





Energy footprint of the Levenshtein distance computing algorithm

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Introduction

Chapter 1

Technical work

1.1 Goal

My goal is to measure energy footprints of different programming languages through execution of programs executing the same task. The results of this experiment can be very beneficial for many people in specific domains, especially in two. First, developers who tend to favour energy-efficiency over performance will know how language to use to produce the most economical program possible. Second, language designers can analyse the metrics and reach conclusions on how a design of a language can be improved or rewritten.

1.2 Overview

1.3 Architecture and design

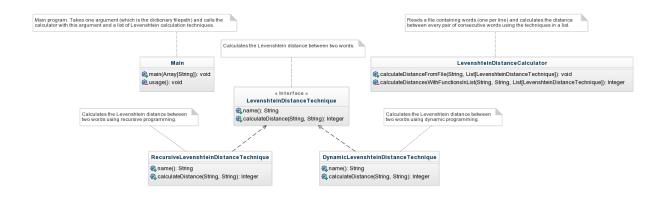


Figure 1.1: UML diagram of the algorithm in Scala

1.4 Algorithm

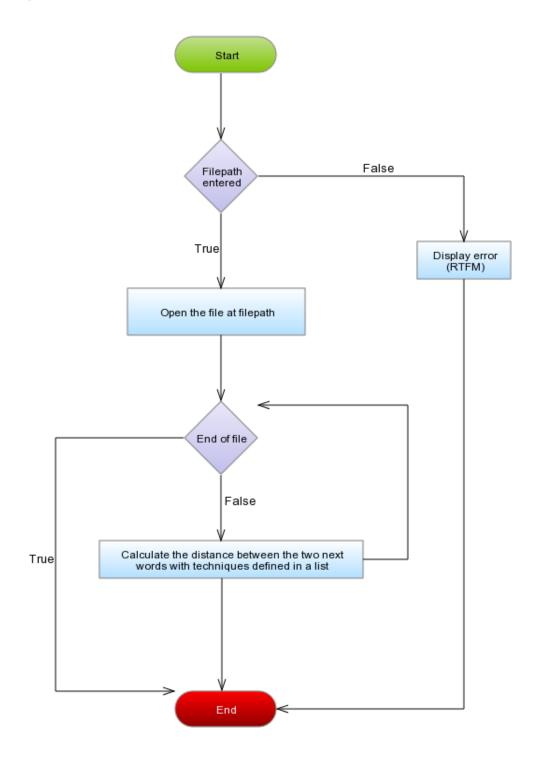


Figure 1.2: Flowchart of the algorithm

Algorithm 1: The recursive version of the Levenshtein distance computing algorithm

Function recursiveLevenshteinDistance(word1 : String, word2 : String)

```
if word1 == word2 then
\perp return \theta
lengthWord1 \leftarrow word1.length();
lengthWord2 \leftarrow word2.length();
if lengthWord1 == 0 then
| return lengthWord2
if length Word2 == 0 then
| return length Word1
firstLetterWord1 \leftarrow lengthWord1;
firstLetterWord2 \leftarrow lengthWord2;
subWord2 \leftarrow word2.substring(1);
subWord2 \leftarrow word2.substring(1);
\mathbf{if} \ firstLetterWord1 == firstLetterWord2 \ \mathbf{then}
return recursiveLevenshteinDistance(subWord1, subWord2)
else
   return 1 + min(recursiveLevenshteinDistance(subWord1, word2),
    recursiveLevenshteinDistance(word1, subWord2),
    recursiveLevenshteinDistance(subWord1, subWord2)
```

Algorithm 2: The dynamic version of the Levenshtein distance computing algorithm

```
Function dynamicLevenshteinDistance(word1 : String, word2 : String)
   if word1 == word2 then
    \perp return \theta
   lengthWord1 \leftarrow word1.length();
   lengthWord2 \leftarrow word2.length();
   table \leftarrow \inf[\operatorname{lengthWord1} + 1, \operatorname{lengthWord2} + 1];
   if lengthWord1 == 0 then
    ∟ return length Word2
   if length Word2 == 0 then
    | return length Word1
   for i \leftarrow 0 to length Word1 do
       for j \leftarrow 0 to length Word2 do
           if i == 0 then
            | table[i][j] = j
           else if j == \theta then
            | table[i][j] = i
           else if word1/i - 1/ == word2/j - 1/ then
            | table[i][j] = table[i - 1][j - 1]
           else
             [table[i][j] = 1 + min(table[i - 1][j], table[i][j - 1], table[i - 1][j - 1]]
   return table [length Word1] [length Word2]
```

1.5 Implementation

PowerAPI [1]

1.6 Usage

Chapter 2

Evaluation

2.1 Performance

Language	Execution Time		
С	1		
C++	1		
Go	1		
Haskell	1		
Java	1		
Ocaml	5m29s		
Perl (?)	1		
Python	1		
Ruby	1		
Rust	1		
Scala	45s		
Smalltalk	1		

Figure 2.1: Execution time of the algorithm in the implemented languages

2.2 Ease of use

2.3 Validation

Conclusion

Bibliography

[1] Inria Spirals Team. Powerapi, a middleware toolkit for software-defined power meters.

Official website: http://powerapi.org

GitHub: https://github.com/Spirals-Team/powerapi

 ${\rm Git} {\rm Hub} \ {\rm Wiki: https://github.com/Spirals-Team/powerapi/wiki.}$