

Hashes, MACs & Authenticated Encryption

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Today's Lecture

- Hashes and Message Authentication Codes
 - Properties of Hashes and MACs
 - CBC-MAC, MAC \rightarrow HASH (slow),
 - SHA1, SHA2, SHA3
 - HASH \rightarrow MAC, HMAC
- Authenticated Encryption
 - CCM

Hashes

- A hash of any message is a short string generated from that message.
- The hash of a message is always the same.
- Any small change makes the hash totally different.
- It is very hard to go from the hash to the message.
- It is very unlikely that any two different messages have the same hash.

Signatures

- Using RSA $E_{\text{pub}}(D_{\text{priv}}(M)) = M$
- This can be used to sign messages.
- Sign a message with the private key and this can be verified with the public key.
- Any real crypto suite will not use the same key for encryption and signing.
 - as this can be used to trick people into decrypting.

Signatures

Alice has a signing key K_s
and wants to sign message M

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Detached Signature: $D_{k_s}(\#(M))$

- Firstly hash message
- Secondly, decrypt with private key k_s



Uses of Hashing

- Download/Message verification
- Tying parts of a message together (hash the whole message)
- Hash message, then sign the hash.
- Protect Passwords
 - Store the hash, not the password

- Hash of VM
- md5 VM
- shasum VM
- shasum -a 512 VM

Attacks on hashes

- Preimage Attack: Find a message for a given hash: very hard.
- Prefix Collision Attack: a collision attack where the attacker can pick a prefix for the message.
- Collision Attack: Find two “random” messages with the same hash.

Birthday Paradox

- How many people do you need to ask before you find 2 that have the same birthday?
- 23 people, gives $(23*22)/2 = 253$ pairs.
- Prob. that two people have a different birthday is: $364/365$
- $(364/365)^{(23*22/2)} = 0.4995$

Message Authentication Codes

- MACs are hashes with a key.
 - Written $\text{MAC}_{\text{key}}(M)$
- You can only make or check the hash, if you know the key.
- Stops guessing attacks.

Message Authentication Codes

- MACs are sometimes used for authentication:
 - E.g. in Alice and Bank share keyA, Alice sends to the bank:
“Pay Bob £10”, $\text{MAC}_{\text{keyA}}(\text{“Pay Bob £10”})$

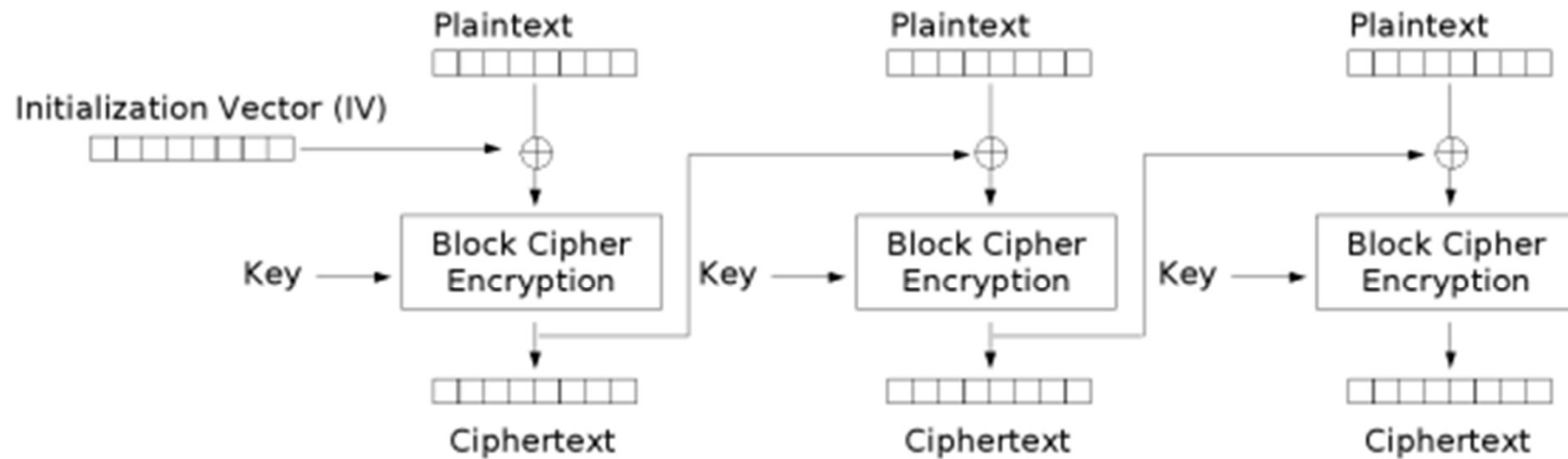
Possible attack on MAC: “Length extension attack” add data to a MAC without knowing the key

How can we make a MAC?

