

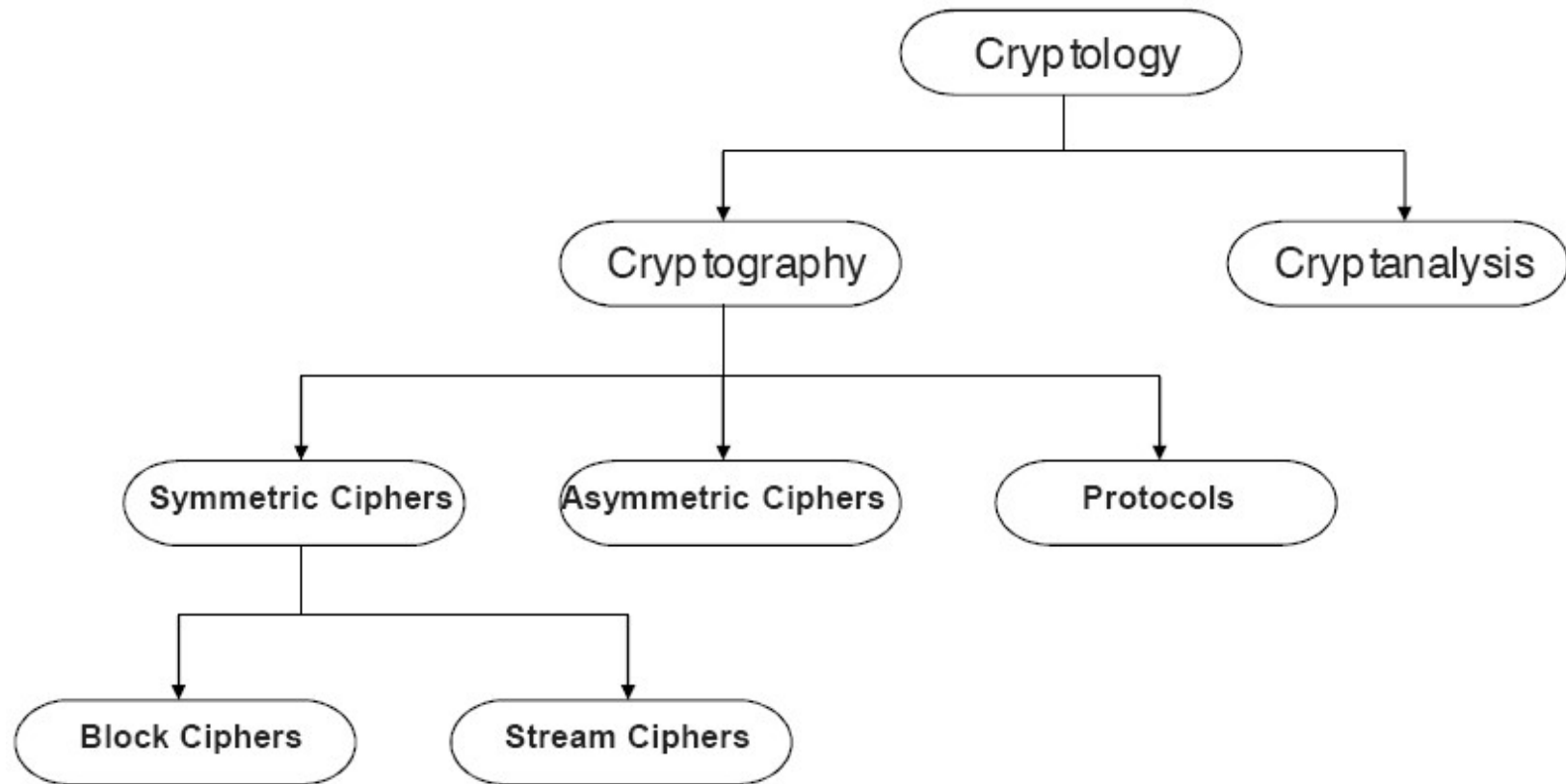
Information Security

CS3002

Lecture 9
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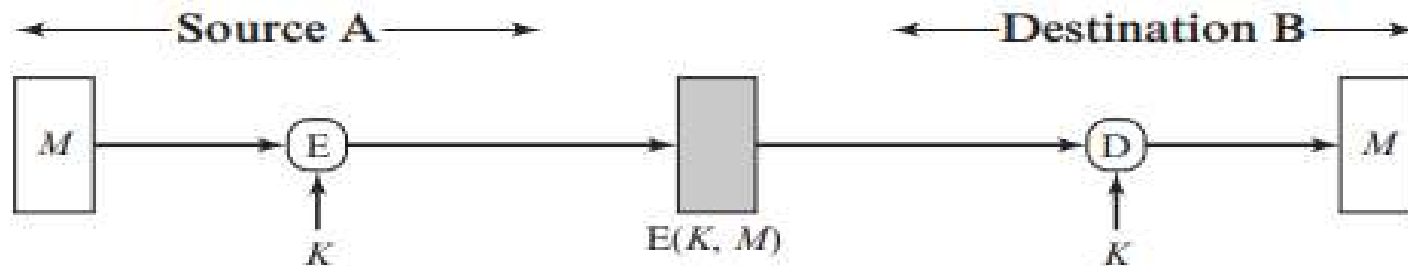
Classification of the Field of Cryptology



