

Fully Homomorphic Encryption I

Zama's FHE compiler

Agenda

FHE	04
TFHE	13
Concrete Optimizer	21
Conclusion	33

FHE

What is FHE ?

Past



Cloud Computing

Past



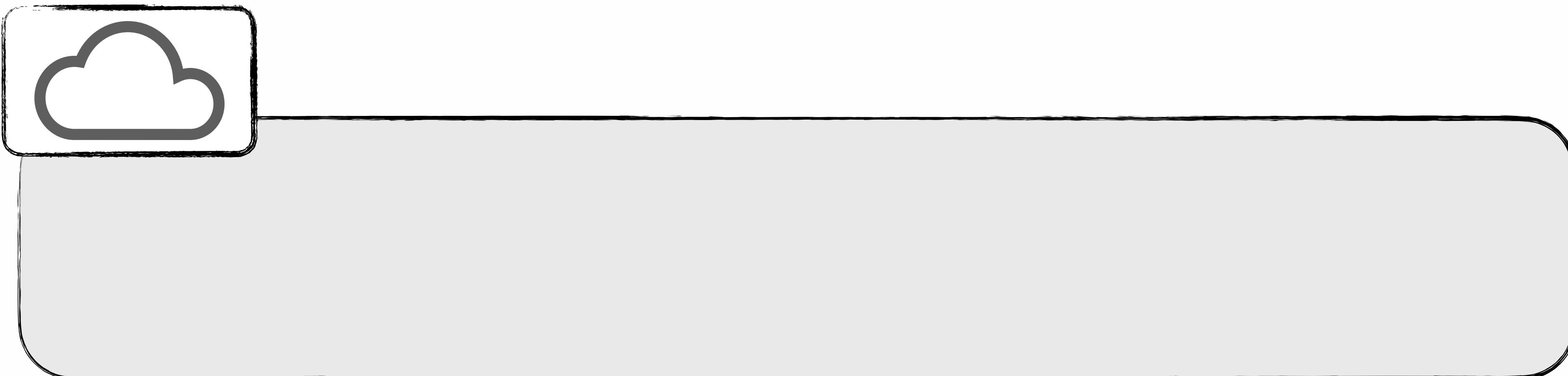
Cloud Computing



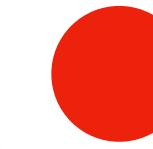
Past



Cloud Computing



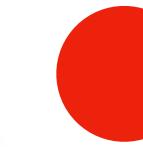
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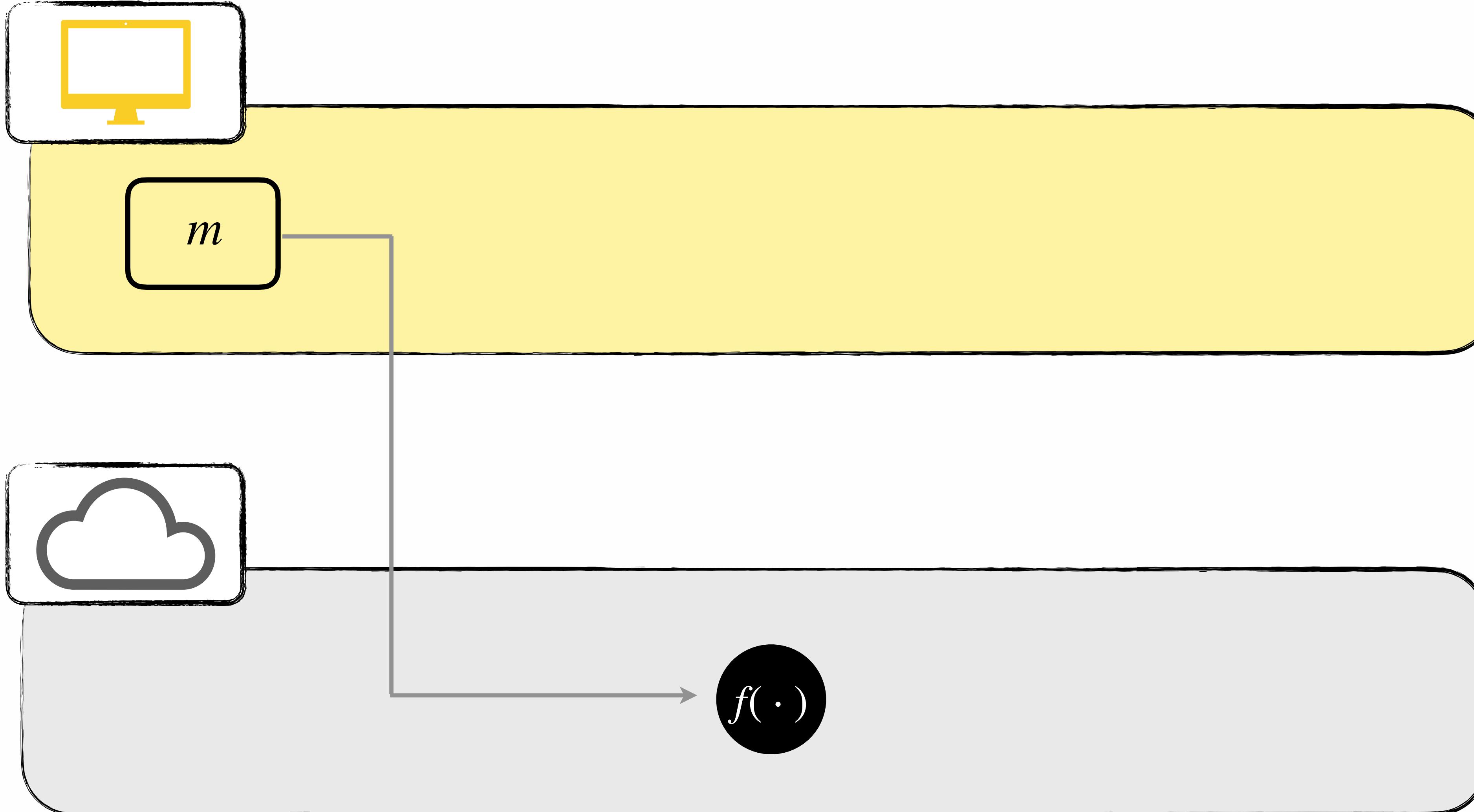
Cloud Computing



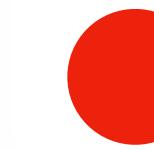
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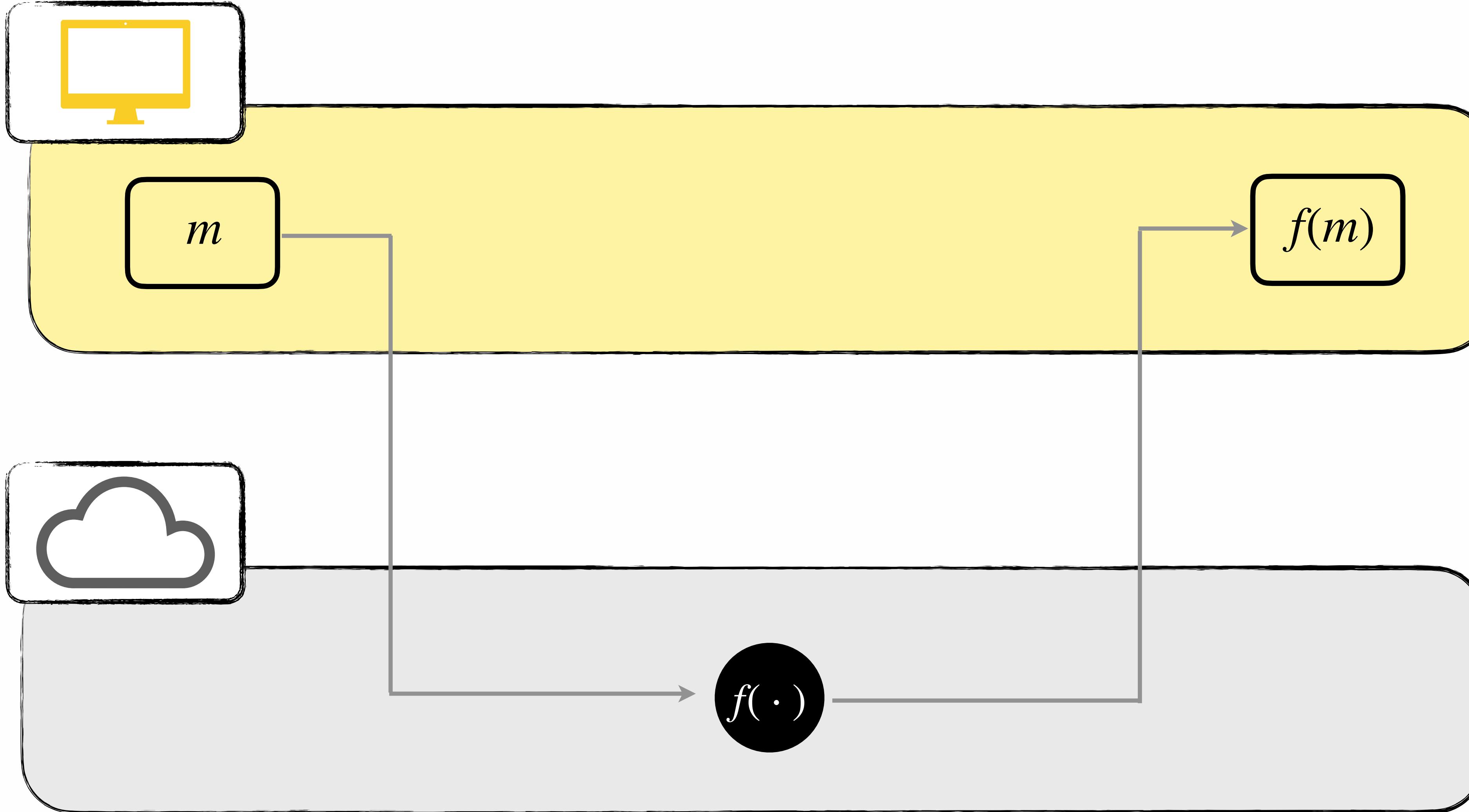
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Past



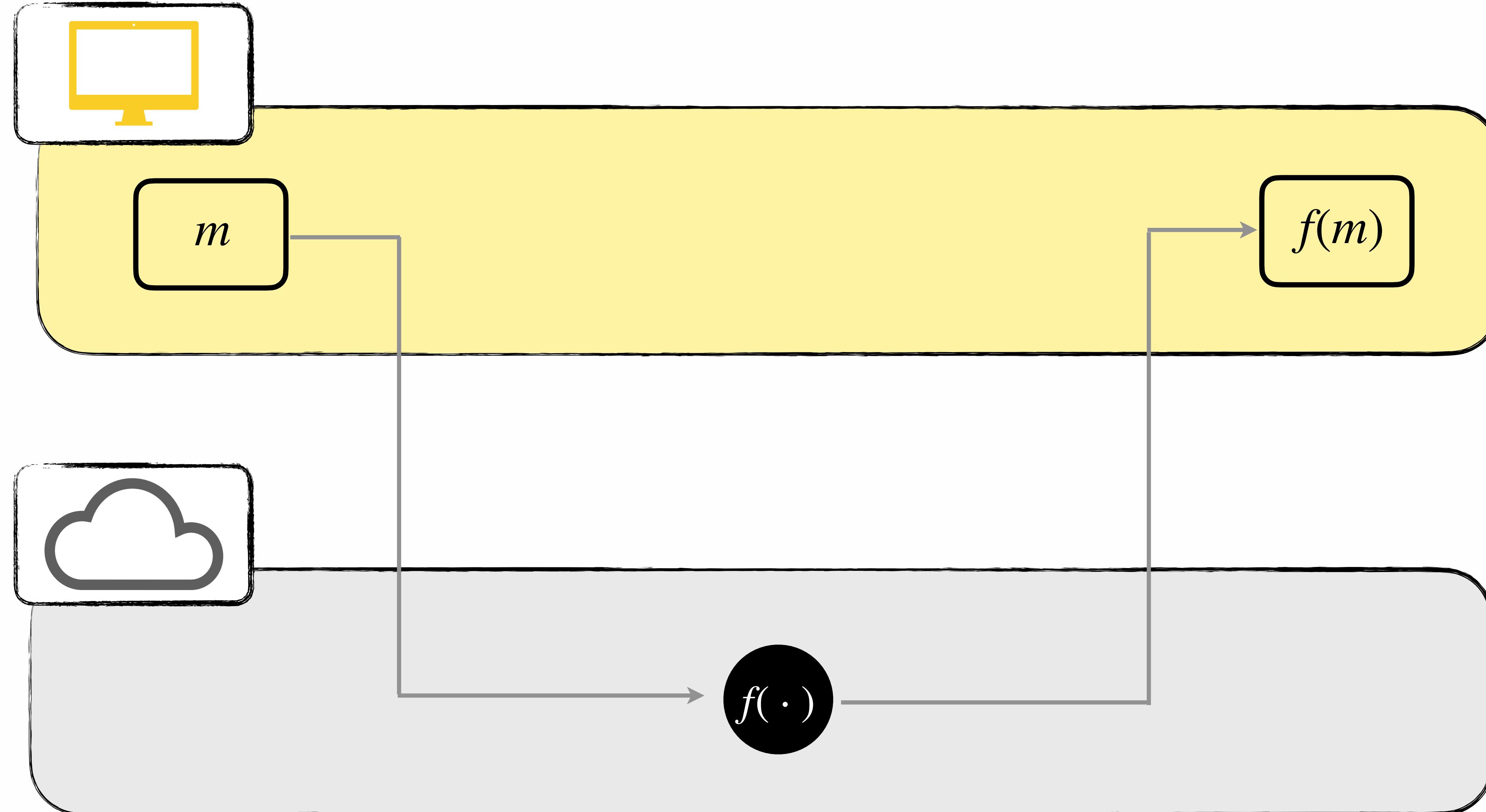
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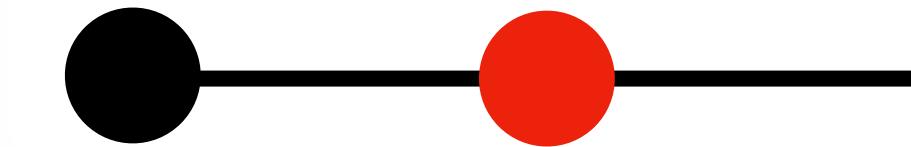
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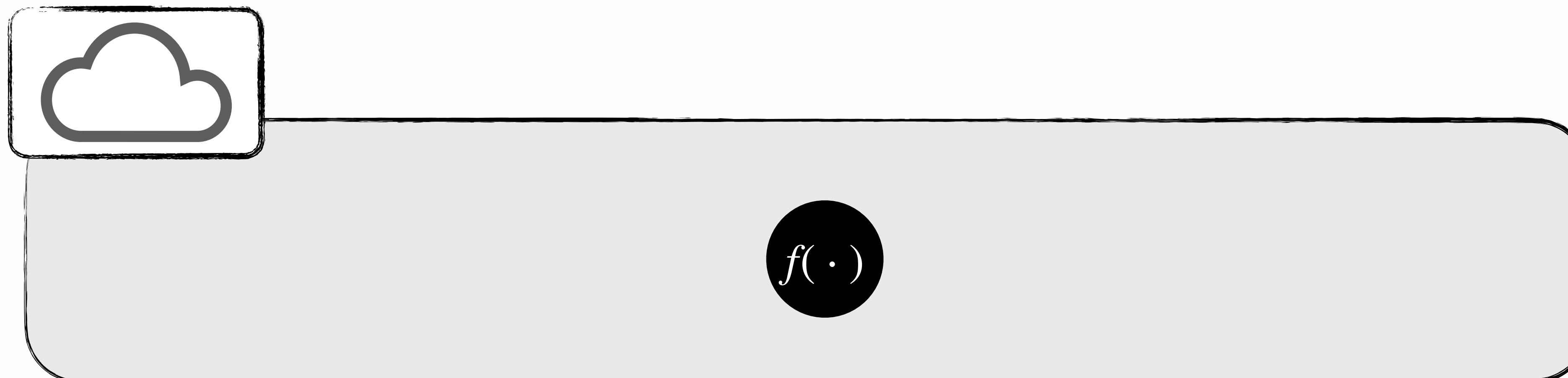
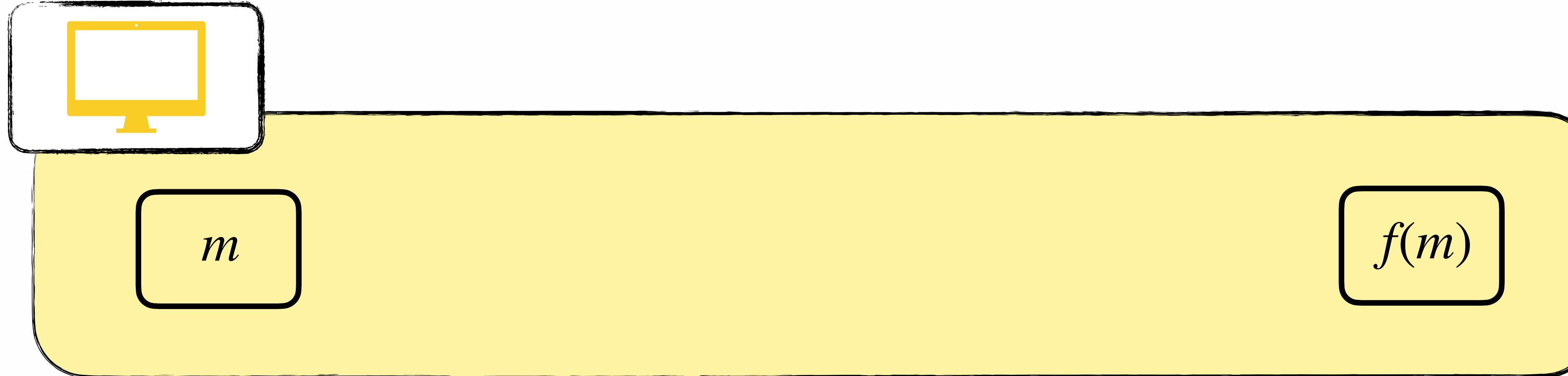
Cloud Computing



Past Present

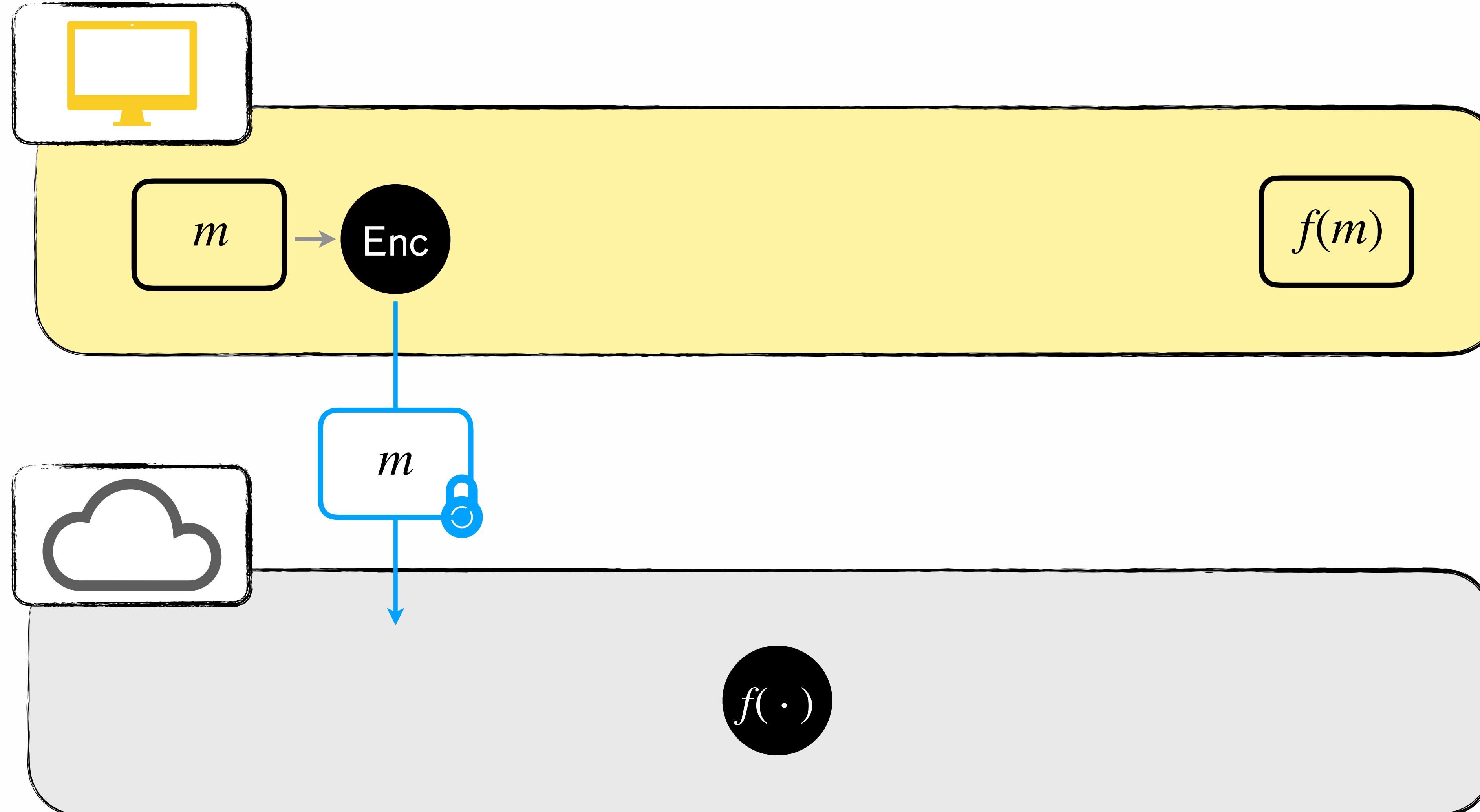


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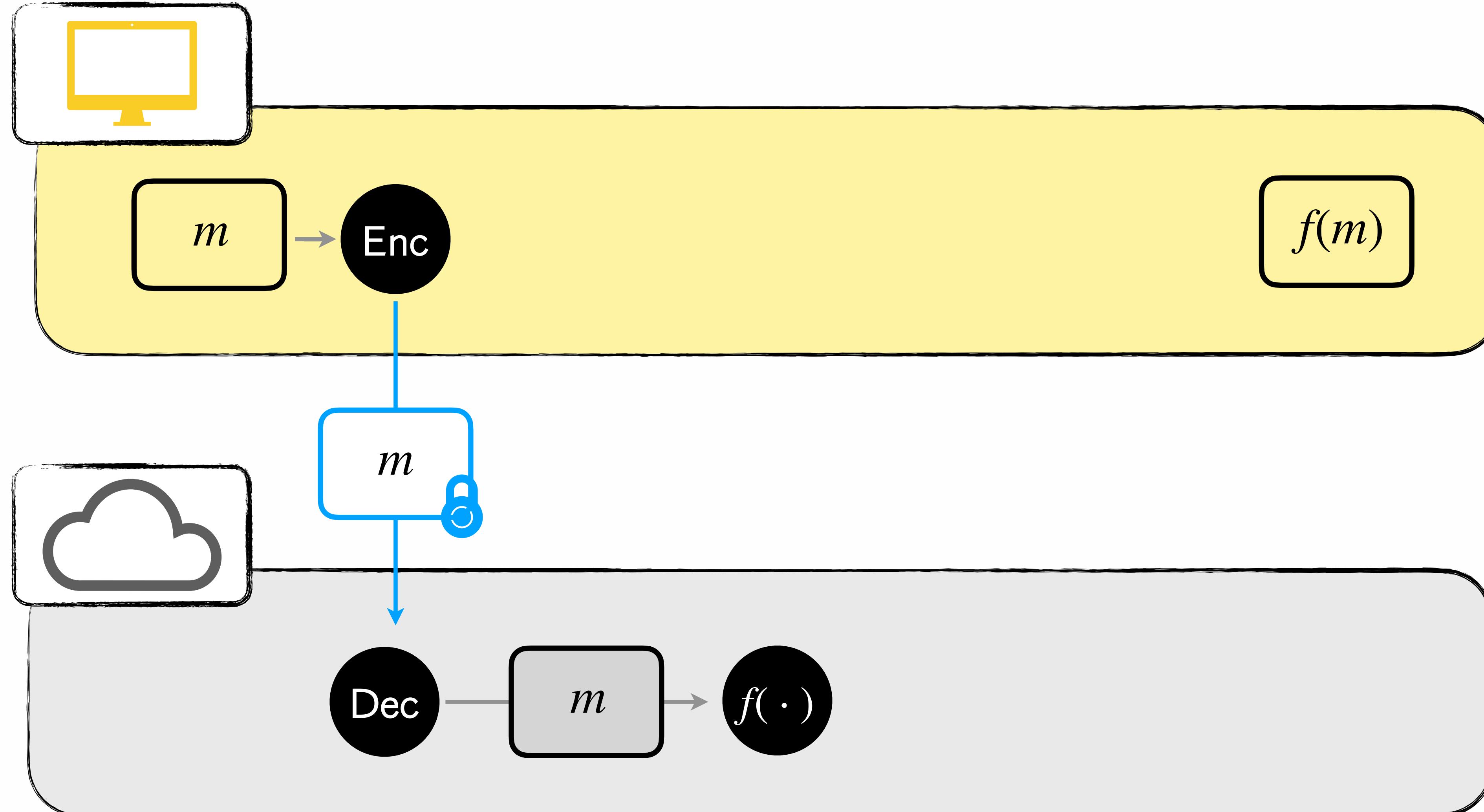
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Cloud Computing



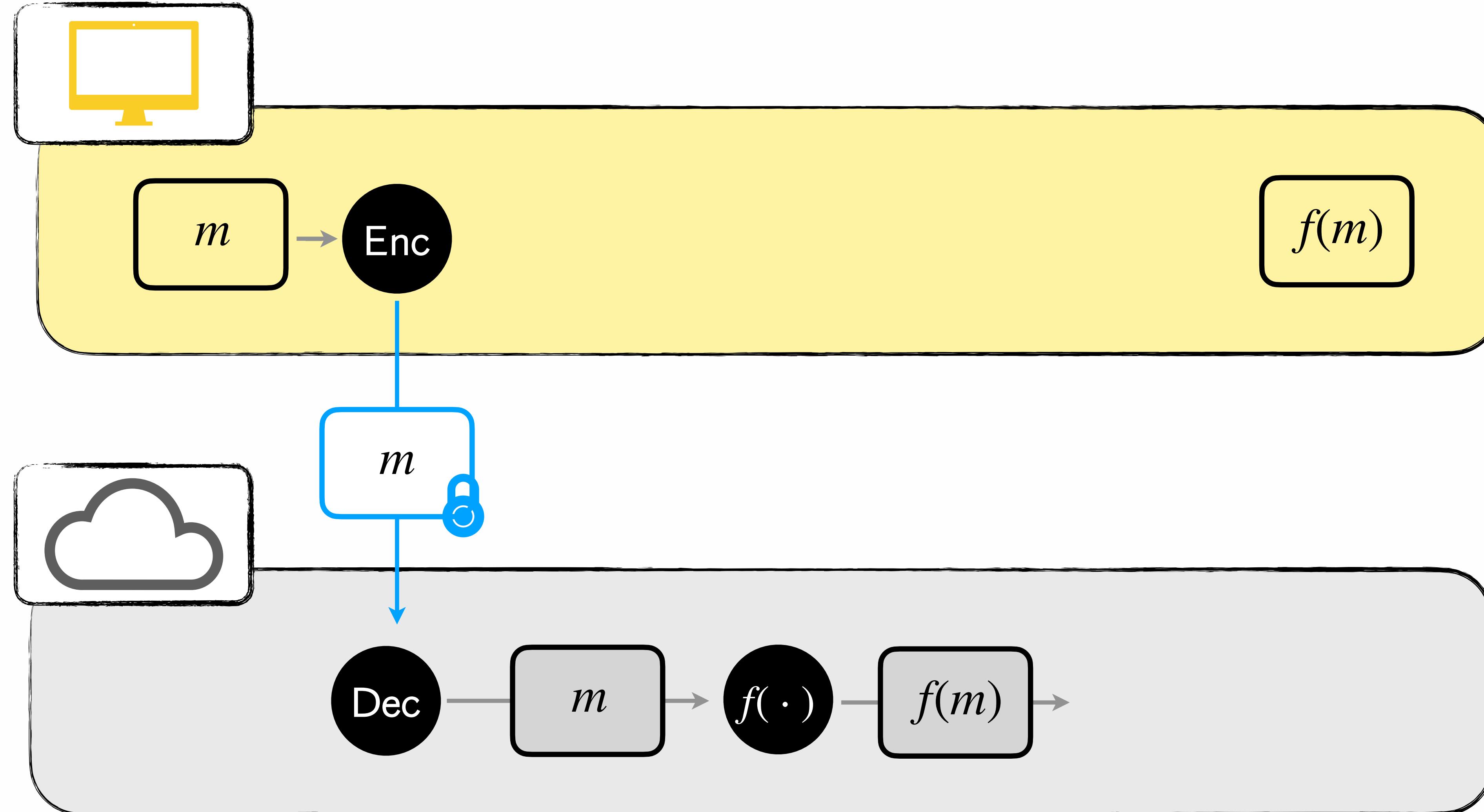
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Cloud Computing



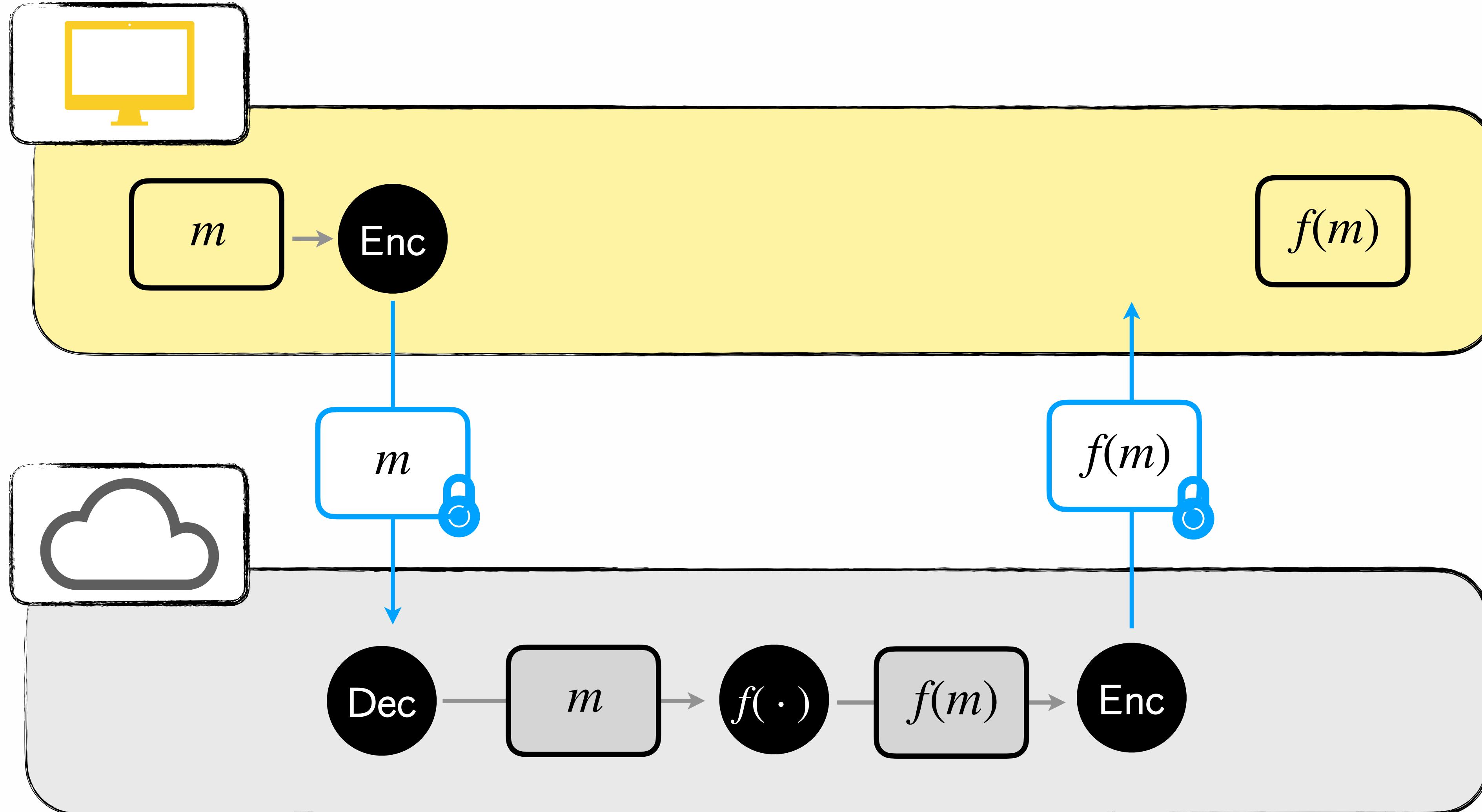
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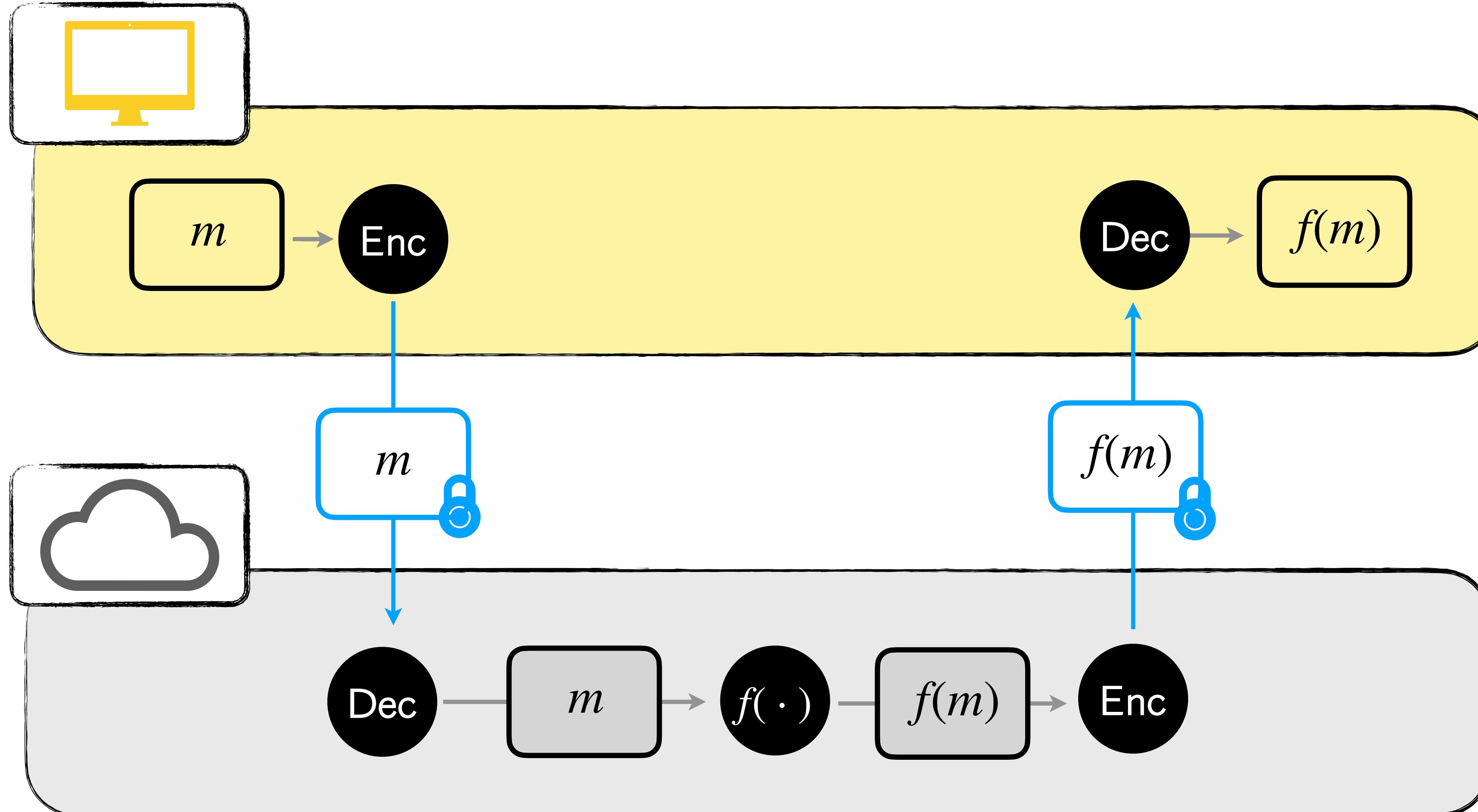
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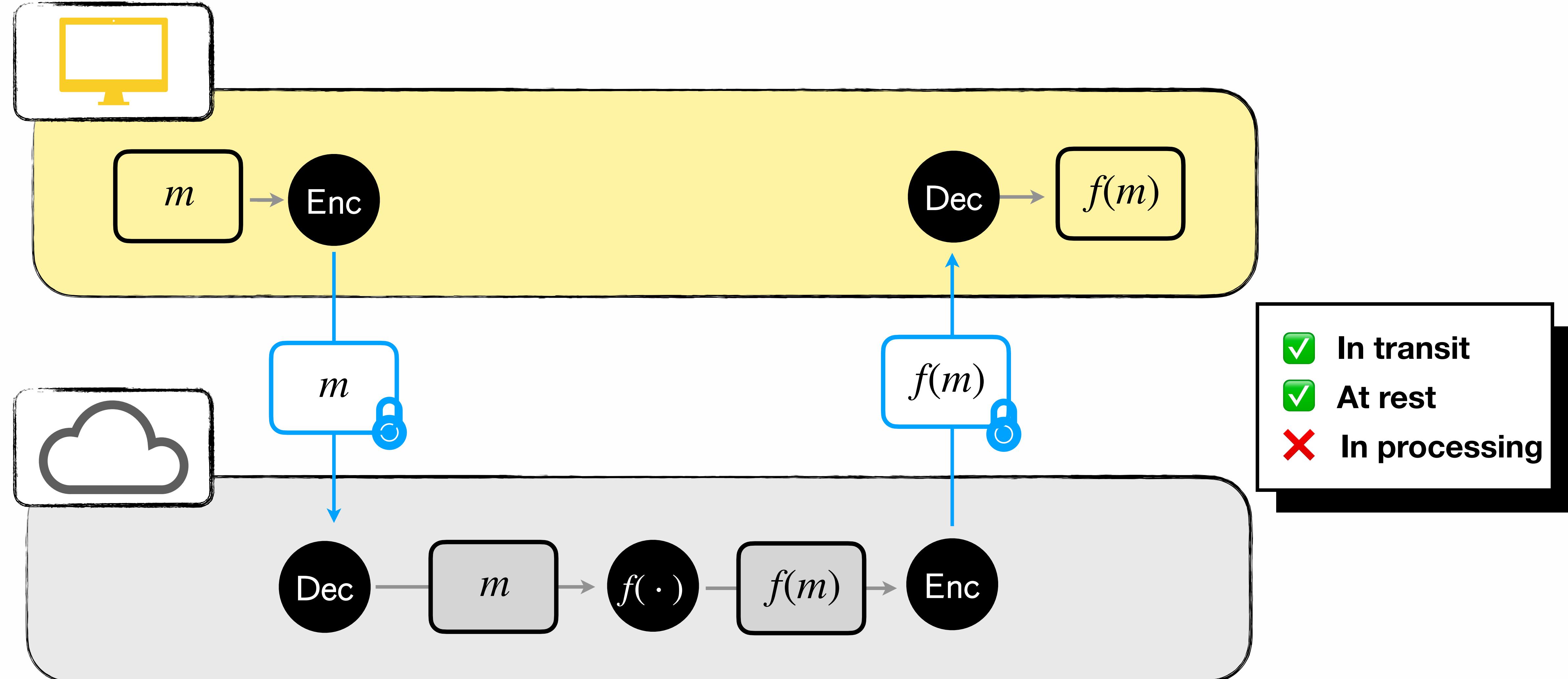
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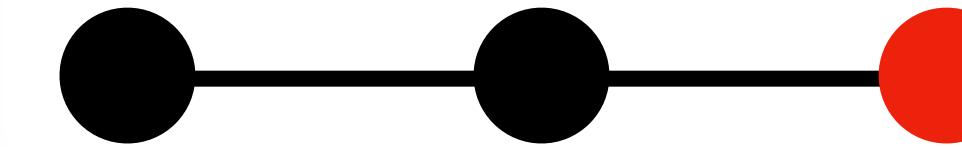


