

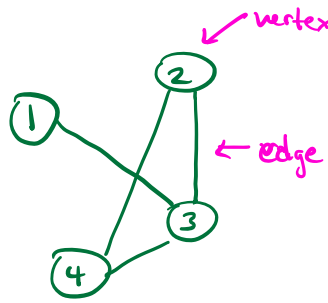
Graphs

Terminology

V - vertices (nodes)

E - edges

G - graph = (V, E)



Undirected graph: edges are bidirectional

→ Degree: number of edges touching a node

Directed graph: edges are unidirectional

→ In-degree: number of edges to a node

→ Out-degree: number of edges from a node

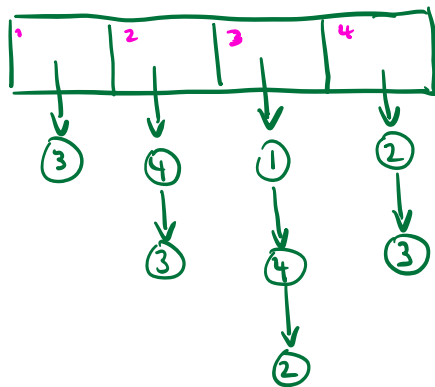
Representation

Adjacency Matrix

		dst			
		1	2	3	4
src	1	0	0	1	0
	2	0	0	1	1
	3	1	1	0	1
	4	0	1	0	0

(does not need to be symmetric for directed graphs)

Adjacency lists (using array or linked list design)



(also doesn't need to be symmetric)

Note: vertices can store other info

(i.e., attributes - name, ID addr, etc)

+ helper info for graph algs

Goal: perform operations

→ Checking for edge membership

→ querying for node neighbors

Tradeoffs:

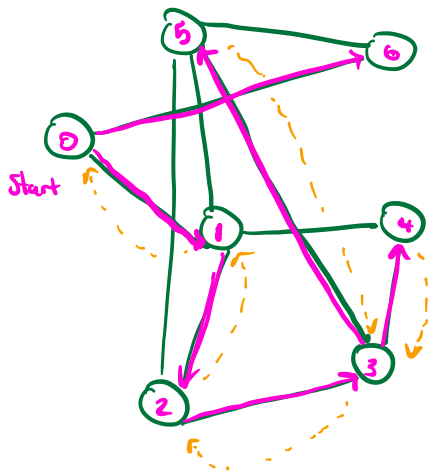
	Adj. mat.	Adj. list
Edge membership	$O(1)$	$O(V)$ ← $O(\deg(x))$ for a node x
Neighbor query	$O(V)$	$O(1)$ ← $O(\deg(x))$ if need to iterate
Space requirements	$O(V ^2)$	$O(V + E)$

Depth-First Search (DFS)

Generic algorithm for graph traversal.

Idea: explore graph w/ "string and chalk"

Example



Pseudocode

Algorithm DFS(v)

$v.visited = \text{True}$

do something with v ← generalizable!

for w in $v.neighbors$:

if $\neg v.visited$:

DFS(w)

note: can also do this iteratively with a stack
(more efficient for larger data)

Runtime

→ visit every vertex once, doing $O(1)$ work at each step

→ loop over neighbors and check visited in $O(1)$: $O(\deg(v))$

→ $\sum_{v \in V} (O(\deg(v)) + O(1)) = O(|V| + |E|) = O(|E|)$

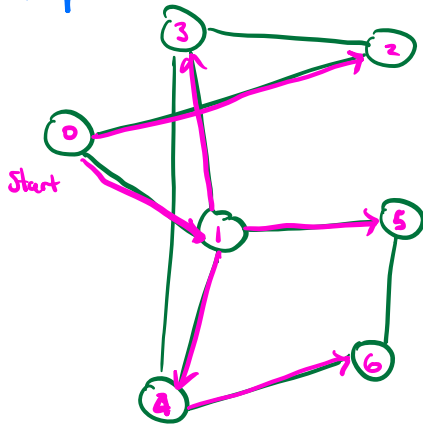
($|V| \leq |E| + 1$ for any graph)

Breadth-First Search (BFS)

Explore graphs in order of steps to reach each node

Main application: finding shortest paths.

Example



Pseudocode

Algorithm BFS (v, E)

Let Q be a queue

$Q.push(v)$ // start node

while $!Q.empty$:

$w = Q.pop()$

Mark w as visited

do something with w

for x in $w.neighbors$:

if $!x.visited$:

$Q.push(x)$

Runtime

For finding shortest path, $O(|E|)$