Must have:

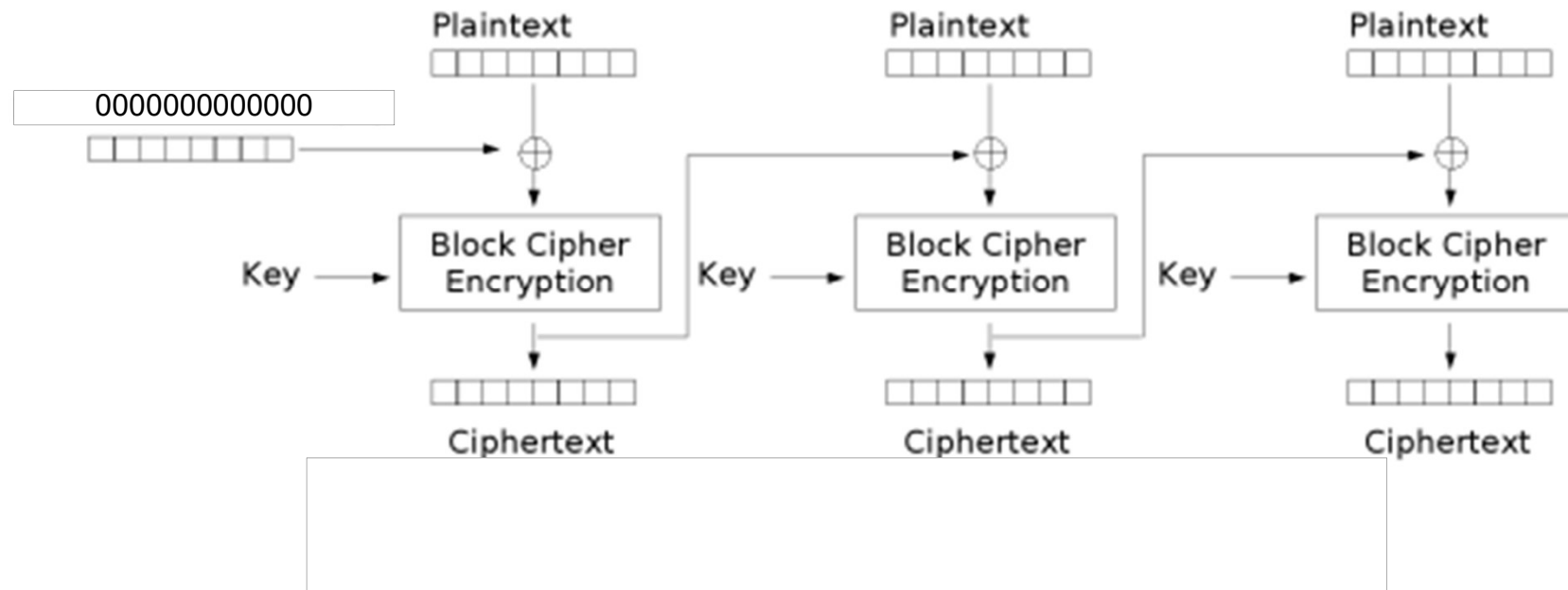
- A Key
- Short string from long message.
- Single bit change in message = totally different hash.
- Hard to go from the hash to the message.
- Unlikely that any two different messages have the same hash.

