



Chance-Constrained Optimization applied to the Optimal Power Flow problem

A MIP approach

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June 10th, 2022

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Introduction to Chance-Constrained Problems

General chance-constrained SAA MIP reformulation

Tightening and Screening

Valid inequalities

Computational Results: OPF

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- **Optimal Power Flow**: minimize the expected operating cost whilst guaranteeing that the system withstands unforeseen peaks of electrical load due to stochastic demand.
- General (linear) formulation:

$$\begin{aligned} \min_x \quad & c^\top x \\ \text{s.t.} \quad & x \in X \\ & \mathbb{P} \left\{ a_j(\omega)^\top x \leq b_j(\omega), \forall j \right\} \geq 1 - \epsilon. \end{aligned}$$

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$$\sum_n B_{ln} \left(\sum_{g \in \mathcal{G}_n} (p_g - \Omega_s \beta_g) - d_n + \omega_{ns} \right) \geq -\bar{f}_l - y_s M_{ls}, \quad \forall l, s$$

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Algorithm 1 Iterative Coefficient Strengthening (κ iterations)

Initialization: $k \leftarrow 0$, $M_{js}^0 \leftarrow \infty$.

while $k < \kappa$ **do**

for $j \in \mathcal{J}$ and $s \in \mathcal{S}$ **do**

if $M_{js}^k > 0$ **then**

 1) Tightening phase: Solve

$$M_{js}^{k+1} \leftarrow \arg \max_{x,y} a_{js}^\top x - b_{js}$$

s.t. $x \in X$

$$x^\top a_j^0 + \Omega_s \hat{a}_j^\top x - b_{js} \leq M_{js}^k y_s, \quad \forall j, s$$

$$\sum_{s \in \mathcal{S}} y_s \leq p$$

$$0 \leq y_s \leq 1, \quad \forall s.$$

end if

if $M_{js}^{k+1} < 0$ **then**

 2) Screening phase: Eliminate constraint (j, s) from the model.

end if

end for

 Set $k \leftarrow k + 1$.

end while

Valid inequalities

For a fixed j :

$$x^\top a_j^0 + \Omega_s \hat{a}_j^\top x \leq b_{js}, \quad \forall s \in \mathcal{S}$$

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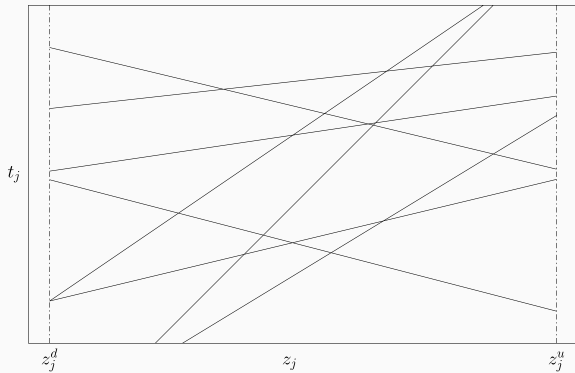
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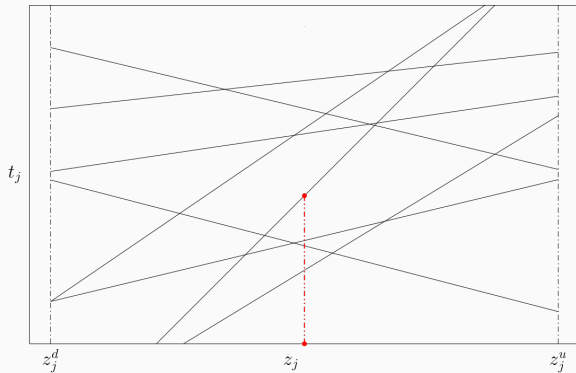
$$z_j := \hat{a}_j^\top x, \quad z_j \in [z_j^d, z_j^u]:$$

$$L_{js} : f_{js}(z_j) = \Omega_s z_j - b_{js}$$

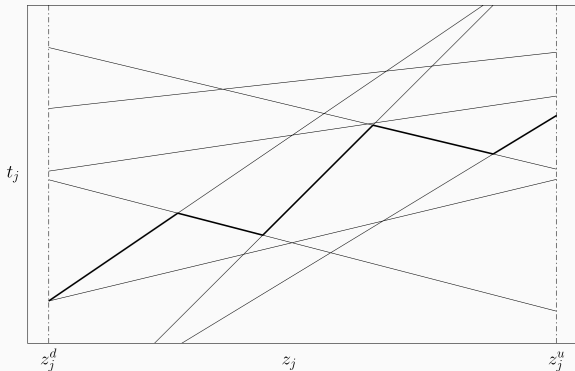
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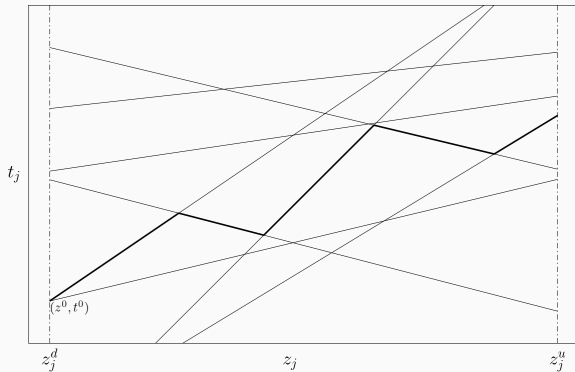
Proposition

For a fixed $j \in \mathcal{J}$, let $U_j^{p+1}(\cdot)$ be the $(p+1)$ -upper envelope of the set of lines \mathcal{L}_j , with $p := \lfloor \epsilon |\mathcal{S}| \rfloor$. Then the inequality

$$U_j^{p+1}(\hat{a}_j^\top x) + x^\top a_j^0 \leq 0, \quad x \in X$$

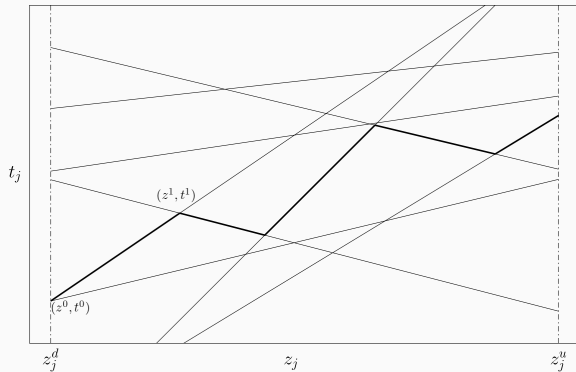
is valid.

Valid inequalities

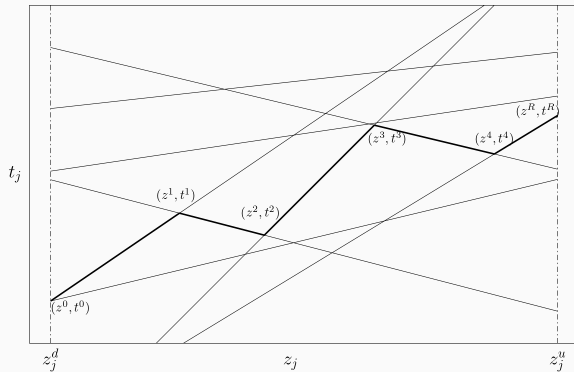


Rider Algorithm

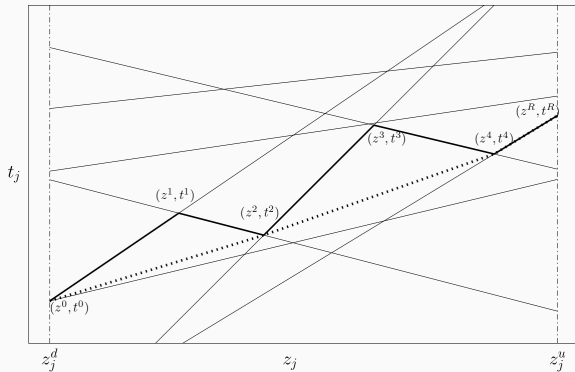
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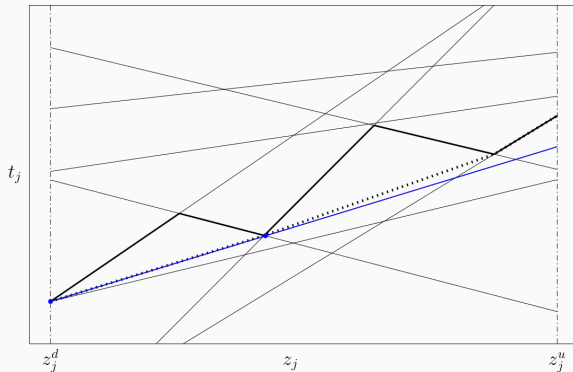


Valid inequalities



Lower hull: *Jarvis March, Graham scan*

Valid inequalities

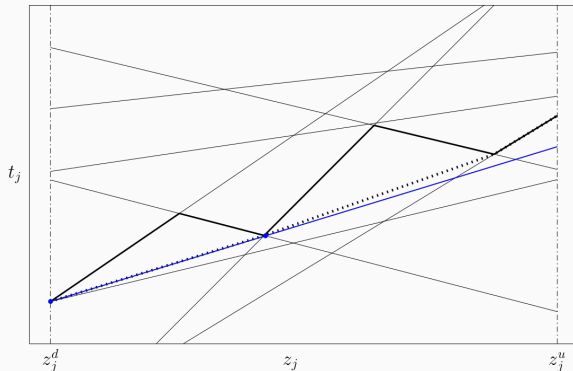


Corollary

Let $\{(z^r, t^r)\}$, $r \in \{0, \dots, R\}$, be an ordered set of vertices, and let $\{(z^{r'}, t^{r'})\}$, $r' \in \{0, \dots, R'\} \subseteq \{0, \dots, R\}$, be the ordered subset of vertices such that the associated polygonal chain is the lower hull. Then the following linear inequalities are valid:

$$\frac{t^{r'+1} - t^{r'}}{z^{r'+1} - z^{r'}} (\hat{a}_j^\top x - z^{r'}) + t^{r'} \leq -x^\top a_j^0, \quad x \in X, r' \in \{0, \dots, R' - 1\}$$

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- Relationship with *quantile cuts*.

Computational Results: OPF

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- Approaches **T**, **TS**, **V** and **TS+V** using five standard power systems.
- Instance: IEEE-300 test system: 300 nodes, 57 generators, 411 lines.
- GUROBI 9.1.2 on a Linux-based server with CPUs clocking at 2.6 GHz, 6 threads and 32 GB of RAM.
- 1000 scenarios, 5% violation of the JCC ($\epsilon = 0.05$, $p = 50$).
- Time limit: 10 hours.
- Results averaged over ten instances.

Computational Results

IEEE-300	BN	T(3)	TS(3)	BN+V	TS(1)+V
#CON	936939	100%	8.0%	101.5%	3.30%
LRgap	1.114%	0.264%	0.264%	0.3192%	0.1603%
MIPgap	0.27% (0)	0.07% (0)	0.01% (4)	0.08% (0)	0.00% (10)
Time	36000	1.0x	1.2x	1.0x	8.5x

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THANK YOU FOR YOUR ATTENTION