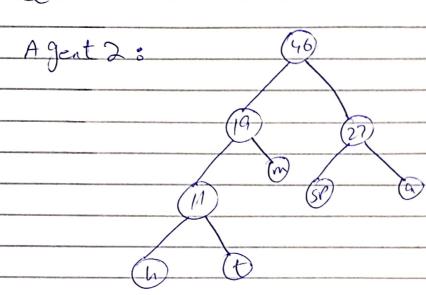
184-1182	
Q # 1:	
Time Complexity:	
	i v No. & itration
	0
$= N^2 - N(N+1)$	1 2
2	2 3
$if loop (largest i^m) \rightarrow N^2 - (N(N+1))$	3 4
$\frac{1}{2} \log \left(\frac{1}{2} \log (1) \log (1) \right) \right) \right) \right) \right) \right) \right)}{\log \log (1)}} \right) \right)} \right) \right)} \right) d d d d d d d d d d d d d d d d d d $	3 4 14 5
Swap l. > N-2	
	N-2 N-1
$= 1 + N^2 - \frac{N}{2} (N + 1) + N^2 - \frac{N}{2} (N + 1) + N^2$	2
	Sn=2 (20,+(n-1)d)
$= 2N^2 - 2N(N+1) + N-1$	d
2	2 N (2(1) +N-1)
Now Total itrations	2
	= N (N+1)
2 2 N2 - ON (N+1) + N-1+3	2
= 2N2-N(N+1)+N+2	
= O(N2). Ans:	
(b) For better sol to improve	running time is
heap sorting.	
Time complexity of heap	corting will be
alogia.	3-1111

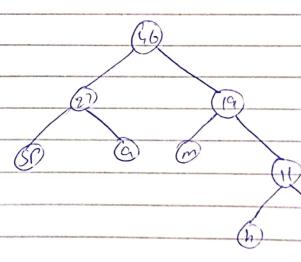
Q# 3

-



a: 11 h: 000 m: 01 t: 001 Sp: 10

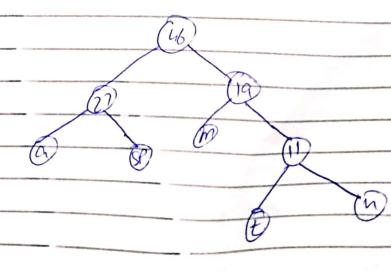
Agent 3:



α : 01 μ : 110 m : 10 t : 111 SP: 00

(F)

Agent 1:



9:00 h:111 m:10 t:110 sp:01 As All code are correct, show in above three diagrams or trees and I also write huffman codes according to trees drawn by me. So all codes are correct.

Q#4:		182-1182
i) Nodes A B C D E F	8 8 8 8	A O A O B 3 B 3
A O B 3 C 5 D 7 E 9 F 9		A 0 A 0 B 3 B 3 C \$5 -> C 5 -> D 7 E 10 E 9 F 013 F 9 Nodes Final-Cost A 0 B 3 C 5 D 7 E 9 F 9

182-1182

Q#4 (b)

Yes, If we add an edge into minimum spanning tree it will be no more mst because new edge can entroduce a cycle & can increase cost as MST already at minimum cost and we are not well known about new edge so we can't add new edge to minimum spanning tree.

But in case of man spanning tree new edge increase the cost overall of a tree.

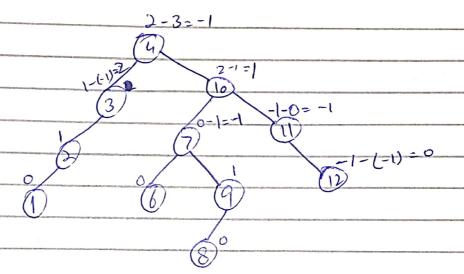


182-1182

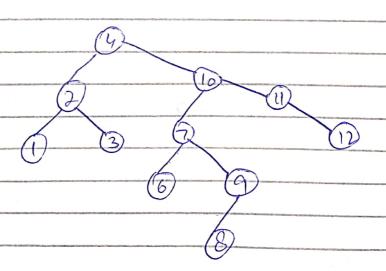
Wert - word werd

0#7

(a) Removing 5



Right Rotate





RL (4,7,10)

