## Dynamical Systems (ODEs)

A continuous-time dynamical system is defined by a system of n ODEs:  $\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x}, t; \theta)$ , where  $\mathbf{x} \in \mathbb{R}^n$ .

- Equilibrium Point  $\mathbf{x}_e$ : A state where the system does not change, i.e.,  $\mathbf{f}(\mathbf{x}_e)=0$ .
- Linear System: A system of the form  $\dot{\mathbf{x}} = A\mathbf{x}$ .

## Solving Linear Systems: $\dot{\mathbf{x}} = A\mathbf{x}$

- 1. Find eigenvalues  $\lambda_i$  by solving the characteristic equation:  $\det(A-\lambda I)=0.$
- 2. For each eigenvalue  $\lambda_i$ , find the corresponding eigenvector  $\mathbf{u}_i$  by solving  $(A \lambda_i I) \mathbf{u}_i = \mathbf{0}$ .
- The general solution is a linear combination of the "straight-line" solutions:

$$\mathbf{x}(t) = \sum_{i=1}^{n} c_i e^{\lambda_i t} \mathbf{u}_i$$

The constants  $c_i$  are determined by the initial conditions  $\mathbf{x}(0)$ .

## Stability of Equilibria

Stability for a linear system's equilibrium at the origin (or for a nonlinear system's equilibrium  $\mathbf{x}_e$  based on its Jacobian J) is determined by the eigenvalues  $\lambda = \alpha \pm i\beta$ .

- Asymptotically Stable: All eigenvalues have real parts  $\alpha < 0$ . All nearby solutions converge to the equilibrium.
- Stable: All eigenvalues have  $\alpha \leq 0$ . No real parts are positive, and any with  $\alpha = 0$  (purely imaginary) are simple. Nearby solutions stay nearby, but don't necessarily converge (e.g., centers).
- Unstable: At least one eigenvalue has a real part  $\alpha > 0$ . Most nearby solutions will move away.

Eigenvalues	Type
$\lambda_1, \lambda_2$ real, distinct	
$\lambda_1, \lambda_2 < 0$	Stable Node (All paths head to origin)
$\lambda_1, \lambda_2 > 0$	Unstable Node (All paths leave origin)
$\lambda_1 \cdot \lambda_2 < 0$	Saddle (Unstable)
$\lambda_1 = \lambda_2$ real, repeated	<b>Node</b> (Stable if $\lambda < 0$ , Unstable if $\lambda > 0$ )
$\lambda = \alpha \pm i\beta$ (where $\beta \neq 0$ )	
$\alpha < 0$	Stable Spiral (Spiral into the origin)
$\alpha > 0$	Unstable Spiral (Spiral away from the origin)
$\alpha = 0$	Center (Neutrally Stable, paths are closed orbits)

## Nonlinear Systems: $\dot{\mathbf{x}} = \mathbf{f}(\mathbf{x})$

- **Linearization**: To analyze the stability of an equilibrium point  $\mathbf{x}_e$ , we linearize the system using the **Jacobian matrix**  $J = J_{\mathbf{f}}(\mathbf{x}_e)$ , where  $J_{ij} = \frac{\partial f_i}{\partial x_j}$ .
- We then analyze the stability of the linear system  $\dot{\mathbf{u}}=J\mathbf{u}$ , where  $\mathbf{u}=\mathbf{x}-\mathbf{x}_e.$
- Hartman-Grobman Theorem: The stability of the linearized system (based on J's eigenvalues) usually determines the stability of the nonlinear system, except for marginal cases (like  $\alpha=0$  centers).

