



COL215 DIGITAL LOGIC AND SYSTEM DESIGN

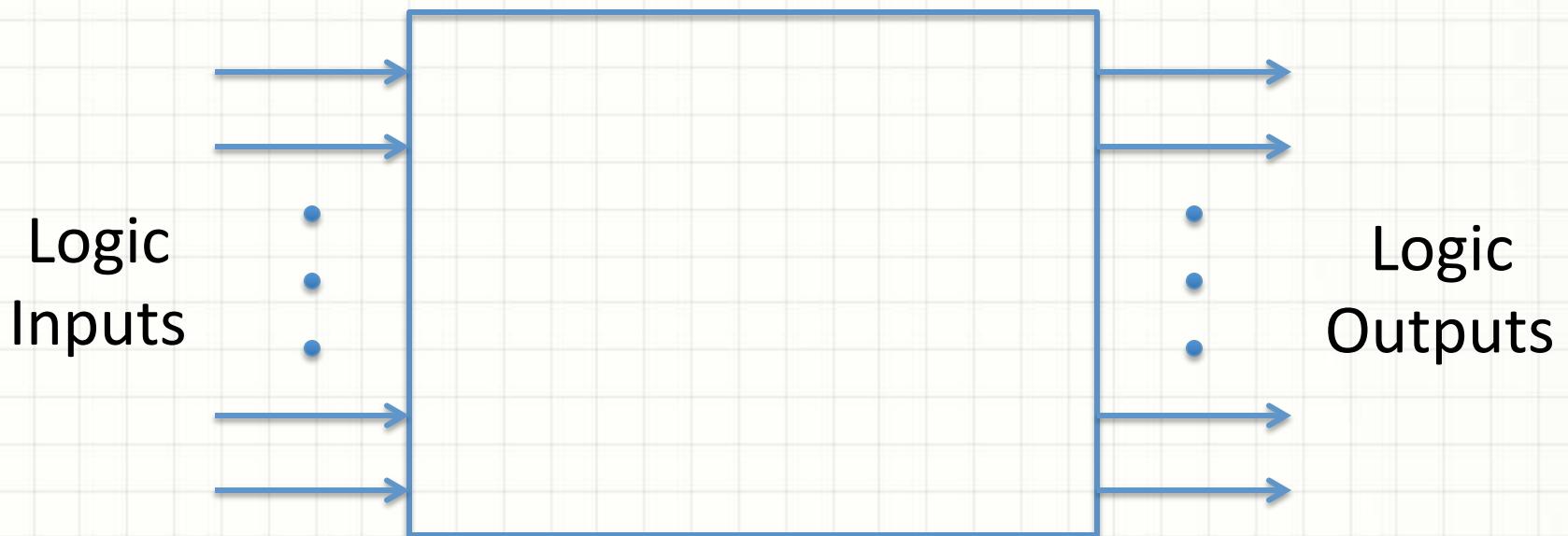
Combinational Logic Circuits

26 July 2017

Lecture Outline

- Logic circuits
- Combinational vs sequential circuits
- Building blocks – Logic gates
- Putting gates together
- Logic simplification
- Boolean Algebra
- Special forms of circuits

