



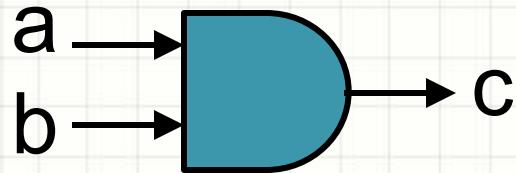
COL215 DIGITAL LOGIC AND SYSTEM DESIGN

Combinational Circuits
continued
28 July 2017

Lecture Outline

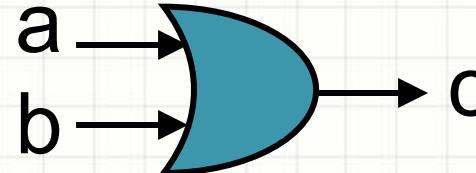
- Logic gates, truth tables, boolean algebra
- Gates
 - All types of gates
 - Minimum set of gates
- SoP, PoS and canonical forms
 - From truth tables
 - From arbitrary circuits or expressions
 - Notations

Building Blocks: Gates



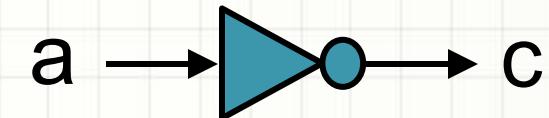
$c = a \text{ AND } b$

$c = a \cdot b$



$c = a \text{ OR } b$

$c = a + b$



$c = \text{NOT } a$

$c = \bar{a}$ or $c = a'$

a	b	c
0	0	0
0	1	0
1	0	0
1	1	1

a	b	c
0	0	0
0	1	1
1	0	1
1	1	1

a	c
0	1
1	0

Boolean Algebra

Set B, operations . & +, identity 0 & 1, inverse '

$$a + a = a$$

$$a \cdot a = a$$

$$a + b = b + a$$

$$a \cdot b = b \cdot a$$

$$a + (b+c) = (a+b) + c$$

$$a \cdot (b.c) = (a.b) \cdot c$$

$$a \cdot (b+c) = (a.b) + (a.c)$$

$$a + (b.c) = (a+b) \cdot (a+c)$$

$$a \cdot 1 = a$$

$$a + 0 = a$$

$$a \cdot a' = 0$$

$$a + a' = 1$$

Derived results

$$(a')' = a$$

$$a \cdot 0 = 0$$

$$(a + b)' = a' \cdot b'$$

$$x + (x \cdot y) = x$$

$$x + (x' \cdot y) = x + y$$

$$a + 1 = 1$$

$$(a \cdot b)' = a' + b'$$

$$x \cdot (x + y) = x$$

$$x \cdot (x' + y) = x \cdot y$$

$$(x \cdot y) + (y \cdot z) + (x' \cdot z) = (x \cdot y) + (x' \cdot z)$$

$$(x+y) \cdot (y+z) \cdot (x'+z) = (x+y) \cdot (x'+z)$$

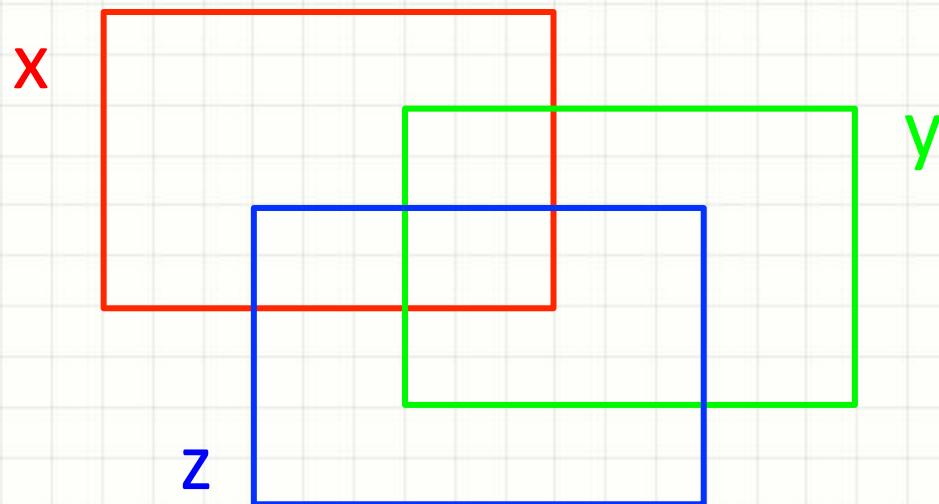
Comparing two circuits/expressions

Is $(x.y) + (y.z) + (x'.z) = (x.y) + (x'.z)$?

1. Transform LHS into RHS (or vice versa) using boolean algebra
2. Compare truth tables of LHS and RHS
3. Use Venn diagrams

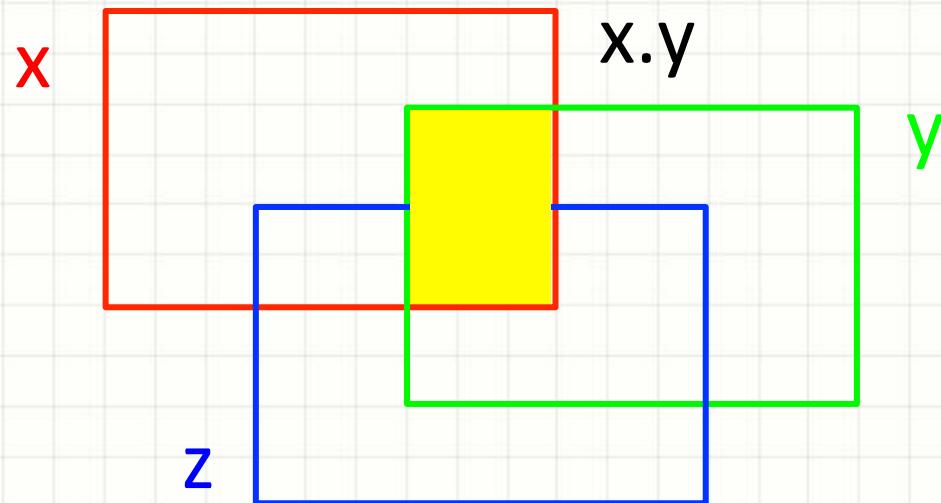
Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$