CBC Encryption

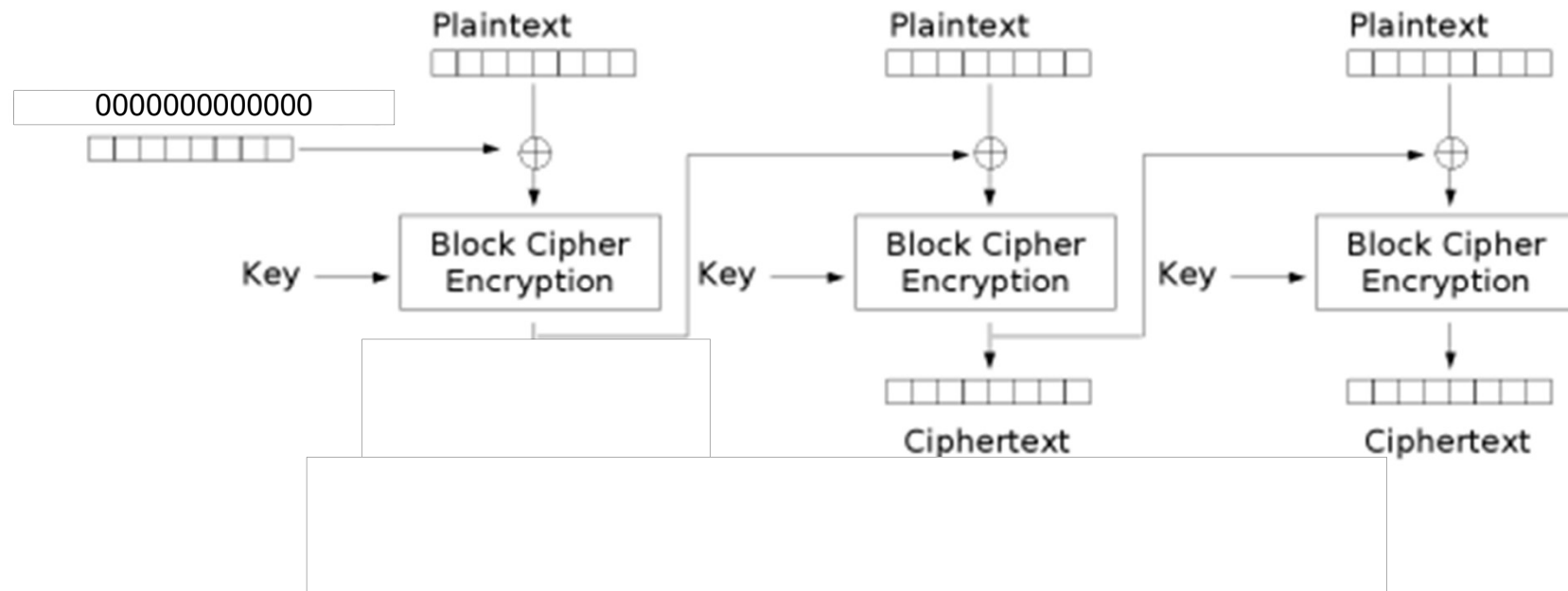


Cipher Block Chaining (CBC) mode encryption

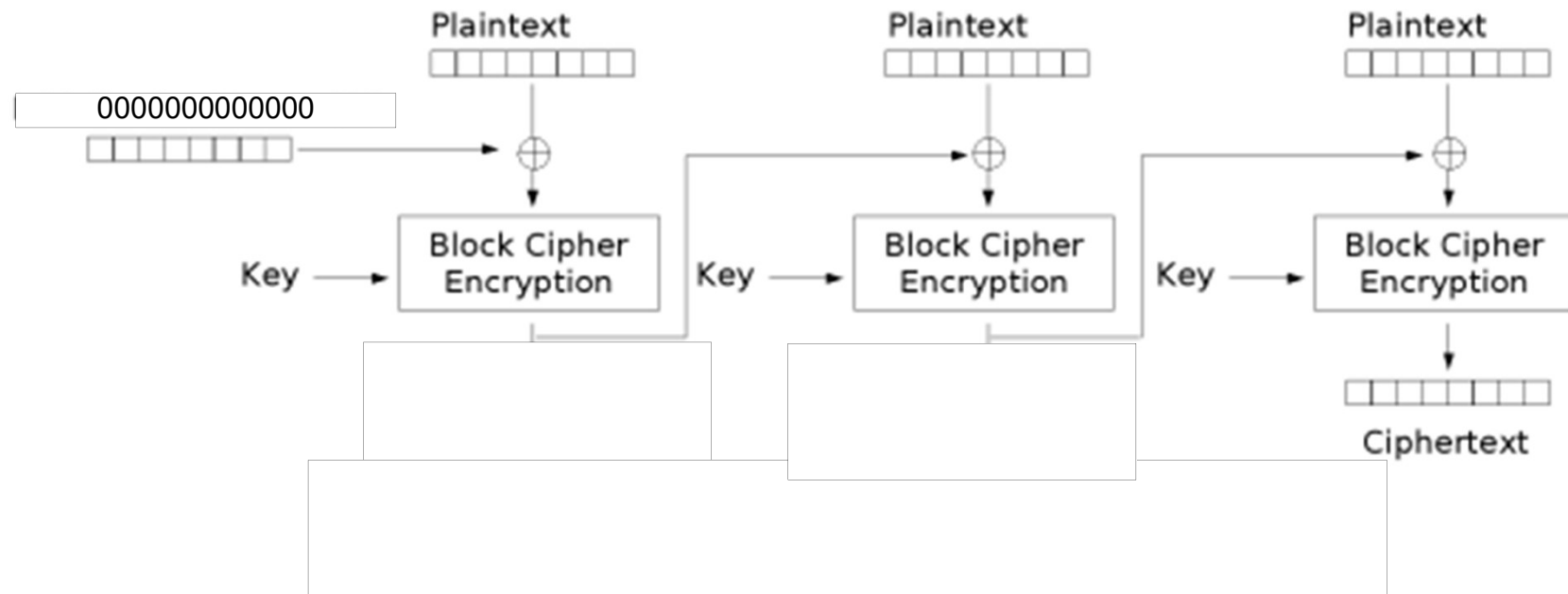
Making a CBC MAC



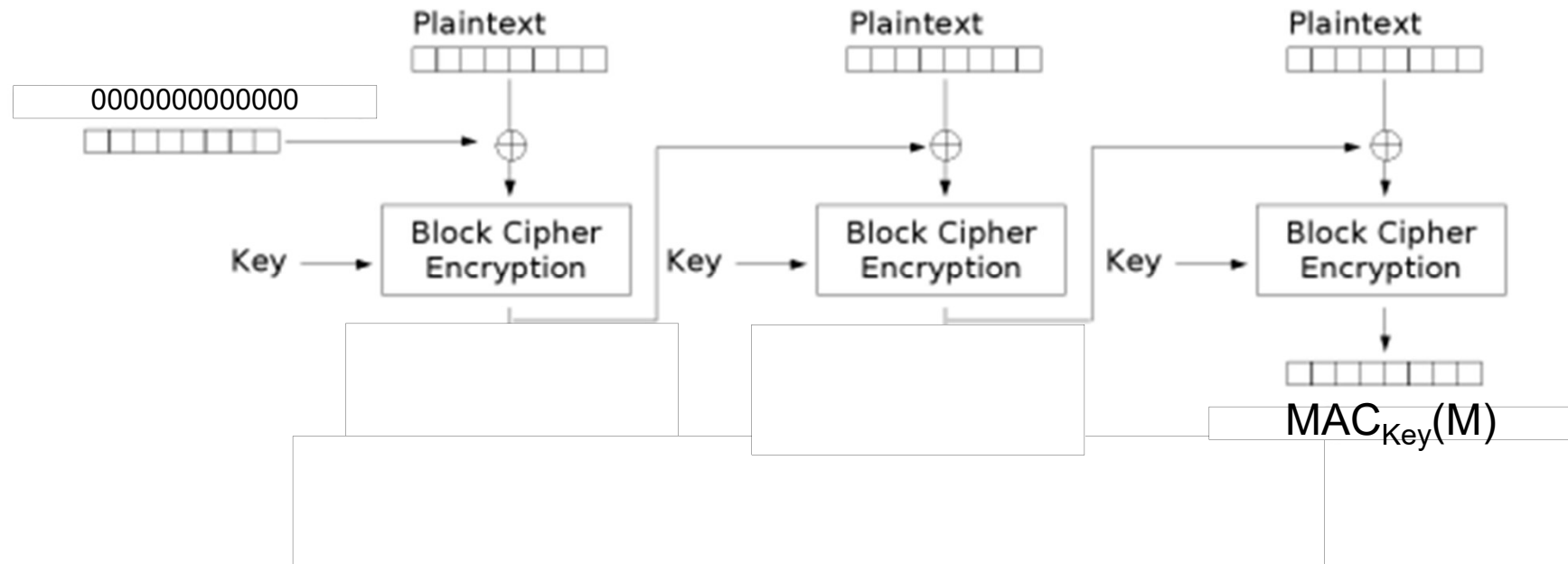
Making a CBC MAC



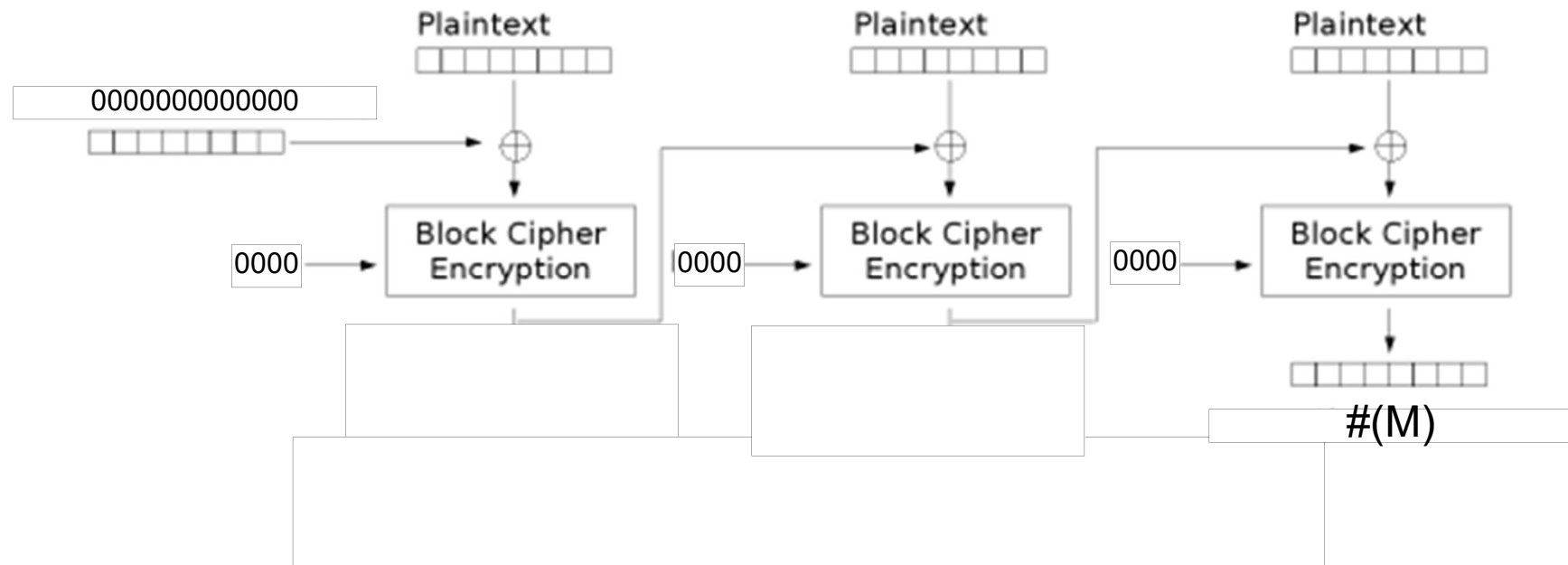
Making a CBC MAC



CBC MAC



An Inefficient Hash Function



The SHA Family of Hash

- The most common (and best) hashes are the SHA hashes.
- 1993, The US National Institute of Standards and Technology (NIST), developed a new hash SHA-0
- 1995, the NSA stepped in and “fixed” it: SHA-1 (160-bit hash).

SHA1

- A birthday attack on SHA-1 should need 2^{80} hash tests
- In 2005 a 2^{63} attack was found.
- Not really practical, but no-one trusts SHA-1 any more.
- So ... SHA-2

SHA2

- SHA2 is an improved version of SHA1 with a longer hash.
- 256 or 512 bits: also called SHA256, SHA512.
- Based on SHA-1 it has some of the same weaknesses. So, even though it seems secure the cryptographers aren't happy.

The SHA-3 Competition

- Submissions opened on October 31, 2008,
- Round 1
 - 13 submissions rejected without comment.
 - 10 withdrawn by authors.
 - 16 rejected for design or performance.
 - Inc. Sony's
- Conference in Feb 2009. 14 scheme picked to go through to round 2.
 - Dropped schemes include
 - Ron Rivest's,
 - Lockheed Martin

The SHA-3 Competition

- Winner announced on October 2, 2012 as Keccak, (Daemen et al. the AES guy)
- Adopted as NIST-standard in 2015

Merkle–Damgård (MD) Hashes

- The MD family of hashes is also popular.
- MD4 & MD5 used, but weak.
 - Only useful when we only care about preimage attacks or Integrity.
- MD6: Ron Rivest's candidate for SHA3.
 - Seems good & fast.