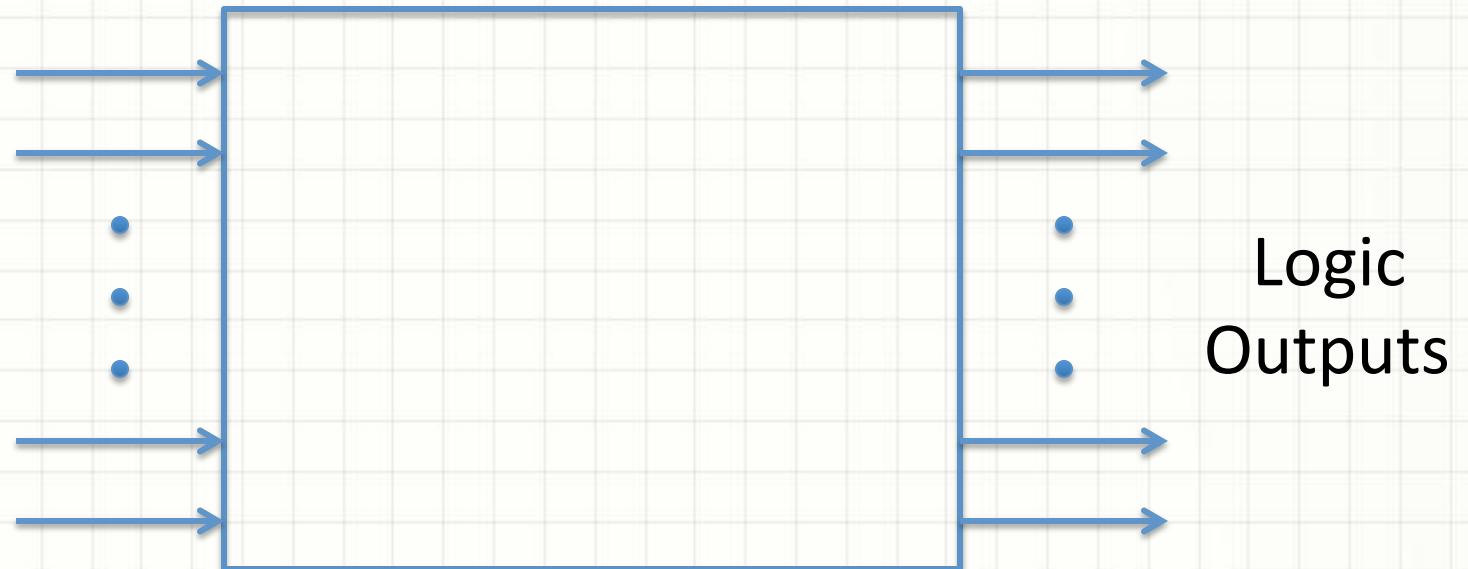
Logic Circuits



Combinational Logic Circuits

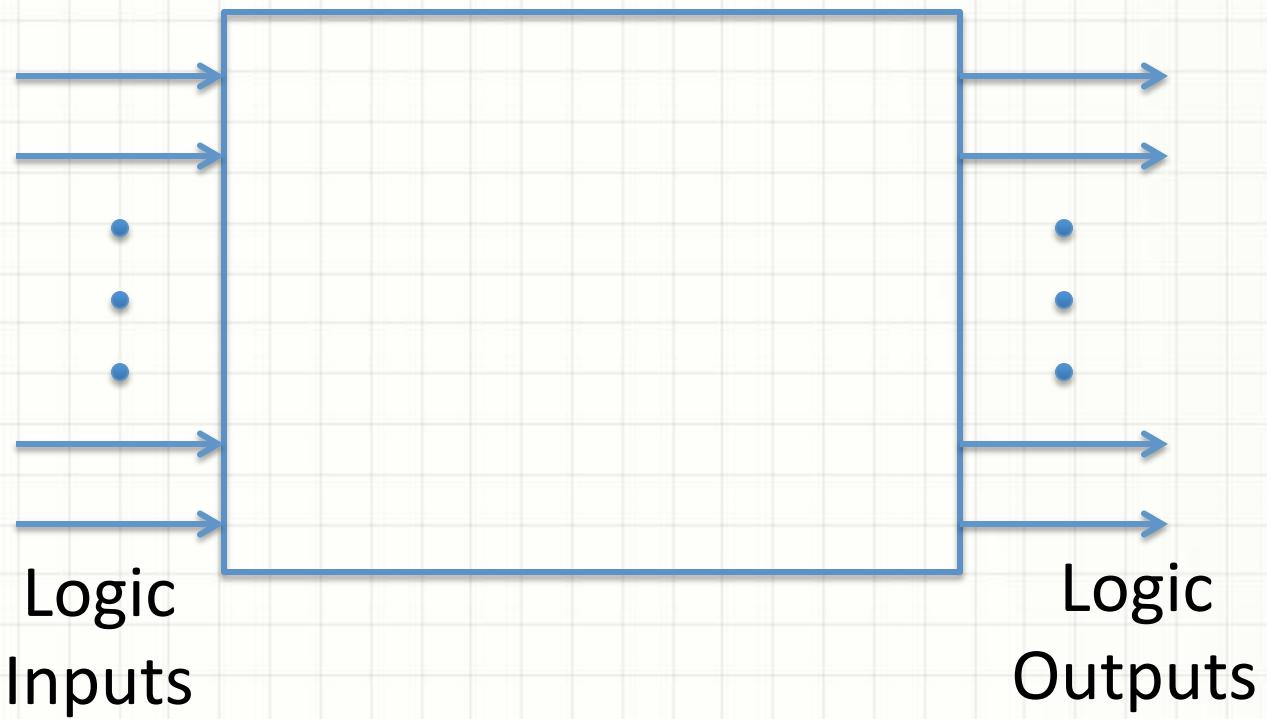
Output values depend on
the current combination of
input values

