Code

```
MODULE main
VAR
  proc1
           : process task(x);
           : process task(x);
  proc2
  x:0..200;
ASSIGN
  init(x) := 0;
---EF p = exists some path (E) that eventually in the future satisfies p.
---EG p = exists some path (E) that condition p is continuously true.
SPEC EG (x = 0)
SPEC EG (x = 1)
SPEC EG (x = 2)
SPEC EF (x = 0 \& proc1.counter = 100 \& proc2.counter = 100)
SPEC EF (x = 1 \& proc1.counter = 100 \& proc2.counter = 100)
SPEC EF (x = 2 & proc1.counter = 100 & proc2.counter = 100)
SPEC EF (x = 100 & proc1.counter = 100 | proc2.counter = 100 )
SPEC EF (x = 200 & proc1.counter = 100 & proc2.counter = 100)
SPEC EF (x = 201)
MODULE task(x)
VAR
  counter: 0.. 100;
ASSIGN
  init(counter) := 0;
  next(x) :=
    case
      (counter < 100) & (x < 200) : x+1;
      TRUE: x;
    esac;
  next(counter) :=
    case
      (counter < 100) : counter + 1;
      TRUE : counter;
    esac;
```

Command

NuSMV.exe hw2.smv

Results

Result shows that in **smv** implementation of the problem, there is a chance of starvation but there is no chance of accessing the shared variable at the same time.

```
-- specification EG x = 0 is true
```

None of processes start running

```
-- specification EG x = 1 is false

-- specification EG x = 2 is false

-- specification EF ((x = 0 & proc1.counter = 100) & proc2.counter = 100) is false

-- specification EF ((x = 1 & proc1.counter = 100) & proc2.counter = 100) is false

-- specification EF ((x = 2 & proc1.counter = 100) & proc2.counter = 100) is false
```

The above lines show that mutual exclusion condition is always preserved.

```
-- specification EF ((x = 100 & proc1.counter = 100) | proc2.counter = 100) is true
```

It shows that starvation condition may happen.

```
-- specification EF ((x = 200 \& proc1.counter = 100) \& proc2.counter = 100) is true
```

It shows that the ideal condition of processing can be achieved.

```
-- specification EF x = 201 is false
```