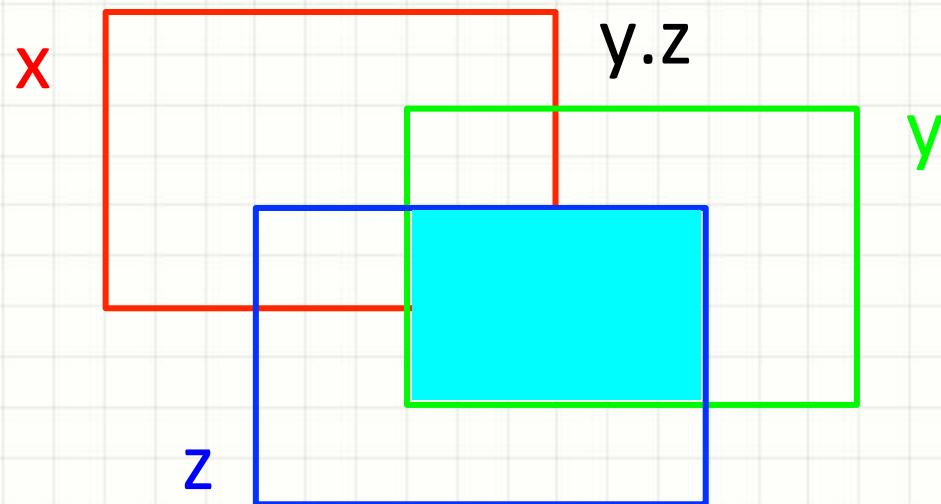
Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$



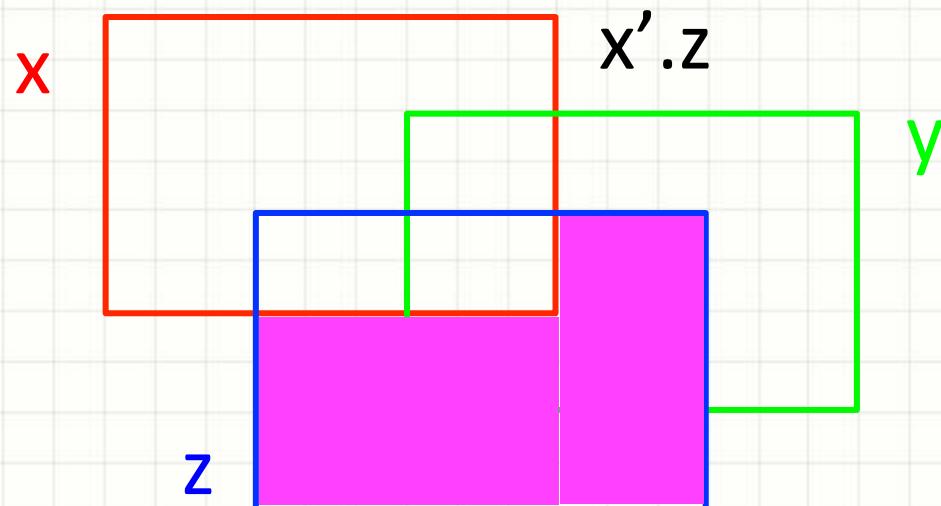
Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$



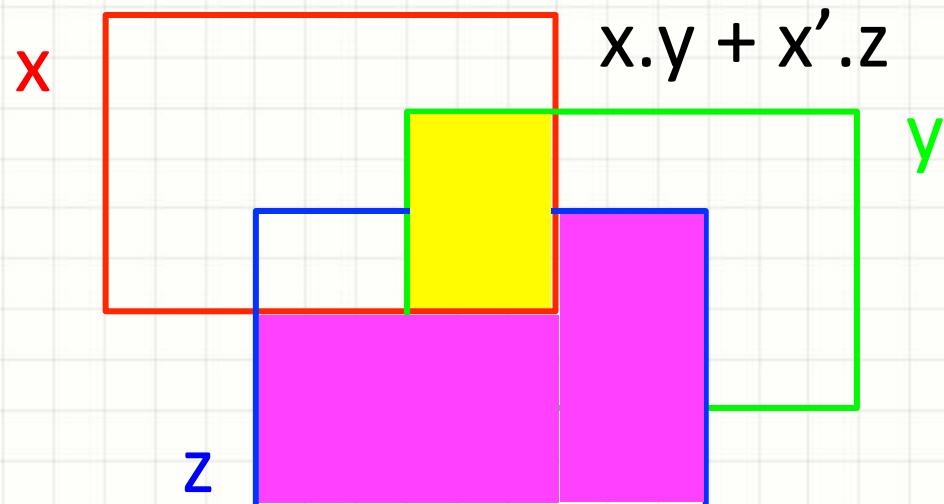
Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$



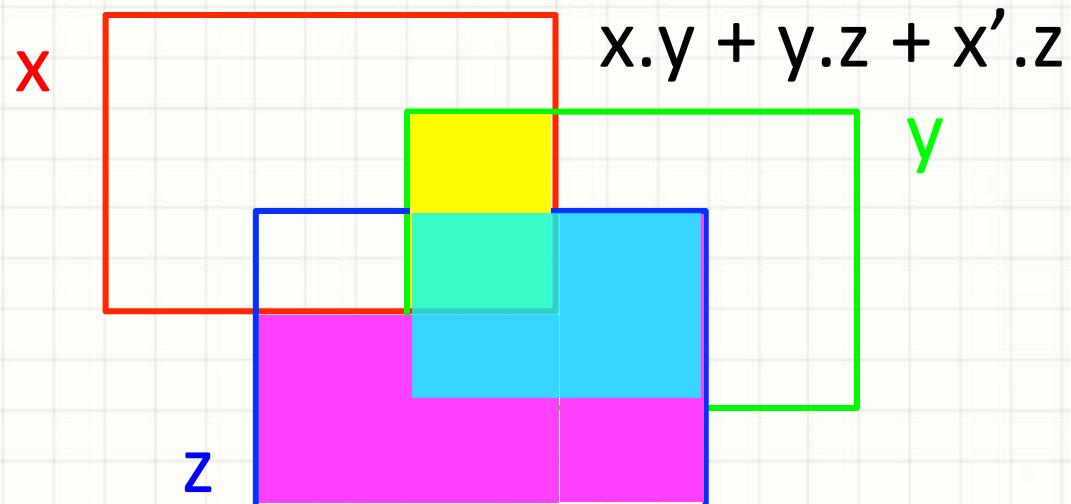
Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$



