SOP optimization with simulate annealing

* Problem Description:

The Sequential Ordering Problem (SOP) with precedence constraints consists of finding a minimum weight Hamiltonian path on a directed graph with weights on the arcs and on the nodes, subject to precedence constraints among nodes.

* Instances Description:

Instances are provided by TSPLIB that it is a library of sample instances for the TSP (and related problems like SOP, ATSP, HCP) from various sources and of various types.

Each instance file consists of two part as **specification part** that contains information about the instance data and **data part**.

* Algorithm Description:

The algorithm designed base on the *LUCA MARIA GAMBARDELLA & MARCO DORIGO* that described in related paper as *(An Ant Colony System Hybridized with a New Local Search*

*for the Sequential Ordering Problem).*

**Constructive heuristic** used for generating initial solution in the way that from the bingeing each time minimum possible length edge based on precedence condition selected.

For neighboring method to move from current solution to another, **Lexicographic Search** using **forwarding and backwarding path-preserving-3-exchange** applied.

The only difference is that in this algorithm lexicographic search doesn’t applied on whole search space by iteratively change the parameters “h, i, j”, instead random “h” generated and according to that random “i,j” created to do the search.

With the use of loop with size half of dimension forward and with same size loop backward exchanging applied.

It means that in each simulated annealing iteration best solution selected from a list of solutions with size of problem dimension.

* + Initial Constructive heuristic
  + O(dimension/2) forward searching with random “h, i, j” parameters.
  + O(dimension/2) backward searching with random “h, i, j” parameters.
  + Selecting the best from search as next solution

Algorithm time complexity:

For updating the temperature 2 methods (**logarithmic** and **exponential**) applied to find the best to work with.

* Algorithm Results:

T = 1

ALPHA = 0.8 (for using in temperature updating)

TEMP\_MODE = EXP (temperature updating method)

INIT\_HEURISTIC = True (using initial heuristic)

NUM\_ITERATIONS = 500

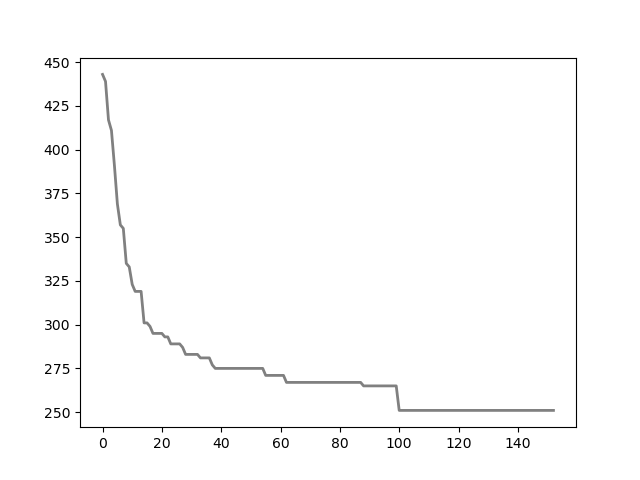
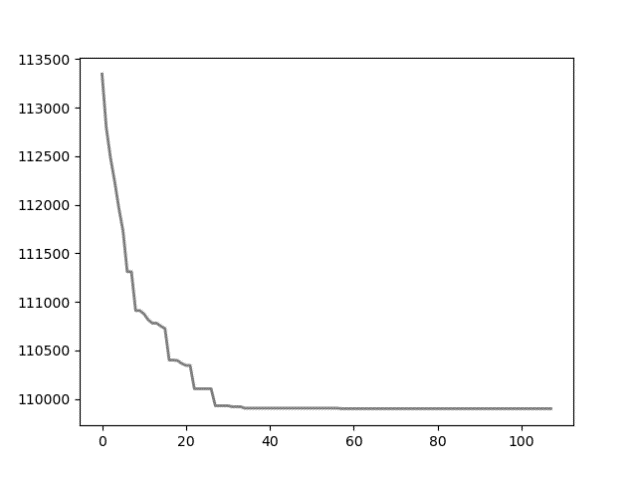
* + Algorithm progress plot for sample instances:



p43.4.sop jpeg.4753.54.sop

The whole results (main, max, avg) came in table in excel file named “Results”.

* Initial methods comparison:
  + Random



*p43.4.sop* *jpeg.4753.54.sop*

* + Heuristic:



*p43.4.sop* *jpeg.4753.54.sop*

* Temperature update methods comparison:
* Comparison with BKSs:
* Algorithm analysis:

**Strength:**

**Weakness:**

Our code has two main methods:

*static int LUPdecompose(int size, Type \*\*A, int \*P);*

which return LU matrix in A and permutation matrix in P.

