

Satisfiability Checking - WS 2016/2017

Series 5

gereon.kremer@cs.rwth-aachen.de
https://ths.rwth-aachen.de/teaching/

Exercise 1

- Give a formula describing the Unequal game instance of Series 3 Exercise 2 in *equality logic with uninterpreted functions*. Remember, that the formula shall be satisfiable iff the game instance has a solution. You must not use propositional variables in your solution!
- Compare the resulting formula to the propositional encoding. More precisely, compare the number of literals and clauses using the big \mathcal{O} notation. Draw a conclusion.
- Does the usage of equality logic with uninterpreted functions improve the runtime complexity of solving the problem? Give an explanation!

2 + 1 + 1 points

Exercise 2

In the lecture “Equalities and Uninterpreted Functions”, given on November 19th, 2015, it was mentioned that for each formula of equality logic with uninterpreted functions φ^{UF} there is an equisatisfiable formula $\hat{\varphi}^{UF}$ without constants. Define a general constant elimination procedure for a formula in equality logic with uninterpreted functions containing constants.

3 points

Exercise 3

Consider the following formula in equality logic:

$$\varphi := a = b \wedge (b = c \vee c = e) \wedge (b = d \vee c = f) \wedge a \neq e$$

- Construct the equality graph with polarity for φ .
- Simplify the constructed equality graph and the formula using the method presented in the lecture (slides 37-38).
- Make the simplified equality graph without polarity chordal. What are the chord-free simple cycles?
- Construct the satisfiability-equivalent propositional logic formula for φ using the previous results.

1 + 1 + 1 + 1 points