Venn diagram example

Compare $(x.y) + (y.z) + (x'.z)$ with $(x.y) + (x'.z)$

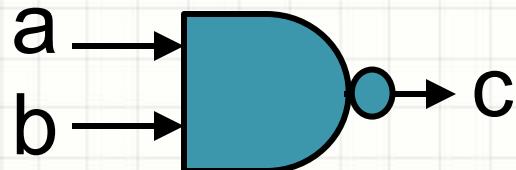


Boolean algebra with $|B| > 2$

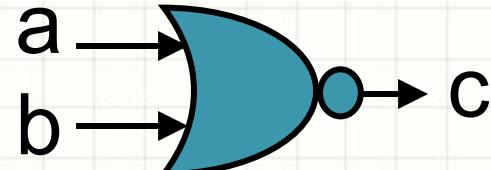
Example:

- Let S be a finite set
- $B = \text{powerset } (S)$
- $0 = \{ \}$ $1 = S$
- $\text{AND} = \text{set intersection}$ $\text{OR} = \text{set union}$
- $\text{NOT} = \text{set complement}$

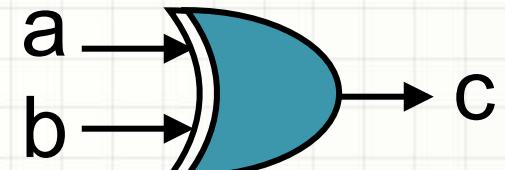
Are there other types of gates?



$c = a \text{ NAND } b$



$c = a \text{ NOR } b$



$c = a \text{ XOR } b$

a	b	c
0	0	1
0	1	1
1	0	1
1	1	0

a	b	c
0	0	1
0	1	0
1	0	0
1	1	0

a	b	c
0	0	0
0	1	1
1	0	1
1	1	0

How many types of gates exist?

Any 2 input gate is defined by a truth table of the form:

Number of possible values of t_3, t_2, t_1, t_0 is
 $2^4 = 16$

a	b	c
0	0	t_0
0	1	t_1
1	0	t_2
1	1	t_3

0	a nor b	a' and b	a'
a and b'	b'	a xor b	a nand b
a and b	a xnor b	b	a' or b
a	a or b'	a or b	1

How many types of gates exist?

Any 2 input gate is defined by a truth table of the form:

Number of possible values of t_3, t_2, t_1, t_0 is
 $2^4 = 16$

a	b	c
0	0	t_0
0	1	t_1
1	0	t_2
1	1	t_3

0	a nor b	$(a \leq b)'$	a'
$(a \Rightarrow b)'$	b'	a xor b	a nand b
a and b	$a \equiv b$	b	$a \Rightarrow b$
a	$a \leq b$	a or b	1

Generalize to n inputs

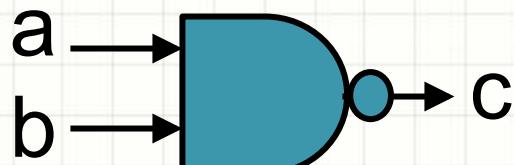
n input truth table :

Number of rows = 2^n

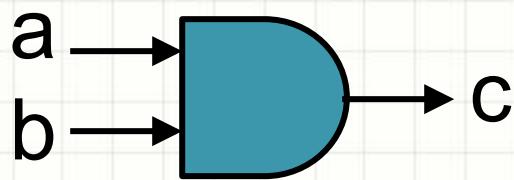
Number of possible values of output column is 2^{2^n}

Do we need all types of gates?

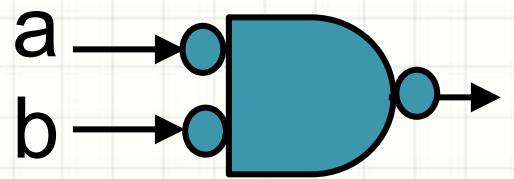
Any function can be expressed as SoP or PoS =>
 $\{\text{AND}, \text{OR}, \text{NOT}\}$ is a functionally complete set.



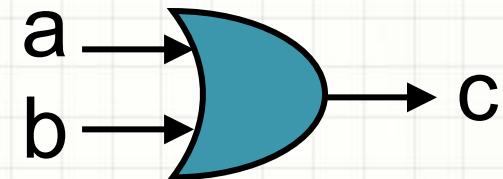
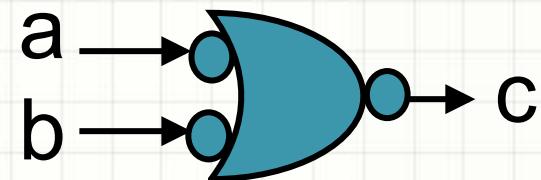
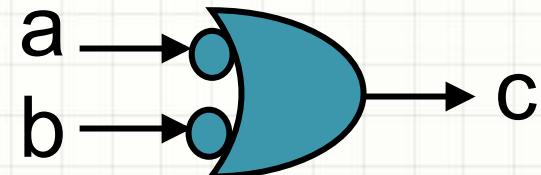
equivalent to



equivalent to

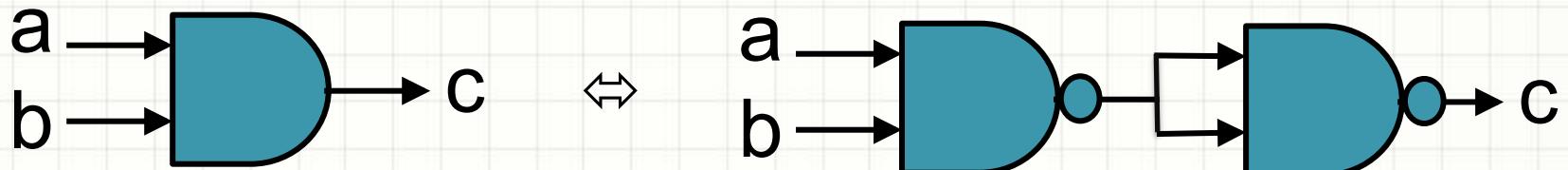


equivalent to

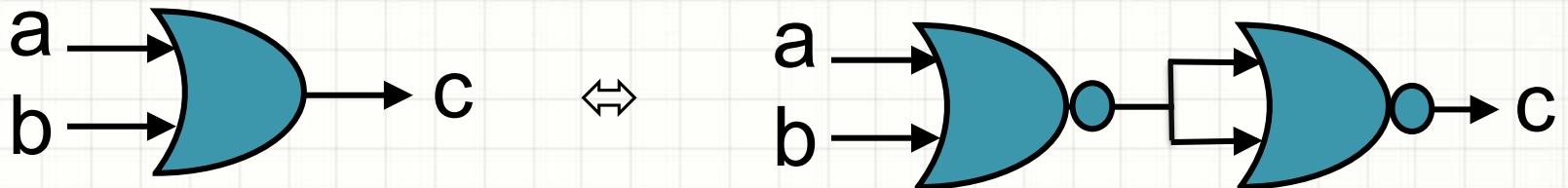


With NOT, only one of {AND, OR} is sufficient. =>
 $\{\text{AND}, \text{NOT}\}$, $\{\text{OR}, \text{NOT}\}$ are functionally complete

Further . . .

