

# Lab Final Exam

Subject : Advanced Algorithm Lab  
Code : ESE 2032

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Section : 01

Semester : Spring - 22  
Batch : 45th

Date : June 21, 2022

Time : 11:00 AM

Answer to the Question NO. 04

Text in my nickname in ABC

where my Nick Name is Salem

So, the updated text will be: mynicknameisaleh  
all lowercase

all lowercase

BWT % (Rotationn)

my nickname is Saleh

§ my nickname is aleh

h § my nickname in sale

enf my nickname is sal

Let's my nickname is sa

alek & mynickname ins

Salem ♀ my nickname is

ssalen & my nickname

Issaleh Bonyekname

*Elssaleh Smymekham*

messing my milk up  
again like you did

name insalch & myrick

Now we have to sort the rotated matrix in lexi. order

Santed:

\$ my nickname is nateh  
aleh \$ my nickname is n  
ame in nateh \$ my nickname  
ekname insaleh \$ my ni  
eh \$ my nickname is nateh  
is nateh \$ my nickname  
n \$ my nickname is nateh  
ickname is nateh \$ my n  
innateh \$ my nickname  
kname insaleh \$ my nie  
leh \$ my nickname is na  
me is nateh \$ my nickname  
my nickname is nateh \$  
name is nateh \$ my nie  
nickname is nateh \$ my  
sakeh \$ my nickname is  
~~osakeh~~

ssaleh & my nickname i  
y nickname innaleh & m

o<sup>o</sup> BWT of the given text in two ways

~~best hsmilmenca~~

=hsnilmene caa & kysim

Fig. 11 shows the new field.

10. *Leucosia* *leucostoma* (Linné) *var.* *leucostoma*

What would I say if it was her?

~~4~~ ~~16517~~ ~~19517~~ 16517 18517

2000-2001-03-10

www.dreamspirit.com

Aris

1000-10000 m.s.

