() function.
    - Finishes the operation with the finish_operation() function.
    - Exits the thread with pthread_exit(NULL).

- aborttx():
    - Receives a struct param as an argument, containing the transaction ID and count.
    - Starts the operation with the start_operation() function.
    - Calls the do_commit_abort() function for aborting the transaction with status 'A'.
    - Finishes the operation with the finish_operation() function.
    - Exits the thread with pthread_exit(NULL).

- committx():

- Receives a struct param as an argument, containing the transaction ID and count.
- Starts the operation with the start_operation() function.
- Calls the do_commit_abort() function for committing the transaction with status 'C'.
- Finishes the operation with the finish_operation() function.
- Exits the thread with pthread_exit(NULL).

- do_commit_abort():
  - Writes a log record for the transaction (commit or abort).
  - Retrieves the transaction using the transaction ID.
  - If the transaction exists, it releases the locks held by the transaction, and either end or removes the transaction based on the status.
  - If there are transactions waiting on the semaphore, it releases them using zgt_v().

- remove_tx():
  - Removes the transaction from the transaction manager.
  - Scans through the list of transactions and updates the nextr value accordingly.
  - If the transaction does not exist, log an error message.

- set_lock():
  - Sets a lock on an object with a specific lock mode (Shared or Exclusive) for a transaction.
  - Considers conditions such as current transaction ownership, lock modes, and waiting transactions.
  - If successful, returns 0, otherwise, return -1.

Main Difference between commit and abort:

- committx finalizes and saves the changes made during a transaction, making them permanent.
- aborttx discards the changes made during a transaction and restores the system to its state before the transaction started.

Logical errors and how we handled them:

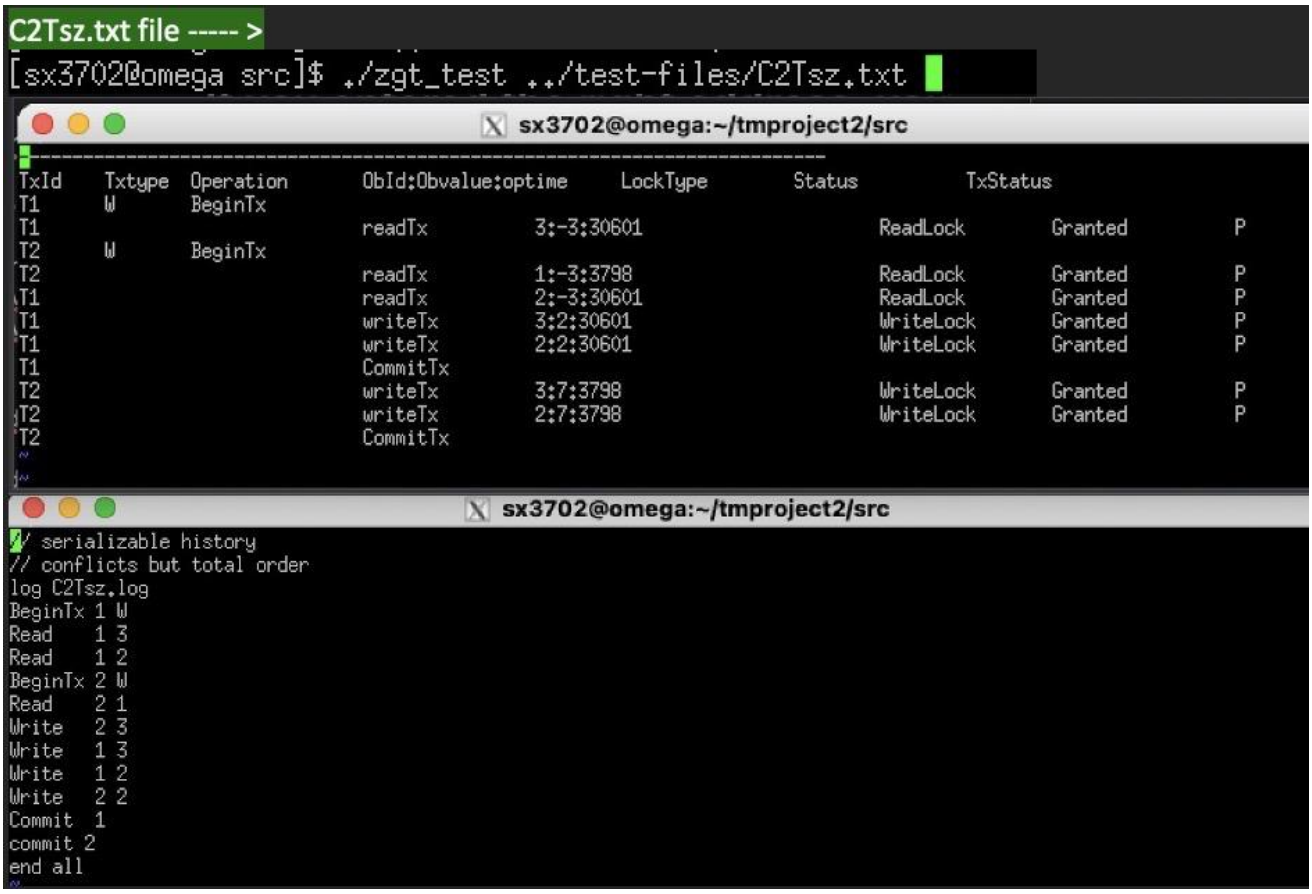
- Deadlock prevention: When a deadlock scenario occurs when two threads try to read/write the same object, the resolution comes by using the semaphore mechanism, which ensures that a transaction waits for a lock only if it can obtain it, preventing deadlocks from occurring.
- Incorrect or non-existent transaction IDs: To avoid a bad (illogical) test scenario, The code checks whether a transaction exists before performing any operation on it (e.g., in the do_commit_abort() function). If a transaction does not exist, it logs an error message and does not proceed with the operation.
- Ensuring lock acquisition: The code ensures that a transaction acquires the appropriate lock (shared or exclusive) before performing read or write operations. This is done in the zgt_tx:set_lock() function, where the lock is acquired if it does not already exist, or if the current transaction already holds the lock.
- Maintaining transaction states: For easier debugging and ease of management, the code keeps track of the state of a transaction, such as active, waiting, or committed. This ensures that operations are performed only when the transaction is in the correct state.
- Proper clean-up after commit or abort: The code makes sure that the locks are released, and the transaction is removed from the transaction manager after a commit or abort operation. This is done in the zgt_tx:free_locks(), zgt_tx:end_tx(), and zgt_tx:remove_tx() functions.
- Ensuring transaction order: The code prevents two operations of the same transaction from following one another by using a sequence number (SEQNUM[tid]). This ensures that the correct operation order is maintained for each transaction.
- Local build errors: Faced some initial local build errors which were resolved after using guards (**reference**). Eventually, we removed them since the error did not show up in the OMEGA server and we were not allowed to touch header files.

