### Machine-Learning-aided Optimal Transmission Switching

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JOINT WORK WITH: Juan Miguel Morales González Salvador Pineda Morente









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XL Congreso Nacional de Estadística e Investigación Operativa (SEIO 2023)

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### Outline

- Motivation
- 2 Methodology
- 3 Computational Experience
- 4 Conclusions and Further Research

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Mathematical Optimization

Machine Learning

#### Combine knowledge from both worlds

Recent reviews: Bengio et al. [2021]; Gambella et al. [2021]

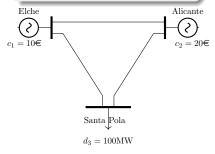
Mathematical Optimization

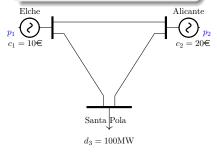
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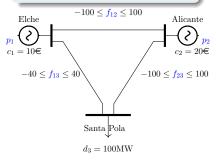
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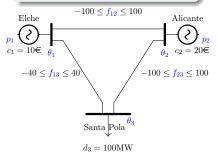
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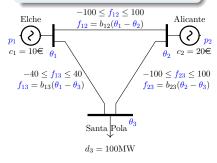
Power Systems Operations

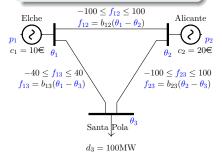






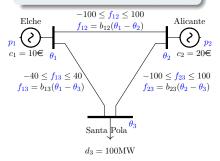






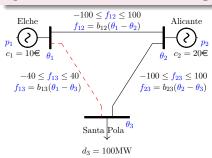
- Balance constraints.
- Assume  $b_{nm} = 1, \forall (n, m)$ .
- Minimizing costs.  $Cost^* = 1800 \in$ .
- $p_1^* = 20, p_2^* = 80.$
- $\theta_1^* = 40, \, \theta_2^* = 60, \, \theta_3^* = 0.$
- $f_{12}^* = -20, f_{13}^* = 40, f_{23}^* = 60.$

#### Optimal Power Flow (OPF)

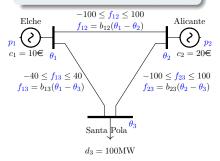


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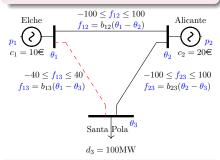


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#### Optimal Transmission Switching



- Switchable lines,  $\mathcal{L}_S$ .
- Binary variables  $x_{nm}, \forall (n,m) \in \mathcal{L}_S$ .
- Minimizing costs.  $Cost^* = 1000 \in$ .
- $p_1^* = 100, p_2^* = 0.$
- $\theta_1^* = 200, \, \theta_2^* = 100, \, \theta_3^* = 0.$

 $\begin{cases} \min \limits_{p_n,f_{nm},\theta_n,x_{nm}} & \sum_n c_n \, p_n \\ \text{s.t.} & \underbrace{p_n \leq p_n \leq \overline{p}_n, \quad \forall n \in \mathcal{N}} \\ & \sum_{(n,m) \in \mathcal{L}_n^-} f_{nm} - \sum_{(n,m) \in \mathcal{L}_n^+} f_{nm} = p_n - d_n, \quad \forall n \in \mathcal{N} \\ & f_{nm} = x_{nm} b_{nm} (\theta_n - \theta_m), \quad \forall (n,m) \in \mathcal{L}_{\mathcal{S}} \\ & -x_{nm} \overline{f}_{nm} \leq f_{nm} \leq x_{nm} \overline{f}_{nm}, \quad \forall (n,m) \in \mathcal{L}_{\mathcal{S}} \\ & f_{nm} = b_{nm} (\theta_n - \theta_m), \quad \forall (n,m) \in \mathcal{L} \setminus \mathcal{L}_{\mathcal{S}} \\ & -\overline{f}_{nm} \leq f_{nm} \leq \overline{f}_{nm}, \quad \forall (n,m) \in \mathcal{L} \setminus \mathcal{L}_{\mathcal{S}} \\ & x_{nm} \in \{0,1\}, \quad \forall (n,m) \in \mathcal{L}_{\mathcal{S}} \\ & \theta_1 = 0 \end{cases}$ 

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- MINLP because  $x_{nm}(\theta_n \theta_m)$ .
- NP-hard problem.

#### Original constraint

$$f_{nm} = x_{nm}b_{nm}(\theta_n - \theta_m)$$

#### Reformulation using big-Ms

$$b_{nm}(\theta_n - \theta_m) - \overline{M}_{nm}(1 - x_{nm}) \le f_{nm} \le b_{nm}(\theta_n - \theta_m) - \underline{M}_{nm}(1 - x_{nm})$$

#### How to compute bigM values?

$$\underline{M}_{nm} \le \underline{M}_{nm}^{\text{OPT}} := b_{nm} \times \min_{x_{nm}=0} (\theta_n - \theta_m)$$

$$\overline{M}_{nm} \ge \overline{M}_{nm}^{\text{OPT}} := b_{nm} \times \max_{x_{nm}=0} (\theta_n - \theta_m)$$

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- Computing optimal bigM values is as hard as solving the original problem, [Fattahi et al., 2019].
- We have to be happy with bounds.
- Objective: Find good bounds for bigMs.

