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RSA Algorithm

Dr. V. K. Pachghare

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RSA Public Key Cryptosystem  
Key Generation Algorithm

**Step 1:** Choose two random large prime numbers  $p$  and  $q$   
For maximum security, choose  $p$  and  $q$  are of about equal length,  
e.g. 512-1024 bits each.

**Step 2:** Compute the product  $n = p \cdot q$

**Step 3:** Choose a random integer  $e < m$  [Where  $m = \Phi(n) = (p-1)(q-1)$ ]  
The numbers  $e$  and  $(p-1)(q-1)$  must be relatively prime, i.e. they should not share common prime factors.  $\text{GCD}(e, m) = 1$

**Step 4:** Compute the unique inverse  $d = e^{-1} \pmod{m}$   
The equation  $d \cdot e \pmod{m} = 1$   
can be solved using the Euclidian algorithm

**Step 5:**  $C = P^e \pmod{n}$

**Step 6:**  $P = C^d \pmod{n}$

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# Encryption

$p = 3, q = 11$

$n = p \cdot q = 33$

$m = \text{Totient function } \Phi(n) = (p-1) \cdot (q-1)$   
 $= 2 \cdot 10 = 20$

The public exponent 'e' must be relatively prime to 20,  
i.e.  $e = 3, 7, 9, 11, 13, 17, 19$   
here we select  $e = 3$

Plaintext x: 9

$y = x^e \pmod{n}$

**Ciphertext:** 3

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Decryption

$$d = e^{-1} \bmod m$$
$$d = 3^{-1} \bmod 20$$
$$\cancel{d \cdot e \bmod 20 = 1}$$

So, for  $e = 3$ ,  $d = 7$

Ciphertext  $y: 3$

$$x = y^7 \bmod 33$$

Plaintext: 9

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$p = 7, q = 19, M = 6$

Find out the minimum value for e and corresponding value for d.

Encrypt the message  $M = 6$  using above e.

Also decrypt using above d to get  $M = 6$

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$p = 7, q = 19, M= 6$

Find out the minimum value for e and corresponding value for d.

Encrypt the message  $M = 6$  using above e.

Also decrypt using above d to get  $M= 6$

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We compute 'd' using Euclids algorithm

We get  $d = 65$  for  $e = 5$

Public Key: (5, 133)  
Private Key: (65, 133)

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Use RSA algorithm to encrypt the message  $M = 123$ . The parameters given are:  $p = 61$ ,  $q = 53$ ,  $e = 17$ . Also find out the value of  $d$  and decrypt the message to get  $M = 123$ .

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$3120 = 17(183) + 9$	$9 = 3120 - 17(183)$
$17 = 9(1) + 8$	$8 = 17 - 9(1)$
$9 = 8(1) + 1$	$1 = 9 - 8(1)$
$8 = 1(8)$	

$1 = 9 - 8(1)$   
 $= 9 - [17 - 9(1)](1)$   
 $= 9(2) - 17$   
 $= [3120 - 17(183)](2) - 17$   
 $= 3120(2) - 17(367)$   
 $= 3120(2) + 17(-367)$  I  
so the inverse of 17 mod 3120 = -367 = 2753

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Use RSA algorithm to encrypt the message  $M = 123$ . The parameters given are:  $p = 61$ ,  $q = 53$ ,  $e = 17$ . Also find out the value of  $d$  and decrypt the message to get  $M = 123$ .

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Decryption:

Ciphertext = 855 and the private key is (2753, 3233)

The decryption function is:

$$\text{decrypt}(c) = c^d \bmod n = c^{2753} \bmod 3233$$

where  $c$  is the ciphertext.

To decrypt the ciphertext value 855, we calculate

$$\text{decrypt}(855) = 855^{2753} \bmod 3233 = 123$$

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# Example - 4

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1 = 7 – 3(2)  
1 = 7 – 3(9-7) = 4(7) – 3(9)  
1 = 4(214-9(23))-3(9) = 4(214) – 95(9)  
1 = 4(214) – 95(223 – 214) = 99(214) – 95(223)  
1 = 99(660 – 2(223))-95(223)  
1 = 99(660)-293(223)  
so the inverse of 223 mod 660 = -293 = 367  
Therefore I  
**Private Key = (367, 713)**

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# Decryption

The ciphertext is 284

Plaintext  $P = C^d \bmod n$

$P = 284^{367} \bmod 713 = 439$

Use modulo arithmetic to compute the value of P

Message  $M = P = 439$

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## Attacks against RSA

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Math-Based Key Recovery Attacks

Three possible approaches:

1. Factor  $n = pq$
2. Determine  $\Phi(n)$
3. Find the private key  $d$  directly

All the above are equivalent to factoring  $n$

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## Knowing $\Phi(n)$ Implies Factorization

- If a cryptanalyst can learn the value of  $\Phi(n)$ , then he can factor  $n$  and break the system.
- Computing  $\Phi(n)$  is no easier than factoring  $n$ .
- In fact, knowing both  $n$  and  $\Phi(n)$ , one knows  $n = pq$
- Therefore  $q = n/p$

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$\Phi(n) = (p-1)(q-1) = pq - p - q + 1$

(Substitute  $pq = n$  and  $q = n/p$ )

$= n - p - n/p + 1$

$p\Phi(n) = np - p^2 - n + p$

shift RHS terms to LHA

$p^2 - np + \Phi(n)p - p + n = 0$

$p^2 - (n - \Phi(n) + 1)p + n = 0 = p^2 - 12p + 12$

There are two solutions of  $p$  in the above equation.

Both  $p$  and  $q$  are solutions.

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- If a cryptanalyst can learn the value of  $\Phi(n)$ , then he can factor  $n$  and break the system.
- Computing  $\Phi(n)$  is no easier than factoring  $n$ .
- In fact, knowing both  $n$  and  $\Phi(n)$ , one knows  $n = pq$
- Therefore  $q = n/p$

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Suppose the cryptanalyst has learned that  $n = 33$  and  $\Phi(n)=20$ . Find out the two factors of  $n$ .

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Solution:

$$p^2 - (n - \Phi(n) + 1)p + n = 0$$
$$p^2 - (33 - 20 + 1)p + 33 = 0$$
$$p^2 - (14)p + 33 = 0$$
$$p = 11 \text{ and } 3$$

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Suppose the cryptanalyst has learned that  $n = 84773093$  and  $\Phi(n)=84754668$ .  
Find out the two factors of  $n$ .

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Decryption attacks on RSA

RSA Problem:

Given a positive integer  $n$  (that is a product of two distinct large primes  $p$  and  $q$ ), a positive integer  $e$  (Public Key exponent) such that  $\gcd(e, (p-1)(q-1))=1$ , and an integer  $c$  (Ciphertext), find an integer  $m$  (plaintext) such that  $m^e \equiv c \pmod{n}$ . It is widely believed that the RSA problem is computationally equivalent to integer factorization; however, no proof is known.

The security of RSA encryption's scheme depends on the hardness of the RSA problem.

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# Finding d: Timing Attacks

Timing Attacks on Implementations of Diffie-Hellman, RSA, DSS, and Other Systems (1996), Paul C. Kocher

By measuring the time required to perform decryption (exponentiation with the private key as exponent), an attacker can figure out the private key

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- Possible countermeasures: –
  - use constant exponentiation time
  - add random delays
  - blind values used in calculations

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- Researchers have discovered a timing attack on RSA keys, to which OpenSSL is generally vulnerable, unless RSA blinding has been turned on.
- RSA blinding: the decryption time is no longer correlated to the value of the input ciphertext
- Instead of computing  $c^d \bmod n$ , choose a secret random value  $r$  and compute  $(r^2c)^d \bmod n$ .
- A new value of  $r$  is chosen for each ciphertext

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1 Diffie-Hellman Key Exchange Dr. V. K. Pachghare

2 Primitive Roots

3 Primitive Roots

4 How to prove that given  $\alpha$  is primitive root of  $\mathbb{Z}_p^*$ ?

5 Roots

6 But  $\alpha^n \equiv 1 \pmod p$  doesn't necessarily mean that  $\alpha$  is primitive root.

7 If  $\alpha^{p-1} \equiv 1 \pmod p$  then  $\alpha$  is primitive root of  $\mathbb{Z}_p^*$ .

8 Example: If  $p = 5$  &  $\alpha = 2$ , then  $\alpha^2 \equiv 1 \pmod 5$  but  $\alpha^4 \not\equiv 1 \pmod 5$ . So  $\alpha$  is not primitive root.

1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23

# Diffie-Hellman Key Exchange

## Dr. V. K. Pachghare

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1 Diffe-Hellman Key Exchange Dr V. A. Pachghare

2 Primitive Roots

3 Primitive Roots

4 How many primitive roots are there in mod n? That is how many primitive roots of n are there? We have to recall that the powers of n form a repeating cycle and that each power is unique. By a primitive root we mean if all roots in the group are generated by it.

5 Fermat's Little Theorem

6 But if there is a primitive root between the powers of n then all powers of n will be unique.

7 For a prime p, show that the elements of  $\mathbb{Z}_p^*$  are the primitive roots of  $\mathbb{Z}_p^*$ . Hint: If  $a$  is a primitive root of  $\mathbb{Z}_p^*$ , then  $a^k \neq 1$  for any  $k < p-1$ .

8 Example: If  $p = 11$ , then  $\mathbb{Z}_{11}^* = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ . Check that  $2$  is a primitive root of  $\mathbb{Z}_{11}^*$ . Hint: If  $a$  is a primitive root of  $\mathbb{Z}_p^*$ , then  $a^{p-1} \equiv 1 \pmod{p}$ .

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1 Diffe-Hellman Key Exchange Dr V. K. Pachghare

2 Primitive Roots

3 Primitive Roots Let  $p$  be a prime. Then  $b$  is a primitive root for  $p$  if  $b^{\frac{p-1}{d}}$  mod  $p$  is 1 for every divisor  $d$  of  $p-1$ . i.e.,  $b^{\frac{p-1}{2}}$ ,  $b^{\frac{p-1}{3}}$ , ...,  $b^{\frac{p-1}{p-1}}$  mod  $p$  include all of the residue classes mod  $p$  (except 0).

4 How many primitive roots are there for a given prime  $p$ ? That is, how many primitive roots does  $\mathbb{Z}_p^*$  have? It can be shown that the group of  $\mathbb{Z}_p^*$  has  $\varphi(p-1)$  elements and that  $\varphi(p-1)$  is the number of primitive roots. By a primitive root, we mean an element whose powers generate all the non-zero residue classes mod  $p$ .

5 Fermat's Little Theorem If  $p$  is a prime, then for any integer  $a$ ,  $a^p \equiv a$  mod  $p$ . This is equivalent to saying that  $a^{p-1} \equiv 1$  mod  $p$  for all  $a \neq 0$  mod  $p$ .

6 But if  $a$  is a primitive root modulo  $p$ , then  $a^{p-1} \equiv 1$  mod  $p$  and  $a^k \not\equiv 1$  mod  $p$  for any  $k < p-1$ . Hence,  $a^{p-1} \equiv a$  mod  $p$  implies  $a^{p-1} \equiv 1$  mod  $p$ .

7 For a prime  $p$ , the number of primitive roots is  $\varphi(p-1)$ . For example, for  $p = 7$ , the primitive roots are 3 and 5.

8 Example: If  $p = 11$ , then  $\mathbb{Z}_{11}^*$  has 4 primitive roots. They are 2, 6, 7, and 8.

# Primitive Roots

Let  $p$  be a prime. Then  $b$  is a *primitive root* for  $p$  if the powers of  $b$ ,

$1, b, b^2, b^3, \dots$

include all of the residue classes mod  $p$  (except 0).

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1 Since there are  $p-1$  residue classes mod  $p$  (not counting 0), that means the first  $p-1$  powers of  $b$  have to be different mod  $p$ .

We have noticed that the powers of  $b$  form a repeating cycle, and that cycle can't be longer than  $p-1$  (because of Fermat's Little Theorem). So,  $b$  is a primitive root if the cycle is as long as it can possibly be.

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1 Diffie-Hellman Key Exchange Dr V. K. Pachhare

2 Primitive Roots

3 Primitive Roots

4 How to prove that 3 is primitive root of 7

5 Roots

6 But 3 is a primitive root because the powers of 3 have a repeating cycle of length 6

7 For a prime number p > 3, the primitive roots of p are given by  $\frac{p-1}{2}$  numbers

8 Example 3 is a primitive root of 7 because its powers mod 7 are 1, 3, 2, 6, 4, 5.....

**Example:**

If  $p=7$ ,

$\text{Mod } 7 = \{0, 1, 2, \dots, 6\}$

Powers of 3 are  $\{1, 3, 9, 27, 81, \dots\}$

(Powers of 3) mod 7 are  $\{1, 3, 2, 6, 4, 5, \dots\}$

Hence 3 is a primitive root for  $p$  because the powers of 3 are 1, 3, 2, 6, 4, 5---that is, every number mod 7 occurs except 0.

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But 2 isn't a primitive root because the powers of 2 are 1, 2, 4, 1, 2, 4, 1, 2, 4...missing several values.

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if  $p = 14$  then the elements of  $\mathbb{Z}_n^*$  are the congruence classes  $\{1, 3, 5, 9, 11, 13\}$ ; there are  $\varphi(14) = 6$  of them.

Here is a table of their powers modulo 14:

$x$	$x, x^2, x^3, \dots \pmod{14}$
1 :	1
3 :	3, 9, 13, 11, 5, 1
5 :	5, 11, 13, 9, 3, 1
9 :	9, 11, 1
11 :	11, 9, 1
13 :	13, 1

The order of 1 is 1, the orders of 3 and 5 are 6, the orders of 9 and 11 are 3, and the order of 13 is 2.  
Thus, 3 and 5 are the primitive roots modulo 14.

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**Example:** If  $p=13$ , then  $2$  is a primitive root because the powers of  $2 \pmod{13}$  are  $1, 2, 4, 8, 3, 6, 12, 11, 9, 5, 10, 7$ ---which is all of the classes mod 13 except 0. There are other primitive roots for 13.



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Alternative method

We can verify any given number  $b$  is a primitive root of a number  $p$  using following method. Calculate  $n = \phi(\phi(p))$  where  $n$  is the total number of primitive roots for  $p$

1. Select any number  $b$  to check whether it is a primitive root of  $p$  or not.
2. Calculate  $b^{(\phi(p)/2)} \pmod{p}$ ,  $b = 3$ ,  $\phi(7) = 6$ ,  $b^{(6/2)} = 3^3 \equiv 27 \pmod{7} \equiv 6 \pmod{7}$
3. If  $b^{(\phi(p)/2)} \equiv -1 \pmod{p}$ , then  $b$  is the primitive root of  $p$

OTHERWISE NOT.

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If  $p = 19$ , state whether  $b = 5$  is a primitive root or not.

Solution:

There are exactly  $\phi(\phi(19))=\phi(18)=\phi(2 \times 3^2)=2 \times 3=6$  primitive roots. So pick one at random and check to see if  $b^9 \equiv -1 \pmod{19} = 18 \pmod{10}$ ; if yes, then  $b$  is a primitive root; if not, then pick something else.

Suppose  $p = 5$ ; then  $5^9 = 1 \pmod{19}$ . Therefore 5 is not a primitive root of 19.

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**Diffie-Hellman Key Exchange**

- First public-key type scheme proposed by Diffie & Hellman in 1976 along with the exposition of public key concepts
  - note: now know that James Ellis (UK CESG) secretly proposed the concept in 1970
- is a practical method for public exchange of a secret key
- used in a number of commercial products

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The idea of public key schemes, and the first practical scheme, which was for key distribution only, was published in 1977 by Diffie & Hellman. The concept had been previously described in a classified report in 1970 by James Ellis (UK CESG) - and subsequently declassified in 1987. See [History of Non-secret Encryption](#).



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• q  
•  $\alpha$   
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Prime number  
 $\alpha < q$  and  $\alpha$  a primitive root of  $q$

User A key generation  
Select private  $x_A$   $x_A < q$   
Calculate public  $Y_A$   $Y_A = \alpha^{x_A} \text{ mod } q$

User B key generation  
Select private  $X_B$   $X_B < q$   
Calculate public  $Y_B$   $Y_B = \alpha^{X_B} \text{ mod } q$

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Generate a Secret Key by User A

$$K = Y_B^{X_A} \bmod q$$

Generate a Secret Key by User B

$$K = Y_A^{X_B} \bmod q$$

Both calculations produce identical results.

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prime  $q = 23$  and  $\alpha = 5$

select random secret keys:

A chooses  $X_A = 6$

B chooses  $X_B = 15$

Compute the Shared session key



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$$Y_A = 5^6 \bmod 23 = 8$$

$$Y_B = 5^{15} \bmod 23 = 19$$

$$K = 19^6 \bmod 23 = 2$$

$$K = 8^{15} \bmod 23 = 2$$



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Users A and B use the Diffie–Hellman key exchange technique with a common prime  $q = 71$  and a primitive root  $a = 7$ .

- (a) If user A has private key  $X_A = 5$ , what is A's public key  $Y_A$ ?
- (b) If user B has private key  $X_B = 12$ , what is B's public  $Y_B$ ?
- (c) What is the shared secret key?



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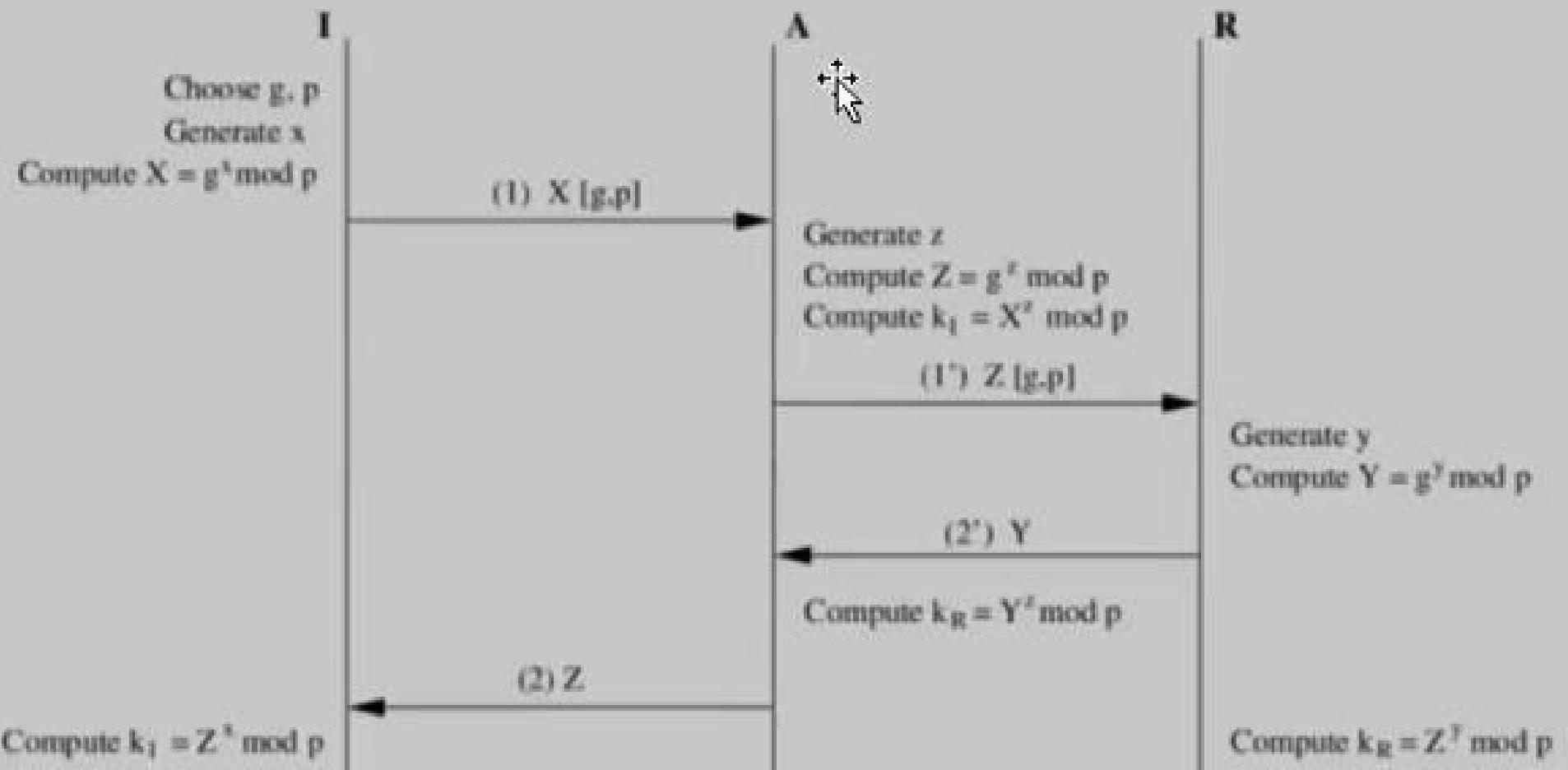


Figure 3: Man-in-the-Middle attack on Diffie-Hellman key exchange protocol

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Unit-VI  
Date: 13 Oct 2021  
Dr. V. K. Pachghare

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1 Cryptography and Network Security  
Unit-VI  
Date: 13 Oct 2021  
Dr. V. K. Pachghare

2 Firewalls  
V. K. Pachghare

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3 Introduction

- Play an important and major role to protect the system.
- Prevents unauthorized access to and from the network.
- Protect the network from the attackers.
- Allow the internal users to access the outside network via Internet and WLAN.

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4 Functions

- It blocks unauthorized traffic.
- It forward the incoming traffic to more reliable internal computer systems.
- It hides internal computers or network which are vulnerable.
- It hides the information about internal network such as name of the computer system, network topology used, types of network device, etc.

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- It provides strong user authentication.
- It can serve as a platform for IPsec.

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# Firewalls

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3 Introduction

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The idea of public key schemes, and the first practical scheme, which was for key distribution only, was published in 1977 by Diffie & Hellman. The concept had been previously described in a classified report in 1970 by James

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# Functions

- It blocks unauthorized traffic
- It forward the incoming traffic to more reliable internal computer systems
- It hides internal computers or network which are vulnerable
- It hides the information about internal network such as names of the computer system, network topology used, types of network device, etc.

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- It provides strong user authentication
- It can serve as a platform for IPSec

**Techniques to control access**

- Service control: It determines the access given to which types of internet services. It filters the traffic on the basis of port number, protocol or IP address.
- Direction control: It decides from where the particular service requests be initiated. It decides whether to allow or not allow the request to flow through the firewall.

**User control:** Depending upon the user access, it controls the access to a service.

**Behavior control:** It controls the behavior of a particular service.

**Classification of Firewalls**

- Packet filtering firewall
- Application level proxy
- Circuit switched gateway

Combination of above is dynamic packet filter

**Packet Filtering Firewall**



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**Techniques to control access**

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**Classification of Firewalls**

- Packet filtering firewall
- Application level proxy
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Combination of above is dynamic packet filter

**Packet Filtering Firewall**

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- Behavior control:** It controls the behavior of a particular service.

**Classification of Firewalls**

- Packet filtering firewall
- Application level gateway
- Circuit gateway gateway

Combination of above is dynamic packet filter

**Packet Filtering Firewall**



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- **User control:** Depending upon the user access, it controls the access to a service.
- **Behavior control:** It controls the behavior of a particular service.

**Classification of Firewalls**

- Packet filtering firewall
- Application level gateway
- Circuit gateway gateway

Combination of above is dynamic packet filter

**Packet Filtering Firewall**

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# Classification of Firewalls

- Packet filtering firewall
- Application level gateways
- Circuit (proxies) gateways

Combination of above is dynamic packet filter

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Slides Outline

**Packet Filtering Firewall**

10. Simplest of components  
User transport layer information only  
- IP Source Address, Destination Address  
- Protocol/Next Header (TCP, UDP, ICMP, etc)  
- TCP or UDP source & destination ports  
- TCP Flags (SYN, ACK, FIN, RST, PSH, etc)  
Examples  
- DNS uses port 53

11. Advantages  
Performance is good  
Very fast  
Relatively inexpensive  
Easy to use  
Traffic management is good  
Simplicity

12. Disadvantages  
Direct connections are allowed between untrusted and trusted hosts  
Vulnerable to spoofing attacks  
Poor scalability  
Large portranges may be opened

13. Most of these firewalls do not support advanced user authentication schemes  
Due to improper configuration it is susceptible to various breaches

**(a) Packet-filtering router**

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Slides Outline

10. Simplest of components

- Uses transport-layer information only
  - IP Source Address, Destination Address
  - Protocol/Next Header (TCP, UDP, ICMP, etc)
  - TCP or UDP source & destination ports
  - TCP Flags (SYN, ACK, FIN, RST, PSH, etc)
- Examples
  - DNS uses port 53

11. Advantages

- Performance is good
- Very fast
- Relatively inexpensive
- Transparent to users
- The traffic management is good
- Simplicity

12. Disadvantages

- Direct connections are allowed between untrusted and trusted hosts
- Vulnerable to spoofing attacks
- Poor scalability
- Large port ranges may be opened

13. Most of these firewalls do not support advanced user authentication schemes.  
Due to improper configuration it is susceptible to security breaches.

14. Attacks Against Packet Filtering

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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Slides Outline

10. **Samplest of components**

- User transport layer information only
  - IP Source Address, Destination Address
  - Protocol/Next Header (TCP, UDP, ICMP, etc)
  - TCP or UDP source & destination ports
  - TCP Flags (SYN, ACK, FIN, RST, PSH, etc)
- Examples
  - DNS uses port 53

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11. **Advantages**

- Performance is good
- Very fast
- Relatively inexpensive
- Transparent to users
- The traffic management is good
- Simplicity

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12. **Disadvantages**

- Direct connections are allowed between untrusted and trusted hosts
- Vulnerable to spoofing attacks
- Poor scalability
- Large port ranges may be opened

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13. **Most of these firewalls do not support advanced user authentication schemes**

- Due to improper configuration it is susceptible to security breaches

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14. **Attacks Against Packet Filtering**

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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The image shows a Microsoft PowerPoint presentation titled "Firewall - Microsoft PowerPoint". The main slide, slide 11, is titled "Advantages" in large blue font. It lists several bullet points: "Performance is good", "Very fast", "Relatively inexpensive", "Transparent to users", "The traffic management is good", and "Simplicity". To the left of the main slide, a navigation pane shows five other slides. Slide 10 discusses the components of a firewall, mentioning IP source and destination addresses, protocols like TCP and UDP, and TCP flags. Slides 12 and 13 discuss disadvantages, mentioning direct connections between untrusted and trusted hosts, vulnerability to spoofing, poor scalability, and security breaches due to configuration. Slide 14 discusses attacks against packet filtering, specifically IP address spoofing, source routing attacks, and tiny fragment attacks. The bottom of the screen shows the Windows taskbar with various icons and the date/time.



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Slides Outline

10. **Sniffing of components**

- Uses transport layer information only
  - IP Source Address, Destination Address
  - Protocol(Next Header: TCP, UDP, ICMP, etc)
  - TCP or UDP source & destination ports
  - TCP Flags (SYN, ACK, FIN, RST, PSH, etc)
- Examples
  - DOS user port 80

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11. **Advantages**

- Performance is good
- Very fast
- Relatively inexpensive
- Transparent to users
- The traffic management is good
- Simplicity

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12. **Disadvantages**

- Direct connections are allowed between untrusted and trusted hosts
- Vulnerable to spoofing attacks
- Poor scalability
- Large port ranges may be opened

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13. **Most of these firewalls do not support advanced user authentication schemes**

- Due to improper configuration it is susceptible to security breaches

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14. **Attacks Against Packet Filtering**

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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## Disadvantages

- Direct connections are allowed between untrusted and trusted hosts
- Vulnerable to spoofing attacks
- Poor scalability
- Large port ranges may be opened

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- Snapshot of components
  - Uses transport layer information only
    - IP Source Address, Destination Address
    - Protocol Next Header (TCP, UDP, ICMP, etc)
    - TCP or UDP source & destination ports
    - TCP Flags (SYN, ACK, FIN, RST, PSH, etc)
  - Examples
    - DOS uses port 80
- Advantages
  - Performance is good
  - Very fast
  - Relatively inexpensive
  - Transparent to users
  - The traffic management is good
  - Simplicity
- Disadvantages
  - Direct connections are allowed between untrusted and trusted hosts
  - Vulnerable to spoofing attacks
  - Poor scalability
  - Large port ranges may be opened
- Most of these firewalls do not support advanced user authentication schemes
- Due to improper configuration it is susceptible to security breaches

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Slides Outline

14 Attacks Against Packet Filtering

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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15 IP Address Spoofing

- The attackers put the internal address as a source address
- He believes that due to this spoof address, the firewall assumes that they belong from trusted internal host and allow to pass that packet.

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16 Countermeasure

- If the packet coming from outside and has the IP address of the internal host, then discard such a packet. This is implemented at the router which is external to the firewall.

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17 Source Routing Attacks

- Attackers assume that the route routing information is not analyzed so the route of a packet across the Internet is specified by the source

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18 Countermeasure

- Discard the packets which use this option

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Slides Outline

15 Attacks Against Packet Filtering

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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16 IP Address Spoofing

- The attackers put the internal address as a source address
- He believes that due to this spoof address, the firewall assumes that this packet is from trusted internal host and allow to pass that packet.

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17 Countermeasure

- If the packet coming from outside and has the IP address of the internal host, then discard such packet. This is implemented at the router which is essential to the firewall.

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18 Source Routing Attacks

- Attackers assume that the route routing information is not analyzed so the route of a packet across the Internet is specified by the source

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19 Countermeasure

- Discard the packets which use this option

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# IP Address Spoofing

- The attackers put the internal address as a source address
- He believes that due to this spoof address, the firewall assumes that this packet is from trusted internal host and allow to pass that packet.

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14 Attacks Against Packet Filtering

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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15 IP Address Spoofing

- The attackers put the internal address as a source address
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- If the packet coming from outside and has the IP address of the internal host, then discard such packet. This is implemented at the router which is external to the firewall.

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17 Source Routing Attacks

- Attackers assume that the route routing information is not analyzed so the route of a packet across the Internet is specified by the source

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18 Countermeasure

- Discard the packets which use this option

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The screenshot shows a Microsoft PowerPoint presentation titled 'Firewall'. The main slide, slide 16, has a blue header and footer. The header contains the title 'Countermeasure' in large purple font. The footer includes the college's name 'Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education' and its logo. The main content area contains a bulleted list of countermeasures for source routing attacks. The slide is part of a larger presentation with other slides visible in the navigation pane on the left, each with a different title and bullet points. The status bar at the bottom shows the slide number, language, date, and time.



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Slides Outline

14 Attacks Against Packet Filtering

- IP ADDRESS SPOOFING
- SOURCE ROUTING ATTACKS
- TINY FRAGMENT ATTACKS

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15 IP Address Spoofing

- The attackers put the internal address as a source address
- He believes that due to this spoof address, the firewall assumes that they packets from trusted internal host and allow to pass that packet.

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16 Countermeasure

- If the packet coming from outside and has the IP address of the internal host, then discard such packet. This is implemented at the router which is external to the firewall.

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17 Source Routing Attacks

- Attackers assume that the source routing information is not analyzed so the route of a packet across the Internet is specified by the source

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18 Countermeasure

- Discard the packets which use this option.

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# Source Routing Attacks

- Attacker assume that the source routing information is not analyzed so the route of a packet across the Internet is specified by the source

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Slides Outline

18 Countermeasure  
Discard the packets which use this option

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19 Tiny Fragment Attacks  
The attacker creates small fragments using the IP fragmentation option.  
The reassembly fragment is used for the TCP header information.  
This design helps to avoid the filtering rules based on TCP header information.

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20 The filtering decision is taken from the first fragment of a packet and on the basis of this first fragment, subsequent fragments of that packet are allowed or discarded.  
The attacker takes advantage of this strategy that only the first fragment of the packet is examined by forwarding the complete packet or discarding it.

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21 Countermeasure  
Enforce a rule that the first fragment must hold a predefined minimum amount of the transport header.  
The filter should reassemble the packet. Then discard the remaining fragments of the packet.

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22 Application Level Gateways  
Packet filtering decisions are based on address information, so it examines the lower layers of the OSI model.

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Slides Outline

1 Countermeasure  
Discard the packets which use this option

2 Tiny Fragment Attacks  
The attacker creates small fragments using the IP fragmentation option.  
The separate fragment is used for the TCP header information.  
This design helps to avoid the filtering rules based on TCP header information.

3 The filtering decision is taken from the first segment of a packet and on the basis of this first fragment, subsequent fragments of that packet are allowed or discarded.  
The attacker takes the advantage of this strategy that only the first fragment of the packet is examined by forwarding the complete packet or discarding it.

4 Countermeasure  
Enforce a rule that the first fragment must hold a predefined minimum amount of the transport header. The filter should reassemble the packet. Then discard the remaining fragments of the packet.

5 Application Level Gateways  
Packet filtering decisions based on address information, so it examines the lower layers of the OSI model.

# Tiny Fragment Attacks

- The attacker creates small fragments using the IP fragmentation option.
- The separate fragment is used for the TCP header information.
- This design helps to avoid the filtering rules based on TCP header information.

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**Slides Outline**

18 Countermeasure  
• Discard the packets which use this option

19 Tiny Fragment Attacks  
• The attacker creates small fragments using the IP fragmentation option.  
• The retransmit fragment is used for the TCP header information.  
• This design helps to avoid the filtering rules based on TCP header information.

20 Countermeasure  
• The filtering decision is taken from the first fragment of a packet and on the basis of this first fragment, subsequent fragments of that packet are allowed or discarded.  
• The attacker takes advantage of this strategy that only the first fragment of the packet is examined for forwarding the complete packet or discarding it.

21 Countermeasure  
• Enforce a rule that the first fragment must hold a predefined minimum amount of the transport header.  
• Routers should reassemble the packet if the first fragment of the packet is received. Then discard the remaining fragments of the packet.

22 Application Level Gateways  
• Packet filtering decisions are based on address information, so it examines the lower layers of the OSI model.

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**Countermeasure**

- Discard the packet which use this option

**Tiny Fragment Attacks**

- The attacker creates small fragments using the IP fragmentation option.
- The reassembly is used for the TCP header information.
- This design helps to avoid the filtering rules based on TCP header information.

The filtering decision is taken from the first fragment of a packet and on the basis of this first fragment, subsequent fragments of that packet are allowed or discarded.

The attacker takes advantage of this strategy that only the first fragment of the packet is examined by discarding the complete packet or discarding it.

**Countermeasure**

Enforce a rule that the first fragment must hold a predefined minimum amount of the transport header. The filter should remember the packet, if the first fragment of the packet is rejected. Then discard the remaining fragments of the packet.

**Application Level Gateways**

Packet filtering decisions are based on address information, so it examines the lower layers of the OSI model.

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<p>Firewall - Microsoft PowerPoint</p> <p>You are sharing your entire screen. Stop Sharing</p> <p>Slides Outline</p> <p>Application Level Gateways</p> <ul style="list-style-type: none"><li>• Packet filtering firewalls based on address information, so it examines the lower layers of the OSI model.</li><li>• Application level gateway firewall provides security to all layers of the OSI model.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p>23</p> <ul style="list-style-type: none"><li>• It uses server based programs known as proxy servers or bastion hosts.</li><li>• It forward a request by ensuring that the protocol specification is correct.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p>24</p> <ul style="list-style-type: none"><li>• It receives a request from the external side, examine the request, and then forward the legitimate and required requests to the destination host or to the other side.</li><li>• It makes decisions at all the seven layers of the OSI model.</li><li>• It acts as a mediator for different applications such as e-mail, FTP, etc. It does not permit the client to directly connect to the destination node.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p>25</p> <p>Advantages</p> <ul style="list-style-type: none"><li>• It is configured so that firewall is the only host addressed that is visible to an outside network.</li><li>• For separate services, separate proxy servers are used.</li><li>• It supports strong user authentication.</li><li>• Application level security.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p>26</p> <ul style="list-style-type: none"><li>• At the application level, it is easy to log and audit all the incoming traffic.</li><li>• It provides strong access controls.</li></ul> <p>Click to add notes</p> <p>Side 22 of 39 Default Design English (United States)</p> <p>110% ENG 4:30 PM IN 10/13/2021 +91</p>				

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<p>Firewall - Microsoft PowerPoint</p> <p>Home Insert Design Animations Slide Show Review View</p> <p>Cut Copy Format Painter New Slide Delete Slides</p> <p>Font Paragraph Drawing Editing</p> <p>You are sharing your entire screen. Stop Sharing</p> <p>Slides Outline</p> <p><b>Advantages</b></p> <ul style="list-style-type: none"><li>• Packet filtering firewall based on address information, as it examines the lower layers of the OSI model.</li><li>• Application level gateway firewall provides security to all layers of the OSI model.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p><b>Advantages</b></p> <ul style="list-style-type: none"><li>• It uses server based programs known as proxy servers or bastion hosts.</li><li>• It forward or reject the packets by ensuring that the protocol specification is correct.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p><b>Advantages</b></p> <ul style="list-style-type: none"><li>• It receives requests from the external side, examine the request, and then forward the legitimate and required requests to the destination host on the other side.</li><li>• It makes decisions at all the seven layers of the OSI model.</li><li>• It acts as a medium for different applications such as e-mail, FTP, etc. It does not permit the client to directly connect to the destination nodes.</li></ul> <p>Department of Computer Engineering and Information Technology College of Engineering Pune (COEP) Forerunners in Technical Education</p> <p>At the application level, it is easy to log and audit all the incoming traffic.</p> <p>Provides strong access controls.</p> <p>Click to add notes</p>				

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**Disadvantages**

- At the application level, it is easy to log and audit all the incoming traffic
- It provides strong access controls
- More secure than packet filtering firewall

**Disadvantages**

- For each application special proxy is required
- Performance is slow
- On each connection there is an additional processing overhead
- Sometimes it is inconvenient faced by the users
- There is lack of transparency

**DMZ (Demilitarized Zone)**

- Two types of firewall
  - Internal firewall
  - External firewall
- Internal firewalls protect the entire network of an organization.
- The external firewall provides basic security to the entire network. An internal firewall is installed at the boundary of local or organization network.

**It is located inside the boundary zone. This region between two firewalls is called demilitarized zone or DMZ.**

- Many network devices are located between these two firewalls.
- It includes devices which are allowed to access from external networks.
- These devices can be accessed externally but protected from vulnerability.

The internal firewall provides two-way protection.

- Initially it protects the entire internal network from attacks launched from DMZ systems originate from viruses, worms, botnets or other malicious software.
- Then the internal firewall protects system located at DMZ.

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**DMZ (Demilitarized Zone)**

- Two types of firewall
  - Internal firewall
  - External firewall

Internal firewalls protect the entire network of an organization.

The external firewall provides basic security to the entire network. An external firewall is installed at the boundary of local or organization network.

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26 At the application level, it is easy to log and audit all the incoming traffic.

It provides strong access controls.

More secure than packet filtering firewall.

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27 Disadvantages

- For each application special proxy is required.
- Performance is slow.
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28 DMZ (Demilitarized Zone)

- Two types of firewall
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29 It is located inside the boundary router. This region between two firewalls is called demilitarized zone or DMZ.

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DMZ

```
graph LR; subgraph IPN [Internal Private Network]; S1[Server]; C1[Client]; R1[Router]; end; subgraph DMZ [DMZ]; WS[Web Server]; FS[FTP Server]; FW[Firewall]; end; subgraph EPN [External Public Network]; C2[Customer]; H1[Hacker]; H2[Hacker]; end; R1 --- WS; R1 --- FS; R1 --- FW; S1 --- R1; C1 --- R1; WS --- FW; FW --- C2; FW --- H1; FW --- H2;
```

**DMZ**

Internal Private Network

DMZ

External Public Network

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- The external firewall provides boundary protection.
- Initially it protects the entire internal network from attacks launched from DMZ systems against the servers, hosts or other malicious software.
- Then the internal firewall protects system located at DMZ area from internal attack. To protect different portions of the large network from each other, multiple internal firewalls are used.

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**Circuit Level Gateways**

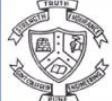
- Stand-alone system
- It validates connections and then allows the data to be exchanged.
- They also work as per defined rules similar to packet filters.
- Circuit level gateways cannot route the packets.

33

- The connections are allowed or discarded based on these rules. So circuit level gateway establishes the connection between the source and the destination. It focuses on the TCP/IP layer.
- This firewall is installed between the internal and the external networks such as Internet. The actual address is hidden from the external users because only the address of the proxy is transmitted.

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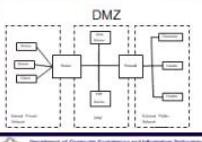
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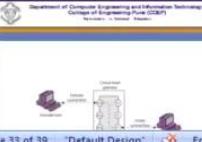
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30 • The internal firewall provides recovery protection.  
Initially it protects the entire internal network from attacks launched from DMZ systems against its users, lots or other malicious software.  
Then the internal firewall protects system located at DMZ area from internal attack. To protect different portions of the large network from each other, multiple internal firewalls are used.

31 

32 **Circuit Level Gateways**  
• Stand-alone system  
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33 • The connections are allowed or discard based on these rules. So circuit level gateway establishes the connection between the source and the destinations. It focuses on the TCP/IP layer.  
• This firewall is installed between the router and the external network such as Internet. The actual address is hidden from the external users because only the address of the proxy is transmitted.

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Slides Outline

35 Advantages

- Transparency to users
- Excellent for relaying external traffic

36 Disadvantages

- Slower than packet filtering firewall
- Inbound traffic is risky

37 Benefits of a Firewall

- Increased ability to enforce network security standards/polices
- Generalization of inter-network audit capability audit of in/out-bound traffic

38 Limitations of a Firewall

- It cannot protect those attacks that bypass the firewall.
- It cannot protect the network against the internal attacks.

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Outside connection

Inside connection

Circuit-level gateway

Outside host

Inside host

(c) Circuit-level gateway

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**Advantages**

- Transparent to users
- Excellent for relaying outbound traffic

**Disadvantages**

- Slower than packet filtering firewall
- Inbound traffic is risky

**Benefits of a Firewall**

- Increased ability to enforce network security standards/polices
- Generalization of inter-network audit capability audit of in/out-bound traffic

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- It cannot protect those attacks that bypass the firewall.
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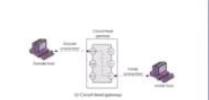
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34. 

35. **Advantages**

- Transparent to users
- Excellent for isolating internal traffic

36. **Disadvantages**

- Slower than packet filtering firewall
- Inbound traffic is risky

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**Benefits of a Firewall**

- Increased ability to enforce network security standards/policies
- Centralization of inter-network audit capability (audit of in/out-bound traffic)

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37 Benefits of a Firewall

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38 Limitations of a Firewall

- It cannot protect those attacks that bypass the firewall.
- It cannot protect the network against the internal attacks.

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An internal firewall that separates the different parts of a network cannot protect against wireless communications among local computer systems on different sides of the internal firewall.

Different devices such as laptop or portable storage device may be used and infected outside the network, and then used internally.

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# Intrusion

- An intrusion is a deliberate, unauthorized attempt to access or manipulate information or system and to render them unreliable or unusable.
- Intruders may be from outside the network or legitimate users of the network.
- Intrusion can be a physical, system or remote intrusion.



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# Intruders

- Significant issue for networked systems is hostile or unwanted access
- either via network or local
- can identify classes of intruders:
  - **Masquerader**- An individual who is not authorized to use the computer and who penetrates a system's access controls to exploit a legitimate user's account



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## Jargon Related to IDS

- False Negative:

False negatives are any alert that should have happened but didn't. Attack → Normal

- False Positive:

False positive is any normal or expected behavior that is identified as anomalous or malicious

Normal → Attack



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# Intrusion Detection Systems (IDS)

Intrusion Detection Systems look for attack signatures, which are specific patterns that usually indicate malicious or suspicious intent.



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## Types of IDS

- Anomaly detection
- Signature based (misuse)
- Host based
- Network based



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## Anomaly based IDS

- Anomaly-Based IDS examines ongoing traffic, activity, transactions and behavior in order to identify intrusions by detecting anomalies.
- This IDS models the normal usage of the network as a noise characterization.
- Anything distinct from the noise is assumed to be an intrusion activity. E.g flooding a host with lots of packet.
- The primary strength is its ability to recognize novel attacks.



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## Drawbacks

- Assumes that intrusions will be accompanied by manifestations that are sufficiently unusual so as to permit detection.
- These generate many false alarms and hence compromise the effectiveness of the IDS.



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# Signature based IDS

- Signature-Based IDS use a rule set to identify intrusions by watching for patterns of events specific to known and documented attacks.
- It is typically connected to a large database which houses attack signatures.
- It compares the information it gathers against those attack signatures to detect a match.



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# Signature based IDS

- These IDS are normally presumed to be able to detect only attacks “known” to its database.
- Thus, if the database is not updated with regularity, new attacks could slip through.
- It can, however, detect new attacks that share characteristics with old attacks,
- In cases of new, uncataloged attacks, this technique is pretty porous.



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## Drawbacks of Signature based IDS

- They are unable to detect novel attacks.
- Suffer from false alarms
- Have to programmed again for every new pattern to be detected.
- Also, it may affect performance in cases when intrusion patterns match several attack signatures



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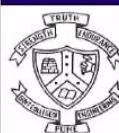
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## Host based IDS

- IDS is installed on a host in the network.
- It collects and analyzes the traffic that is originated or is intended to that host.
- HIDS leverages their privileged access to monitor specific components of a host that are not readily accessible to other systems.



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- These audit information includes events like the use of identification and authentication mechanisms (logins etc.), file opens and program executions, admin activities etc.
- This audit is then analyzed to detect trails of intrusion.



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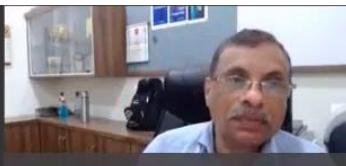
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## Strengths of the host based IDS

- Attack verification
- System specific activity
- Encrypted and switch environments
- Monitoring key components
- Near Real-Time detection and response.
- No additional hardware



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## Drawbacks

- Can not see all network activities
- Running audit mechanisms adds overload to system, performance may be an issue
- Audit trails can take lots of storage
- OS vulnerabilities can undermine the effectiveness of agents
- Escalation of false positive
- Greater deployment and maintenance cost



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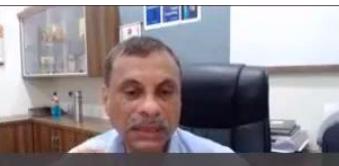
Network based IDS

- Network IDSs (NIDS) are placed in key areas of network infrastructure and monitors the traffic as it flows to other host.
- Unlike HIDS, NIDS have the capability of monitoring the network and detecting the malicious activities intended for that network.

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## Network based IDS

- This IDS looks for attack signatures in network traffic via a promiscuous interface.
- A filter is usually applied to determine which traffic will be discarded or passed on to an attack recognition module. This helps to filter out known un-malicious traffic.



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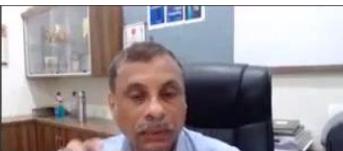
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## Strengths

- Can get information quickly without any reconfiguration of computers or need to redirect logging mechanisms
- Does not affect network or data sources
- Monitor and detects in real time networks attacks or misuses
- Does not create system overhead



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## Disadvantages

- Cannot scan protocols if the data is encrypted
- Can infer from network traffic what is happening on host but cannot tell the outcome
- Hard to implement on fully switched networks
- Has difficulties sustaining network with a very large bandwidth



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## Commercial ID Systems

- ISS – Real Secure from Internet Security Systems:
  - Real time IDS.
  - Contains both host and network based IDS.
- Tripwire – File integrity assessment tool.
- Bro and Snort – open source public-domain system.



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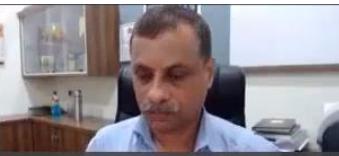
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# Vulnerabilities in TCP/IP model

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## Introduction

- The TCP/IP protocol suite was created
  - as an internetworking solution (with little or no regard to security aspects.)
  - The development of TCP/IP protocol suite was focused on creating a communication protocol standard that can interoperate between different hardware devices and software independently.



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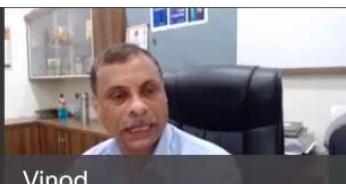
- Other major goals included:
  - failure recovery and the ability to handle high error rates,
  - efficient protocol with low overhead, routable data
  - the ability to add new networks to an already existing network without disrupting the existing network

**Its main emphasis was providing a suite not the security**

Hence, by default, TCP/IP has security flaws at both the protocol level and implementation.



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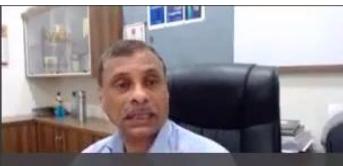
## TCP/IP Protocol Model

- Four Layers
  - Application
  - Transport
  - Internetwork
  - Network access

At each layer, there are some security weaknesses that can be exploited by attackers



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## APPLICATION LAYER

Some of application layer protocols are:

- Web Application and Browser Security Vulnerabilities
- Domain Name System (DNS)
- Dynamic Host Control Protocol (DHCP)

**Each of these protocols has vulnerabilities.**



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## Caching

- Web browsers perform caching of the web pages
- The contents in the cache are saved temporarily
- The cache can contain images, passwords and user names.
- If the user's computer was compromised, an attacker can view all the contents and the user's browsing habits without any need of being authenticated and thus can be a privacy concern



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## Counter-measure

- It is important to clear the cache once in a while
- Disable the auto saving feature in the browser of passwords and user names in the cache



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## Session Hijacking

- Hijacking is possible when the attacker steals an HTTP session after observing and capturing the packets using a packet sniffer.
- It enables the attacker to have full access to the HTTP session and the communication changes from the client to the web server to attacker to the web server.
- Hijacking is possible when there is weak authentication between the client and the web server during the initializing of the session



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## Cookie Poisoning

- Cookies are used to maintain the state of a session to avoid the user to retype their credentials every time they visit a site or change web pages.
- Cookies are used by many web applications (including browsers) to save information. This information is stored permanently or temporarily on the client machine.



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- Cookie poisoning is the modification or theft of cookie in a user's machine by an attacker in order to release personal information.
- If the attacker gets hold of a cookie containing a password and username, they can use the cookie on their machine and the web server will not request any authentication because the cookie will issue out the username and password automatically.
- With cookie poisoning, an attacker can gain access to unauthorized information about a user and possibly steal their identity

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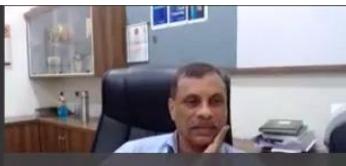
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## Counter Measure

- Web Application Firewalls (WAF) are able to detect and block cookie poisoning attacks.
- Web Application Firewalls are able to inspect the HTTP sessions and can trace down the parameters set in the cookies that have been issued by the web server



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## Replay Attack

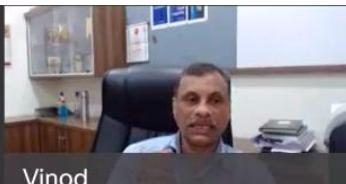
- This involves man-in-the-middle attack in which the sent data is repeatedly sent to the server.
- This is more than a hijack. The data resent can be modified and can bring different results.
- Furthermore, the attacker can spoof the IP address of a client and thus redirect his/her machine.

### Counter measure:

- The web browser should be able to tell that replayed traffic is not legitimate.



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## Cross-Site Scripting

- Hacker injecting malicious code in a web application or browser and is executed at the client side.
- The essence of this attack is to perform a session hijack by stealing session tokens and cookies of a legitimate user's session.

**Counter measure:**

- Disable scripts to run on the website. However, this control means that some functions and features on the website will not be available.
- Another alternative is to enhance the security controls when dealing with cookie based user authentication

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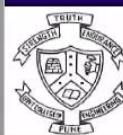
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## DNS Protocol Attacks

- There are following common attacks:
  - DNS cache poisoning
  - DNS spoofing



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## DNS Cache Poisoning

- It is an integrity attack that involves manipulating the information saved in the DNS cache giving it wrong information.
- This false information will offer a name to IP; mapping it to a wrong IP address.
- The objective is to divert the requests to another web site.
- This new web site might be bogus and offers the same or similar products and services as the real web site.
- If the user does not notice anything, he enters the user name and password.
- Then the attacker can steal the user credentials.



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# DNS Spoofing

- DNS spoofing refers to faking the IP address of a computer to match the IP address of the DNS server so that requests can be directed to that wrong computer.
- In this attack, the hacker's computer is considered to be a legitimate DNS server by clients and other servers.
- It will impersonate the DNS server and reply to all incoming requests from the clients thus misdirecting them
- DNS server spoofing attacks are often used to spread computer worms and viruses

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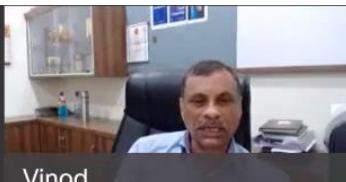
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## Counter Measure

- Avoid trust relationships:** Organizations should develop protocols that rely on **trust relationships as little as possible**. It is significantly easier for attackers to run spoofing attacks when trust relationships are in place because trust relationships only use IP addresses for authentication.



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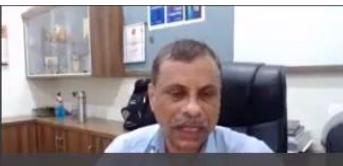
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## DHCP Starvation Attack

- It is the consuming of IP address space allocated by the DHCP server.
- An attacker broadcasts large number of DHCP REQUEST messages with spoofed source MAC addresses.
- If the legitimate DHCP Server in the network start responding to all these bogus DHCP REQUEST messages, then available IP Addresses in the DHCP server scope will be depleted within a very short span of time.



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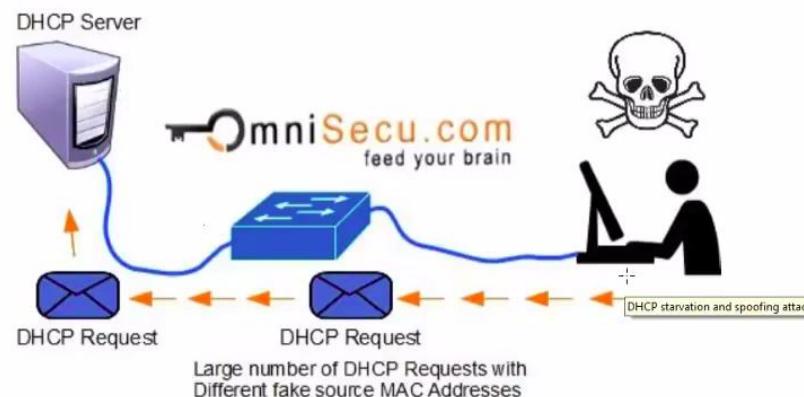
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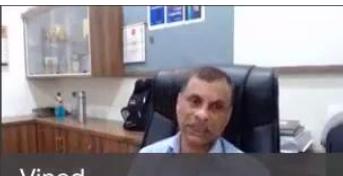
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- When a genuine user wants to access the network, the server will not offer an IP address automatically and the user will not be granted access into the network. This is a denial of service attack



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- Once the available number of IP Addresses in the DHCP server is depleted, network attackers can then set up a rogue DHCP server and respond to new DHCP requests from network DHCP clients. By setting up a rogue DHCP server, the attacker can now launch DHCP spoofing attack.

### Counter Measure:

- To prevent DHCP starvation, port security can be used because it only allows a specified number of MAC addresses per port



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## DHCP spoofing attack

- After a DHCP starvation attack and setting up a rogue DHCP server, the attacker can start distributing IP addresses and other TCP/IP configuration settings to the network DHCP clients.
- TCP/IP configuration settings include Default Gateway and DNS Server IP addresses.
- Network attackers can now replace the original legitimate Default Gateway IP Address and DNS Server IP Address with their own IP Address.

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