



Today's agenda

↳ $\text{Pow}(a, n)$

↳ TC & SC of recursion



AlgoPrep



Q) Given a and n , calculate a^n .

Ex:

a n

2 3 \rightarrow 8

3 5 \rightarrow 243

```
int Pow (int a, int n) {  
  1 if (n == 1) { return a; }  
}
```

Faith: Given a and n , calculate and return a^n .

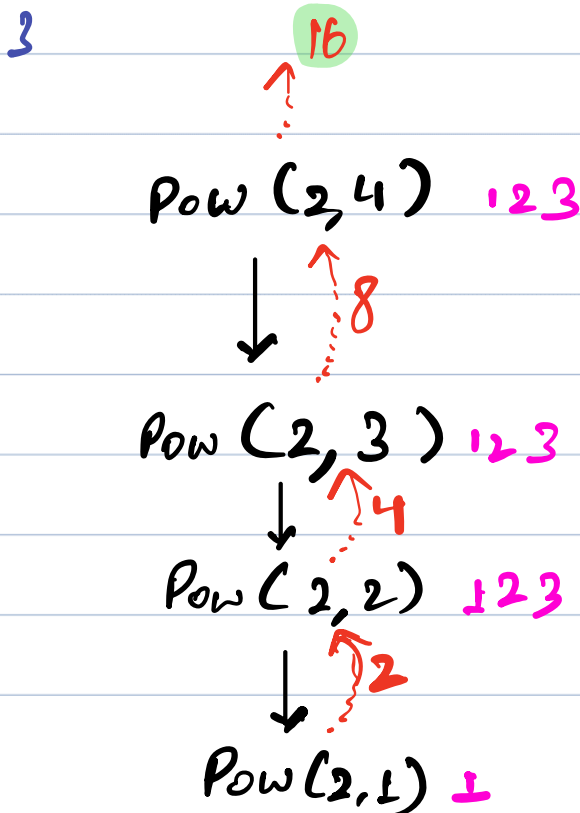
```
2 int temp = Pow(a, n-1);  
3 return temp * a;
```

Main logic:

$a^n \rightarrow temp * a$
 \downarrow
 $a^{n-1} \rightarrow temp$

Base case:

if (n == 1) { return a; }



T.C. of 1 function: $O(1) \rightarrow x$
No. of function: $N \rightarrow y$

T.C of recursion: $x * y = O(1) * N = O(N)$

S.C of recursion: $O(1) * N = O(N)$



// do something better than $O(n)$.

$$\rightarrow a^n = a^{n-1} * a$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} \quad \text{if } n \text{ is even}$$

$$\text{ex: } 2^8 = 2^4 * 2^4$$

$$2^9 = 2^4 * 2^4 * 2$$

$$\hookrightarrow a^n = a^{n/2} * a^{n/2} * a \quad \text{if } n \text{ is odd}$$

$$2^9 = 2^4 * 2^4 * 2$$

```
int Pow (int a, int n) {
    if (n == 1) { return a; }
```

Faith: Given a and n , calculate and return a^n .

```
    int temp = Pow(a, n/2);
```

```
    if (n % 2 == 0) { return temp * temp; }
```

```
    else { return temp * temp * a; }
```

Main logic:

a^n

\rightarrow n is even $= \text{temp} * \text{temp}$

\rightarrow n is odd $= \text{temp} * \text{temp} * a$

$a^{n/2} \rightarrow \text{temp}$

Base case:

```
    if (n == 1) { return a; }
```



```
int Pow (int a, int n) {
1 if (n == 1) { return a; }
```

$a = 2$

$n = 37$

```
2 int tmp = Pow(a, n/2);
3 { if (n % 2 == 0) { return tmp * tmp; }
  else { return tmp * tmp * a; }
```

}

overall T.C:

T.C of 1 function: $O(1)$

No. of function: $O(\log n)$

$\approx O(\log n)$

overall s.c: $O(1) * \log n \approx O(\log n)$

$Pow(2, 37) \rightarrow 123$

$Pow(2, 18) \rightarrow 123$

$Pow(2, 9) \rightarrow 123$

$Pow(2, 4) \rightarrow 123$

$Pow(2, 2) \rightarrow 123$

$Pow(2, 1) \rightarrow 1$

if n is negative

$2^{-37} \Rightarrow 2^{37} \rightarrow$ recursive soln \rightarrow ans

\Downarrow
return $\frac{1}{ans}$

$\rightarrow a^n = a^{n/2} * a^{n/2} \rightarrow$ most optimal here.