18. A new group of 30 students, 23 male and 7 female

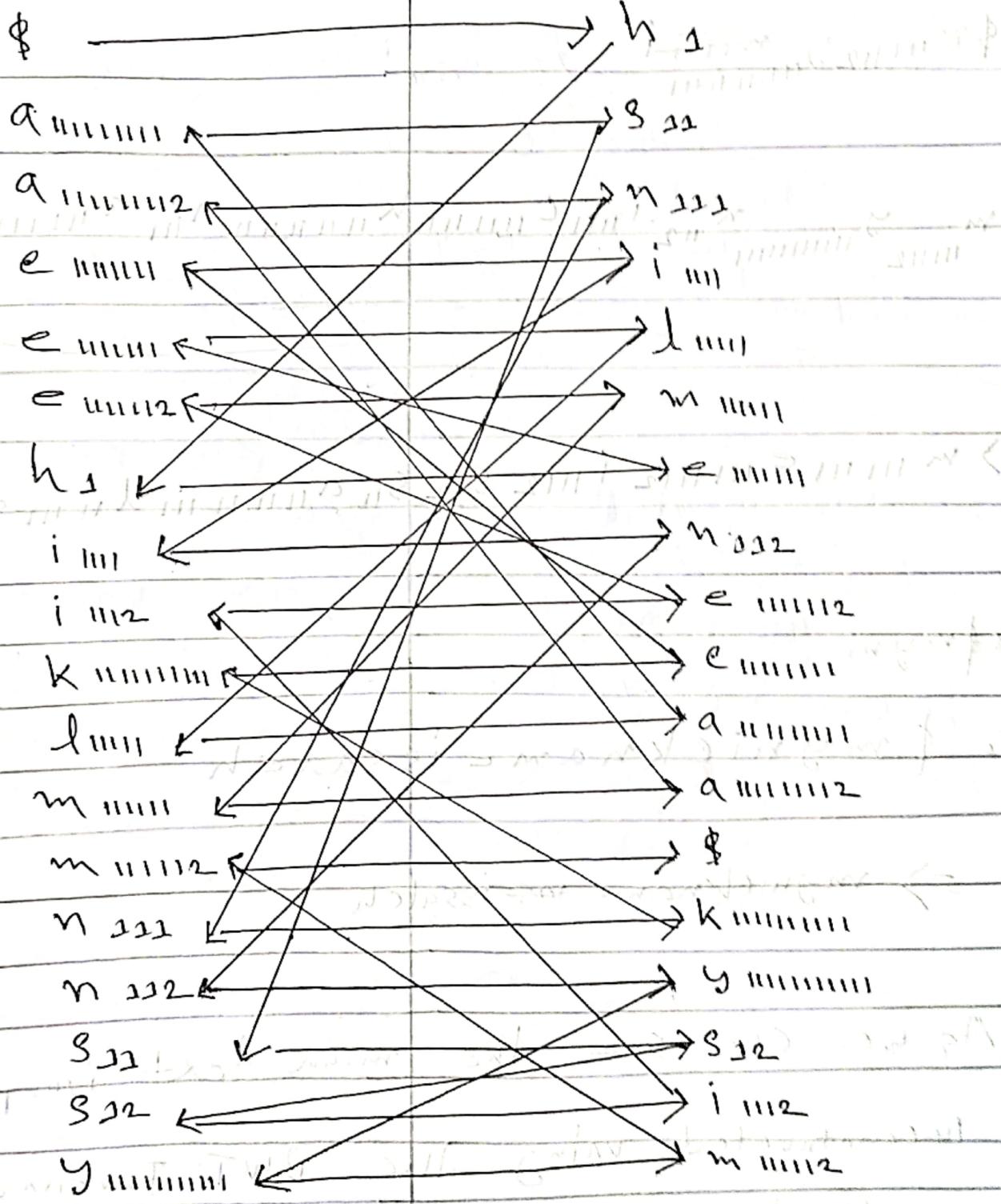
Large & well built man

*P. 1*

# Inverting BWT using LF mapping

first col.

Last col.



$\phi_m$   $\approx$   $1/2$   $y$   $\approx$   $m$

$\frac{1}{2} m_2 y_{n_2}^2 + \frac{1}{2} m_2 i_{n_2}^2 + C_{n_2} K_{n_2} + \frac{1}{2} m_2 a_{n_2}^2 )$

~~is it my mi~~

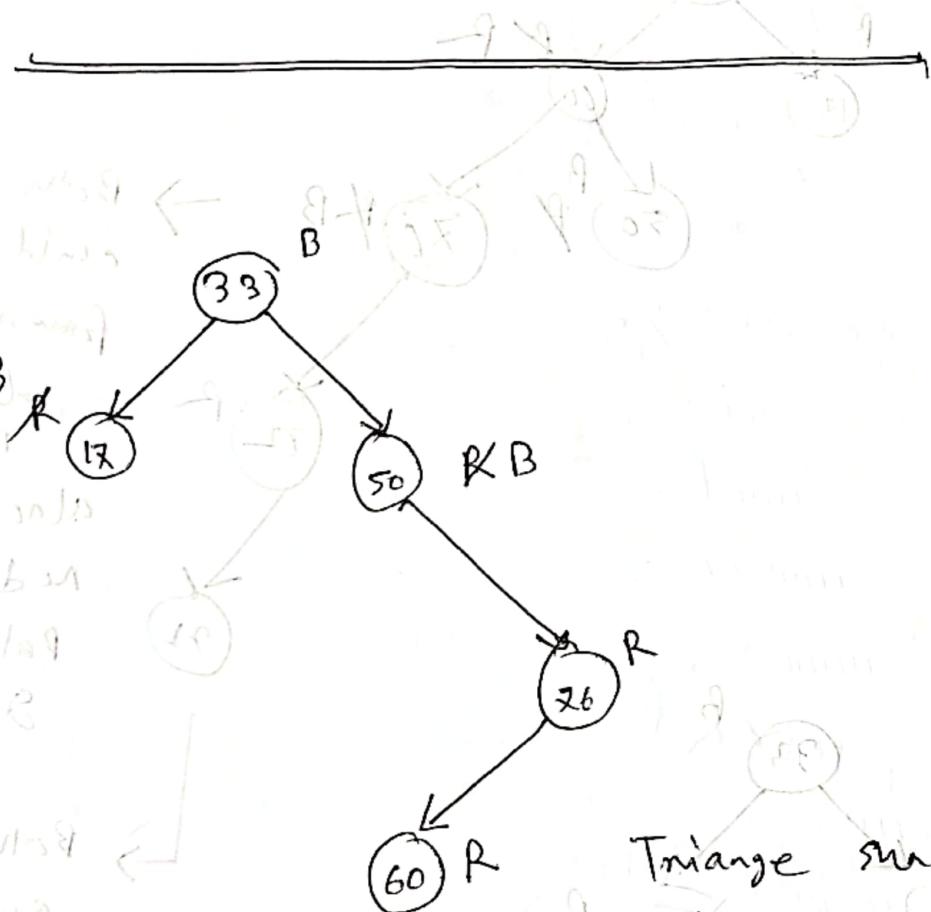
% mynickname issaleh

$\Rightarrow$  my nickname is issaleh

As we can see the main text is perfectly reconstructed using the BWT Inverse mechanism.