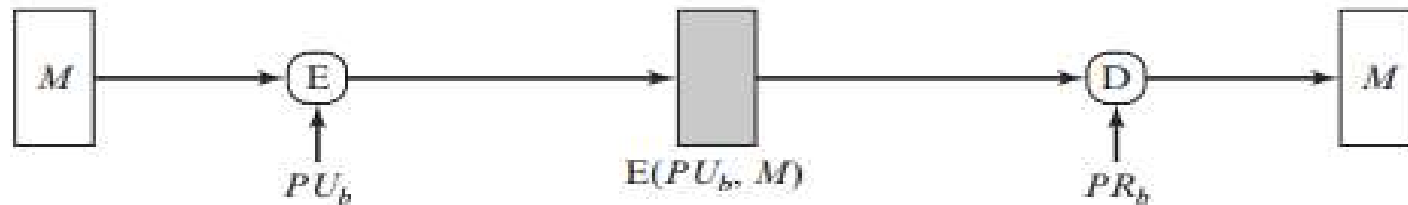
Adopted with thanks from: Chapter 1 of Understanding Cryptography by Christof Paar and Jan Pelzl

MESSAGE AUTHENTICATION CODE (MAC)

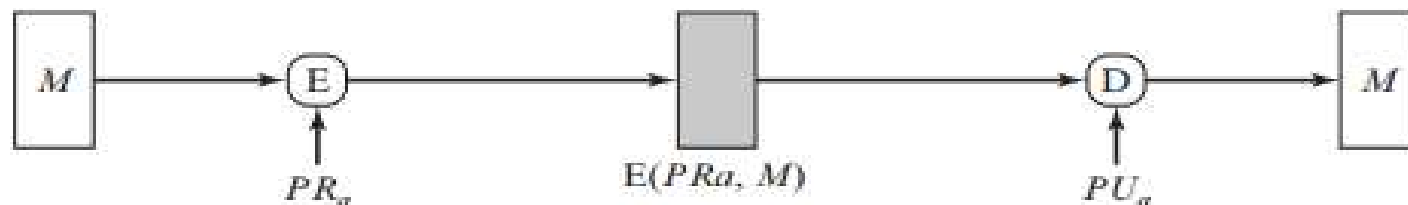
Authentication by Using of Encryption



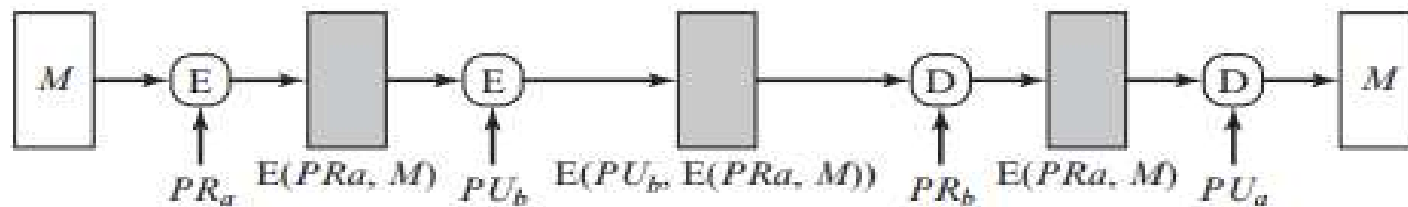
(a) Symmetric encryption: confidentiality and authentication



