

**CSCI/MATH 2113 Discrete Structures**  
Chapter 8 The Principle of Inclusion and Exclusion

Alyssa Motas

April 7, 2021

## Contents

<b>1</b>	<b>8.1 The Principle of Inclusion and Exclusion</b>	<b>3</b>
1.1	Theorem . . . . .	5
1.2	Corollary: At least one condition . . . . .	6
1.3	Notation . . . . .	6
1.4	Examples . . . . .	7
<b>2</b>	<b>8.3 Derangements: Nothing Is in Its Right Place</b>	<b>10</b>
2.1	Definition . . . . .	10
2.2	Euler's totient function . . . . .	12

## 1 8.1 The Principle of Inclusion and Exclusion

Let  $S$  represent the set of 100 students. Now let  $c_1, c_2$  denote the following conditions satisfied by some of the elements of  $S$ :

$c_1$  : a student is enrolled in Writing

$c_2$  : a student is enrolled in Economics

Suppose that 35 students are enrolled in Writing and 30 of them are enrolled in Economics. We shall denote this by

$$N(c_1) = 35 \quad \text{and} \quad N(c_2) = 30.$$

If nine of the students are enrolled in both Writing and Economics, we write  $N(c_1 c_2) = 9$ . Furthermore, there are  $100 - 35 = 65$  who are *not* taking Writing and we denote this by writing

$$N(\overline{c_1}) = N - N(c_1) = 65.$$

Similarly,

$$N(\overline{c_2}) = N - N(c_2) = 100 - 30 = 70.$$

The number who are taking Writing and who are *not* taking Economics is

$$N(c_1 \overline{c_2}) = N(c_1) - N(c_1 c_2) = 35 - 9 = 26.$$

Conversely, we also have

$$N(\overline{c_1} c_2) = N(c_2) - N(c_1 c_2) = 30 - 9 = 21.$$

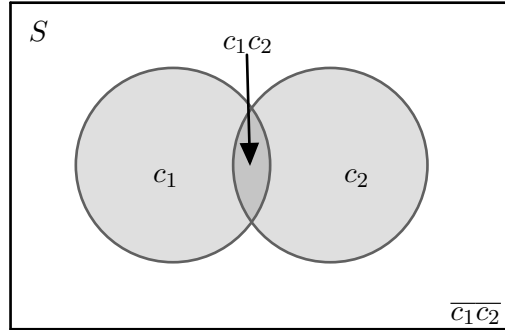
For students who are not taking Writing and Economics, we have

$$N(\overline{c_1 c_2}) = N(\overline{c_1}) - N(\overline{c_1} c_2) = 65 - 21 = 44$$

because we have  $N(\overline{c_1}) = N(\overline{c_1} c_2) + N(\overline{c_1} \overline{c_2})$ . Note that

$$\begin{aligned} N(\overline{c_1 c_2}) &= N(\overline{c_1}) - N(\overline{c_1} c_2) = [N - N(c_1)] - [N(c_2) - N(c_1 c_2)] \\ &= N - N(c_1) - N(c_2) + N(c_1 c_2) = N - [N(c_1) + N(c_2)] + N(c_1 c_2) \\ &= 100 - [35 + 30] + 9 = 44, \text{ as we saw above.} \end{aligned}$$

In diagrams, we have



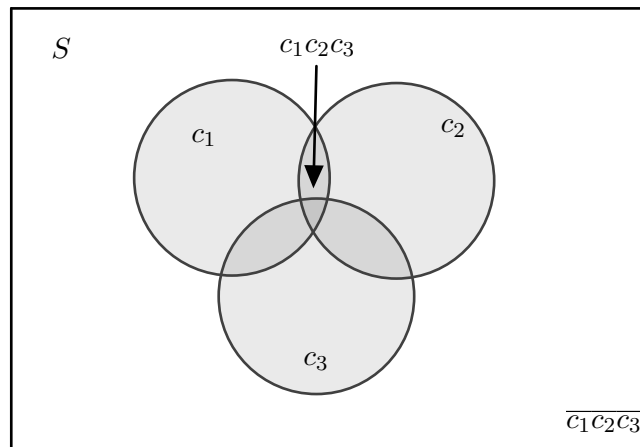
Suppose that we introduce a new condition:

$c_3$  : a student enrolled in Programming

We then have

$$N(\overline{c_1c_2c_3}) = N - [N(c_1) + N(c_2) + N(c_3)] + [N(c_1c_2) + N(c_1c_3) + N(c_2c_3)] - N(c_1c_2c_3).$$

