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COE 301: Computer Architecture

LAB 07: MIPS Functions and Stack Segment

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Agenda

- Caller vs. Callee
- Functions: Declaration, Execute (Call), Return Back
- Registers Use
- Stack Segment
- Recursive Function Example
- Live Examples
- Tasks

Caller vs. Callee

- The function that initiates the call to another function is known as Caller.
- The function that receives and executes the call is known as the Callee.
- To execute a function, the program must follow these steps:
 - The caller must put the parameters in a place where the callee function can access them
 - Transfer control to the callee function
 - Execute the callee function
 - The callee function must put the results in a place where the caller can access them
 - Return control to the caller (point of origin) next to where the call was made

Functions: Declaration, Execution(Call), Return Back

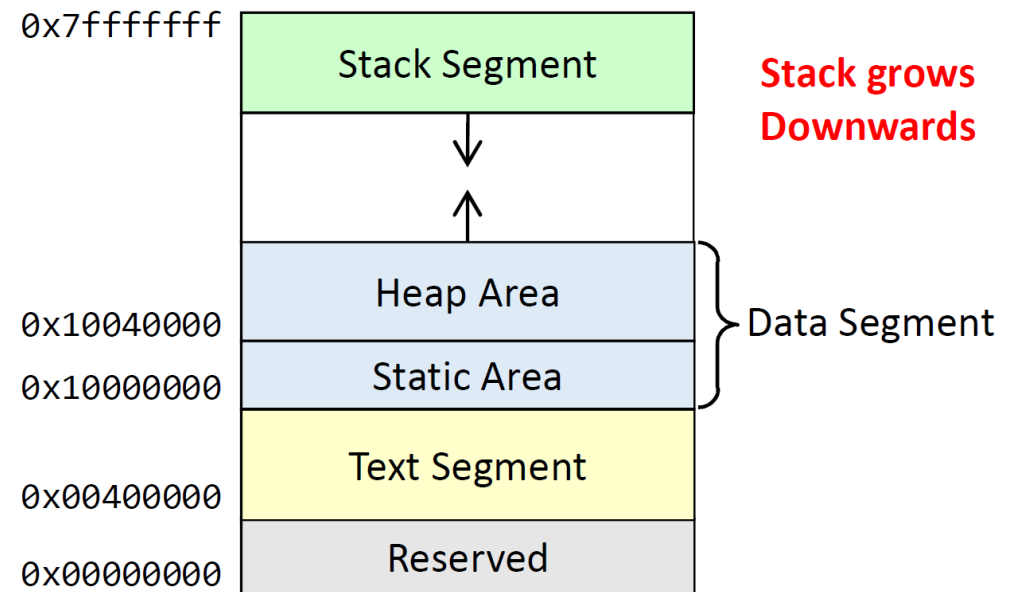
- Declaration:
 - Define a label similar to if statements and loops
 - Write the body of the function after the label
- Execution:
 - Prepare the arguments in \$a0-\$a3 registers
 - Call the function using the jal instruction (e.g. jal function)
- Return Back
 - Prepare the results if any in \$v0-\$v1 registers
 - Return to the caller using jr instruction (jr \$ra)

Registers Use

Register Name	Register Number	Register Usage
\$zero	\$0	Always zero, forced by hardware
\$at	\$1	Assembler Temporary register, reserved for assembler use
\$v0 - \$v1	\$2 - \$3	Results of a function
\$a0 - \$a3	\$4 - \$7	Arguments of a function
\$t0 - \$t7	\$8 - \$15	Registers for storing temporary values
\$s0 - \$s7	\$16 - \$23	Registers that should be saved across function calls
\$t8 - \$t9	\$24 - \$25	Registers for storing more temporary values
\$k0 - \$k1	\$26 - \$27	Registers reserved for the OS kernel use
\$gp	\$28	Global Pointer register that points to global data
\$sp	\$29	Stack Pointer register that points to top of stack
\$fp	\$30	Frame Pointer register that points to stack frame
\$ra	\$31	Return Address register used to return from a function call

Stack Segment

- Stack Segment provides an area that can be allocated and freed by functions. The programmer has no control over where these segments are located in memory.
- The stack segment can be used by functions for passing many parameters, for allocating space for local variables, and for saving and preserving registers across calls.
- Without the stack segment in memory, it would be impossible to write recursive functions, or pure functions that have no side effects.



Recursive Function Example

```
int fact (int n) {  
    if (n<2) return 1;  
    else return (n*fact(n-1));  
}
```

```
fact:  
    bge $a0, 2, else    # branch if (n >= 2) to else  
    li $v0, 1           # $v0 = 1  
    jr $ra              # return to caller  
  
else:  
    addi $sp, $sp, -8   # allocate a stack frame of 8 bytes  
    sw $a0, 0($sp)      # save the argument n  
    sw $ra, 4($sp)      # save the return address  
    addi $a0, $a0, -1   # argument $a0 = n-1  
    jal fact            # call fact(n-1)  
    lw $a0, 0($sp)      # restore $a0 = n  
    lw $ra, 4($sp)      # restore return address  
    mul $v0, $a0, $v0   # $v0 = n * fact(n-1)  
    addi $sp, $sp, 8    # free stack frame  
    jr $ra              # return to the caller
```

Live Examples