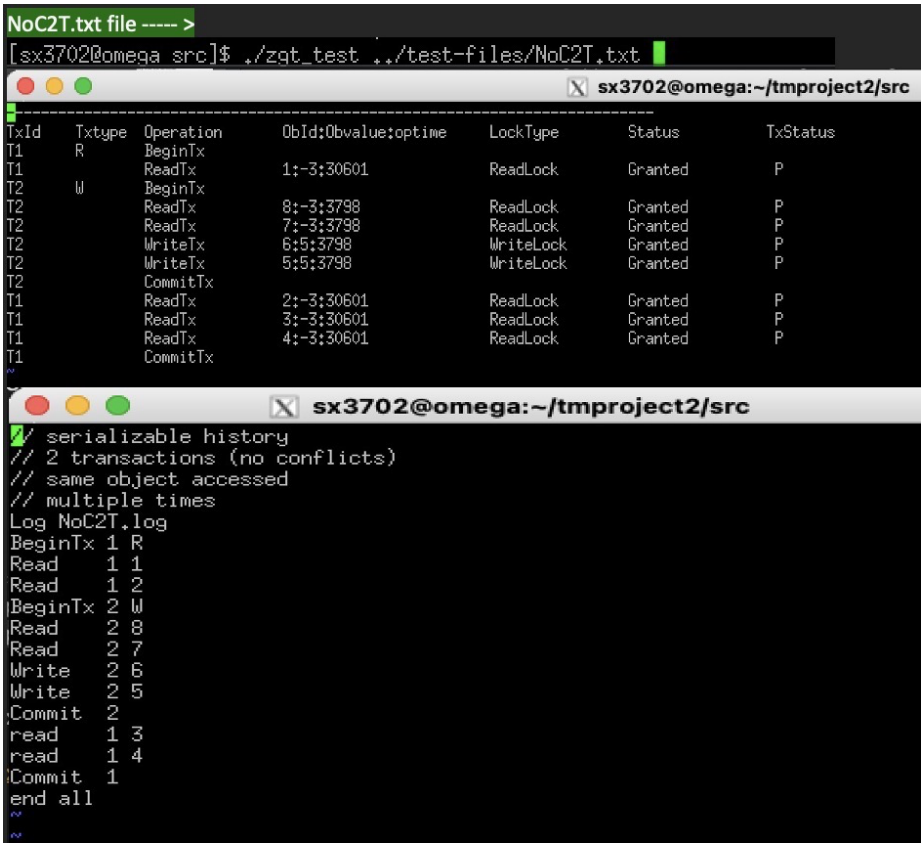
Future Improvements

- Deadlock detection and resolution: Enhance the system's deadlock detection and resolution capabilities. Consider implementing a more advanced algorithm, such as the Wait-for-Graph (WFG) or the edge-chasing algorithm, to efficiently detect and resolve deadlocks in the system.
- Distributed transaction support: Extend the system to support distributed transactions across multiple nodes or clusters. Implement a distributed transaction coordinator to manage transactions that span multiple nodes and ensure atomicity, consistency, isolation, and durability (ACID) properties are maintained.
- Object versioning: Introduce object versioning to enable Multi-Version Concurrency Control (MVCC). This can help reduce contention by allowing transactions to work with different versions of objects without needing to lock them.
- Enhanced logging and recovery: Improve the logging and recovery mechanisms by implementing techniques like Write-Ahead Logging (WAL) or the ARIES recovery algorithm. These can help ensure faster and more efficient recovery after a system crash or failure.
- Monitoring and diagnostics: Develop a comprehensive monitoring and diagnostics framework that provides insights into the system's performance, resource usage, and transaction statistics. This can help identify potential bottlenecks, optimize performance, and aid in troubleshooting issues.
- Transaction prioritization: Implement transaction prioritization, allowing certain transactions to be prioritized over others based on criteria like age, resource requirements, or user-defined priorities. This can help improve the responsiveness of the system for critical transactions and reduce the impact of long-running or resource-intensive transactions.
- Security enhancements: Integrate security features such as authentication, authorization, and encryption to protect sensitive data and ensure that only authorized users can perform transactions. Implementing role-based access control (RBAC) can also help manage user permissions more effectively.

Test on omega Screenshots:

Sirtestfiles-given. ----------




NoC2T.txt file ----- >
[sx3702@omega src]$ ./zgt_test ../test-files/NoC2T.txt
sx3702@omega:~/tmproject2/src
TxId Txtype Operation ObId:Obvalue:optime LockType Status TxStatus
T1 R BeginTx
T1 ReadTx 1:-3:30601 ReadLock Granted P
T2 W BeginTx
T2 ReadTx 8:-3:3798 ReadLock Granted P
T2 ReadTx 7:-3:3798 ReadLock Granted P
T2 WriteTx 6:5:3798 WriteLock Granted P
T2 WriteTx 5:5:3798 WriteLock Granted P
T2 CommitTx
T1 ReadTx 2:-3:30601 ReadLock Granted P
T1 ReadTx 3:-3:30601 ReadLock Granted P
T1 ReadTx 4:-3:30601 ReadLock Granted P
T1 CommitTx
sx3702@omega:~/tmproject2/src
// serializable history
// 2 transactions (no conflicts)
// same object accessed
// multiple times
Log NoC2T.log
BeginTx 1 R
Read 1 1
Read 1 2
BeginTx 2 W
Read 2 8
Read 2 7
Write 2 6
Write 2 5
Commit 2
read 1 3
read 1 4
Commit 1
end all

