Reverse Engineering For Beginners



Linux x86_64





Intro

5 - 7 min.



02.

x86_64 Basics

(Theory) approx. 40 min.



CTF





RE Basics

Lab & Demos approx. 1 hour.









#Whoami

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- Organizer BSides Singapore
- 9+ years in Infosec
- Interested in vulnerability research (OpenSource)
- Learning solidity smart contract security
- Tweet me @samanl33t







WHAT - WHY - HOW ? RE

At the core: To dismantle an object to see how it works.

For software: (roughly) To disassemble the compiled code to understand it's logic with the intention to circumvent or replicate it.

WHY?

- Military use cases
- Vulnerability Research
- Malware Analysis
- CTFs and curiosity
- Historically:
 - Cracks
 - > Keygens

In the context of Linux ELF binaries:

• Programming/ Development process:

Code (mainly C/C++) > Compile > Assemble

Generates an ELF Binary.

Reverse Engineering process:

Disassemble (ELF Binary) > Decompile > Lots and lots of Reading > 5 stages of grief > Voila!!







Theory







X86_64 Basics (Theory)



- The lowest level of programming language for CPUs
- 64-bit version of x86 instruction set
- Consists of
 - Registers (rbp, rsp etc.)
 - Instructions (mov, cmp, jmp etc.)
 - Memory (stack, program data, heap etc.)







REGISTERS

- Special kind of super fast memory to work with instructions.
- 64-bit/8 byte wide
- Can be accessed partially
 - For e.g: Lower WORD (2 bytes), DWORD (4 bytes)
- Must know types:
 - General Purpose Registers (16)
 - Pointer Registers (1)
 - RFlags (3)







General Purpose Registers

<u>64-bit</u>	<u>32-bit</u>	<u> 16-bit</u>	<u>8H</u>	<u>8L</u>	<u>Description</u>	
rax	eax	ax	ah	al	Accumulator Register	
rbx	ebx	bx	bh	bl	Base Register	
rcx	есх	СХ	ch	cl	Counter Register	
rdx	edx	dx	dh	dl	Data Register	
rsp	esp	sp		spl	Stack Pointer	
rbp	ebp	bp		bpl	Base Pointer	
rsi	esi	si		sil	Source Index	
rdi	edi	di		dil	Destination index	





General Purpose Registers

<u>64-bit</u>	32-bit	<u> 16-bit</u>	<u>8H</u>	<u>8L</u>	<u>Description</u>	
r8	r8D	r8W		r8B	General Purpose	
r9	r9D	r9W		r9B	General Purpose	
r10	r10D	r10W		r10B	General Purpose	
r11	r11D	r11W		r11B	General Purpose	
r12	r12D	r12W		r12B	General Purpose	
r13	r13D	r13W		r13B	General Purpose	
r14	r14D	r14W		r14B	General Purpose	
r15	r15D	r15W		r15B	General Purpose	





Pointer Registers & RFlags

<u>64-bit</u>	<u>32-bit</u>	<u> 16-bit</u>	<u>8H</u>	<u>8L</u>	<u>Description</u>
rip	eip	ip			Instruction Pointer

RFlag	<u>Description</u>			
CF	Carry Flag			
SF	Sign FLag			
ZF	Zero Flag			





INSTRUCTIONS

		OPERAND .	OPERAND
	 Operation/What to do		with/on
Fa	mov	rax r	sn







INSTRUCTIONS

Different Types:

- Data Manipulation:
 - o mov rsp,rax
 - o inc [rax]
 - Arithmetic ops.: add rax, r9 / sub rax,r9 / mul rax, r9
- Control Flow & Conditions:
 - Jumps: je, jne, jg,jge, jle etc.
 - o cmp, test etc.
- System Calls:
 - To interact with the OS.
 - System call (syscall) numbers are used (rax)
 - Eg: Read, write, open etc.







MEMORY

Stack

- Last-in First-Out (LIFO)
- Grows toward lower memory addresses/backwards.
- Used to store data temporarily during program execution. Eg:
 - Storing and using function return values (push & pop instructions)
 - Stores function local variables
 - Sometimes, used to pass function arguments.

Heap/Others

- Dynamically mapped regions in the memory:
 - Eg. using malloc(), mmap() etc.