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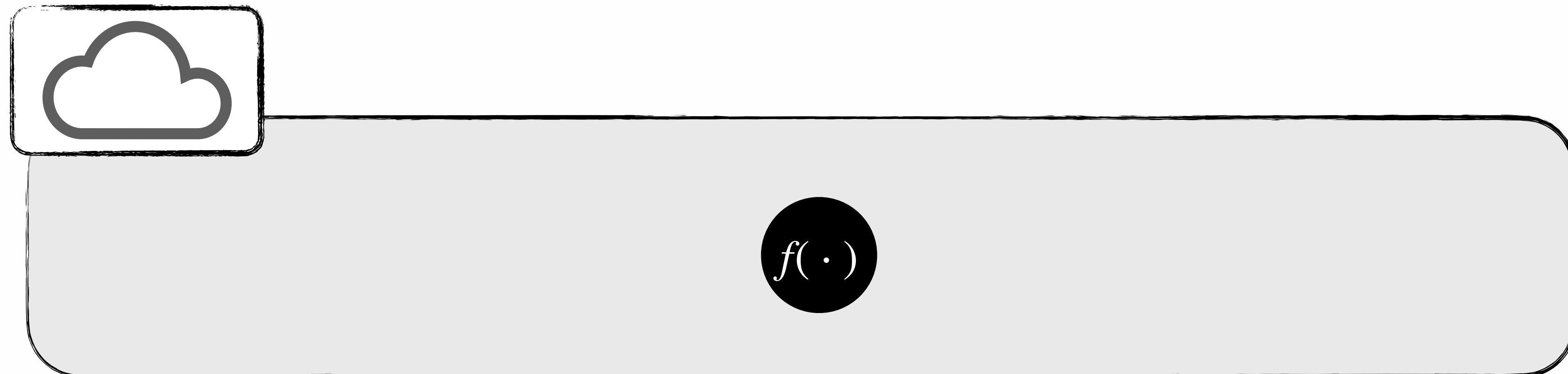
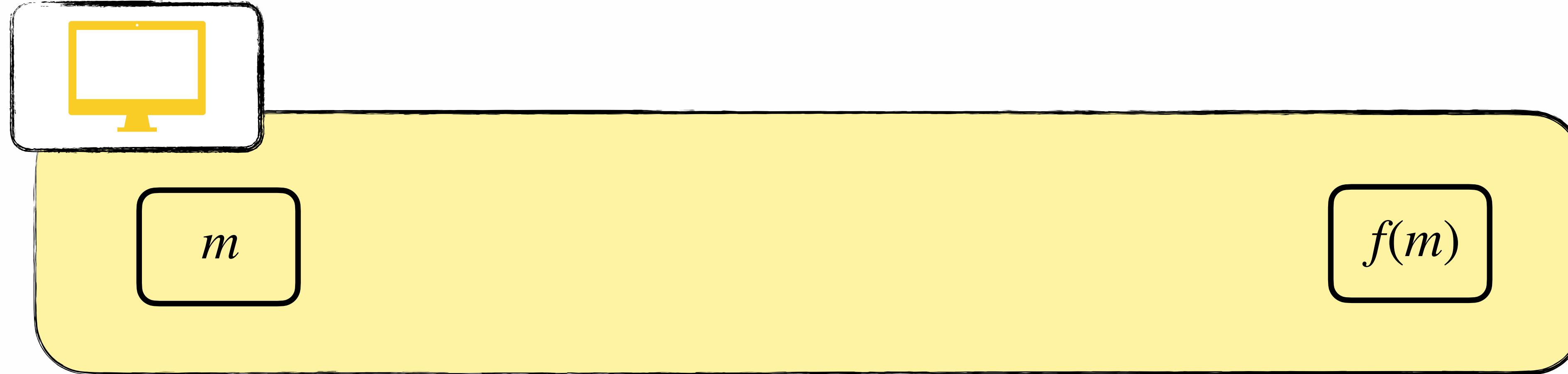
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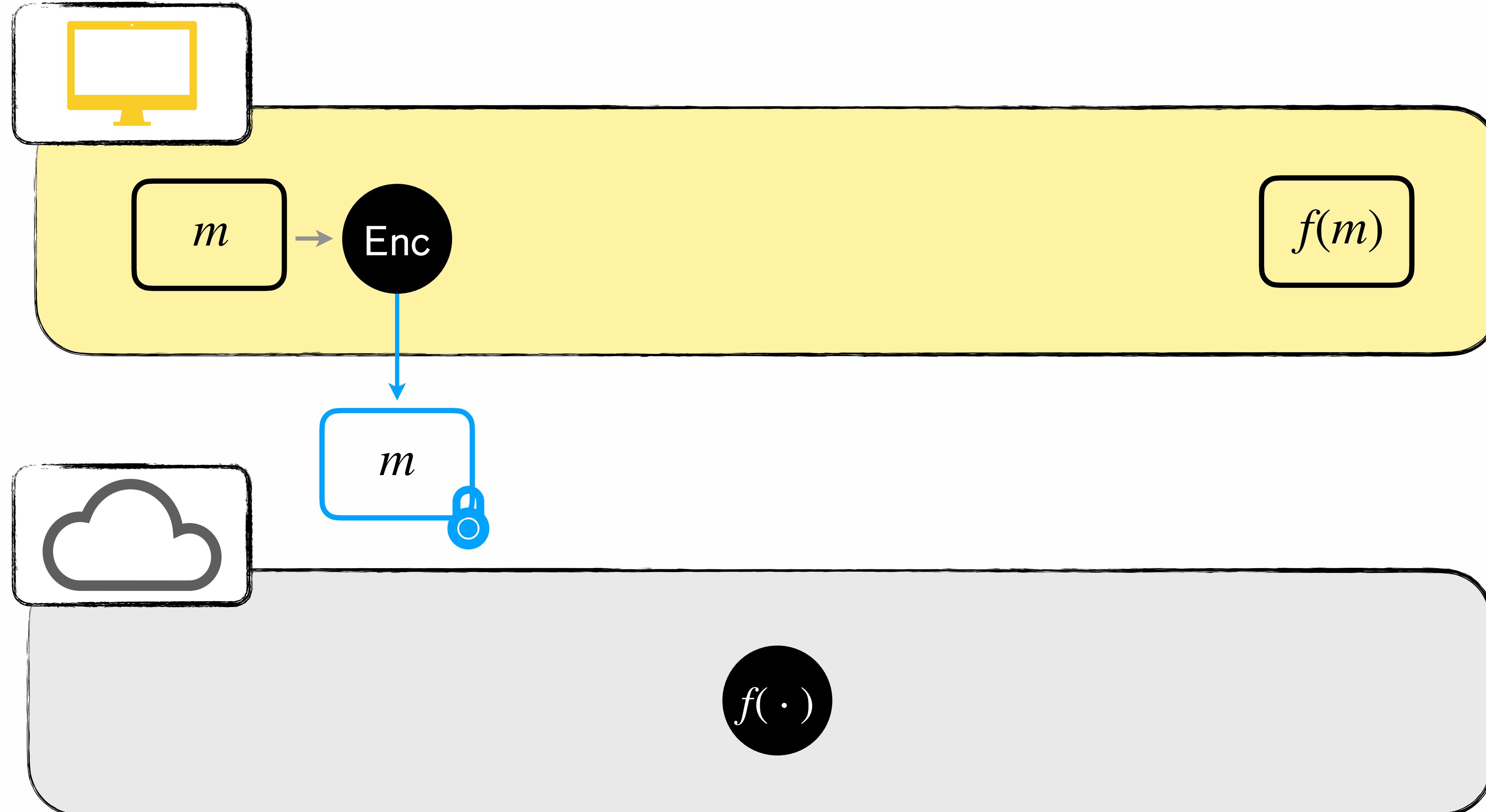
Past Present FHE



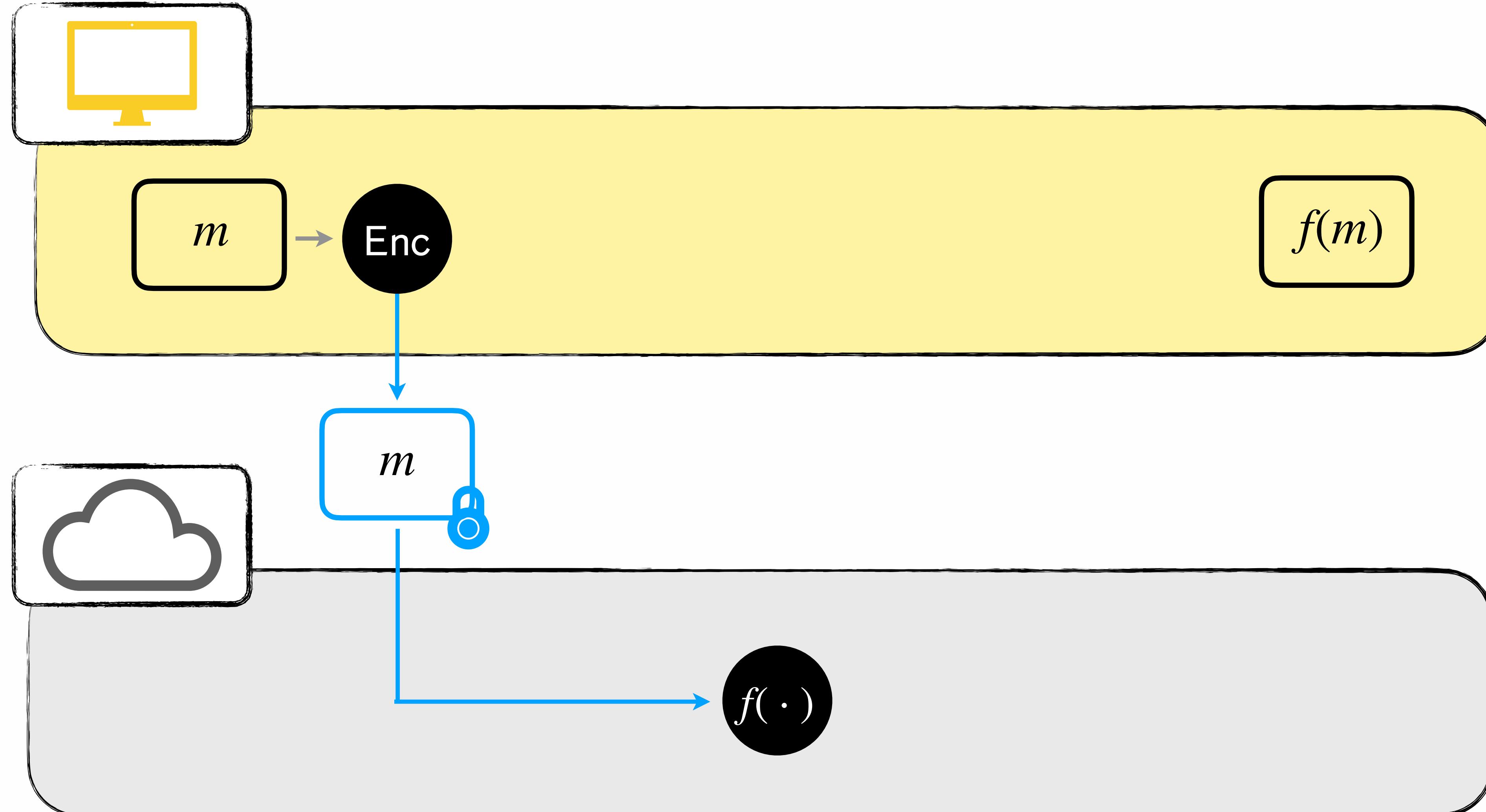
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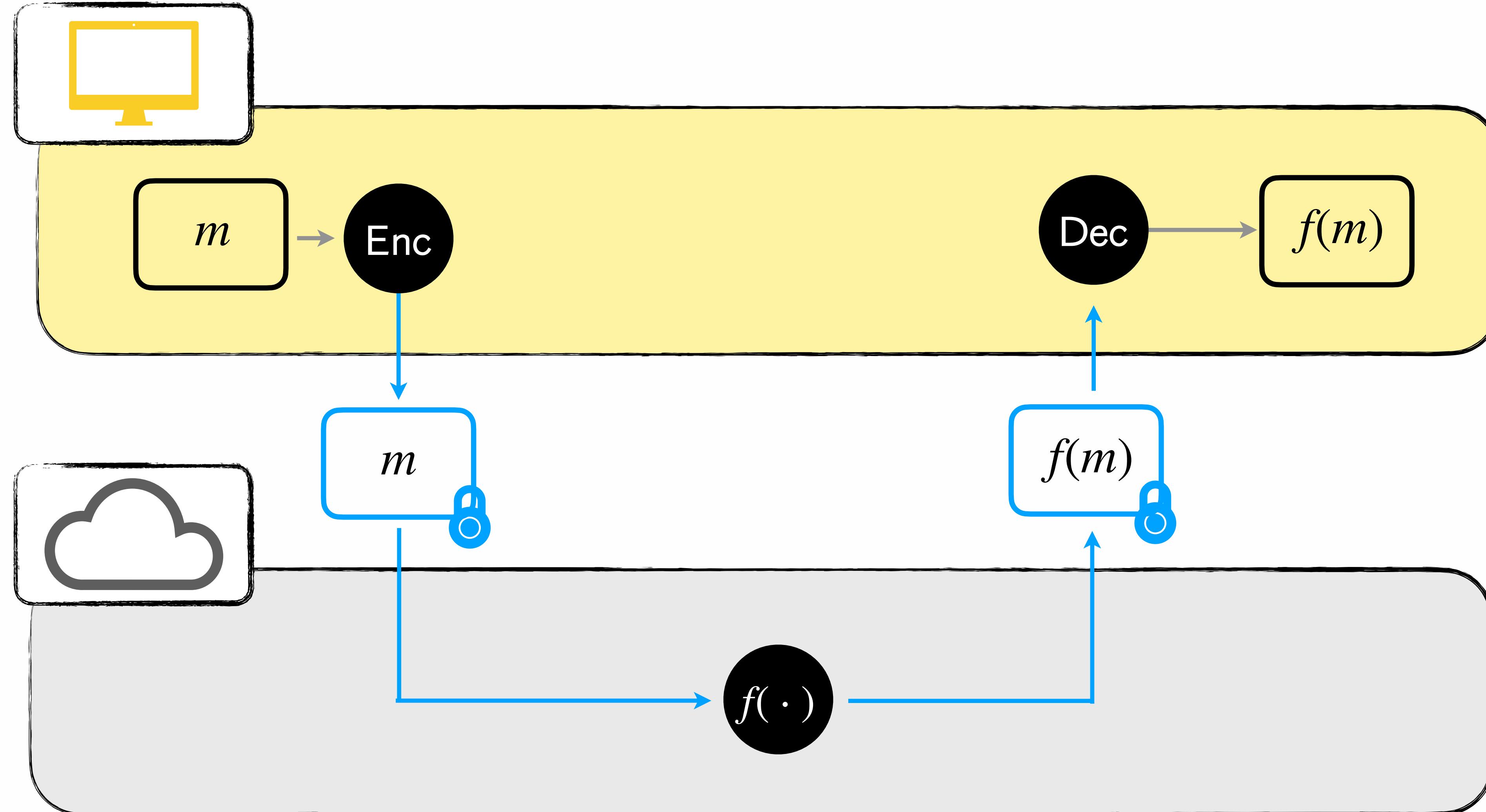
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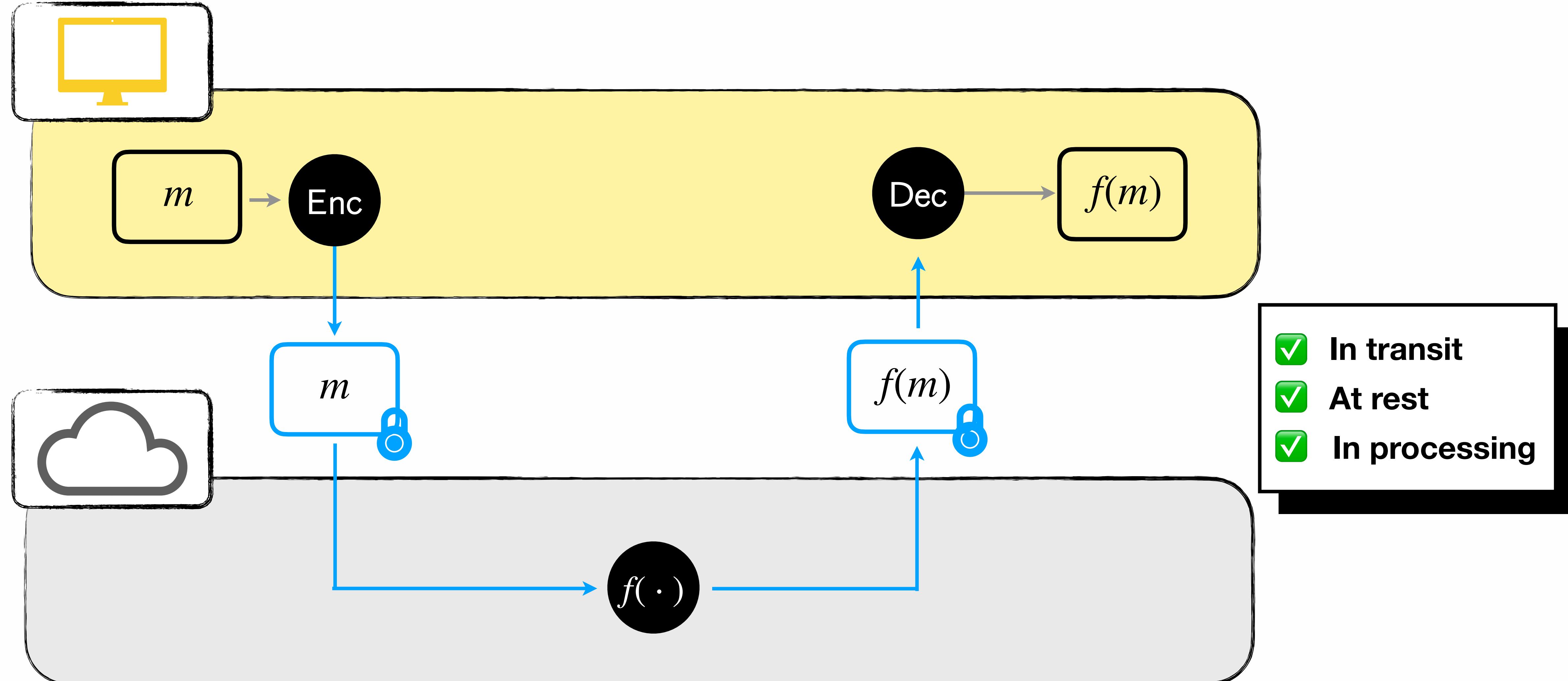
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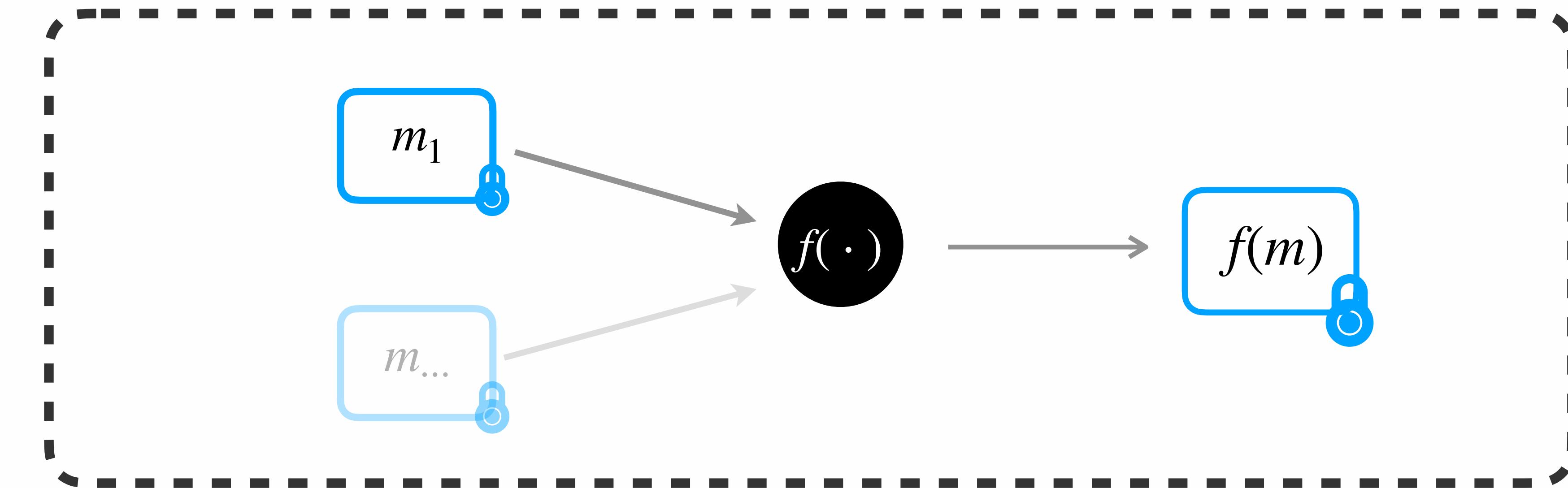
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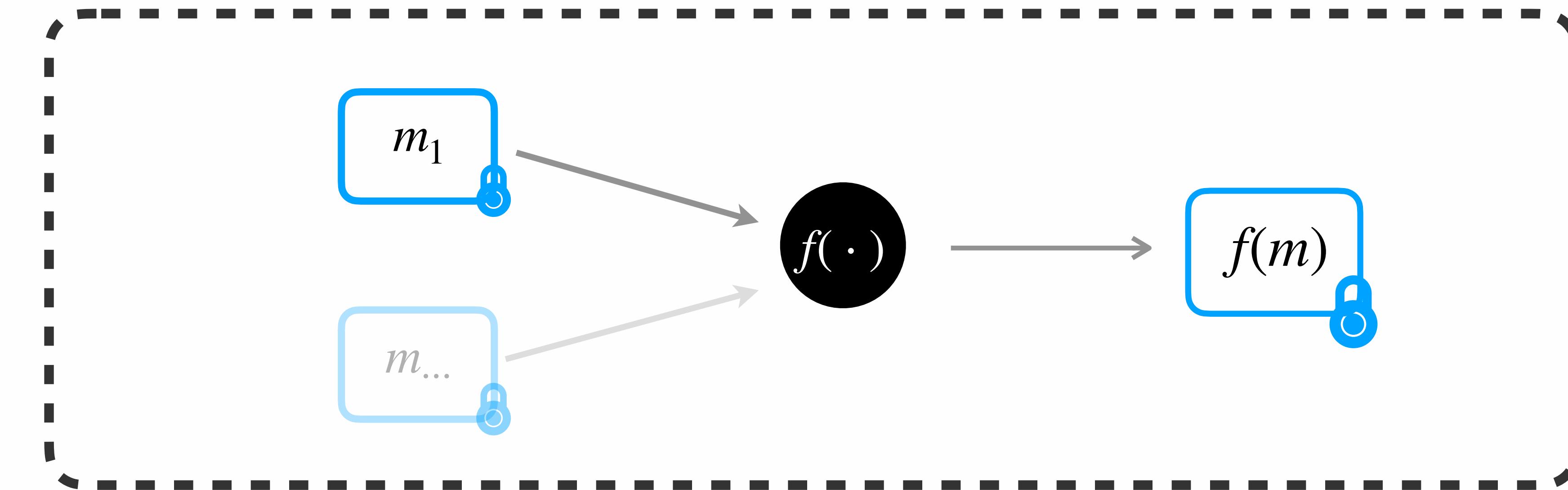
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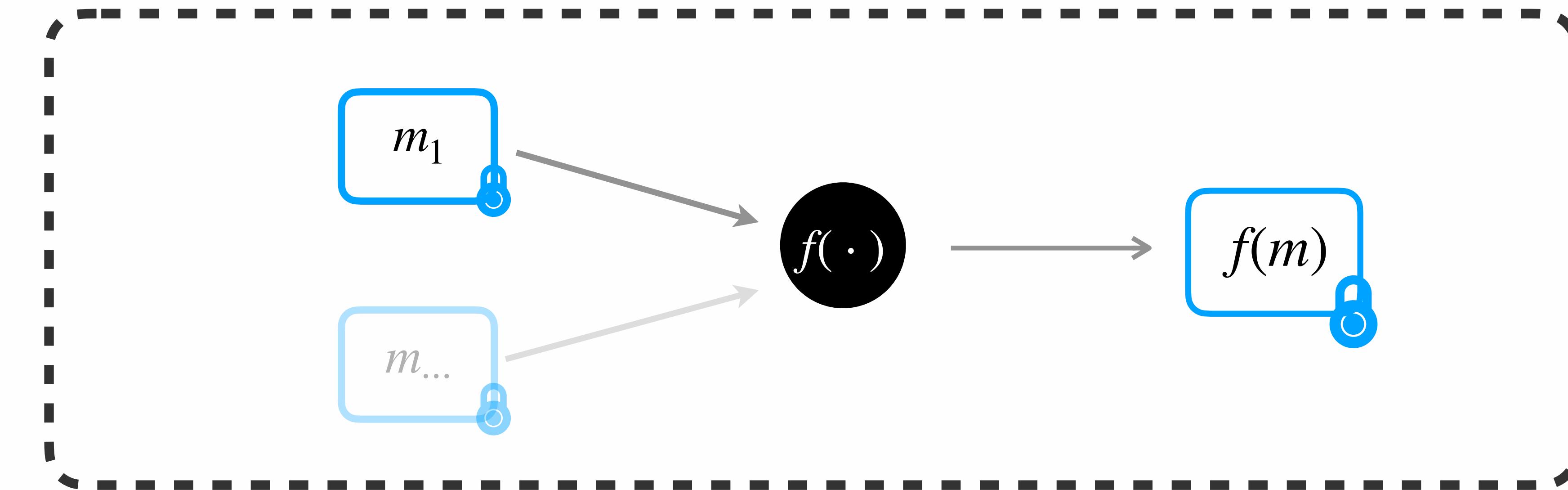
FHE



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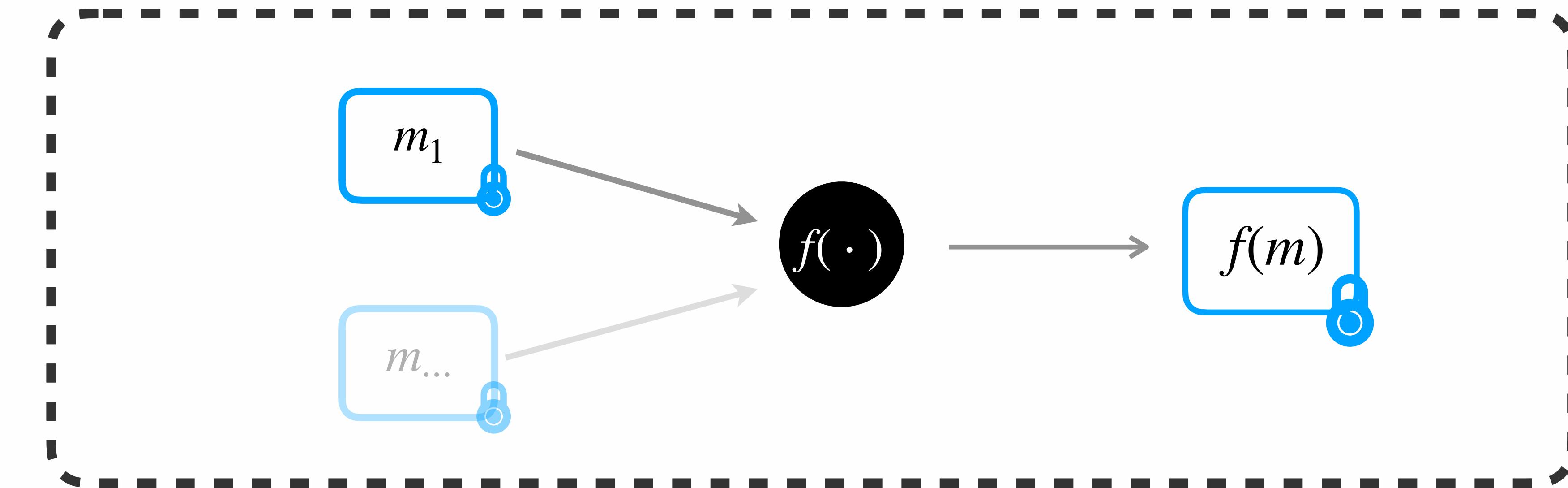