Deadlock.txt file

```
[sx3702@omega src]$ ./zgt_test ../test-files/deadlock.txt

creating TxRead thread for Tx: 1

exiting TxRead thread create for Tx: 1
commit 2
Commit : 1


creating TxRead thread for Tx: 2

exiting TxRead thread create for Tx: 2
end all
Release all resources and exit:


Entering End of schedule thread with thrNum: 8
Wait for threads and cleanup
Thread 0 completed with ret value: 0
Thread 1 completed with ret value: 0
Thread 2 completed with ret value: 0
Thread 3 completed with ret value: 0
```

sx3702@omega:~/tmproject2/src

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T2	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2		ReadTx	2:-3:3798	ReadLock	Granted	P

```
// 2 transactions
// generates a deadlock
// will hang w/o deadlock resolution
Log deadlock.log
BeginTx 1 W
BeginTx 2 W
Read    1 1
Read    2 2
Write   1 2
Write   2 1
Commit  1
commit 2
end all
```

RR.txt file

```
[sx3702@omega src]$ ./zgt_test ../test-files/RR.txt
+-------------------------------------------------------------------------------
TxId    Txtype  Operation       ObId:Obvalue:optime     LockType        Status          TxStatus
T1      R       BeginTx
T1              ReadTx          1:-3:30601              ReadLock        Granted          P
T2      R       BeginTx
T2              ReadTx          1:-6:3798               ReadLock        Granted          P
T2              ReadTx          2:-3:3798               ReadLock        Granted          P
T2              ReadTx          3:-3:3798               ReadLock        Granted          P
T2              CommitTx
T1              ReadTx          3:-6:30601              ReadLock        Granted          P
T1              ReadTx          2:-6:30601              ReadLock        Granted          P
T1              CommitTx
~
~
```

```
// read read history
// 2 transactions
// same object accessed
// multiple times
Log RR.log
BeginTx 1 R
Read    1 1
Read    1 3
begintx 2 R
read    2 1
Read    2 2
read    1 2
Read    2 3
commit 1
commit  2
end all
~
```

RRW.txt file

```
[sx3702@omega src]$ ./zgt_test ../test-files/RR.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	R	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T3	R	BeginTx				
T1		CommitTx				
T2		WriteTx	1:2:3798	WriteLock	Granted	P
T3		ReadTx	1:-1:19913	ReadLock	Granted	P
T2		CommitTx				
T3		CommitTx				