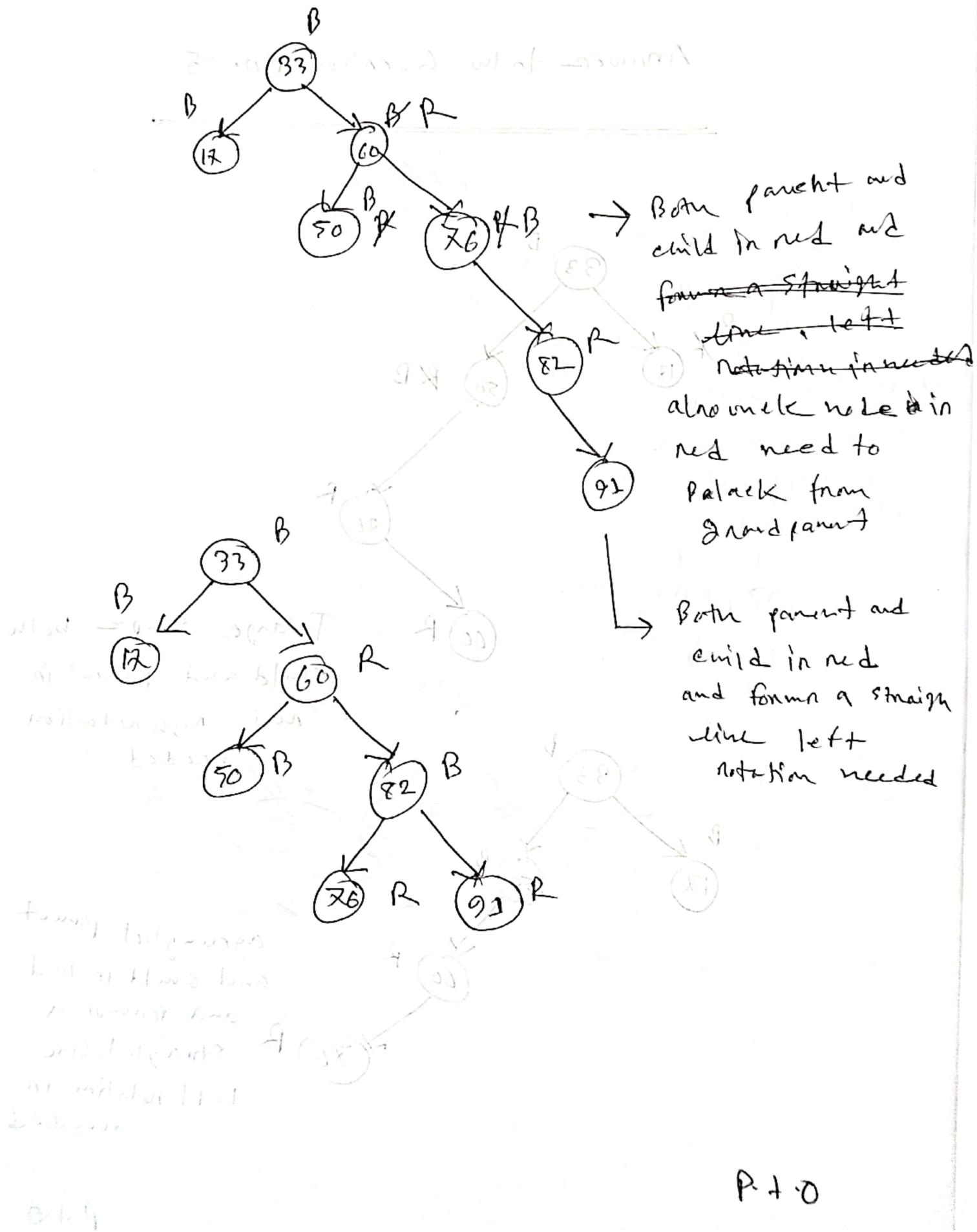
Ann:

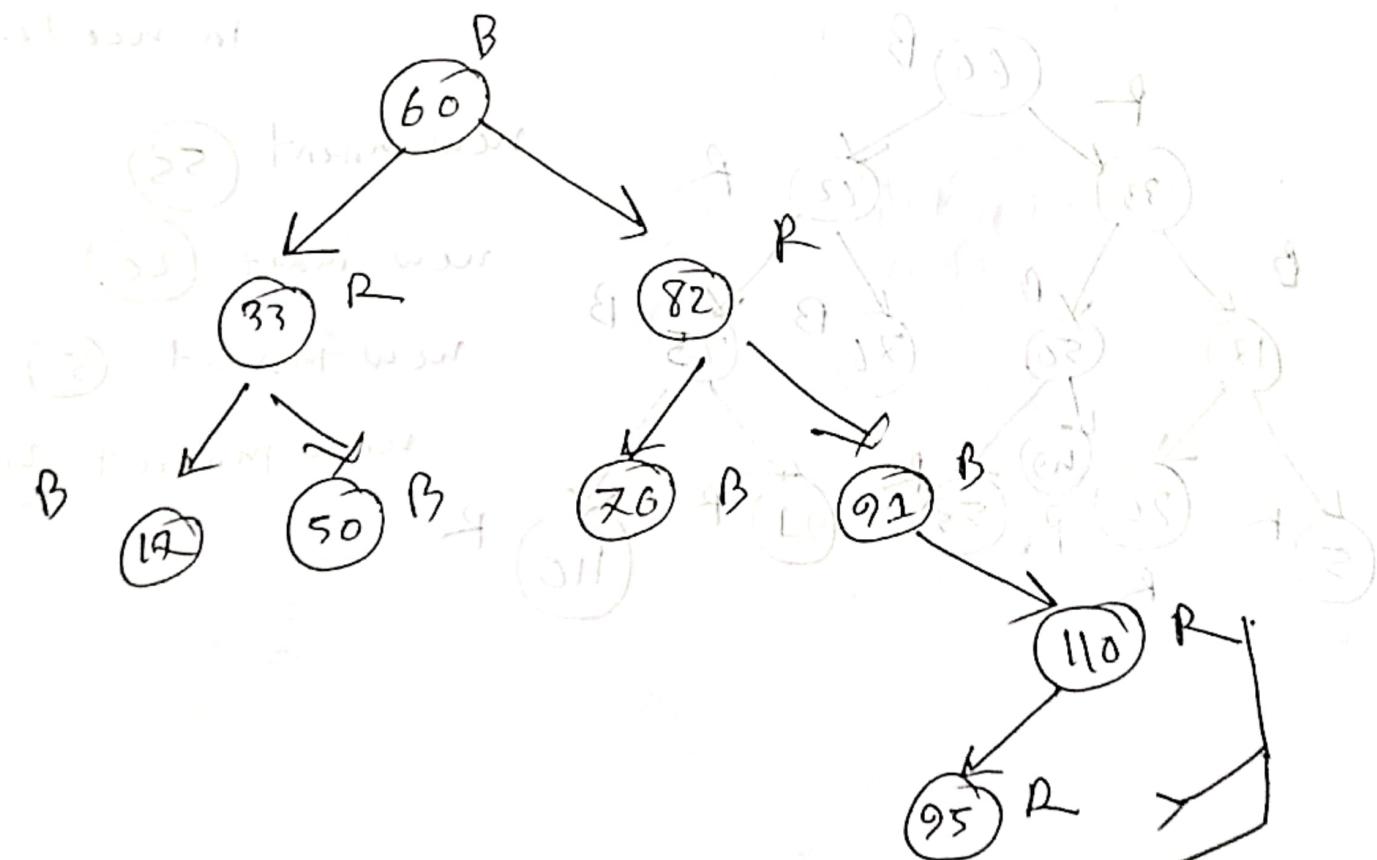
