

# Computability Theory

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## Computability Theory.

**Motivation.** What kind of problems can we solve with a computer?

**Automata** An automaton is a mathematical model of a computer.

**Why build models?**

- Mathematical simplicity
- Intellectual robustness  
(i.e., generalization)

**Types of models**

- Finite automata  
(finite resources, actually buildable machines)
- Turing machines  
(upper bounds for what could be built)

## Formal Language Theory

**Alphabet**  $\Sigma$ , set of possible characters that are valid

**String** Finite sequence of characters drawn from  $\Sigma$

The empty string  $\epsilon$  has no characters.

## Language

A language is a set of strings

We say  $L$  is a language over  $\Sigma$

iff it is a set of strings over  $\Sigma$

The set of all strings composed from

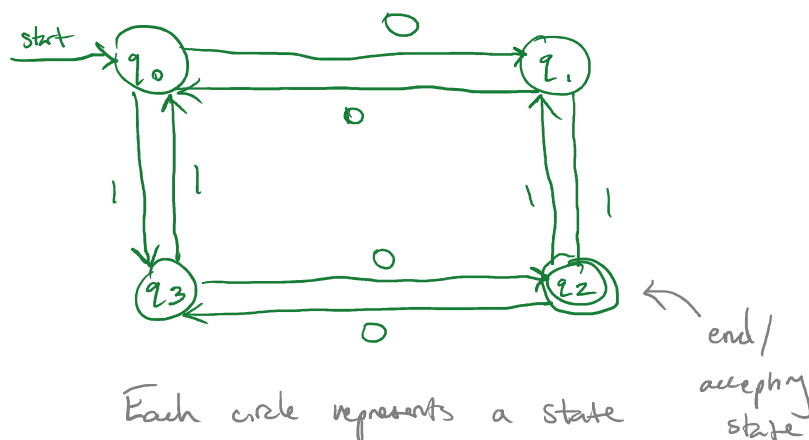
$\Sigma$  is  $\Sigma^*$ .

$$L \subseteq \Sigma^*$$

## Finite Automata

**Defn.** A finite automaton is a simple type of mathematical machine for determining whether a string is contained within some language.

**Ex.**



Each circle represents a state

Each arrow is a transition

Automaton operates on string and  
returns yes or no

(Creating chars from string sequentially)

Outputs yes if last char lands on accepting state

0 1 0 1 1 0

$q_0 \rightarrow q_1 \rightarrow q_2 \rightarrow q_3 \rightarrow q_0 \rightarrow q_3 \rightarrow q_2$

Language

The language of an automaton is the strings it accepts.

If  $D$  is an automaton over the alphabet  $\Sigma$ ,  
then  $L(D)$  is defined as

$$L(D) = \{w \in \Sigma^* \mid D \text{ accepts } w\}$$

DFA

A DFA is a Deterministic Finite Automaton

For each state in the DFA, there must be exactly one transition defined for each symbol in the alphabet

There is a unique start state

There are zero more accepting states

At each point in execution, a DFA can only remember what state it is in.