



Today's agenda

- ↳ understanding sorting
- ↳ Problems on sorting
- ↳ Sorting techniques



AlgoPrep



Sorting: Arranging data in increasing/decreasing.

↳ on the basis of which parameter

Ex1: 2 4 10 15 27 → true

Ex2: 20 7 7 -5 -8 → true

Ex3: 1 2 3 7 4 9 6

*factors: 1 2 2 2 3 3 4

→ not sorted on the basis val.

→ sorted on the basis of factor count.

→ How to sort an array

↳ bubble sort ↴

↳ selection sort

↳ insertion sort

↳ quick sort

↳ merge sort

↳ :

:

:

dec order

Arrays.sort(arr, reverseOrder());

arr[7]: 1 2 3 7 4 9 6

↓
Arrays.sort(arr);
inc. order

↓
T.C: O(n log n)

1 2 2 4 6 7 9



Q) Order of Removal

Given N elements at every step remove an array element. Cost to remove element = Sum of array elements
Problem. Find min cost to remove all elements.

Note: Add cost first and then remove.

$$\text{Ex1: } \text{arr}[3] = \{3 \ 2 \ 5\}$$

$\begin{matrix} x & x & x \\ 0 & 1 & 2 \end{matrix}$

remove 2: 10

remove 5: 8

remove 3: 3

21

$$\text{arr}[3] = \{3 \ 2 \ 5\}$$

$\begin{matrix} x & x & x \\ 0 & 1 & 2 \end{matrix}$

remove 2: 10

remove 3: 8

remove 5: 5

23

$$\text{Ex2: } \text{arr}[4] = \{ \cancel{4} \ \cancel{3} \ \cancel{2} \ \cancel{1} \}$$

remove 7: 4 + 6 + 2 + 7

remove 2: 4 + 6 + 2

remove 6: 4 + 6

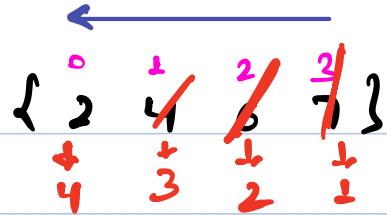
remove 4: 4

max Contain' in total cost = 0th index = min ele

2nd max Contain' in total cost = 1st index = 2nd min ele

2nd index = 3rd min ele.

→ Array should be in inc. order.

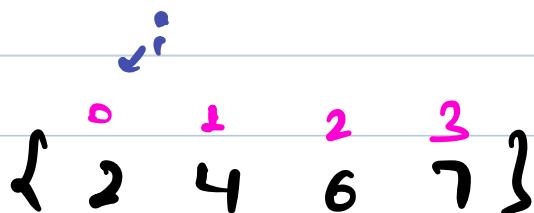


$$\text{ans} += (\text{Removal freq}) = (N - i)$$

remove 7: $2 + 4 + 6 + 7$

remove 6: $2 + 4 + 6$

remove 4: $2 + 4$



$$\text{ans} = 0 + (7 * 1) + (6 * 2) + (4 * 3) + (2 * 4)$$

// Pseudo code

```
int removalOrder (int arr[N]) {
```

`Arrays.sort (arr);`

T.C: $O(n \log n) + O(n)$
 $\approx O(n \log n)$

`int ans = 0;`

`for (int i = N-1; i >= 0; i--) {`

`ans = ans + [arr[i] * (N-i)];`

`}`

`return ans;`



```
int removalOrder (int arr[n]) {
```

```
    Arrays.sort (arr);
```

```
    int ans = 0;
```

```
    for (int i = n - 1; i >= 0; i--) {
```

```
        ans = ans + [arr[i] * (n - i)];
```

```
}
```

```
    return ans;
```

↑
0 1 2 3
{ 2 4 6 7 }

$$ans = 0 + 7 + 12 + 12 + 8$$

remove 7 : 2 + 4 + 6 + 7

remove 6 : 2 + 4 + 6

remove 4 : 2 + 4

remove 2 : 2



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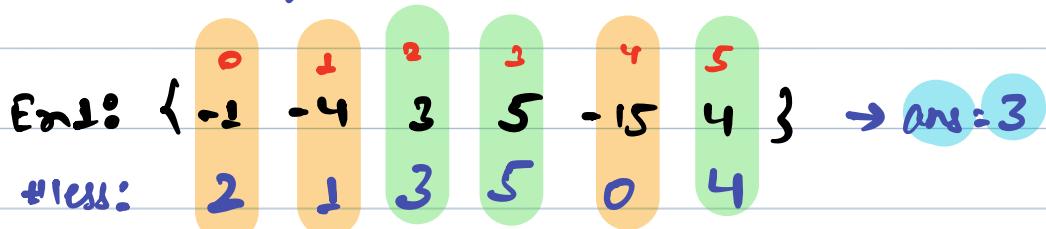


Q) Good Integer { only distinct }

Given $arr[n]$, calculate no. of good integers.

An element is said to be good if

{No. of element $< ele == ele$ itself}



Idea 1

for every number, count smaller no. and check
if count \leq no.