(b) Public-key encryption: confidentiality

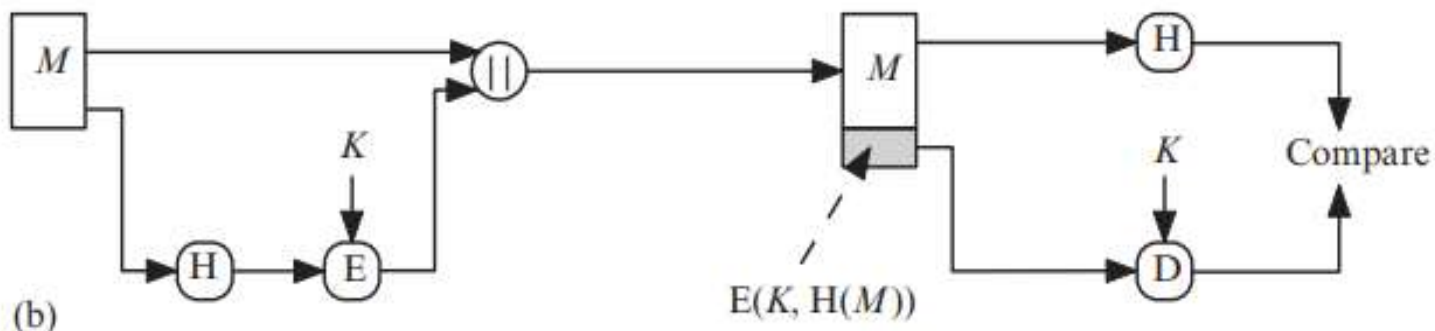
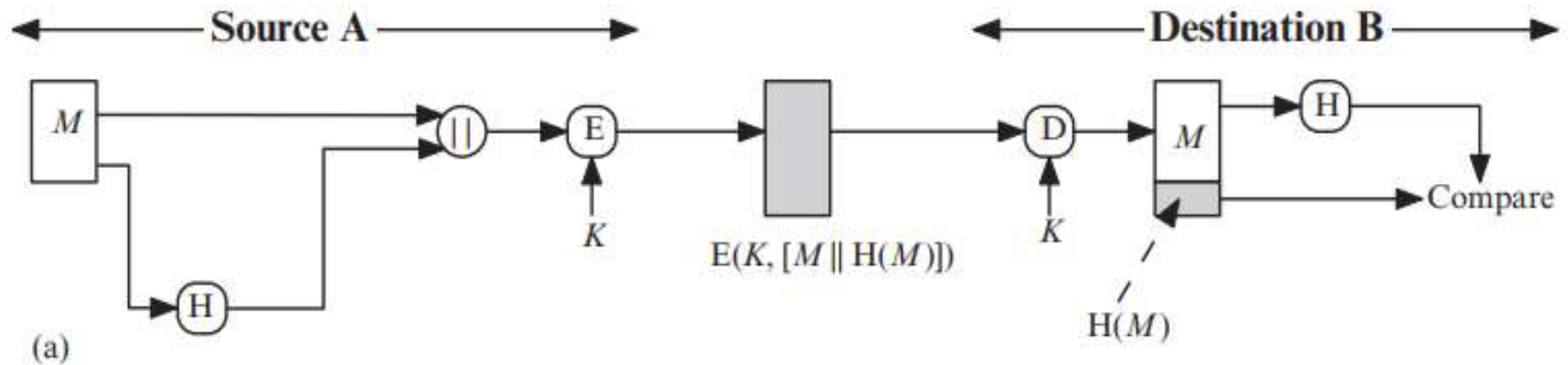


(c) Public-key encryption: authentication and signature

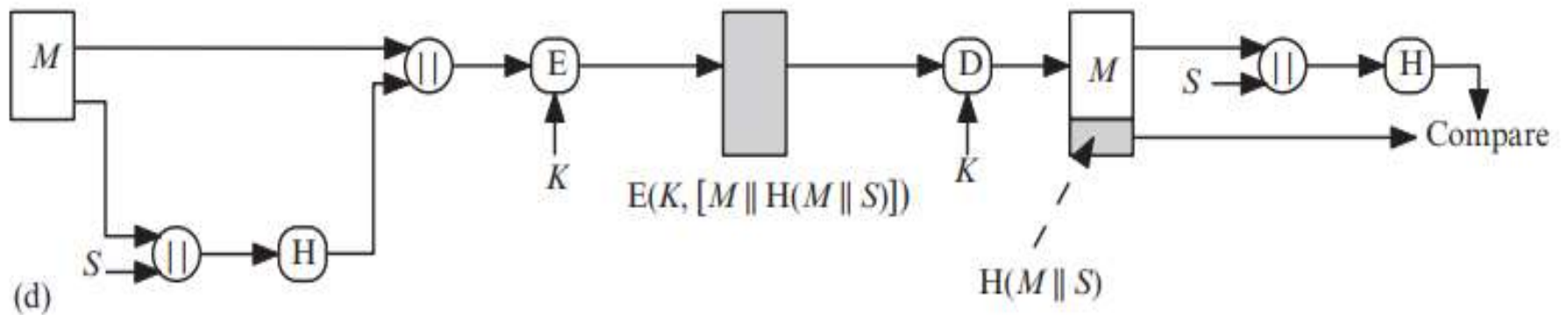
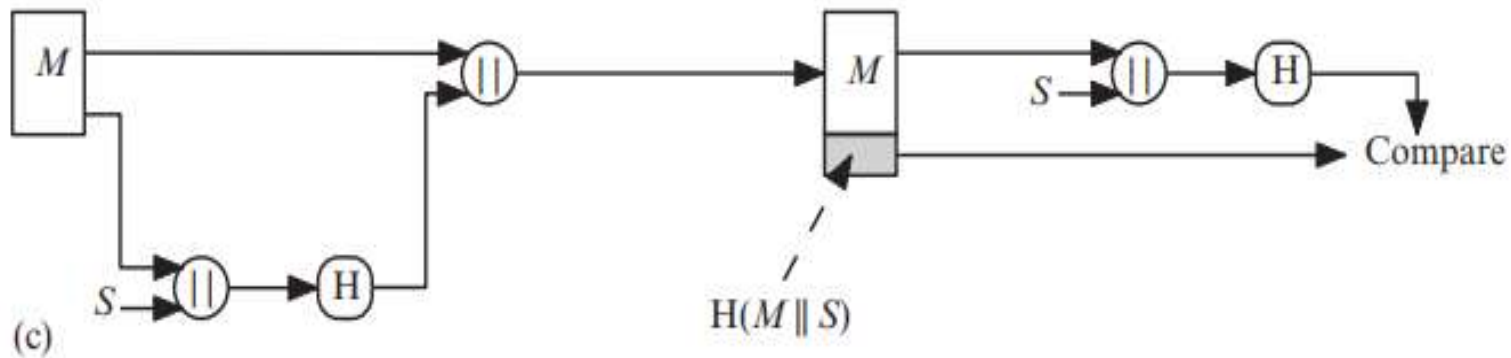