# Other 2D+ Concepts

- Nullclines: Curves in the phase space where one component of the vector field is zero (e.g.,  $\dot{x}=0$  or  $\dot{y}=0$ ). Equilibrium points occur at the intersections of \*all\* nullclines.
- Limit Cycle: A stable, isolated, closed-orbit trajectory. Solutions nearby may spiral into or away from a limit cycle.
- Poincaré-Bendixson Theorem: In a 2D continuous system, the only possible long-term behaviors are approaching an equilibrium point, approaching a limit cycle, or diverging to infinity. This theorem implies that chaos cannot occur in 2D continuous systems.
- Chaos: Aperiodic, long-term behavior in a deterministic system that
  exhibits sensitive dependence on initial conditions. Associated with
  Strange Attractors (e.g., Lorenz Attractor) and requires 3 or more
  dimensions for continuous systems.

# Iterated Maps (Discrete-Time Systems)

A discrete-time system:  $x_{n+1} = f(x_n)$ . The sequence  $x_0, x_1, x_2, ...$  is the **orbit**.

- Fixed Point  $x^*$ : A point that maps to itself,  $f(x^*) = x^*$ .
- Stability of Fixed Points: Determined by the multiplier  $\lambda = f'(x^*)$ .
  - $|\lambda| < 1$ : **Stable** (attracting). Orbits starting near  $x^*$  converge to it.
  - $-\ |\lambda|>1$ : Unstable (repelling). Orbits starting near  $x^*$  (but not exactly on it) move away.
  - $|\lambda| = 1$ : Marginal case (e.g.,  $f'(x^*) = -1$  leads to a bifurcation).
- Cobweb Plot: A graphical method to visualize orbits. Draw a line from  $(x_n, x_n)$  to  $(x_n, f(x_n))$ , then horizontally to  $(f(x_n), f(x_n)) = (x_{n+1}, x_{n+1})$ , and repeat.

## Logistic Map

The canonical example of a route to chaos:  $x_{n+1} = rx_n(1 - x_n)$ .

- $\bullet$  As the parameter r increases (from 0 to 4), the system's long-term behavior changes.
- It moves from a single stable fixed point, to a stable 2-cycle, then a 4-cycle, 8-cycle, etc. This is the **period-doubling bifurcation** route to chaos.
- Past a certain r, the system becomes chaotic, with regions of stability ("islands") interspersed.

### Lyapunov Exponent

Measures the average exponential rate of divergence or convergence of nearby orbits, quantifying sensitivity to initial conditions.

$$\lambda = \lim_{n \to \infty} \frac{1}{n} \sum_{k=0}^{n-1} \ln |f'(x_k)|$$

- $\lambda > 0 \implies$  Chaos. Nearby orbits diverge exponentially.
- $\lambda < 0 \implies$  **Stable**. Nearby orbits converge.
- $\lambda = 0 \implies \text{Marginal}$ .

## Cellular Automata (CA)

Discrete-time, discrete-space, discrete-state systems where cells update their state based on a **local rule** applied to their **neighborhood**.

- Neighborhood: Set of cells that influence a cell's next state.
- Von Neumann: 4 neighbors (N, S, E, W).
- Moore: 8 neighbors (all adjacent cells).
- Update Schedule:
- Synchronous: All cells update simultaneously.
- Asynchronous: Cells update one at a time (e.g., in a random order).
- Combinatorics: For k states and a neighborhood of size |N|, the total number of possible rules is  $k^{(k^{|N|})}$ .
- Wolfram's 1D CA: k=2, |N|=3 (cell + left/right neighbors)  $\Rightarrow 2^{(2^3)}=256$  possible rules.

# Wolfram's 4 Classes of Behavior

- Class 1 (Stable): Evolves to a stable, homogeneous state (e.g., all black or all white).
- 2. Class 2 (Periodic): Evolves to simple periodic structures (e.g., stable patterns or simple oscillators).
- 3. Class 3 (Chaotic): Exhibits chaotic, aperiodic, fractal-like patterns (e.g., Rule 30).
- 4. Class 4 (Complex): Exhibits complex, localized structures ("gliders") that move and interact. Can support computation (e.g., Rule 110 is Turing complete).