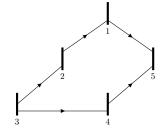
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- (4) Conclusions and Further Research

- Exact approach.
- Benchmark.
- Shortest path problem (Dijkstra).
- [Fattahi et al., 2019].

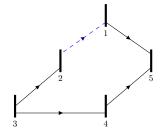
$$-\underline{M}_{nm} = \overline{M}_{nm} = b_{nm} \sum_{(k,l) \in \mathrm{SP}_{nm}} \frac{\overline{f}_{kl}}{b_{kl}}, \quad \forall (n,m) \in \mathcal{L}_{\mathcal{S}}$$

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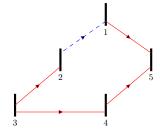
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$$\overline{M}_{21} \ge b_{21}(\theta_2 - \theta_1) \text{ if } x_{21} = 0$$

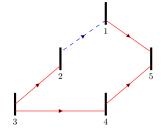
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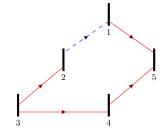
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$$f_{23} \le \overline{f}_{23} \implies b_{23}(\theta_2 - \theta_3) \le \overline{f}_{23} \implies (\theta_2 - \theta_3) \le \overline{f}_{23}/b_{23}$$

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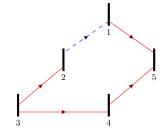
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$$(\theta_2 - \theta_1) \le \overline{f}_{23}/b_{23} + \overline{f}_{34}/b_{34} + \overline{f}_{45}/b_{45} + \overline{f}_{51}/b_{51}$$

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$$\overline{M}_{21} \ge b_{21}(\overline{f}_{23}/b_{23} + \overline{f}_{34}/b_{34} + \overline{f}_{45}/b_{45} + \overline{f}_{51}/b_{51})$$

## Strategy 2, Knn-D

- $\bullet$  Data-driven approach (Knn).
- Naive approach. Learning this problem is a challenge.

# Strategy 2, Knn-D

- Data-driven approach (Knn).
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#### Algorithm

- 1) Training set  $\mathcal{T} = \{(\mathbf{d}^t, \mathbf{x}^t)\}$  for  $\forall t$ .
- 2) For a given test demand  $\mathbf{d}^{\hat{t}}$ , compute K closest neighbors,  $\mathcal{T}_K$ .
- 3) Binary  $\mathbf{x}^{\hat{t}}$  as the rounded mean of the binary decision values from  $\mathcal{T}_K$  to the closest integer.
- 4) Solve an LP from the OTS by fixing variables.

### Strategy 3, Knn-BM

- Data-driven approach (Knn).
- One of our proposals.

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- 3) Fixing the binary variables if there unanimity in the value among the instances in  $\mathcal{T}_K$ .
- **4)** The bigM values of the remaining variables are updated using the shortest path.
- 5) Solve the resulting MILP.

# Strategy 4, Knn-B $\widehat{M}$

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- 4) The bigM values of the remaining variables are updated using the angles information.

$$\overline{M}_{nm} = b_{nm} \times \max_{t \in \mathcal{T}: x_{nm}^t = 0} (\theta_n^t - \theta_m^t)$$

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**5)** Solve the resulting MILP.

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# Experimental Setup

- Realistic network, [Blumsack, 2006].
- 118 buses 186 lines.
- $|\mathcal{L}_S| = 69$ .
- 500 instances.
- Demand follows uniform distribution in  $[0.9d_n, 1.1d_n]$ .
- Gurobi 9.1.2.
- Gap = 0.01%.
- Time limit: 1 hour.
- Github.

	# opt	# sub	gap-max	time (s)
BEN	500	0	-	145.44
Knn-D	0	500	14.13	0.0
Knn-BM	500	0	-	12.33
$K$ nn-B $\widehat{\mathbf{M}}$	495	5	0.39	0.7

Table: All approaches (K = 50)

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#### BEN vs Knn-D

- Knn-D has no optimal instances.
- $\bullet$  Knn-D is faster.
- Max gap: 14%

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#### BEN vs Knn-BM

- All optimal instances
- $\bullet$  Knn-BM is faster.
- No gap.

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Table: All approaches (K = 50)

## BEN vs Knn-B $\widehat{M}$

- Knn-B $\widehat{M}$  has almost all optimal instances
- Knn-B $\widehat{M}$  is very fast.
- Small gap.

	# opt	# sub	gap-max	time (s)
BEN	500	0	-	145.44
Knn-D	0	500	14.13	0.0
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Table: All approaches (K = 50)

## Knn-BM vs Knn-B $\widehat{M}$

- Knn-BM is more conservative than Knn-B $\widehat{M}$  but slower.
- Trade-off decision.
- Competitive against existing approaches.

### More details

#### Learning-Assisted Optimization for Transmission Switching

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#### Available at:

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#### Conclusions

- OTS is a challenging problem.
- Useful to reduce costs.
- Find good bigM values for the reformulation.
- Learning strategies.
- Tested on a real-world network.

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#### Further research

- All switchable lines.
- Other ML approaches.

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#### Thank you very much for your attention!









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