(d) Public-key encryption: confidentiality, authentication, and signature

Authentication by Using of Hash Function

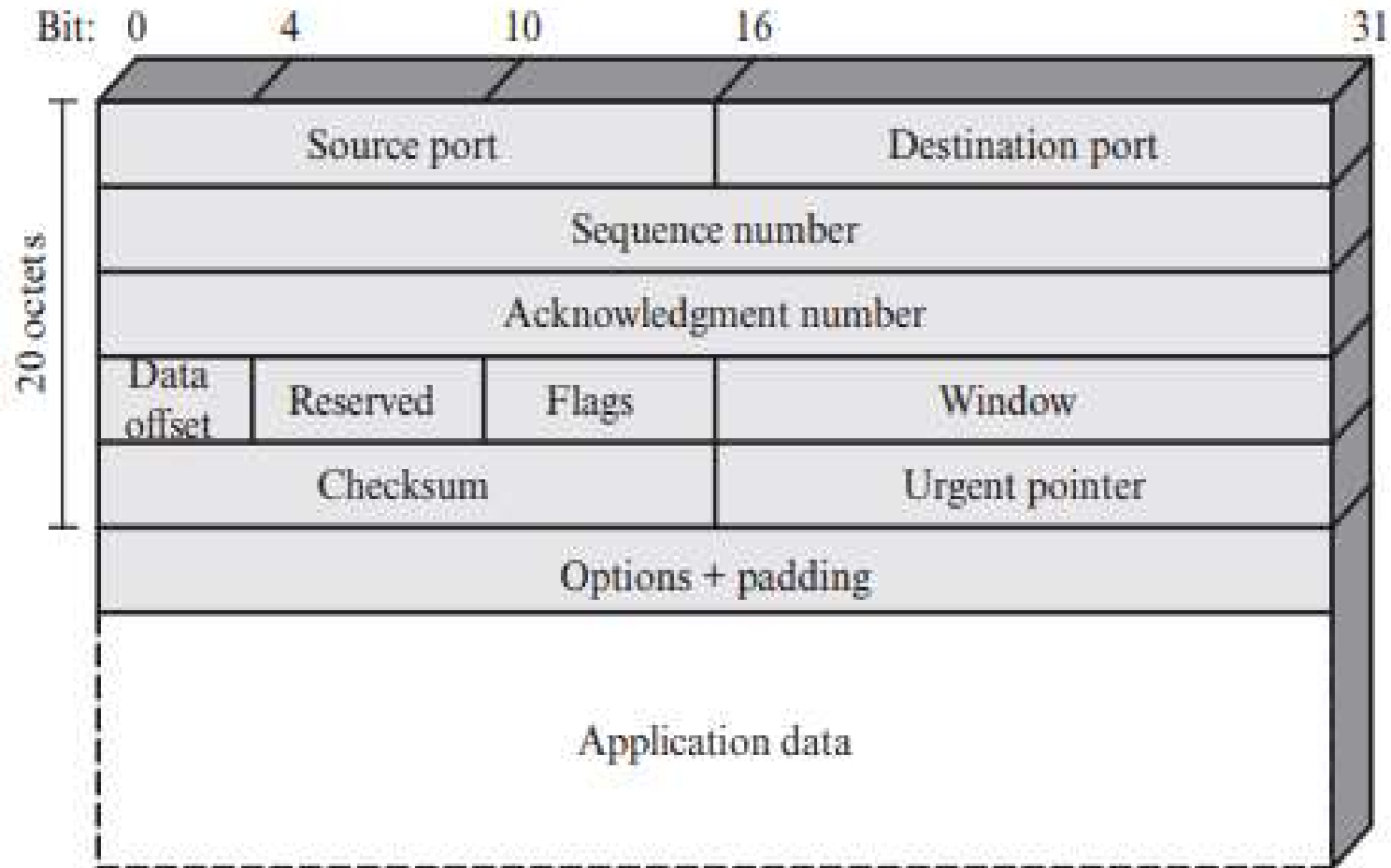


Authentication by Using of Hash Function



MESSAGE AUTHENTICATION CODE (MAC)

TCP Segment



Message Authentication Code (MAC)

- Also known as a *keyed hash function*.
- Typically used between two parties that share a secret key to authenticate information exchanged between those parties.