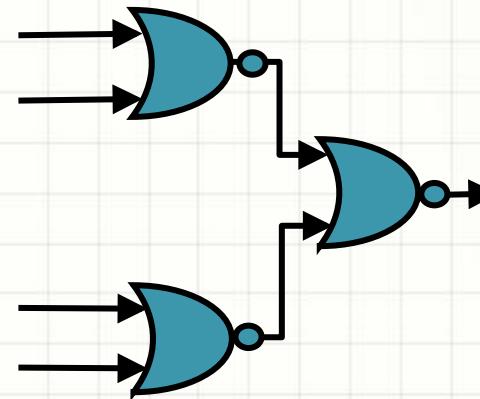
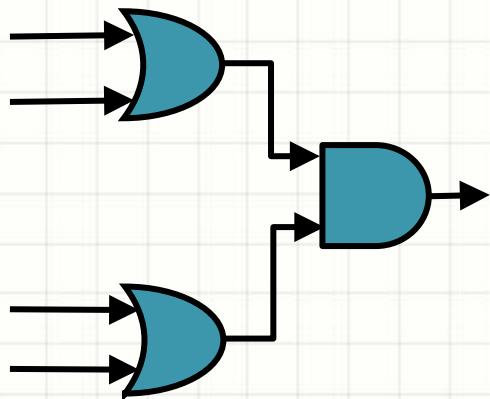
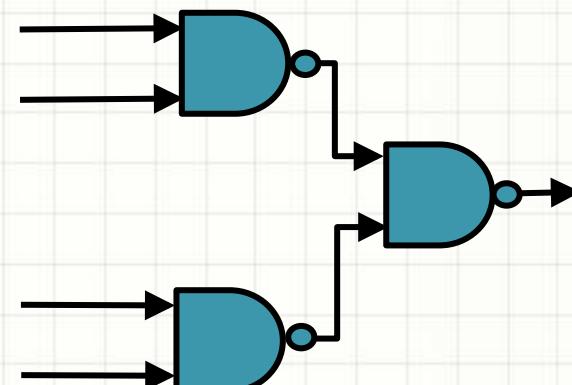
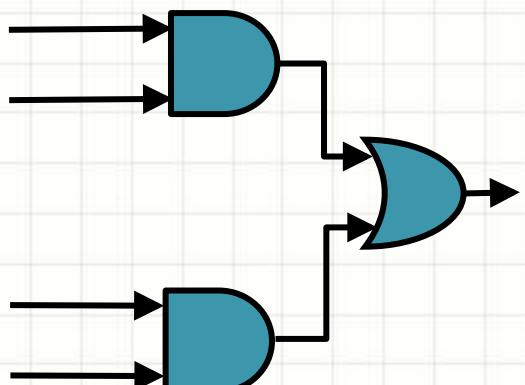


=> {NAND} is functionally complete set.

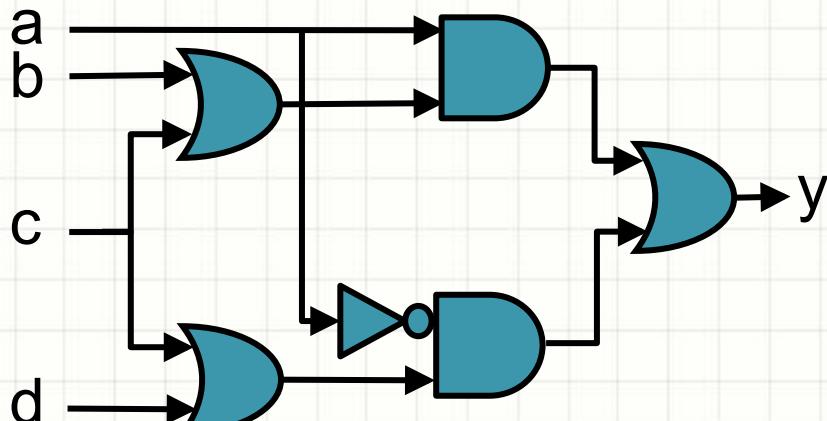


=> {NOR} is functionally complete set.

