



## Today's agenda

→ Priority Queue

↳ Introduction to Heap / PQ.

↳ K smallest element. +1

↳ median of an array. → Hard



# AlgoPrep



## 11 Introduction

insert( $n$ )

getmin()

deletemin()

LinkedList

$O(1) | O(n)$

$O(n)$

$O(n)$

Queue

$O(1)$

$O(n)$

$O(n)$

HashMap

$O(1)$

$O(n)$

$O(n)$

PQ

$O(\log n)$   
size of PQ

$O(1)$

$O(\log n)$

min PQ

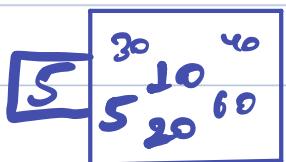
name

max PQ

name

PriorityQueue<Integer> PQ = new

PriorityQueue<?>();



PQ.add(5); → add in Priority Queue.

PQ.add(10);

PQ.add(20);

PQ.add(2);

PQ.peek(); → get the min element

2

PQ.remove(); → remove and return the min ele

2

PQ.add(5);

PQ.size();

↳ no. of elements in PQ.



## (Q) Kth Smallest Element

↳ Given  $n$  distinct elements, Point  $K$  Smallest elements.

Ex:  $\text{arr}[10] = \{ \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ 8 & 3 & 10 & 4 & 11 & 2 & 7 & 6 & 14 & 1 \end{matrix} \}$

$K=4:$  1 2 3 4

$\text{arr}[9] = \{ \begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 \\ -3 & 6 & 2 & 0 & 8 & -3 & 10 & 4 \end{matrix} \}$

$K=3:$  -3 -3 0

### 1/Idea 1

↳ Sort the array and return the first  $K$  elements.

T.C:  $O(n \log n + k) \approx O(n \log n)$

### 1/Idea 2

↳ Add all the elements to min PQ and get the first  $K$  elements.

T.C:  $O(n \log n + k \log n) = O(n \log n)$

S.C:  $O(n)$



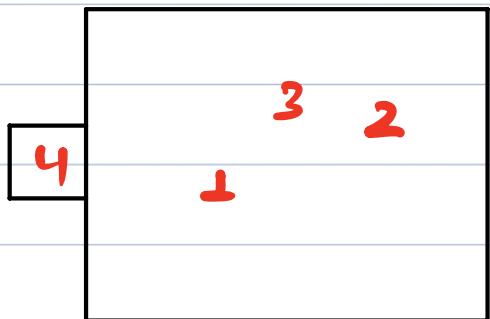
## II Idea 3

$\text{arr}[10] = \{8, 3, 10, 4, 11, 2, 7, 6, 14, 1\}$

$K=4; \quad 4 \quad 3 \quad 2 \quad 1$

$v_i^R$

$\log K \leftarrow$



## II Pseudo code

```
void Kthsmallest ( int arr[N], int k) {
```

```
    maxHeap<Integer> mh;
```

```
    for (int i=0; i<k; i++) {  
        mh.add (arr[i]);
```

3

T.C:  $O(n \log k)$

S.C:  $O(k)$

```
    for (int i=k; i<n; i++) {  
        if (arr[i] < mh.Peek()) {  
            mh.remove();  
            mh.add (arr[i]);
```

3

```
    while (mh.size() > 0) {  
        point (mh.remove());
```

3

3



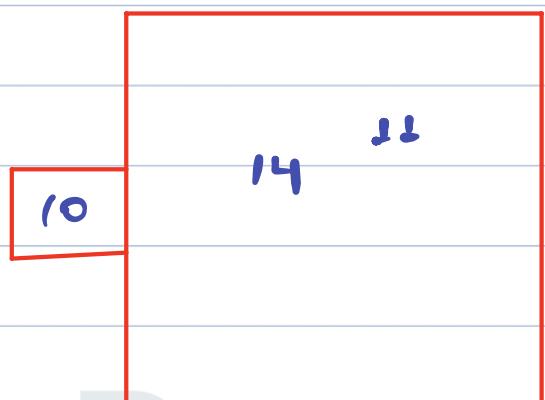
## Q) K largest elements

arr[10]: {<sup>0</sup> 8 <sup>1</sup> 3 <sup>2</sup> 10 <sup>3</sup> 4 <sup>4</sup> 11 <sup>5</sup> 2 <sup>6</sup> 7 <sup>7</sup> 6 <sup>8</sup> 14 <sup>9</sup> 1}

K=3

↳ 10 11 14

✓ min PQ



AlgoPrep



//median

↳ middle element of sorted number.

