



# **COL215 DIGITAL LOGIC AND SYSTEM DESIGN**

Number Representation and  
Arithmetic Operations

05 September 2017

# Values and Representations

- Values : Elements of a domain
  - integers
  - real numbers
  - letters of an alphabet
  - characters
  - logic values
- Representations : Arrangements of symbols
  - bits  $\{0,1\}$  can be used to represent values
  - from a variety of domains

# Positional representation of numbers

- Consider only non-negative integers at present (unsigned)
- Radix or base =  $r$

Representation :  $d_{n-1} d_{n-2} \dots d_2 d_1 d_0$

$$0 \leq \text{val}(d_i) \leq r-1$$

Example:

Representation                      “1110” , radix 2

Value                                      14

# From representation to value

Representation :  $d_{n-1} d_{n-2} \dots d_2 d_1 d_0$

$$0 \leq \text{val}(d_i) \leq r-1$$

Value :  $\sum_{i=0}^{n-1} \text{val}(d_i) \cdot r^i$

or simply  $\sum_{i=0}^{n-1} d_i \cdot r^i$  [if it is clear that  $d_i$  denotes value of  $d_i$ ]

# Another way: from representation to value

$$((\dots(d_{n-1} \times r + d_{n-2}) \times r + \dots + d_1) \times r) + d_0$$

# From Value to Representation

Given value =  $V$ , radix =  $s$

- $\text{val}(d_0) = V \bmod s$
  - $\text{val}(d_1) = (V \text{ div } s) \bmod s$
  - $\text{val}(d_2) = ((V \text{ div } s) \text{ div } s) \bmod s$
  - ...
  - $\text{val}(d_i) = ( \dots (V \text{ div } s) \dots ) \bmod s$
  - ...
- $\underbrace{\hspace{10em}}$   
i times

# Conversion from radix $r$ to $s$

Do computation in radix  $r$  (division method)

- start with value in radix  $r$  (same as representation in  $r$ )
- compute the digit values with radix  $= s$  (from value to representation)
- put together these digits (representations, not values) to form the representation in radix  $s$



# Conversion from radix $r$ to $s$

Do computation in radix  $s$  (multiplication method)

- start with representation in radix  $r$
- this gives the digits
- from digits compute the value (from representation to value)
- the value expressed in radix  $s$  is the representation in radix  $s$



# Example to illustrate the methods

- Conversion from decimal (radix 10) to binary (radix 2) and vice versa
- Use subscript D for decimal and B for binary
- $213_D \Rightarrow ???????_B$
- $11010101_B \Rightarrow ???_D$

# Decimal to Binary, division method

- $213_D \bmod 2_D = 1_D$
- $213_D \div 2_D = 106_D$
- $106 \bmod 2 = 0$
- $106 \div 2 = 53$
- $53 \bmod 2 = 1$
- $53 \div 2 = 26$
- $26 \bmod 2 = 0$
- $26 \div 2 = 13$
- $13 \bmod 2 = 1$
- $13 \div 2 = 6$
- $6 \bmod 2 = 0$
- $6 \div 2 = 3$
- $3 \bmod 2 = 1$
- $3 \div 2 = 1$
- $1 \bmod 2 = 1$
- $1 \div 2 = 0$

Result = 1 1 0 1 0 1 0 1<sub>B</sub>

# Decimal to Binary, multiplication method

$2\ 1\ 3_D \Rightarrow$  digit values are  $10_B, 1_B, 11_B$   
radix value is  $1010_B$

Do this in binary:  $(2_D \times 10_D + 1_D) \times 10_D + 3_D$

$(10_B \times 1010_B + 1_B) \times 1010_B + 11_B$

$= (10100 + 1) \times 1010 + 11$

$= 10101 \times 1010 + 11$

$= 10101000 + 101010 + 11$

$= 11010101$

# Binary to Decimal, division method

- $11010101_B \bmod 1010_B = 11_B$
- $11010101_B \div 1010_B = 10101_B$
- $10101 \bmod 1010 = 1$
- $10101 \div 1010 = 10$
- $10 \bmod 1010 = 10$
- $10 \div 1010 = 0$

$$\begin{aligned}\text{Result} &= 10_B \ 1_B \ 11_B = 2_D \ 1_D \ 3_D \\ &= 213_D\end{aligned}$$

