

2 4 6  
10 0 0  
11 0 0

2<sup>2</sup> 2<sup>2</sup>  
= 4

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23BCSI1290

KRG - 2B

→ Input:  $A = [1, 3, 5]$

Output: 8

Input:  $[2, 3]$

Output: 2

Q. We define  $f(x, y)$  as no. of different corresponding bits in the binary representation of  $x$  and  $y$ . For example,  $f(2, 7) = 2$  since binary representation of 2 and 7 are 010 and 111, respectively. The first and third bit differ so  $f(2, 7) = 2$ .

You are given an array of  $N$  integers,  $A_1, A_2, \dots, A_N$ . Find sum of  $f(A_i, A_j)$  for all pairs  $i, j$  such that  $1 \leq i, j \leq N$ . Return the answer modulo  $10^9 + 7$ .

# approach 1 (unoptimized)

just iterate two loops and add  $arr[i] \otimes arr[j]$  count of bits

→

~~int arr[n]~~

~~for~~



Pseudocode :

loop loop cnt = 0

int n; input n;

int arr[n];

for (int i = 0; i < n; i++) {

for (int j = 0; j < n; j++) {

if (i == j) continue;

cnt += ~~arr[i]~~ pop-count(arr[i] ^ arr[j]);

}

}

Print cnt / ( (int) (1e9 + 7) );

Time Complexity :  $O(n^2)$

Space Complexity :  $O(1)$

# Approach 2

for every bit for each number,  
count set bit,

now for each (i, j) differing  
bit,

set count \* unset count is the  
count of differing bits pair for  
that bit



```

pseudocode: int n; input int arr[n]; input arr;
long long cnt = 0;
for (int b = 0; b < 32; b++) {
    long long k = 0;
    for (int i = 0; i < n; i++) {
        if (arr[i] & (1 << b)) {
            k++;
        }
    }
    cnt += (k) * (n - k);
}
cnt *= 2;
print cnt % (int)(1e9 + 7);

```

Time Complexity :  $O(n)$   
Space Complexity :  $O(1)$

Shubh