```
// read read write history
// 3 transactions
// same object accessed
// multiple times
Log RRW.log
BeginTx 1 R
Read    1 1
BeginTx 2 W
Write    2 1
BeginTx 3 R
Read    3 1
commit 1
commit 3
commit 2
end all
```

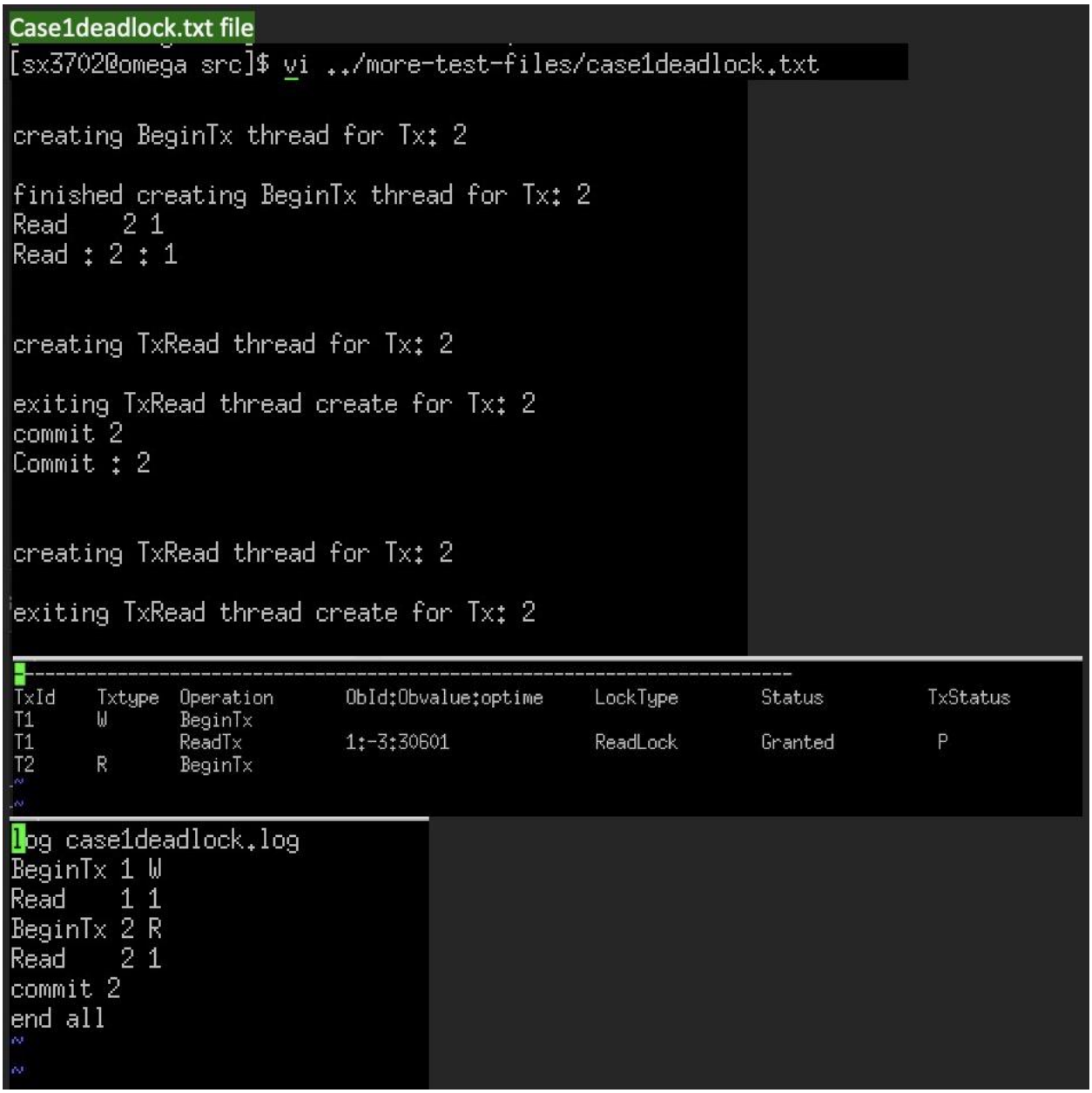
S2T.txt file

```
[sx3702@omega src]$ ./zgt_test ../test-files/S2T.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T2		ReadTx	5:-3:3798	ReadLock	Granted	P
T2		WriteTx	5:2:3798	WriteLock	Granted	P
T2		WriteTx	6:5:3798	WriteLock	Granted	P
T2		ReadTx	6:2:3798	ReadLock	Granted	P
T2		CommitTx				
T1		ReadTx	2:-3:30601	ReadLock	Granted	P
T1		WriteTx	3:5:30601	WriteLock	Granted	P
T1		WriteTx	4:5:30601	WriteLock	Granted	P
T1		ReadTx	1:-6:30601	ReadLock	Granted	P
T1		WriteTx	2:2:30601	WriteLock	Granted	P
T1		WriteTx	4:10:30601	WriteLock	Granted	P
T1		WriteTx	4:15:30601	WriteLock	Granted	P
T1		CommitTx				

```
// serial history
// 2 transactions
// same object accessed
// multiple times
Log S2T.log
BeginTx 1 W
Read    1 1
Read    1 2
Write   1 3
Write   1 4
read    1 1
write   1 2
write   1 4
write   1 4
commit 1
begintx 2 W
read    2 5
write   2 5
write   2 6
read    2 6
commit  2
end all
```

New test cases added.