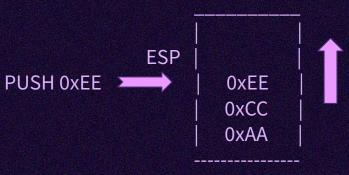






Stack Layout

0x00000000 (Lower Memory address)



0xFFFFFFFF (Higher Memory Address)

- PUSH: Pushes 8 bytes to the top of stack
- POP: Pops 8 bytes from the top of stack
- RSP: Stack pointer gets updated based on the instructions. (+/- 8 bytes)





Endianess

- The order in which data bytes are arranged into larger numerical values
- Intel x86_64 CPUs follow Little Endian format
 - i.e the bytes are ordered from least significant bit.

	Low a	ddress	High address					
Address	0	1	2	3	4	5	6	7
Little-endian	Byte 0	Byte 1	Byte 2	Byte 3	Byte 4	Byte 5	Byte 6	Byte 7
Big-endian	Byte 7	Byte 6	Byte 5	Byte 4	Byte 3	Byte 2	Byte 1	Byte 0
Memory content	0x22	0x33	0x44	0x55	0x66	0x77	0x88	
64 bit value on Little-endian 64 bit value on Big-endian								
0x8877665544332211 0x1122334455667788								





x86_64 Calling Conventions

How a function call is made in Assembly?

- Passing arguments:
 - Using registers: rdi, rsi, rdx, rcx, r8, r9
 - Uses stack: When there are many arguments.
- Return values:
 - returned in rax
- Callee-saved registered: rbx, rbp, r12, r13, r14, r15

```
int sum (a,b) {
Int c;
c = a + b;
return c;
int result:
result = sum(1,5);
// rdi = 1;
// rsi = 5
// rax = 6 = result
```





X86_64 Linux Syscalls

- For the program to communicate with outside world (OS)
- Approx. 330 syscalls in linux
- Calling convention:
 - Syscall number passed to rax
 - Pass arguments in registers (rdi, rsi,rdx etc.)
 - o call 'syscall' instruction.
 - Eg: system call: exit(1);

Syscall Table: https://x64.syscall.sh/

```
mov rax, 0x3C
mov rdi, 1
syscall
```







x86_64 Functions & Frames

Simply put:

- Every program consists of multiple functions (within different modules)
 - Modules can be internal or external
- Each function is a combination of multiple blocks of instructions
 - Instructions => Assembly instructions
- The function usually has a specific goal. Eg:
 - read/write/update data
 - call other functions etc.
- The combination of blocks can be represented as Control Flow Graph.
 - Where each instruction is executed one after another
 - And the flow is defined by different conditions





Control Flow Graph

```
loc 1340:
                                                      eax, [rbp+var 18]
                                                      [rbp+var_10], eax
short loc_131E
                                              ile
                                                                     eax, [rbp+var C]
mov
        esi, eax
                                                                     loc 131E:
mov
lea
        rdi, aTotalSumOfTheN ; " Total Sum of the numbers = %d
                                                                     lea
                                                                             rax, [rbp+var 14]
mov
                                                                     mov
                                                                             rsi, rax
call
        printf
                                                                     lea
                                                                             rdi, aD 0
mov
        eax, 0
                                                                     mov
                                                                             eax, 0
        rdx, [rbp+var 8]
                                                                     call
                                                                                isoc99 scanf
mov
        rdx, fs:28h
                                                                             eax, [rbp+var_14]
xor
                                                                     mov
                                                                             [rbp+var_C], eax
[rbp+var_10], 1
jz
        short locret 1377
                                                                     add
                                                                     add
                   call
                                stack chk fail
                           M M
                           locret 1377:
                           leave
                           retn
                           addition endp
```





x86_64 Functions & Frames

Before a function is called:

- The address to return to is pushed on the stack.
- Function arguments are moved to registers (sometimes on stack if required)

```
loc_1229:
mov eax, 0
call addition
jmp short loc_1281
```



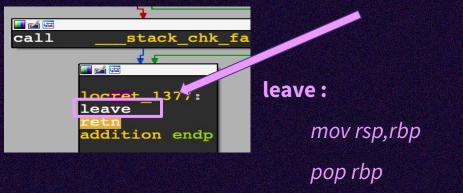




X86_64 Functions & Frames

When a function is called:

- Sets the stack frame Prologue
- Executes the Instructions inside
- Tears down the stack frame Epilogue