Logic
Inputs



Sequential Logic Circuits

Output values depend on
the current and previous
input values





Dependence on previous input values

- The circuit “remembers” the previous input values !
- Practical circuits have “finite” memory
- Previous input sequences can be represented by a finite number of states

Example: Car Interior Light Control

- Switch OFF/DOOR/ON
- Doors open/close (4)

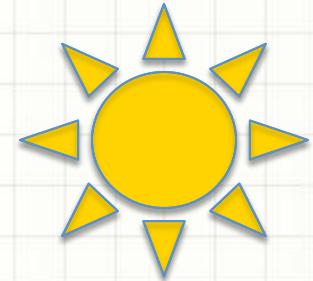


More sophisticated control

- Switch OFF/DOOR/ON
- Doors open/close (4)
- Door key
- Ignition key
- Gradual fading
- Time-out

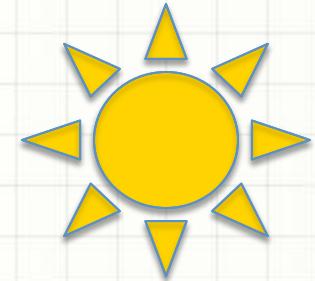
Typical sequence: Entry

- Unlock the door
- Open the door
- Enter the car
- Close the door
- Put key in ignition switch
- Switch on ignition