RW_Disjoint.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/RW_disjoint.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T2		WriteTx	4:5:3798	WriteLock	Granted	P
T2		WriteTx	5:5:3798	WriteLock	Granted	P
T3	W	BeginTx				
T3		WriteTx	6:5:19913	WriteLock	Granted	P
T2		CommitTx				
T5	R	BeginTx				
T5		ReadTx	9:-3:16916	ReadLock	Granted	P
T3		WriteTx	7:5:19913	WriteLock	Granted	P
T5		ReadTx	10:-3:16916	ReadLock	Granted	P
T3		ReadTx	8:-3:19913	ReadLock	Granted	P
T5		ReadTx	11:-3:16916	ReadLock	Granted	P
T1		WriteTx	2:5:30601	WriteLock	Granted	P
T3		CommitTx				
T5		ReadTx	12:-3:16916	ReadLock	Granted	P
T5		ReadTx	13:-3:16916	ReadLock	Granted	P
T1		ReadTx	3:-3:30601	ReadLock	Granted	P
T1		CommitTx				
T5		ReadTx	1:-6:16916	ReadLock	Granted	P
T5		CommitTx				

```
// Multiple RW Txs test case with no deadlock
log RW_disjoint.log
// op   Tx#     type
// op   Tx#     Obj
BeginTx 1 W
Read 1 1
Write 1 2
Read 1 3
BeginTx 2 W
Write 2 4
Write 2 5
BeginTx 3 W
Write 3 6
Write 3 7
read 3 8
Commit 3
commit 2
Commit 1
begintx 5 R
read 5 9
read 5 10
read 5 11
read 5 12
read 5 13
read 5 1
commit 5
end all
```

disj_multi_accesses.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/disj_multi_accesses.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T2		ReadTx	5:-3:3798	ReadLock	Granted	P
T2		WriteTx	5:2:3798	WriteLock	Granted	P
T2		WriteTx	6:5:3798	WriteLock	Granted	P
T2		ReadTx	6:2:3798	ReadLock	Granted	P
T2		CommitTx				
T1		ReadTx	2:-3:30601	ReadLock	Granted	P
T1		WriteTx	3:5:30601	WriteLock	Granted	P
T1		WriteTx	4:5:30601	WriteLock	Granted	P
T1		ReadTx	1:-6:30601	ReadLock	Granted	P
T1		WriteTx	2:2:30601	WriteLock	Granted	P
T1		WriteTx	4:10:30601	WriteLock	Granted	P
T1		WriteTx	4:15:30601	WriteLock	Granted	P
T1		CommitTx				

```
// serial history
// 2 transactions
// same disjoint objects accessed
// multiple times
Log disj_multi_accesses.log
BeginTx 1 W
Read    1 1
Read    1 2
Write   1 3
Write   1 4
read    1 1
write   1 2
write   1 4
write   1 4
commit 1
begintx 2 W
read    2 5
write   2 5
write   2 6
read    2 6
commit  2
end all
```

Ddlk_3Tx file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/ddlk_3Txs.txt
--------------------------------------------------------------------------------
TxId    Txtype  Operation       ObId:Obvalue:optime     LockType        Status          TxStatus
T1      W       BeginTx
T1              ReadTx          1:-3:30601              ReadLock        Granted          P
T2      W       BeginTx
T2              ReadTx          2:-3:3798               ReadLock        Granted          P
T3      R       BeginTx
~
~
~
```