FHE



f , a boolean circuit

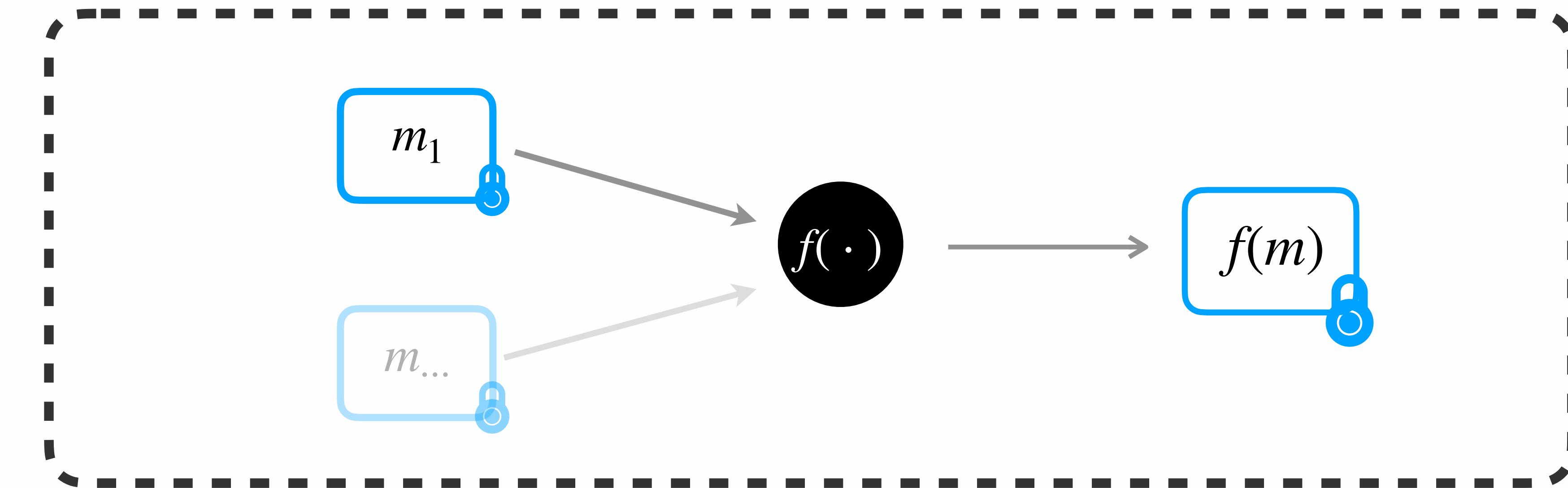
FHE



f , a boolean circuit

f , a neural network

FHE

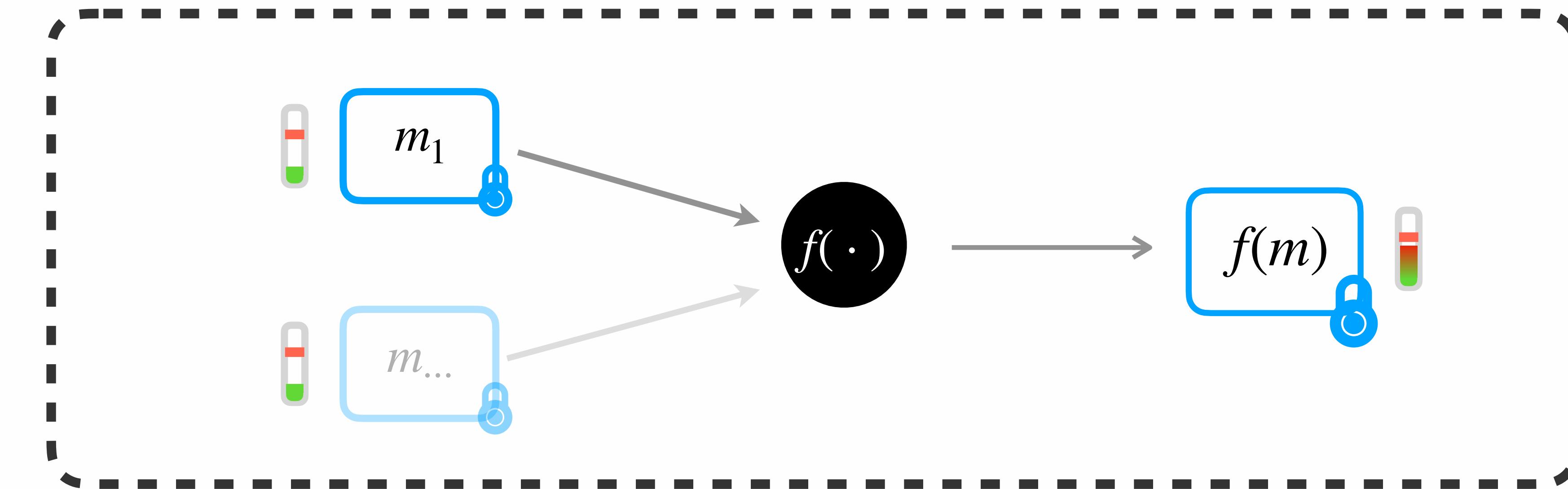


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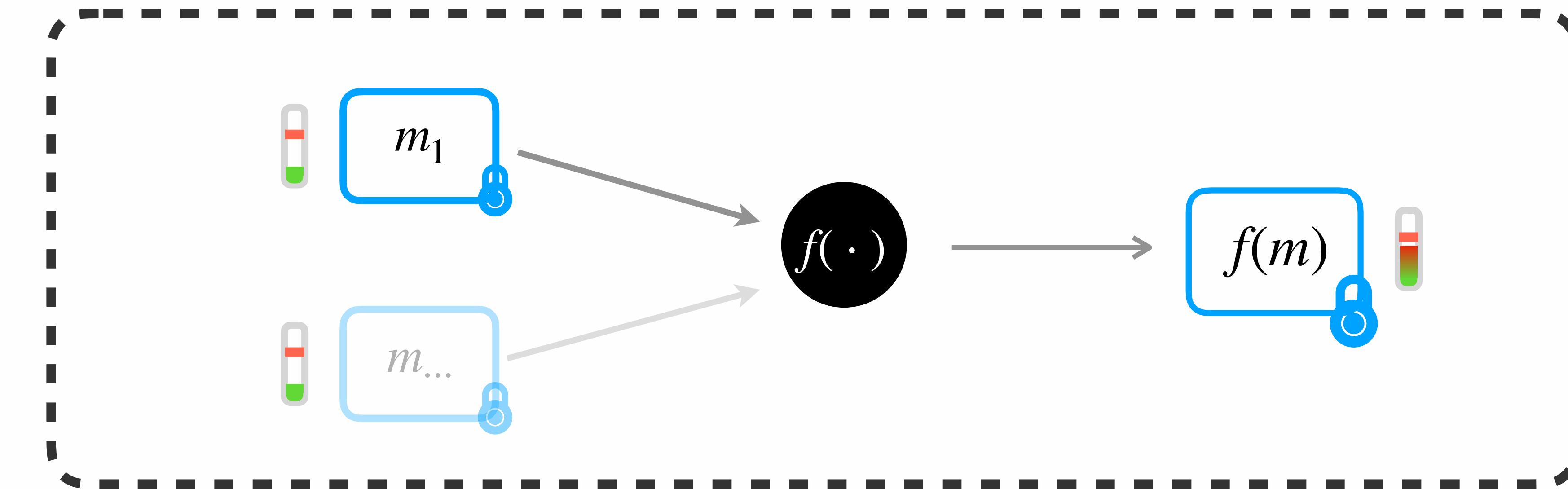
f , a neural network

f , any function

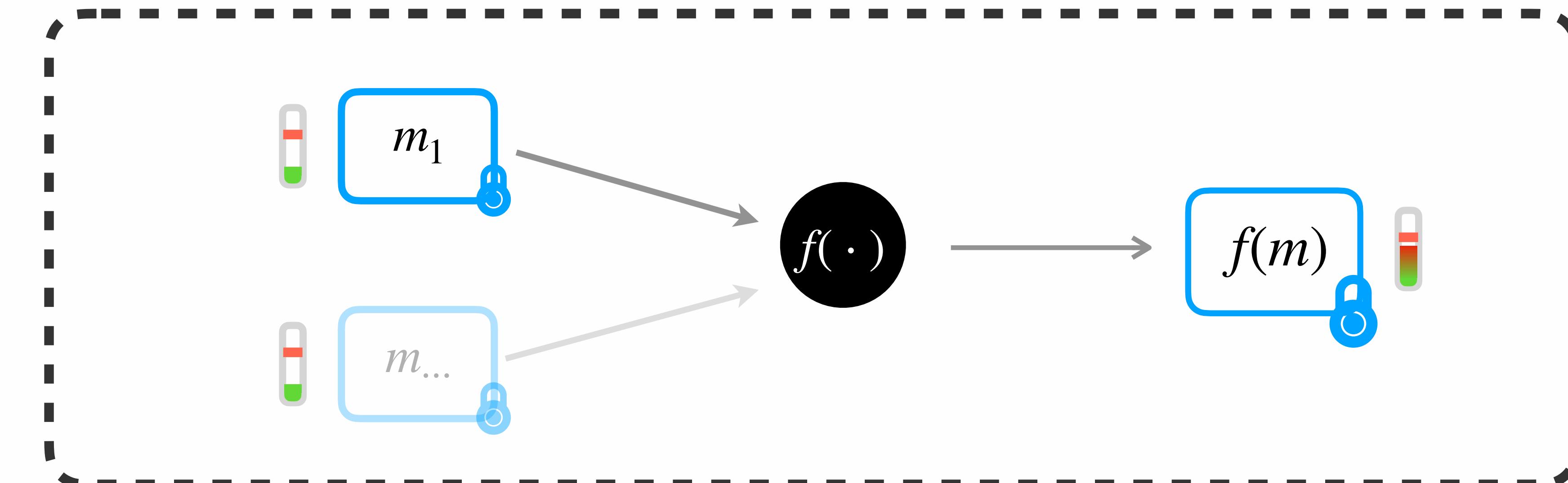
FHE



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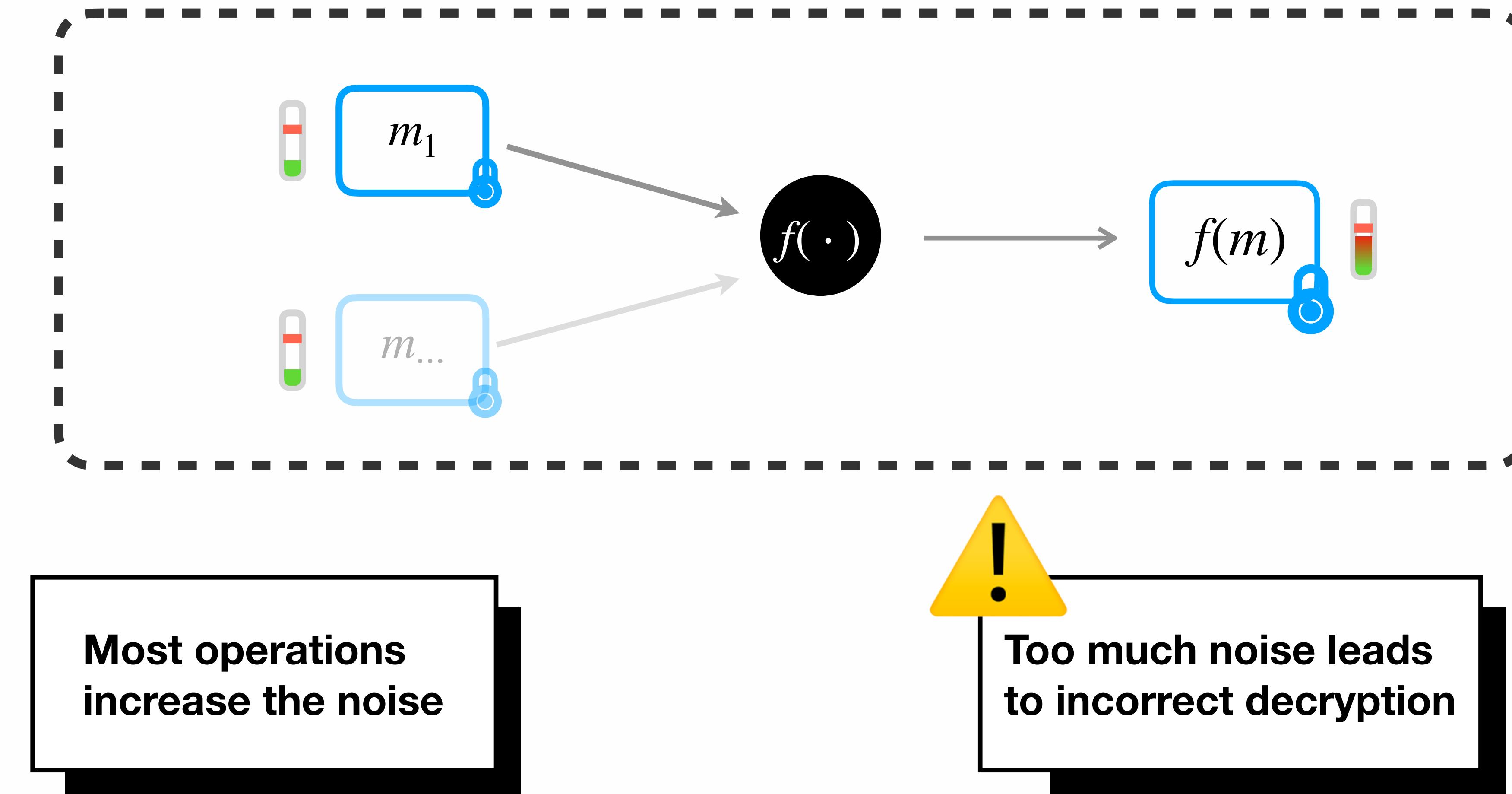


FHE



Most operations
increase the noise

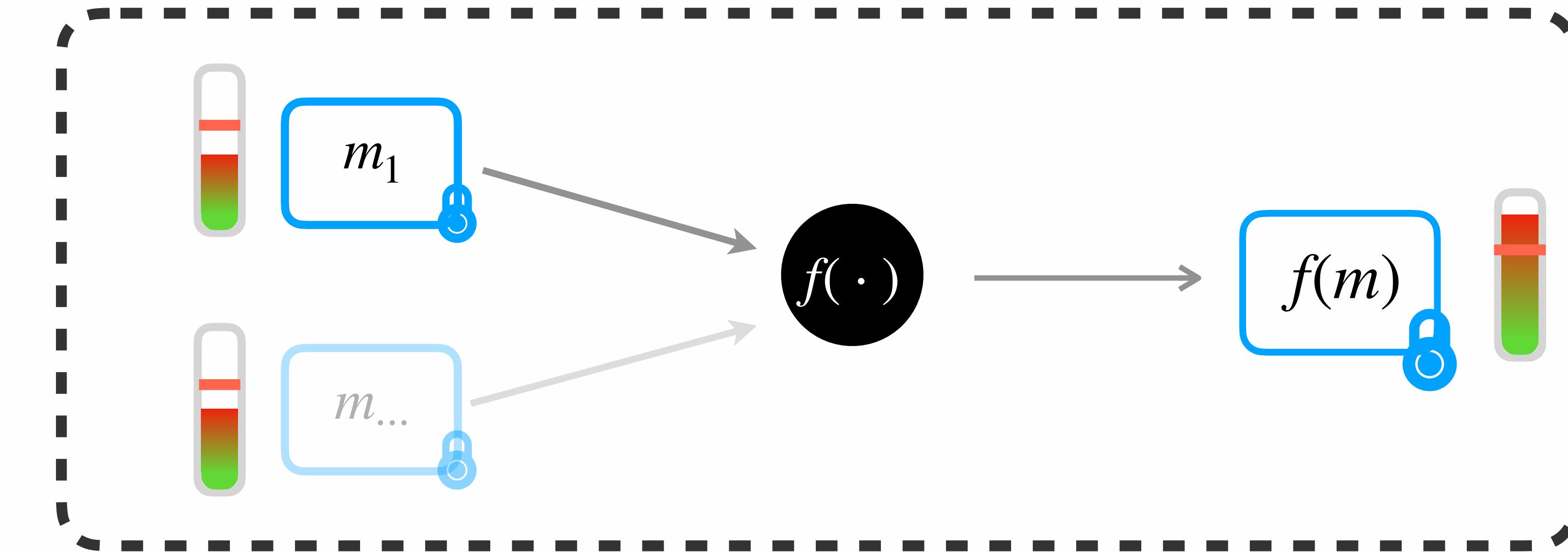
FHE



Leveled Homomorphic Encryption

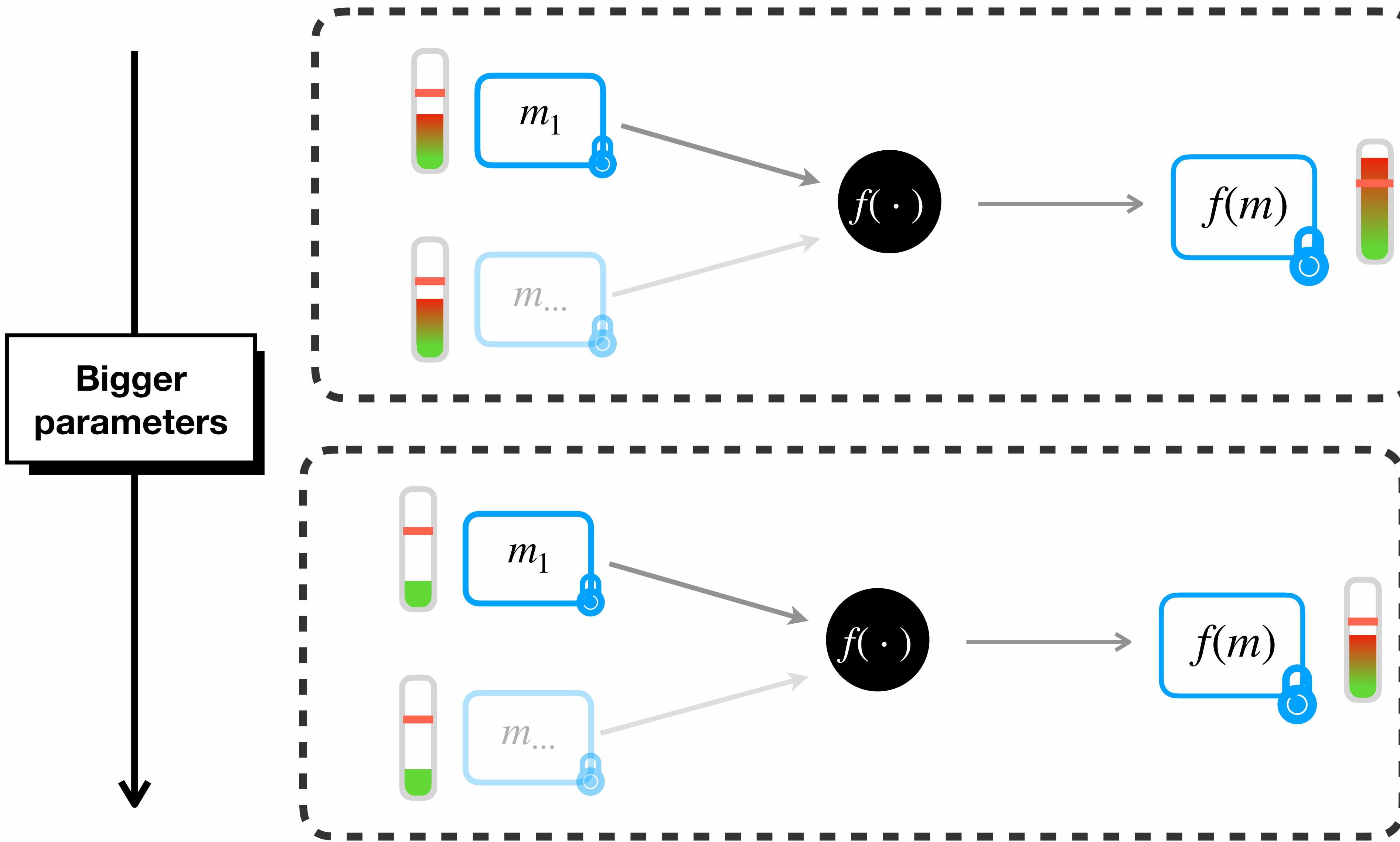
Leveled Homomorphic Encryption

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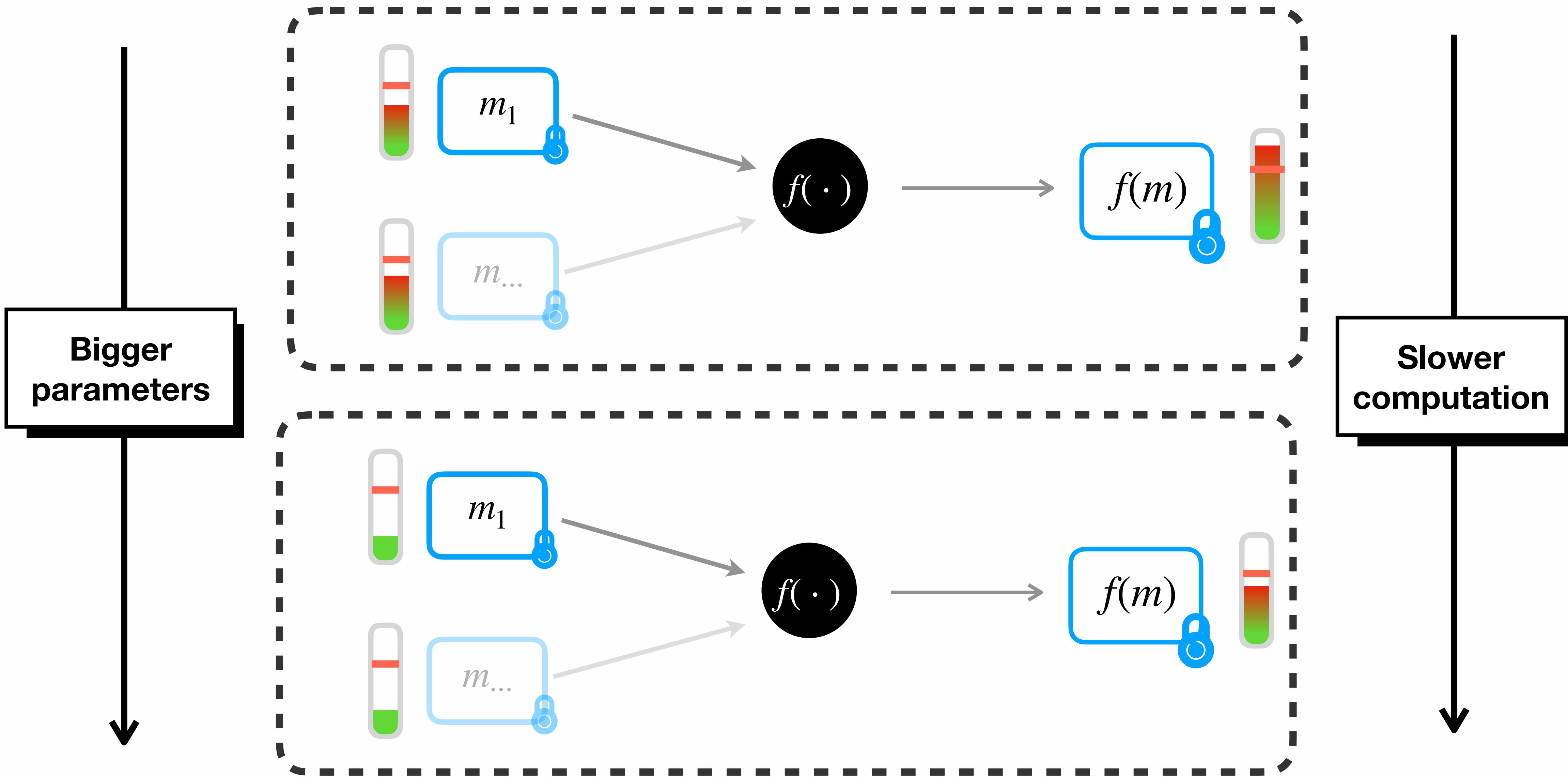
Leveled Homomorphic Encryption

Fully Homomorphic Encryption I



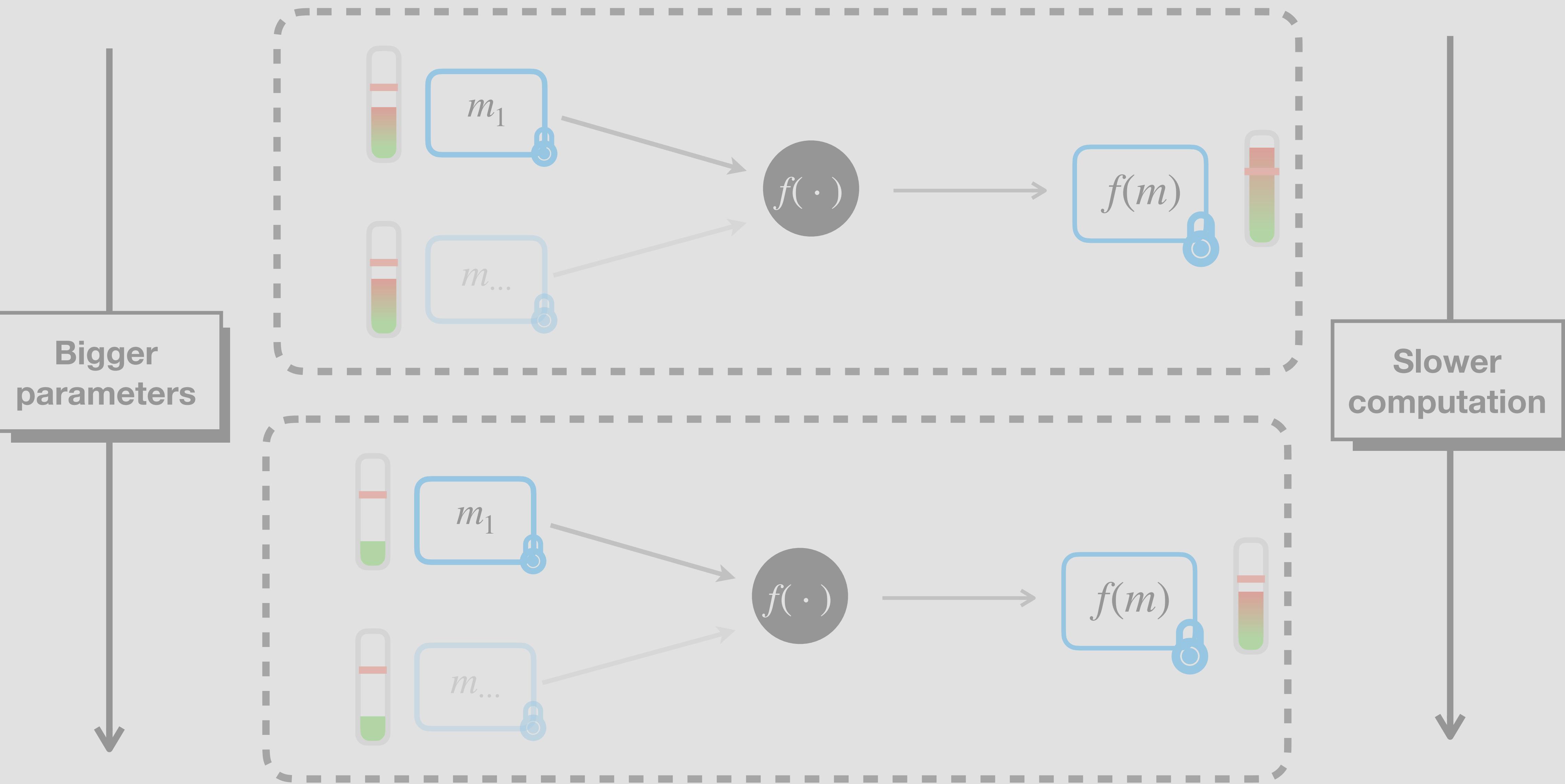
Leveled Homomorphic Encryption

Fully Homomorphic Encryption I



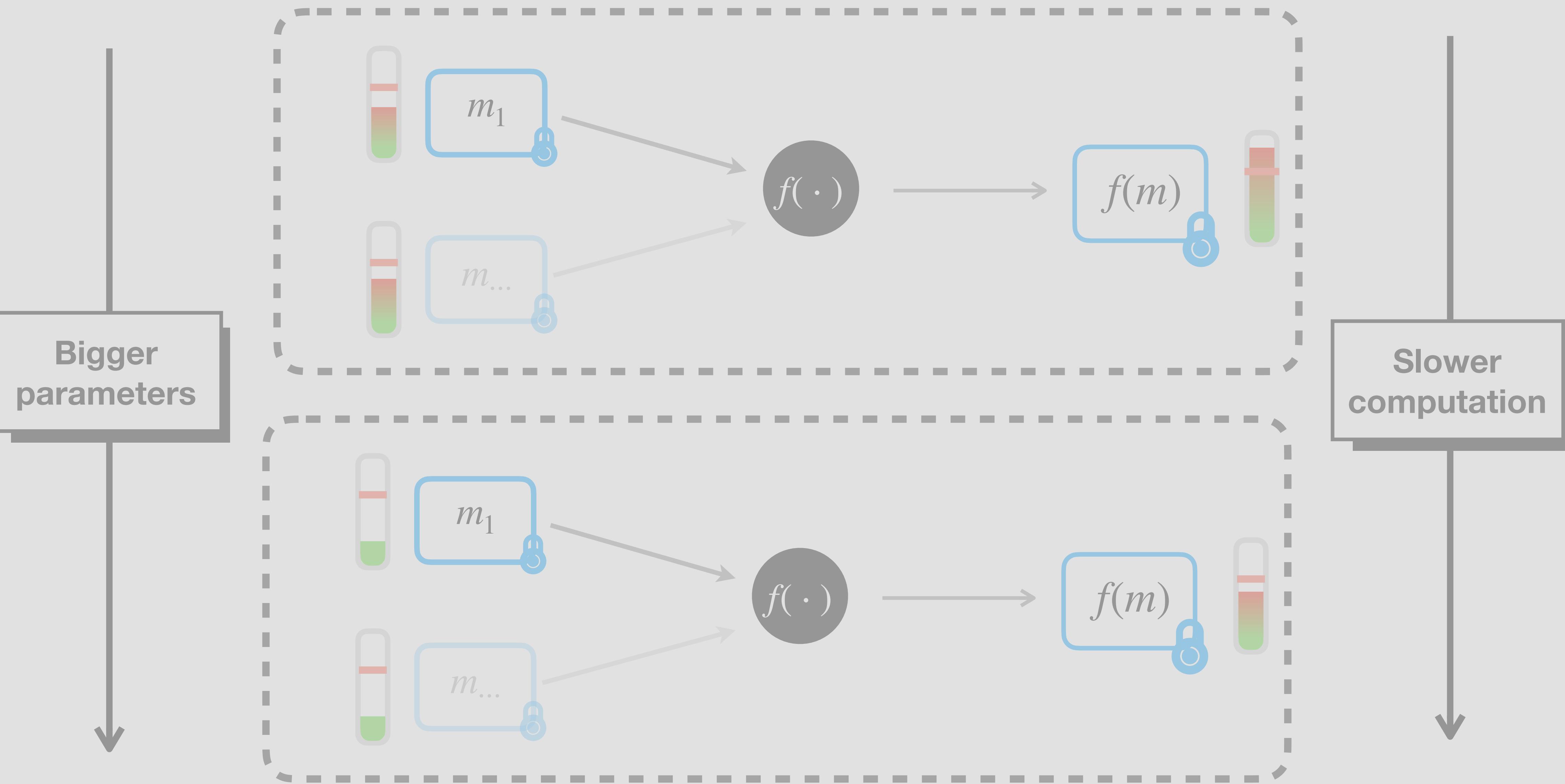
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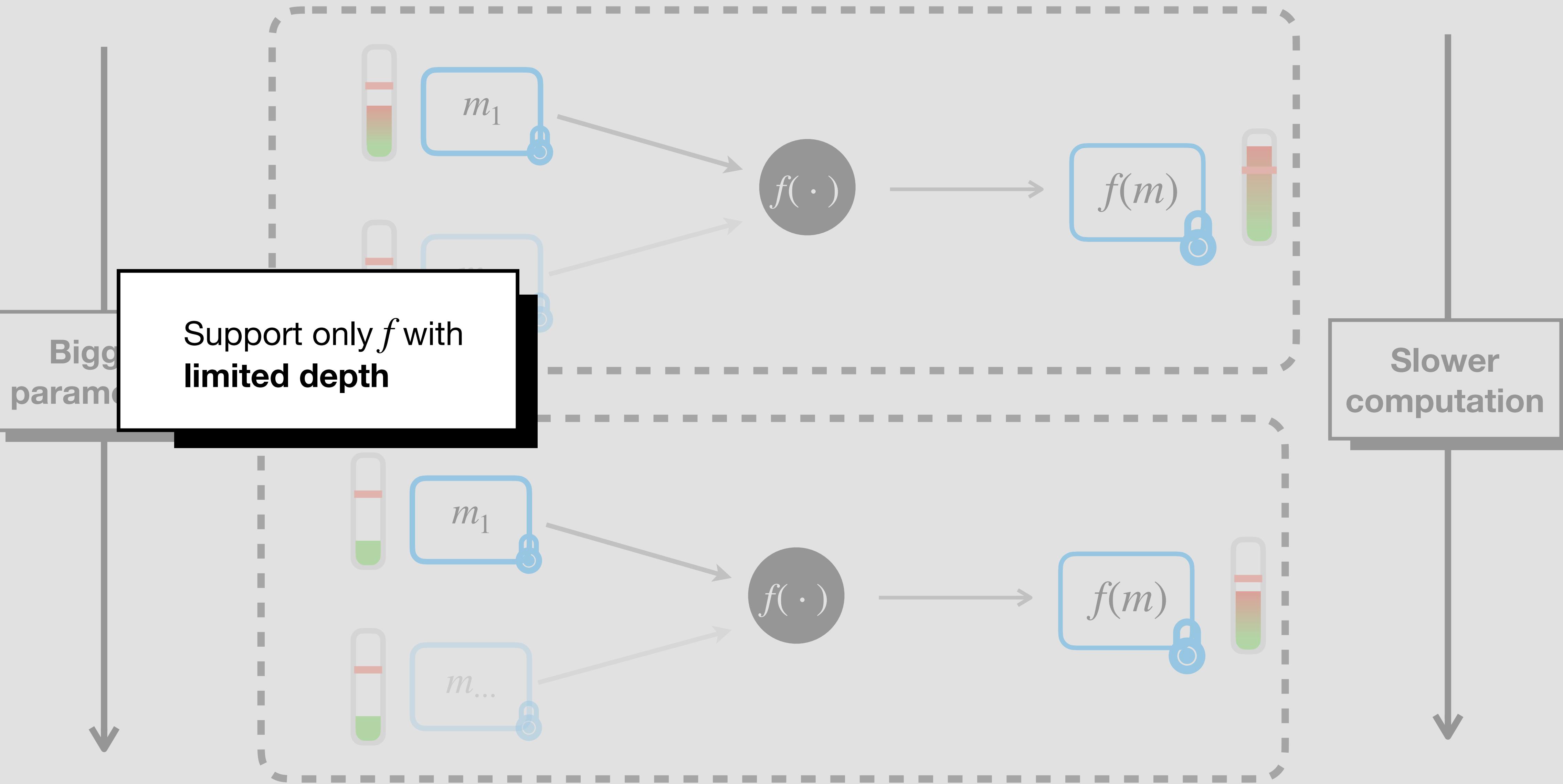
Leveled Homomorphic Encryption