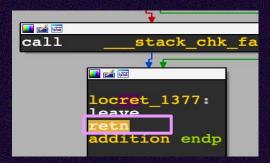
```
endbr64
        rbp
oush
mov
        rbp, rsp
sub
mov
        rax, is:28
        [rbp+var 8], rax
mov
xor
        eax, eax
        [rbp+var C], 0
mov
lea
        rdi, aHowManyNumbers ; " How many numbers you w
        eax, 0
mov
call
        printf
        rax, [rbp+var 18]
lea
mov
        rsi, rax
lea
        rdi, aD
mov
        eax, 0
call.
           isoc99 scanf
        rdi, aEnterTheNumber; " Enter the numbers: \n
lea
mov
call
        printf
mov
        [rbp+var_10], 1
jmp
        short loc 1340
                       loc 1340:
                               eax, [rbp+var 18]
                               [rbp+var 10], eax
                               short loc 131E
```



X86_64 Functions & Frames

When all instructions in the function are executed

- A return instruction is called, which:
 - Pops the previously stored return address from stack









Function Stack Frame

0x00000000

Function variables etc.

Base pointer for original stack frame

Address to return to after function execution

STACK

local data

local data

rbp (original)

retn address

args (if any)

env data

0xFFFFFFF





Reverse Engineering -Basics (Lab)



RE Basies Lab Intro/Setup

Lab Instructions

- Clone the repository:
 - https://github.com/samanL33T/x86_64_RE_For_Beginners_STANDCON_
 2022.git
- System requirements:
 - Any linux system (Ubuntu Desktop latest build preferred)
- Tools:
 - o gcc, file, strings, objdump, ltrace, strace installed
 - o IDA Free/Binary Ninja/Ghidra Installed (We'll use IDA Free for workshop)
 - Binary Ninja Cloud can also be used (<u>https://cloud.binary.ninja/</u>)
 - or pwndbg (https://github.com/hugsy/gef)
 or pwndbg (https://github.com/pwndbg/pwndbg/pwndbg/pwndbg)



RE Basies Lab Intro/Setup

Lab Instructions & Goals

- CTF styled target
 - Practice target available in git repo (/CTF_Binary/KnightInit)
 - CTF target available as hosted binary later.
- Go through each phase of Reverse Engineering process (each Lab)
- Find and submit flags for each stage

Lab worksheet provided in the git repo







RE Basies Lab - Basic Analysis

Stage 0

Analyse the target file for following:

- Binary type
- Architecture
- Linking
- Stripped/non-stripped
- Hard-coded secrets

Tools

- file
- strings







RE Basies Lab - Static Analysis

Stage 1 & 2

- Disassemble the target using objdump
- Familiarize with objdump output
- Disassemble the target using IDA Free (or Ghidra/Binary Ninja)
- Understand the control flow.

Tools:

- objdump
- IDA Free/Ghidra/Binary Ninja







RE Basics Lab - Dynamic Analysis

Stage 3

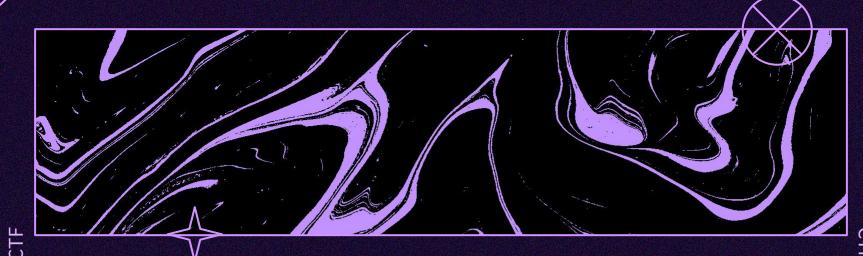
- Tracing the target libraries and system calls
- Understanding the control flow of target during runtime
- Quick primer on and familiarizing with gdb

Tools:

- ltrace
- strace
- gdb (gef/pwndbg)











Reverse Engineering - CTF



STC_50f5a4c7355fb82a67b3c18624e837e2



RE CTF

Stage 4 onward..

• The CTF Binary is hosted on:

<IP>:1337

- Only the first hardcoded flag is same for local and hosted binary.
- Get the real flags from hosted binary.

Total Flags: 5

- Use the gained knowledge to find more flags in the target
- Submit the flags via Twitter or Discord DM.







More Learning RESOURCES & References

- PWN College https://pwn.college/
- @mytechnotalent: https://0xinfection.github.io/reversing/

https://github.com/mytechnotalent/Reverse-Engineering

Practice RE with crackmes: https://crackmes.one/











THANK YOU!







Reach out to me at:

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THANKS!



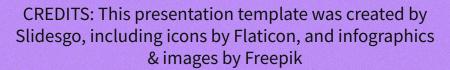








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Marketing plan