```
// possible deadlock test case
// Two write transactions
log ddlk_3Tx.log
// op   Tx#     type
BeginTx 1 W
// op   Tx#     Obj
Read 1 1
Write 1 2
Read 1 6
BeginTx 2 W
Read 2 2
Write 2 1
Read 2 7
commit 2
Commit 1
begintx 3 R
read 3 2
write 3 1
read 3 2
end all
~
~
```

RW_pot_ddlk.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/RW_pot_ddlk.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T2		WriteTx	4:5:3798	WriteLock	Granted	P
T2		WriteTx	5:5:3798	WriteLock	Granted	P
T3	W	BeginTx				
T3		WriteTx	6:5:19913	WriteLock	Granted	P
T2		CommitTx				
T5	R	BeginTx				
T3		WriteTx	7:5:19913	WriteLock	Granted	P
T3		ReadTx	9:-3:19913	ReadLock	Granted	P
T1		WriteTx	2:5:30601	WriteLock	Granted	P
T3		CommitTx				
T1		ReadTx	3:-3:30601	ReadLock	Granted	P
T1		WriteTx	8:5:30601	WriteLock	Granted	P
T1		CommitTx				
T5		ReadTx	1:-6:16916	ReadLock	Granted	P
T5		ReadTx	2:2:16916	ReadLock	Granted	P
T5		ReadTx	3:-6:16916	ReadLock	Granted	P
T5		ReadTx	8:2:16916	ReadLock	Granted	P
T5		ReadTx	6:2:16916	ReadLock	Granted	P
T5		ReadTx	7:2:16916	ReadLock	Granted	P
T5		CommitTx				

sx3702@omeg

```
// Multiple RW Txs test case with no deadlock
log RW_pot_ddlk.log
// op   Tx#     type
// op   Tx#     Obj
BeginTx 1 W
Read 1 1
Write 1 2
Read 1 3
Write 1 8
BeginTx 2 W
Write 2 4
Write 2 5
BeginTx 3 W
Write 3 6
Write 3 7
Read 3 9
Commit 3
commit 2
Commit 1
begintx 5 R
read 5 1
read 5 2
read 5 3
read 5 8
read 5 6
read 5 7
commit 5
end all
```

test_abort.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/test_abort.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	6:-3:30601	ReadLock	Granted	P
T2	W	BeginTx				
T3	R	BeginTx				
T3		ReadTx	4:-3:19913	ReadLock	Granted	P
T2		ReadTx	8:-3:3798	ReadLock	Granted	P
T2		WriteTx	7:5:3798	WriteLock	Granted	P
T2		AbortTx				
T3		WriteTx	5:5:19913	WriteLock	Granted	P
T3		ReadTx	9:-3:19913	ReadLock	Granted	P
T1		WriteTx	7:10:30601	WriteLock	Granted	P
T3		CommitTx				
T1		WriteTx	7:15:30601	WriteLock	Granted	P
T1		ReadTx	6:-6:30601	ReadLock	Granted	P
T1		CommitTx				

```
// simple deadlock test case
// Two write transactions
log test_abort.log
// op   Tx#     type
// op   Tx#     Obj
BeginTx 1 W
read 1 6
write 1 7
write 1 7
read 1 6
begintx 2 W
read 2 8
write 2 7
abort 2
begintx 3 R
read 3 4
write 3 5
read 3 9
commit 3
commit 1
end all
```

Scenario1.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/Scenario1.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	R	BeginTx				
T1		CommitTx				
T2		ReadTx	1:-6:3798	ReadLock	Granted	P
T2		CommitTx				