Takes as input a secret key and a data block and produces a hash value (MAC) which is associated with the protected message

- If the integrity of the message needs to be checked, the MAC function can be applied to the message and the result compared with the associated MAC value.
- An attacker who alters the message will be unable to alter the associated MAC value without knowledge of the secret key.

Message Authentication Code (MAC)

- Generated by an algorithm that creates a small fixed-sized block.
 - Depending on both message and some key.
 - Like encryption though need not be reversible.
- Appended to message as a **signature**.
- Receiver performs same computation on message and checks it matches the MAC.
- Provides assurance that message is unaltered and comes from sender.

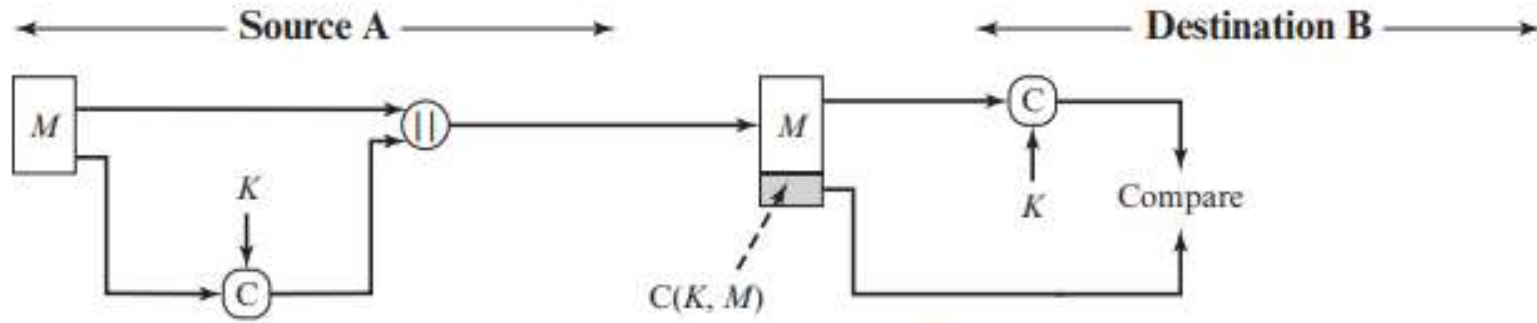
MAC Properties

- A MAC is a cryptographic checksum

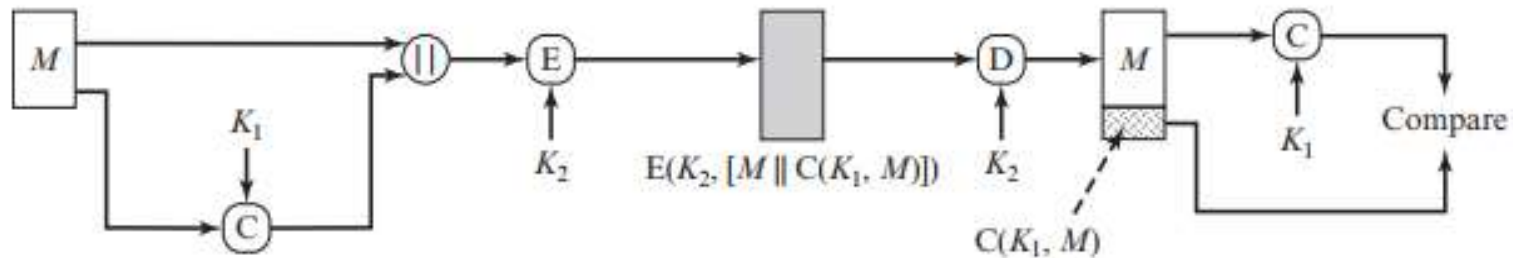
$$\text{MAC} = C(K, M)$$

- Condenses a variable-length message M .
 - Using a secret key K .
 - To a fixed-sized authenticator.
- Is a many-to-one function
 - Potentially many messages have same MAC.
 - But finding these needs to be very difficult.