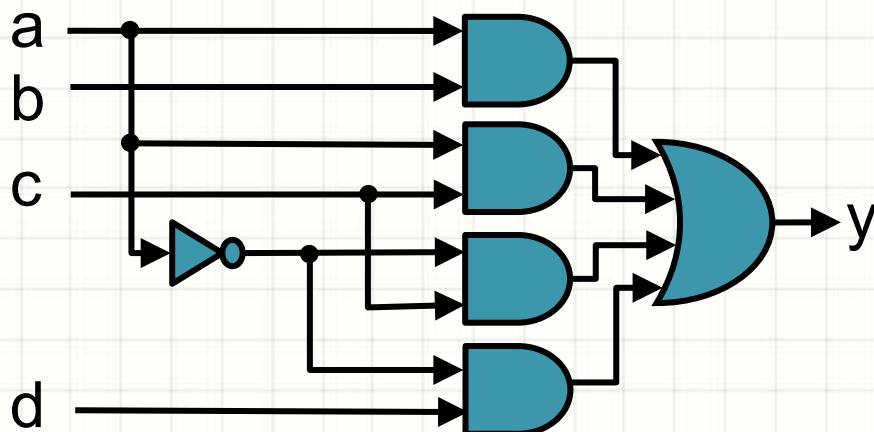
SOP to NAND-NAND, POS to NOR-NOR



Multi-level and SoP forms



$$y = (a \cdot (b + c)) + (a' \cdot (c + d))$$

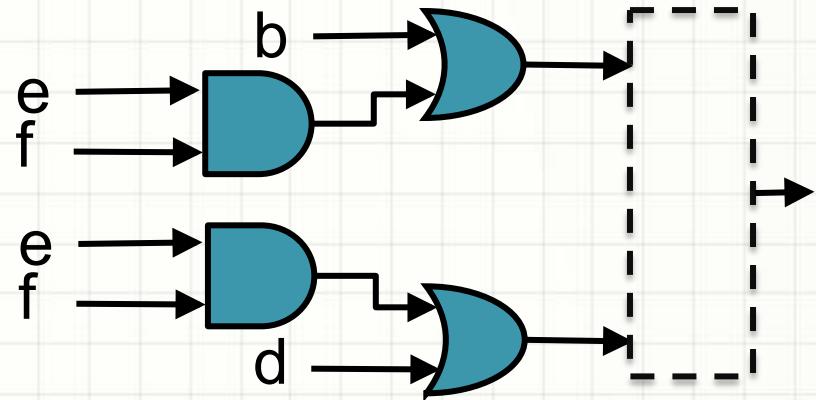
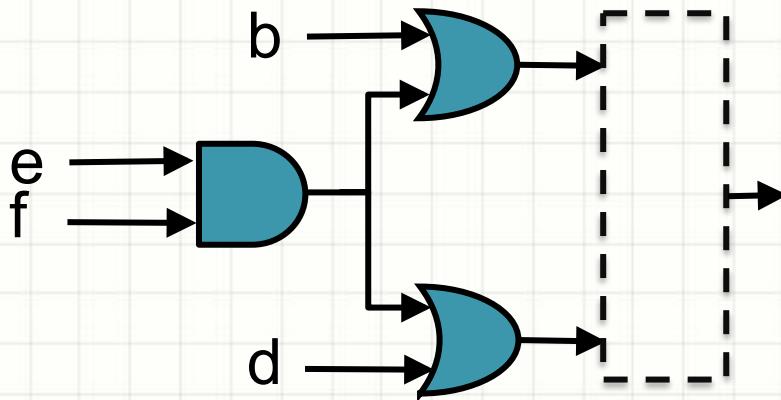
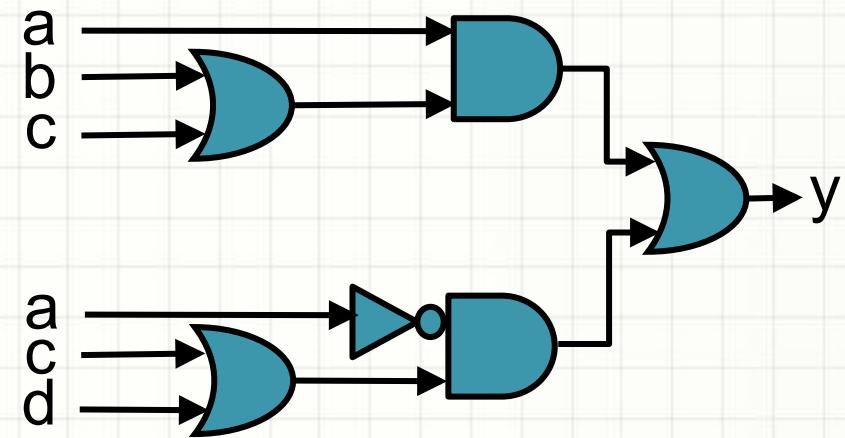
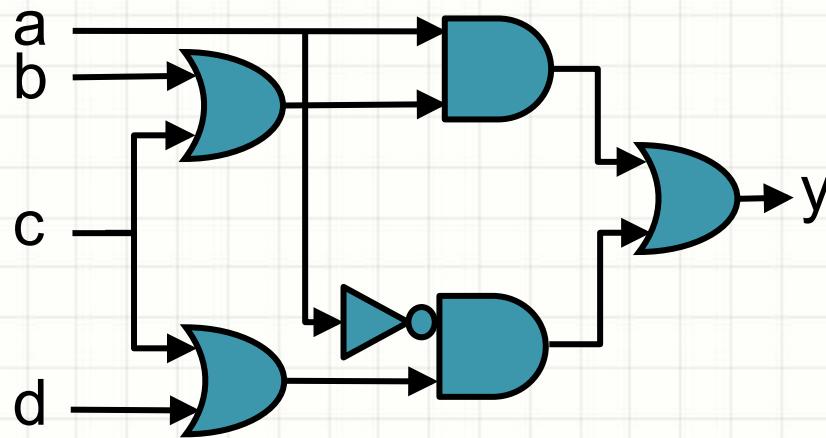


$$y = (a \cdot b) + (a \cdot c) + (a' \cdot c) + (a' \cdot d)$$

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

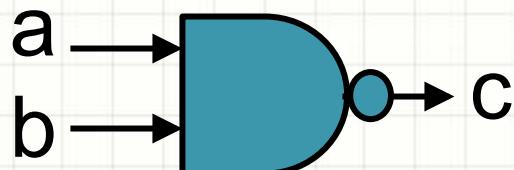
Arbitrary Circuit to SoP/PoS

Convert to tree

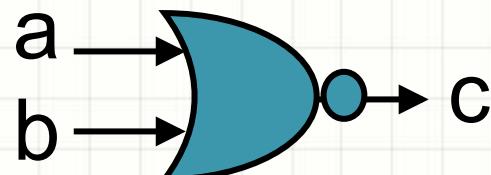
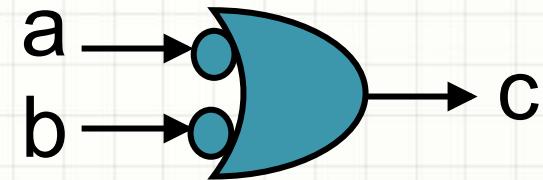