```
log Scenario1.log
BeginTx 1 W
Read    1 1
BeginTx 2 R
Read    2 1
Commit  1
commit 2
end all
//readlock on object 1 by T2 must wait for readlock on 1 by T1 to end
```

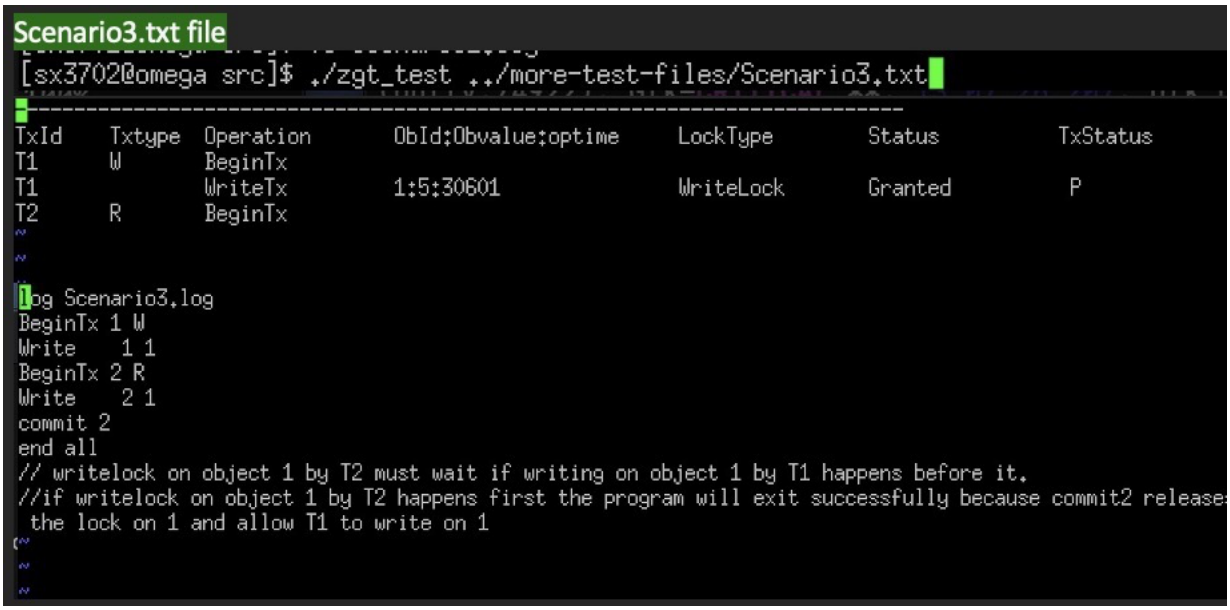
Scenario2.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/Scenario2.txt
```

TxId	Txtype	Operation	ObId:Obvalue:optime	LockType	Status	TxStatus
T1	W	BeginTx				
T1		ReadTx	1:-3:30601	ReadLock	Granted	P
T2	R	BeginTx				

```
log Scenario2.log
BeginTx 1 W
Read    1 1
BeginTx 2 R
Write   2 1
commit 2
end all
// writelock on object 1 by T2 must wait for the readlock by T1 on 1 to end
```

Scenario3.txt file



Scenario4.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/Scenario4.txt
+-------------------------------------------------------------------------
TxId    Txtype  Operation       ObId:Obvalue:optime     LockType        Status          TxStatus
T1      R       BeginTx
T1              ReadTx          1:-3:30601              ReadLock        Granted          P
T2      R       BeginTx
T2              ReadTx          1:-6:3798               ReadLock        Granted          P
T2              CommitTx
~
~
log Scenario4.log
BeginTx 1 R
Read    1 1
BeginTx 2 R
Read    2 1
commit 2
end all
//beginTx 1 R begins transaction 1 where all the reads have shared locks and all the write has exclusive l
ocks.
~
~
```

Scenario5.txt file

```
[sx3702@omega src]$ ./zgt_test ../more-test-files/Scenario5.txt
+-------------------------------------------------------------------------
TxId    Txtype  Operation       ObId:Obvalue:optime     LockType        Status          TxStatus
T1      W       BeginTx
T1              ReadTx          1:-3:30601              ReadLock        Granted          P
T2      R       BeginTx
~
log Scenario5.log
BeginTx 1 W
Read    1 1
BeginTx 2 R
Read    2 1
commit 2
end all
//beginTx 1 W begins T1 where all the reads have exclusive locks, in this case, reading on 1 by T1 is excl
usive Hence, T2 must wait before it can read 1
~
~
~
```