In diagrams, we have



## 1.1 Theorem

Let  $S$  be a set with  $|S| = N$ , and conditions  $c_i, 1 \leq i \leq t$ , each of which may be satisfied by some of the elements of  $S$ . The number of elements of  $S$  that satisfy *none* of the conditions  $c_i, 1 \leq i \leq t$ , is denoted by  $\overline{N} = N(\overline{c_1 c_2 c_3} \dots \overline{c_t})$  where

$$\begin{aligned} \overline{N} = & N - [N(c_1) + N(c_2) + N(c_3) + \dots + N(c_t)] \\ & + [N(c_1 c_2) + N(c_1 c_3) + \dots + N(c_1 c_t) + N(c_2 c_3) + \dots + N(c_{t-1} c_t)] \\ & - [N(c_1 c_2 c_3) + N(c_1 c_2 c_4) + \dots + N(c_1 c_2 c_t) + N(c_1 c_3 c_4) + \dots \\ & + N(c_1 c_3 c_t) + \dots + N(c_{t-2}, c_{t-1}, c_t)] + \dots + (-1)^t N(c_1 c_2 c_3 \dots c_t), \end{aligned}$$

or

$$\begin{aligned} \overline{N} = & N - \sum_{1 \leq i \leq t} N(c_i) + \sum_{1 \leq i < j \leq t} N(c_i c_j) - \sum_{1 \leq i < j < k \leq t} N(c_i c_j c_k) + \dots \\ & + (-1)^t N(c_1 c_2 c_3 \dots c_t). \end{aligned}$$

*Proof.* We give a combinatorial proof. For  $x \in S$ , we have:

- If  $x$  satisfies none of the conditions then  $x$  contributes 1 to each side of the equality.
- If  $x$  satisfies exactly  $r$  of the conditions ( $1 \leq r \leq t$ ) then  $x$  contributes 0 to the LHS of the equality.

Considering the RHS, we get that  $x$  adds to

- 1 time in  $N$
- $r$  times in  $\sum N(c_i)$
- $\binom{r}{2}$  times in  $\sum N(c_i c_j)$
- $\binom{r}{3}$  times in  $\sum N(c_i c_j c_k)$
- $\vdots$
- $\binom{r}{r}$  times in  $\sum N(c_{i_1} c_{i_2} \dots c_{i_r})$ .

In total, the contribution of  $x$  is

$$1 - r + \binom{r}{2} - \binom{r}{3} + \dots + (-1)^r \binom{r}{r} \Rightarrow (1 + (-1))^r = 0.$$

This completes the proof. □

## 1.2 Corollary: At least one condition

Under the hypotheses of the previous theorem, the number of elements in  $S$  that satisfy at least one of the conditions  $c_i$ , where  $1 \leq i \leq t$ , is given by

$$N(c_1 \text{ or } c_2 \text{ or } \dots \text{ or } c_t) = N - \overline{N}.$$

## 1.3 Notation

To simplify the theorem above, we write

$$\begin{aligned} S_0 &= N, \\ S_1 &= [N(c_1) + N(c_2) + \dots + N(c_t)], \\ S_2 &= [N(c_1 c_2) + N(c_1 c_3) + \dots + N(c_1 c_t) + N(c_2 c_3) + \dots + N(c_{t-1} c_t)], \end{aligned}$$

and, in general,

$$S_k = \sum N(c_{i_1} c_{i_2} \dots c_{i_k}), 1 \leq k \leq t,$$

where the summation is taken over all selections of size  $k$  from the collection of  $t$  conditions. Hence  $S_k$  has  $\binom{t}{k}$  summands in it. We can also rewrite the equation above as

$$\overline{N} = S_0 - S_1 + S_2 - S_3 + \dots + (-1)^t S_t.$$

*Remark:* Let  $U_1, \dots, U_t$  be sets. Then we have

$$\left| \bigcup_{i=1}^t U_i \right|$$

is the number of elements that belong to at least one of the  $U_i$ . Thinking of  $x \in U_i$  as “ $x$  satisfies condition  $c_i$ ” then we have

$$\begin{aligned} \left| \bigcup_{i=1}^t U_i \right| &= N - \overline{N} = N - (S_0 - S_1 + S_2 - \dots + (-1)^t S_t) \\ &= S_1 - S_2 + \dots + (-1)^{t+1} S_t. \end{aligned}$$

Alternatively, we could write

$$\left| \bigcup_{i=1}^t U_i \right| = \sum_{1 \leq i \leq t} |U_i| - \sum_{1 \leq i < j \leq t} |U_i \cap U_j| + \dots + (-1)^{t+1} |U_1 \cap U_2 \cap \dots \cap U_t|,$$

or more concisely as

$$\sum_{\emptyset \neq J \subseteq 1, \dots, t} (-1)^{|J|+1} \left| \bigcap_{j \in J} U_j \right|.$$