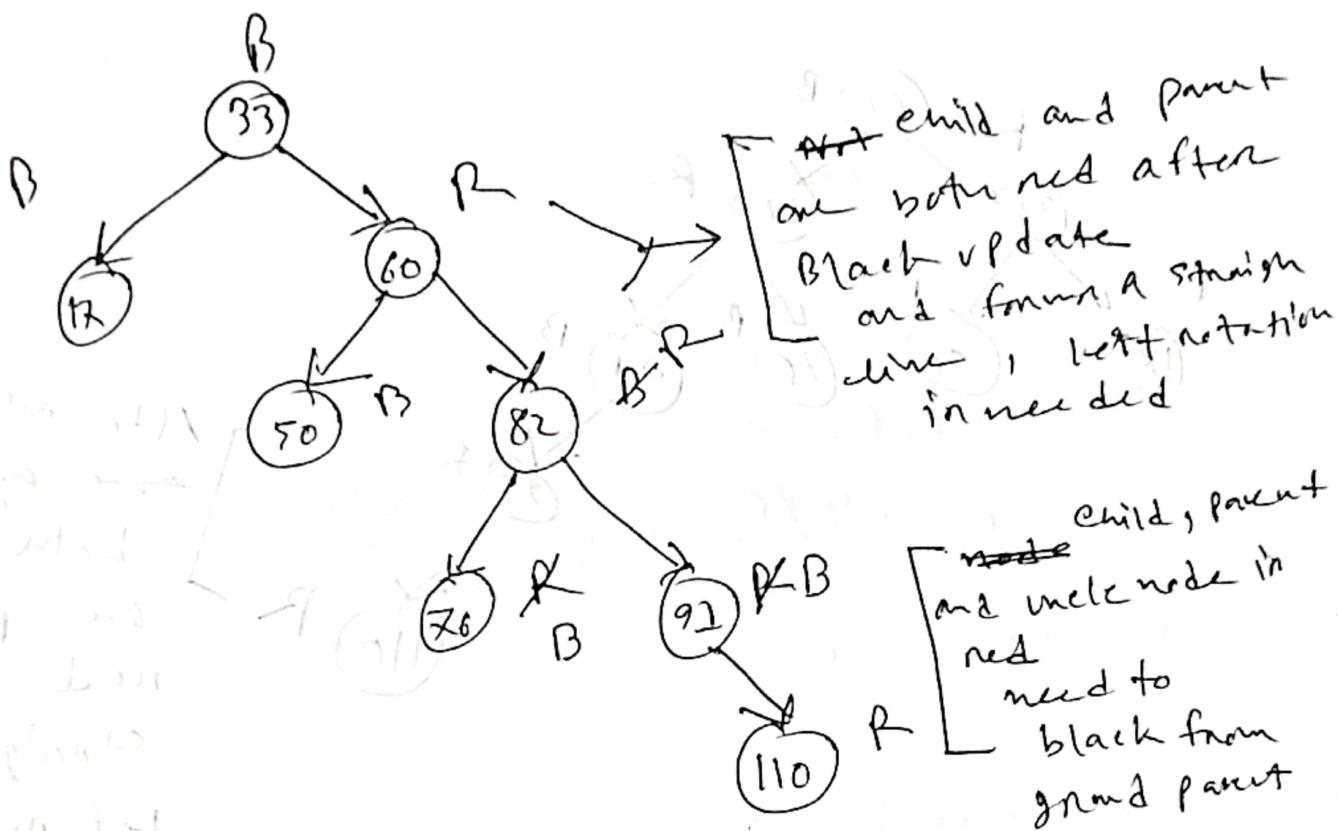
## Answer to the Question No. 05