$$arr[3] = \{2, 5, 3\}$$

$$\hookrightarrow \{2, 3, 5\} \rightarrow 3$$

$$arr[5] = \{4, 3, 6, 8, 5\}$$

$$\hookrightarrow \{3, 4, 5, 6, 8\} \rightarrow 5$$

$$arr[6] = \{4, 3, 9, 5, 12, 2\}$$

$$\hookrightarrow \{2, 3, 4, 5, 9, 12\} \xrightarrow{\frac{4+5}{2}} 4.5$$

$$arr[4] = \{4, 6, 10, 14\}$$

$$\hookrightarrow \frac{6+10}{2} = 8$$

Break till 10:17 pm



Q) Point median after each insertion.

arr[5]: 9 6 3 10 4  
↓ ↓ ↑ ↓ ↓  
3 7.5 6 7.5 6

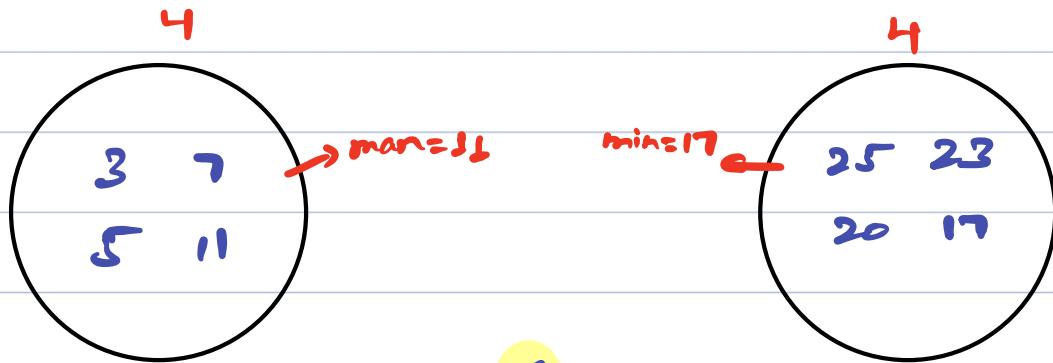
Idea 1

After every insertion, sort the array and return the middle one.

T.C:  $O(n \log n * n) \approx O(n^2 \log n)$

Idea 2

even element 3 5 11 23 20 25 17 7

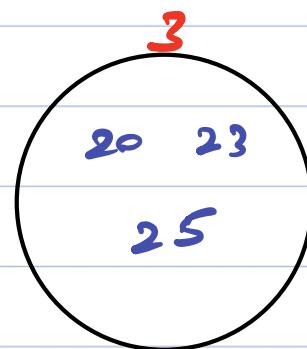
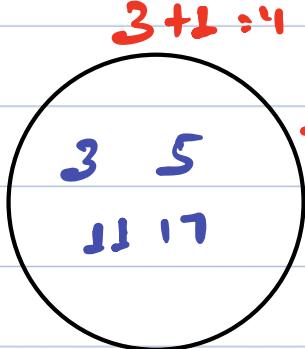


median: man of left bucket + min of right bucket  
2

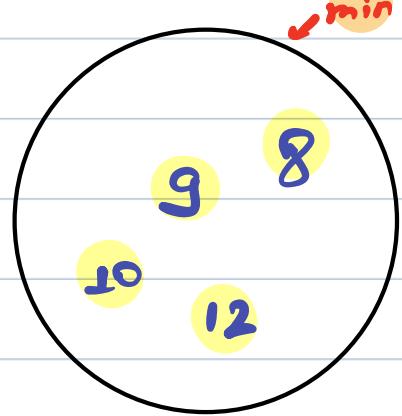
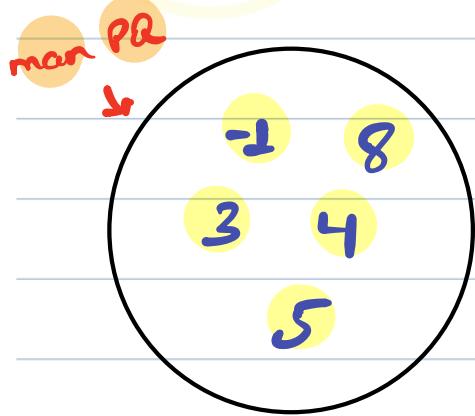


odd element

3 5 11 23 20 25 17



median : man of left



left

right

if (`left.size() == right.size()`) { new element is  
(odd)th element }



## IIIP Suedo Code

```
Class medianFinder {
```

```
    maxHeap<Integer> left;
```

```
    minHeap<Integer> right;
```

```
    Public medianFinder () {
```

```
}
```

```
    Public void addnum (int num) {
```

```
        if (left.size() == right.size()) {
```

```
            right.add (num);
```

```
            left.add (right.remove());
```

```
        } else {
```

```
            left.add (num);
```

```
            right.add (left.remove());
```

```
}
```

```
}
```

```
    Public double findmedian () {
```

```
        if (left.size() == right.size()) {
```

```
            return left.peek() + right.peek() / 2.0;
```

```
        } else {
```

```
            return (left.peek() + 100);
```

```
}
```

T.C:  $O(n \log n)$

S.C:  $O(n)$

$3d \log n$

$O(1)$