$a^n = a^{n/3} * a^{n/3} * a^{n/3}$

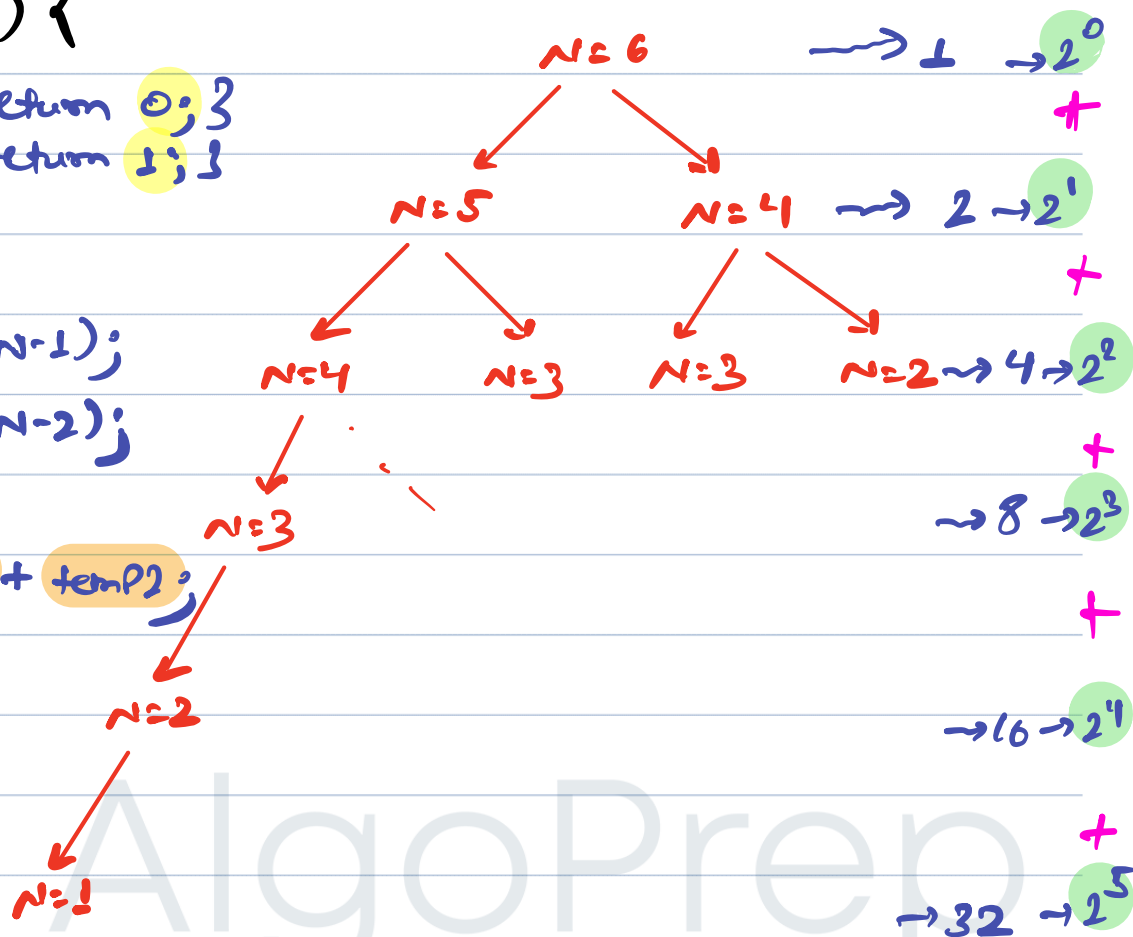
$a^n = a^{n/4} * a^{n/4} * a^{n/4} * a^{n/4}$

$a^n = a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5} * a^{n/5}$

$a^n = a^{n/n} * a^{n/n} * a^{n/n} * a^{n/n} \dots a^{n/2} \rightarrow$ ideal $\rightarrow O(n)$



```
int fib(int n) {  
1  if (n == 0) { return 0; }  
   if (n == 1) { return 1; }  
  
2  int temp1 = fib(n-1);  
3  int temp2 = fib(n-2);  
  
4  return temp1 + temp2;  
}
```



$$2^0 + 2^1 + 2^2 + 2^3 + \dots + 2^{n-1}$$

$$\text{Sum of G.P.} = a * \frac{(2^n - 1)}{2 - 1}$$

$$= 1 * \frac{2^n - 1}{2 - 1} = 2^n - 1$$

$$\text{Overall T.C: } O(1) * 2^n \approx O(2^n)$$



*

No. of Calls = n

Count of levels = n

no. of functions = 2^n

Break till 9:49 PM.