# Binary to Decimal, multiplication method

$11010101_B \Rightarrow$  digit values are  $1_D, 1_D, 0_D, 1_D, 0_D, 1_D, 0_D, 1_D$

radix value is  $2_D$

Compute (in decimal):

$$\begin{aligned} & ((((((1 \times 2 + 1) \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \\ &= ((((((3 \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \\ &= ((((((6 \times 2 + 1) \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \\ &= ((((((13 \times 2 + 0) \times 2 + 1) \times 2 + 0) \times 2 + 1) \\ &= ((((((26 \times 2 + 1) \times 2 + 0) \times 2 + 1) \end{aligned}$$

# Binary to Decimal, multiplication method - continued

$$= ((26 \times 2 + 1) \times 2 + 0) \times 2 + 1$$

$$= (53 \times 2 + 0) \times 2 + 1$$

$$= 106 \times 2 + 1$$

$$= 213_D$$

# Binary numbers : 4 bits to 64 bits

4 bits : 0 .. 15

nibble

8 bits : 0 .. 255

byte

16 bits : 0 .. 65,535

half word

32 bits : 0 .. 4,294,967,295

word

64 bits : 0 .. 18,446,744,073,709,551,615

double word



# 4 bit binary numbers

Binary	Decimal	Binary	Decimal
0 0 0 0	0 0	1 0 0 0	0 8
0 0 0 1	0 1	1 0 0 1	0 9
0 0 1 0	0 2	1 0 1 0	1 0
0 0 1 1	0 3	1 0 1 1	1 1
0 1 0 0	0 4	1 1 0 0	1 2
0 1 0 1	0 5	1 1 0 1	1 3
0 1 1 0	0 6	1 1 1 0	1 4
0 1 1 1	0 7	1 1 1 1	1 5

# Radix 8 (octal) numbers

Binary	Decimal/Octal	Binary	Decimal/Octal
0 0 0 0	0 0 / 0 0	1 0 0 0	0 8 / 1 0
0 0 0 1	0 1 / 0 1	1 0 0 1	0 9 / 1 1
0 0 1 0	0 2 / 0 2	1 0 1 0	1 0 / 1 2
0 0 1 1	0 3 / 0 3	1 0 1 1	1 1 / 1 3
0 1 0 0	0 4 / 0 4	1 1 0 0	1 2 / 1 4
0 1 0 1	0 5 / 0 5	1 1 0 1	1 3 / 1 5
0 1 1 0	0 6 / 0 6	1 1 1 0	1 4 / 1 6
0 1 1 1	0 7 / 0 7	1 1 1 1	1 5 / 1 7