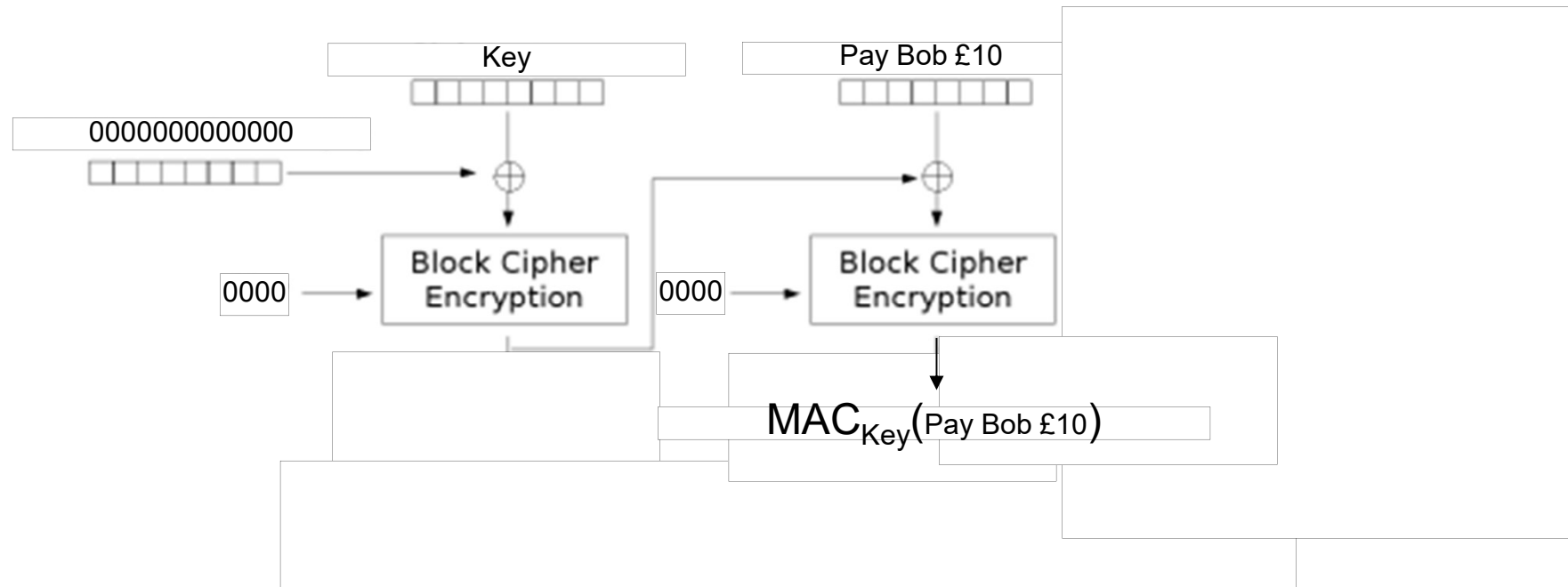
Broken Hash to MAC

- If we had a Hash we could try to make a MAC by:

$$\text{MAC}_{\text{Key}}(M) = H(\text{Key}, M)$$

- But this might allow a length extension attack.

Broken Hash to MAC

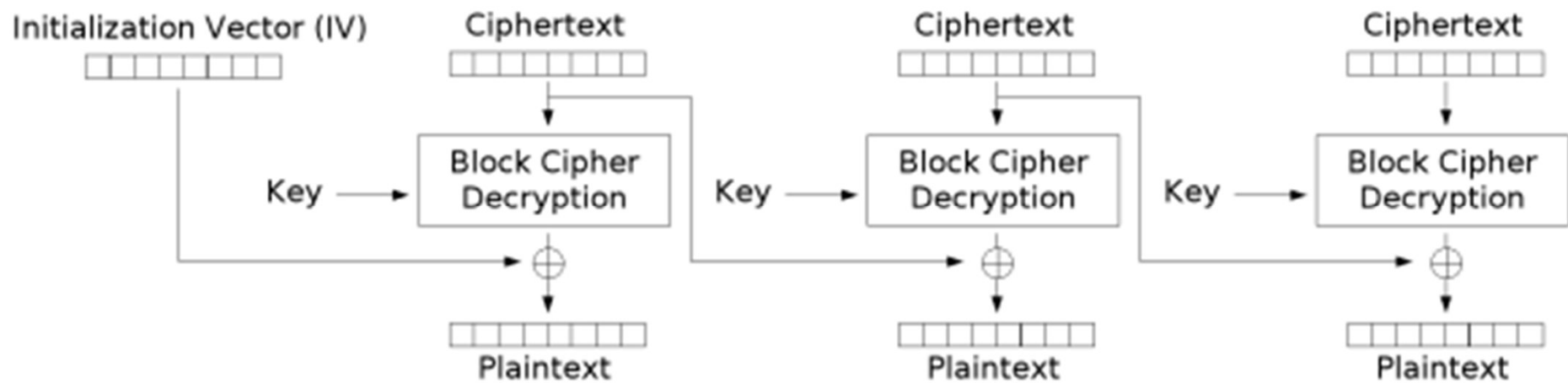


Only Alice and the Bank know “Key”
Alice sends “Pay Bob £10, $\text{MAC}_{\text{Key}}(\text{Pay Bob } \pounds 10)$ ” to Bank