Triangle shape both  
child and parent id  
and right notation  
needed

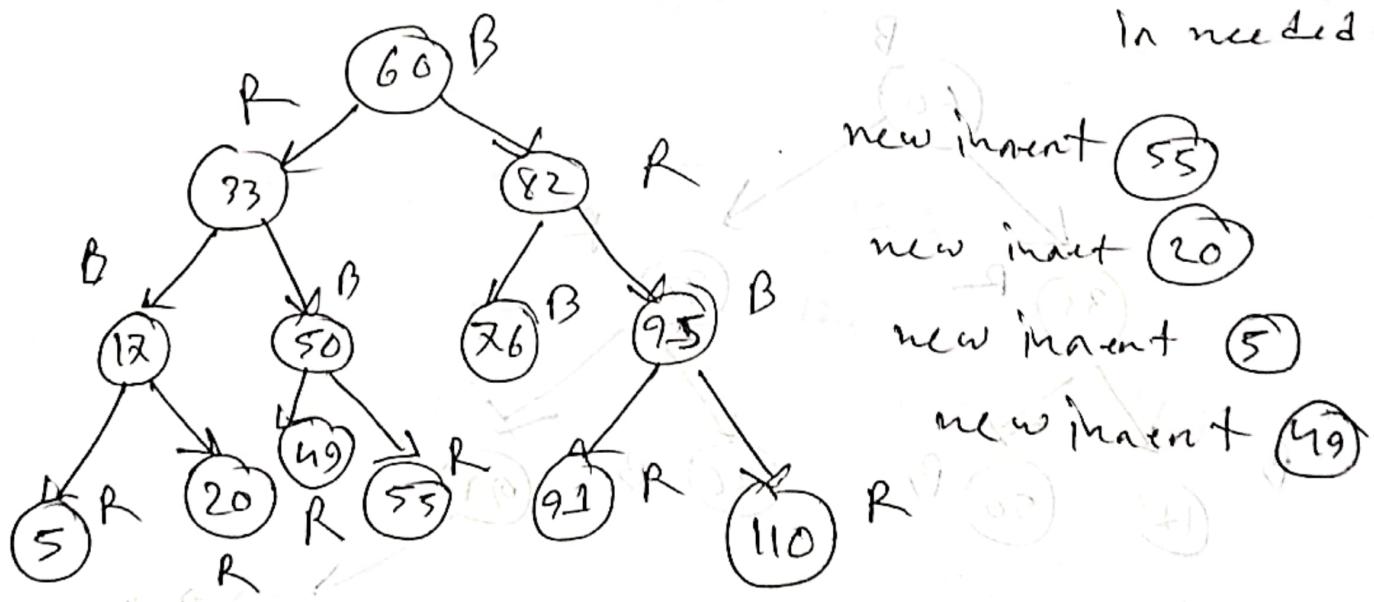
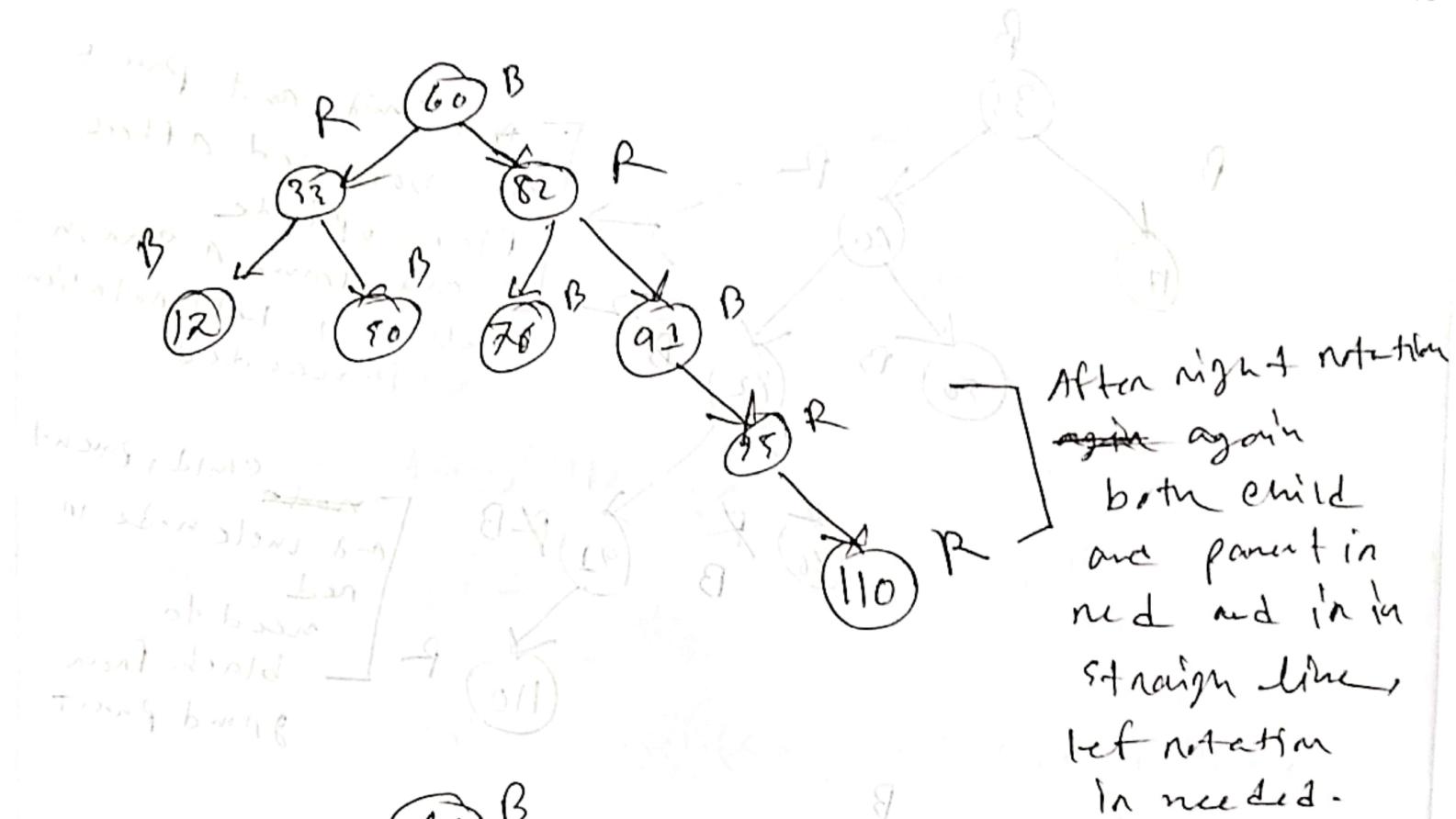
against Parent  
and child in red  
and forms a  
straight line  
left rotation in  
needed