Arbitrary Circuit to SoP/PoS

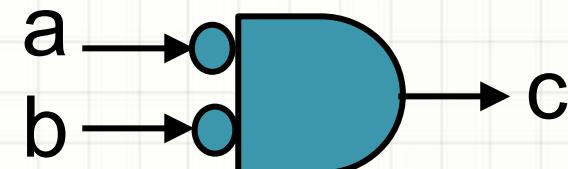
Move inverters to leaves



=>

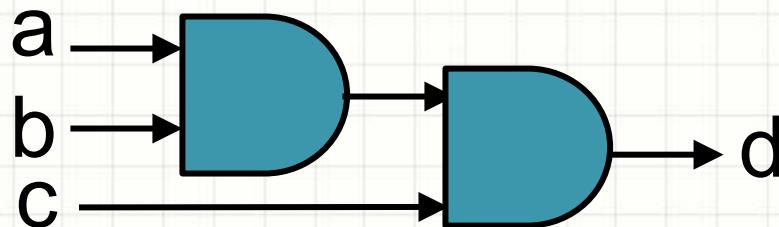


=>

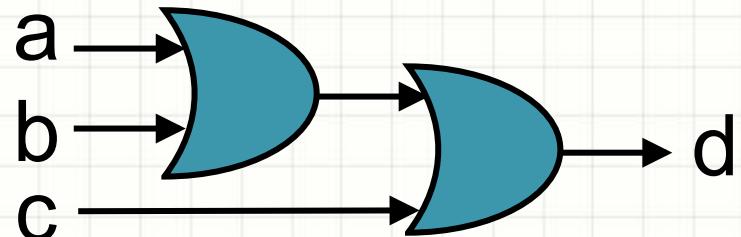
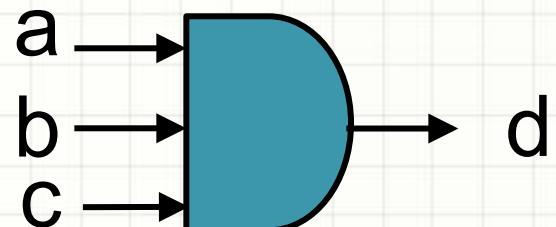


Arbitrary Circuit to SoP/PoS

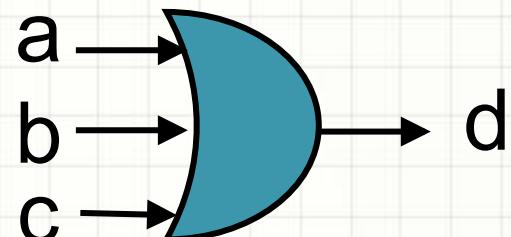
Merge gates



=>

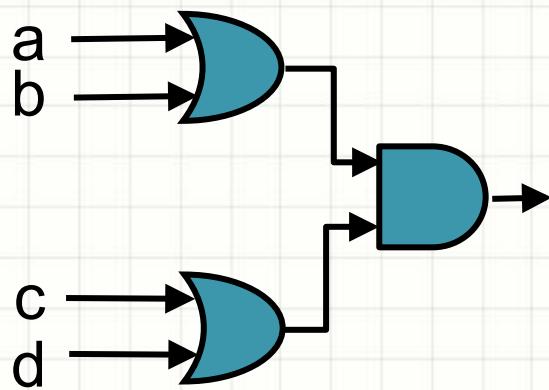


=>

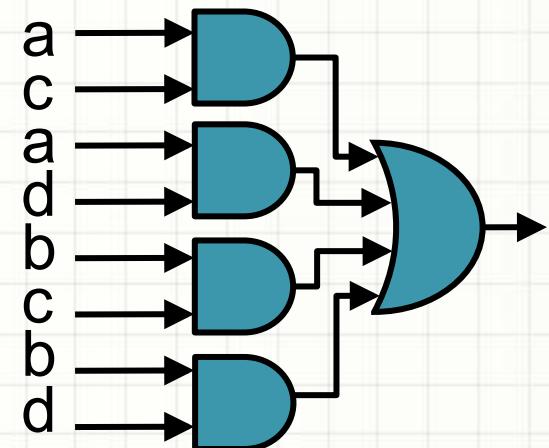


Arbitrary Circuit to SoP/PoS

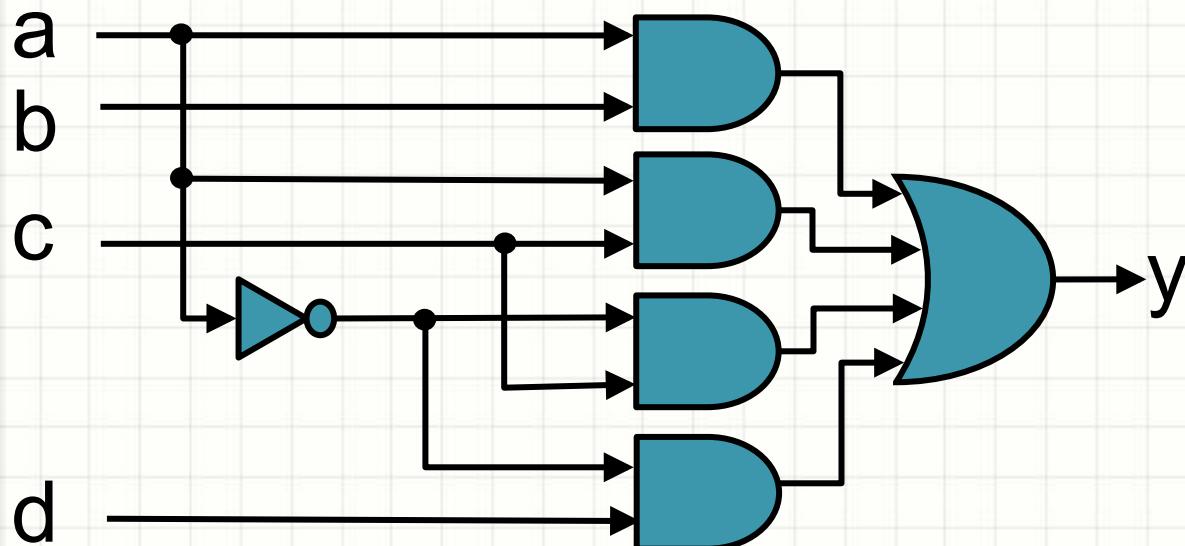
PoS to SoP and vice versa (locally)