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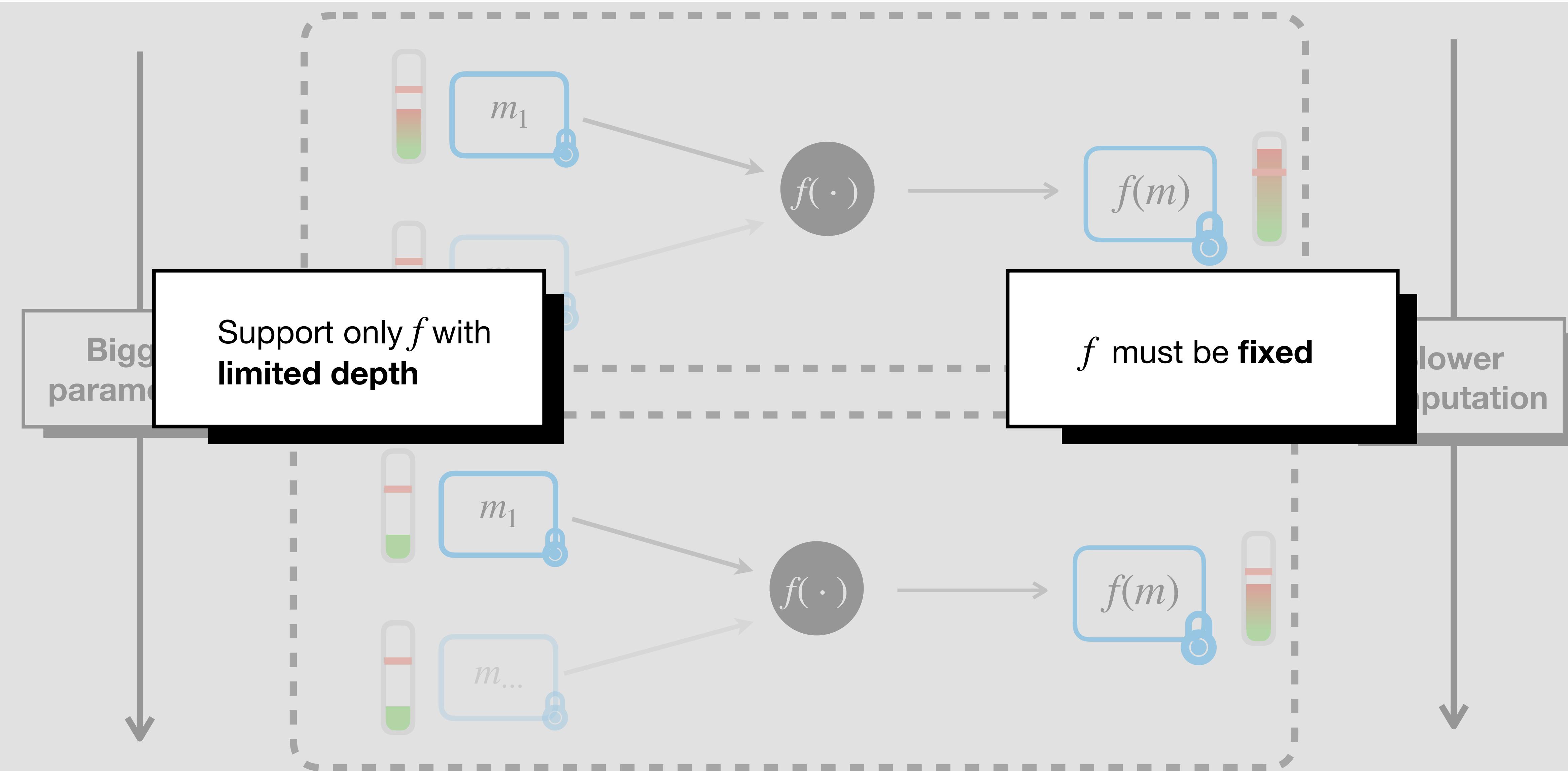
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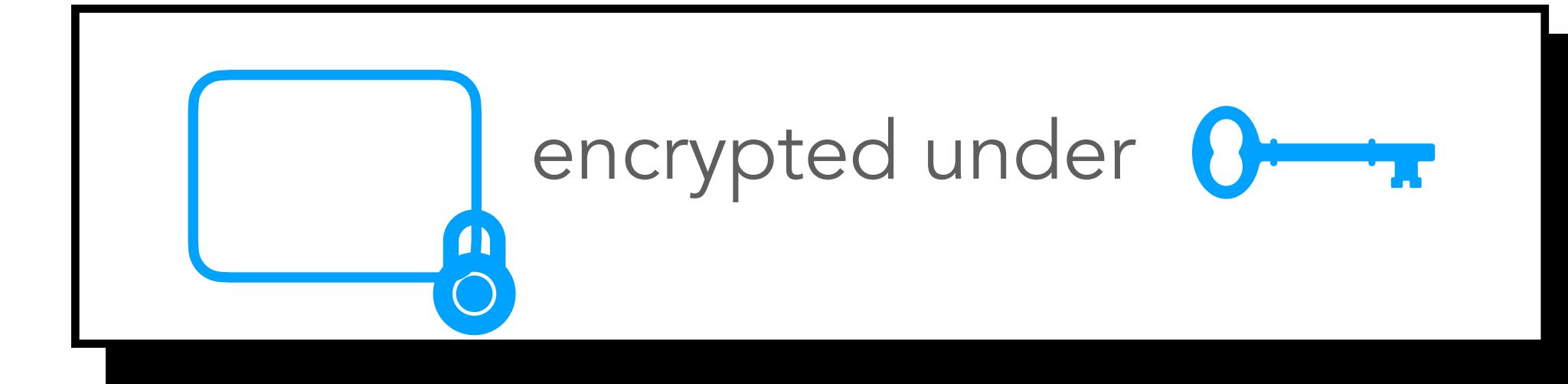
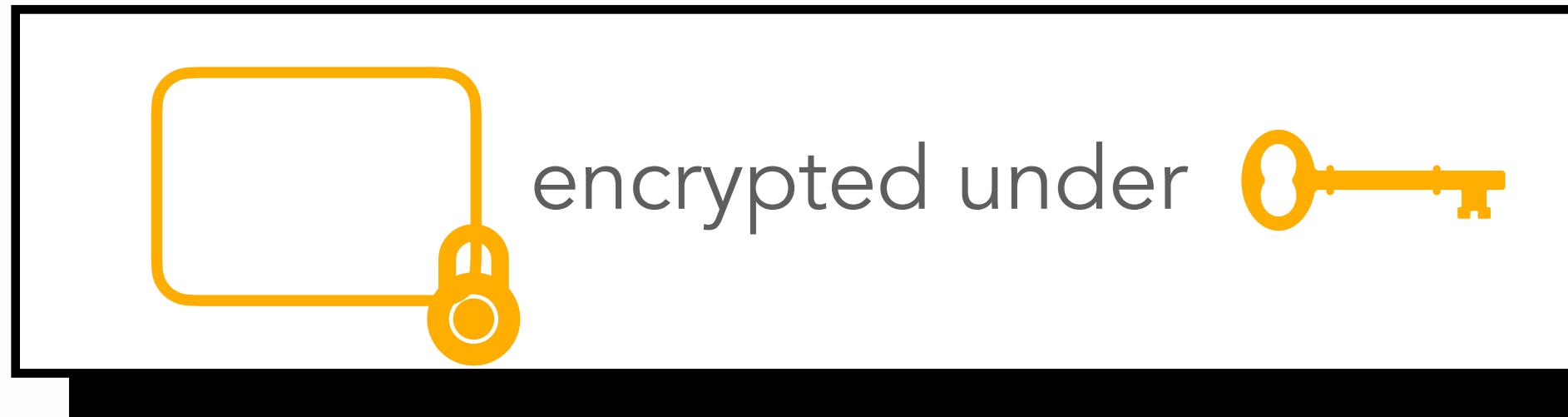
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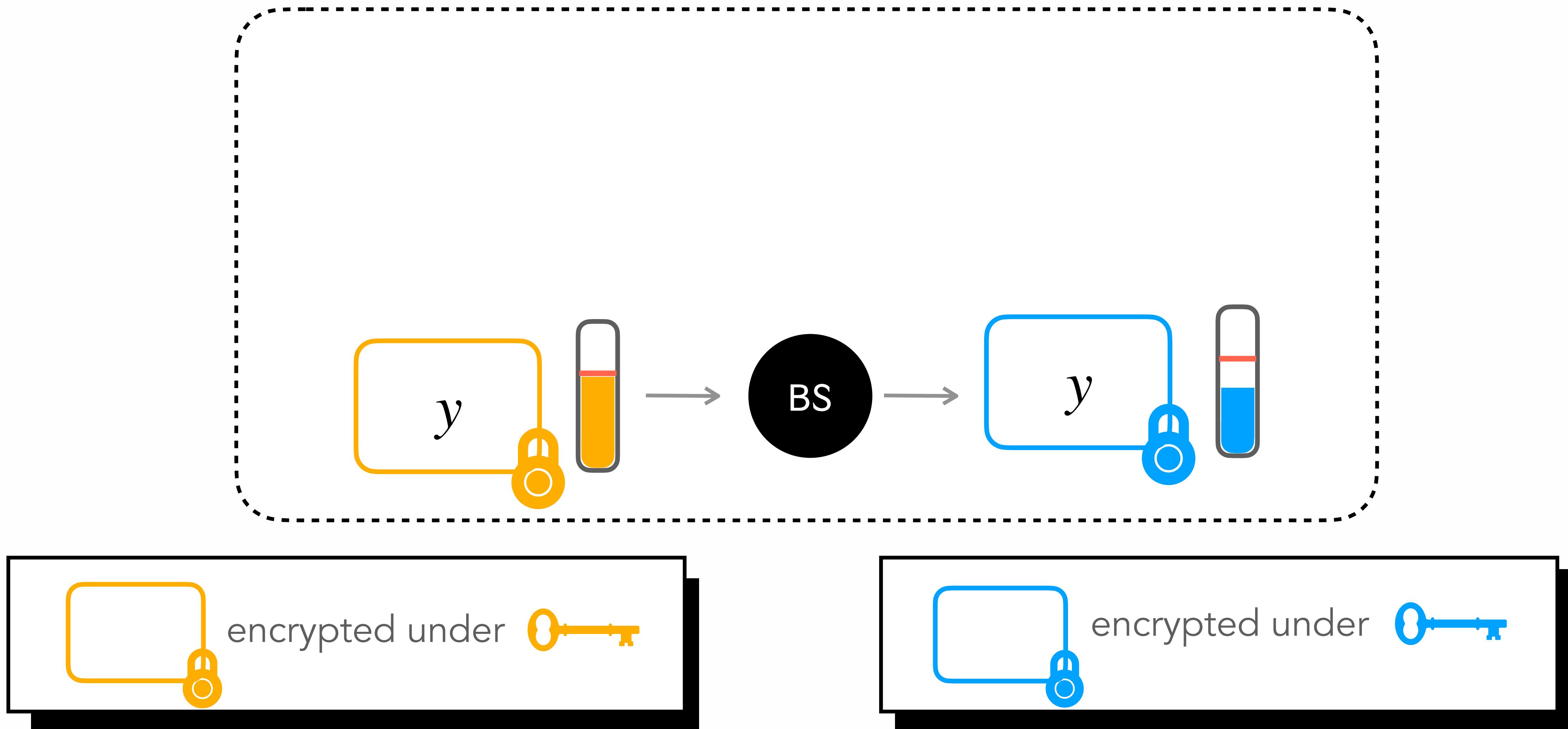
Bootstrapping [Gen09]

Fully Homomorphic Encryption I



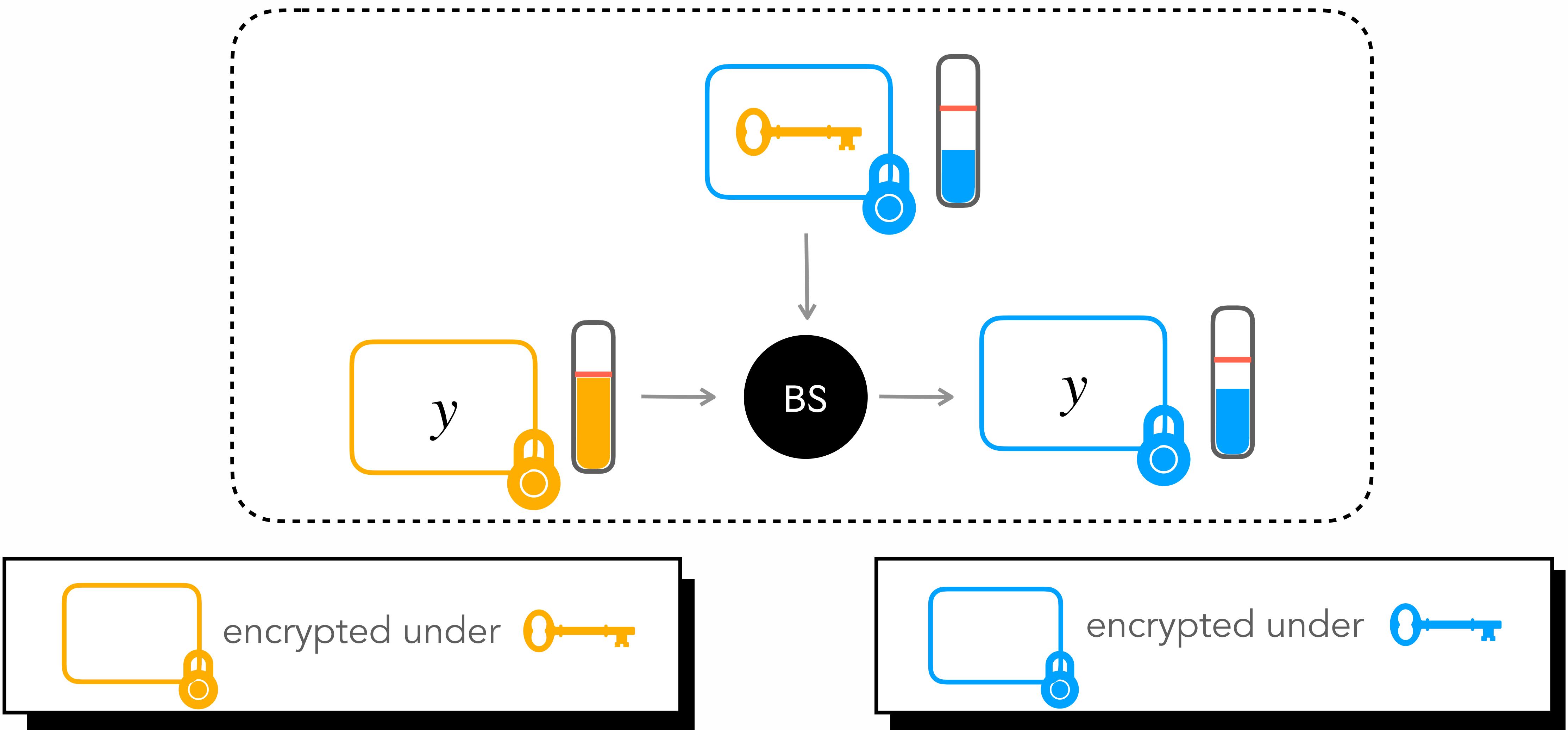
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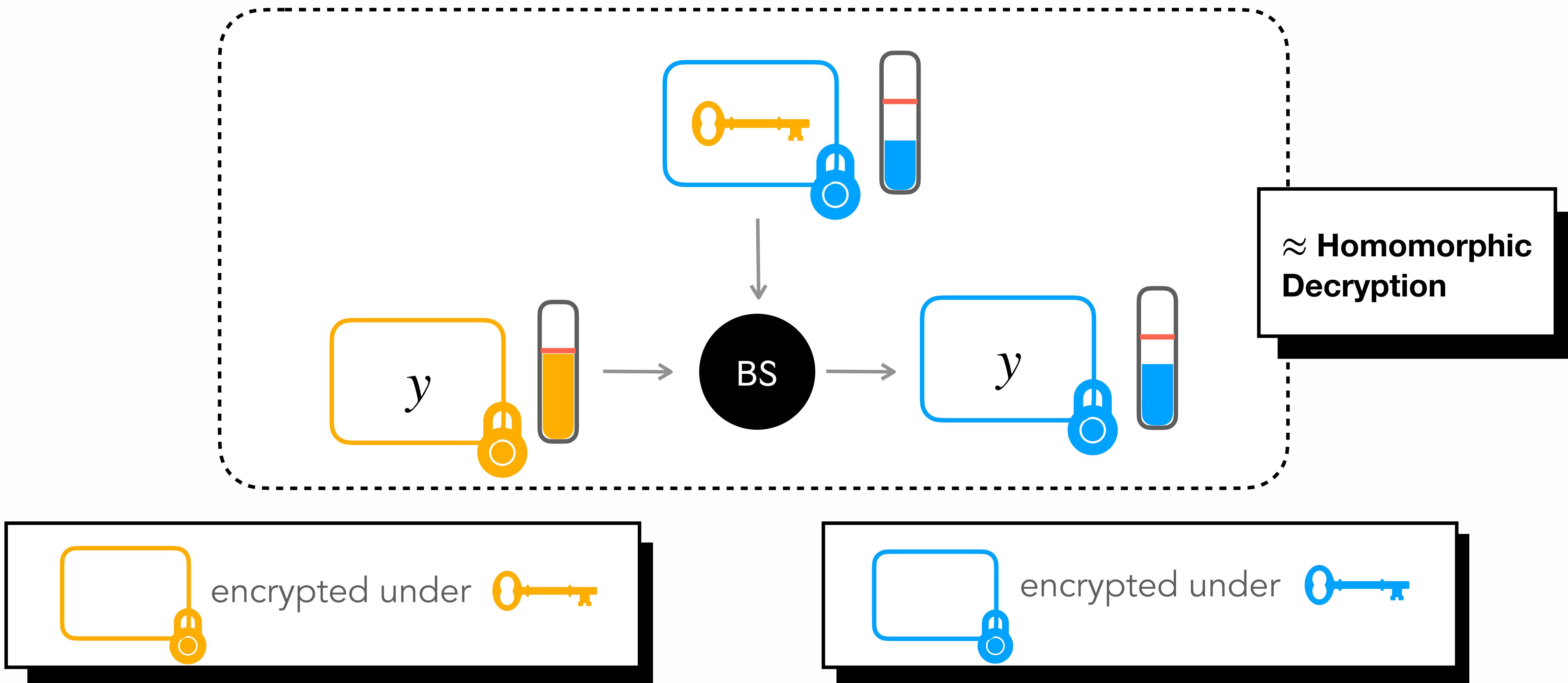
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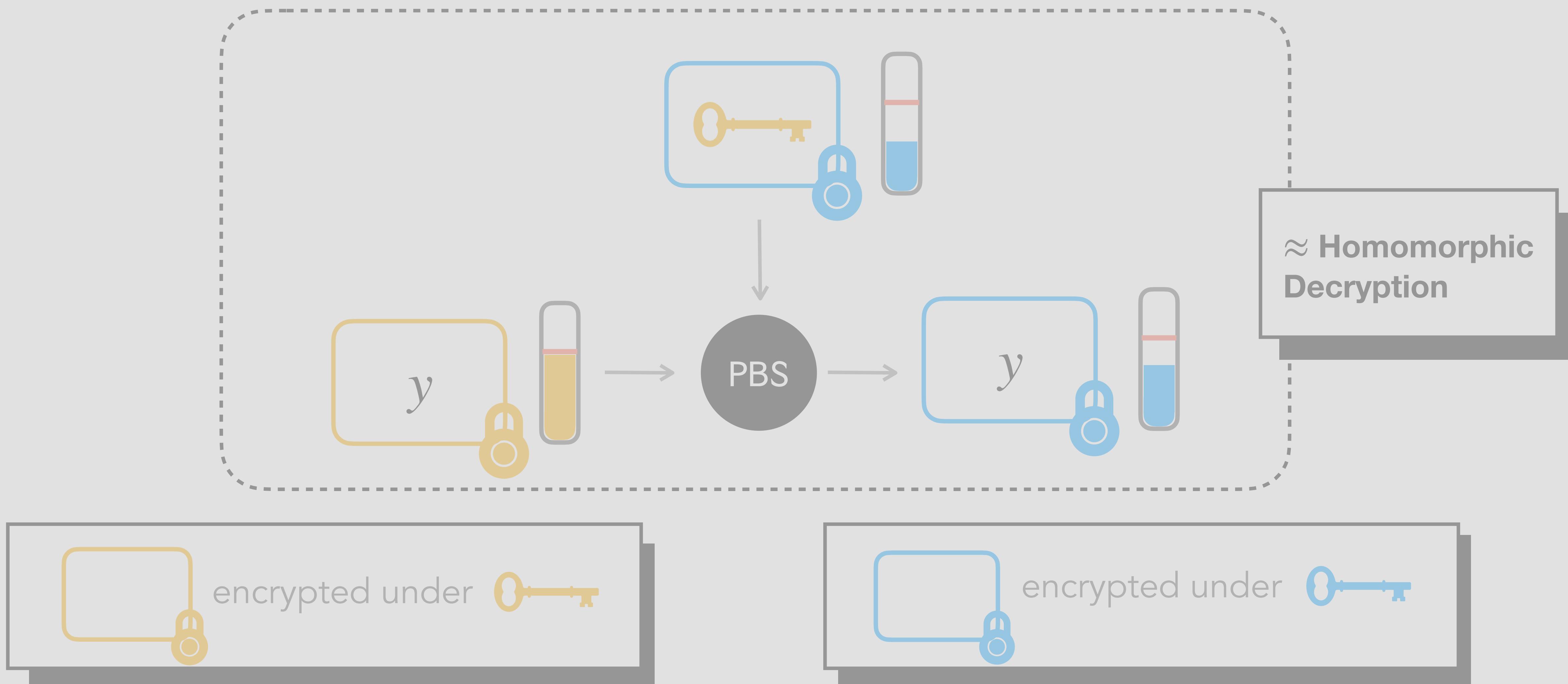
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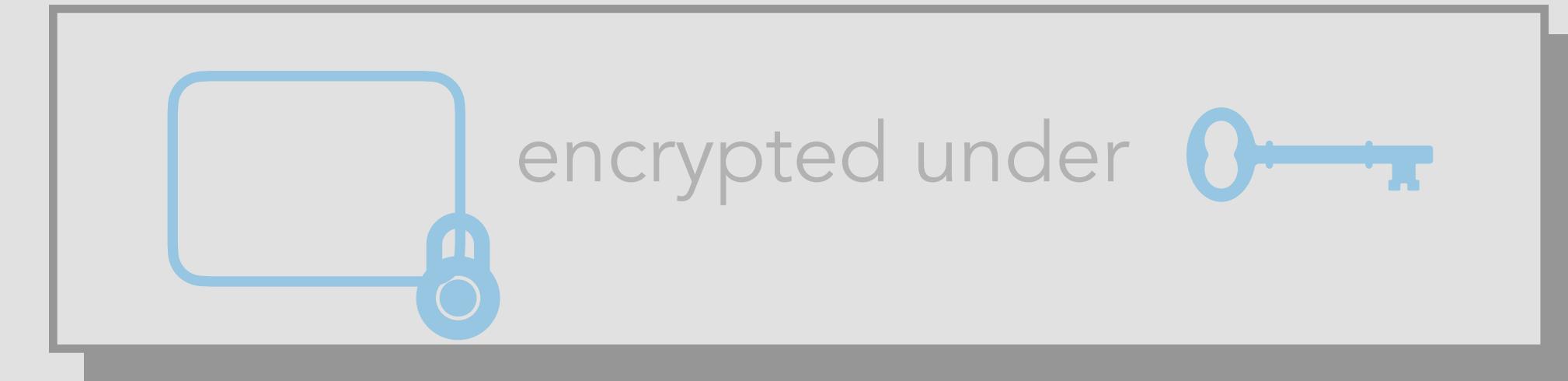
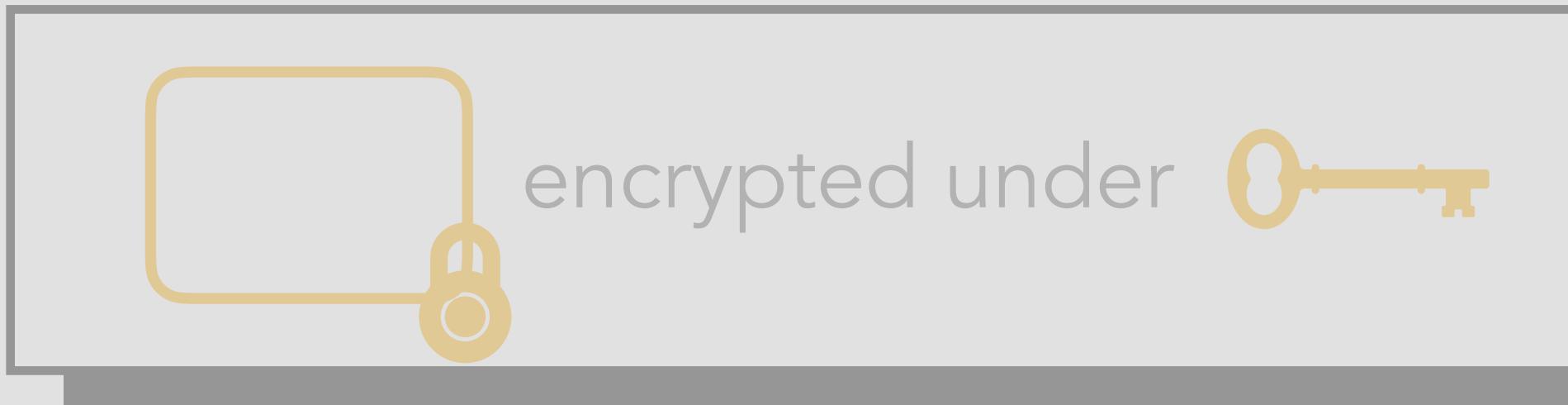
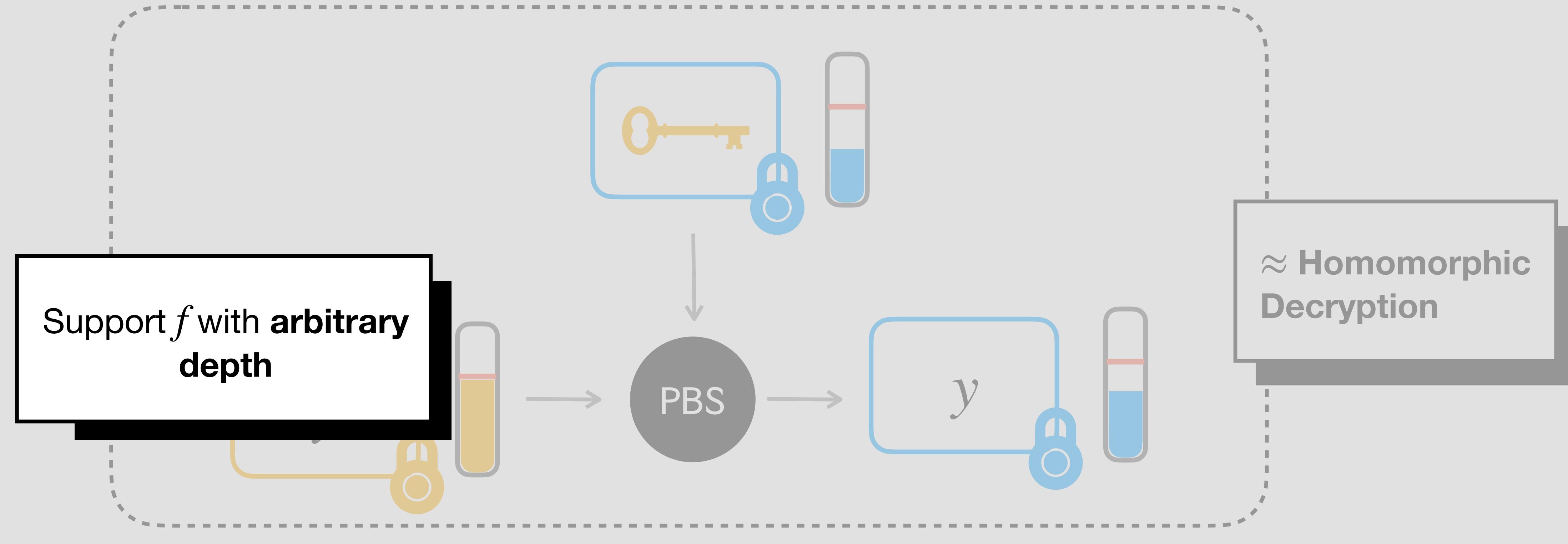
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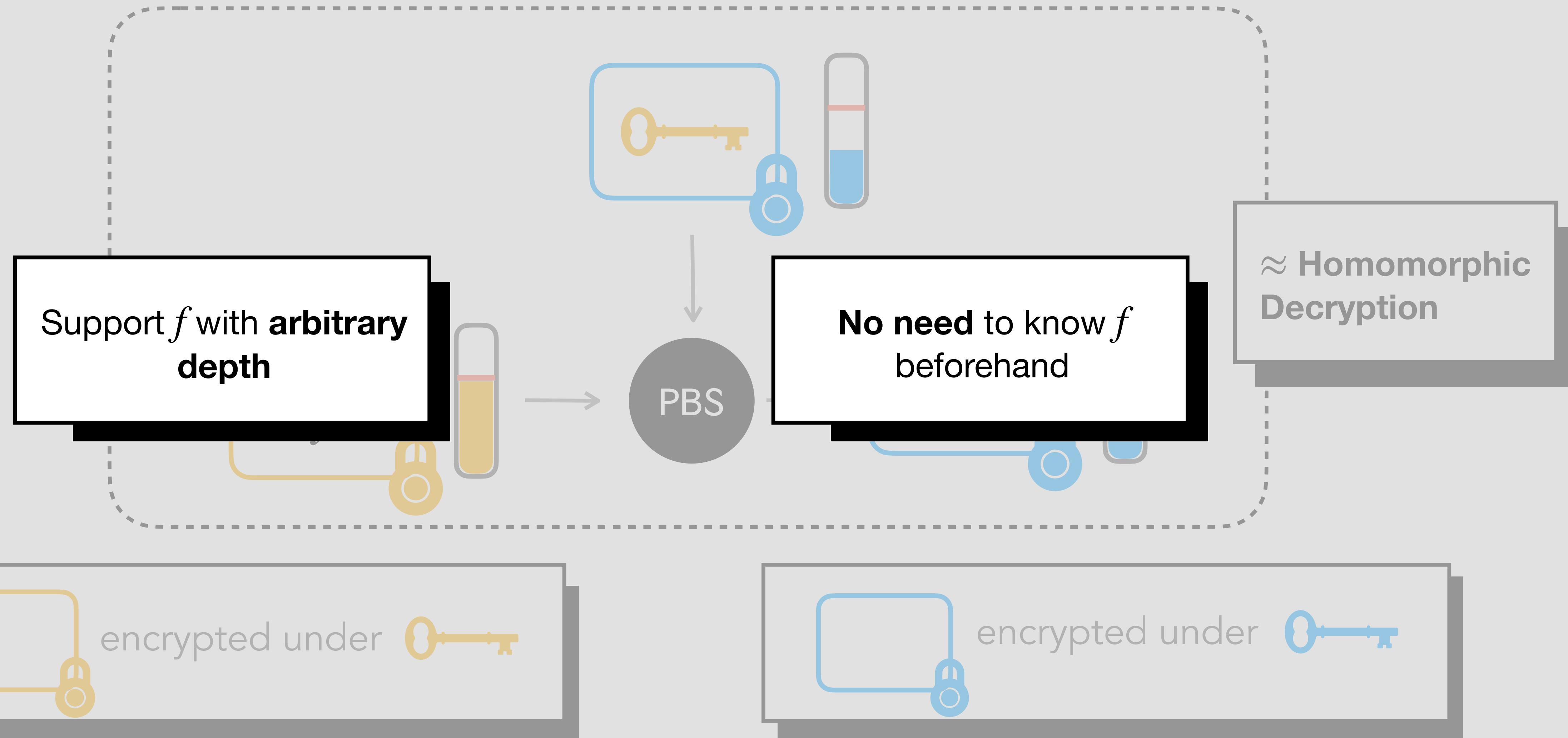
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Fully Homomorphic Encryption I



Bootstrapping [Gen09]

Fully Homomorphic Encryption I



TFHE

Intuition and main algorithms

Encryption

 Δ

MSB

LSB

Encryption

1

Draw a mask $\vec{a} = (a_0, \dots, a_{n-1})$ from $\mathbb{U}(\llbracket 0; q - 1 \rrbracket)$ and
a secret key $\vec{s} = (s_0, \dots, s_{n-1})$ from $\mathbb{U}(\{0,1\})$
a noise e from $\mathcal{N}(0, \sigma^2)$

Δ

MSB

LSB

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- 2 $b = \langle \vec{a}, \vec{s} \rangle + \Delta m + e \pmod{q}$

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$$m_1 = (\vec{a}, b)$$

 Δ

MSB

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$$b = \langle \vec{a}, \vec{s} \rangle + \Delta m + e \pmod{q}$$

LWE Dimension

$$m_1 = (\vec{a}, b)$$

 Δ

MSB

LSB

Encryption

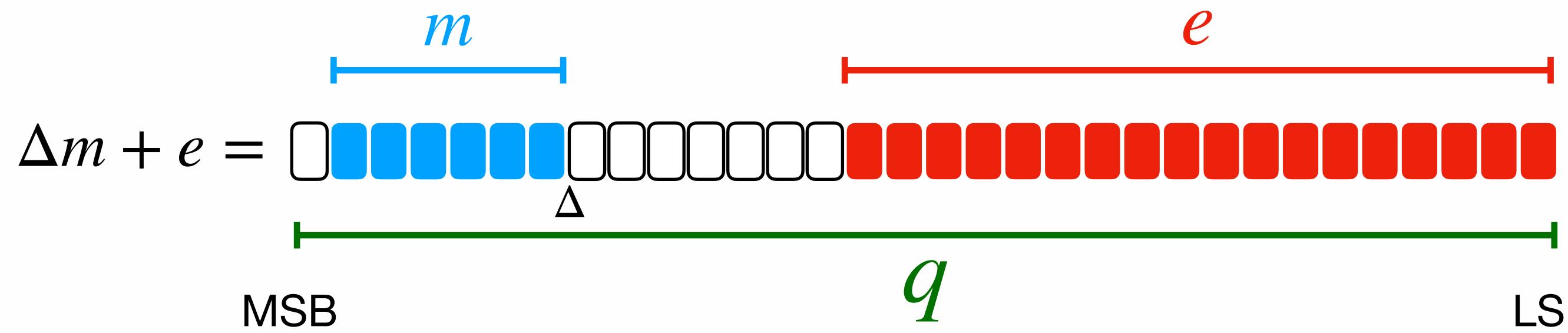
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Encryption

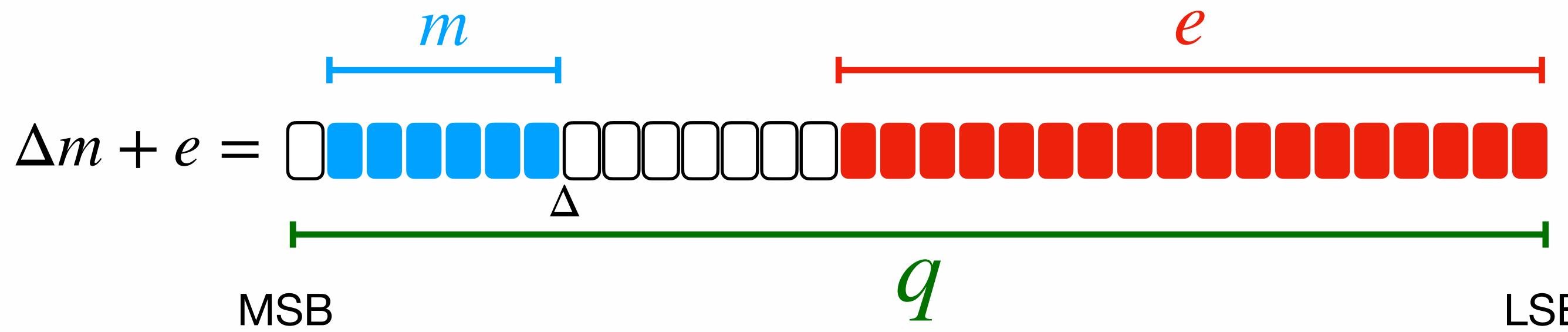
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 a noise e from $\mathcal{N}(0, \sigma^2)$

2

$$b = \langle \vec{a}, \vec{s} \rangle + \Delta m + e \pmod{q}$$

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For a security level λ ,

$$n \nearrow \implies \sigma \searrow$$

Decryption

Decryption

1

$$b - \langle \vec{a}, \vec{s} \rangle = \cancel{\langle \vec{a}, \vec{s} \rangle} + \Delta m + e - \cancel{\langle \vec{a}, \vec{s} \rangle} \pmod{q}$$

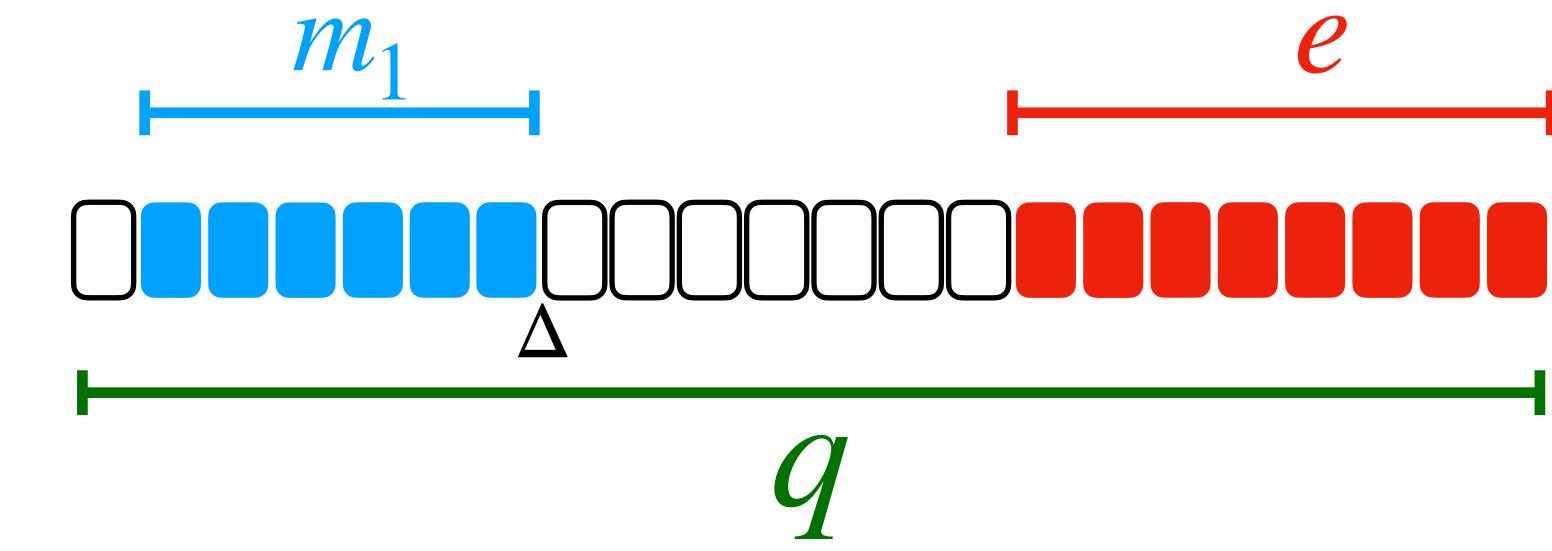
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2a

$$|e| < \frac{\Delta}{2} \implies \left\lceil \frac{\Delta m + e}{\Delta} \right\rceil = m$$



Decryption

1

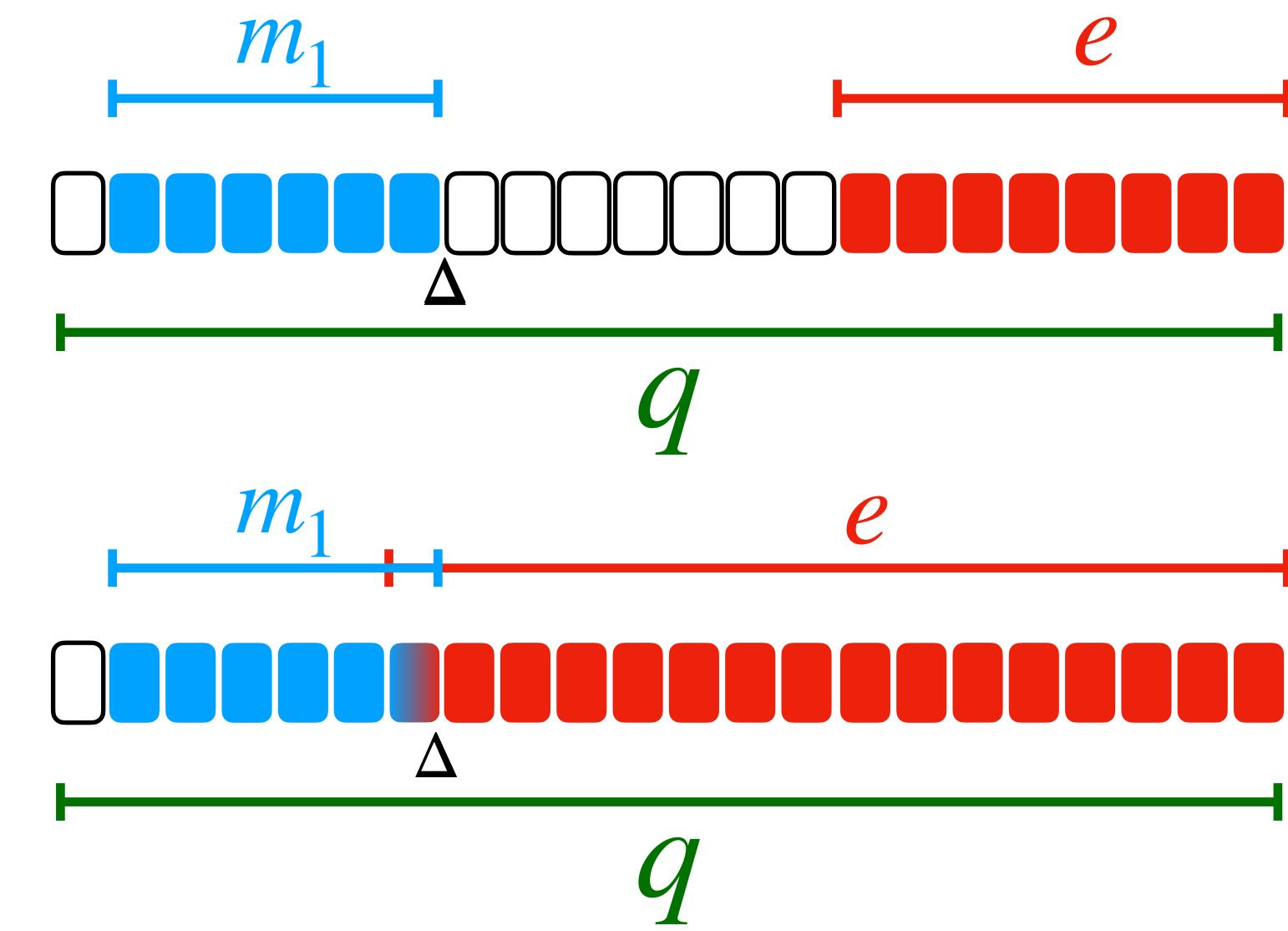
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$$|e| < \frac{\Delta}{2} \implies \left\lceil \frac{\Delta m + e}{\Delta} \right\rceil = m$$

