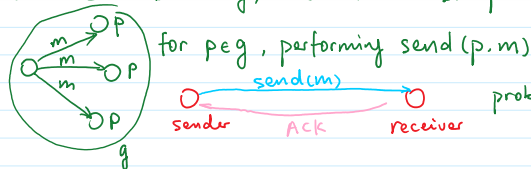


Basic / Reliable Multicast

Basic Multicast - B-delivery.

Guaranteed Delivery, unless multicast process crashes.



problem: ACK received on sender, at almost same time. if buff overflow, cause of dropping of ACKs retransmit would create more ACKs

Reliable Multicast \rightarrow R-multicast
R-delivery.

Properties: a correct proc p delivers a msg at most once \rightarrow no duplication / retransmission
if a correct proc p multicast a msg, it would eventually deliver msg. (validity).
if a correct (on receiving) proc p delivers msg m , then every other proc in group g would deliver msg m as well (equality)

For proc p to R-multicast to group g :

init: $\text{recv queue} := \{\}$

send: for q in $\text{group}(g)$:

B-multicast $\langle g, m \rangle + \{p \text{ itself}\}$

recv: if $\{m\} \notin \text{recv queue} \rightarrow \text{recv queue} += \{m\}$ (non-duplication).

then reply p with Ack (Agreement), send m to any other proc in g , if $(q \neq p)$.

pass msg m to process on g .

$\text{msg class} + \text{isEqual}$.

Sequencer TO

* for proc in group g :

sending: B-multicast $\langle \text{msg}, i \rangle$ to $g \cup \{\text{sequencer}(g)\}$ \rightarrow id-proc

receiving: $\text{msg} \langle m, i \rangle$ from peer: hold-back $Q += \langle m, i \rangle$

seq msg $\langle \text{order}, i, S \rangle$ from sequencer, wait until $S = r_g \rightarrow$ deliver m in holdback Q (TO-deliver)

$r_g = S + 1$

* for sequencer of g :

init: $S_g = 0$

receiving: $\langle m, i \rangle \rightarrow$ B-multicast $\langle \text{order}, i, S_g \rangle$,

S_g++ .

msg + id, class. \downarrow direction
sequencer2, g