Ohm: To Study phases of compiler.

Objectiva:

-To study and understood the conject of compiler

- To study different phones of compiler

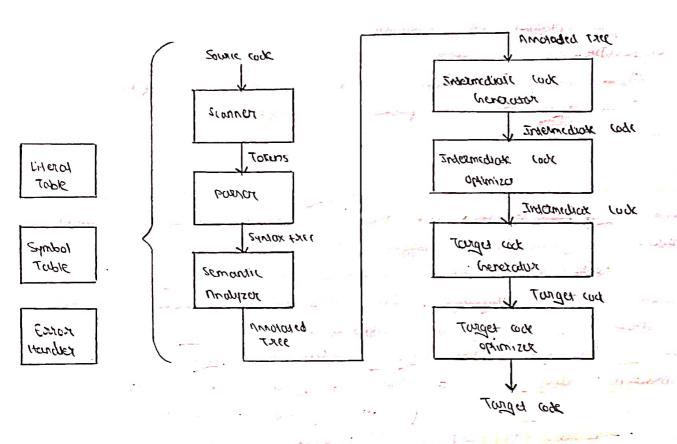
High level language (ompiler Torges program)

Exion merrage

Dudked No.1 Aim: To study phones of compiler OBJECTIVES: - To Study by underestand the concept of complex. - To study different proves of compiler In order to reduce the compexity of designing and hailding computers, nearly all of their are made to execute dekiningly simple removerings & byodraw for a combined wint pe exercise pring ph lambining these very simple commands into a program in what in could machine language, since this is a tedious and eventprince must programming in instead done wing a high-level programming language. This language can be very different from the machine language that the computer can execute , so some means of bridging the gap in required. This is where the compiler comes in a compiler transacs a program willen in a pigh-level brodzamning languages that in sinappe for primary bradyamusks into we part-level wagning randrade that in sequised by computer country this is process, the complete will also attempt to spot and report opinion by Johnson mind a yilly pend produce for brodrammind you a pride impact on how loss programs can be developed. The main recover for this one: compared to machine language the notation used by programming forguages in closer to the every human think about problems, · The compiler and spot same devious programming mistakes · Programs wither in a high-level vanguage tend to be shorter than edinimient byodyane myther is warpine pridrater.

Another advantages of using a high-level language is that the

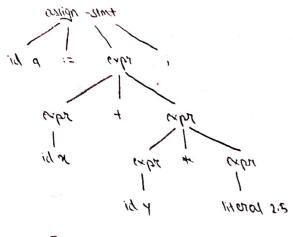
sums program can be composed to many different marking languages

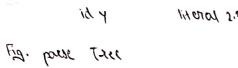




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	and, here, he brought to sun on many different martines on other
	hand programs that are unition in a high-level language
	and answapianh paintents to marine induste with sin
	Somewhard Slower than programs that are mand- wated in
	marking language were, some time - without program are still
	THE STATE OF THE THINK - CONTROL PROGRAM OR SHILL
	hurilles partly in machine language. A good compiler will,
	however, be able to get very close to the speed of hand.
	willer machine lade when translating well- translating
	bjod jans
	The Translation Process:
_	A compiler performs two major tanks:
	margara simus soft la ilylana -
	-> synatheric of the torget - language inorductions.
	3 3
	Phases of a compiler:
	Stanning
	Passing
	semantic Analysis
	Intermediate lade Invertation
	Intermediate look optimized
	Tangel Code Chenetatur
	Torge upinizer
	Thre auxiliary compenent interact with some or au phones:
	literal larghil
	symbol table
	rother korts





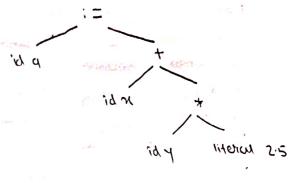


Fig. Syntax Tace

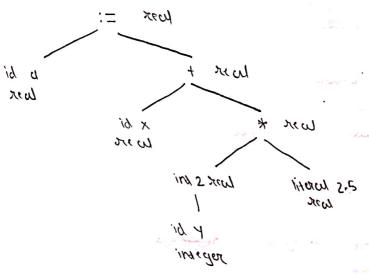


Fig: rundoded Syndon the

	Scanner:
	The scanner begins the analysis of the source program by:
•	Reading file character by character.
	Coronging characters into tokens.
	Eliminating unneeded info comments and white space.
•	Entering preliminary info. into literal or symbol table.
•	processing compiles directives by setting flags.
	Tokens represent basic program entities such as:
	Identifica literals Reserved words operators, delimiter, etc.
	Example:
	a:= x + y + 2.5; as scanned as a identifier y identifier
	:= assignment uperator, * multiplication operature, x
_	identifier, 2.5 real literal, + plus operator,
	's semicolon,
	Passes:>
	Receivers taken from the scanner.
è	Recognize the structure of program as a paint tree
•	Park true is recognized at to a writer free grammar.
0	syntax exars are reported is the program is motactically
	Inlasted.
6	A prove take us inelficient to represent the exactoric of program
	A system tree is a more condensed newsian of the purse
	\$75¢
	A spot or true in usually generated as output by the purser,
	Sementic Analyze:
•	Sementic of pragram are its meaning as opposed to syntax or structure
•	the semantic country of:
	Runting semantics - penavior of program at runting
	static semantics - enecured by the compiler. Page No.
	, , , , , , , , , , , , , , , , , , , ,

real 1 real p bi re al templ = indexicos (4) t run × bi temps := temps real + 25 $\omega_{2\kappa}$ temp 3 := > real + temp 2interal 2.5 in 2 real a := tem 3 min 1 r bi

Fig: Imamediate look generates

brouph look (HAbatyetica) Thric - addies 10ck LORDI RI, γ ;; $\epsilon_1 \leftarrow \gamma$ temp1:= ind 2 year(4) temps := temp1 xxxxx * 2.5 if fi = ind 2 xcm (RI) F1,81 mouf 11 FZ - F1 * 25 1emp 3 := x + temps MULF 62,11,25 chi= temp 3 LOADF F3 14 : F3 ← X LOOL F4, F3, F2 :: F4 - F31F2 ii a = Fy 9.44 79072

Fig: code honorodor.

	· Storic semantics Sucretes:
	· Declarations of variables and constants before use.
	e calling functions that exists (predefined in library or defined by
	the mes)
	Passing persumeters properly.
	· Type charing.
	Static semantics are difficult to check by the passer.
	Semantic analyzer does the following:
7	- checks the static semantics of languages.
	- Monotates the syntax tree with type information,
	Intermediate lade beneatation: -
j'e	comes after Syntax and smartic analysis.
	separates the compiler from end from its backend.
	To be well use sent stone of the stone of th
	Intermediate representation should have 2 into properties
	- should be easy to produce
	- sporty po sand to translate just his tardet byodran
•	intermediate representation can have a writing of forms:
	- Three address code, P-code for an abstract machine,
-	Take or DAW representation
- 15	
	(ade heneratox:->
	(nenerates code for the torget machine, typically:
	- Assembly lade, ar
	- Rejocutable machin coll
	Properties of the torget machine become a major factor
, 0	code deverator refert abbiditions warping justination
•	Alloxuses memory locations for variables.
9	phonodes registers for entermediate computations.
	Dana Ma
_	Page No. ———————————————————————————————————

conclusion:

Thus, we how studies and understand the concept of computer and phases of compiles.

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	tode Improvement:
1	lode improvement techniques can be applied to:
	- Intermebox vole - independent of the target machine.
	- Tangar rode - dependent on the tanget machine
	Intermediare rode improvement Includes:
	Constant folding
1	Elimination of common sub - expressions
	Identification and dimination of universaliable code,
	Jutroning pobs
	Improving function alls
	Targets code improvement includes:
	Allocation and use of registers
•	solution of bester (fuster) instructions and addressing mades.
	Conclusion:
	Thus we have sudies and understood the concept of
	compiler and phases of compiler.
_	
•	LING YOLF GREATION:
	which in compiler?
->	to combiles in a broducin . Hurs saige o broducin matter in
	one language - the source language and framlates it into
	equivalent program in another language - the target language.
	The compiler reports to its wer the presence of errors
	in the course program.
(3)	what are the two parts of the compilation? explain briefly.
_	Analysis and Synthesis are two parts of compilation.
	Analysis por breaks up the source program into completed
	piaces and creases on insermediate septementation of the
	Sauxe Program Page No. —
	1 ago No.
	0



11	The synthesis part constructs the desire target program from
	Define symbol Table. The is a dotal thructure used by the complex to keep trade of semantics of the variables. It stores the information about supe and building info. about names.
•	I ist the Various phones of a computer? Lexical Analyses Syntax Analyses Sometic Analyses Lude aptimises Lude aptimises Lude aptimises
- 0	List the phases that constitute the front end of a lamplete. Lexical and Syndactric analysis The creation of Symbol take semantic analysis cheneration of intermedicate look, the certain amount of look optimisation can be done by front end as well. Also includes obtain handling that goes along with each of these phases.
->	The backend of compiler includes those portions that depend on the larger machine and generally those partion do not defend on the source language. Just intermediate language. These includes— - code optimization - code generation, along with error handling and posternol-table operation