2b

$$|e| \geq \frac{\Delta}{2} \implies \left\lceil \frac{\Delta m + e}{\Delta} \right\rceil \neq m$$



Condition to guarantee
the correctness of a
computation

Decryption

1

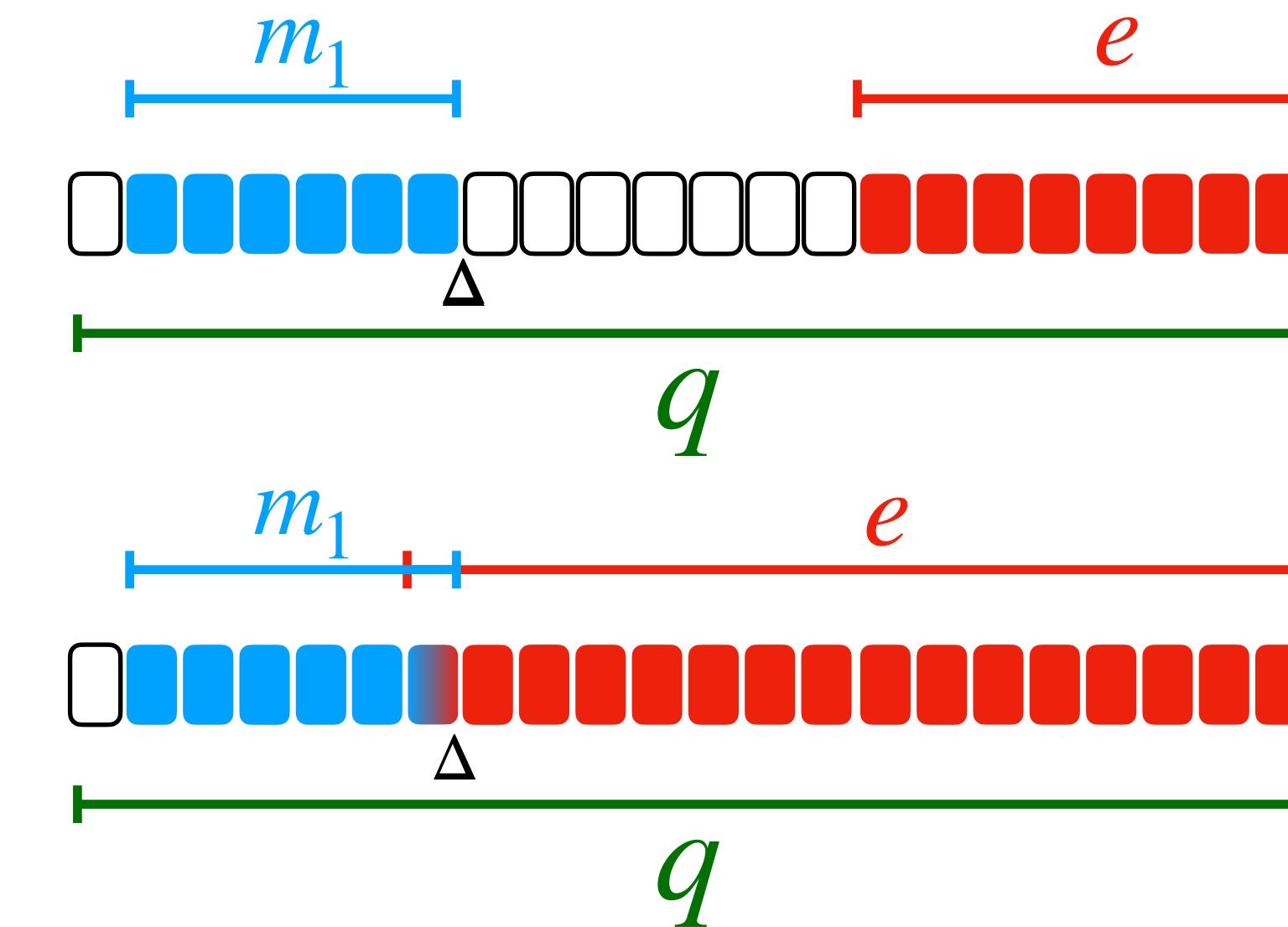
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Addition

Consider $\text{ct}_1 = (\vec{a}_1, b_1)$ and $\text{ct}_2 = (\vec{a}_2, b_2)$

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Define ct_3 such that

$$\text{ct}_3 \leftarrow \text{ct}_1 \oplus \text{ct}_2 = \left(\vec{a}_1 + \vec{a}_2 \pmod{q}, b_1 + b_2 \pmod{q} \right)$$

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Decrypt ct_3

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Decrypt ct_3

$$b_1 + b_2 - \langle \vec{a}_1 + \vec{a}_2, \vec{s} \rangle \pmod{q}$$

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$$\begin{aligned} & b_1 + b_2 - \langle \vec{a}_1 + \vec{a}_2, \vec{s} \rangle \pmod{q} \\ &= \underbrace{\langle \vec{a}_1 + \vec{a}_2, \vec{s} \rangle}_{\text{message}} + \underbrace{m_1 + m_2}_{\text{message}} + e_1 + e_2 - \underbrace{\langle \vec{a}_1 + \vec{a}_2, \vec{s} \rangle}_{\text{noise}} \pmod{q} \end{aligned}$$

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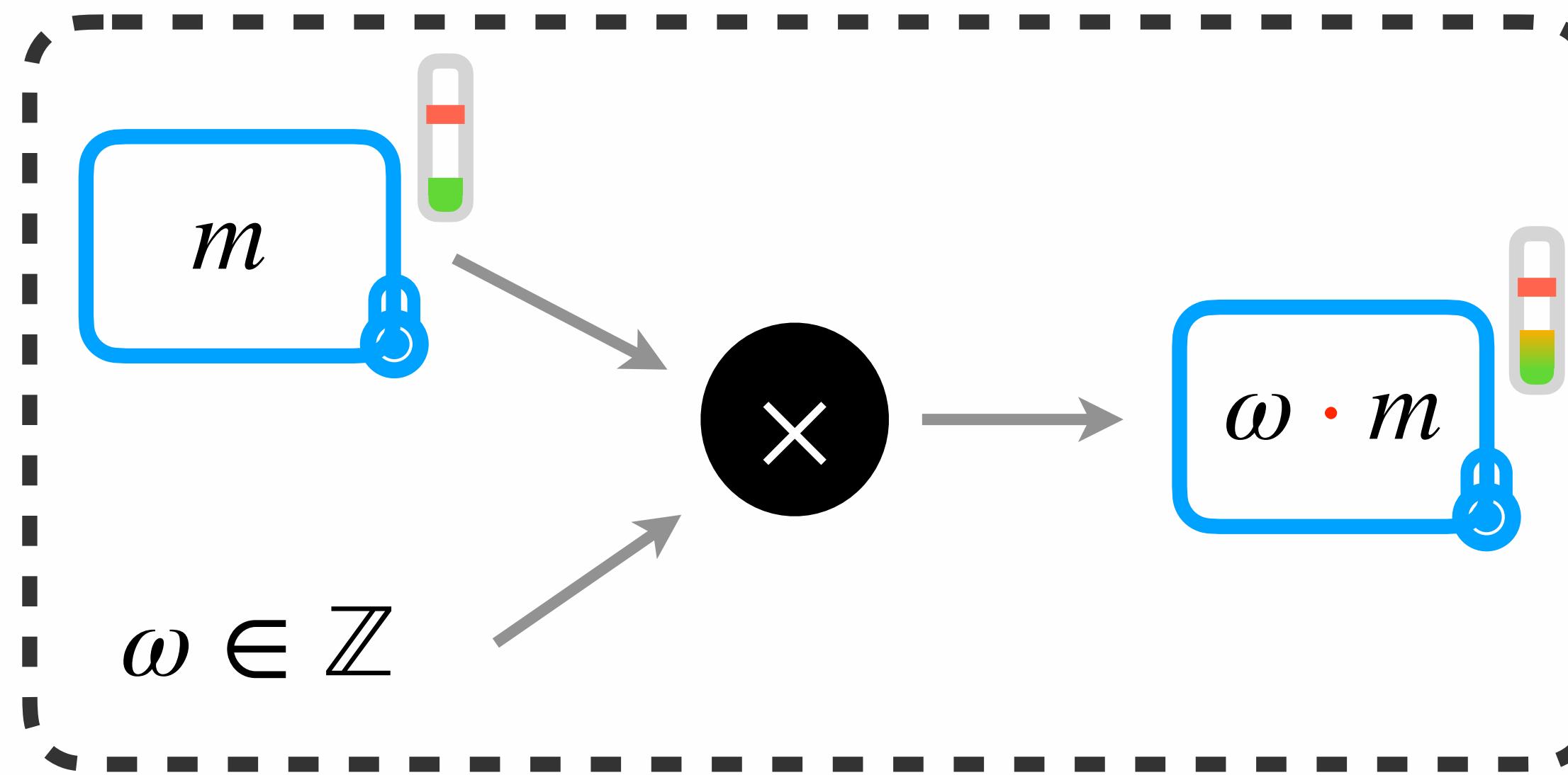
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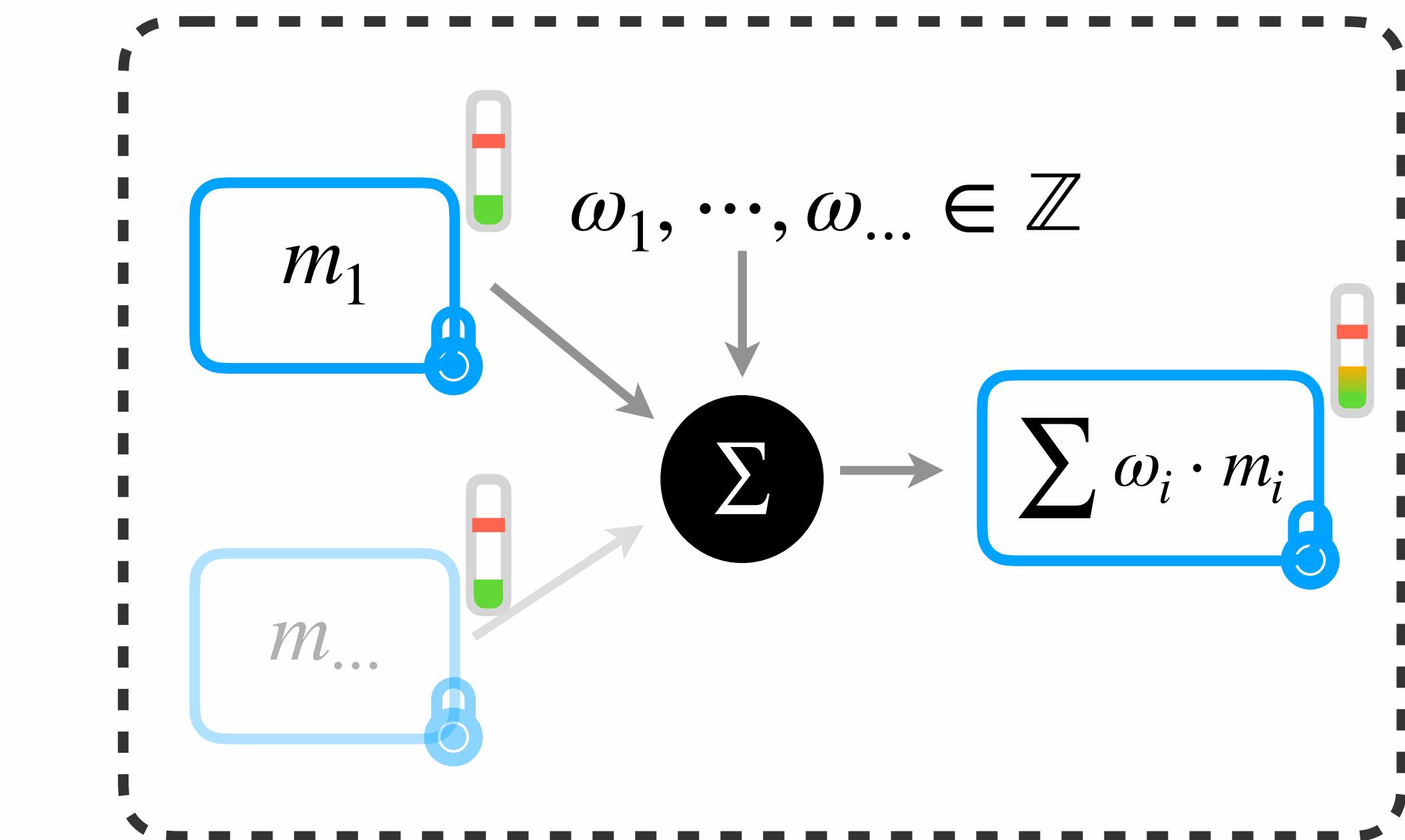
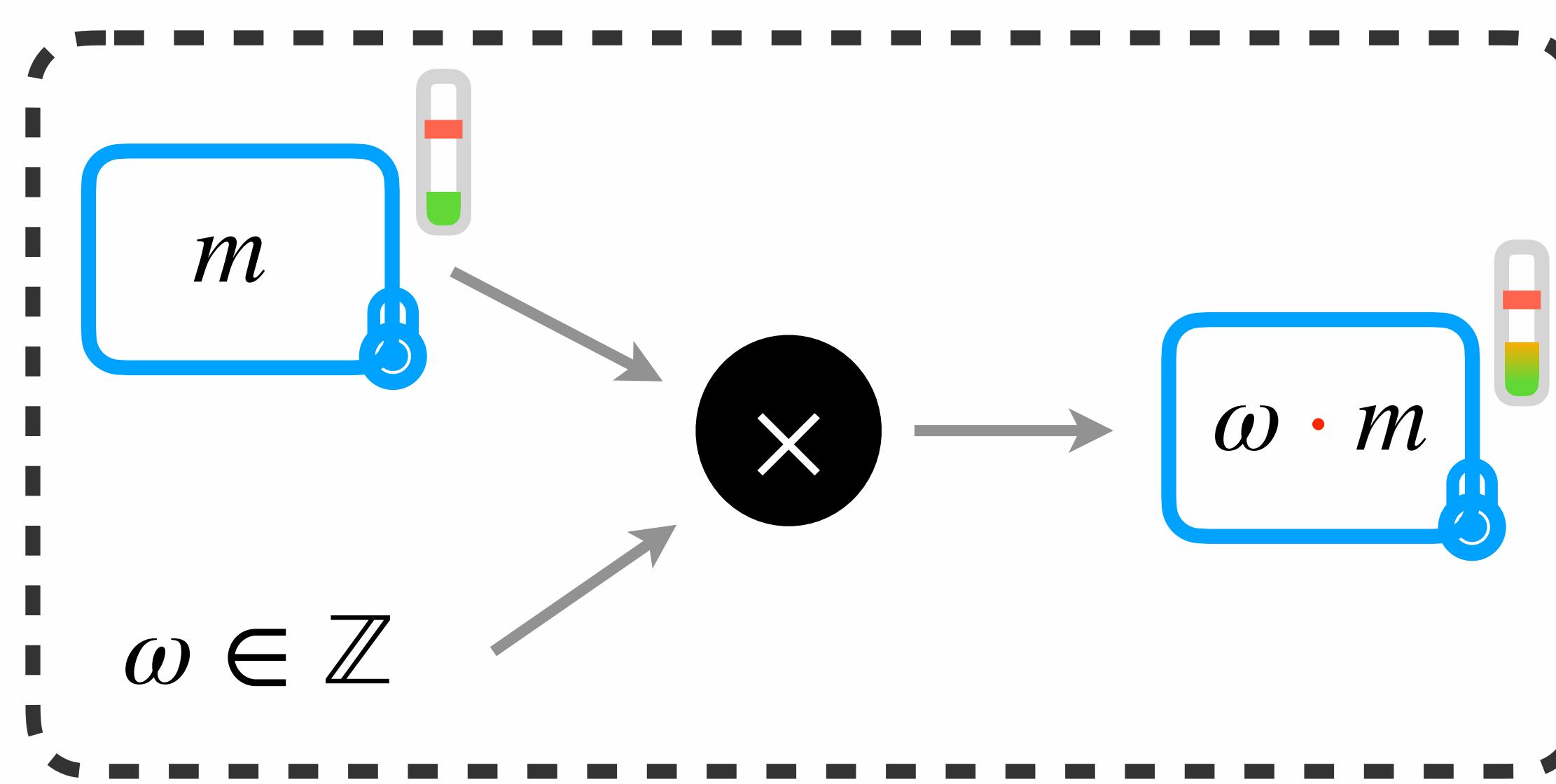


Dot Product

Dot Product

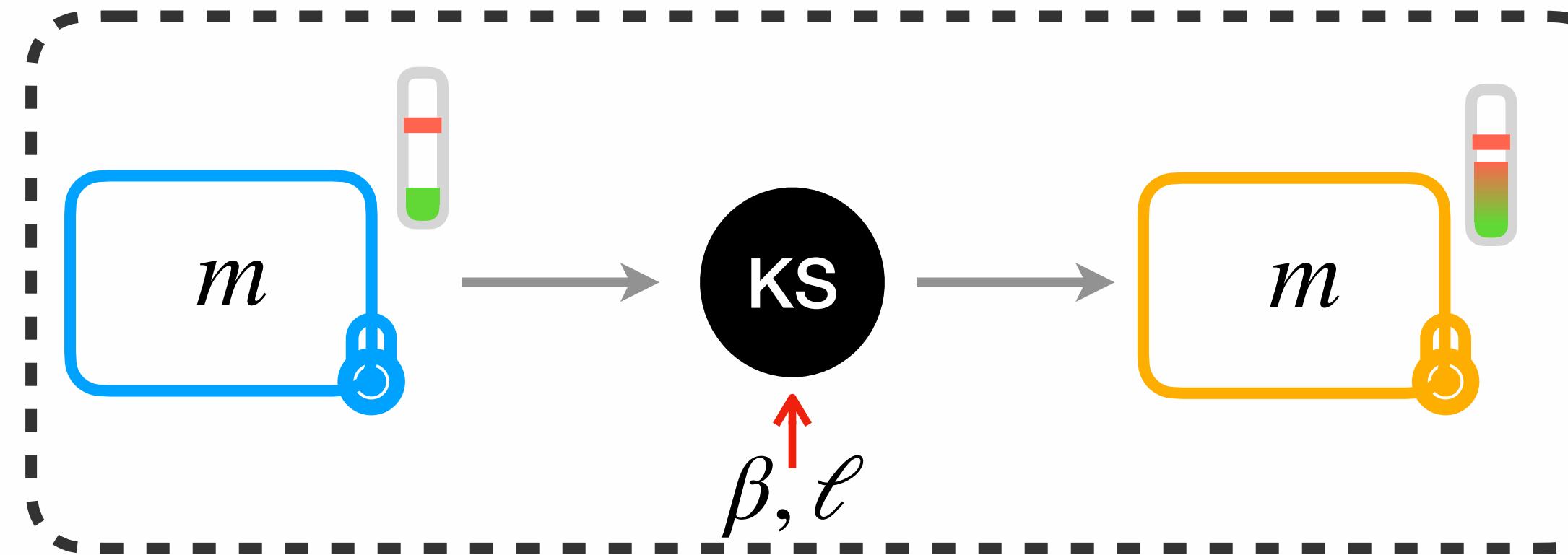


Dot Product



Keystwitch

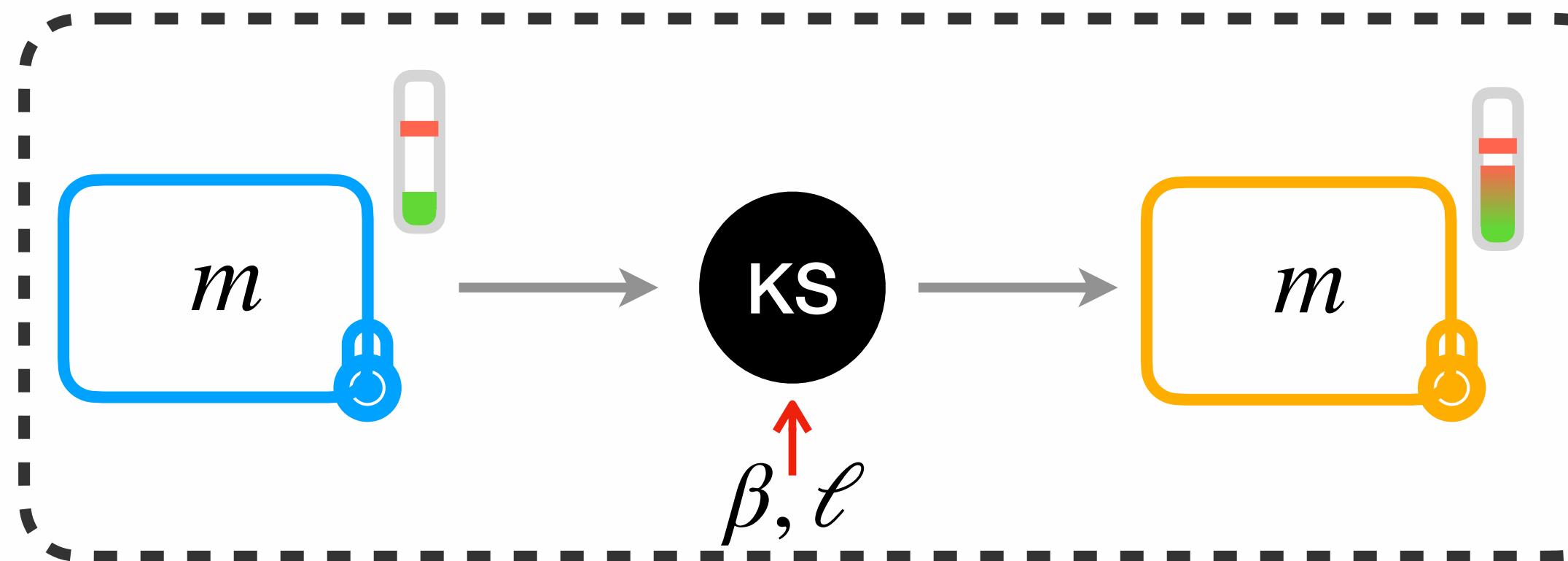
Keyswitch



Keyswitch

1 $b - \langle \vec{a}, \vec{s} \rangle = \Delta m + e \pmod{q}$

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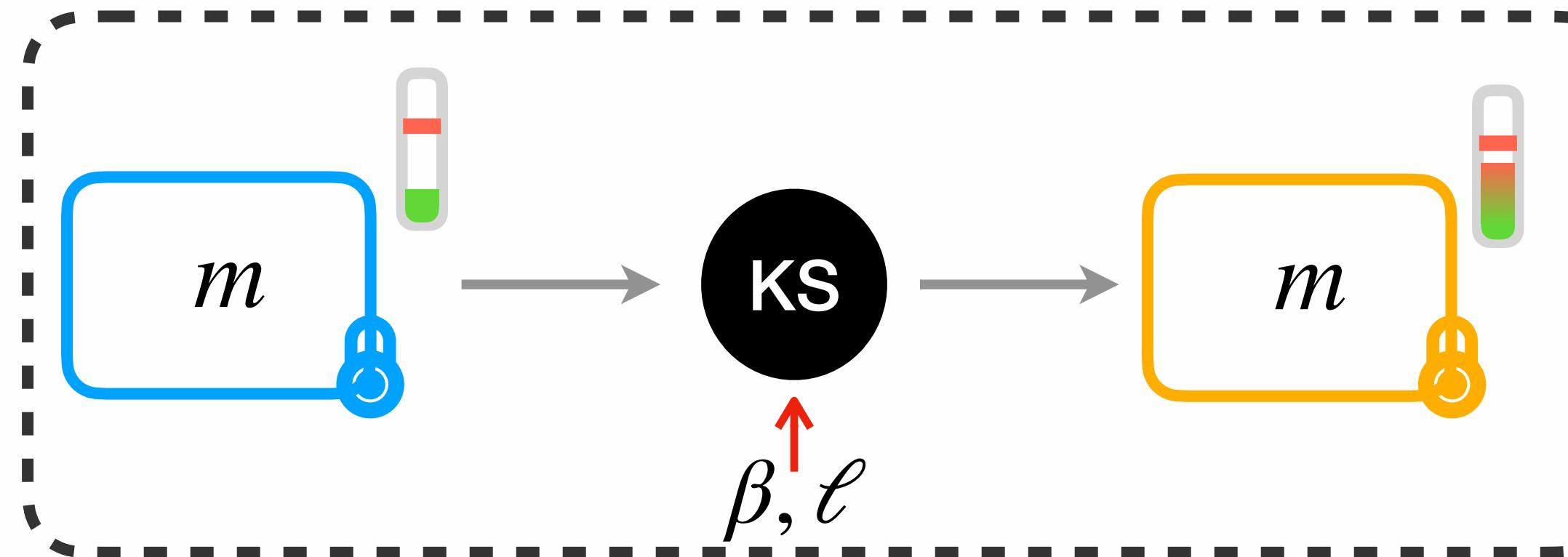


Keyswitch

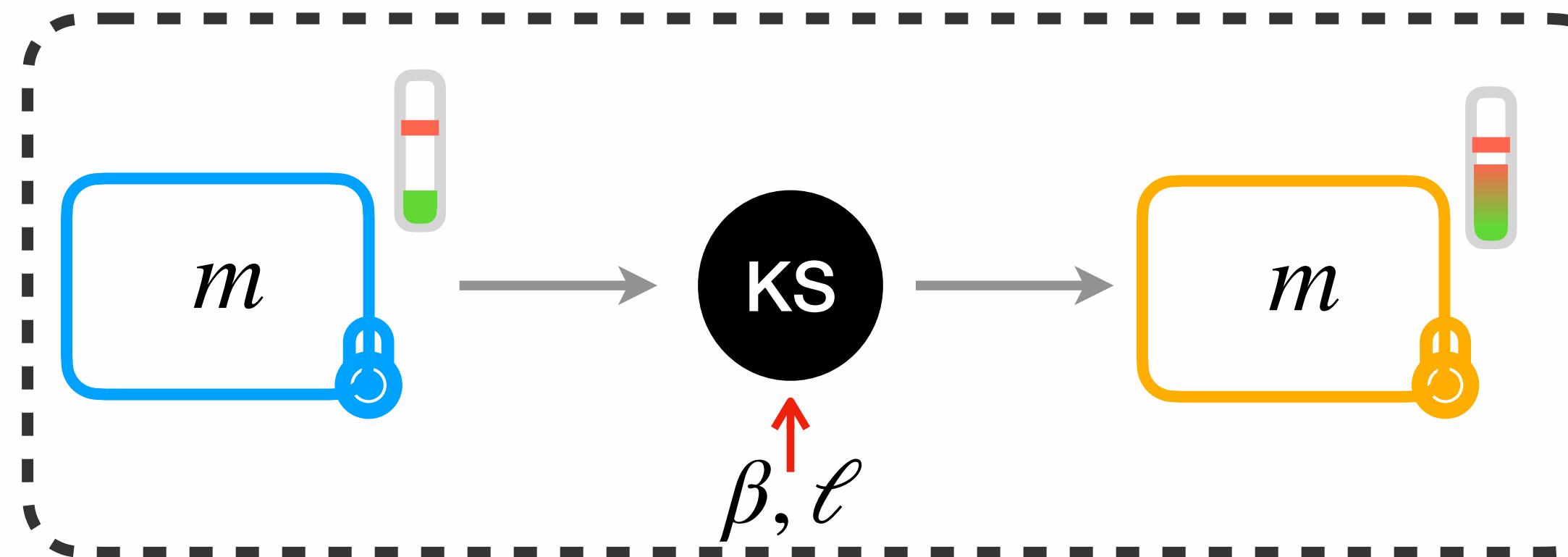
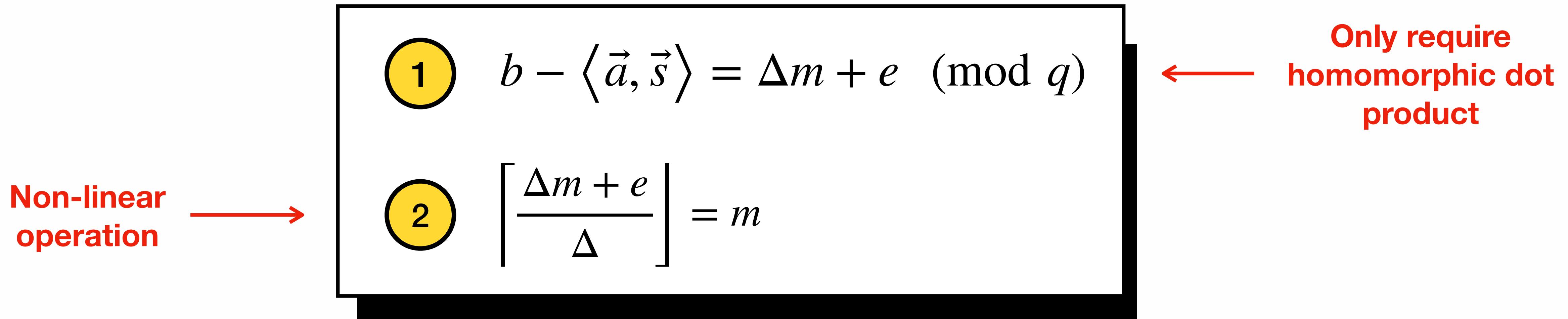
1 $b - \langle \vec{a}, \vec{s} \rangle = \Delta m + e \pmod{q}$