Typical sequence: Exit

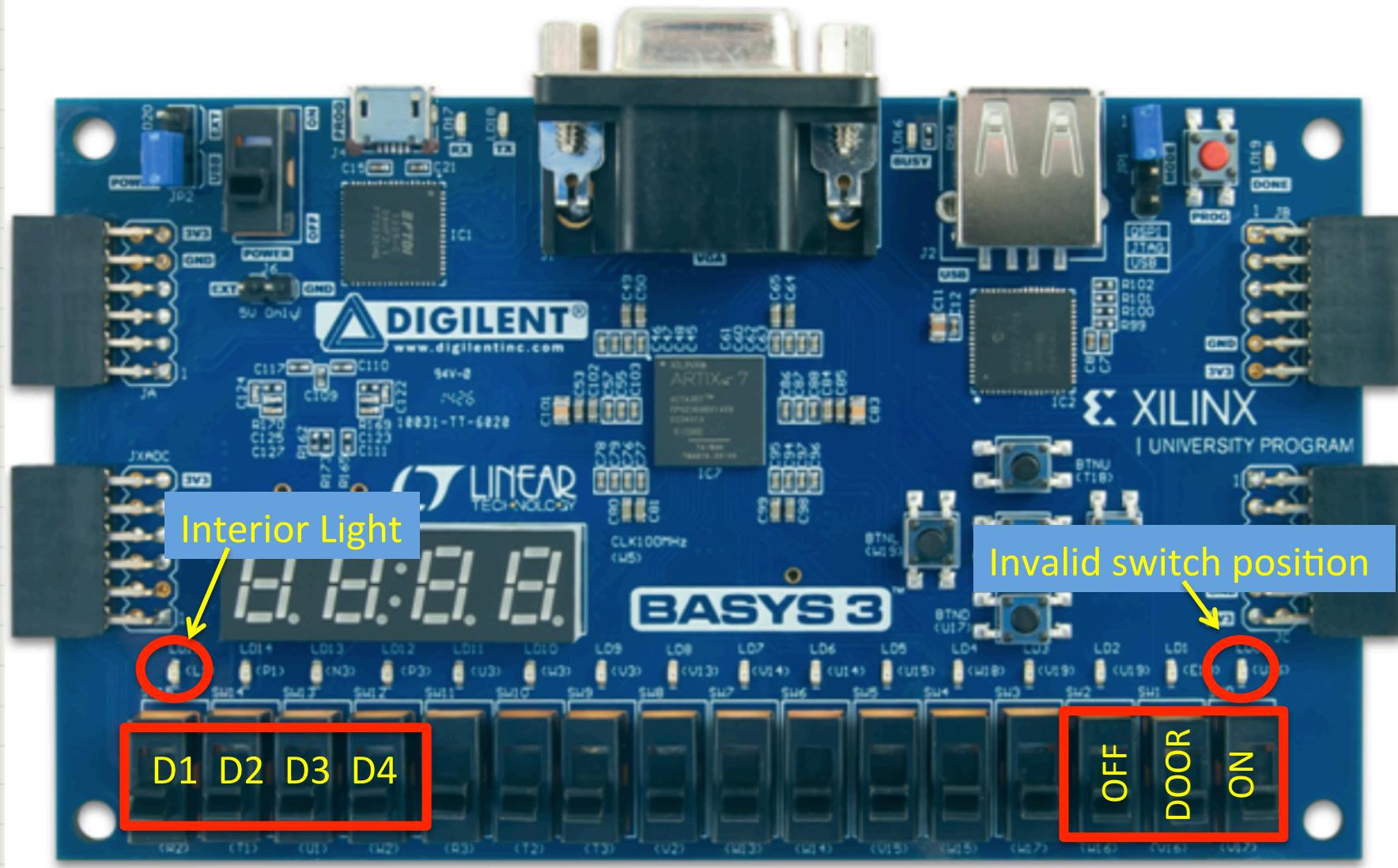
- Switch off ignition
- Take ignition key out
- Open the door
- Exit the car
- Close the door
- Lock the door



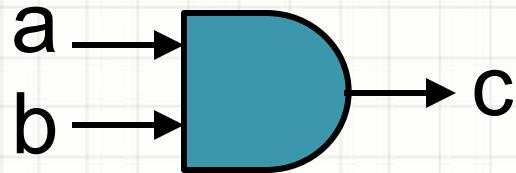
Additional Logic

- When the light is off, opening any door turns it on and closing the door makes it turn off gradually
- When the light is on, it times out in 8-10 sec if there is no event