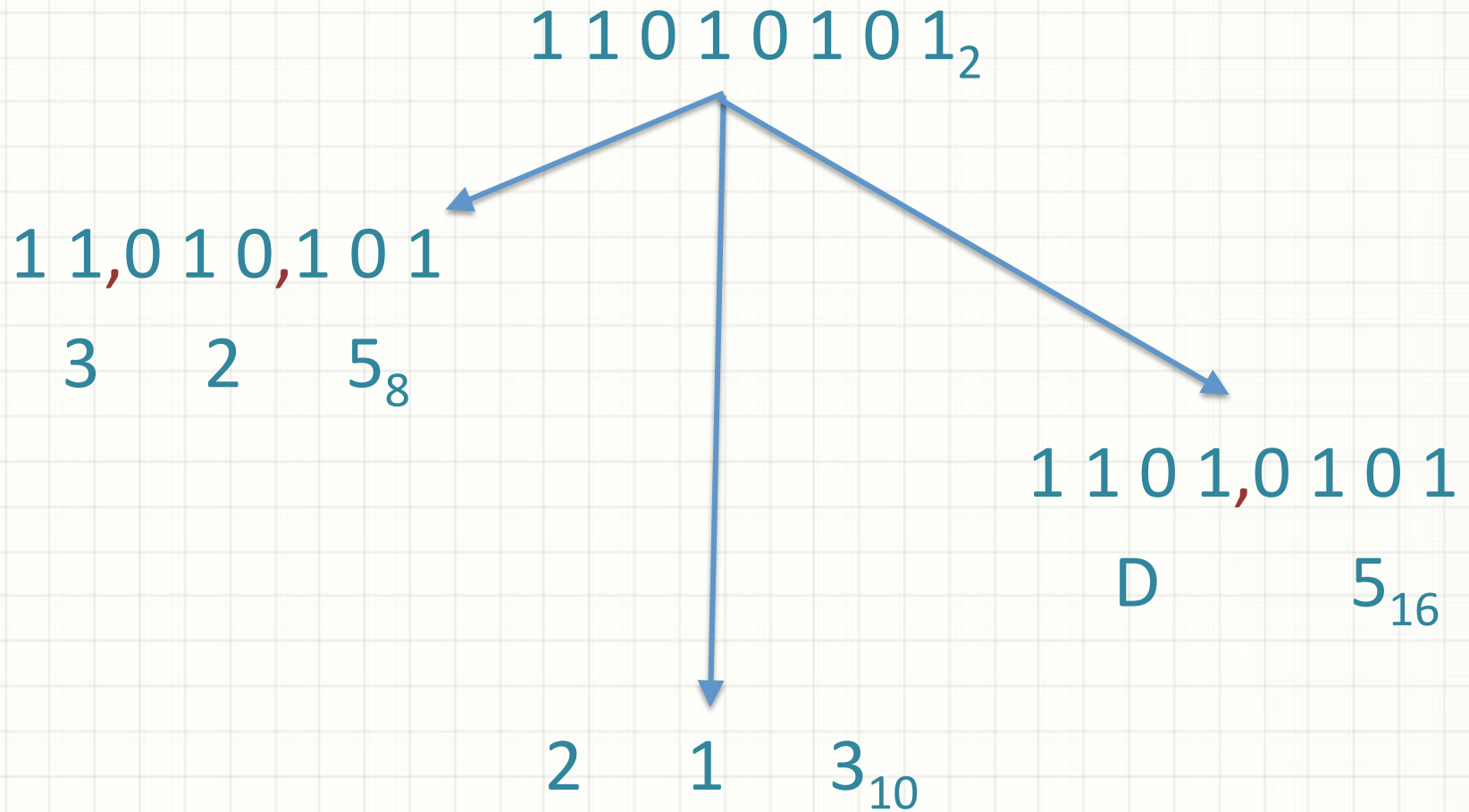
# Radix 16 (hexadecimal) numbers

Binary	Decimal/Hex	Binary	Decimal/Hex
0 0 0 0	0 0 / 0	1 0 0 0	0 8 / 8
0 0 0 1	0 1 / 1	1 0 0 1	0 9 / 9
0 0 1 0	0 2 / 2	1 0 1 0	1 0 / A
0 0 1 1	0 3 / 3	1 0 1 1	1 1 / B
0 1 0 0	0 4 / 4	1 1 0 0	1 2 / C
0 1 0 1	0 5 / 5	1 1 0 1	1 3 / D
0 1 1 0	0 6 / 6	1 1 1 0	1 4 / E
0 1 1 1	0 7 / 7	1 1 1 1	1 5 / F

# BCD (Binary Coded Decimal) numbers

Binary	Decimal/BCD	Binary	Decimal/BCD
0 0 0 0	0 0 / 0000 0000	1 0 0 0	0 8 / 0000 1000
0 0 0 1	0 1 / 0000 0001	1 0 0 1	0 9 / 0000 1001
0 0 1 0	0 2 / 0000 0010	1 0 1 0	1 0 / 0001 0000
0 0 1 1	0 3 / 0000 0011	1 0 1 1	1 1 / 0001 0001
0 1 0 0	0 4 / 0000 0100	1 1 0 0	1 2 / 0001 0010
0 1 0 1	0 5 / 0000 0101	1 1 0 1	1 3 / 0001 0011
0 1 1 0	0 6 / 0000 0110	1 1 1 0	1 4 / 0001 0100
0 1 1 1	0 7 / 0000 0111	1 1 1 1	1 5 / 0001 0101

# Radix 2, 8, 10, 16



# Binary Addition

Just like in primary school

0 0 1 1 [3]	0 1 0 1 [5]	0 1 1 0 [6]
+ 0 1 1 0 [6]	+ 0 1 1 1 [7]	+ 1 1 0 1 [13]
<hr/>		
1 0 0 1 [9]	1 1 0 0 [12]	0 0 1 1 [3]

Overflow will be discussed later.

