

Chapter 11 Cryptology

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Computing Modular Powers

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RSA Cryptosystem

Composite and Prime Numbers

Composite Numbers have a divisor other than itself and one.

For example $4|20$ means that $20 = 5 * 4$

The divisors of 12 are 1,2,3,4,6 and 12

Prime numbers have no divisors but 1 and itself

First 10 Primes

2, 3, 5, 7, 11, 13, 17, 19, 23, 29

Greatest Common Divisor

If $h|m$ and $h|n$ then h is called a common divisor

A common divisor is a number that is a factor of both numbers

The greatest common divisor is the largest factor for both numbers

This is denoted $\gcd(n,m)$

For example $\gcd(12,15) = 3$

For any two integers n and m where $m \neq 0$ the quotient is n divided by m is given by

$$q = \lfloor n/m \rfloor$$

The remainder r of dividing n by m is given by

$$r = n - qm$$

Greatest Common Divisor (cont)

Let n and m be integers, not both 0 and let

$$d = \min \{ in + jm \text{ such that } i, j \in \mathbb{Z} \text{ and } in + jm > 0 \}$$

That is, d is the smallest positive linear combination of n and m

For example we know $\gcd(12, 8) = 4$,

the smallest linear combination is

$$4 = 3(12) + (-4)8$$

Now suppose we have $n \geq 0$ and $m > 0$ and $r = n \bmod(m)$ then

$$\gcd(n, m) = \gcd(m, r)$$

$$\text{so } \gcd(64, 24) = \gcd(24, 16)$$

$$= \gcd(16, 8)$$

$$= \gcd(8, 0)$$

$$= 8$$

Least Common Multiple

For n and m where they are both nonzero, the least common multiple is denoted $\text{lcm}(n,m)$

For example $\text{lcm}(6,9) = 18$ because $6|18$ and $9|18$

The $\text{lcm}(n,m)$ is a product of primes that are common to m and n , where the power of each prime in the product is the larger of its orders in n and m

So $12 = 2^2 3^1$ and $45 = 3^2 5^1$

so $\text{lcm}(12,45) = 2^2 3^2 5^1 = 180$

Prime Factorization

Two integers are relatively prime because the gcd of them is 1

For example $\gcd(12, 25) = 1$ so they are relatively prime

If h and m are relatively prime and h divides nm , then h divides m .

That is $\gcd(h, m) = 1$ and $h \mid nm$ implies $h \mid n$

Prime Factorization (cont)

Every integer $X > 1$ can be written as a unique product of primes

That is $X = p_1^{k_1} * p_2^{k_2} * \dots * p_n^{k_n}$

Where $p_1 < p_2 < \dots < p_n$ and this representation of n is unique

Example being $22,275 = 3^4 * 5^2 * 11$

To solve $\gcd(3,185,325, 7,276,500)$ we know

$$3,185,325 = 3^4 5^2 11^2 13^1$$

$$7,276,500 = 2^2 3^3 5^3 7^2 11^1$$

We then take the common divisors and take the lower power to create the gcd

$$\text{so } \gcd(3,185,325, 7,276,500) = 3^3 5^2 11^1 = 7,425$$

Euclid's Algorithm

Euclid's Algorithm gives us a straight forward way to find the gcd of two numbers

```
int gcd(int n, int m)
{
    if(m == 0)
        return n;

    else
        return gcd(m, n mod m);
}
```

Extension to Euclid's Algorithm

```
void Euclid (int n, int m, int gcd, int i, int j){  
    if (m == 0) {  
        gcd = n; i = 1; j = 0;  
    }  
    else {  
        int iprime, jprime, gcdprime;  
        Euclid (m, n mod m, gcdprime, iprime, jprime);  
        gcd = gcdprime;  
        i = jprime;  
        j = iprime -  $\lfloor n/m \rfloor$  jprime ;  
    }  
}
```

Why Use the Other Algorithm?

This other algorithm will give us integers i and j as well

So, $\text{gcd} = in + jm$

For Example $\text{Euclid}(42, 30, \text{gcd}, i, j)$ outputs

$\text{gcd} = 6, i = -2$ and $j = 3$

$$6 = -2(42) + 3(30)$$

Proof Extended Algorithm

Induction Base: In the last recursive call $m = 0$, which means $\gcd(n, m) = n$

Since the values of i and j are assigned values 1 and - respectively we have

$$in + jm = 1n + 0m = n = \gcd(n, m)$$

Induction Hypothesis: Assume in the k th recursive call the values determined for i and j are such that

$$\gcd(n, m) = in + mj$$

Then the values returned by that call for i' and j' are values such that

$$\gcd(m, n \bmod m) = i'm + j'n \bmod m$$

Proof Extended Algorithm (cont)

Induction Step: We have for the $(k - 1)$ st call that

$$\begin{aligned}in + mmj &= j'n + (i' - \lfloor n/m \rfloor j')m \\&= i'm + j'(n - \lfloor n/m \rfloor m) \\&= i'm + j'n \bmod m \\&= \gcd(m, n \bmod m) \\&= \gcd(n, m)\end{aligned}$$

The second to last equality is due to the induction hypothesis

Group Theory

A closed binary operation $*$ on a set S is a rule for combining two elements of S to yield another element of S .

This operation is associative, has an identity element and inverse element

For example with integers $\in \mathbb{Z}$ with addition constitute a group.

The identity element is 0 and the inverse of a is $-a$

Every element in the group only has one inverse

A group is said to be finite if S contains a finite number of elements

A group is said to be commutative (or abelian) if for all $a, b \in S$,

Congruency Modulo n

Let m and k be integers and n be a positive integer. If $n|(m - k)$ we say m is congruent to k modulo n , and this is written by

$$m \equiv k \pmod{n}$$

For Example

Since $5|(33 - 18)$, $33 \equiv 18 \pmod{5}$

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