Cipher Texts Can Be Altered

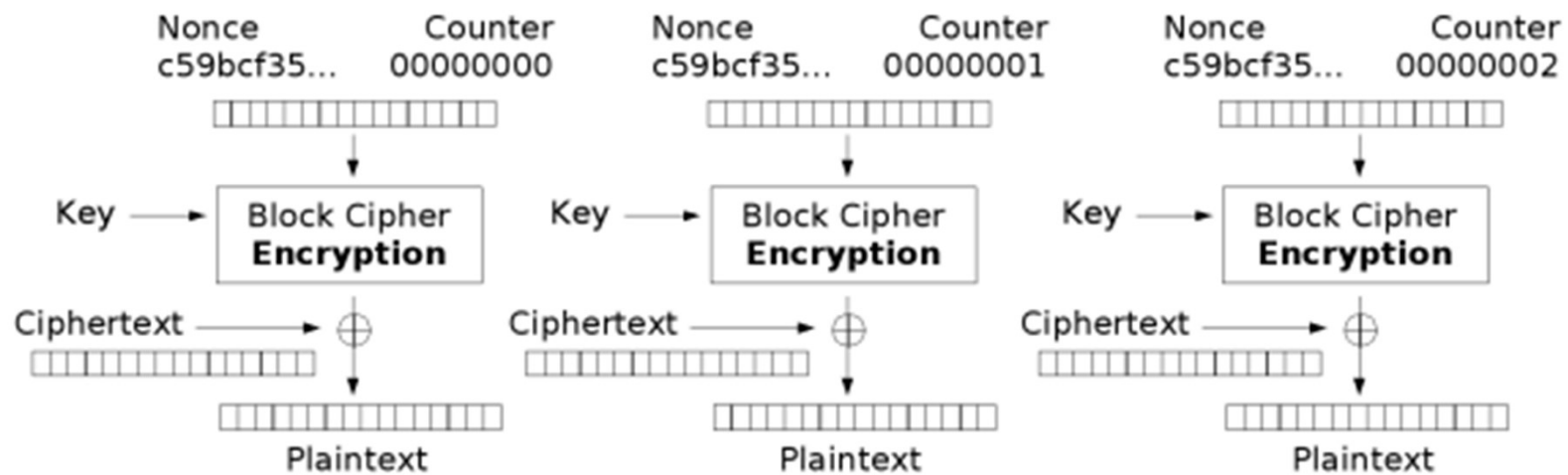
- AES encryption with a particular key maps any 128-bit block to a 128-bit block (or 256)
- AES decrypt also maps any 128-bit block to a 128-bit block.
- Decrypt can be run on any block (not just encryptions).

