



Today's agenda

- ↳ Reverse an array
- ↳ Reverse a given part of array.
- ↳ Rotate array by k.
- ↳ greater than itself.
- ↳ Two Sum.



AlgoPrep



Q) Reverse array

↳ Given array of length N , Reverse the whole array.

ex: arr[5]: {⁰10 ¹20 ²30 ³40 ⁴50}
50 40 30 20 10

arr[8]: {⁰10 ¹20 ²30 ³40 ⁴50 ⁵60 ⁶70 ⁷80}
80 70 60 50 40 30 20 10

arr[8]: {⁰10 ¹20 ²30 ³40 | ⁴50 ⁵60 ⁶70 ⁷80}
80 70 60 50 40 30 20 10

1P 2P
↓ ↓
Swap (0, 7)
↓ ↓
Swap (1, 6)
↓ ↓
Swap (2, 5)
↓ ↓
Swap (3, 4)



//Pseudo code

```
int main() {  
    //input → arr, 2 idn  
    reverse(arr);  
    for(int i=0; i< arr.length; i++) { s.o.p(arr[i]); }  
}
```

```
public static void reverse (int [] arr) {  
    int n = arr.length;
```

```
    int sp = 0;  
    int ep = n-1;
```

```
    while (sp < ep) {
```

```
        int temp = arr[sp];  
        arr[sp] = arr[ep];  
        arr[ep] = temp;  
        sp++;  
        ep--;
```

```
    }
```

```
}
```

T.C: $O(N)$

S.C: $O(1)$



n=7

```
Public static void reverse (int[] arr){
```

```
    int n = arr.length;
```

```
    int sp = 0;
```

```
    int ep = n-1;
```

```
    while (sp < ep) {
```

```
        int temp = arr[sp];
```

```
        arr[sp] = arr[ep];
```

```
        arr[ep] = temp;
```

```
        sp++;
```

```
        ep--;
```

```
    }
```

arr[7]: {⁰10, ¹20, ²30, ³40, ⁴50, ⁵60, ⁶70}

sp

ep

sp < ep

0

6

T

1

5

T

2

4

T

3

3

F
4 iterations



AlgoPrep



Q) Reverse a Part of array

↳ Given N array element and $[s, e]$, reverse the array from $[s, e]$.

$[3, 7]$

arr[10] = {⁰-3 ¹4 ²2 ³8 ⁴3 ⁵9 ⁶6 ⁷2 ⁸8 ⁹10}

-3 4 2 $[2$ 6 9 3 8] 8 10

swap(3, 7)

swap(4, 6)

```
public static void reverse (int [] arr, int s, int e) {
```

```
    int n = arr.length;
```

```
    int sp = x; s;
```

```
    int ep = n-1; e;
```

```
    while (sp < ep) {
```

```
        int temp = arr[sp];
```

```
        arr[sp] = arr[ep];
```

```
        arr[ep] = temp;
```

```
        sp++;
```

```
        ep--;
```

```
    }
```

```
}
```

Worst case
time complexity
↑

T.C: $O(\frac{e-s}{2}) \approx O(\frac{e-s}{2})$ int

S.C: $O(1)$ $O(n)$



Q) Rotate the array \rightarrow {google, meta, amazon}

↳ Given n elements, rotate array from left to right by K times.