# Conway's Game of Life (2D CA)

A famous 2D CA with a Moore neighborhood and 2 states (live/dead).

- Survival: A live cell with exactly 2 or 3 live neighbors survives.
- **Death**: A live cell with < 2 neighbors (loneliness) or > 3 neighbors (overcrowding) dies.
- Reproduction: A dead cell with exactly 3 live neighbors becomes alive.

# Networks

Networks (graphs) consist of nodes (vertices) and edges (links).

# Key Properties of Real Networks

- Small-World: Low average path length h. The average distance between any two nodes is short, typically  $h \sim O(\log N)$ .
- **High Clustering**: High clustering coefficient C. A node's neighbors are also likely to be neighbors of each other.
- Scale-Free: The degree distribution  $p_k$  (probability a node has k links) follows a power-law,  $p_k \sim k^{-\gamma}$ . This implies a few "hubs" with many links and many nodes with few links.

#### **Network Models**

Model	Scale-Free?	Small-World?	High C?
Erdos-Renyi (ER)	No	Yes	No
Watts-Strogatz (WS)	No	Yes	Yes
Barabasi-Albert (BA)	Yes	Yes	Kinda

- Erdos-Renyi (ER): A random graph. Start with N nodes, connect every pair with probability p. Degree distribution  $p_k$  is Poisson (peaked).
- Watts-Strogatz (WS): The "small-world" model. Start with a regular ring lattice (high C, high h), and "rewire" each edge with probability p. This quickly lowers h while retaining high C.
- Barabasi-Albert (BA): The "scale-free" model. Built via two mechanisms:
  - 1. Growth: Start with a small network, add one node at a time.
  - 2. **Preferential Attachment**: New nodes prefer to link to existing nodes that already have a high degree ("rich get richer").

### Centrality Measures

Ways to quantify a node's "importance" in a network.

- Degree: Number of connections. Simple count of a node's "popularity".
- Betweenness: Fraction of all shortest paths in the network that pass through this node. Identifies "bridges" or "bottlenecks".
- Closeness: Inverse of the average shortest-path distance to all other nodes. Measures how "central" a node is or how fast it can reach everyone.
- Eigenvector: A node's importance is determined by the importance of its neighbors. A node is important if it's connected to other important nodes. Solved by finding the principal eigenvector of the adjacency matrix A: Ac = λc.
- Alpha-Centrality: A generalization that includes a node's intrinsic importance:  $\mathbf{c} = \beta (I \alpha A)^{-1} \mathbf{e}$ .

## Optimization

### Gradient Descent (GD)

An iterative algorithm to find a local minimum of a function  $f(\mathbf{w})$  by repeatedly moving in the direction of the negative gradient,  $-\nabla f(\mathbf{w})$ .

$$\mathbf{w}_{k+1} = \mathbf{w}_k - \alpha \nabla f(\mathbf{w}_k)$$

Where  $\alpha$  is the **learning rate**.

- Batch GD: Calculate  $\nabla f$  using the entire dataset. Very accurate gradient, but computationally slow for large datasets.
- Stochastic GD (SGD): Calculate ∇f using only \*one\* data point.
   Very fast updates, but the gradient is noisy, leading to a volatile convergence path.
- Mini-batch SGD: A compromise. Calculate  $\nabla f$  using a small batch of data. Offers a balance between the stability of Batch GD and the speed of SGD.

Adam Optimizer: An advanced optimization algorithm popular for deep learning. It adaptively adjusts the learning rate for each parameter, combining the ideas of Momentum (using a moving average of the 1st moment/gradient) and RMSprop (using a moving average of the 2nd moment/squared gradient).

## Local Search

An iterative improvement heuristic. Start with a candidate solution and repeatedly move to a "neighboring" solution if it is better.