Lab Exercise 1

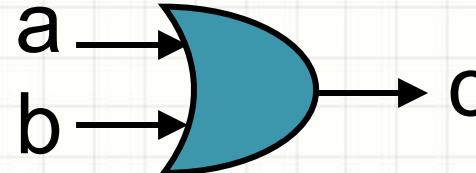


Building Blocks: Gates



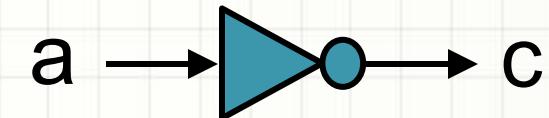
$c = a \text{ AND } b$

$c = a \cdot b$



$c = a \text{ OR } b$

$c = a + b$



$c = \text{NOT } a$

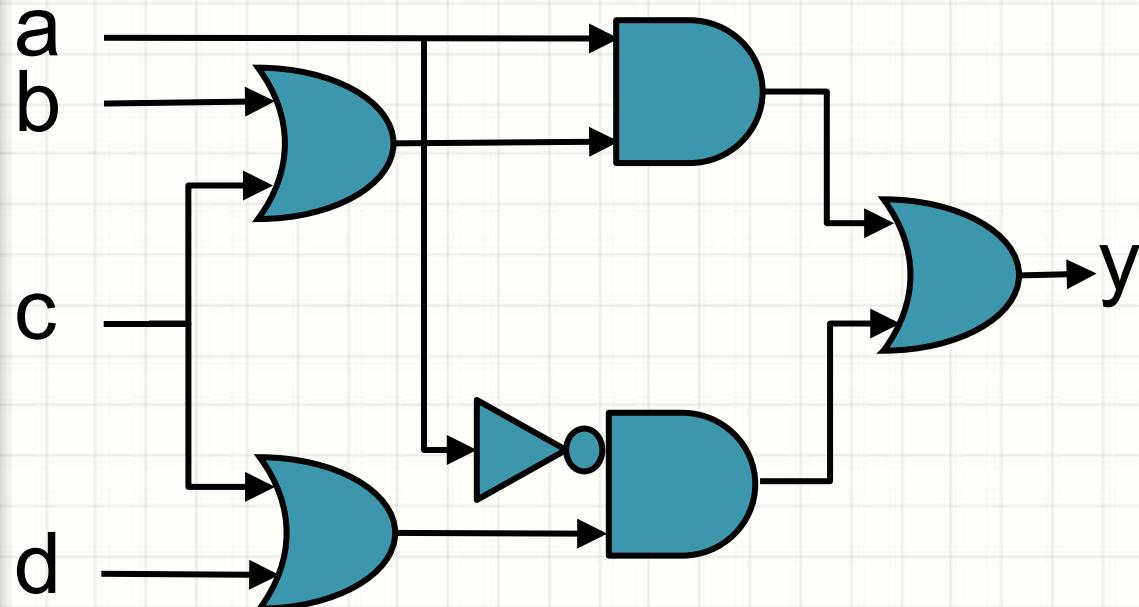
$c = \bar{a}$ or $c = a'$

a	b	c
0	0	0
0	1	0
1	0	0
1	1	1

a	b	c
0	0	0
0	1	1
1	0	1
1	1	1

a	c
0	1
1	0

Putting Gates together

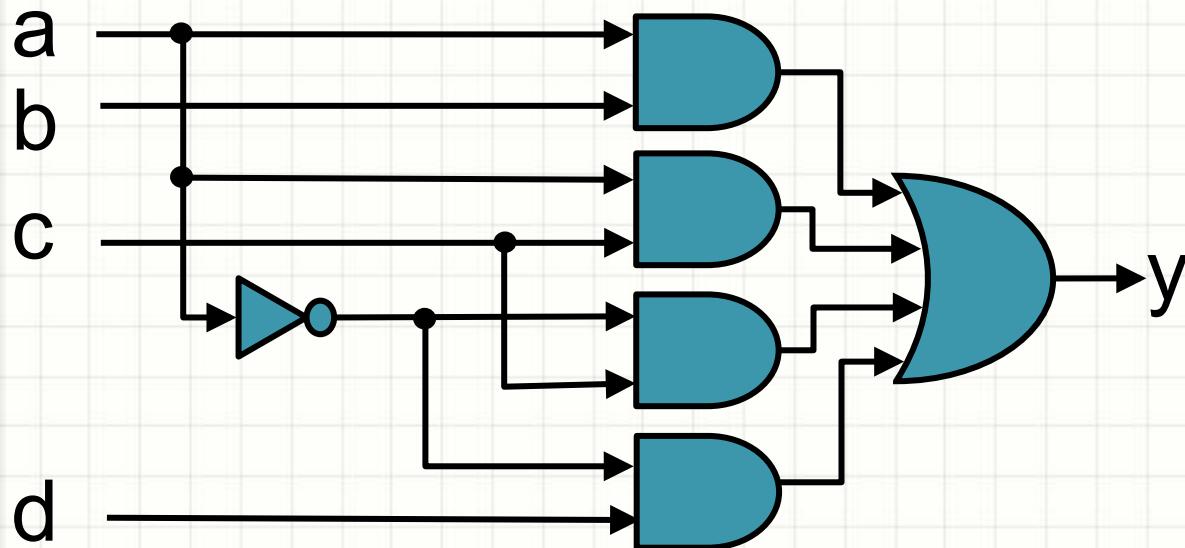


$$y = (a \text{ AND } (b \text{ OR } c)) \text{ OR } (\text{NOT } a \text{ AND } (c \text{ OR } d))$$

