



Today's agenda

↳ Pow(a,n)

↳ TC & SC of recursion



AlgoPrep



Q) Given  $a$  and  $n$ , calculate  $a^n$ .

Ex:  $a \ n$

$$2 \ 3 \rightarrow 8$$

$$3 \ 5 \rightarrow 243$$

```
int Pow (int a, int n) {
    if (n == 1) { return a; }
}
```

Faith: Given  $a$  and  $n$ , calculate and return  $a^n$ .

Main logic:

```
int temp = Pow(a, n-1);
return temp * a;
```

$$\begin{array}{l} a^n \rightarrow \text{temp} * a \\ \downarrow \\ a^{n-1} \rightarrow \text{temp} \end{array}$$

Base case:

$\text{if } (n == 1) \{ \text{return } a; \}$

3  
↑  
16

Pow(2, 4) 123

↓  
8

Pow(2, 3) 123

↓  
4

Pow(2, 2) 123

↓  
2

Pow(2, 1) 1

T.C. of 1 function:  $O(1) \rightarrow n$

No. of function:  $N. \rightarrow y$

T.C of recursion:  $n * y = O(1) * N$   
 $= O(N)$

S.C of recursion:  $O(1) * N = O(N)$



II do something better than  $O(n)$ .

$$\rightarrow a^n = a^{n-1} * a$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} \quad \text{if } n \text{ is even}$$

$$\text{ex: } 2^8 = 2^4 * 2^4$$

$$2^9 = 2^4 * 2^4 * 2$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} * a \quad \text{if } n \text{ is odd}$$

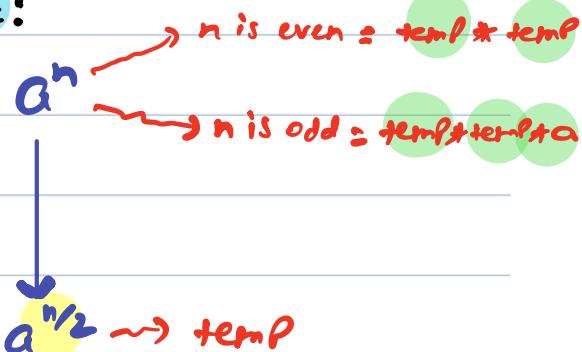
$$2^9 = 2^4 * 2^4 * 2$$

```
int Pow (int a, int n) {
    if (n == 1) { return a; }
```

Faith: Given  $a$  and  $n$ , calculate and return  $a^n$ .

Main logic:

```
int temp = Pow(a, n/2);
if (n%2 == 0) { return temp * temp; }
else { return temp * temp * a; }
```



3

Base case:

```
if (n == 1) { return a; }
```



```
int Pow (int a, int n) {
    1 if (n == 1) { return a; }
```

2 int temp = Pow (a, n/2);

3 { if (n%2 == 0) { return temp + temp; }  
else { return temp \* temp \* a; }}

3

Overall T.C:

T.C of 1 function:  $O(1)$

No. of function:  $O(\log n)$

$\approx O(\log n)$

Overall S.C:  $O(1) \times \log n \approx O(\log n)$

if n is negative

$2^{-37} \Rightarrow 2^{37} \rightarrow$  recursive soln  $\rightarrow$  any

↓

return  $\frac{1}{ans}$

$$\rightarrow a^n = a^{n/2} * a^{n/2} \rightarrow O(\log n) \rightarrow \text{most optimal here.}$$

$$a^n = a^{n/3} * a^{n/3} * a^{n/3}$$

$$a^n = a^{n/4} * a^{n/4} * a^{n/4} * a^{n/4}$$

$$a^n = a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5}$$

$$a^n = a^{n/n} * a^{n/n} * a^{n/n} * a^{n/n} * \dots - \rightarrow a^{n/2} \rightarrow \text{ideal} \rightarrow O(n)$$



```
int fib(int n) {
```

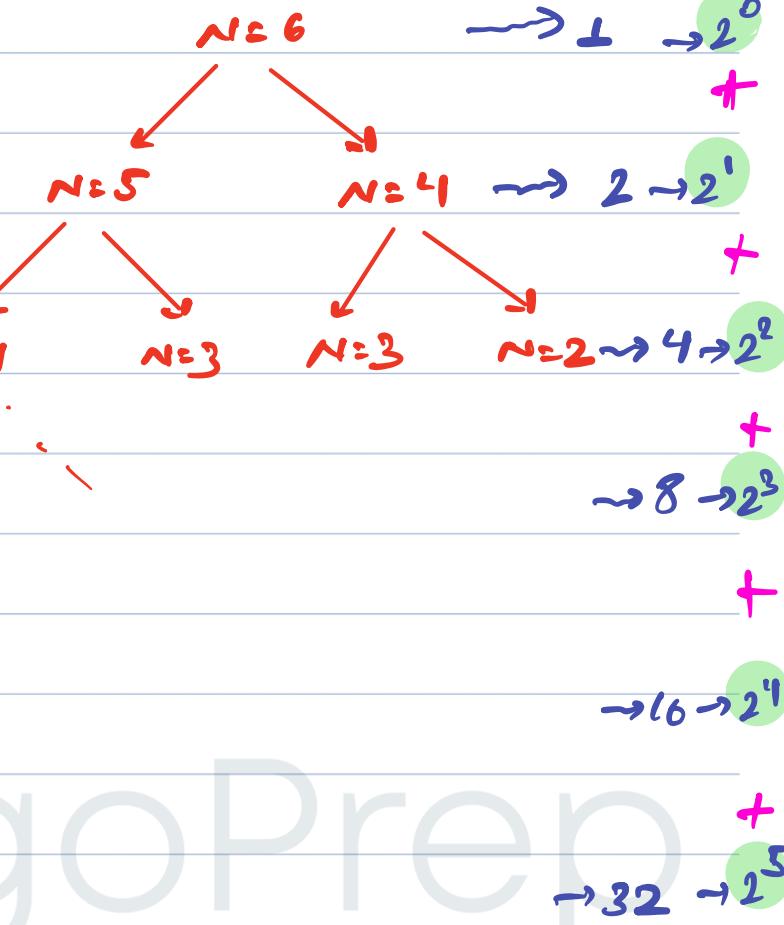
1 if ( $n == 0$ ) { return 0; }  
 2 if ( $n == 1$ ) { return 1; }

