Porter Stemmer

The Porter Stemmer (Porter, 1980)

- A simple rule-based algorithm for stemming
- An example of a HEURISTIC method
- Based on rules like:
 - ATIONAL -> ATE (e.g., relational -> relate)
- The algorithm consists of seven sets of rules, applied in order

The Porter Stemmer: definitions

- Definitions:
 - CONSONANT: a letter other than A, E, I, O, U, and Y preceded by consonant
 - VOWEL: any other letter
- With this definition, all words are of the form: (C)(VC)^m(V)

C=string zero or more consonants V=string of one or more vowels

- E.g.,
 - Troubles
 - C V CVC

The Porter Stemmer: rule format

The rules are of the form:

(condition) S1 -> S2

Where S1 and S2 are suffixes

Conditions:

m	The measure of the stem
*5	The stem ends with S
* _V *	The stem contains a vowel
*d	The stem ends with a double consonant
*0	The stem ends in CVC (second C not W, X, or Y)

The Porter Stemmer: Step 1

- 1. SSES -> SS
 - 1. expresses -> express
- 2. IES -> I
 - 1. ponies -> poni
 - 2. ties -> ti
- 3. SS -> SS
 - 1. process -> process
- 4. $S \rightarrow \epsilon$
 - 1. cats -> cat

The Porter Stemmer: Step 2a (past tense, progressive)

- 1. (m>1) EED -> EE
 - 1. Condition verified: agreed -> agree
 - 2. Condition not verified: feed -> feed
- 2. (*V*) ED -> ϵ
 - 1. Condition verified: plastered -> plaster
 - 2. Condition not verified: bled -> bled
- 3. (*V*) ING -> ϵ
 - Condition verified: motoring -> motor
 - Condition not verified: sing -> sing

The Porter Stemmer: Step 2b (cleanup)

- (These rules are ran if second or third rule in 2a apply)
- 4. AT-> ATE
 - conflat(ed) -> conflate
- 5. BL -> BLE
 - Troubl(ing) -> trouble
- 6. (*d & ! (*L or *S or *Z)) -> single letter
 - Condition verified: hopp(ing) -> hop, tann(ed) -> tan
 - Condition not verified: fall(ing) -> fall
- 7. (m>1 & *o) -> E
 - Condition verified: fil(ing) -> file
 - Condition not verified: fail -> fail

The Porter Stemmer: Steps 3 and 4

- Step 3: Y Elimination (*V*) Y -> I
 - Condition verified: happy -> happi
 - Condition not verified: sky -> sky
- Step 4: Derivational Morphology, I
 - 8. (m>0) ATIONAL -> ATE
 - Relational -> relate
 - 9. (m>0) IZATION -> IZE
 - generalization-> generalize
 - 10. (m>0) BILITY-> BLE
 - sensibility -> sensible

The Porter Stemmer: Steps 5 and 6

- Step 5: Derivational Morphology, II
 - (m>0) ICATE -> IC
 - triplicate -> triplic
 - (m>0) FUL -> ε
 - hopeful -> hope
 - (m>0) NESS -> ε
 - goodness -> good
- Step 6: Derivational Morphology, III
 - (m>0) ANCE -> ε
 - allowance-> allow
 - (m>0) ENT -> ε
 - dependent-> depend
 - (m>0) IVE -> ε
 - effective -> effect
 - (m>0) IZE -> ε
 - generalize -> general
 - (m>0) ANT -> ϵ
 - reluctant-> reluct
 - (m>0) r -> ε
 - computer -> compute

The Porter Stemmer: Step 7 (cleanup)

- Step 7a
 - $(m>1) E -> \epsilon$
 - probate -> probat
 - $(m>1 \& !*o) NESS -> \epsilon$
 - goodness -> good
- Step 7b
 - (m>1 & *d & *L) -> single letter
 - Condition verified: controll -> control
 - Condition not verified: roll -> roll

Examples

- computers
 - Step 1, Rule 4: -> computer
 - Step 6, Rule 4: -> compute
- singing
 - Step 2a, Rule 3: -> sing
- controlling
 - Step 2a, Rule 3: -> controll
 - Step 7b : -> control
- generalizations
 - Step 1, Rule 4: -> generalization (noun)
 - Step 4, Rule 9: -> generalize (verb)
 - Step 6, last rule: -> general (adjective)

Problems

- elephants -> eleph
 - Step 1, Rule 4: -> elephant
 - Step 6, Rule 7: -> eleph
- doing > do
 - Step 2a, Rule 3: -> do

References

- The Porter Stemmer home page (with the original paper and code):
 http://www.tartarus.org/~martin/PorterStemmer/
- Jurafsky and Martin, chapter 3.4
- The original paper: Porter, M.F., 1980, An algorithm for suffix stripping, *Program*, 14(3):130-137.

SPELL CHECKER

Detecting and Correcting Spelling Error

Detecting and Correcting Spelling Error

- Detection and correction of spelling errors is an integral part of modern word processors and search engines.
- It is important in correcting errors in the recognition of human printed or cursive handwriting as the user is writing.

Categorization

- Categorized into 3 broader problems, as per Kukich(1992):
- 1. **non-word error detection:** detecting spelling errors that result in non-words (like *graffe for giraffe*).
- 2. **isolated-word error correction:** correcting spelling errors that result in nonwords, for example correcting *graffe to giraffe, but looking only at the word in* isolation.
- 3. context-dependent error detection and correction: using the context to help detect and correct spelling errors even if they accidentally result in an actual REALWORD word of English (real-word errors). This can happen from typographical errors (insertion, deletion, transposition) which accidentally produce a real word (e.g., there for three), or because the writer substituted the wrong spelling of a homophone or near-homophone (e.g., dessert for desert, or piece for peace).

Error detection

- Detecting non-word errors is generally done by marking any word that is not found in a dictionary.
- For example, the misspelling graffe would not occur in a dictionary.

Error Correction

- Algorithms for isolated-word error correction operate by finding words which are the likely source of the error-ful form.
- For example, correcting the spelling error *graffe* requires searching through all possible words like giraffe, graf, craft, grail, etc, to pick the most likely source.
- To choose among these potential sources we need a distance metric between the source and the surface error.
- Intuitively, giraffe is a more likely source than grail for graffe, because giraffe is closer in spelling to graffe than grail is to graffe.

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MINIMUM EDIT DISTANCE (Wagner and Fischer 1974)

- The non-probabilistic algorithm for this solution is the minimum edit distance algorithm.
- The distance between two strings is a measure of how alike two strings are to each other.
- The minimum edit distance between two strings is the minimum number of editing operations (insertion, deletion, substitution) needed to transform one string into another.

Minimum Edit Distance

- For example the gap between the words intention and execution is five operations, shown in Fig. 3.23 as an alignment between the two strings.
- Given two sequences, an alignment is a correspondence between substrings of the two sequences. Thus I aligns with the empty string, N with E, T with X, and so on.
- Beneath the aligned strings is another representation; a series
 of symbols expressing an operation list for converting the top
 string into the bottom string; d for deletion, s for substitution,
 i for insertion.



Figure 3.23 Representing the minimum edit distance between two strings as an **alignment**. The final row gives the operation list for converting the top string into the bottom string; d for deletion, s for substitution, i for insertion.

MINIMUM EDIT DISTANCE

- Cost or weight can be assigned to each of these operations.
- The **Levenshtein distance** between two sequences is the simplest weighting factor in which each of the three operations has a cost of 1 (Levenshtein, 1966).
- Thus the Levenshtein distance between intention and execution is 5.
- Levenshtein also proposed an alternate version of his metric in which each insertion or deletion has a cost of one, and substitutions are not allowed (equivalent to allowing substitution, but giving each substitution a cost of 2, since any substitution can be represented by one insertion and one deletion).
- Using this version, the Levenshtein distance between intention and execution is 8.

MINIMUM EDIT DISTANCE

- The minimum edit distance is computed by dynamic programming.
- The intuition of a dynamic programming problem is that a large problem can be solved by properly combining the solutions to various sub-problems.

How to find the Min Edit Distance?

- Searching for a path (sequence of edits) from the start string to the final string:
 - Initial state: the word we're transforming
 - Operators: insert, delete, substitute
 - Goal state: the word we're trying to get to
 - Path cost: what we want to minimize: the number of edits

Defining Min Edit Distance

- For two strings
 - X of length n
 - Y of length m
- We define D(*i*, *j*)
 - the edit distance between X[1..i] and Y[1..j]
 - i.e., the first i characters of X and the first j characters of Y
 - The edit distance between X and Y is thus D(n,m)

MINIMUM EDIT DISTANCE

- Dynamic programming algorithms work by creating a distance matrix with one column for each symbol in the source sequence and one row for each symbol in the target sequence (i.e., source along the top, target along the side).
- For minimum edit distance, this matrix is the edit-distance matrix.
 Each cell edit-distance[i, j] contains the distance between the first i characters of the target and the first j characters of the source.
- Each cell can be computed as a simple function of the surrounding cells; thus starting from the beginning of the matrix it is possible to fill in every entry.
- The value in each cell is computed by taking the minimum of the three possible paths through the matrix which arrive there:

$$\begin{aligned} \textit{distance}[i,j] &= \min \left\{ \begin{array}{l} \textit{distance}[i-1,j] + \text{ del-cost}(\textit{source}_{i-1}) \\ \textit{distance}[i-1,j-1] + \text{subst-cost}(\textit{source}_{i-1},\textit{target}_{j-1}) \\ \textit{distance}[i,j-1] + \text{ ins-cost}(\textit{target}_{j-1}) \end{array} \right. \end{aligned}$$

Min-Edit distance Algorithm

```
function MIN-EDIT-DISTANCE(target, source) returns min-distance
     m ← LENGTH(target)
      n \leftarrow LENGTH(source)
     Create a distance matrix -distance[n+1, m+1]
     Initialize the zeroth row and column to be the distance
     from empty string
       distance[0,0] = 0
        for each row i from 1 to n do
           distance[i, 0] \leftarrow distance[i-1, 0] + ins-cost(target[i])
           for each column j from 1 to m do
           distance[0, j] \leftarrow distance[0, j-1] + del-cost(source[j])
       for each row i from 1 to n do
          for each column j from 1 to m do
          if ( i != j)
              distance[i, j] \leftarrow MIN (distance[i-1, j] + ins-cost(target_i),
                       distance[i-1, j-1] + subst-cost(source;, target;),
                       distance[i, j-1] + del-cost(source<sub>i</sub>))
         distance[i, j] \leftarrow distance[i-1, j-1]
      return distance[n, m]
```

Convert string abcdef to string azced

	#	a	b	С	d	е	f				
#	0	1	2	3	4	5	6				
a	1										
Z	2				(D	/i-1 i) :	∔ 1				
С	3		D($ D(i,j) = \min \begin{cases} D(i-1,j) + 1 \\ D(i,j-1) + 1 \end{cases} $							
е	4		D(i,j) = min D(i,j-1) + 1 D(i-1,j-1) + 1								
d	5		_		ا	(' +/) +	, . <u>.</u>				

Given source string: abcdef

	Source →											
Т	# a b c d e											
a	#											
r	a											
g	Z											
e	C											
t	е											
\	d											

Given source string: abcdef

	Source →											
Т		#	a	b	С	d	е	f				
a	#	0										
r	a											
g	Z											
e	С											
t	е											
1	d											

Given source string: abcdef

	Source →											
Т		#	a	b	С	d	е	f				
a	#	0	1									
r	а											
g	Z											
e	С											
t	е											
1	d											

Given source string: abcdef

	Source →											
T		#	а	b	С	d	е	f				
a	#	0	1	2								
r	a											
g	Z											
e	C											
t	е											
1	d											

Given source string: abcdef

	Source →											
Т		#	a	b	С	d	е	f				
a	#	0	1	2	3							
r	a											
g	Z											
e	C											
t	е											
1	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4						
r	a											
g	Z											
e	С											
t	е											
	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5					
r	a											
g	Z											
e	С											
t	е											
	d											

Given source string: abcdef

	Source →											
Т		# a	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	a											
g	Z											
e	С											
t	е											
\downarrow	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	a	1										
g	Z											
e	С											
t	е											
\	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	a	1										
g	Z	2										
e	C											
t	е											
	d											

Given source string: abcdef

	Source →											
Т		#	a	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	a	1										
g	Z	2										
e	С	3										
t	е											
	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1										
g	Z	2										
e	С	3										
t	е	4										
	d											

Given source string: abcdef

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1										
g	Z	2										
e	С	3										
t	е	4										
1	d	5										

Given source string: abcdef

Convert into target string: azced

			S	ource	\rightarrow			
Т		#	а	b	С	d	е	f
а	#	0	1	2	3	4	5	6
r	а	1	0					
g	Z	2						
e	С	3						
t	e	4						
	d	5						

Source: a

Given source string: abcdef

Convert into target string: azced

			S	ource	\rightarrow			
Т		#	а	b	С	d	е	f
a	#	0	1	2	3	4	5	6
r	а	1	0	1				
g	Z	2						
e	С	3						
t	е	4						
	d	5						

Source: a b

Given source string: abcdef

Convert into target string: azced

			S	ource	$e \rightarrow$			
Т		#	а	b	С	d	е	f
a	#	0	1	2	3	4	5	6
r	а	1	0	1	2			
g	Z	2						
e	С	3						
t	е	4						
1	d	5						

Source: abc

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
а	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3						
g	Z	2										
e	С	3										
t	е	4										
1	d	5										

Source: abcd

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4					
g	Z	2										
e	С	3										
t	е	4										
	d	5										

Source: abcde

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2										
e	С	3										
t	е	4										
1	d	5										

Source: abcdef

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2	1									
e	С	3										
t	е	4										
1	d	5										

Source: a

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2	1	1								
e	С	3										
t	е	4										
1	d	5										

Source: ab

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
а	#	0	1	2	3	4	5	6				
r	a	1	0	1	2	3	4	5				
g	Z	2	1	1	2							
e	С	3										
t	е	4										
1	d	5										

Source: abc

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3					
e	С	3									
t	е	4									
	d	5									

Source: abcd

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
а	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4				
e	С	3									
t	е	4									
1	d	5									

Source: abcde

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3									
t	е	4									
	d	5									

Source: abcdef

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	а	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2	1	1	2	3	4	5				
e	С	3	2									
t	е	4										
1	d	5										

Source: a

Given source string: abcdef

Convert into target string: azced

	Source →										
Т	T # a b c d										
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2							
t	е	4									
1	d	5									

Source: ab

Given source string: abcdef

Convert into target string: azced

			S	ource	$e \rightarrow$			
Т		#	a	b	С	d	е	f
a	#	0	1	2	3	4	5	6
r	а	1	0	1	2	3	4	5
g	Z	2	1	1	2	3	4	5
e	С	3	2	2	1			
t	е	4						
1	d	5						

Source: abc

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2					
t	е	4									
1	d	5									

Source: abcd

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		е	f								
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3				
t	е	4									
1	d	5									

Source: abcde

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
а	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4									
1	d	5									

Source: abcdef

Given source string: abcdef

Convert into target string: azced

	Source →											
Т		#	a	b	С	d	е	f				
а	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2	1	1	2	3	4	5				
e	С	3	2	2	1	2	3	4				
t	е	4	3									
1	d	5										

Source: a

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
а	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3							
1	d	5									

Source: ab

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2						
1	d	5									

Source: abc

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2					
1	d	5									

Source: abcd

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2				
1	d	5									

Source: abcde

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
\downarrow	d	5									

Source: abcdef

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
1	d	5	4								

Source: a

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
1	d	5	4	4							

Source: ab

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	а	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
е	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
<u> </u>	d	5	4	4	3						

Source: abc

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
\downarrow	d	5	4	4	3	2					

Source: abcd

Given source string: abcdef

Convert into target string: azced

			S	ource	$e \rightarrow$			
Т		#	a	b	С	d	е	f
a	#	0	1	2	3	4	5	6
r	а	1	0	1	2	3	4	5
g	Z	2	1	1	2	3	4	5
e	С	3	2	2	1	2	3	4
t	е	4	3	3	2	2	2	3
	d	5	4	4	3	2	3	

Source: abcde

Given source string: abcdef

Convert into target string: azced

	Source →										
Т		#	a	b	С	d	е	f			
a	#	0	1	2	3	4	5	6			
r	а	1	0	1	2	3	4	5			
g	Z	2	1	1	2	3	4	5			
e	С	3	2	2	1	2	3	4			
t	е	4	3	3	2	2	2	3			
\downarrow	d	5	4	4	3	2	3	3			

Source: abcdef

Given source string: abcdef

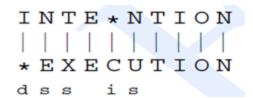
Convert into target string: azced

	Source →											
Т		#	a	b	С	d	е	f				
a	#	0	1	2	3	4	5	6				
r	а	1	0	1	2	3	4	5				
g	Z	2	1	1	2	3	4	5				
e	С	3	2	2	1	2	3	4				
t	е	4	3	3	2	2	2	3				
	d	5	4	4	3	2	3	3				

Source: abcdef

Example 2:

• Find the minimum number of edit operations required to convert **intention to execution**



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	Ι	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2									
Е	3									
С	4									
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: IN

Tar: E

Delete I, Subs: $N \leftarrow E$



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	Ι	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	2	3	4	5	6	7	8	9	10
X	2									
Е	3									
С	4									
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: INTENTION

Tar: E

Delete source



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	I	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3	4	5	6	7	8	9	10
X	2	2								
Е	3									
С	4									
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: I Tar: EX

Delete I, Insert E, X



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	Ι	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3									
С	4									
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: INT Tar: EX

Delete I, Subs: $T \leftarrow X$



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	I	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4									
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: INTE Tar: EXE

No subs cost



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: INTE Tar: EXEC Insert C



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5									
Т	6									
I	7									
0	8									
N	9									

Src: INTEN
Tar: EXECU

Subs:N← U



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5					8				
Т	6									
I	7									
0	8									
N	9									

Src: INTEN
Tar: EXECU
Subs:N← U



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5					8				
Т	6						8			
I	7									
0	8									
N	9									

Src: INTENT Tar: EXECUT



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5					8				
Т	6						8			
I	7							8		
0	8									
N	9									

Src: INTENTI Tar: EXECUTI



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	I	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5					8				
Т	6						8			
I	7							8		
0	8								8	
N	9									

Src: INTENTIO Tar: EXECUTIO



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

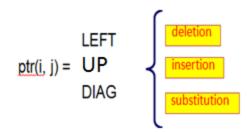
	#	Ι	N	Т	Е	N	Т	Ι	0	N
#	0	1	2	3	4	5	6	7	8	9
Е	1	1	3							
X	2			5						
Е	3				5					
С	4				6					
U	5					8				
Т	6						8			
I	7							8		
0	8								8	
N	9									8

Src: INTENTION Tar: EXECUTION



Insertion cost=1
Deletion Cost=1
Substitution Cost=2

	#	Ι	N	Т	Е	N	Т	Ι	0	N
#	0 ←	—1 _K	2	3	4	5	6	7	8	9
Е	1	1	3,							
X	2			5 ~						
Е	3				5					
С	4				6					
U	5					8,				
Т	6						8 ~			
I	7							8		
0	8								86	
N	9									8



Example 3:

 Find the minimum number of edit operations required to convert load to lord, graffe to giraffe