# Binary Subtraction

Just like in primary school

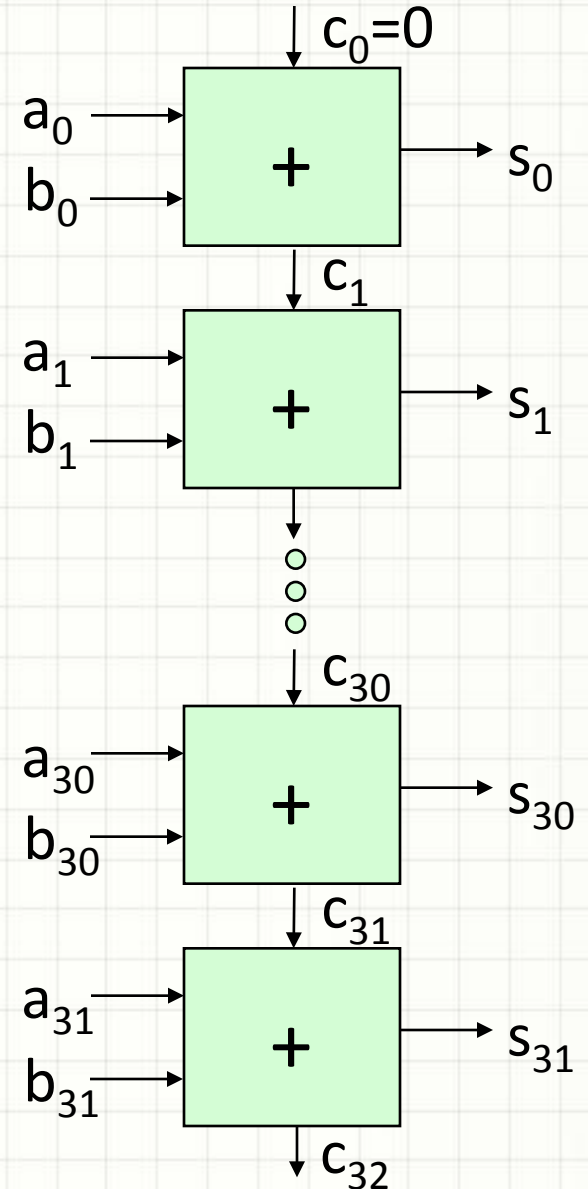
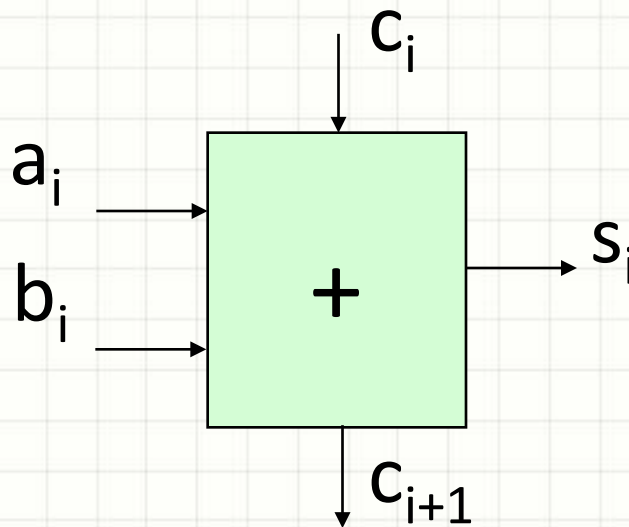
1 1 0 1 [13]	0 1 1 1 [7]	0 0 1 1 [3]
- 0 1 1 0 [6]	- 0 1 0 1 [5]	- 0 1 1 0 [6]
<hr/>		
0 1 1 1 [7]	0 0 1 0 [2]	1 1 0 1 [13]

Overflow will be discussed later.



# Adder circuit

- Perform bit by bit addition - similar to the “paper and pencil” way



# One bit adder truth table

$a_i$	$b_i$	$c_i$	$c_{i+1}$	$s_i$
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	1

# Boolean expressions for adder


$$s_i = a_i' b_i' c_i + a_i' b_i c_i' + a_i b_i' c_i' + a_i b_i c_i$$

alternatively,

$$s_i = a_i \oplus b_i \oplus c_i$$

$$c_{i+1} = a_i b_i + a_i c_i + b_i c_i$$

# One bit subtractor truth table

$a_i$	$b_i$	 $c_i$	$c_{i+1}$	$s_i$
0	0	0	0	0
0	0	1	1	1
0	1	0	1	1
0	1	1	1	0
1	0	0	0	1
1	0	1	0	0
1	1	0	0	0
1	1	1	1	1