T.C: $O(n^2)$

Idea 2

$$\{ -2, -4, 3, 5, -15, 4 \}$$



$$\{ -15, -4, -2, 3, 4, 5 \}$$

$\frac{1}{1}, \frac{2}{2}, \frac{5}{5}$

After sorting,

;

$arr[i]$

{No. of element $< ele == ele$ itself}



II Pseudo code

```
int goodInteger (int arr[n]) {  
    Arrays.sort (arr);
```

```
    int ans = 0;
```

```
    for (int i=0; i<n; i++) {  
        if (arr[i] == i) {  
            ans++;  
        }  
    }  
    return ans;
```

}

Break till 10: 29 pm

{No. of element < ele == ele itself}



Good Integers : { Data can repeat }

Ex1: { 0 1 2 2 3 3 3 5 } → ans = 3
Hence: 0 1 1 3 3 5

After sorting

Ex2: { -4 -4 2 2 4 4 4 4 8 8 8 8 10 11 12 }
0 0 2 2 4 4 4 4 8 8 8 8 10 11 12

Obs1: All the same elements will be either good or bad (count of no. < ele will remain same)

Obs2:

if ele is 1st occ. → arr[i] != arr[i-1]

No. of ele < ele == ?



II Pseudo Code

```
int goodintegerduplicate (int arr[n]) {
```

Arrays sort (arr);

int ans=0;

→ //Handle for 0th index → H.W

int lessCount=0;

```
for (int i=1; i<n; i++) {
```

if (arr[i] == arr[i-1]) { "first occ.

lessCount = i;

else {

if (arr[i] == lessCount) { ans++;

}

}

T.C: O(nlogn)

S.C: O(1)

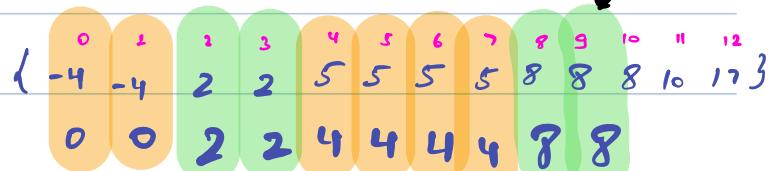
ans = 628884



```
int goodintegerduplicate (int arr[n]) {
```

 Arrays sort (arr);

 int ans=0;



 freqCount = 62888

 int lessCount=0;

```
    for (int i=1; i<n; i++) {  
        if (arr[i] == arr[i-1]) { //first occ.  
            lessCount = i;  
        }  
        else { //next occ.  
            j  
        }  
        if (arr[i] == lessCount) { ans++; }  
    }
```

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11 Sorting techniques

① Bubble Sort

↳ Sort the array in asc. order but we can swap

adjacent elements only.

$$arr[8] = \{ 5 \ 7 \ 5 \ 4 \ 10 \ -2 \ 6 \ 3 \}$$

iter 0: i $(j, j+1)$ j

$$\{ 5 \ 7 \ 8 \ 1 \ 10 \ -2 \ 6 \ 3 \} \rightarrow \{ 0, \sim 2 \}$$

iter 1: i j

$$\{ 5 \ 8 \ 7 \ 1 \ 4 \ 10 \ -2 \ 6 \ 3 \} \rightarrow \{ 0, \sim 3 \}$$

iter 2: i j

$$\{ 8 \ 5 \ 4 \ 1 \ 7 \ 10 \ 6 \ 3 \} \rightarrow \{ 0, \sim 4 \}$$

Total no. of iter = $7 = \sim 1$

$$[0, \sim 2] / [1, \sim 1]$$

iter



II Pseudo code

```
void bubblesort (int arr[n]) {
```

```
    for (int i = 0; i <= n - 1; i++) {
```

```
        for (int j = 0; j <= n - 2 - i; j++) {
```

```
            if (arr[j] > arr[j + 1]) {
```

```
                int temp = arr[j];
```

```
                arr[j] = arr[j + 1];
```

```
                arr[j + 1] = temp;
```

T.C: $O(n^2)$

S.C: $O(1)$



}

}

}

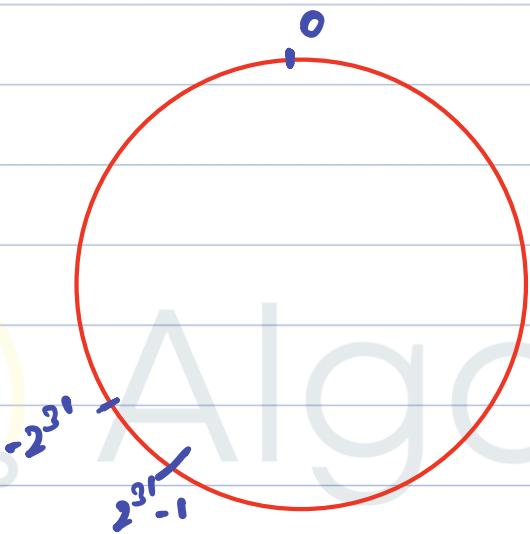
}

}



$$\text{int} \rightarrow -2^{31} \quad 2^{31}-1$$

$$-(-2^{31}) \rightarrow 2^{31} = -2^{31}$$



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