*Remark:* When  $t = 2$ , we get the “usual” equality:

$$|A \cup B| = |A| + |B| - |A \cap B|.$$

## 1.4 Examples

- (a) How many integers  $n$  ( $1 \leq n \leq 100$ ) are not divisible by 2, 3, or 5? Let

$c_1 : x$  is divisible by 2

$c_2 : x$  is divisible by 3

$c_3 : x$  is divisible by 5

We are looking for  $N(\overline{c_1 c_2 c_3})$ . We have the following:

- $N(c_1) = \lfloor \frac{100}{2} \rfloor = 50$ ;
- $N(c_2) = \lfloor \frac{100}{3} \rfloor = 33$ ;
- $N(c_3) = \lfloor \frac{100}{5} \rfloor = 20$ ;
- $N(c_1 c_2) = \lfloor \frac{100}{6} \rfloor = 16$ ;
- $N(c_1 c_3) = \lfloor \frac{100}{10} \rfloor = 10$ ;
- $N(c_2 c_3) = \lfloor \frac{100}{15} \rfloor = 6$ ;
- $N(c_1 c_2 c_3) = \lfloor \frac{100}{30} \rfloor = 3$ .

Hence, we have

$$N(\overline{c_1 c_2 c_3}) = S_0 - S_1 + S_2 - S_3 = 100 - (50 + 33 + 20) + (16 + 10 + 6) - 3 = 26.$$

- (b) Let  $A$  and  $B$  be sets with  $|A| = m \geq n = |B|$ . Let  $A = \{a_1, a_2, \dots, a_m\}$  and  $B = \{b_1, b_2, \dots, b_n\}$ . Let us count the number of surjective functions from  $A$  to  $B$ . For  $1 \leq i \leq n$ , we define the condition  $c_i$  by

$b_i$  is not in the range of  $f$ .

Note that  $c_i$  is a condition on functions  $A$  to  $B$ . If  $f$  satisfies *none* of the conditions, then  $f$  is onto. So we are looking for  $N(\overline{c_1 c_2 \dots c_n})$ . We have

- $N(c_i) = (n-1)^m \Rightarrow S_1 = n \cdot (n-1)^m$
- $N(c_i c_j) = (n-2)^m \Rightarrow S_2 = \binom{n}{2} (n-2)^m$
- In general, for  $1 \leq k \leq n$  we have  $S_k = \binom{n}{k} (n-k)^m$ .

Therefore, we have

$$\begin{aligned}
N(\overline{c_1 c_2} \dots \overline{c_n}) &= S_0 - S_1 + S_2 + S_3 + \dots + (-1)^n S_n \\
&= \sum_{i=0}^n (-1)^i S_i \\
&= \sum_{i=0}^n (-1)^i \binom{n}{i} (n-i)^m \\
&= \sum_{i=0}^n (-1)^i \binom{n}{n-i} (n-i)^m,
\end{aligned}$$

which is the number of surjections from  $A$  to  $B$ .

- (c) In how many ways can the 26 letters of the alphabet be arranged so that none of the patterns “car,” “dog,” “pun,” or “byte” appear? Let  $c_1$  be the condition “the arrangement does contain the pattern car.” Similarly  $c_2, c_3$ , and  $c_4$  are defined for dog, pun, or byte, respectively. Then, we are looking for  $\overline{N} = N(\overline{c_1 c_2 c_3 c_4})$ . By the Inclusion-Exclusion Principle, we have

$$\begin{aligned}
\overline{N} &= N - (N(c_1) + N(c_2) + N(c_3) + N(c_4)) \\
&\quad + (N(c_1 c_2) + N(c_1 c_3) + N(c_1 c_4) + N(c_2 c_3) + N(c_2 c_4) + N(c_3 c_4)) \\
&\quad - (N(c_1 c_2 c_3) + N(c_1 c_2 c_4) + N(c_1 c_3 c_4) + N(c_2 c_3 c_4)) \\
&\quad + N(c_1 c_2 c_3 c_4).
\end{aligned}$$

Then we have the following:

- $N = 26!$
- $N(c_1) = 24! = N(c_2) = N(c_3)$  which is the number of ways we can arrange the “letter” CAR and the 23 remaining letters
- $N(c_4) = 23!$  since “byte” has 4 letters
- $N(c_1 c_2) = 22! = N(c_1 c_3) = N(c_2 c_3)$
- $N(c_i c_4) = 21!$
- $N(c_1 c_2 c_3) = 20!$