=>



Canonical SoP form

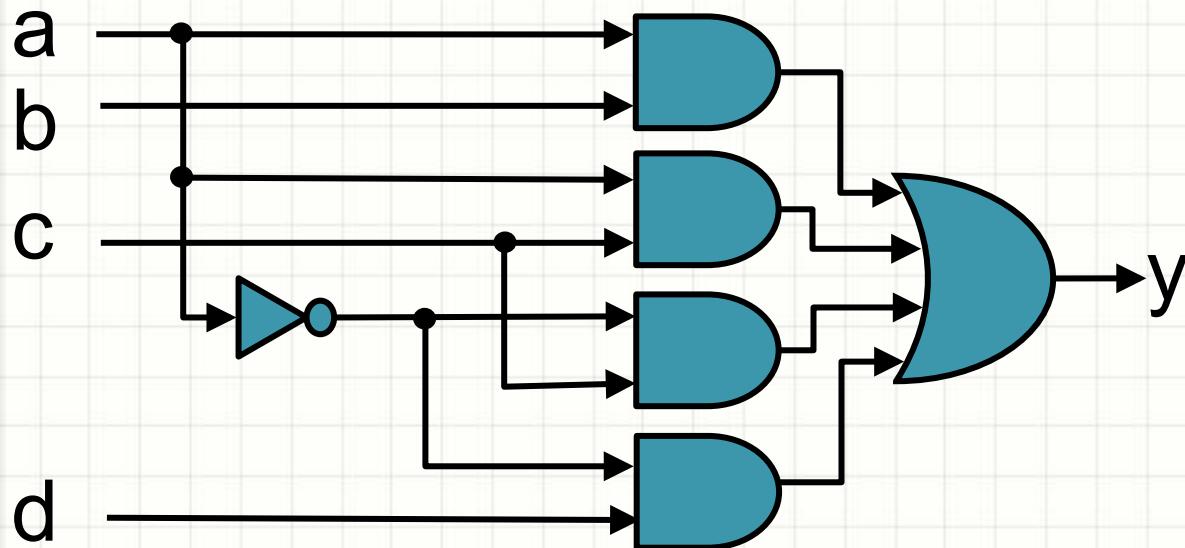


Canonical Sum of Product (SOP) form

$$y = (a'.b'.c'.d) + (a'.b'.c.d') + (a'.b'.c.d)$$

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Canonical SoP form

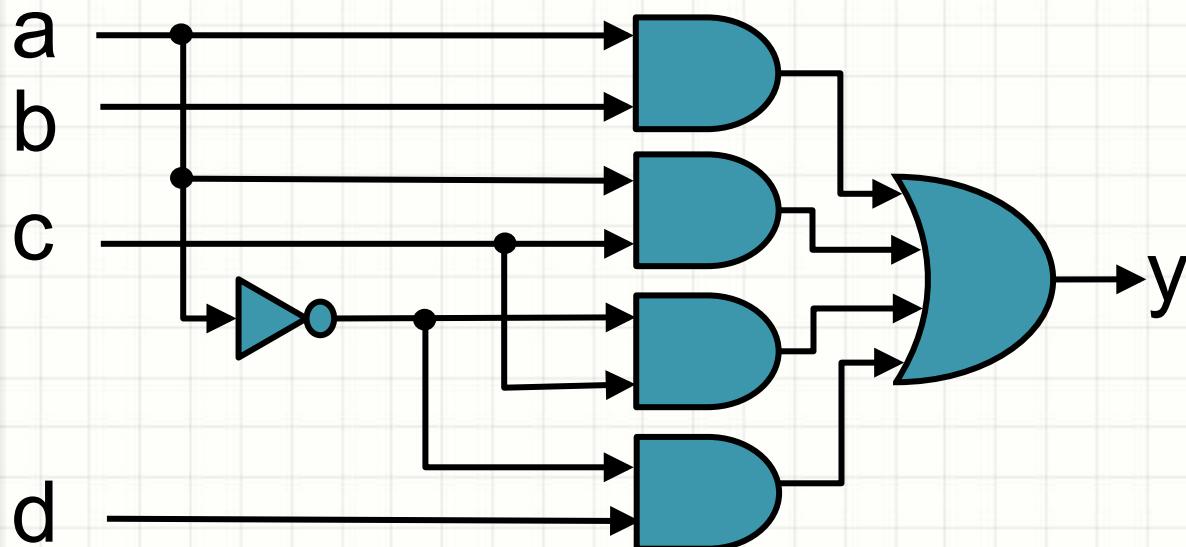


Canonical Sum of Product (SOP) form

$$\begin{aligned}y = & (a'.b'.c'.d) + (a'.b'.c.d') + (a'.b'.c.d) \\& + (a'.b.c'.d) + (a'.b.c.d') + (a'.b.c.d)\end{aligned}$$

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Canonical SoP form

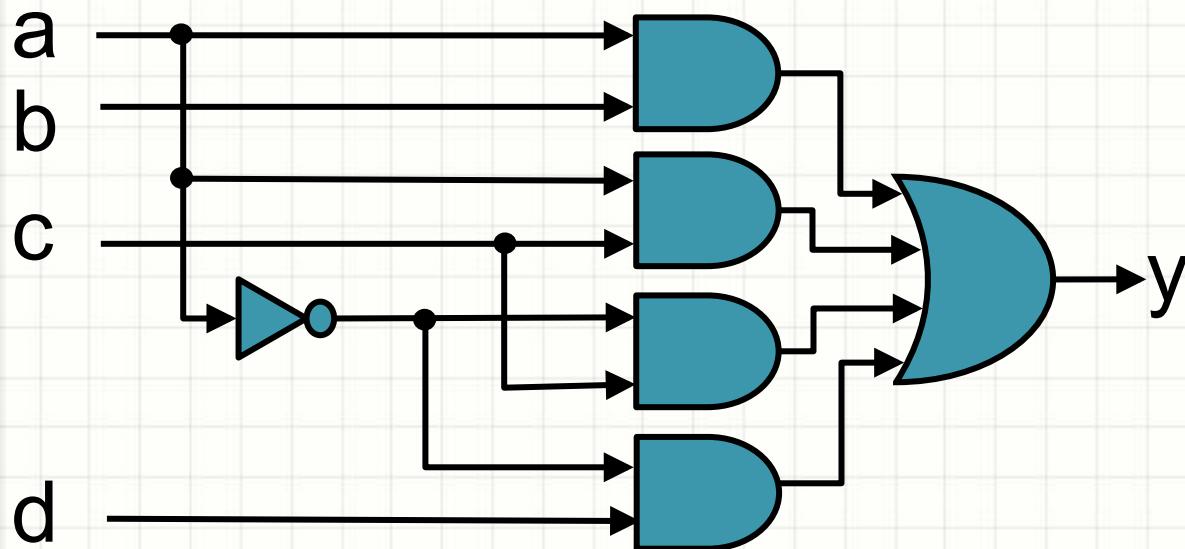


Canonical Sum of Product (SOP) form

$$\begin{aligned} y = & (a'.b'.c'.d) + (a'.b'.c.d') + (a'.b'.c.d) \\ & + (a'.b.c'.d) + (a'.b.c.d') + (a'.b.c.d) \\ & + (a.b'.c.d') + (a.b'.c.d) + (a.b.c'.d') \end{aligned}$$