$K < n$

$K=3$ arr[7]: { 3 -2 1 4 6 9 8 }

expected r.c: $O(n)$
expected s.c: $O(1)$

↓ 1st rot.
{ 8 3 -2 1 4 6 9 }

↓ 2nd rot.
{ 9 8 3 -2 1 4 6 }

↓ 3rd rot.
{ 6 9 8 3 -2 1 4 }

$K=3$ arr[7]: { 3 -2 1 4 6 9 8 }

↓
{ 6 9 8 3 -2 1 4 }



$k=3$
arr[7]: { 3 -2 1 4 6 9 8 }



Reverse the array



{ 8 9 6 | 4 1 -2 3 }



Reverse the first k elements



{ 6 9 8 | 4 1 -2 3 }



Reverse the elements after kth element



{ 6 9 8 | 3 -2 1 4 }

{ 6 9 8 | 3 -2 1 4 }



// Pseudo code

```
int main() {
```

```
    // Input  $\rightarrow$  arr[N], K
```

```
    K = K % N;
```

```
    // Step 1: Reverse the array
```

```
    reverse(arr, 0, N-1);
```

```
    // Step 2: Reverse the first K elements
```

```
    reverse(arr, 0, K-1);
```

```
    // Step 3: Reverse the after K elements
```

```
    reverse(arr, K, N-1);
```

```
}
```

```
public static void reverse (int[] arr, int s, int e) {
```

```
    int n = arr.length;
```

```
    int sp = s;
```

```
    int ep = e;
```

```
    while (sp < ep) {
```

```
        int temp = arr[sp];
```

```
        arr[sp] = arr[ep];
```

```
        arr[ep] = temp;
```

```
        sp++;
```

```
        ep--;
```

```
}
```

```
}
```

T.C: $O(N)$

S.C: $O(1)$

T.C: $O(N)$

S.C: $O(1)$



$K = 7$

$arr[4] = \{ 4^0, 1^1, 6^2, 9^3 \}$

$\downarrow \text{rot}+1$

9 4 1 6

$\downarrow \text{rot}+2$

6 9 4 1

$\downarrow \text{rot}+3$

1 6 9 4

$\downarrow \text{rot}+4$

4 1 6 9

$\downarrow \text{rot}+5$

$K = 10$

$arr[4] = \{ 4^0, 1^1, 6^2, 9^3 \}$

$\downarrow \text{rot}+1$

9 4 1 6

$\downarrow \text{rot}+2$

6 9 4 1

$\downarrow \text{rot}+3$

1 6 9 4

$\downarrow \text{rot}+4$

4 1 6 9

$\text{rot}+5$

$\text{rot}+6$

$\text{rot}+7$

$\text{rot}+8$

4 1 6 9

effective no. of rotation = $K \% N$



→ non length

N

K

effective rotation

$$22 - 7 = 15 - 7 = 8 - 7 = 1$$

$$22 \div 7 = 3$$

7

31

$$31 \div 7 = 4$$

7

3

$$3 \div 7 = 0$$

Break till 10:50 PM



AlgoPrep



Q) Given N array elements, count total no. of elements having atleast 1 element greater than itself.

ex: arr[7]: {⁰-4 ¹-3 ²7 ³9 ⁴3 ⁵9 ⁶4}
↳ ans = 5

arr[8]: {3 4 11 8 2 10 9 11}
↳ ans = 6

arr[5]: {7 7 7 7 7}
↳ ans = 0

arr[6]: {1 2 2 3 3 4}
↳ ans = 5

Idea

obs1: max elements of the array are not valid.

obs2: except for max elements, all the elements are valid.

→ find max no. & count it. → maxcount

↳ ans = no. of elements - maxcount.



// Pseudo Code

```
int CountGreater (int arr[N]) {  
    int max = arr[0];  
    for (int i = 1; i < N; i++) {  
        if (arr[i] > max) { max = arr[i]; }  
    }  
    int maxcount = 0;  
    for (int i = 0; i < N; i++) {  
        if (arr[i] == max) { maxcount++; }  
    }  
    return N - maxcount;  
}
```

T.C: $O(2N) \approx O(N)$

S.C: $O(1)$



manCount = 0 ~~1~~ 2

```
int man = arr[0];
```

arr[8]: { 3 4 11 8 2 10 9 11 }

```
for (int i = 1; i < N; i++) {  
    if (arr[i] > man) { man = arr[i]; }  
}
```

man = 3 4 11

6

```
int manCount = 0;
```

```
for (int i = 0; i < N; i++) {  
    if (arr[i] == man) { manCount++; }  
}
```



AlgoPrep



Q) Two Sum

↳ Given n array elements, check if there exists a pair (i, j) such that $arr[i] + arr[j] = k$ and $i \neq j$

Note: i and j are index value, k is given sum.

ex: $arr[7] = \{ 2 \ -1 \ 0 \ 3 \ 2 \ 5 \ 7 \}$
 $k = 8$

$arr[4] = \{ 1 \ 3 \ -2 \ 6 \}$
 $k = 5$

$arr[5] = \{ \overset{0}{2} \ \overset{1}{4} \ \overset{2}{-3} \ \overset{3}{7} \ \overset{4}{10} \}$
 $k = 8$

$arr[6] = \{ 3 \ 5 \ 1 \ 8 \ 3 \ 7 \}$
 $k = 6$