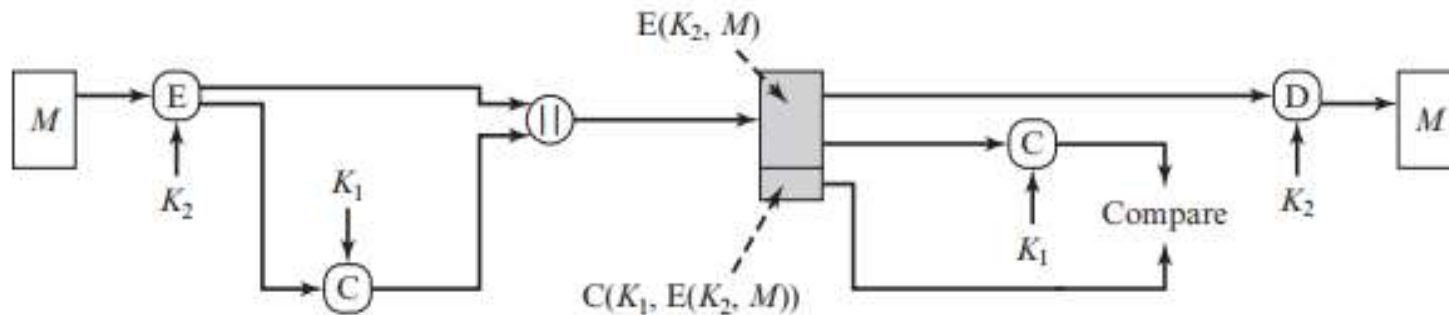
Message Authentication Code (cont.)



(a) Message authentication



(b) Message authentication and confidentiality; authentication tied to plaintext



(c) Message authentication and confidentiality; authentication tied to ciphertext

Requirements for MACs

- Taking into account the types of attacks need the MAC to satisfy the following:
 1. Knowing a message and MAC, is infeasible to find another message with same MAC.
 2. MACs should be uniformly distributed.
 3. MAC should depend equally on all bits of the message.

Brute-Force Attack

- Requires known message-tag pairs
 - A brute-force method of finding a collision is to pick a random bit string y and check if $H(y) = H(x)$

Two lines of attack:

- Attack the key space
 - If an attacker can determine the MAC key then it is possible to generate a valid MAC value for any input x .
- Attack the MAC value
 - Objective is to generate a valid tag for a given message or to find a message that matches a given tag.

