



Today's agenda

- ↳ understanding sorting
- ↳ Problems on sorting
- ↳ Sorting techniques



AlgoPrep



Sorting: Arranging data in increasing/decreasing.

↳ on the basis of which parameter

Ex1: 2 4 10 15 27 → true

Ex2: 20 7 7 -5 -8 → true

Ex3: 1 2 3 7 4 9 6

#factors:

1 2 2 2 3 3 4

Not sorted on the basis val.

Sorted on the basis of factor count.

→ How to sort an array

↳ bubble sort

↳ selection sort

↳ insertion sort

↳ Quick sort

↳ merge sort

↳ :

:

arr[]: 1 2 3 7 4 9 6

Arrays.sort(arr);

↳ T.C: $O(n \log n)$

1 2 2 4 6 7 9

dec order

Arrays.sort(arr, reverseOrder());



Q) Order of Removal

↳ Given n elements at every step remove an array element. Cost to remove element = Sum of array elements Present. Find min cost to remove all elements.

Note: Add cost first and then remove.

Ex1: $arr[3] = \{3 \ 2 \ 5\}$

remove 2: 10

remove 5: 8

remove 3: 3

21

$arr[3] = \{3 \ 2 \ 5\}$

remove 2: 10

remove 3: 8

remove 5: 5

23

Ex2: $arr[4] = \{ \overset{0}{\cancel{4}} \overset{1}{\cancel{6}} \overset{2}{\cancel{2}} \overset{3}{\cancel{7}} \}$

remove 7: $4 + 6 + 2 + 7$

remove 2: $4 + 6 + 2$

remove 6: $4 + 6$

remove 4: 4

max Contribⁿ in total cost = 0th index = min ele

2nd max Contribⁿ in total cost = 1st index = 2nd min ele

2nd index = 3rd min ele.

→ Array should be in inc. order.



←
 $\{ \overset{0}{2} \overset{1}{4} \overset{2}{6} \overset{3}{7} \}$
 $\begin{matrix} +4 \\ +3 \\ +2 \\ +1 \end{matrix}$

$$ans += (\text{removal cost}) = (N - i)$$

remove 7: $2 + 4 + 6 + 7$

remove 6: $2 + 4 + 6$

remove 4: $2 + 4$

✓
 $\{ \overset{0}{2} \overset{1}{4} \overset{2}{6} \overset{3}{7} \}$

$$ans = 0 + (7 * 1) + (6 * 2) + (4 * 3) + (2 * 4)$$

// Pseudo code

```
int removalOrder (int arr[N]) {
```

```
    Arrays.sort (arr);
```

```
    int ans = 0;
```

```
    for (int i = N-1; i >= 0; i--) {
```

```
        ans = ans + (arr[i] * (N-i));
```

```
    }
```

```
    return ans;
```

```
}
```

T.C: $O(N \log N) + O(N)$
 $\approx O(N \log N)$

S.C: $O(1)$



```
int removalOrder (int arr[N]) {
```

```
    Arrays.sort (arr);
```

```
    int ans = 0;
```

```
    for (int i = N-1; i >= 0; i--) {
```

```
        ans = ans + (arr[i] * (N-i));
```

```
    }
```

```
    return ans;
```

```
}
```

i

0 1 2 3

{ 2 4 6 7 }

$ans = 0 + 7 + 12 + 12 + 8$

remove 7: 2 + 4 + 6 + 7

remove 6: 2 + 4 + 6

remove 4: 2 + 4

remove 2: 2



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Q) Good Integer (only distinct)

↳ Given $arr[N]$, Calculate no. of good integers.

An element is said to be good if

{No. of element $< arr[i] == arr[i]$ }

Ex: { $\overset{0}{-1}$ $\overset{1}{-4}$ $\overset{2}{3}$ $\overset{3}{5}$ $\overset{4}{-15}$ $\overset{5}{4}$ } $\rightarrow ans = 3$
#less: $\overset{0}{2}$ $\overset{1}{1}$ $\overset{2}{3}$ $\overset{3}{5}$ $\overset{4}{0}$ $\overset{5}{4}$

Idea 1

↳ for every number, count smaller no. and check if count $==$ no.