CBC decrypt



Cipher Block Chaining (CBC) mode decryption

CTRdecrypt



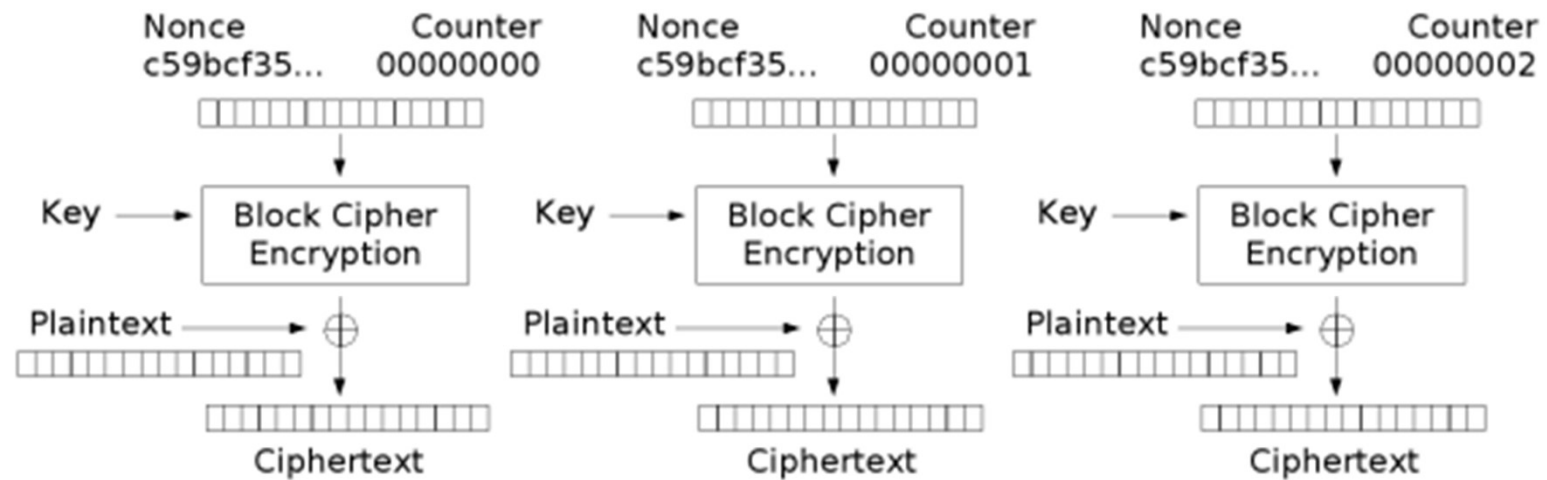
Counter (CTR) mode decryption

Block mode

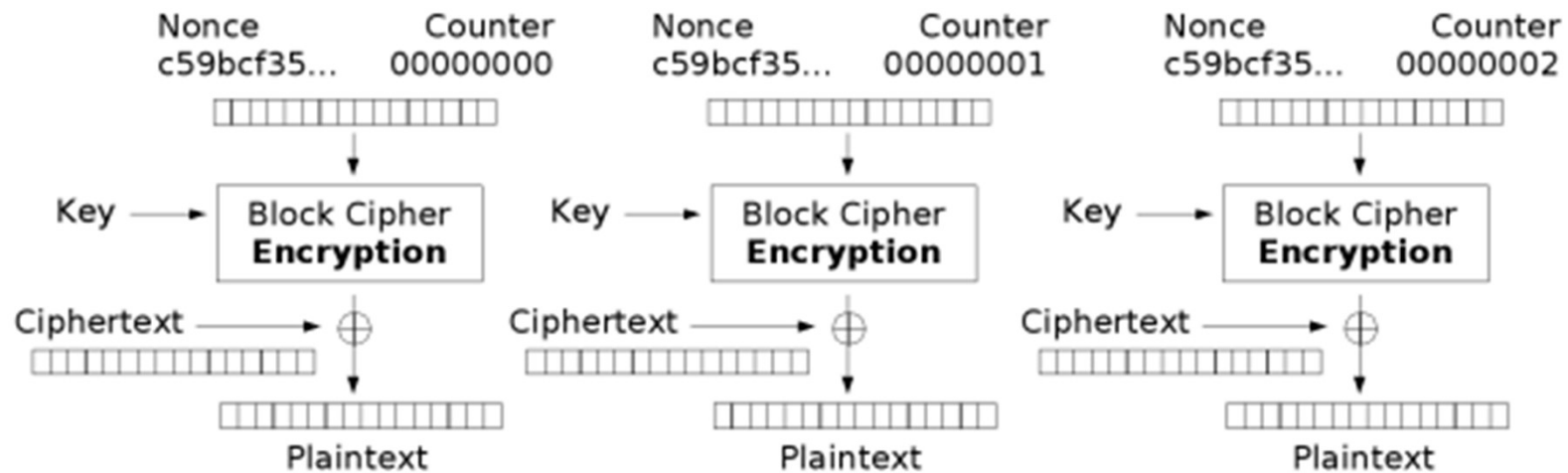
- CBC mode: any change affects all of the rest of the message.
- ECB mode: any change affects only the block.
- CTR mode: any change affects only the bits altered.

Known Plain Text Attacks

- If I know the plaintext I can change CTR encrypted messages. (see previous lecture)



Counter (CTR) mode encryption



Counter (CTR) mode decryption

Authenticated Encryption Modes

- Authenticated encryption modes stop this.
- With Authenticated Encryption you can only form a valid ciphertext if you know the key.
- Most common way to do this is to add a MAC to the ciphertext.

CCM mode encryption

- First calculate an AES CBC-MAC on the data.
- Then encrypt the message followed by the MAC using the same key and CTR mode.
- Not rocket science, but proven secure
 - Fully defined as RFC 3610

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