$$y = (a \cdot (b + c)) + (\bar{a} \cdot (c + d))$$

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

Another Circuit



$$y = (a \text{ AND } b) \text{ OR } (a \text{ AND } c) \text{ OR} \\ (\text{NOT } a \text{ AND } c) \text{ OR } (\text{NOT } a \text{ AND } d)$$

$$y = (a \cdot b) + (a \cdot c) + (\bar{a} \cdot c) + (\bar{a} \cdot d)$$

This is Sum of Products (SOP) form

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1

How to compare two circuits?

- Compare truth tables, row by row, exhaustively!
- Check if two expressions are “equivalent”
- Apply rules, for example,
$$a \cdot (b + c) \equiv (a \cdot b) + (a \cdot c)$$

Among equivalent expressions/circuits, can we identify the best one? Criteria: cost/delay

Transforming logic expressions: Boolean Algebra

Set B, operations . & +, identity 0 & 1, inverse '

$$a + a = a \cdot a = a$$

idempotent

$$a + b = b + a$$

commutative

$$a \cdot b = b \cdot a$$

commutative

$$a + (b + c) = (a + b) + c$$

associative

$$a \cdot (b \cdot c) = (a \cdot b) \cdot c$$

associative

$$a \cdot (b + c) = (a \cdot b) + (a \cdot c)$$

distributive

$$a + (b \cdot c) = (a + b) \cdot (a + c)$$

distributive

Boolean Algebra – contd.

$$a \cdot 1 = a \quad a + 0 = a \quad \text{def: identity elements}$$

$$a \cdot a' = 0 \quad a + a' = 1 \quad \text{def: inverse/complement}$$
$$(a')' = a \quad (\text{by symmetry})$$