# Boolean expressions for subtractor

$$s_i = a_i' b_i' c_i + a_i' b_i c_i' + a_i b_i' c_i' + a_i b_i c_i$$

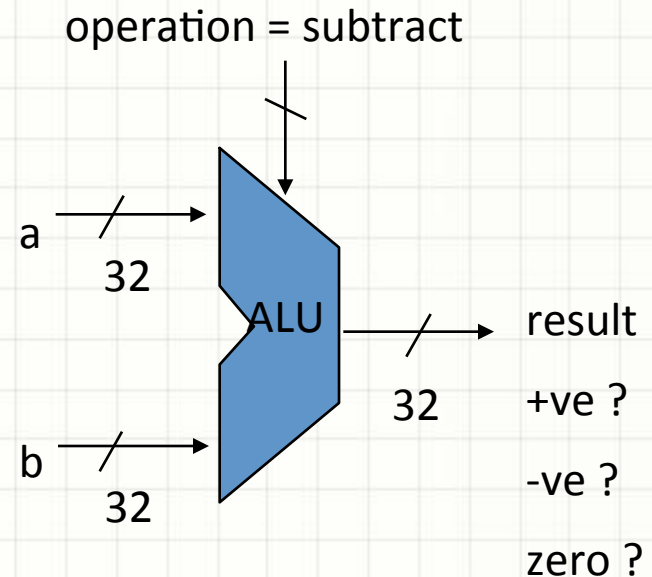
alternatively,

$$s_i = a_i \oplus b_i \oplus c_i$$

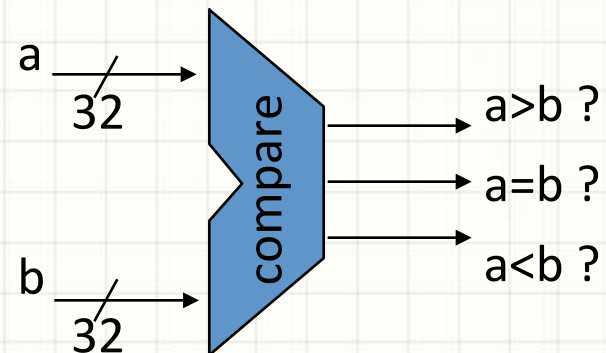
$$c_{i+1} = a_i' b_i + a_i' c_i + b_i c_i$$

# Comparing two integers

- Subtract and check the result



- Compare directly



# Compare (>) directly (unsigned)

Method 1:

$$a_{0..i} > b_{0..i} \equiv (a_{0..i-1} > b_{0..i-1})(a_i + \bar{b}_i) + a_i \bar{b}_i$$

Method 2:

$$a_{i..31} > b_{i..31} \equiv (a_{i+1..31} > b_{i+1..31}) + (a_{i+1..31} = b_{i+1..31})a_i \bar{b}_i$$



# Compare (=) directly (unsigned)

Method 1:

$$a_{0..i} = b_{0..i} \equiv (a_{0..i-1} = b_{0..i-1})(a_i \oplus \overline{b_i})$$

Method 2:

$$a_{i..31} = b_{i..31} \equiv (a_{i+1..31} = b_{i+1..31})(a_i \oplus \overline{b_i})$$

Method 3:

$$a = b \equiv \prod_{i=0}^{31} (a_i \oplus \overline{b_i})$$

# Comparator circuit – method 1

$$a_{0..i} > b_{0..i} \equiv (a_{0..i-1} > b_{0..i-1})(a_i + \bar{b}_i) + a_i \bar{b}_i$$