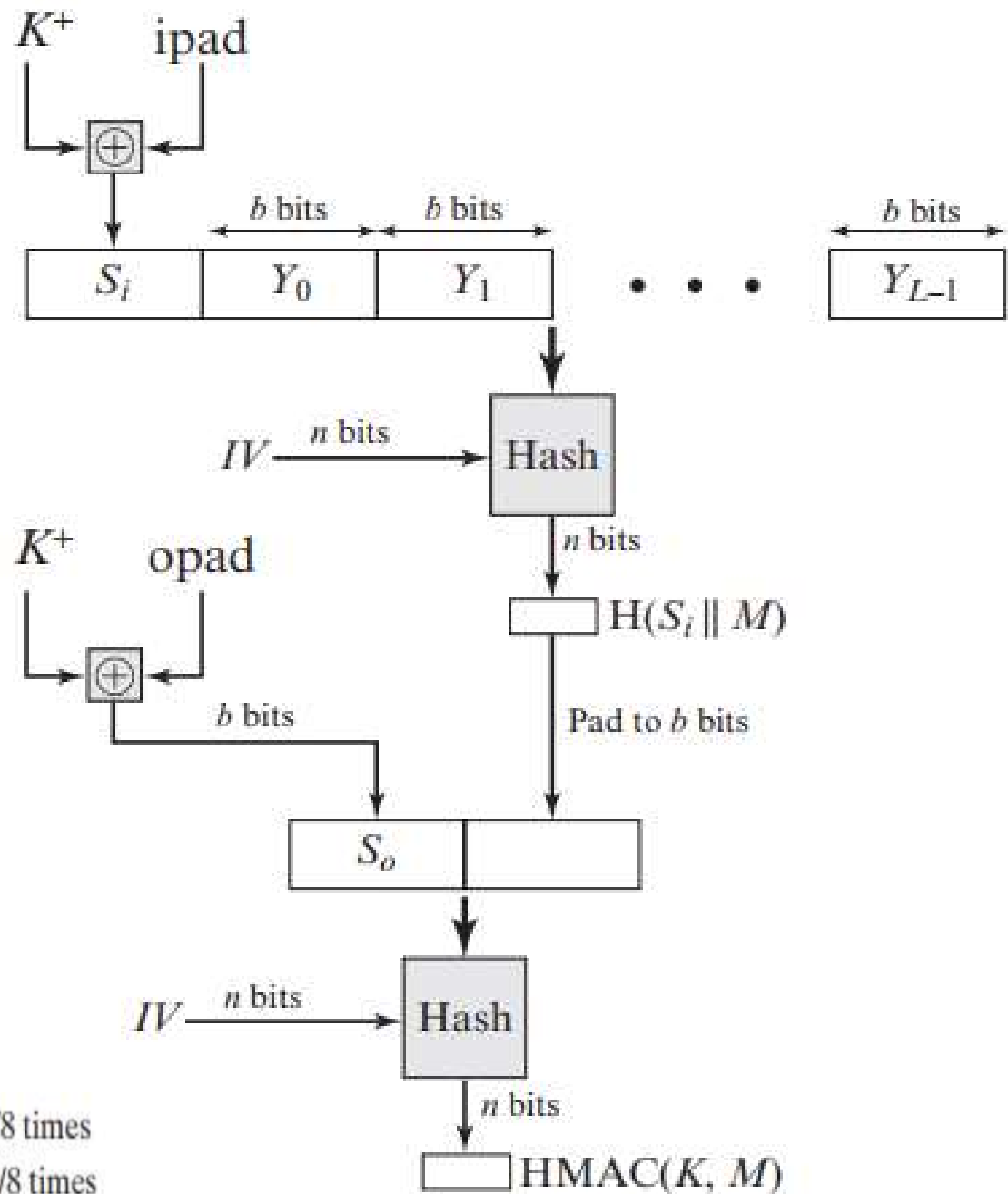
Cryptanalysis

- Cryptanalytic attacks seek to exploit some property of the algorithm to perform some attack other than an exhaustive search.
- An ideal MAC algorithm will require a cryptanalytic effort greater than or equal to the brute-force effort.
- There is much more variety in the structure of MACs than in hash functions, so it is difficult to generalize about the cryptanalysis of MACs.

MACs Based on Hash Functions: **HMAC**

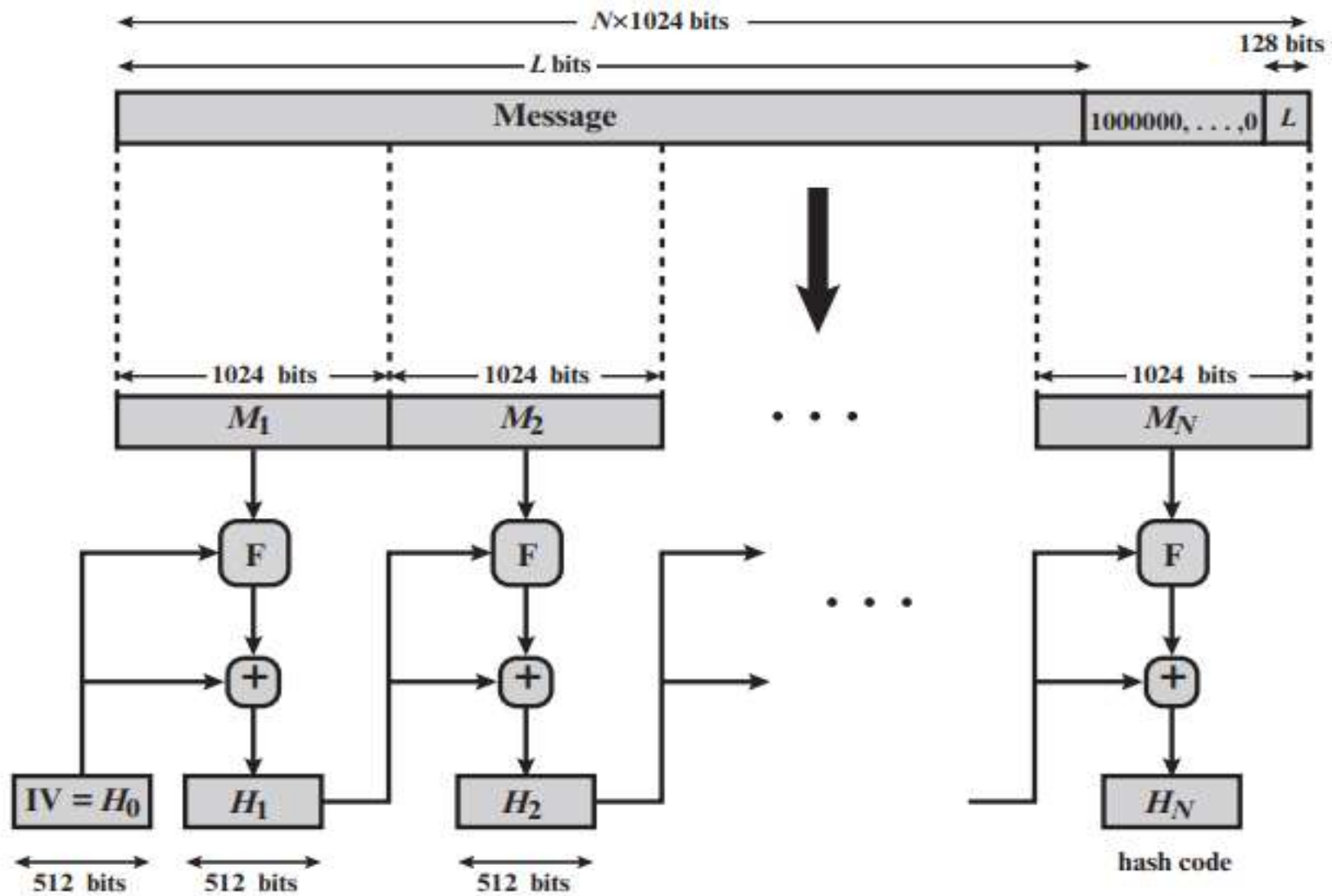
- There has been increased interest in developing a MAC derived from a cryptographic hash function.
- Motivations:
 - Cryptographic hash functions such as MD5 and SHA generally execute faster in software than symmetric block ciphers such as DES.
 - Library code for cryptographic hash functions is widely available.
- HMAC has been chosen as the mandatory-to-implement MAC for IP security.
- Has also been issued as a NIST standard (FIPS 198).