2 $\left\lceil \frac{\Delta m + e}{\Delta} \right\rceil = m$

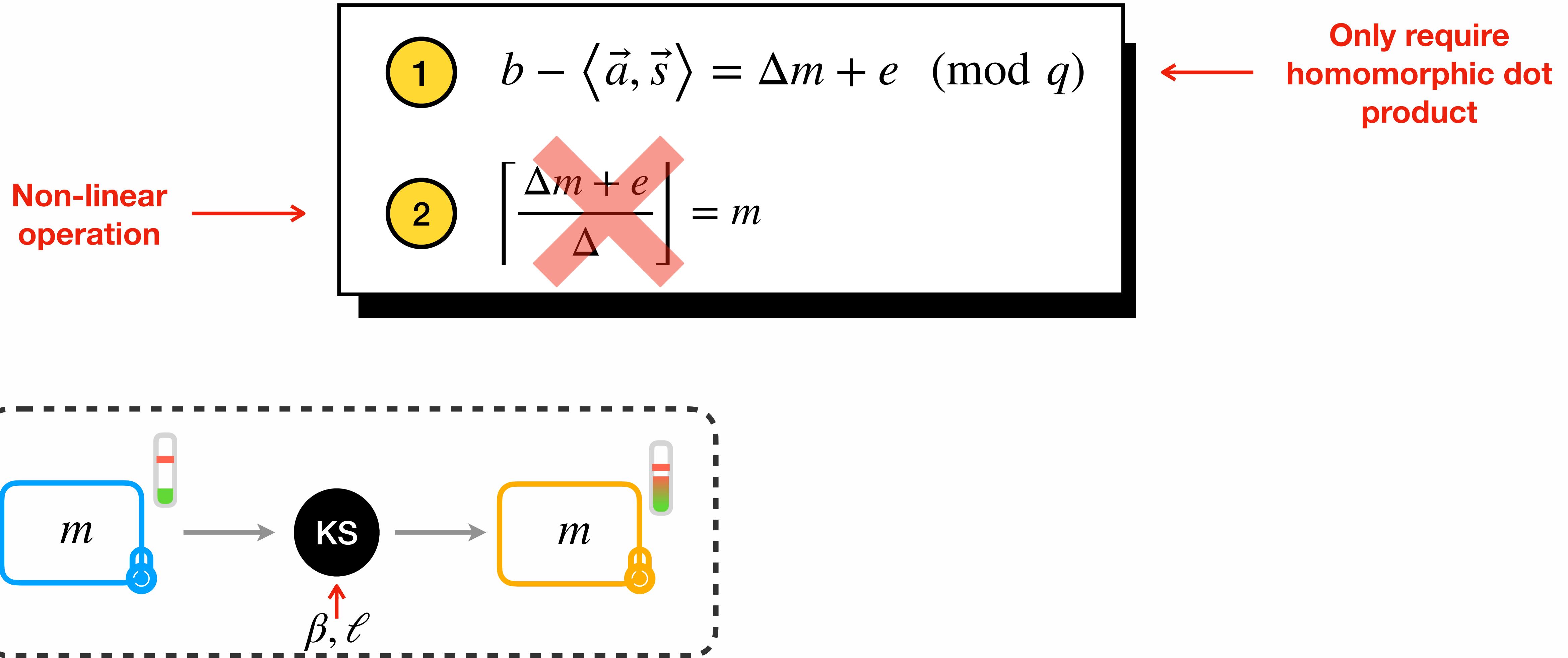
Only require
homomorphic dot
product



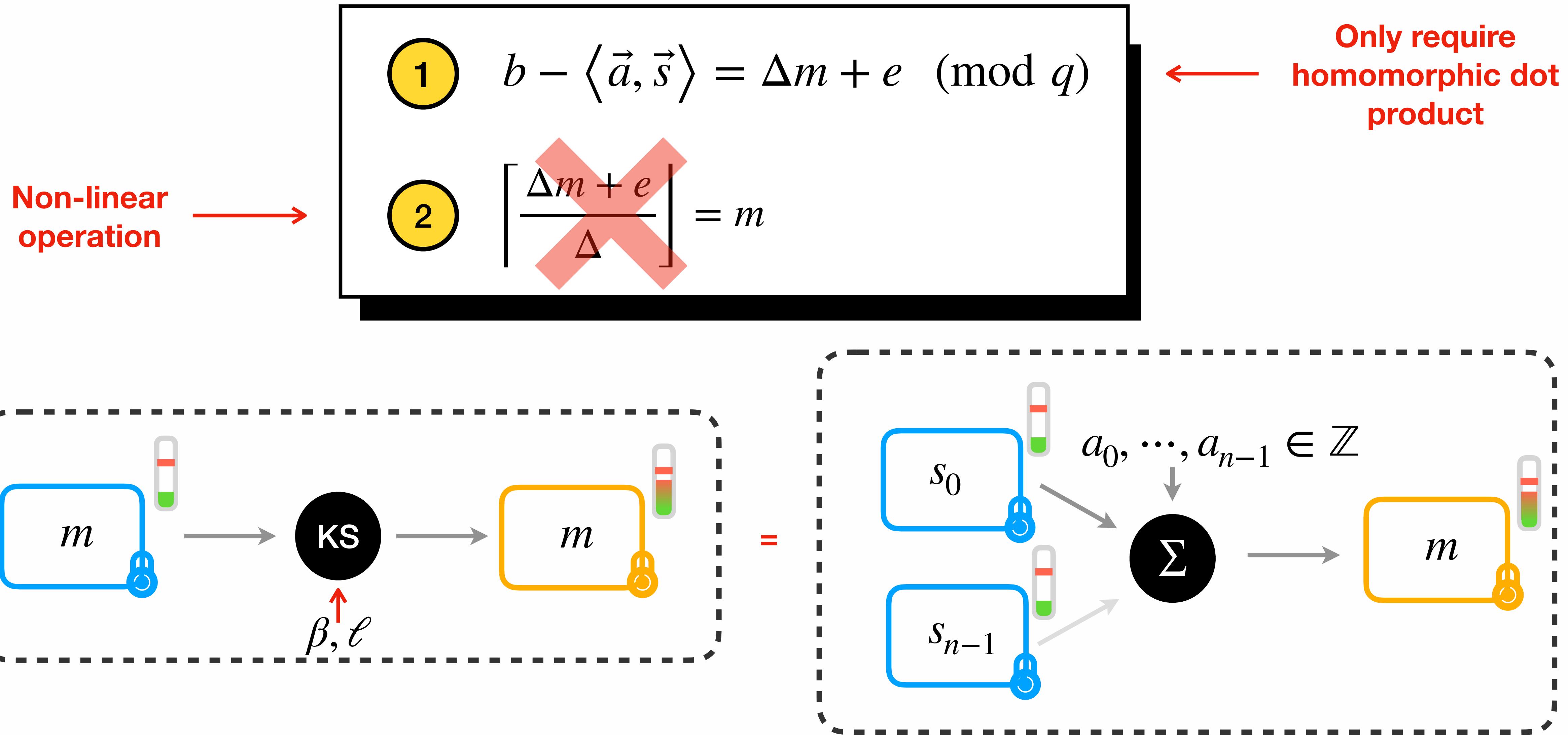
Keyswitch



Keyswitch

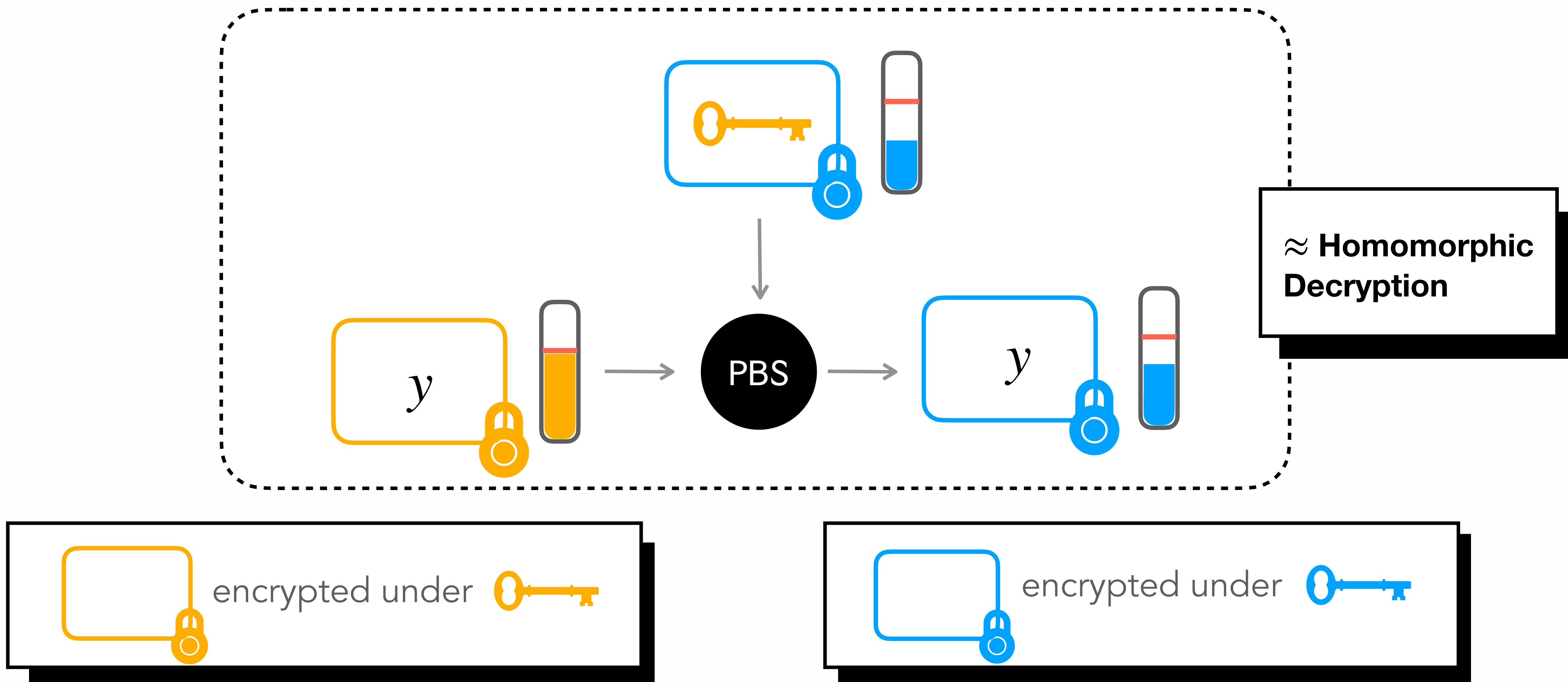


Keyswitch



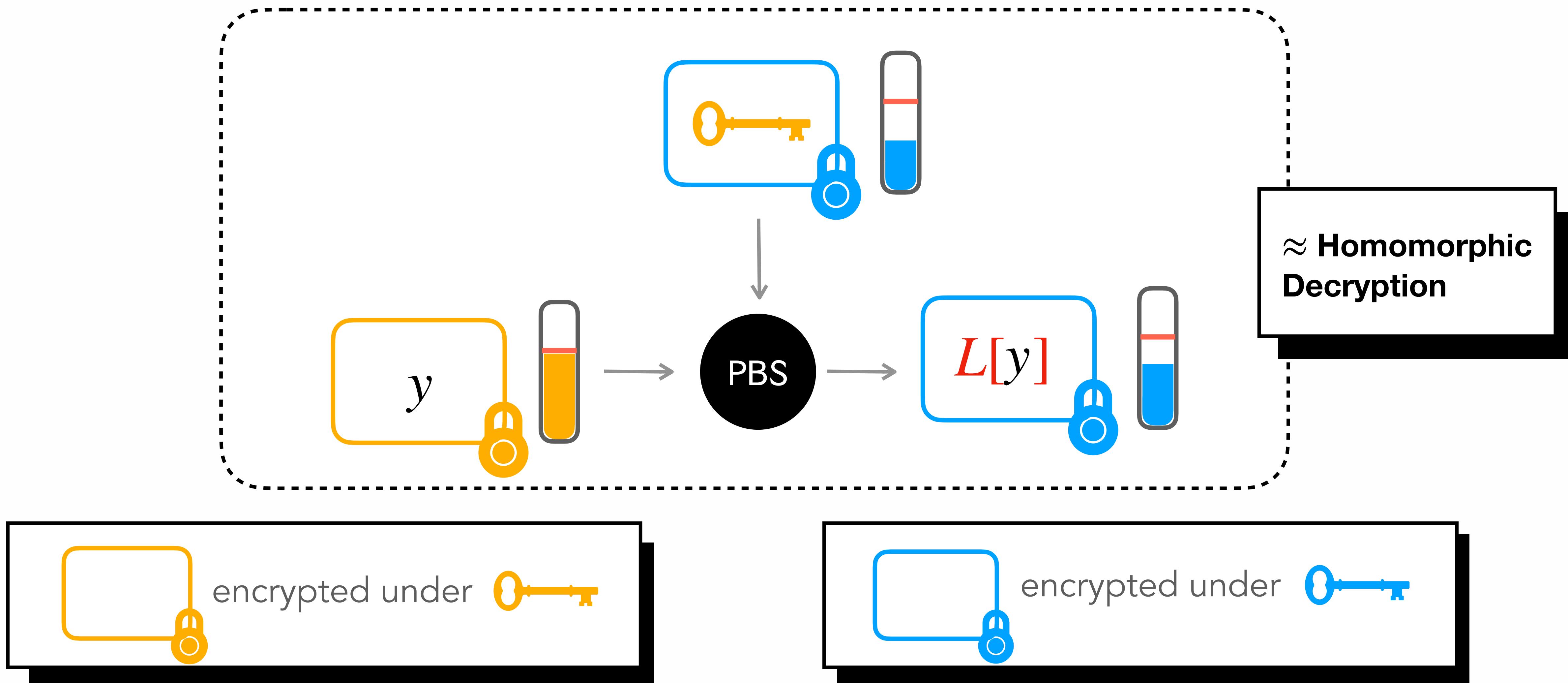
Programmable Bootstrapping (PBS)

Fully Homomorphic Encryption I



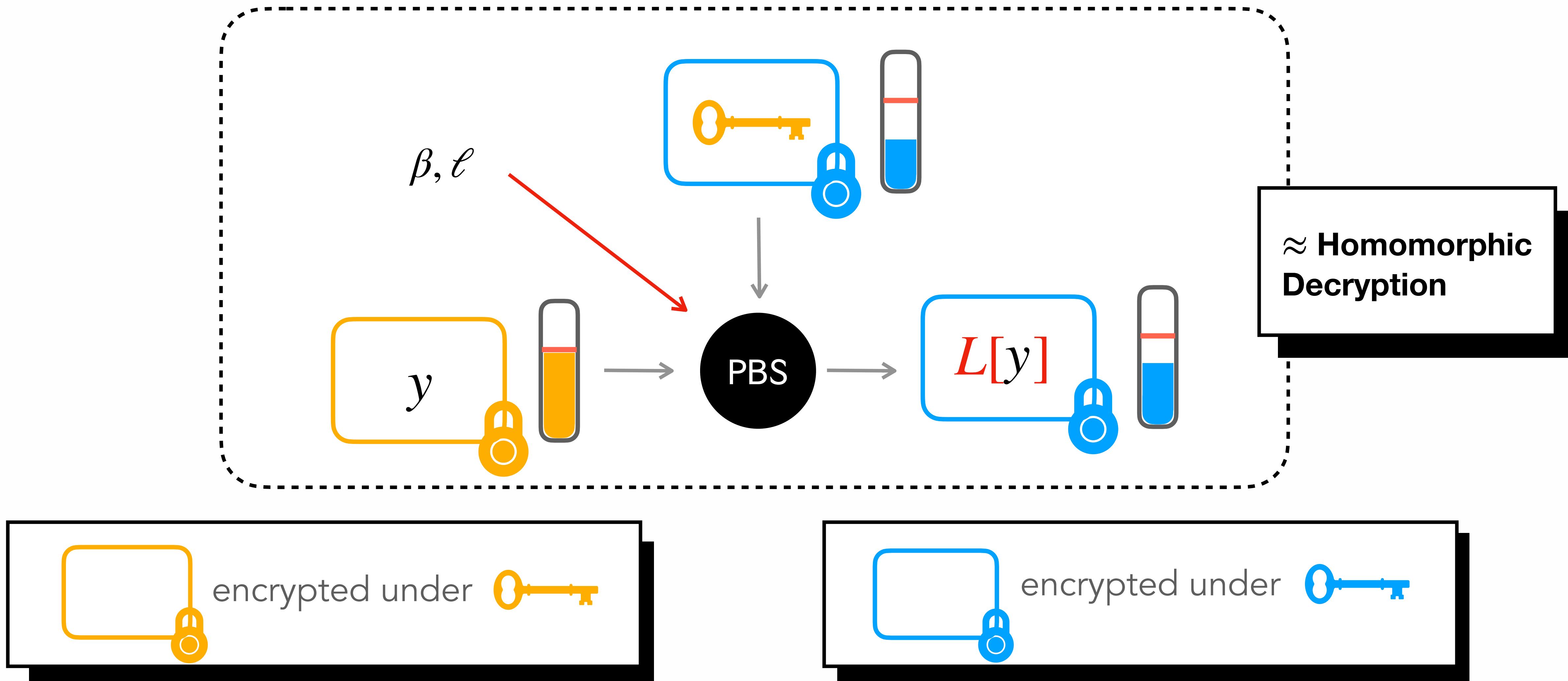
Programmable Bootstrapping (PBS)

Fully Homomorphic Encryption I

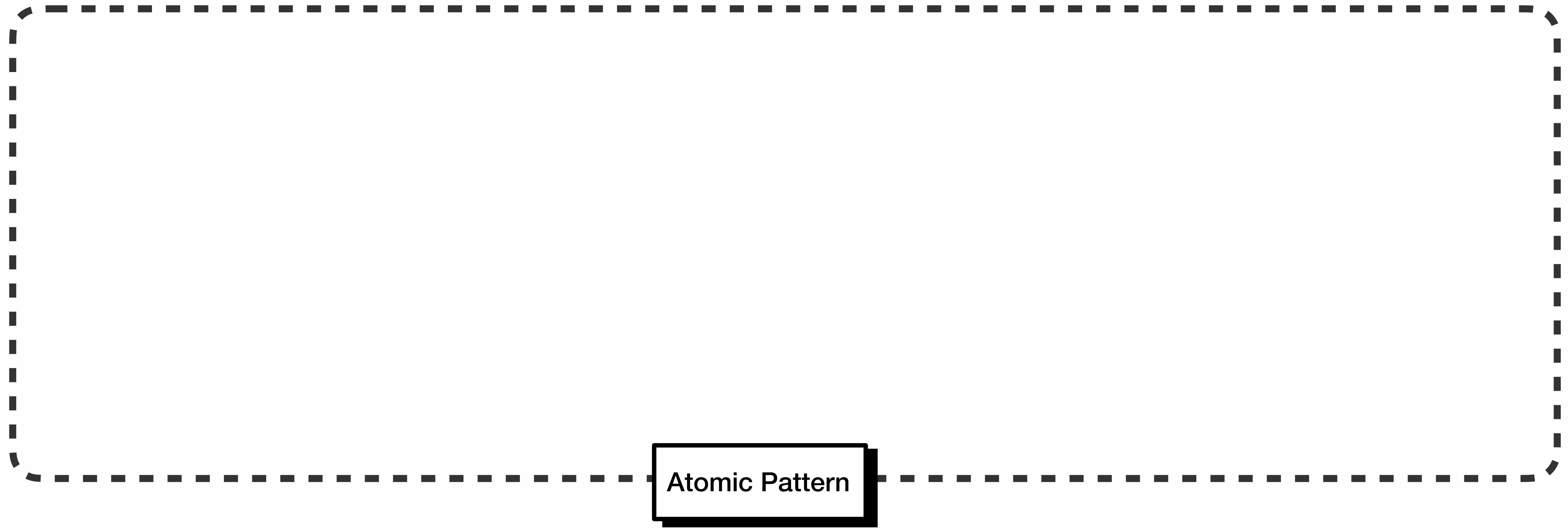


Programmable Bootstrapping (PBS)

Fully Homomorphic Encryption I

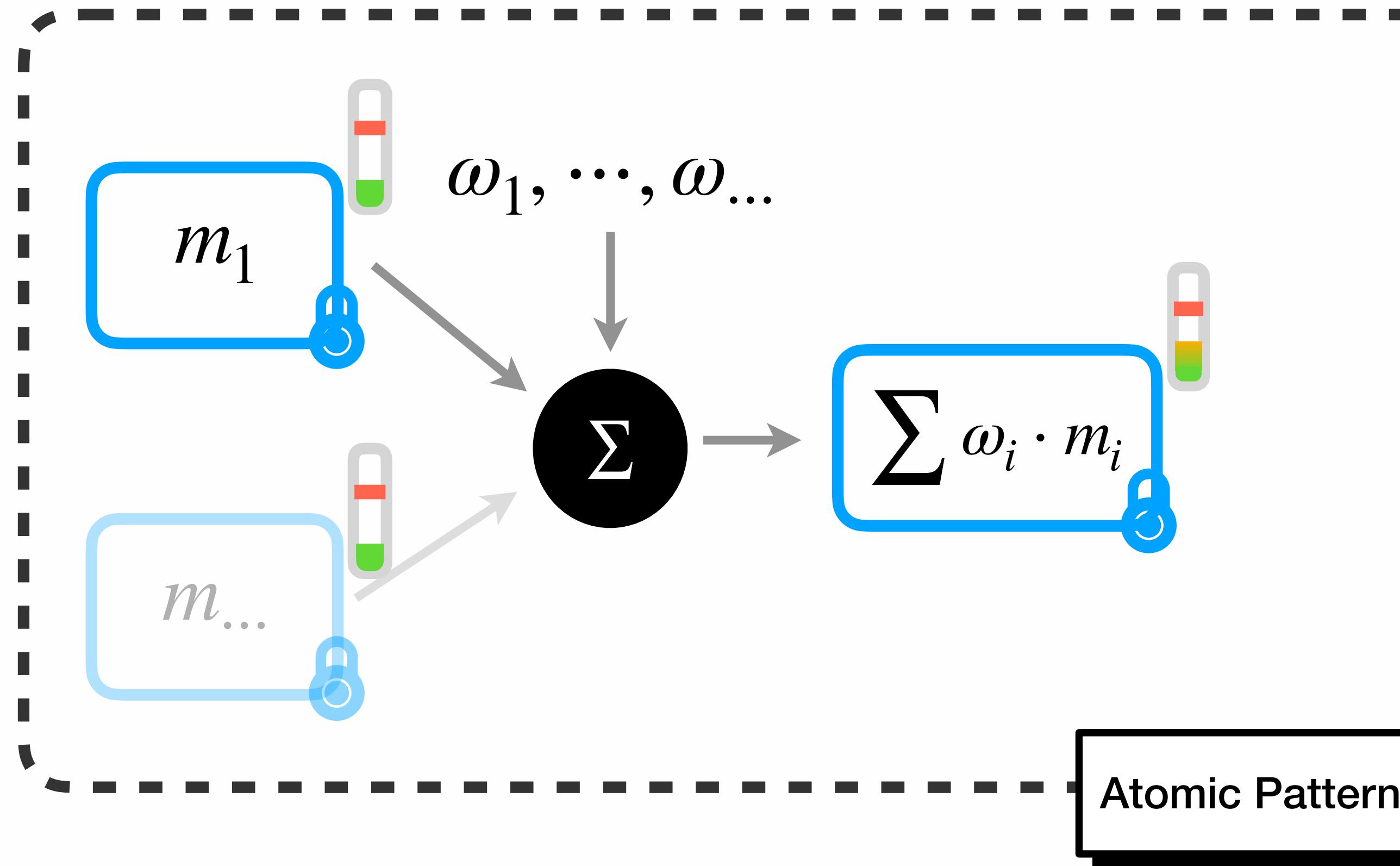


Optimal arrangement of TFHE building blocks



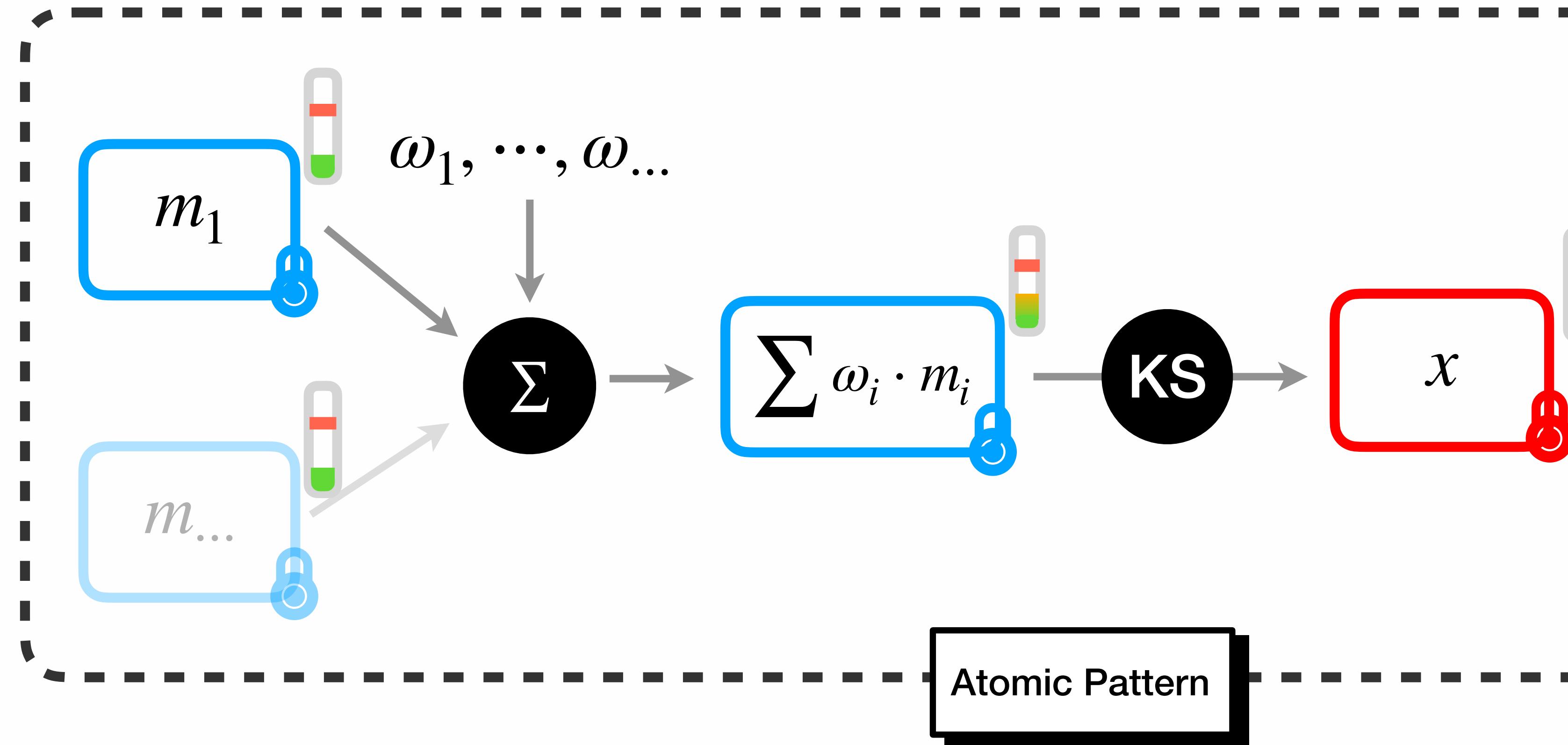
Optimal arrangement of TFHE building blocks

Fully Homomorphic Encryption I



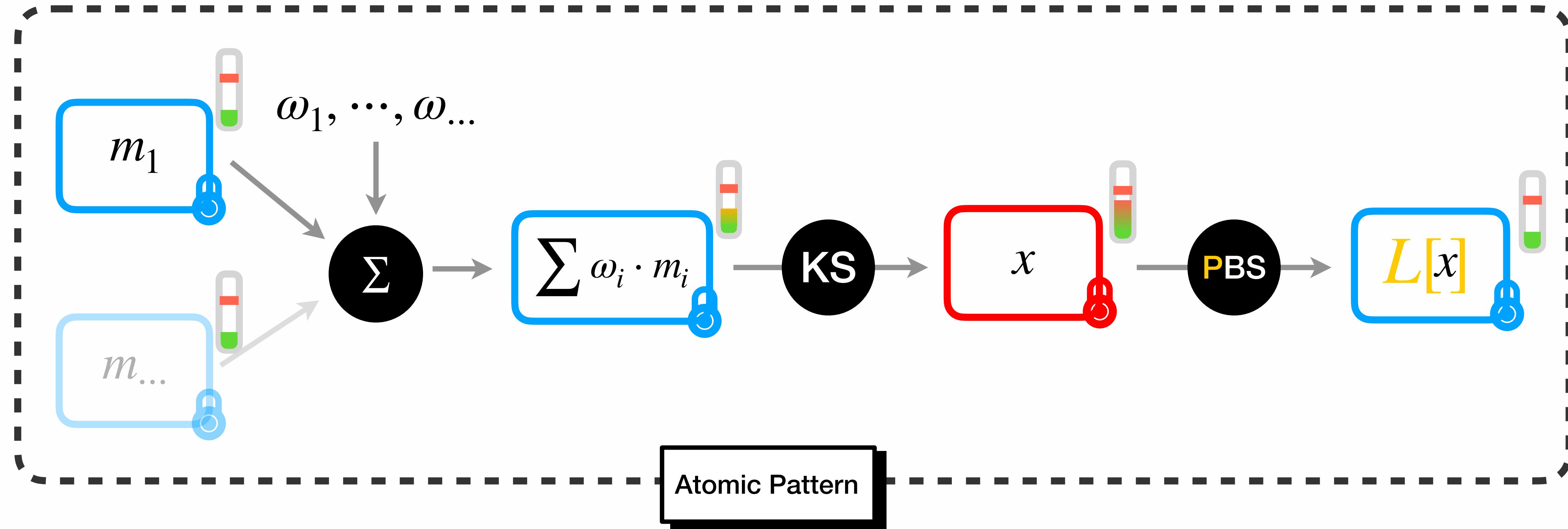
Optimal arrangement of TFHE building blocks

Fully Homomorphic Encryption I



Optimal arrangement of TFHE building blocks

Fully Homomorphic Encryption I



Concrete Optimizer

An optimizer for TFHE

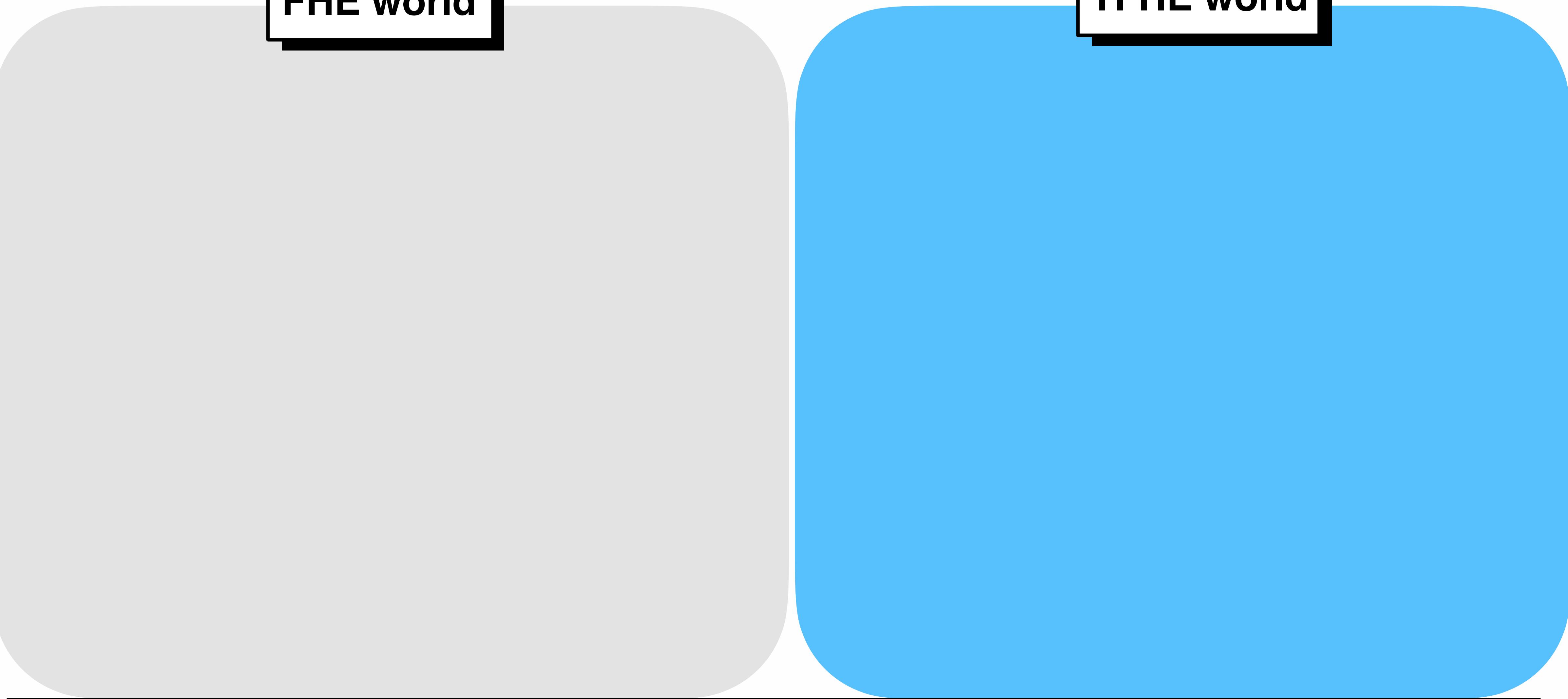
Goals

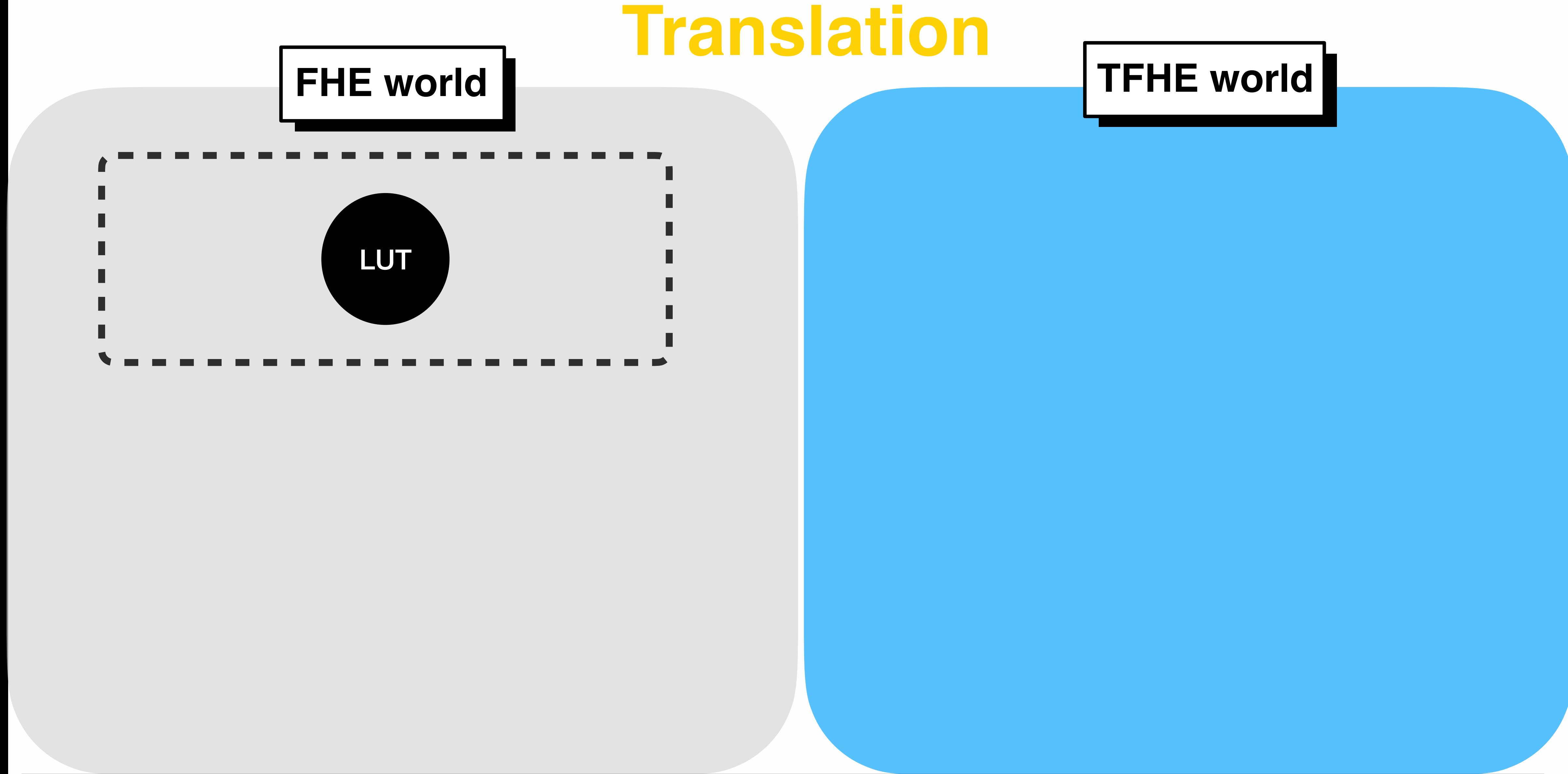
Translate a function
 f into a graph of
TFHE operators