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Canonical SoP form



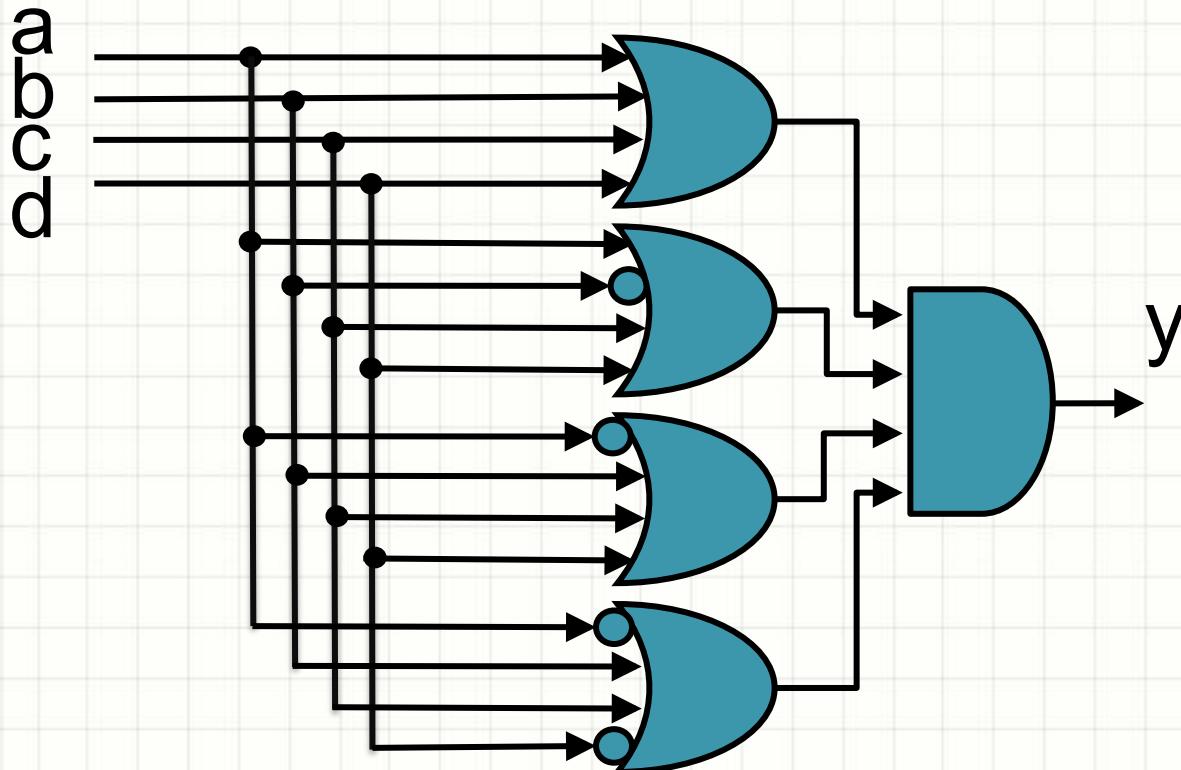
Canonical Sum of Product (SOP) form

$$\begin{aligned}y = & (a'.b'.c'.d) + (a'.b'.c.d') + (a'.b'.c.d) \\& + (a'.b.c'.d) + (a'.b.c.d') + (a'.b.c.d) \\& + (a.b'.c.d') + (a.b'.c.d) + (a.b.c'.d') \\& + (a.b.c.d') + (a.b.c.d') + (a.b.c.d)\end{aligned}$$

These are “min terms”

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Canonical POS form



$$y' = a'b'c'd' + a'b c'd' + a b'c'd' + a b'c'd$$

$$\begin{aligned} y = & (a + b + c + d) \cdot (a + b' + c + d) \\ & \cdot (a' + b + c + d) \cdot (a' + b + c + d') \end{aligned}$$

These are “max terms”.

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Numbering min/max terms

	a	b	c	d	y
0	0	0	0	0	0
1	0	0	0	1	1
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	0
5	0	1	0	1	1
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	1
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1

Notation for min/max terms

Min terms

$$m_0 = a'.b'.c'.d'$$

$$m_1 = a'.b'.c'.d$$

$$m_2 = a'.b'.c.d'$$

$$m_3 = a'.b'.c.d$$

and so on

Max terms

$$M_0 = a + b + c + d$$

$$M_1 = a + b + c + d'$$

$$M_2 = a + b + c' + d$$

$$M_3 = a + b + c' + d'$$

and so on

	a	b	c	d	y
0	0	0	0	0	0
1	0	0	0	1	1
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	0
5	0	1	0	1	1
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	1
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1

Notation for canonical SOP form

Canonical SOP form

$$y = (a'.b'.c'.d) + (a'.b'.c.d') + (a'.b'.c.d) \\ + (a'.b.c'.d) + (a'.b.c.d') + (a'.b.c.d) \\ + (a.b'.c.d') + (a.b'.c.d) + (a.b.c'.d') \\ + (a.b.c.d') + (a.b.c.d) + (a.b.c.d)$$

$$y = f(a, b, c, d) = m_1 + m_2 + m_3 + \\ m_5 + m_6 + m_7 + \\ m_{10} + m_{11} + m_{12} + \\ m_{13} + m_{14} + m_{15}$$

$$= \Sigma (m_1, m_2, m_3, m_5, m_6, m_7, \\ m_{10}, m_{11}, m_{12}, m_{13}, m_{14}, m_{15})$$

$$= \Sigma m (1, 2, 3, 5, 6, 7, 10, 11, 12, 13, 14, 15)$$

	a	b	c	d	y
0	0	0	0	0	0
1	0	0	0	1	1
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	0
5	0	1	0	1	1
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	1
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1

Notation for canonical POS form

Canonical POS form

$$y = (a + b + c + d) \cdot (a' + b' + c + d) \\ \cdot (a' + b + c + d) \cdot (a' + b + c + d')$$

$$y = f(a, b, c, d) = M_0 \cdot M_4 \cdot M_8 \cdot M_9$$

$$= \prod (M_0, M_4, M_8, M_9)$$

$$= \prod M(0, 4, 8, 9)$$

	a	b	c	d	y
0	0	0	0	0	0
1	0	0	0	1	1
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	0
5	0	1	0	1	1
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	1
13	1	1	0	1	1
14	1	1	1	0	1
15	1	1	1	1	1



QUESTIONS?