HMAC Structure



$\text{ipad} = 00110110$ (36 in hexadecimal) repeated $b/8$ times
 $\text{opad} = 01011100$ (5C in hexadecimal) repeated $b/8$ times

Figure 12.5 HMAC Structure



$+$ = word-by-word addition mod 2^{64}

Figure 11.9 Message Digest Generation Using SHA-512

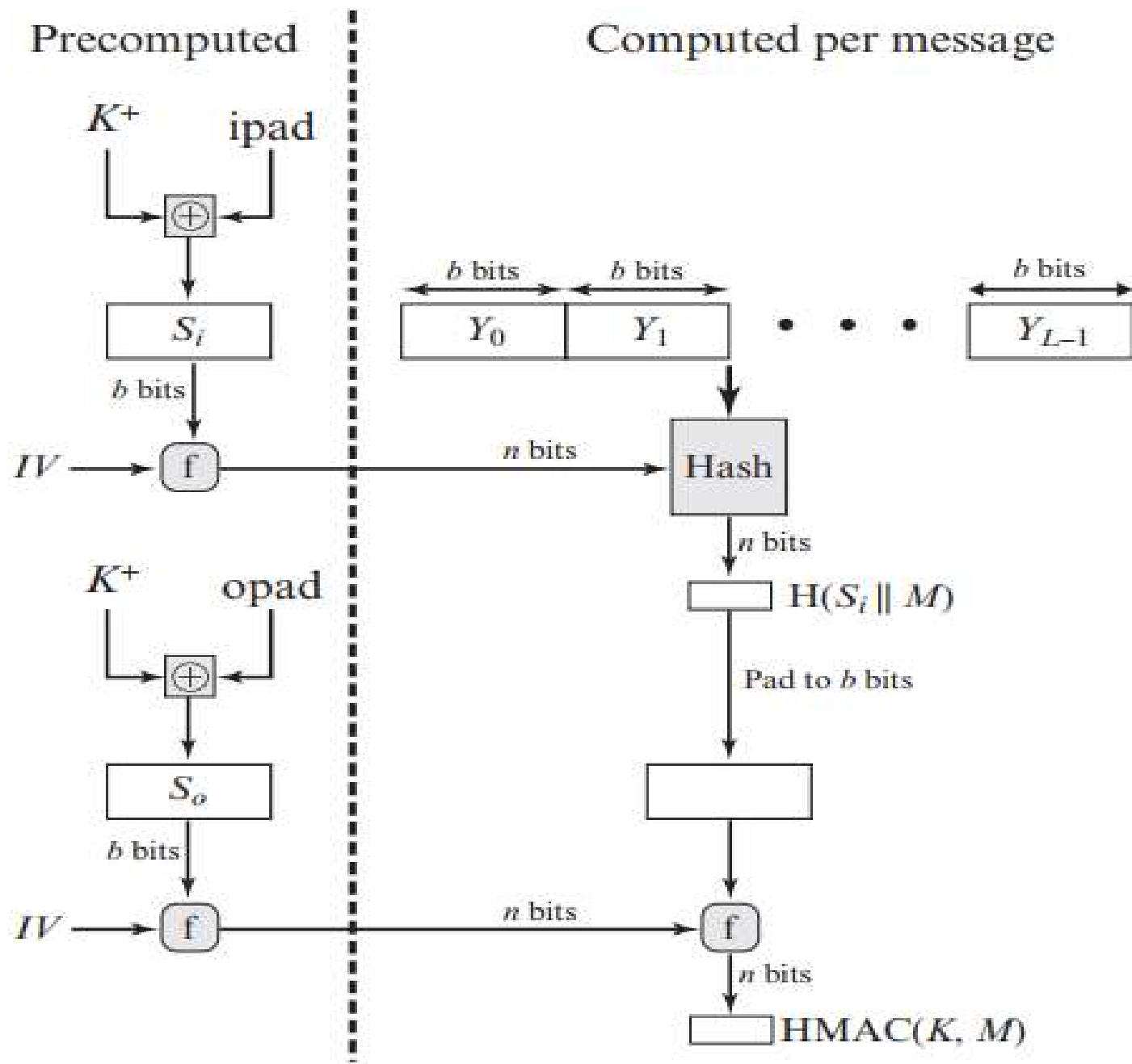


Figure 12.6 Efficient Implementation of HMAC