



Today's agenda

↳ Reverse an array

↳ Reverse a given Part of array.

↳ Rotate array by K.

↳ greater than itself.

↳ Two Sum.



AlgoPrep



Q) Reverse array

In Given array of length n , Reverse the whole array.

ex: $\text{arr}[5]: \{10 \ 20 \ 30 \ 40 \ 50\}$

0 1 2 3 4
50 40 30 20 10

$\text{arr}[8]: \{10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80\}$

0 1 2 3 4 5 6 7
80 70 60 50 40 30 20 10

$\text{arr}[8]: \{10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80\}$

80 70 60 50 40 30 20 10

$i^P \quad eP$
 $\downarrow \quad \downarrow$
Swap ($0 < 7$)
 $\downarrow \quad \downarrow$
Swap ($1 < 6$)
 $\downarrow \quad \downarrow$
Swap ($2 < 5$)
 $\downarrow \quad \downarrow$
Swap ($3 < 4$)



II Pseudo code

```
int main() {
    // Input → arr[], 2 ida
    reverse(arr);
    for(int i=0; i<arr.length; i++) { System.out.println(arr[i]); }
}
```

3

```
public static void reverse(int[] arr) {
```

```
    int n = arr.length;
```

```
    int SP = 0;
    int EP = n-1;
```

T.C: $O(N)$

S.C: $O(1)$

```
    while (SP < EP) {
```

```
        int temp = arr[SP];
```

```
        arr[SP] = arr[EP];
```

```
        arr[EP] = temp;
```

```
        SP++;
        EP--;
    }
```

3



Public static void reverse (int [] arr) {

int n = arr.length;

int SP = 0;

int EP = n-1;

```
while (SP < EP) {
    int temp = arr[SP];
    arr[SP] = arr[EP];
    arr[EP] = temp;
    SP++;
    EP--;
}
```

$n=7$

arr[7]: {⁰ ¹ ² ³ ⁴ ⁵ ⁶
~~10~~ ~~20~~ ~~30~~ ~~40~~ ~~50~~ ~~60~~ ~~70~~
70 60 50 40 30 20 10}

SP

0

1

2

3

EP

6

5

4

3

SP < EP

T

T

T

b4exit



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Q) Reverse a Part of array

Given N array element and $[s, e]$, reverse the array from $[s, e]$.

$\boxed{[3, 7]}$

$arr[10] = \{ -3 \ 4 \ 2 \boxed{8 \ 3 \ 9 \ 6 \ 2} \ 8 \ 9 \ 10 \}$

$swap(3, 7)$
 $swap(4, 6)$

$-3 \ 4 \ 2 \boxed{2 \ 6 \ 9 \ 3 \ 8} \ 8 \ 10$

Public static void reverse (int [] arr, int s, int e) {

int n = arr.length;

worst case time complexity
T.C: $O(\frac{e-s}{2}) \approx O(e-s)$

int SP = ~~s~~; S;

T.C: $O(\frac{e-s}{2}) \approx O(e-s)$
S.C: $O(1)$ $O(n)$

while ($SP < eP$) {

int temp = arr[SP];

arr[SP] = arr[eP];

arr[eP] = temp;

SP++;

eP--;

3

3



Q) Rotate the array → {google, meta, amazon}

In Given N elements, Rotate array from last to first by K times.

$K \leq N$

$K=3$ $\text{arr}[7]: \{ 3 \ 2 \ 1 \ 4 \ 6 \ 9 \ 8 \}$

$\{ 8 \ 3 \ -2 \ 1 \ 4 \ 6 \ 9 \}$
↓
1st rot.
+2nd rot.

$\{ 9 \ 8 \ 3 \ -2 \ 1 \ 4 \ 6 \}$
↓
3rd rot.

$K=3$ $\text{arr}[7]: \{ 3 \ 2 \ 1 \ 4 \ 6 \ 9 \ 8 \}$

$\{ 6 \ 9 \ 8 \ 3 \ -2 \ 1 \ 4 \}$
↓



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$k=3^2$ arr[7]: { 3 0 -2 2 1 3 4 5 6 9 8 }



Reverse the array

{ 8 9 6 | 4 1 -2 3 }



Reverse the first k elements

{ 6 9 8 | 4 1 -2 3 }



Reverse the elements after k-th element

{ 6 9 8 | 3 4 1 -2 3 }



{ 6 9 8 | 3 4 1 -2 3 }



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11) Sseudo Code

```
int main() {
    // Input → arr[n], K
    K = K%N;
    // Step 1: Reverse the array
    reverse(arr, 0, N-1);
    // Step 2: Reverse the first K elements
    reverse(arr, 0, K-1);
```