$$a \cdot 0 = 0 \quad a + 1 = 1 \quad (\text{derived})$$

Derived results

$$(a + b)' = a' \cdot b'$$

DeMorgan's theorem

$$(a \cdot b)' = a' + b'$$

$$x + (x \cdot y) = x$$

Absorption

$$x \cdot (x + y) = x$$

$$x + (x' \cdot y) = x + y$$

$$x \cdot (x' + y) = x \cdot y$$

$$(x \cdot y) + (y \cdot z) + (x' \cdot z) = (x \cdot y) + (x' \cdot z)$$

$$(x + y) \cdot (y + z) \cdot (x' + z) = (x + y) \cdot (x' + z)$$

Duality Principle

Take any Boolean Algebra equality

Replace . with + and vice versa

Replace 1 with 0 and vice versa

The equality will still hold

Prove $[x + (x \cdot y) = x]$

$$[x + (x \cdot y) = x]$$

$$\text{LHS} = x + (x \cdot y)$$

$$= (x \cdot 1) + (x \cdot y)$$

$$= x \cdot (1 + y)$$

$$= x \cdot 1$$

$$= x = \text{RHS}$$

Prove $[x + (x' \cdot y) = x + y]$

$$[x + (x' \cdot y) = x + y]$$

$$\text{LHS} = x + (x' \cdot y)$$

$$= x + (x \cdot y) + (x' \cdot y)$$

$$= x + (x + x') \cdot y$$

$$= x + 1 \cdot y$$

$$= x + y = \text{RHS}$$

Prove $(x.y) + (y.z) + (x'.z) = (x.y) + (x'.z)$

$$[(x \cdot y) + (y \cdot z) + (x' \cdot z)]$$

$$\text{LHS} = (x \cdot y) + (y \cdot z) + (x' \cdot z)$$

$$= (x \cdot y) + (1 \cdot y \cdot z) + (x' \cdot z)$$

$$= (x \cdot y) + ((x + x') \cdot y \cdot z) + (x' \cdot z)$$

$$= (x \cdot y) + (x \cdot y \cdot z) + (x' \cdot y \cdot z) + (x' \cdot z)$$

$$= (x \cdot (y + (y \cdot z))) + (x' \cdot ((y \cdot z) + z))$$

$$= (x \cdot (y)) + (x' \cdot (z))$$

$$= (x \cdot y) + (x' \cdot z) = \text{RHS}$$

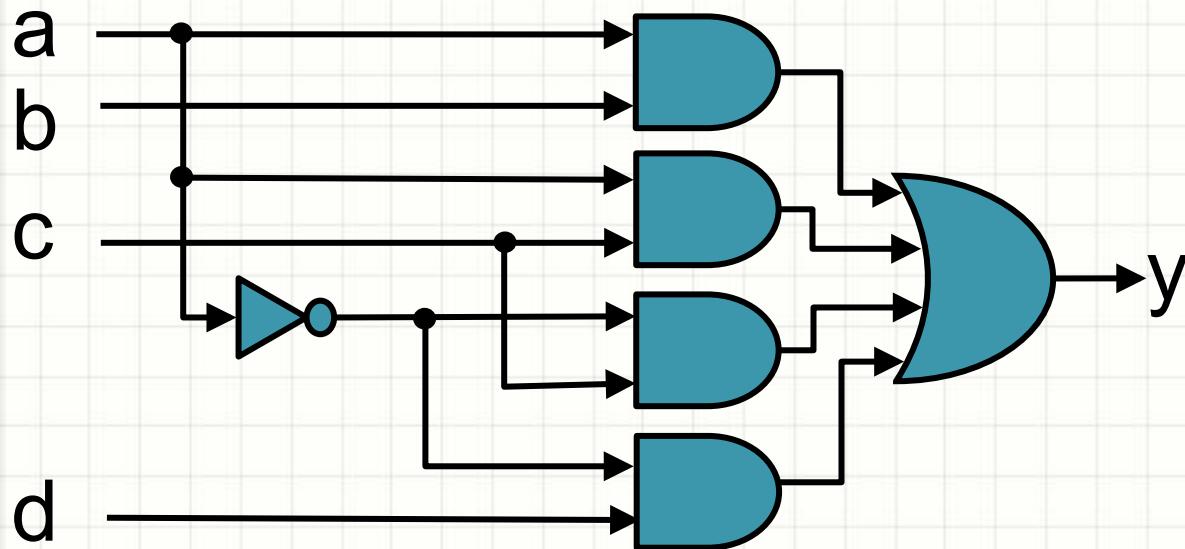
Prove $[(a + b)' = a' \cdot b']$

$$[(a + b)' = a' \cdot b']$$

$y = x'$ iff $[x \cdot y = 0 \text{ and } x + y = 1]$

- $(a + b) \cdot (a' \cdot b')$
 $= (a \cdot a' \cdot b') + (b \cdot a' \cdot b') = (0 \cdot b') + (a' \cdot 0)$
 $= 0 + 0 = 0$
- $(a + b) + (a' \cdot b')$
 $= a + (b + (a' \cdot b')) = a + (b + a')$
 $= (a + a') + b = 1 + b = 1$

Finding POS form



Sum of Product (SOP) form

$$y = (a \cdot b) + (a \cdot c) + (a' \cdot c) + (a' \cdot d)$$

$$y' = a'b'c'd' + a'b\bar{c}d' + a\bar{b}c'd' + a\bar{b}c'd$$

$$= a'c'd' + a\bar{b}c'$$

$$y = (a'c'd' + a\bar{b}c')' = (a'c'd')' \cdot (\bar{a}\bar{b}c')'$$

$$= (a + c + d) \cdot (a' + b + c)$$

This is Product of Sums (POS) form

a	b	c	d	y
0	0	0	0	0
0	0	0	1	1
0	0	1	0	1
0	0	1	1	1
0	1	0	0	0
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	1
1	1	0	0	1
1	1	0	1	1
1	1	1	0	1
1	1	1	1	1



QUESTIONS?