Both child and parent is red and forms triangular shape, right rotation needed

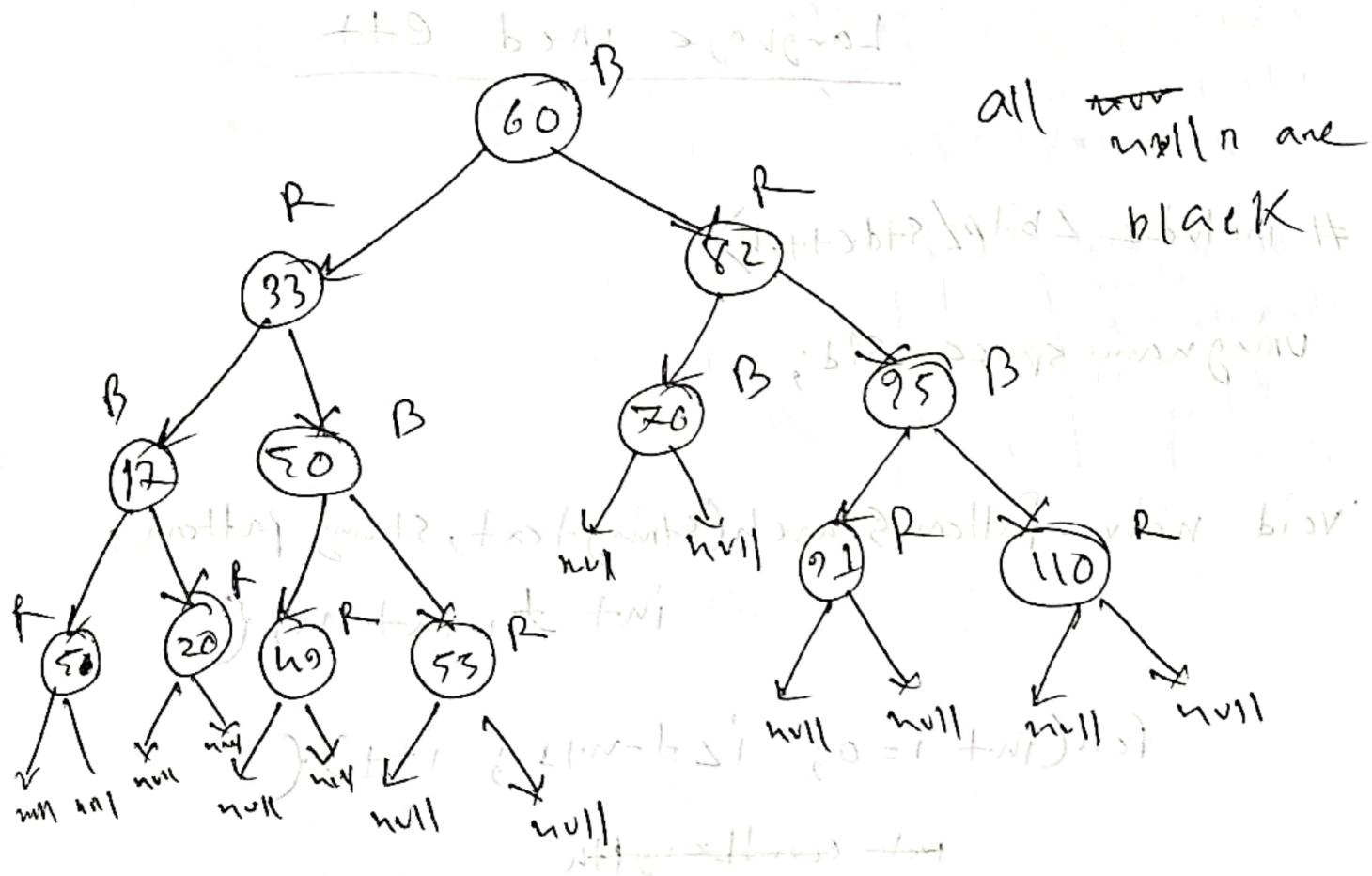
2018000000040



P. + O

# Final RB T in

Ex. on insertion with minA



$$\underline{\text{Black height}} : \text{bh}(82) = 2$$

$$\text{bh}(33) = 2$$

$$\text{bh}(17) = 1$$

$$\text{bh}(50) = 1$$

$$\text{bh}(70) = 1$$

$$\text{bh}(60) = 2$$

$$\text{bh}(95) = 1$$

$$\text{bh}(110) = 1$$

$$\text{bh}(55) = 1$$

$$\text{bh}(20) = 1$$

$$\text{bh}(5) = 1$$

$$\text{bh}(49) = 1$$

Ans of

## Answer to the Question no. 03

---

### Language used C++

```
#include <bits/stdc++.h>
```

```
using namespace std;
```

```
void naivePatternSearch(string text, string pattern,
    int t, int n) {
```

```
for (int i = 0; i < t - n + 1; i++) {
```

```
    int countLength;
```

```
    int countMatchedLength = 0;
```

```
    }
```

```
    for (int j = 0; j < n; j++) {
```

```
        if (text[i + j] == pattern[j]) {
```

```
            countMatchedLength++;
```

```
        }
```

```
    }
```

```
    break;
```

```
}
```