T.C: O(n)

S.C: O(1)

// Step 3: Reverse the after K elements
reverse(arr, K, N-1);

3

```
Public static void reverse (int [] arr, int s, int e) {
    int n = arr.length;
```

T.C: O(n)

int SP = s;

S.C: O(1)

int EP = e;

```
while (SP < EP) {
    int temp = arr[SP];
    arr[SP] = arr[EP];
    arr[EP] = temp;
    SP++;
    EP--;
}
```

3

3



$K = 7$

$$\text{arr[4]} = \{ \begin{matrix} 0 & 1 & 2 & 3 \\ 4 & \perp & 6 & 9 \end{matrix} \}$$

$\downarrow \cancel{\text{loop+1}}$
 9 4 \perp 6

$\downarrow \cancel{\text{loop+2}}$
 6 9 4 \perp

$\downarrow \cancel{\text{loop+3}}$

$\perp 6 9 4$

$\downarrow \cancel{\text{loop+4}}$

4 \perp 6 9

$\downarrow \text{loop+5}$

$K = 10$

$$\text{arr[4]} = \{ \begin{matrix} 0 & 1 & 2 & 3 \\ 4 & \perp & 6 & 9 \end{matrix} \}$$

$\downarrow \text{loop+1}$
 9 4 \perp 6

$\downarrow \text{loop+2}$
 6 9 4 \perp

loop+5

loop+6

$\downarrow \text{loop+3}$

loop+7

loop+8

$\perp 6 9 4$

$\downarrow \text{loop+4}$

4 \perp 6 9

4 \perp 6 9

effective no. of rotation = $K \% N$



→ con. length

N

7

K

22

effective rotation

$$22 - 7 = 15 - 7 = 8 - 7 = 1$$

$$22 \% 7 = 1$$

7

31

$$31 \% 7 = 3$$

7

3

$$3 \% 7 = 3$$

Break till 10:50 PM



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Q) Given N array elements, count total no. of elements having atleast 1 element greater than itself.

ex: $\text{arr}[7]: \{ \overset{0}{-4}, \overset{1}{-3}, \overset{2}{1}, \overset{3}{9}, \overset{4}{3}, \overset{5}{9}, \overset{6}{4} \}$
↳ ans = 5

$\text{arr}[8]: \{ \overset{0}{3}, \overset{1}{4}, \overset{2}{11}, \overset{3}{8}, \overset{4}{2}, \overset{5}{10}, \overset{6}{9}, \overset{7}{11} \}$
↳ ans = 6

$\text{arr}[5]: \{ \overset{0}{7}, \overset{1}{7}, \overset{2}{7}, \overset{3}{7}, \overset{4}{7} \}$
↳ ans = 0

$\text{arr}[6]: \{ \overset{0}{1}, \overset{1}{2}, \overset{2}{2}, \overset{3}{3}, \overset{4}{3}, \overset{5}{4} \}$
↳ ans = 5

II idea

Obs1: man elements of the array are not valid.

Obs2: except for man elements, all the elements are valid.

→ find man no. & count it. $\rightarrow \text{mancount}$

↳ ans = No. of Elements - mancount.



II Pseudo Code

```
int Countgreater (int arr[n]) {
```

```
    int man = arr[0];
```

```
    for (int i=1; i<n; i++) {  
        if (arr[i] > man) { man = arr[i]; }  
    }
```

T.C: $O(2n) \approx O(n)$

S.C: $O(1)$

```
int mancount = 0;  
for (int i=0; i<n; i++) {  
    if (arr[i] == man) { mancount++; }  
}
```

```
return n - mancount;
```

3



manCount = 0 & 2

0 1 2 3 4 5 6 7

arr[8]: { 3 4 11 8 2 10 9 11 }

int man = arr[0];

```
| for (int i=1; i<n; i++) {  
|   if (arr[i] > man) { man = arr[i]; }  
| }
```

man = 3 & 11

6(6)

int manCount = 0;

```
| for (int i=0; i<n; i++) {  
|   if (arr[i] == man) { manCount++; }  
| }
```



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Q) Two sum

Given N array elements, check if there exists a pair (i, j) such that $\text{arr}[i] + \text{arr}[j] = k$ and $i \neq j$

Note: i and j are index values, k is given sum.

ex: $\text{arr}[7]: \{ 2 -1 0 3 2 5 7 \}$
 $k=8$

$\text{arr}[4]: \{ 1 3 -2 6 \}$
 $k=5$

$\text{arr}[5]: \{ 2 4 -3 7 10 \}$
 $k=8$

$\text{arr}[6]: \{ 3 5 1 8 3 7 \}$
 $k=6$