2 int temp1 = fib( $n - 1$ );

3 int temp2 = fib( $n - 2$ );

4 return temp1 + temp2;

3



$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1}$$

$$\text{Sum of G.P} = a * \frac{(2^n - 1)}{2 - 1}$$

$$= 1 * \frac{2^n - 1}{2 - 1} = 2^n - 1$$

Overall T.C:  $O(1) * 2^n \approx O(2^n)$



\*

No. of Calls =  $n$

Count of levels =  $n$

No. of functions =  $n^n$

Break till 9:49 PM.



AlgoPrep



Q) Given an array, check if it is Palindrome or not?  
↳ recursion

Ex: MALAYALAM → true

Ex: a a bba → false

Idea 1

→ MALAYALAM  
→ MALAYALAM → true

a a bba → false  
a b baa

Idea 2

$\downarrow^0$        $\downarrow^{n-1}$   
MALAYALAM

T.C:  $O(1) * N/2$   
 $\approx O(N)$   
S.C:  $O(1) * N/2 \approx O(N)$

boolean isPalindrome (char[] ch, int s, int e) {  
    if ( $s == e$ ) { return true; }  
    if ( $s > e$ ) { return true; }  
    if ( $ch[s] == ch[e]$ ) {

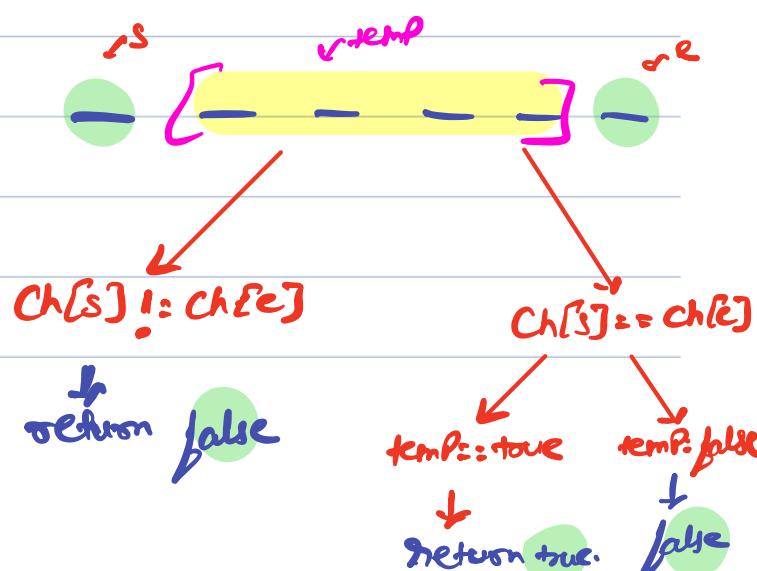
        return temp;

}

    else { return false; }

}

base case:





boolean isPalindrome(char[] ch, int s, int e) ch: 

```

if (s == e) { return true; }
if (s > e) { return true; }
if (s < e) {
    if (ch[s] == ch[e]) {
        isPalindrome(s+1, e-1);
    } else {
        return false;
    }
}

```

```

1 if (ch[s] == ch[e]) {
2     boolean temp = isPalindrome(ch, s+1, e-1);
3     return temp;
}

```

```

3 else { return false; }
3

```

e



s  
4 5 6 7

ch: 

isPalin (0, 7)



true

isPalind (1, 6)



true

isPalind (2, 5)



true

isPalin (3, 4)



true

isPalin (4, 3)

isPalindrome (1, 7) ↘  
true



isPalindrome (2, 6) ↘  
true



isPalindrome (3, 5) ↘  
true



isPalindrome (4, 4) ↘  
true



Ch: M A L A      A L B M

isPalin(0, 7)  
↓  
isPalin(1, 6)

write iterative code →

```
for (int i=0; i<n; i++) {  
    3
```

write recursive code: → within 1 function Space Complexity  
is like iteration.  
↳ you create multiple instance of same function.

pow {  
pow {  
pow {  
pow {  
pow {  
 }

Space will be Considered  
in Space Complexity.

Stack



S.C.: (Space used in 1 function) & (Max no. of function  
in Stack at any Point  
of time.)



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