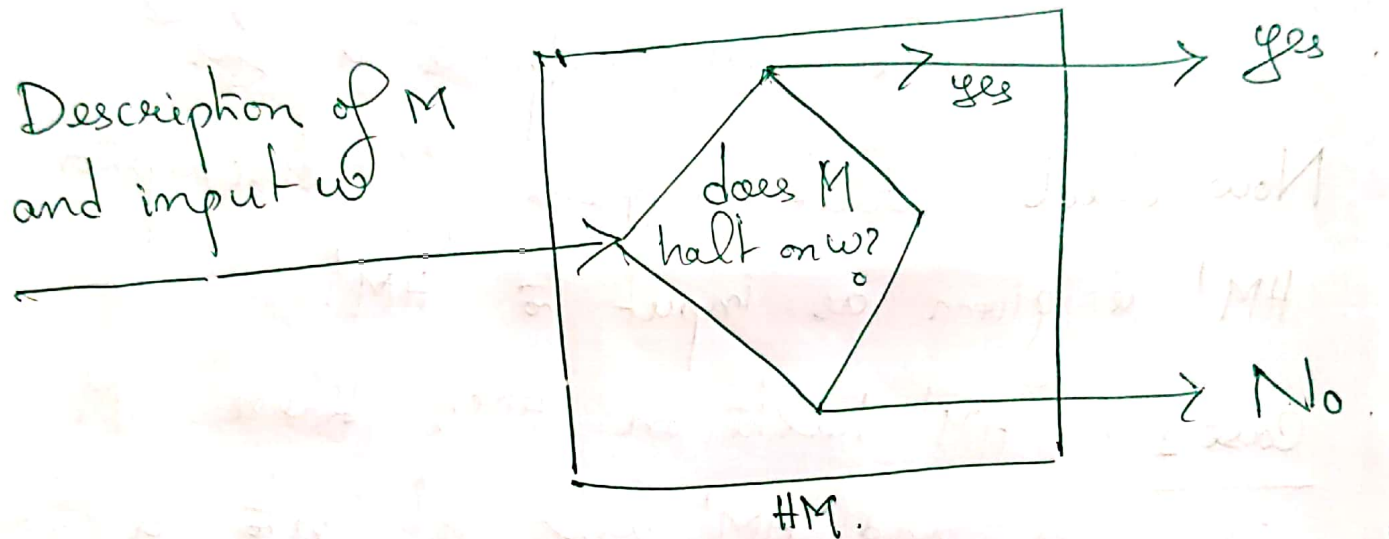


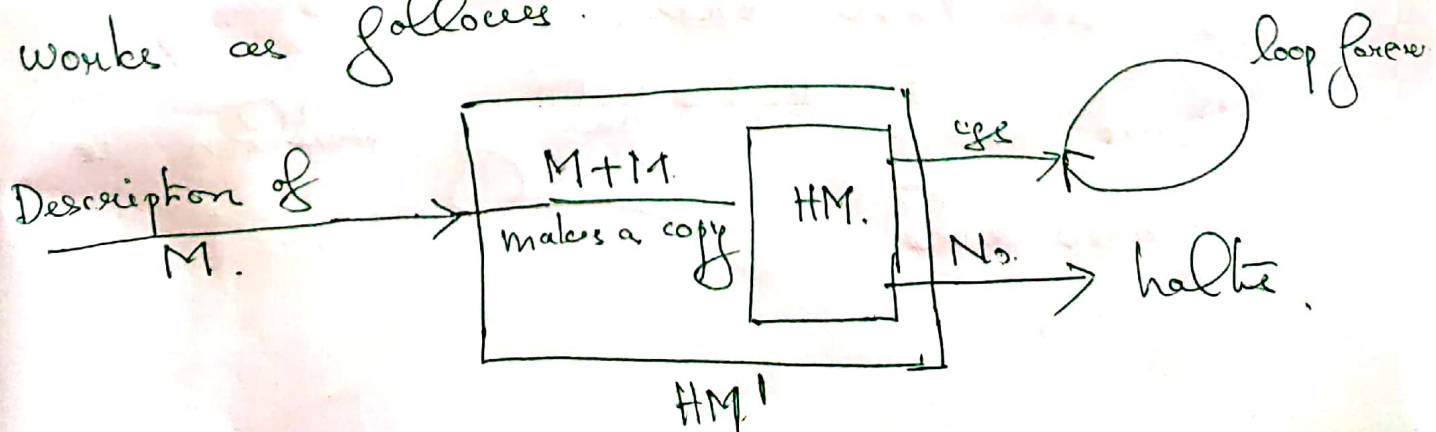
# Halting Problem Theorem

Let us assume that halting problem is decidable. Then there exist a TM HM, that solves the halting problem. The halting machine HM works as follows.



- 1) Input to HM : Description of a TM M and input string w.
- 2) Output of HM : HM says yes, if M halts on w. and HM says No, if M does not halt on w.

Now we can construct a TM  $HM'$  which works as follows.



1) Input to  $HM'$  : Description of a TM,  $M$ .

2) processing :-  $HM'$  makes a copy of its input & gives both copies to  $HM$  as input.

3) output of  $HM'$  :  $HM'$  halts when  $HM$  outputs a 'No' and  $HM'$  loops forever when  $HM$  outputs a 'Yes'.

Now check what happens if description of  $HM'$  is given as input to  $HM'$ .

Case 1 :  $HM'$  halts on  $HM'$ . Then  $HM$  will say yes and  $HM'$  will get into a loop forever.

Case 2 :  $HM'$  does not halt on  $HM'$ . Then  $HM$  will say no &  $HM'$  will halt.

In both cases, we arrive at a contradiction & hence our assumption was incorrect, i.e.,  $HM$  never exists.

Hence proved that halting problem is undecidable.