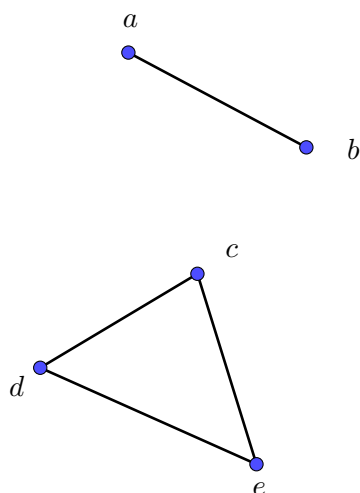


- $N(c_i c_j c_4) = 19!$
- $N(c_1 c_2 c_3 c_4) = 17!$

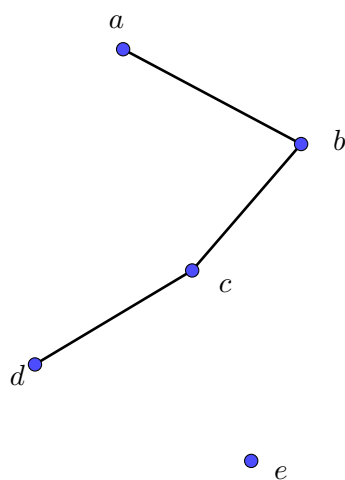
Hence, we have

$$\overline{N} = 26! - (3(24!) + 23!) + (3(22!) + 3(21!)) + (20! + 3(19!)) + 17!$$

- (d) How many arrangements contain the words “bald” and “blad”? Aside, fun problem to look at: superpermutations.
- (e) There are 5 villages. You want to devise a system of roads connecting the villages such that no village is completely isolated. In how many ways can you do this?



No village is isolated.



Village  $e$  is isolated.

Let  $S$  be the set of all (undirected, loop-free) graphs on the vertices  $\{a, b, c, d, e\}$ . We know that

$$|S| = 2^{\binom{5}{2}} = 2^{10}.$$

Now, for  $i \in \{1, \dots, 5\}$ , the condition  $c_i$  is “the system of roads isolates the  $i$ -th village.” Then for  $N(c_1)$ , we have the roads  $\{b, c\}, \{b, d\}, \{b, e\}, \{c, d\}, \{c, e\}, \{d, e\}$ .

Thus

$$N(c_1) = 2^6.$$

Similarly,  $N(c_i) = 2^6$  for  $1 \leq i \leq 5$ . Reasoning in the same way, we find that  $N(c_1c_2) = 2^3$  and similarly  $N(c_ic_j) = 2^3$ . Finally,  $N(c_ic_jc_k) = 2^1$  and  $N(c_ic_jc_kc_l) = 2^0 = N(c_1c_2c_3c_4c_5)$ . In total, we have

$$N(\overline{c_1c_2c_3c_4c_5}) = 2^{10} - \binom{5}{1}2^6 + \binom{5}{2}2^3 - \binom{5}{3}2^1 + \binom{5}{4}2^0 - \binom{5}{5}2^1.$$

## 2 8.3 Derangements: Nothing Is in Its Right Place

A professor wants the students to grade the assignments. In this case, we would like every student to receive a single assignment and not their own. This is a derangement.

Write  $[n] = \{1, 2, \dots, n\}$ . A *derangement* of  $[n]$  is a permutation of  $[n]$  such that no element is left in place.

$[2] \Rightarrow 21$  since the “standard” order is 12

$[3] \Rightarrow 312, 231$

$[4] \Rightarrow 4123 \quad 3142 \quad 2143$

$3412 \quad 4312 \quad 2413$

$2341 \quad 3421 \quad 4321$

### 2.1 Definition

We write  $d(n)$  for the number of derangements of  $[n]$ .

*Proposition:* We have

$$d(n) = \sum_{k=0}^n (-1)^k \cdot \frac{n!}{k!}$$

*Proof.* Write  $T$  for the set of all permutations of  $[n]$ . Then we have  $|T| = n!$ . Now, let  $T_i, 1 \leq i \leq n$ , be the collection of permutations that fix  $i$ . For example,  $[3] \Rightarrow 213 \in T_3$ . Then

$$d(n) = n! - \underbrace{\left| \bigcup_{i=1}^n T_i \right|}_{\text{all the perm. that fix } i}$$

By Inclusion-Exclusion, we have

$$\left| \bigcup_{i=1}^n T_i \right| = \sum_i |T_i| - \sum_{i < j} |T_i \cap T_j| + \sum_{i < j < k} |T_i \cap T_j \cap T_k| - \dots + (-1)^{n+1} \left| \bigcap_i T_i \right|.$$