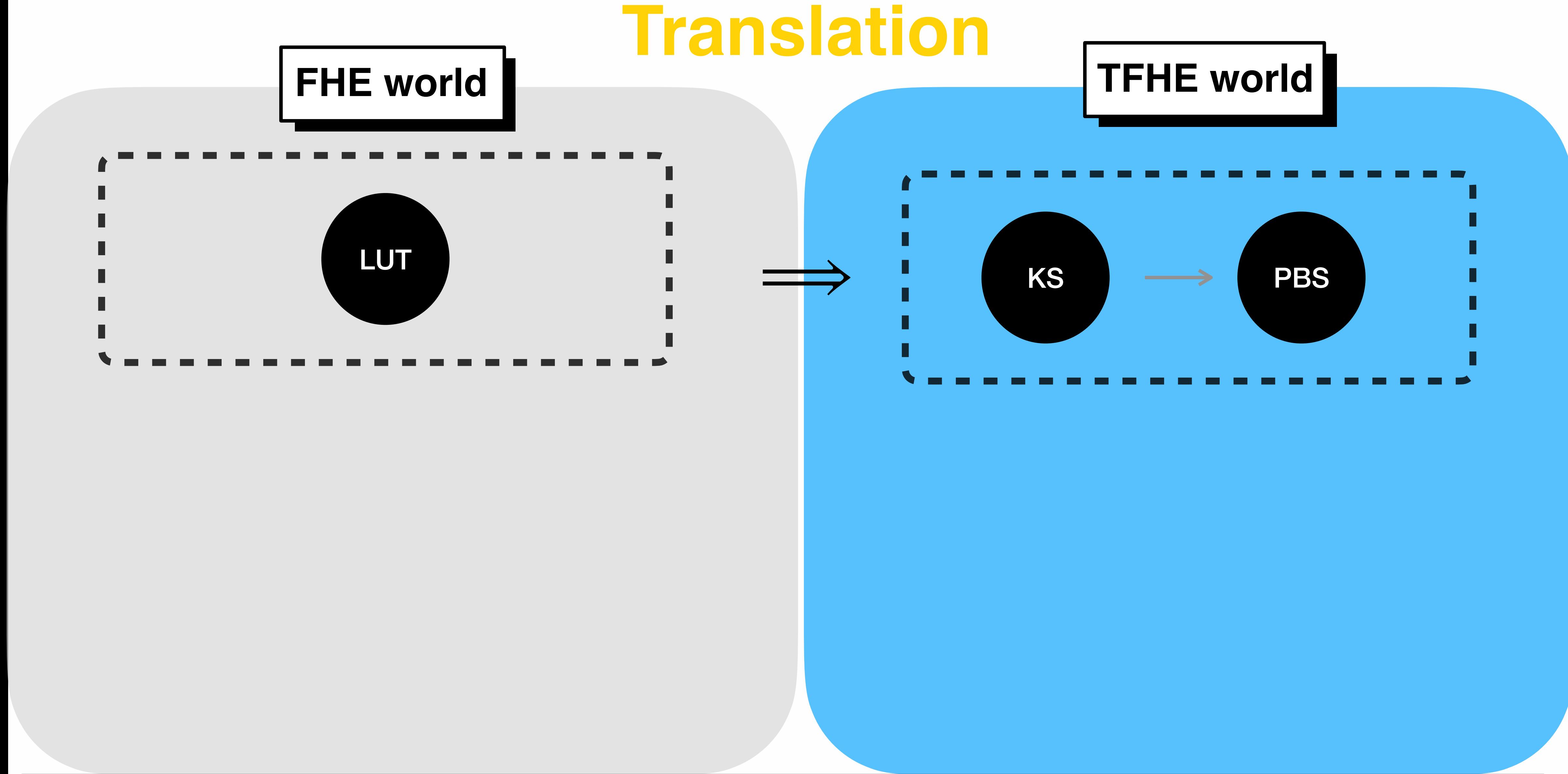
Translation

FHE world

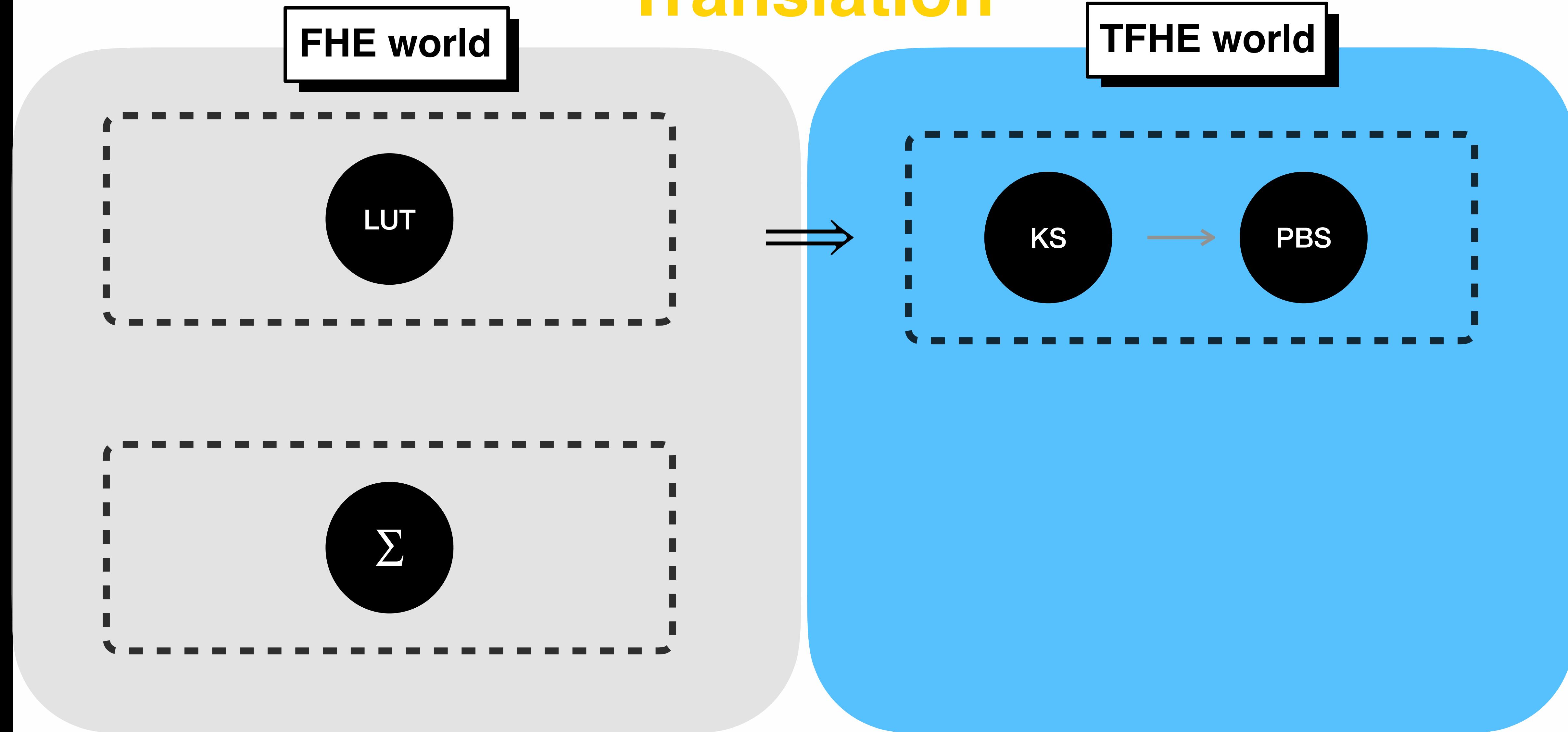
TFHE world



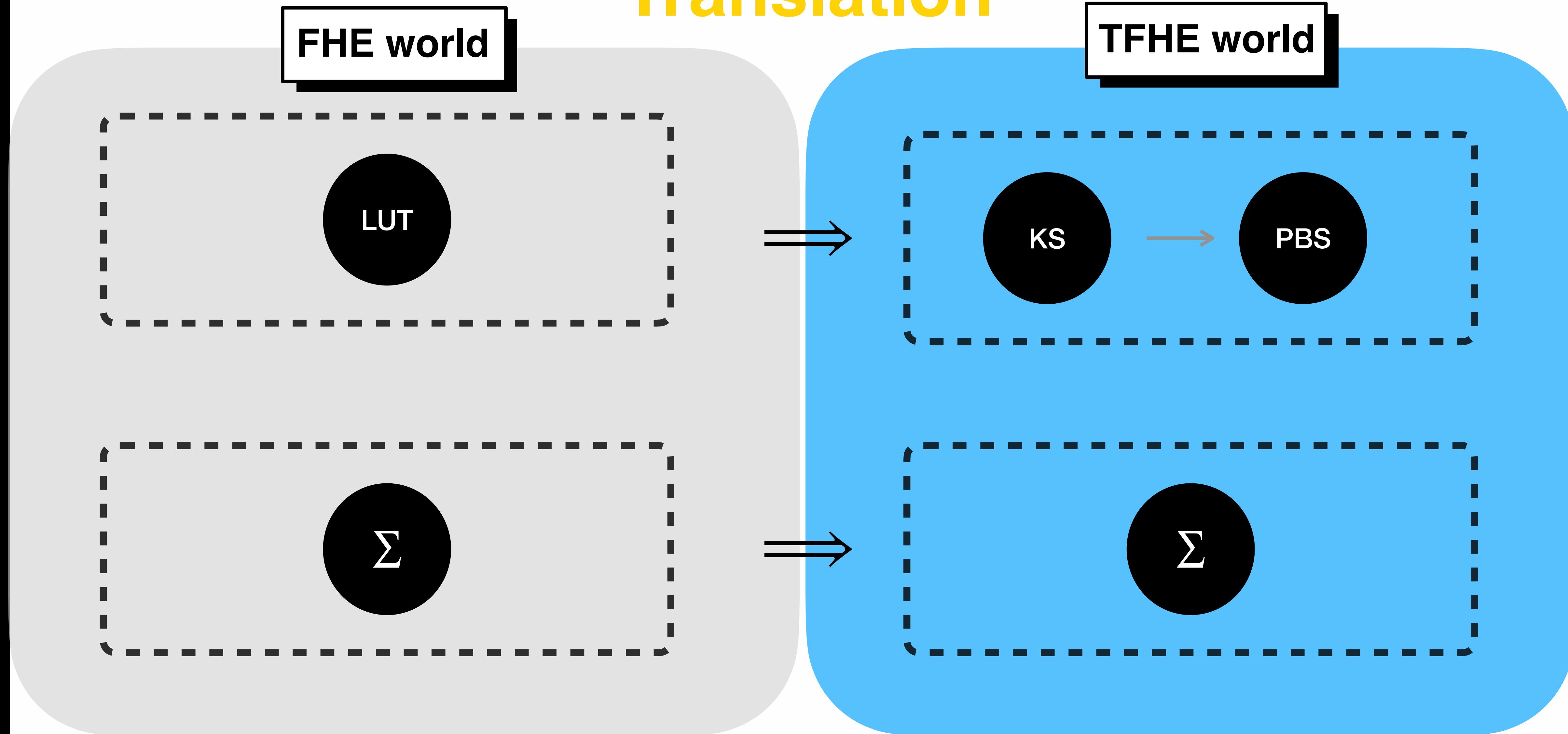




Translation

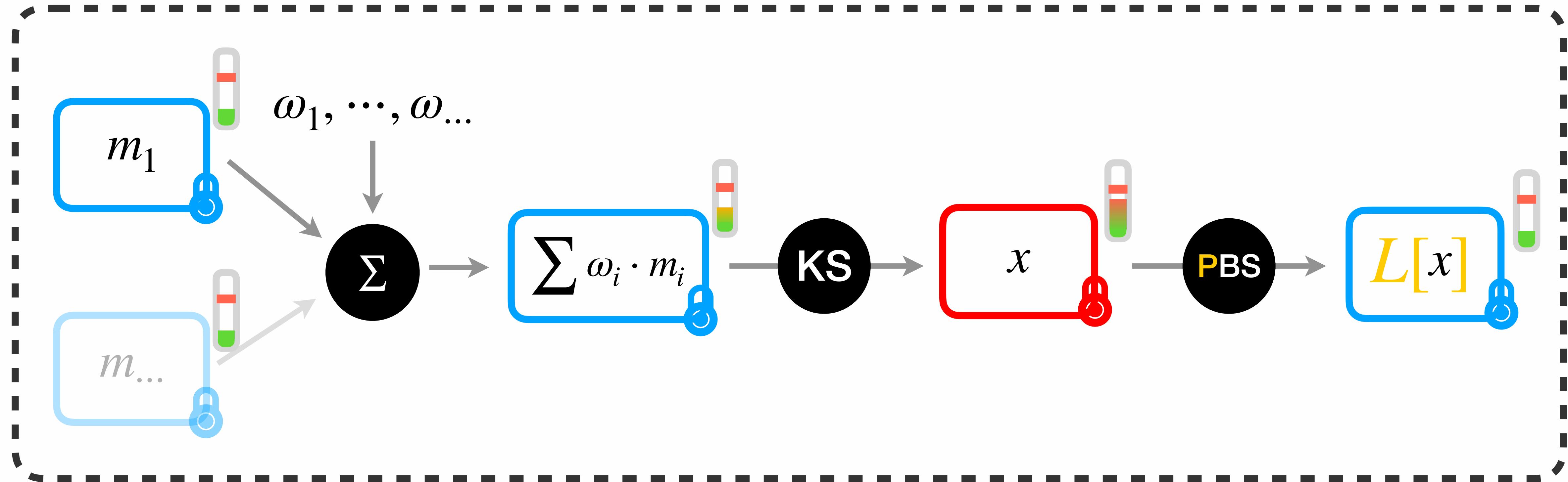


Translation



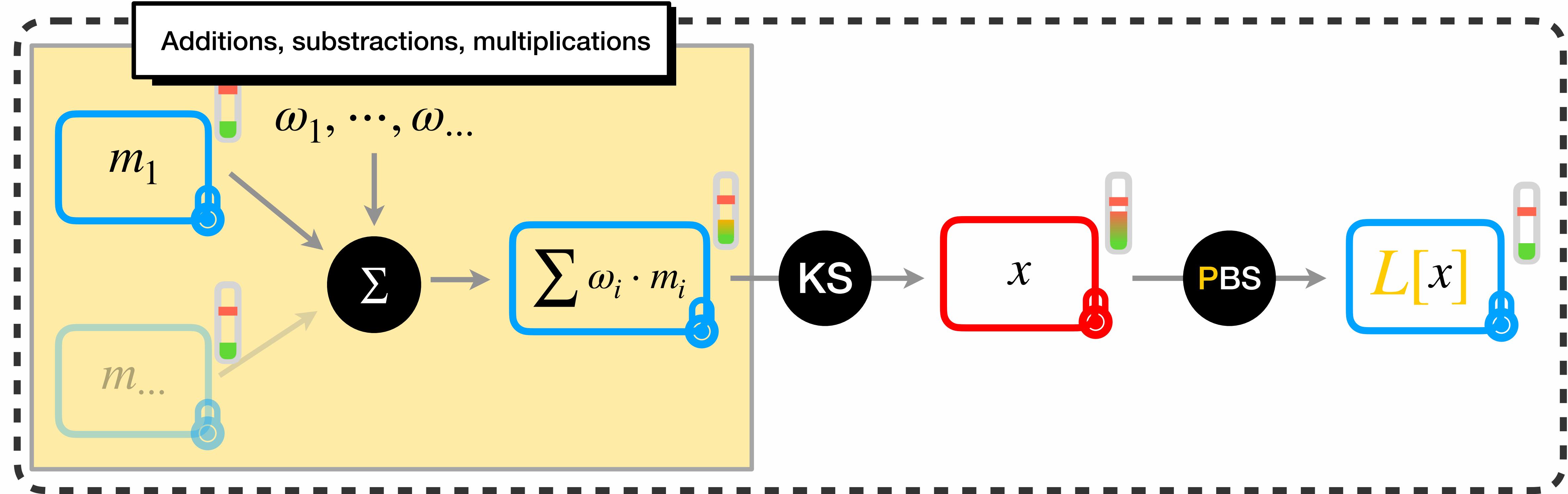
Translation

Fully Homomorphic Encryption I



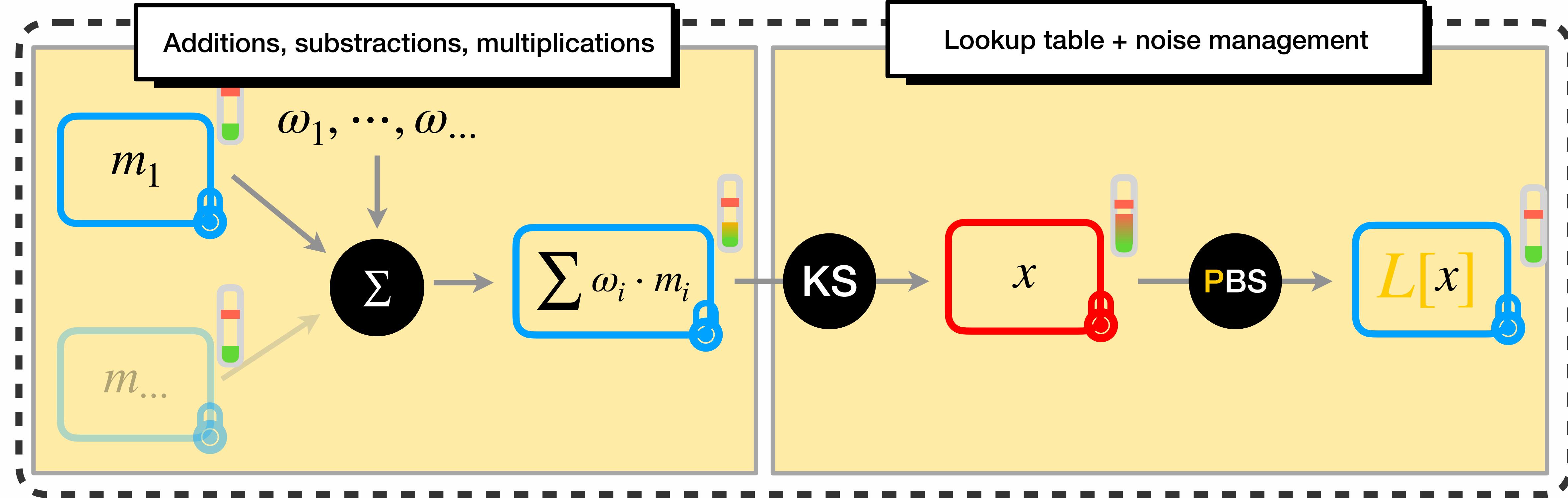
Translation

Fully Homomorphic Encryption I



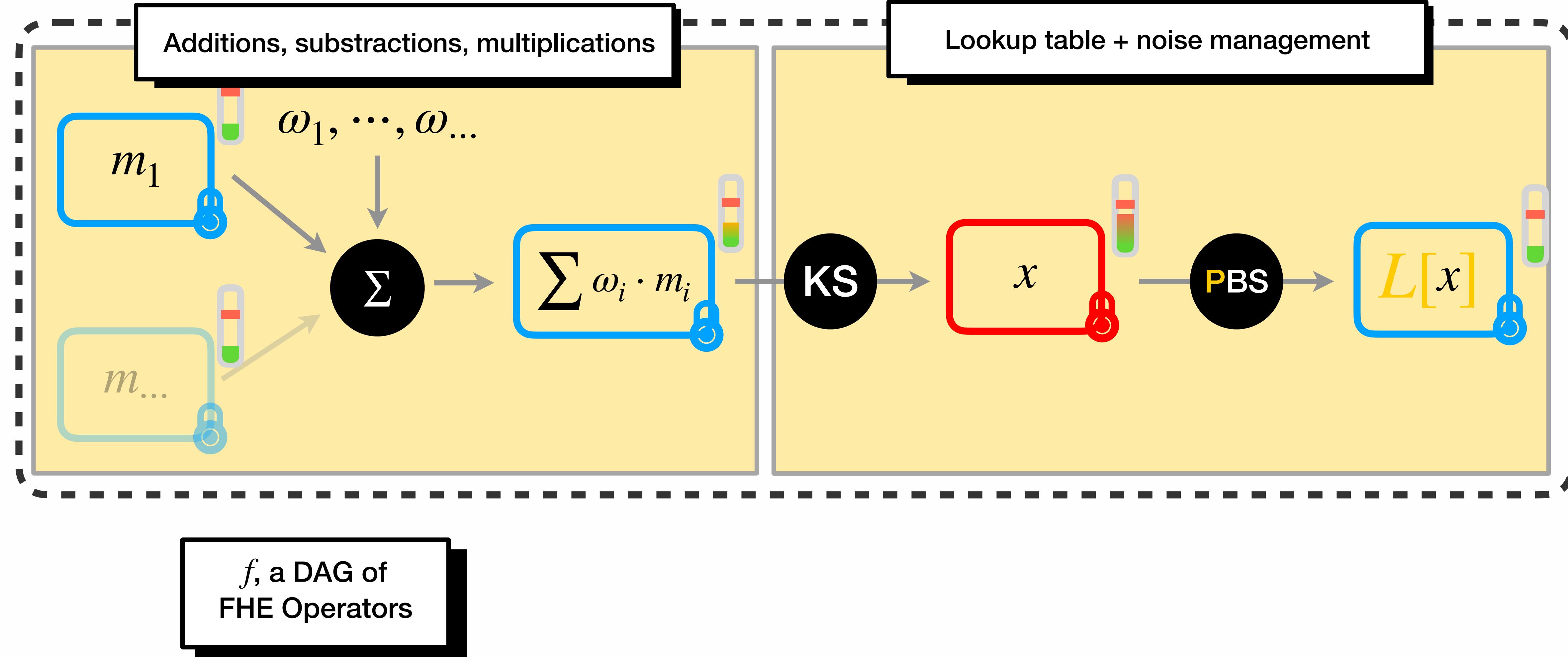
Translation

Fully Homomorphic Encryption I



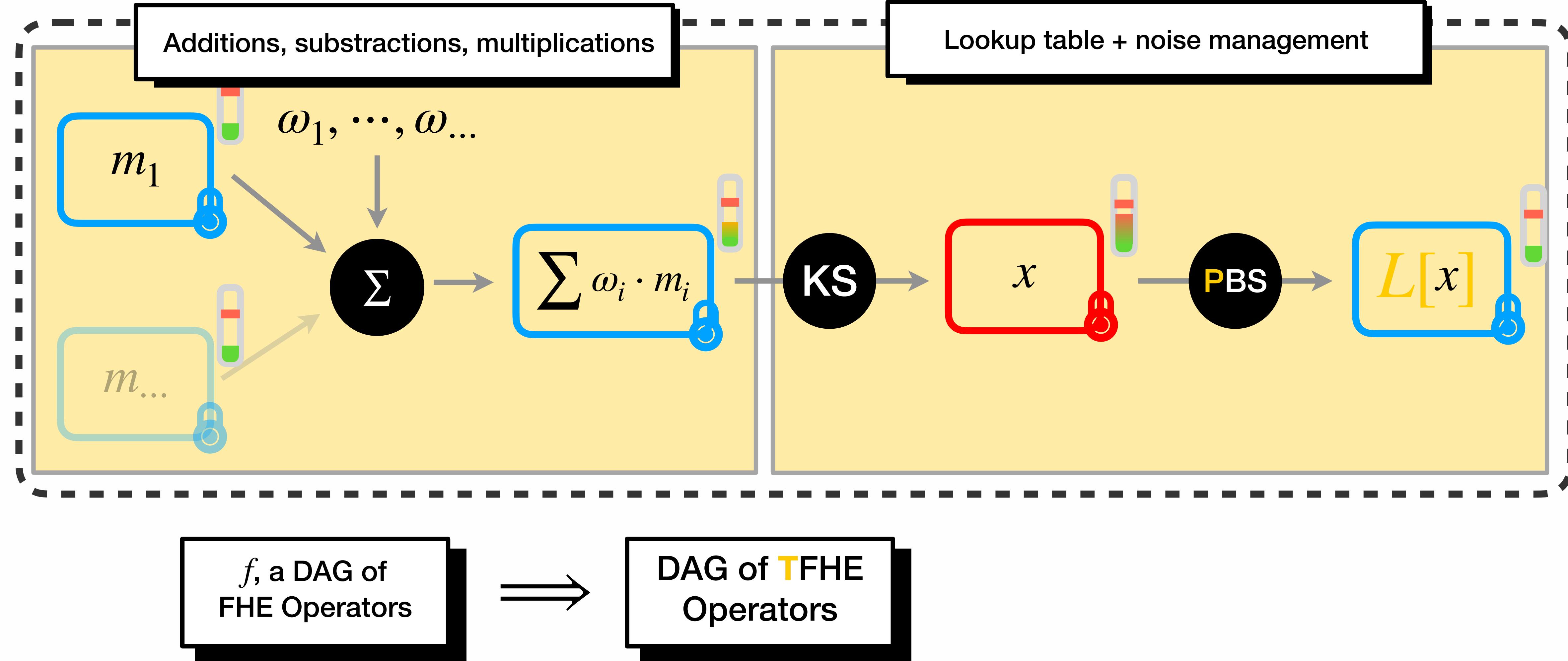
Translation

Fully Homomorphic Encryption I



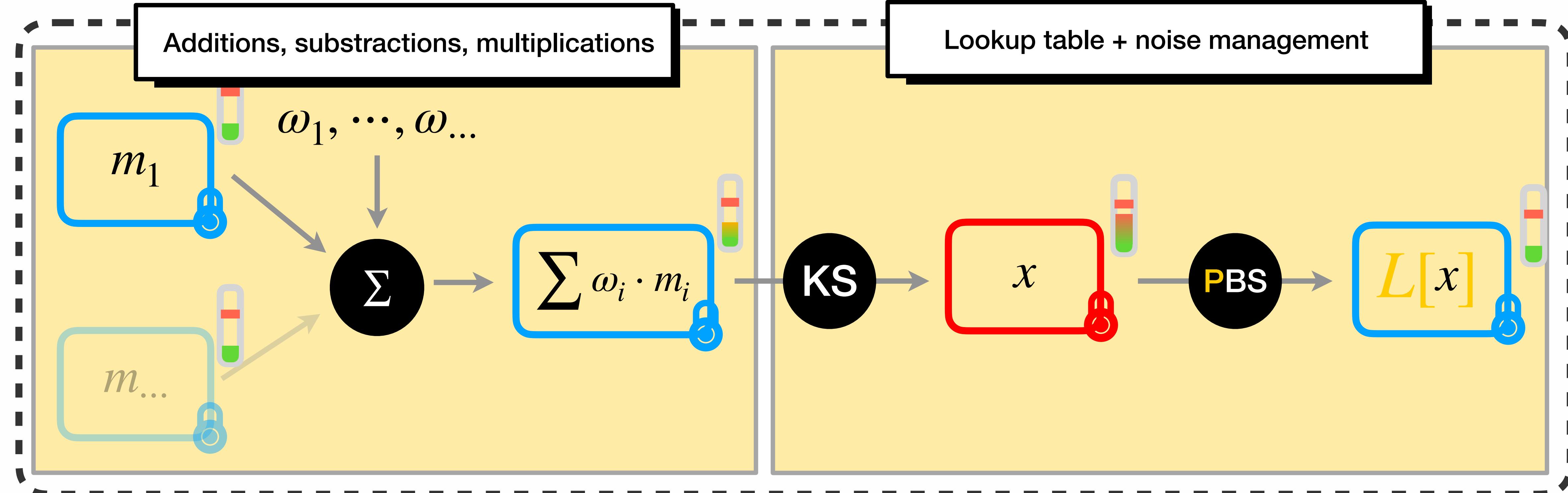
Translation

Fully Homomorphic Encryption I



Translation

Fully Homomorphic Encryption I



f , a DAG of
FHE Operators



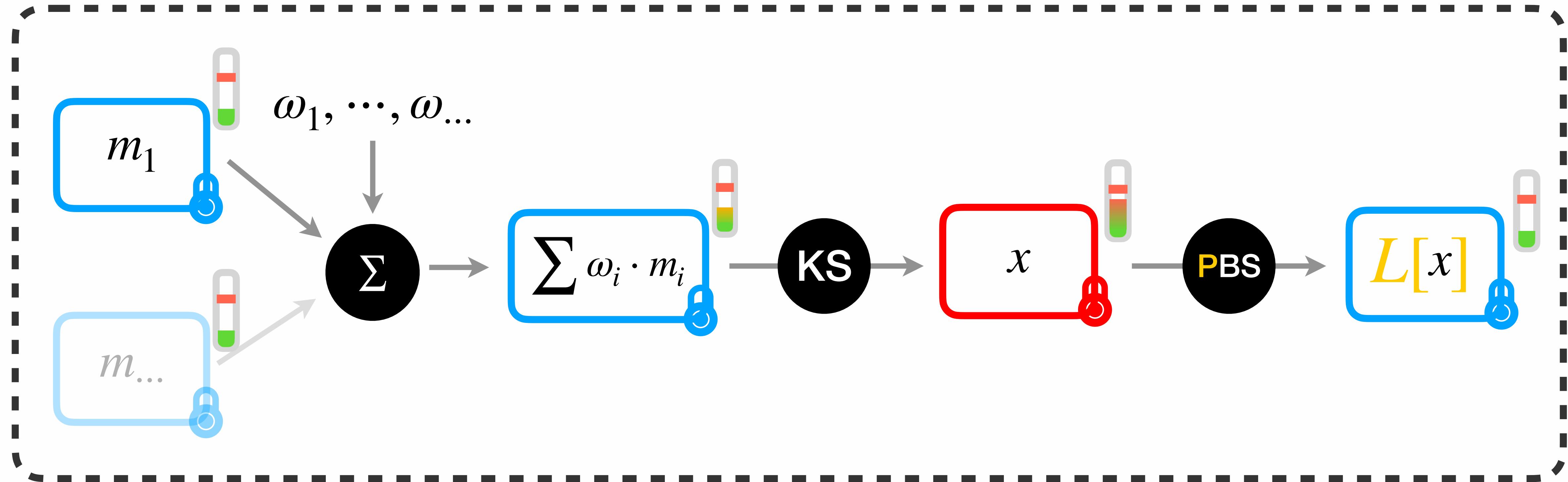
DAG of **TFHE**
Operators



DAG of atomic
patterns

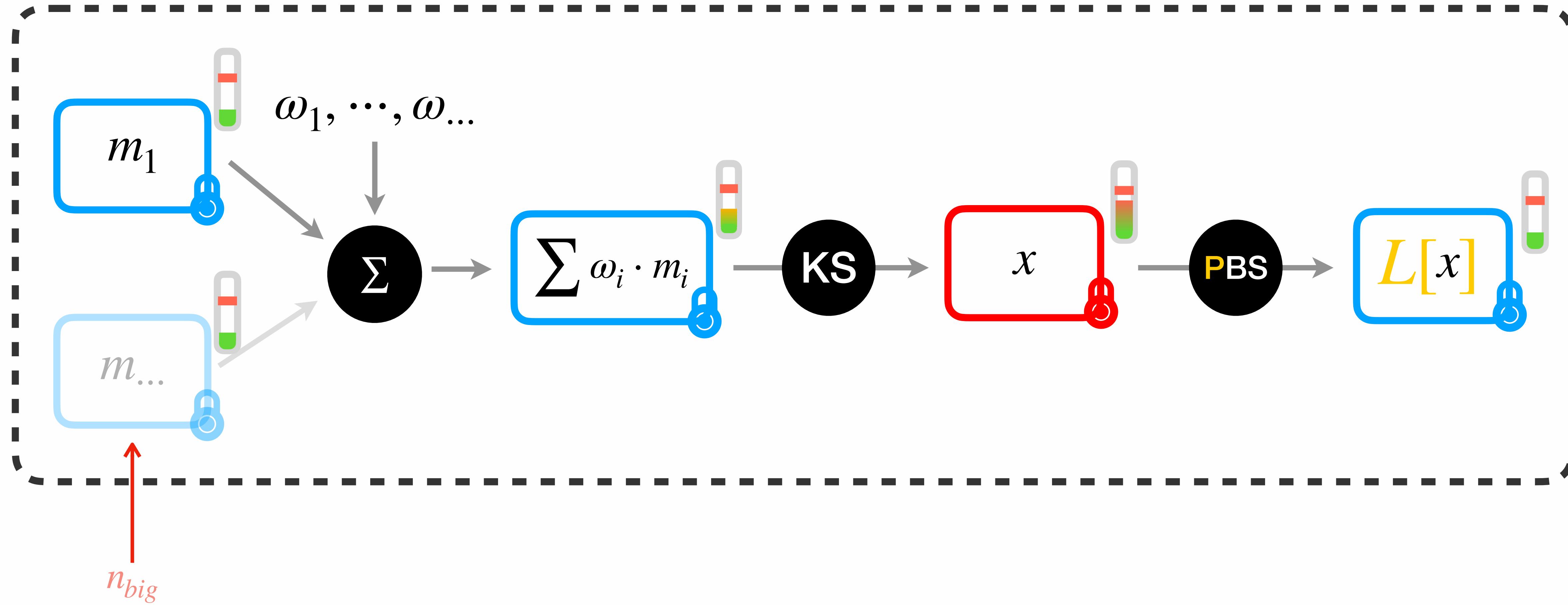
Translation

Fully Homomorphic Encryption I

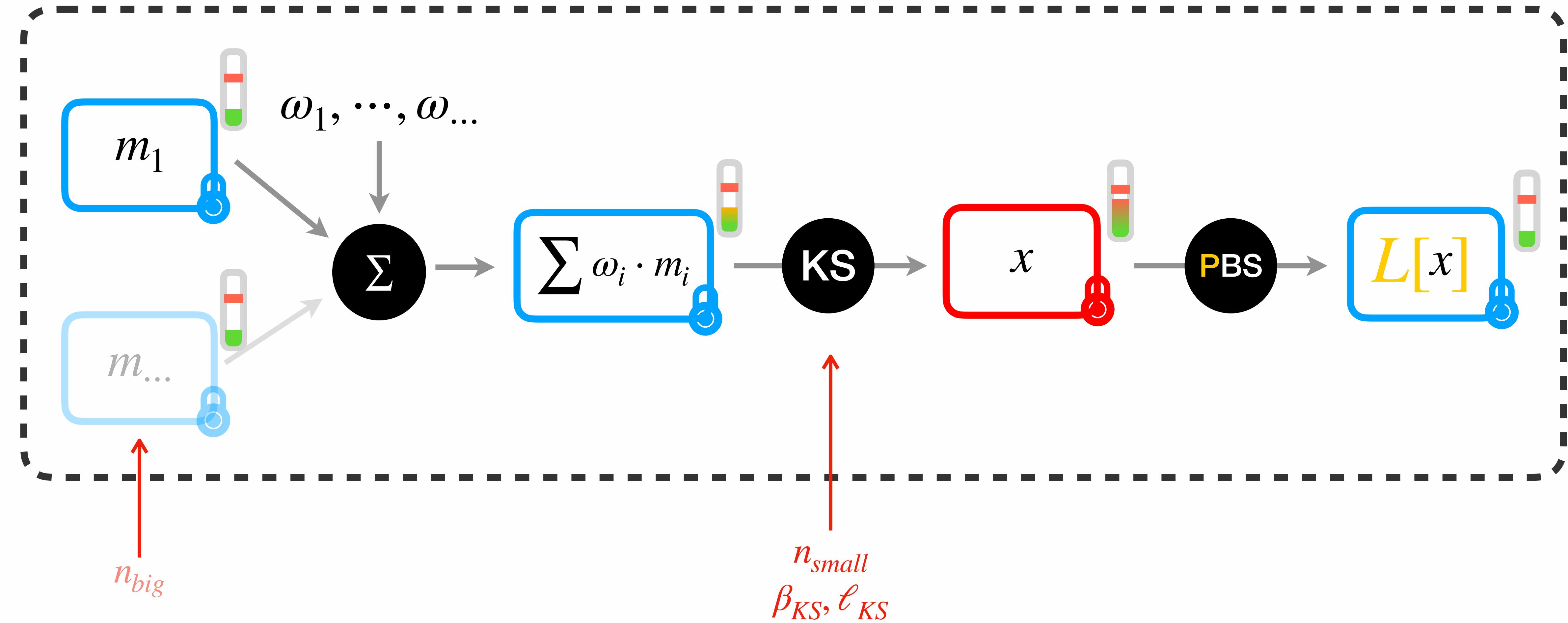


Translation

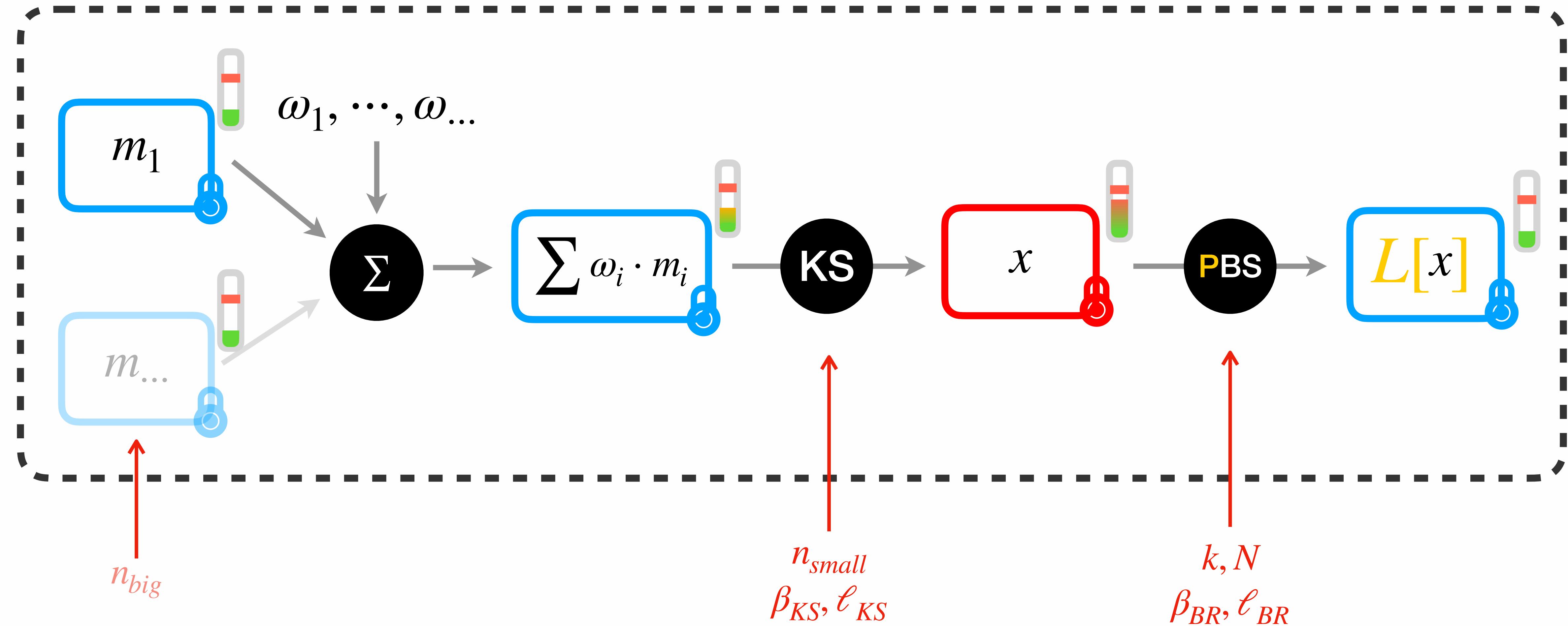
Fully Homomorphic Encryption I



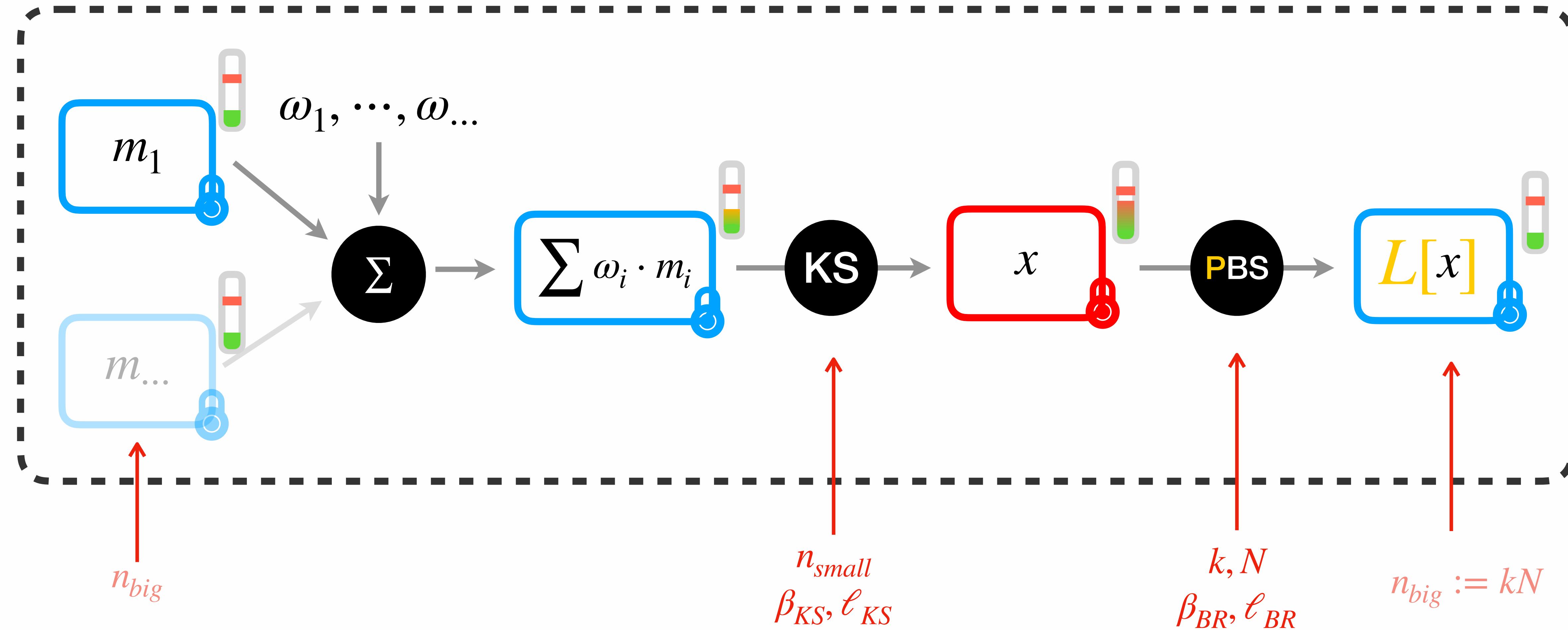
Translation



Translation



Translation



Goals

Translate a function
 f into a graph of
TFHE operators

with

Goals

Translate a function
 f into a graph of
TFHE operators

with



Security

Goals

Translate a function
 f into a graph of
TFHE operators

with



Security

Correctness

Goals

Translate a function
 f into a graph of
TFHE operators

with



Security



Correctness

Efficiency

Security



Security



Lattice Estimator [APS]