- Neighborhood: The set of solutions accessible from the current solution by a small change.
- Example: TSP Neighborhoods:
  - 2-opt: Remove two edges from the tour, and reconnect the four resulting endpoints in the only other possible way (which "uncrosses" the paths).
  - 3-opt: Remove three edges, reconnect in a way that improves the tour.

## Particle Swarm Optimization (PSO)

A population-based stochastic optimization algorithm inspired by social behavior (e.g., bird flocking). Used for black-box optimization.

- A "swarm" of particles (candidate solutions) "fly" through the search space.
- Each particle i has a position  $\mathbf{x}_i$  (its solution) and a velocity  $\mathbf{v}_i$ .
- Each particle remembers its **personal best** position found so far: pbest<sub>i</sub>.
- The swarm tracks the global best position found by \*any\* particle: qbest.

## Core Update Equations

$$\mathbf{v}_{i}(t+1) = \underbrace{\mathbf{v}_{i}(t)}_{\text{Inertia}} + \underbrace{c_{1}r_{1}(pbest_{i} - \mathbf{x}_{i}(t))}_{\text{Cognitive/Personal}} + \underbrace{c_{2}r_{2}(gbest - \mathbf{x}_{i}(t))}_{\text{Social}}$$

$$\mathbf{x}_{i}(t+1) = \mathbf{x}_{i}(t) + \mathbf{v}_{i}(t+1)$$

- $\omega$ : Inertia weight. Balances exploration (high  $\omega$ ) and exploitation (low  $\omega$ ).
- $c_1, c_2$ : Acceleration coefficients. Control the "pull" towards the personal best  $(c_1, \text{cognitive})$  and global best  $(c_2, \text{social})$ .
- $r_1, r_2$ : Random numbers in [0, 1] to add stochasticity.

#### **Topologies**

Defines how information (the *qbest*) is shared among particles.

- **gbest (Global Best)**: All particles are connected. The *gbest* is the best of the entire swarm. Converges very fast, but can get stuck in local optima.
- **lbest (Local Best)**: Particles are in a smaller neighborhood (e.g., a ring). The *gbest* in the equation is replaced with the *lbest* (best in the neighborhood). Slower convergence, but more robust to local optima.

## Multi-Robot Task Allocation (MRTA)

The problem of assigning a set of tasks T to a set of robots R to optimize a collective objective (e.g., minimize time, maximize utility).

### MRTA Taxonomy

Problems are classified by:

- ST/MT: Single-Task / Multi-Task robots (robots can handle one vs. many tasks at a time).
- SR/MR: Single-Robot / Multi-Robot tasks (tasks require one vs. a team of robots).
- IA/TA: Instantaneous Assignment / Time-Extended Assignment (tasks are just assigned vs. tasks involve durations and travel, requiring scheduling/routing).

## Mapping MRTA to Optimization Models

MRTA Type	Optimization
ST-SR-IA (One robot per task, one task per robot) MT-SR-IA (Robots can take multiple tasks) ST-SR-TA (One robot per task, includes routing) MT-SR-TA (Robots take multiple tasks, includes routing)	LAP (Linear A GAP (Generali mTSP (Multipl VRP (Vehicle I

# Set-Based Formulations (for MR tasks)

Used for coalition formation, where tasks require multiple robots (MR).

• Set Covering (MT-MR-IA): Find the minimum cost set of coalitions (subsets of robots) such that \*every task is covered at least once\*. Models MT-MR-IA, as robots can be in multiple coalitions.

$$\min \sum c_j x_j \quad \text{s.t.} \quad \sum a_{ij} x_j \ge 1$$

• Set Packing (ST-MR-IA): Find the maximum profit set of coalitions such that \*each task is covered at most once\*. Models ST-MR-IA, where task allocation must be exclusive.

$$\max \sum p_j x_j$$
 s.t.  $\sum a_{ij} x_j \le 1$ 

• Set Partitioning: Find the coalitions such that \*every task is covered exactly once\*. This is a very common base for routing problems.