AlgoPrep



Q) Given an array, check if it is Palindrome or not?
 ↳ Recursion

Ex: MALAYALAM → true

Ex: a a b b a → false

// Idea 1

→ MALAYALAM

→ MALAYALAM

→ true

a a b b a

→ false

a b b a a

// Idea 2

0 ↙ MALAYALAM ↘ N-1

T.C: $O(1) * N/2$

$\approx O(N)$

S.C: $O(1) * N/2 \approx O(N)$

boolean isPalindrome(char[] ch, int s, int e) { Faith: Check and

if (s == e) { return true; }

if (s > e) { return true; }

return whether

ch array is Palindrome

betⁿ 0 to n-1.

if (ch[s] == ch[e]) {

boolean temp = isPalindrome(ch, s+1, e-1);

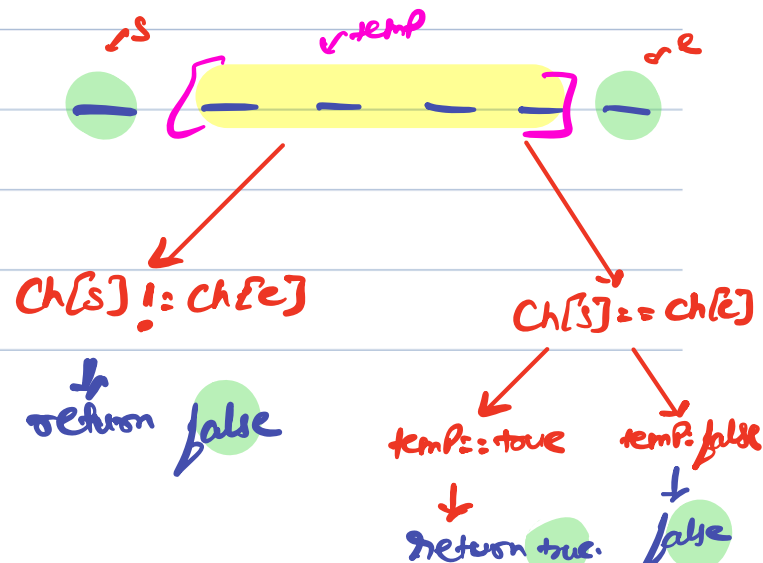
return temp;

main logic:

}

else { return false; }

}



base case:



boolean isPalindrom(char[] ch, int s, int e)

ch: M A L A Y A L A M

```
1 if (s == e) { return true; }  
2 if (s > e) { return true; }  
3 if (s < e) {  
4     if (ch[s] == ch[e]) {  
5         boolean temp = isPalindrom(ch, s+1, e-1);  
6         return temp;  
7     }  
8     else { return false; }  
9 }
```

isPalindrom(0, 8) → true

isPalindrom(1, 7) → true

isPalindrom(2, 6) → true

isPalindrom(3, 5) → true

isPalindrom(4, 4) → true

isPalin(0, 7)

isPalind(1, 6)

isPalind(2, 5)

isPalin(3, 4)

isPalin(4, 3)



ch: M A L A A L B M

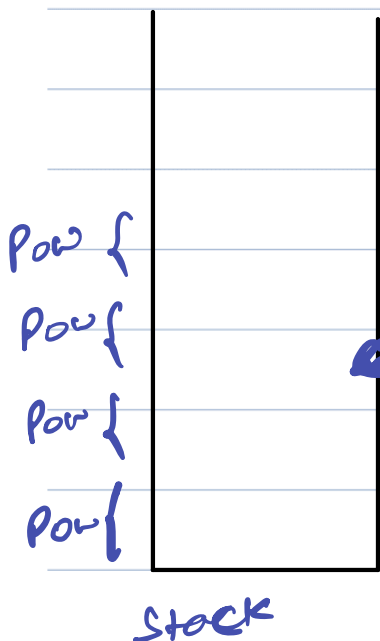
isPalin(0, 7)

isPalin(1, 6)

write iterative code → for (int i=0; i<N; i++) {

}

write recursive code: → within 1 function space complexity is like iteration.
↳ you create multiple instance of same function.



Space will be considered in space complexity.



S.C: (space used in 1 function) or (max no. of function in stack at any point of time.)



AlgoPrep