Reference tool to estimate
the security of lattice-based
encryption scheme **[APS]**

Security



Lattice Estimator [APS]

Reference tool to estimate
the security of lattice-based
encryption scheme **[APS]**

$$g : (n, \lambda, q) \mapsto \sigma_{\text{enc}}$$

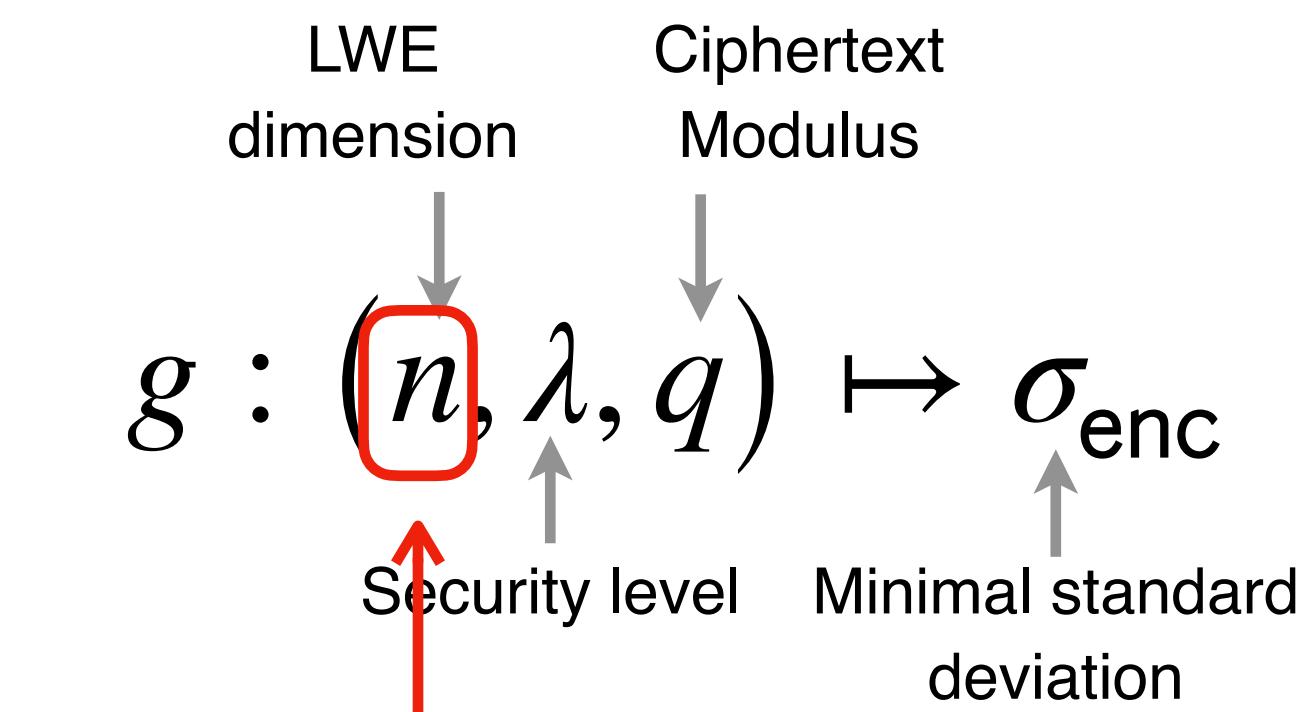
↓ ↓
LWE dimension Ciphertext Modulus
↓ ↑
Security level Minimal standard deviation

Security



Lattice Estimator [APS]

Reference tool to estimate
the security of lattice-based
encryption scheme [APS]



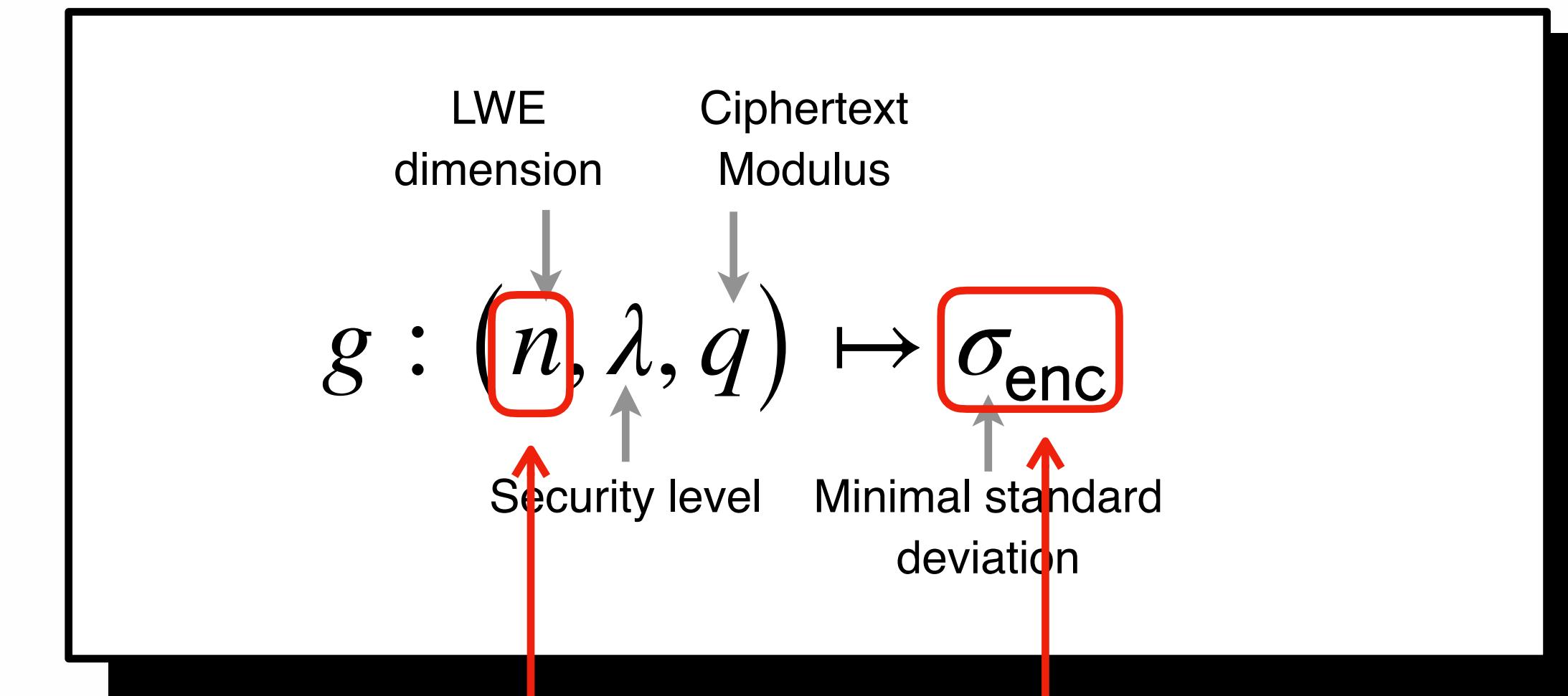
**Degree of
freedom**

Security



Lattice Estimator [APS]

Reference tool to estimate the security of lattice-based encryption scheme [APS]



Degree of freedom

Fully defined by g and n

Efficiency



Efficiency



Rule of thumb

The smaller the parameters,
the faster the computation

Efficiency

Rule of thumb

The smaller the parameters,
the faster the computation



How can we rank the
feasible parameters ?

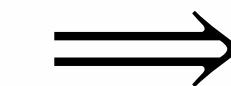
Efficiency

Rule of thumb

The smaller the parameters,
the faster the computation



How can we rank the
feasible parameters ?



Cost Model as a surrogate of the execution time
Ex: algorithmic complexities

Correctness 

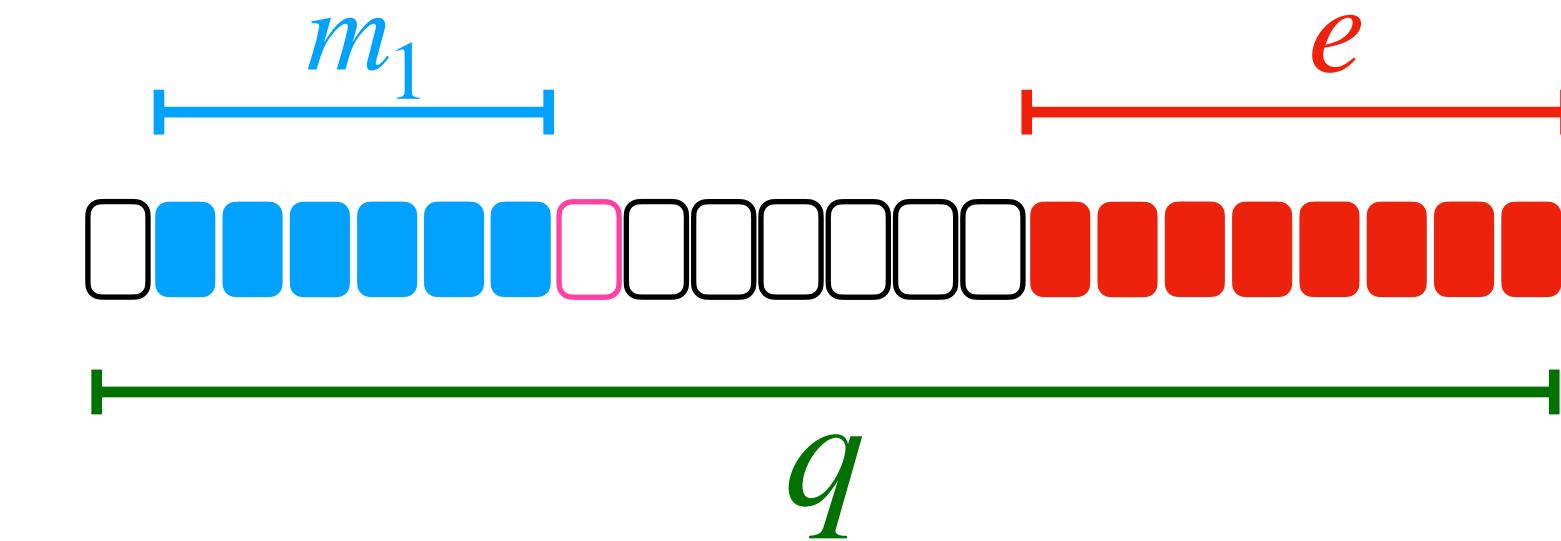
Correctness 

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2}$$



Correctness

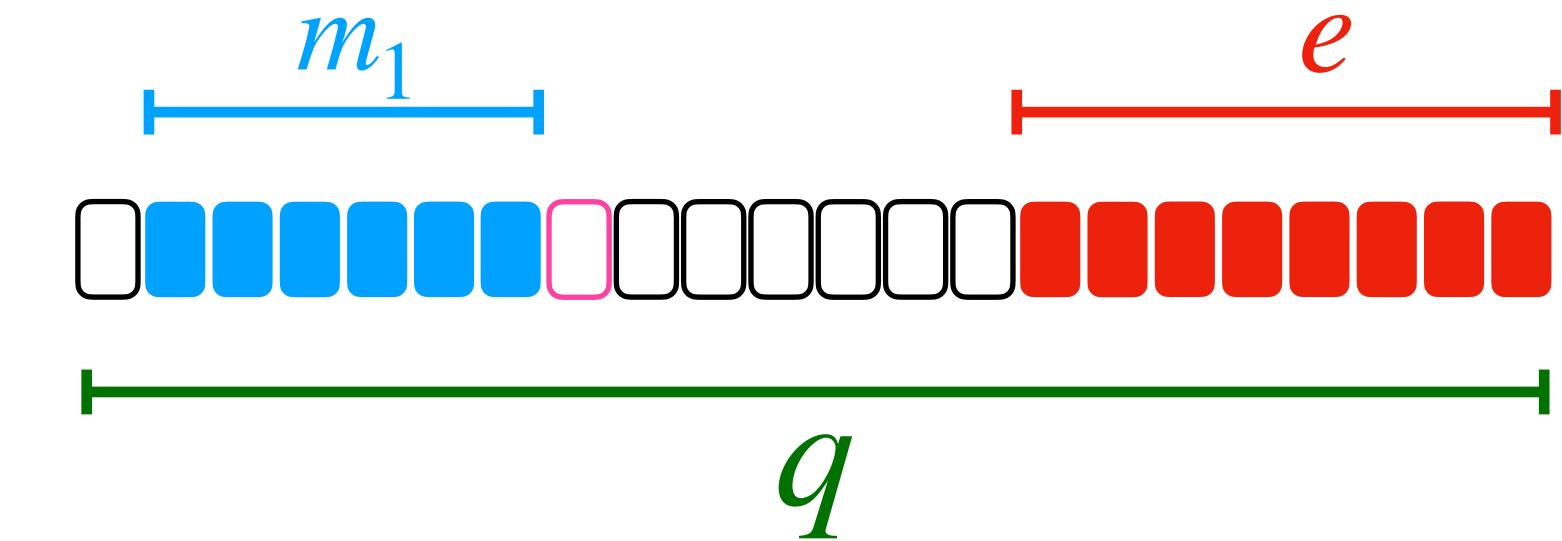


e can not be publicly known without compromising the security



Correctness Condition

$$|e| < \frac{\Delta}{2}$$



Correctness



e can not be publicly known without compromising the security

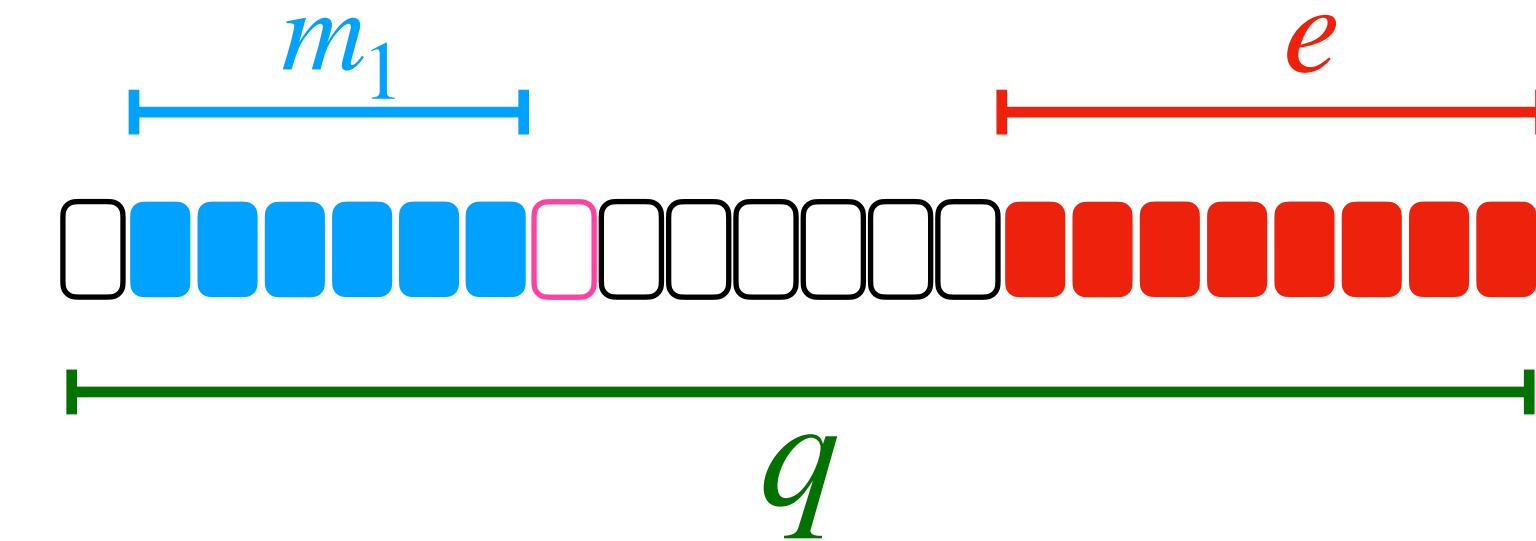


The distribution of e can be publicly known without compromising the security



Correctness Condition

$$|e| < \frac{\Delta}{2}$$



Correctness



e can not be publicly known without compromising the security

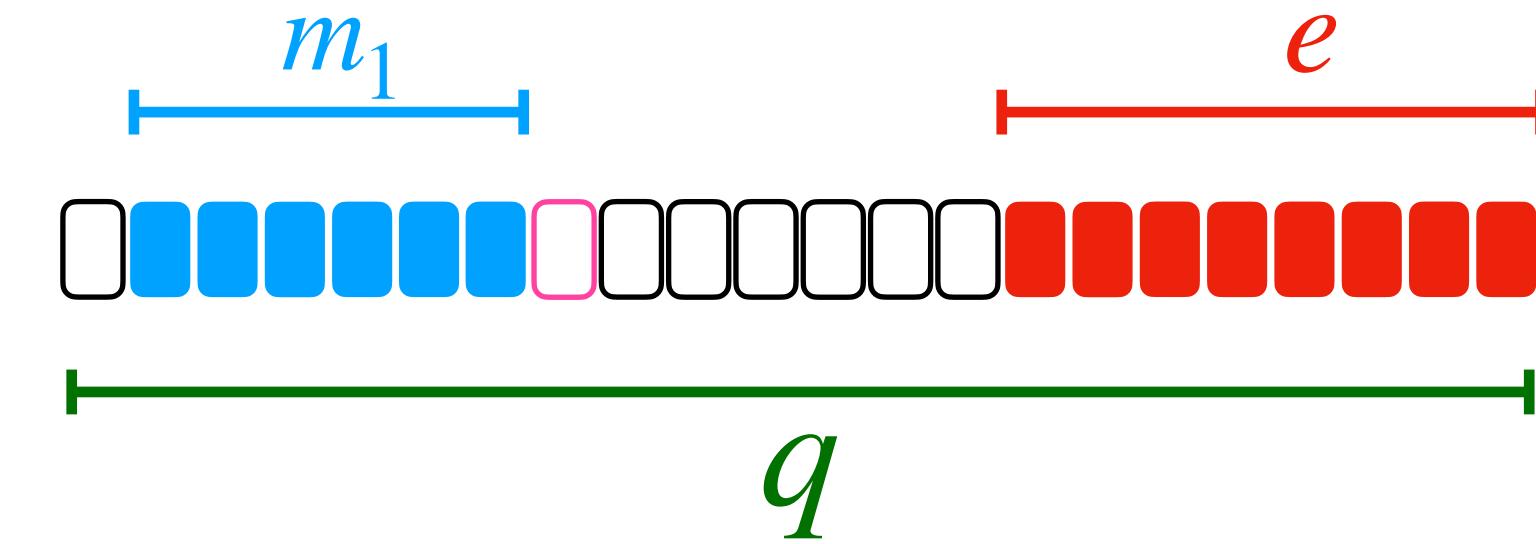


The distribution of e can be publicly known without compromising the security



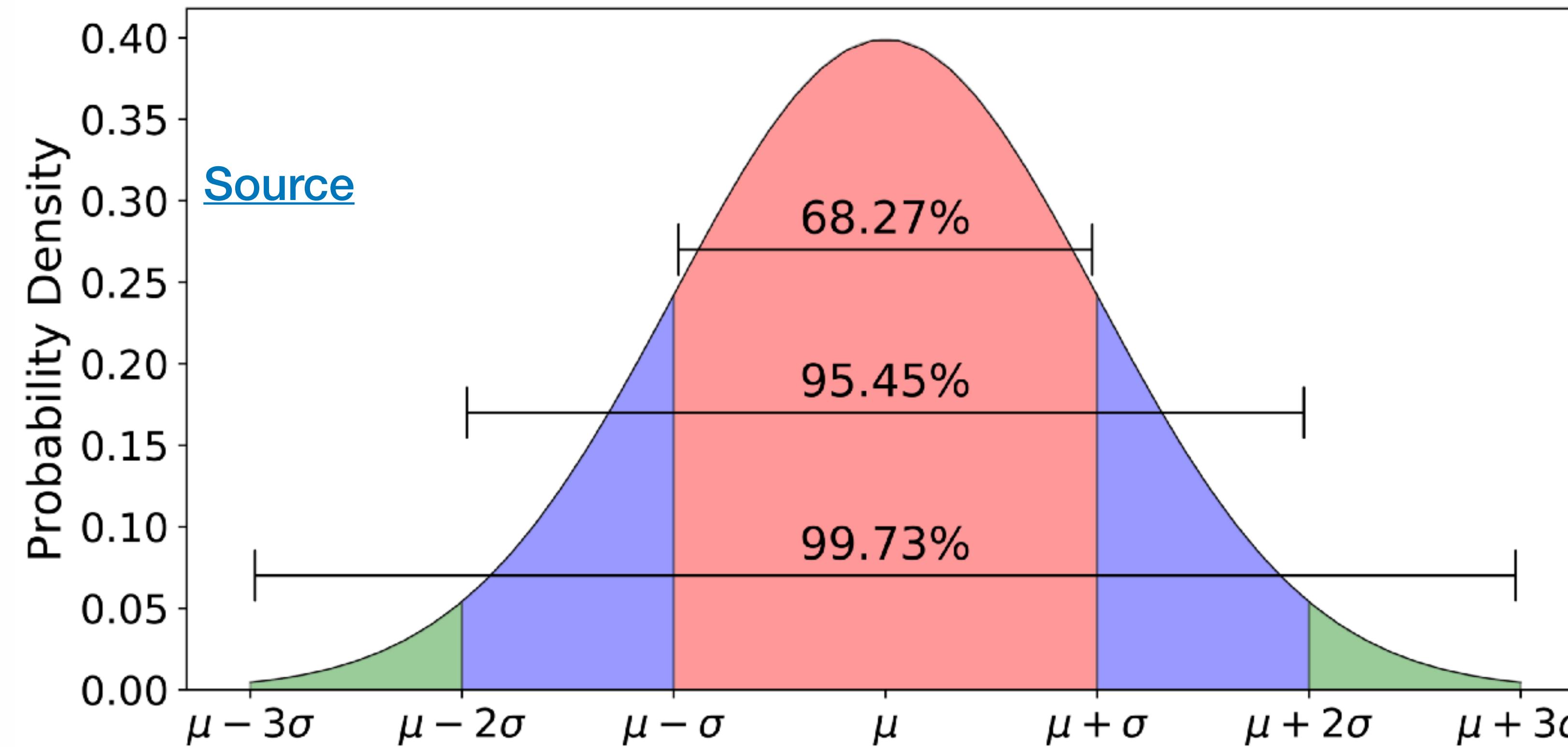
Correctness Condition

$$|e| < \frac{\Delta}{2}$$



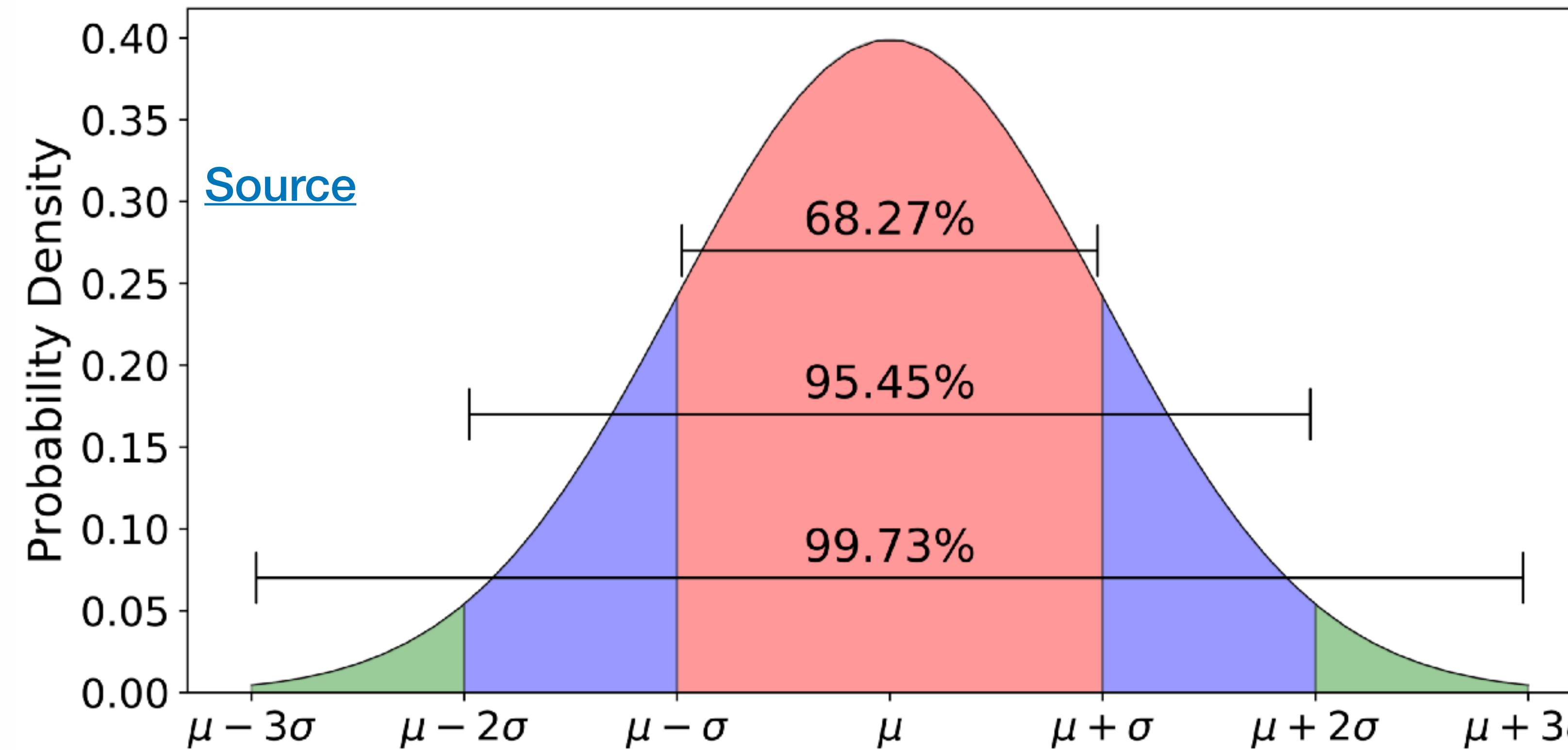
⇒ **Noise Model** to track the noise distribution throughout the computation

Correctness ✓



e is drawn from $\mathcal{N}(\mu = 0, \sigma^2)$

Correctness ✓



e is drawn from $\mathcal{N}(\mu = 0, \sigma^2)$

⇒

$\Pr [e \notin [-\kappa \cdot \sigma, \kappa \cdot \sigma]] = p_{\text{fail}}$

Correctness



Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2}$$

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$



Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$

↔

$$\sigma \leq \frac{\Delta}{2 \cdot \kappa}$$

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$

|*e*|



Noise
(private)



$$\sigma \leq \frac{\Delta}{2 \cdot \kappa}$$

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$

↑
Noise
(private)

↔

$$\sigma \leq \frac{\Delta}{2 \cdot \kappa}$$

Noise Distribution
(public)

Correctness



Correctness Condition

$$|e| < \frac{\Delta}{2} \quad \text{with probability } 1 - p_{\text{fail}}$$

Noise
(private)



Noise bound

$$\sigma \leq \frac{\Delta}{2 \cdot \kappa}$$

**Noise Distribution
(public)**

Optimization Problem

\min **Cost**

Optimization Problem

Optimization Problem

$$\min \quad \text{Cost} \quad \mathcal{G} \quad \text{s.t.} \quad \left\{ \begin{array}{l} \forall i \in I, \sigma_i \leq \frac{\Delta}{2 \cdot \kappa} \end{array} \right.$$

Optimization Problem

$$\min \quad \text{Cost} \quad \mathcal{G} \quad \text{s.t.} \quad \left\{ \begin{array}{l} \forall i \in I, \sigma_i \leq \frac{\Delta}{2 \cdot \kappa} \\ \sigma_{\text{enc}} = g(n, \lambda, q) \end{array} \right.$$

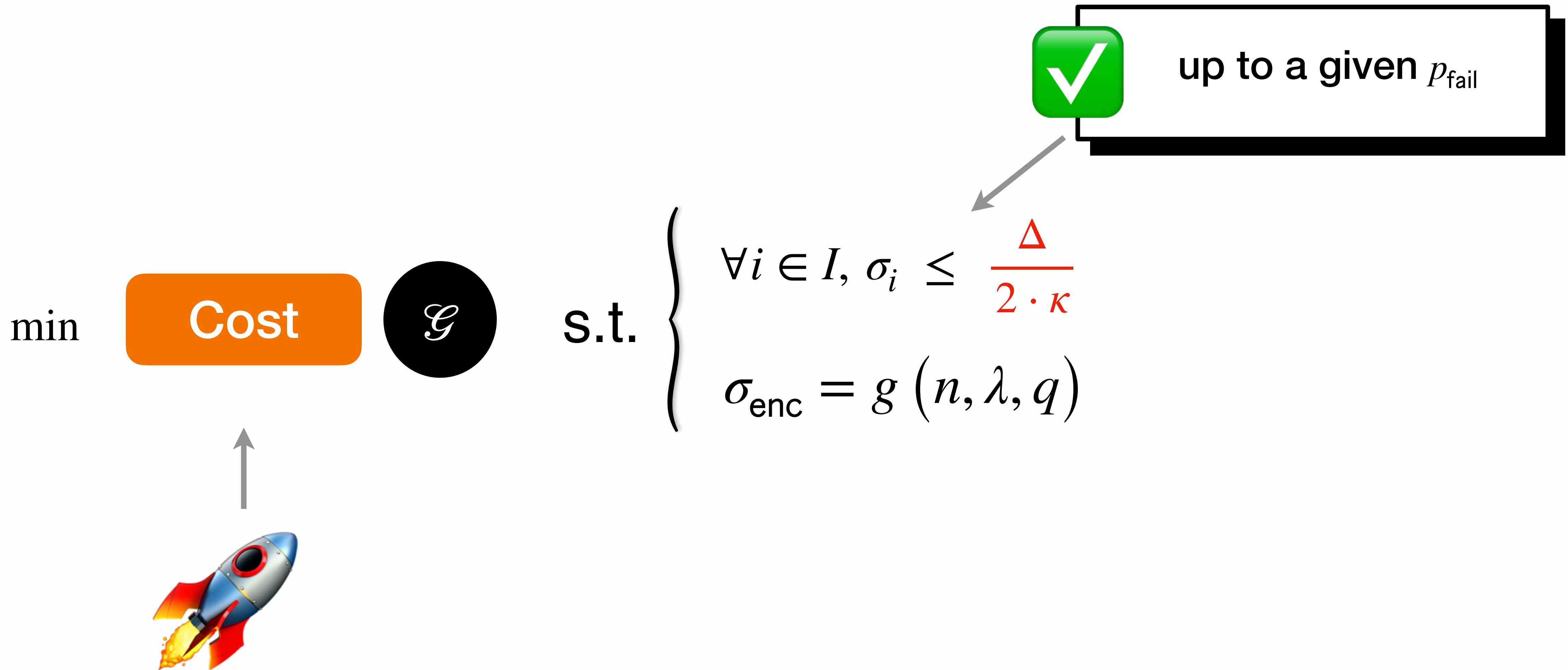
Optimization Problem

$$\min \quad \text{Cost} \quad \mathcal{G}$$

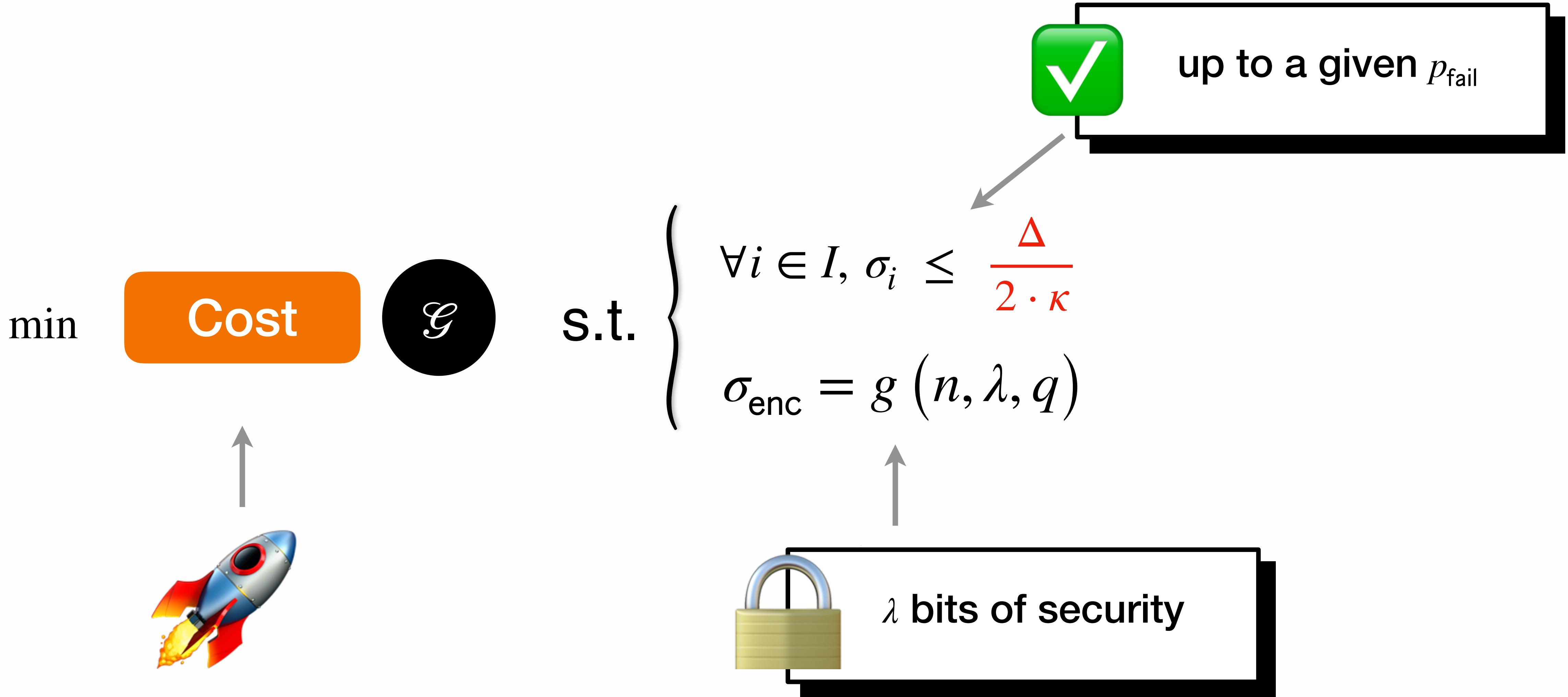
s.t. $\left\{ \begin{array}{l} \forall i \in I, \sigma_i \leq \frac{\Delta}{2 \cdot \kappa} \\ \sigma_{\text{enc}} = g(n, \lambda, q) \end{array} \right.$



Optimization Problem



Optimization Problem



Conclusion

What's next ?

Translation

More than one possible translation: how to find the best one ?

Multi Parameter Set

One parameter set
vs several parameter set

Failure Probability

At operator level vs at graph level

Optimization Problem

Simplification with domination reasoning

Part II

How to translate an arbitrary DAG into an FHE DAG i.e., a DAG only composed of
linear operations and  ?

How to convert a TFHE DAG into a binary executable ?

Contact and Links

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github.com/zama-ai/concrete

community.zama.ai

ZAMA