T.C: $O(N^2)$

Idea 2

{ $\overset{0}{-1}$ $\overset{1}{-4}$ $\overset{2}{3}$ $\overset{3}{5}$ $\overset{4}{-15}$ $\overset{5}{4}$ }

↓

{ $\overset{0}{-15}$ $\overset{1}{-4}$ $\overset{2}{-1}$ $\overset{3}{3}$ $\overset{4}{4}$ $\overset{5}{5}$ }

↓ ↓ ↓

1 2 5

After sorting,

{No. of element $< arr[i] == arr[i]$ }



// Pseudo code

```
int goodInteger (int arr[N] ) {  
    Arrays.sort (arr);  
  
    int ans = 0;  
  
    for (int i = 0; i < N; i++) {  
        if (arr[i] == i) {  
            ans++;  
        }  
    }  
  
    return ans;  
}
```

T.C: $O(N \log N)$
S.C: $O(1)$

Break till 10: 29 PM

{No. of element \leq ele $::$ ele itself}



Good Integers : {Data can repeat}

Ex1: { 0 2 2 3 3 8 } \rightarrow ans = 3
#128: 0 1 1 3 3 5

After sorting

Ex2: { -4 -4 2 2 5 5 5 5 8 8 8 10 12 }
0 0 2 2 4 4 4 4 8 8 8 11 12

obs1: All the same elements will be either good or bad (count of no. \leq ele will remain same)

obs2:

if ele is 1st occ. \rightarrow ans[i] != ans[i-1]

No. of ele \leq ele $::$ i



// Pseudo Code

```
int goodintegerduplicate (int arr[n]) {
```

```
    Arrays Sort (arr);
```

```
    int ans = 0;
```

```
    // Handle for 0th index → H.W
```

```
    int lesscount = 0;
```

T.C: $O(n \log n)$

S.C: $O(1)$

```
    for (int i = 1; i < n; i++) {
```

```
        if (arr[i] != arr[i-1]) { // first occ.
            lesscount = i;
```

```
        }  
        else {
```

```
        if (arr[i] == lesscount) { ans++; }
```

```
    }
```

ans = 0 2 2 3 4



```
int goodIntegerDuplicate (int arr[n]) {
```

0	1	2	3	4	5	6	7	8	9	10	11	12
-4	-4	2	2	5	5	5	5	8	8	8	10	17
0	0	2	2	4	4	4	4	8	8			

Arrays sort (arr);

int ans = 0;

Integer Count = 0 2 2 3 4

int lessCount = 0;

for (int i = 1; i < n; i++) {

if (arr[i] != arr[i-1]) { // first occ.
lessCount = i;

}

~~else { // next occ.~~

}

if (arr[i] == lessCount) { ans++; }

}

}



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// Sorting techniques

① Bubble sort

↳ Sort the array in asc. order but we can swap adjacent elements only.

arr[8] = { 5 7 5 4 10 -2 6 3 }

(j, j+1)

iter 0: { 5 7 5 4 10 -2 6 3 } → { 0, n-2 }

iter 1: { 5 8 4 7 -2 6 3 10 } → { 0, n-3 }

iter 2: { 4 5 -2 5 6 3 7 10 } → { 0, n-4 }

Total no. of iter = 7 = n-1

[0, n-2] / [1, n-1]

iter



11Pseudo code

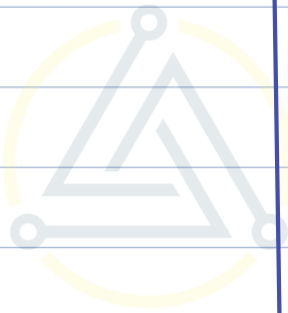
```
void bubblesort (int arr[n]) {
```

```
    for (int i = 0; i <= n-2; i++) {
```

```
        for (int j = 0; j <= n-2-i; j++) {  
            if (arr[j] > arr[j+1]) {  
                int temp = arr[j];  
                arr[j] = arr[j+1];  
                arr[j+1] = temp;  
            }  
        }  
    }
```

T.C: $O(n^2)$

S.C: $O(1)$



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$$in \rightarrow -2^{31} \quad 2^{31}-1$$

$$-(2^{31}) \rightarrow 2^{31} = -2^{31}$$