- For  $\sum |T_i|$ , we have  $|T_i| = (n-1)!$  hence

$$\sum_i |T_i| = n \cdot (n-1)!$$

- For  $\sum_{i < j} |T_i \cap T_j|$ , we have  $|T_i \cap T_j| = (n-2)!$  hence

$$\sum_{i < j} |T_i \cap T_j| = \binom{n}{2} (n-2)!$$

Hence we have

$$\begin{aligned} \left| \bigcup_{i=1}^n T_i \right| &= n \cdot (n-1)! - \binom{n}{2} (n-2)! + \binom{n}{3} (n-3)! - \dots \\ &= \sum_{k=1}^n (-1)^{k+1} \binom{n}{k} (n-k)! \end{aligned}$$

Therefore, we have

$$\begin{aligned} d(n) &= n! - \left| \bigcup_i T_i \right| \\ &= n! - \sum_{k=1}^n (-1)^{k+1} \binom{n}{k} (n-k)! \\ &= n! + \sum_{k=1}^n (-1)^k \binom{n}{k} (n-k)! \\ &= \sum_{k=0}^n (-1)^k \binom{n}{k} (n-k)! \\ &= \sum_{k=0}^n (-1)^k \cdot \frac{n!}{k!} \end{aligned}$$

where  $\binom{n}{k} = \frac{n!}{k!(n-k)!}$ .

□

## 2.2 Euler's totient function

Let  $n \geq 2$ , then we have

$$\phi(n) = |\{x \in [n] \mid \gcd(x, n) = 1\}|$$

is *Euler's totient function*. For example, we have

- For  $n = 2$ , we have  $\phi(2) = 1$  which is 1.
- For  $n = 3$ , we have  $\phi(3) = 2$  which is 1 and 2.
- For  $n = 4$ , we have  $\phi(4) = 2$  which is 1 and 3.
- For  $n = 5$ , we have  $\phi(5) = 4$  which is 1, 2, 3 and 4.

*Remark:* If  $n$  is prime, then  $\phi(n) = n - 1$ .

*Proposition:* We have

$$\phi(n) = n \cdot \prod_{p|n} \left(1 - \frac{1}{p}\right)$$

where the product is taken over the primes that divide  $n$ .

*Proof.* Given  $n$ , we write it as

$$n = P_1^{e_1} P_2^{e_2} \dots P_t^{e_t}$$

where  $P_i$  are primes and  $e_i \geq 1$ . Suppose, for simplicity, that  $t = 4$ . Let  $S = [n]$  and let  $c_i$  be the condition “ $x$  is divisible by  $P_i$ ” for  $1 \leq i \leq 4$ . Then

$$\phi(n) = N(\overline{c_1 c_2 c_3 c_4}).$$

Now,

- $N(c_i) = \frac{n}{P_i}$
- $N(c_i c_j) = \frac{n}{P_i P_j}$
- $N(c_i c_j c_k) = \frac{n}{P_i P_j P_k}$
- $N(c_1 c_2 c_3 c_4) = \frac{n}{P_1 P_2 P_3 P_4}$

Then, we have

$$\begin{aligned}
\phi(n) &= S_0 - S_1 + S_2 - S_3 + S_4 \\
&= n - \left( \frac{n}{P_1} + \cdots + \frac{n}{P_4} \right) + \left( \frac{n}{P_1 P_2} + \cdots + \frac{n}{P_3 P_4} \right) \\
&\quad - \left( \frac{n}{P_1 P_2 P_3} + \cdots + \frac{n}{P_2 P_3 P_4} \right) + \frac{n}{P_1 P_2 P_3 P_4} \\
&= \frac{n}{P_1 P_2 P_3 P_4} (P_1 P_2 P_3 P_4 - (P_2 P_3 P_4 + \cdots + P_1 P_2 P_3) \\
&\quad - (P_4 + \cdots + P_1) + 1) \\
&= \frac{n}{P_1 P_2 P_3 P_4} ((P_1 - 1)(P_2 - 1)(P_3 - 1)(P_4 - 1)) \\
&= n \left( \frac{P_1 - 1}{P_1} \cdot \frac{P_2 - 1}{P_2} \cdot \frac{P_3 - 1}{P_3} \cdot \frac{P_4 - 1}{P_4} \right) \\
&= n \cdot \prod_{i=1}^4 \left( 1 - \frac{1}{P_i} \right),
\end{aligned}$$

which finishes the proof. □