*static int LUPinverse(int size, int \*P, Type\*\* LU, Type \*\*B, Type \*X, Type \*Y);*

which return invers of matrix in A in LU.

* Compiling:

Space complexity of algorithm is ***O(n2)*** which for large size of ***n*** may cause problem due to default stack size per application as ***2MB*** in my OS and compiler base config.

Because of that I preserve more space for stack size to prevent segment fault of code that cause sudden execution termination at the start of running.

|  |
| --- |
| gcc -Wall -pg lup\_matrix\_inverse.c -o0 -o int\_500\_out.exe |

-o: specify output exe file name

-o0: without optimizing

-pg: Generate extra code to write profile information suitable use gprof.

* System Information’s:

CPU: core i5 8th generation

RAM: 8GB

OS: windows 10

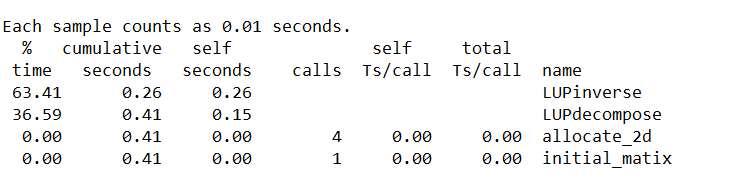
Cache: 1L = 256KB, 2L = 1MB, 3L = 6MB

* Performance profiling with Gprah:

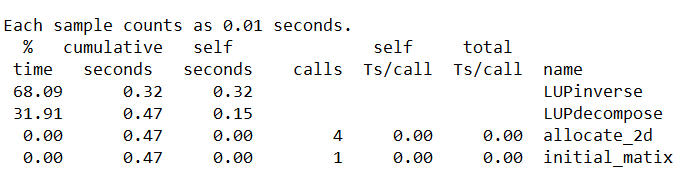
|  |
| --- |
| gprof int\_500\_out.exe > int\_500\_profile-data.txt |

(500 \* 500) Matrix

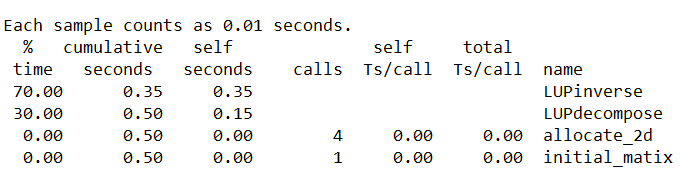
**int data type:**



**float data type:**

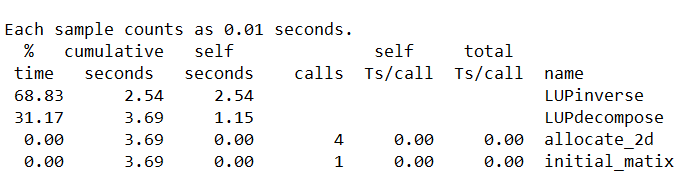


**double data type:**

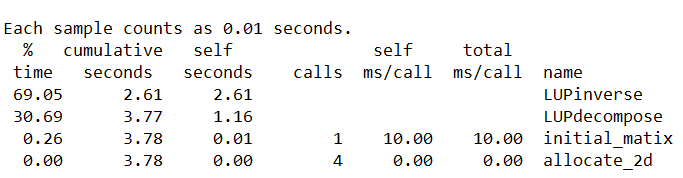


(1000 \* 1000) Matrix

**float data type:**



**double data type:**

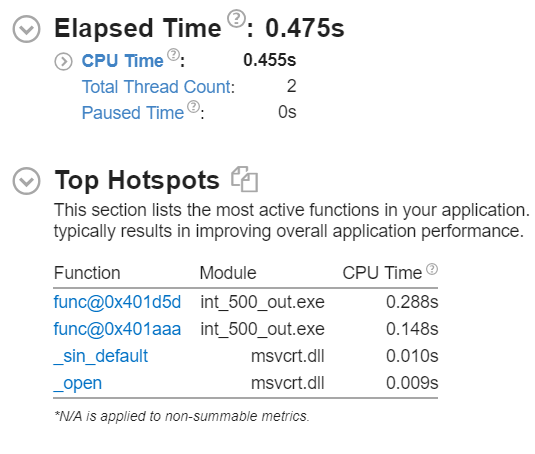


* Performance profiling with Vtune:

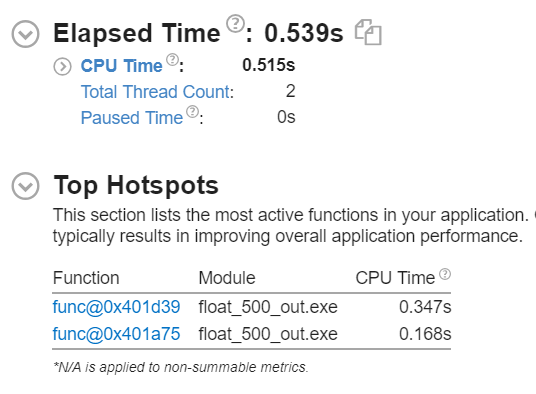
1st function address always is ***LUPinverse*** , 2nd function address always is ***LUPdecompose***

(500 \* 500) Matrix

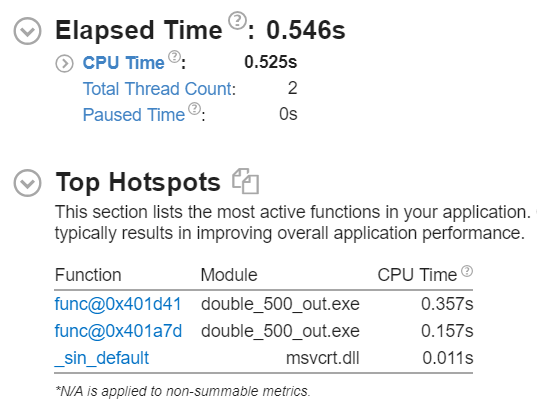
**int data type:**



**float data type:**

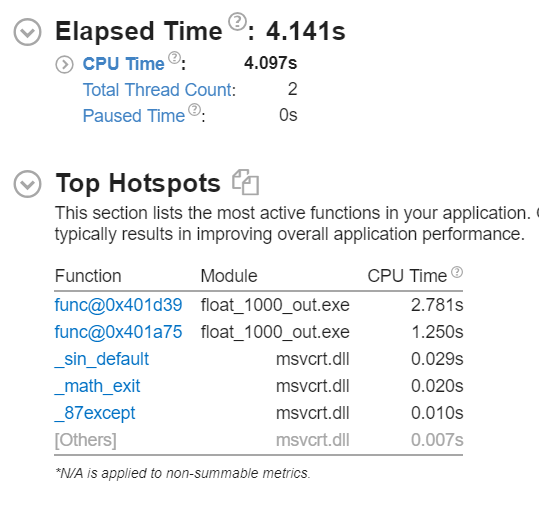


**double data type:**

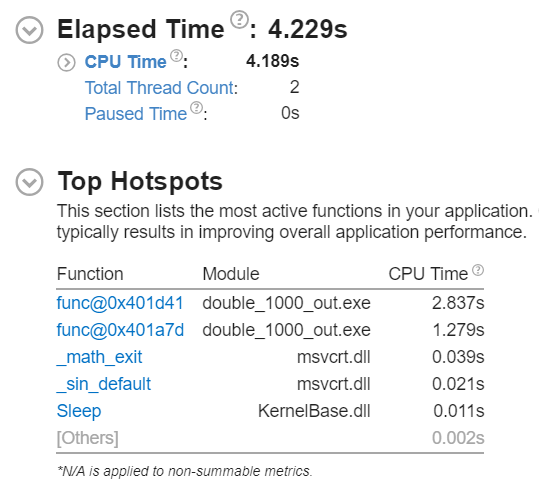


(1000 \* 1000) Matrix

**float data type:**

****

**double data type:**



* Execution analysis:

Program has two main method as:

* ***LUPdecompose:***

With time complexity as O(n2) based on code reviewing and space complexity as O(n+n2) =O(n2).

But according to the algorithm documents, LU decomposition can be computed in time O(M(n)). M(n) ≥ na where a > 2. It means O (n2.376)

* ***LUPinverse:***

With time complexity as O(n3) and space complexity as O(3n+2n2) =O(n2)

And according to code result on the used machine, size (in bytes) and precision (in number of decimal digits) of

float: 4 and 6,

double: 8 and 15

* Proposed improvement:

As matrix size and comparison accuracy increases the execution times growth.

Matrix space always is a good application for parallelism due to high ILP property of matrix-based problems.

As the result shows ***LUPinverse*** takes the most of execution time.

This method solving mathematical equationiteratively over each column.

We can divide this work over multi thread tasks that independently solve equation for specific column vector and in this way make the code much faster.

For another method ***LUPdecompose*** that take O(n2) time we can split the outer loop in specific sizes and pass them to some thread and make it faster.in other word the partial pivoting process that is comparison-based process across each column could be done in parallel manner.