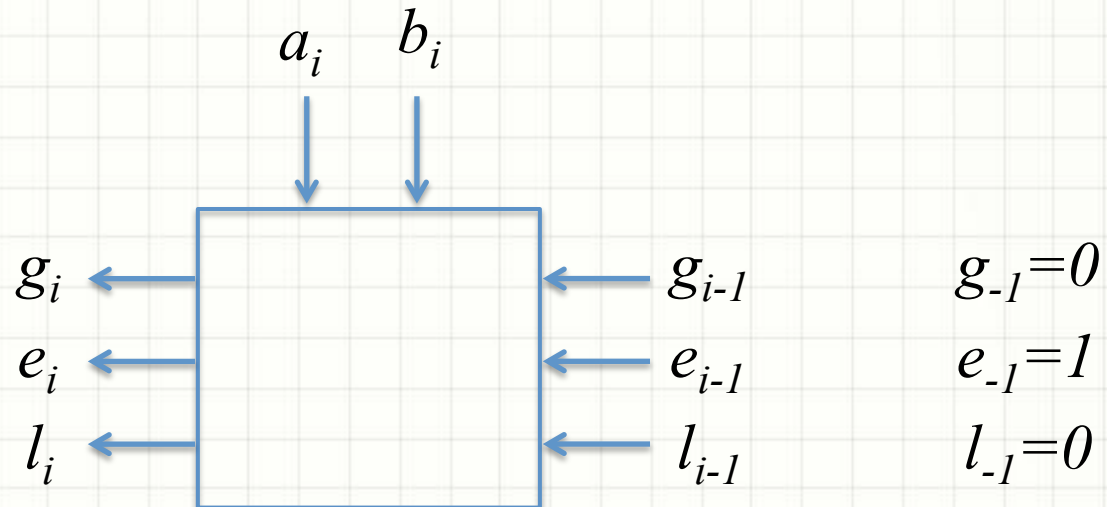
$$a_{0..i} = b_{0..i} \equiv (a_{0..i-1} = b_{0..i-1})(a_i \oplus \bar{b}_i)$$

$$a_{0..i} < b_{0..i} \equiv (a_{0..i-1} < b_{0..i-1})(\bar{a}_i + b_i) + \bar{a}_i b_i$$

$$g_i \equiv g_{i-1}(a_i + \bar{b}_i) + a_i \bar{b}_i$$

$$e_i \equiv e_{i-1}(a_i \oplus \bar{b}_i)$$

$$l_i \equiv l_{i-1}(\bar{a}_i + b_i) + \bar{a}_i b_i$$



# Comparator circuit – method 2

$$a_{i..31} > b_{i..31} \equiv (a_{i+1..31} > b_{i+1..31}) + (a_{i+1..31} = b_{i+1..31}) a_i \bar{b}_i$$

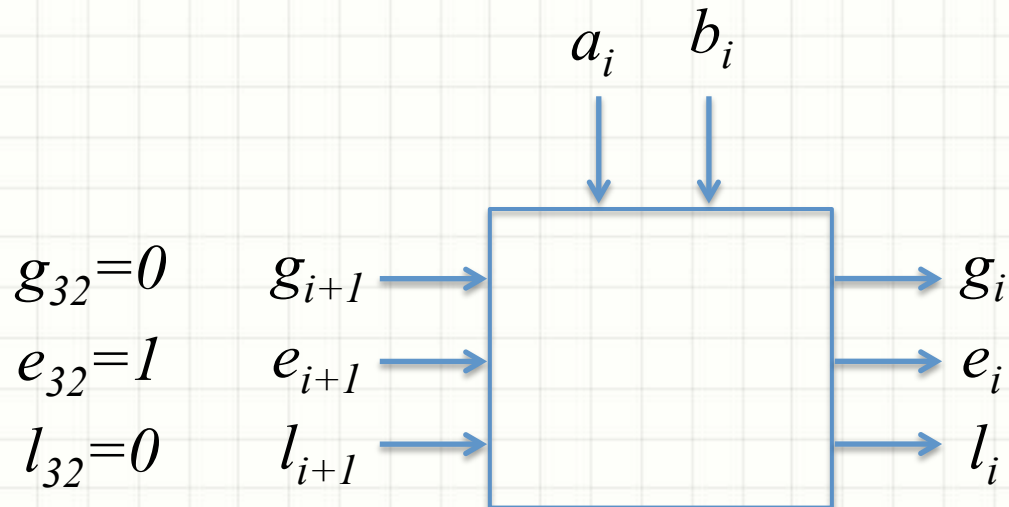
$$a_{i..31} = b_{i..31} \equiv (a_{i+1..31} = b_{i+1..31}) (a_i \oplus b_i)$$

$$a_{i..31} < b_{i..31} \equiv (a_{i+1..31} < b_{i+1..31}) + (a_{i+1..31} = b_{i+1..31}) \bar{a}_i b_i$$

$$g_i \equiv g_{i+1} + e_{i+1} a_i \bar{b}_i$$

$$e_i \equiv e_{i+1} (a_i \oplus b_i)$$

$$l_i \equiv l_{i+1} + e_{i+1} \bar{a}_i b_i$$





**THANKS**