P.T.O.

```
if(countMatchedLength == n){
```

~~if(b[i] matching (text) n times matching given~~

```
cout << "Pattern found at position " << i << endl;
```

}

```
else{
```

cout << "Pattern doesn't match in position".  
 << i << endl;

}

}

```
int main(){
```

string text, pattern;

getline(cin, text);

cin >> pattern;

//length of text

```
int t = text.length();
```

//length of pattern

```
int n = pattern.length();
```

$\} (N = \text{number of characters})$

naivePatternSearch(text, pattern,  $i, n$ );

processes "text" by hand, finds first "pattern".  
Thus, ~~return~~

return 0;

$\}$  Big overhead from obtrusive "looping".

Thus, ~~return~~

return 0;

else

return 0;

else

return 0;

else

Am: —————  $\} O(n^2)$  time but

— : x other obtrusive part?

— : (text[i])  $\neq$  pattern[i]

— : matching of i

— : matching of i,  $\neq$  pattern[i]

— : (text[i])  $\neq$  pattern[i]  $\rightarrow$  fail

— : matching of i,  $\neq$  pattern[i]

— : (text[i])  $\neq$  pattern[i],  $\neq$  pattern[i+1]

Answer to the Question No.: 01

C++

```
#include <bits/stdc++.h>
```

```
using namespace std;
```

```
int main() {
```

```
    int i, j, n, e, val;
```

```
    cout << "Enter the value of n:" << endl;
```

```
    cin >> n;
```

```
    cout << "Enter the value of e:" << endl;
```

```
    cin >> e;
```

```
    int cm[n][e];
```

```
    for (int i = 0; i < n; i++) {
```

```
        for (int j = 0; j < e; j++) {
```

```
            cin >> cm[i][j];
```

```
        }
```

```
}
```

```

for(int i=0; i<n; i++){
    for(int j=0; j<i; j++){
        cout << cm[i][j] << " ";
    }
    cout << endl;
}

```

$VAL = CM[0][0]$   
 $i = 0;$   
 $j = 0;$

```

while((i != (n-1)) || (j != (e-1))){
    if(cm[i][j] <= cm[i+1][j])
        i++;
    else{
        j++;
    }
}

```

```

else if (i == (n-1) {
    cout << endl;
}
else {
    char m[i];
    cout << @i ("\\\"\\\"", "\\\"\\\"") ;
    cout << endl;
    VAL = cm[i][j];
    cout << endl;
}
cout << VAL << endl;
}
return 0;
}

```

- # The pseudo code indicates minimum cost path problem.

$$\text{Ans: } \underline{\underline{\text{INV}}}$$

Ann: \_\_\_\_\_

Business

## Answer to the Question No. 02

```
#include <iostream>
#include <bits/stdc++.h>
```

using namespace std;

```
int main() {
```

```
    int n, m;
    cin >> n >> m;
```

```
    int arr2D[n][m];
```

```
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++) {
            cin >> arr2D[i][j];
        }
    }
```

```
    cout << arr2D[0][0] << endl;
```

```
    int ans = arr2D[0][0];
```

```
    for (int i = 0; i < n; i++) {
```

```
        for (int j = 0; j < m; j++) {
```

```
            if (arr2D[i][j] > ans) {
                ans = arr2D[i][j];
            }
        }
    }
```

```
    cout << ans << endl;
```

```
    return 0;
}
```

Ques. If we have a 2D matrix of size  $n \times m$ , find the maximum element in it.

Ans. We can use nested loops to iterate through all elements of the matrix and keep track of the maximum value found so far.

Time Complexity:  $O(n \times m)$

Space Complexity:  $O(1)$

Optimized Solution: We can use a single pass through the matrix to find the maximum element, which reduces the time complexity to  $O(n + m)$ .

SO - 07 method of solving

for (int i = n-2; i >= 0; i--) {

    for (int j = 1; j < c; j++) {

        if (i % 2 != 0) { // ODD

            arr2D[i][j] = arr2D[n][i-1]

            + arr2D[n+1][j];

        } else if (i % 2 == 0) { // EVEN

            arr2D[i][j] = arr2D[n][i+1]

            + arr2D[n+1][j];

        } else {

            arr2D[i][j] = arr2D[n][i+1]

            + arr2D[n+1][j];

    }

// Top right most cell of the 2D array

cout << "Number of ways to reach the

top right cell from the bottom right

cell " << cout << endl; } // T.O

201200000000

return 0;  
}

Amp: