



GOBIERNO  
DE ESPAÑA

MINISTERIO  
DE UNIVERSIDADES

# Tight and Compact Sample Average Approximation for Joint Chance Constrained Optimal Power Flow

---

**Álvaro Porras**, Concepción Domínguez, Juan Miguel Morales, Salvador Pineda  
October 17<sup>th</sup>, 2022

OASYS Group, Universidad de Málaga

Introduction to Joint Chance Constrained Optimal Power Flow (JCC-OPF)

Sample Average Approximation MIP reformulation

Tightening and screening

Valid inequalities

Computational Results: OPF

# **Introduction to Joint Chance Constrained Optimal Power Flow (JCC-OPF)**

---

# Introduction to JCC-OPF

- **OPF under uncertainty:** minimize the expected operating cost whilst guaranteeing that the system withstands unforeseen peaks of electrical load due to stochastic demand.
- Chance constraints ensure feasibility of the system with a **tolerable probability of constraint violation**.
- Cost-effective decisions can be taken by discarding extreme events or unexpected random circumstances.
- General (linear) formulation:

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^\top \mathbf{x} \\ \text{s.t.} \quad & \mathbf{x} \in X \\ & \mathbb{P} \left\{ \mathbf{a}_j(\omega)^\top \mathbf{x} \leq b_j(\omega), \forall j \right\} \geq 1 - \epsilon. \end{aligned}$$

# Introduction to JCC-OPF

Assumptions considered:

- **Net demand uncertainty:**  $\tilde{d}_n = d_n - \omega_n$ , where  $d_n$  is the predicted value and  $\omega_n$  is the forecast error with a change of sign.
- **Generation:** To cope with the forecast errors  $(\omega_n)_{n \in \mathcal{N}}$ , generators' power outputs are adjusted according to the following affine control policy:

$$\tilde{p}_g = p_g - \beta_g \Omega, \quad \forall g \in \mathcal{G},$$

where  $\Omega := \sum_{n \in \mathcal{N}} \omega_n$  is the system-wise aggregated forecast error, and  $p_g$  and  $\beta_g$  are the power output dispatch and the participation factor of generating unit  $g$ .

- **Power Balance:**

$$\sum_{g \in \mathcal{G}} \tilde{p}_g - \sum_{n \in \mathcal{N}} \tilde{d}_n = \sum_{g \in \mathcal{G}} (p_g - \beta_g \Omega) - \sum_{n \in \mathcal{N}} (d_n - \omega_n) = 0.$$

$$\sum_{g \in \mathcal{G}} p_g - \sum_{n \in \mathcal{N}} d_n = 0 \text{ and } \sum_{g \in \mathcal{G}} \beta_g = 1.$$

- **DC Power Flow.**

# Introduction to JCC-OPF

The mathematical formulation is expressed as follows:

$$\begin{aligned} \min_{p_g, \beta_g} \quad & \mathbb{E} \left[ \sum_{g \in \mathcal{G}} C(p_g, \beta_g) \right] \\ \text{s.t.} \quad & \sum_{g \in \mathcal{G}} \beta_g = 1 \\ & \sum_{g \in \mathcal{G}} p_g - \sum_{n \in \mathcal{N}} d_n = 0 \\ & \mathbb{P} \left( \begin{aligned} & \underline{p}_g \leq p_g - \Omega \beta_g \leq \bar{p}_g, \quad \forall g \in \mathcal{G} \\ & -\bar{f}_l \leq \sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} (p_g - \Omega \beta_g) + \omega_n - d_n \right) \leq \bar{f}_l, \quad \forall l \in \mathcal{L} \end{aligned} \right) \geq 1 - \epsilon, \\ & p_g, \beta_g \geq 0, \forall g \in \mathcal{G}. \end{aligned}$$

# Chance-constrained SAA MIP reformulation (with Big-Ms)

- Sample Average Approx.: assume a finite discrete distribution  $\Rightarrow$  MIP reformulation
- $s \in \mathcal{S}$  scenarios (with equal probabilities)
- $y_s \in \{0, 1\}$  such that  $y_s = 0 \Leftrightarrow$  the scenario  $s$  is satisfied

$$\begin{aligned} \min_x \quad & c^\top x \\ \text{s.t.} \quad & x \in X \\ & \mathbb{P} \left\{ a_j(\omega)^\top x \leq b_j(\omega), \forall j \right\} \geq 1 - \epsilon. \\ & \sum_{s \in \mathcal{S}} y_s = p \\ & y_s \in \{0, 1\} \quad \forall s. \end{aligned}$$

$$\begin{aligned} \min_x \quad & c^\top x \\ \text{s.t.} \quad & x \in X \\ & a_{js}^\top x \leq b_j(\omega) + M_{js} y_s, \quad \forall j, s \\ & \sum_{s \in \mathcal{S}} y_s \leq p \\ & y_s \in \{0, 1\}, \quad \forall s. \end{aligned}$$

# JCC-OPF via Sample Average Approximation

The MIP reformulation of JCC-OPF writes as follows:

$$\begin{aligned} \min_{p_g, \beta_g} \quad & \mathbb{E} \left[ \sum_{g \in \mathcal{G}} C(p_g, \beta_g) \right] \\ \text{s.t.} \quad & \sum_{g \in \mathcal{G}} \beta_g = 1 \\ & \sum_{g \in \mathcal{G}} p_g - \sum_{n \in \mathcal{N}} d_n = 0 \\ & -y_s M_{gs}^1 + \underline{p}_g \leq p_g - \Omega_s \beta_g \leq \bar{p}_g + y_s M_{gs}^2, \quad \forall g, s \\ & -y_s M_{ls}^3 - \bar{f}_l \leq \sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} (p_g - \Omega_s \beta_g) - d_n + \omega_{ns} \right) \leq \bar{f}_l + y_s M_{ls}^4, \quad \forall l, s \\ & p_g, \beta_g \geq 0, \forall g \in \mathcal{G}. \\ & \sum_{s \in \mathcal{S}} y_s \leq p \\ & y_s \in \{0, 1\}, \quad \forall s. \end{aligned}$$



## Tightening and screening

---

# Tightening and screening

---

## Algorithm 1 Iterative Coefficient Strengthening ( $\kappa$ iterations)

---

**Initialization:**  $k \leftarrow 0$ ,  $M_{js}^0 \leftarrow \infty$ .

**while**  $k < \kappa$  **do**

**for**  $j \in \mathcal{J}$  and  $s \in \mathcal{S}$  **do**

**if**  $M_{js}^k > 0$  **then**

            1) Tightening phase: Solve

$$\begin{aligned} M_{js}^{k+1} &\leftarrow \arg \max_{x,y} a_{js}^\top x - b_{js} \\ \text{s.t. } &x \in X \\ &a_{js}^\top x - b_{js} \leq M_{js}^k y_s, \quad \forall j, s \\ &\sum_{s \in \mathcal{S}} y_s \leq p \\ &0 \leq y_s \leq 1, \quad \forall s. \end{aligned}$$

**end if**

**if**  $M_{js}^{k+1} < 0$  **then**

            2) Screening phase: Eliminate constraint  $(j, s)$  from the model.

**end if**

**end for**

    Set  $k \leftarrow k + 1$ .

**end while**

## Valid inequalities

---

# Valid inequalities I: generators

For a given  $g$  (the process is analogous for the upper bound constraint)

- $p_g - \Omega_s \beta_g \geq \underline{p}_g \Rightarrow p_g - \underline{p}_g \geq \Omega_s \beta_g$
- Let us consider a set of scenarios  $\Omega := \{2, -4, 4, 5, 0\}$
- Acceptable violation probability is 40%. Then, if we have 5 scenarios, a constraint can be violated at most in 2 scenarios.

$$p_g - \underline{p}_g \geq 2 \beta_g$$

$$p_g - \underline{p}_g \geq -4 \beta_g$$

$$p_g - \underline{p}_g \geq 4 \beta_g$$

$$p_g - \underline{p}_g \geq 5 \beta_g$$

$$p_g - \underline{p}_g \geq 0 \beta_g$$

\*It is like ordering 1-dimensional affine functions without intercept.

# Valid inequalities I: generators

For a given  $g$  (the process is analogous for the upper bound constraint)

- $p_g - \Omega_s \beta_g \geq \underline{p}_g \Rightarrow p_g - \underline{p}_g \geq \Omega_s \beta_g$
- Let us consider a set of scenarios  $\Omega := \{2, -4, 4, 5, 0\}$
- Acceptable violation probability is 40%. Then, if we have 5 scenarios, a constraint can be violated at most in 2 scenarios.

$$p_g - \underline{p}_g \geq 2 \beta_g$$

$$p_g - \underline{p}_g \geq -4 \beta_g$$

$$p_g - \underline{p}_g \geq 4 \beta_g$$

$$p_g - \underline{p}_g \geq 5 \beta_g$$

$$p_g - \underline{p}_g \geq 0 \beta_g$$

- In a descending order, the constraint with the third largest  $\Omega$  must be satisfied.
- Therefore, the following constraint is a valid inequality.

$$p_g - \underline{p}_g \geq 2 \beta_g$$

\*It is like ordering 1-dimensional affine functions without intercept.

## Valid inequalities II: lines

For a given  $l$  (the process is analogous for the upper bound constraint)

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} (p_g - \Omega_s \beta_g) + \omega_{ns} - d_n \right) \geq -\bar{f}_l$$

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq \Omega_s \sum_{n \in \mathcal{N}} B_{ln} \sum_{g \in \mathcal{G}_n} \beta_g - \sum_{n \in \mathcal{N}} B_{ln} \omega_{ns}$$

Per each line constraint, the right-hand side is an 1-dimensional affine function per scenario as follows:

$$t_{ls} = \Omega_s z_l + b_{ls}$$

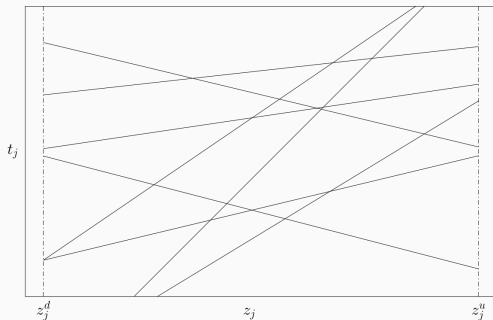
where:

$$z_l = \sum_{n \in \mathcal{N}} B_{ln} \sum_{g \in \mathcal{G}_n} \beta_g \quad \text{and} \quad \underline{z}_l \leq z_l \leq \bar{z}_l$$

$$b_{ls} = \sum_{n \in \mathcal{N}} B_{ln} \omega_{ns}$$

## Valid inequalities II: lines

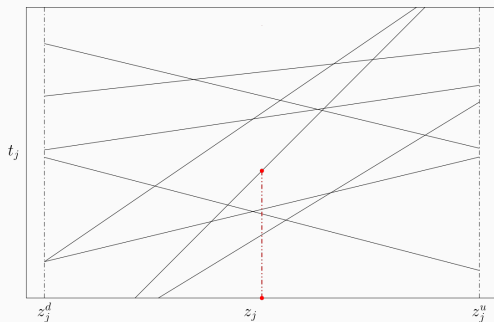
$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



Let us consider an example of 8 scenarios where the constraints can be violated in at most 4 scenarios.

## Valid inequalities II: lines

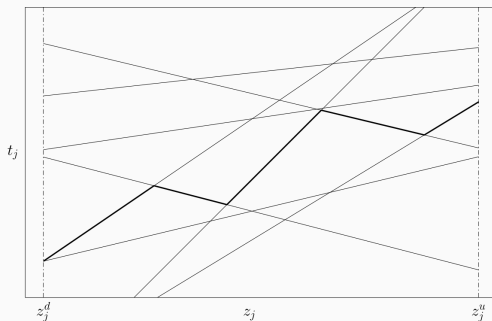
$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$





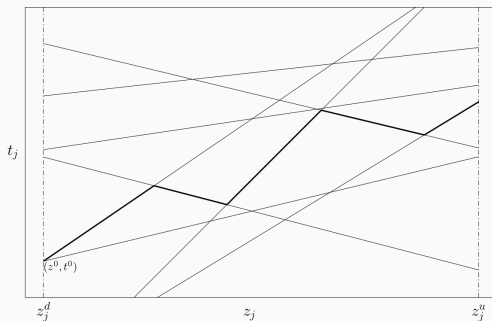
## Valid inequalities II: lines

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



## Valid inequalities II: lines

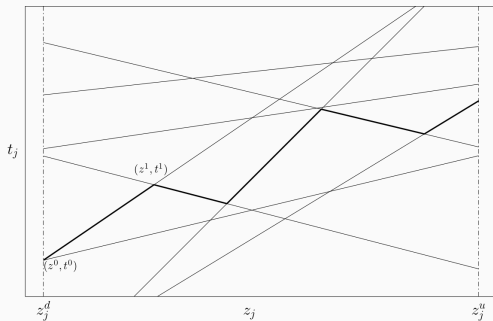
$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



Rider Algorithm

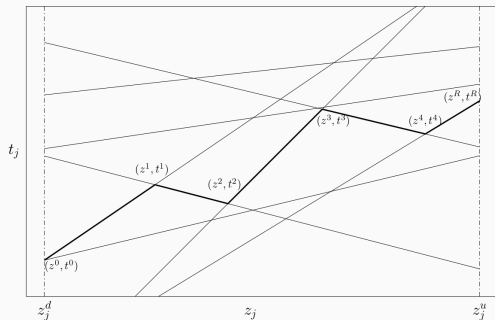
## Valid inequalities II: lines

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



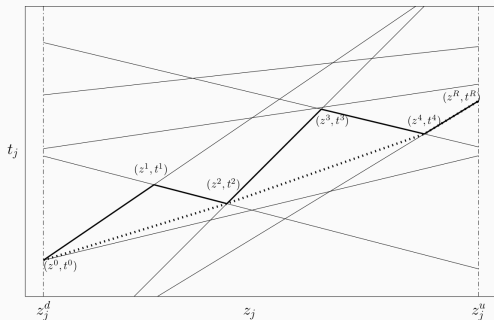
## Valid inequalities II: lines

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



# Valid inequalities II: lines

$$\sum_{n \in \mathcal{N}} B_{ln} \left( \sum_{g \in \mathcal{G}_n} p_g - d_n \right) + \bar{f}_l \geq t_{ls} = \Omega_s z_l + b_{ls}$$



Lower hull: *Jarvis March, Graham scan*

## Computational Results: OPF

---

- Approaches **T**, **TS**, **V** and **TS+V** using five standard power systems.
- Instance: IEEE-118 test system: 118 nodes, 19 generators, 186 lines.
- GUROBI 9.1.2 on a Linux-based server with CPUs clocking at 2.6 GHz, 6 threads and 32 GB of RAM.
- 1000 scenarios, 5% violation of the JCC ( $\epsilon = 0.05$ ,  $p = 50$ ).
- Time limit: 10 hours.
- Results averaged over ten instances.

# Computational Results

IEEE-118	BN	T(3)	TS(3)	BN+V	TS(1)+V
#CON	410413	100%	5.6%	101.6%	2.68%
LRgap	0.956%	0.434%	0.434%	0.4784%	0.2821%
MIPgap	0.29% (0)	0.12% (0)	0.01% (6)	0.03% (1)	0.00% (10)
Time	36000	1.0x	1.4x	1.1x	23.1x



# Comparison

	Methods	IEEE-118
Average cost increase	<b>TS+V</b>	0.00%
	CVaR	0.57%
	ALSO-X	0.08%
	ALSO-X+	0.05%
Speedup factor	<b>TS+V</b>	23.1x
	CVaR	4045.5x
	ALSO-X	148.6x
	ALSO-X+	49.3x

# The End

Á. Porras, C. Domínguez, J.M. Morales, and S. Pineda. (2022) Tight and compact sample average approximation for joint chance-constrained optimal power flow. *arXiv preprint arXiv:2205.03370*.



THANK YOU FOR YOUR ATTENTION