

$$2^2 - 4 = 2$$

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KRn - 2B

→ Input: $A = [1, 3, 5]$

Output: 8

Input: $[2, 3]$

Output: 2

Q. We define $f(x, y)$ as no. of different corresponding bits in the binary representation of x and y . For example, $f(2, 7) = 2$ since binary representation of 2 and 7 are 010 and 111, respectively.

The first and third bit differ so $f(2, 7) = 2$.

You are given an array of N integers, A_1, A_2, \dots, A_N . Find sum of $f(A_i, A_j)$ for all pairs (i, j) such that $1 \leq i, j \leq N$. Return the answer modulo $10^9 + 7$.

approach 1 (unoptimized)

just iterate two loops and add $\text{arr}[i] \otimes \text{arr}[j]$ count of bits

~~not done~~

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Pseudocode :

```
1. for loop cnt = 0  
2.   int n; or input n;  
3.   int arr[n];  
4.   for (int i = 0; i < n; i++) {  
5.     for (int j = 0; j < n; j++) {  
6.       if (i == j) continue;  
7.       cnt += arr[i] popCount(arr[i] & arr[j]);  
8.     }  
9.   }
```

print **cnt** % (1000000000 + 7);
Time Complexity : $O(n^2)$
Space Complexity : $O(1)$

Approach 2

for every bit, for each number,
count set bit,

now for each (i, j) differing
bit,

SetCount + unsetCount is the
Count of differing bits pair for
that bit.

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```
pseudocode: int n; input int arr[n]; input out;
loop loop (cnt = 0)
for (int b = 0; b < 32; b++) {
    loop loop (k = 0)
        for (int i = 0; i < n; i++) {
            if (arr[i] & (1 << b)) {
                k++;
            }
        }
        cnt += (k)(n - k);
    }
    cout * = 0;
    print cout % (int)(1e9 + 7);
}
```

time complexity: $O(n)$
Space complexity: $O(1)$

Ans