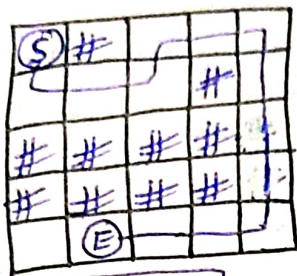
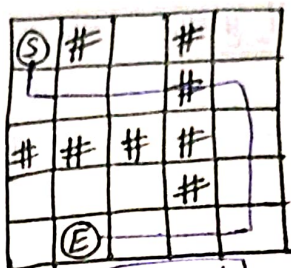


Graph Formulation 1

1. Given a grid with starting point and the ending point. Find:
 (a) Find the min. no of walls you need to break to go from pt S to E.



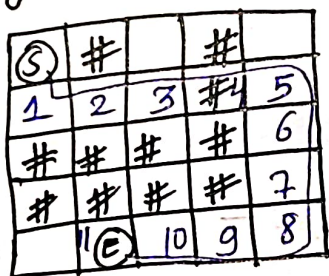
Ans = 0
 No walls to break



Ans = 1
 1 wall needed to break

(b) If you can break $\leq k$ walls, Find min moves needed.

Ex:-



$k = 1$

Ans = 11

(c) Find all ans for (b) for k in range: $1 \leq k \leq N \cdot M$

Constraints: $N, M \leq 100$.

(a) While traversing the cell to go from $S \rightarrow E$:

1. If cell is empty, then take cost = 0

2. If cell has a wall, then take cost = 1.

* Keep track of broken walls as when tracking these walls again will make cost = 0.

→ given code below: for printing path from S to E



(b) We have to minimise the no. of moves

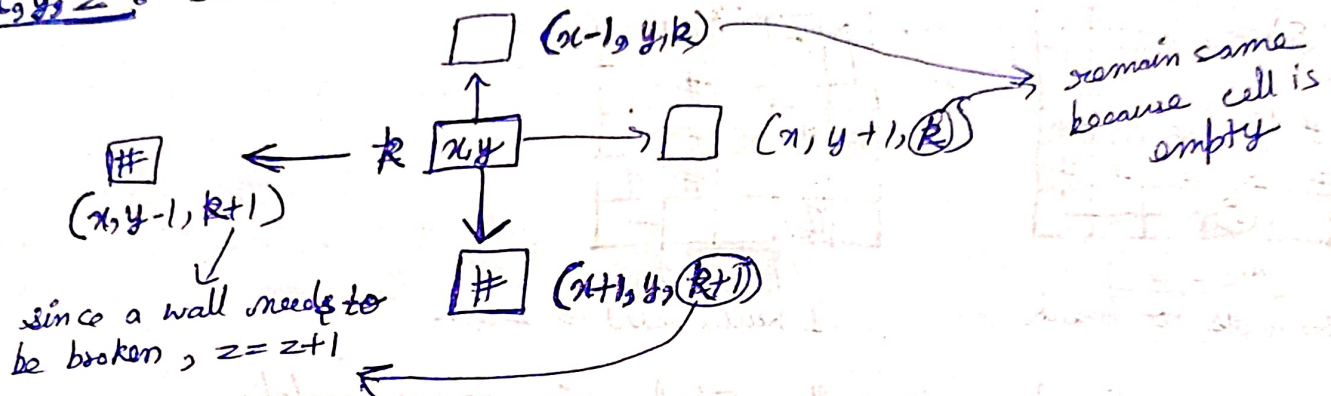
No. of walls broken $\leq k \rightarrow \text{Node}(v)$

Node is now (x, y, z) . It will store

\rightarrow Min cost to (x, y) with z wall broken : This is our state.

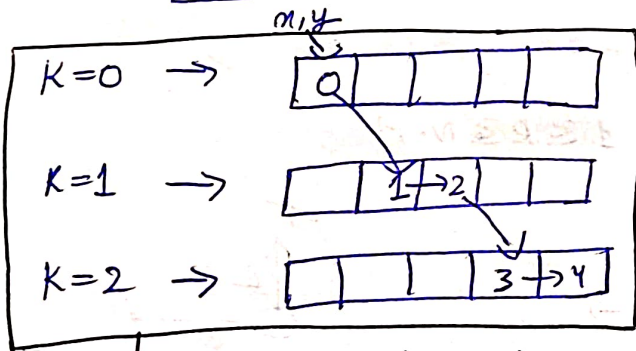
Our real graph is now 3-D graph.

x, y, z : state at which we will run BFS Algo



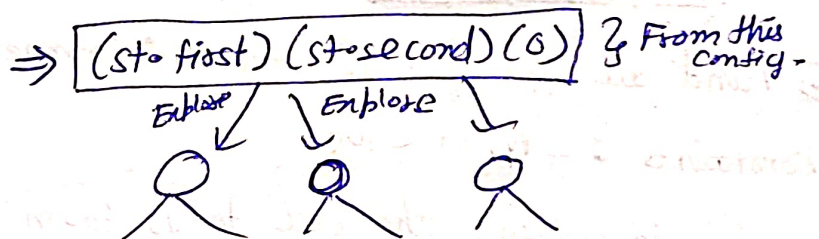
Suppose,

S	#	#	E
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This is the type of graph we're currently moving in.

Starting config :-



$E \leq k$

After end point, $\leq k$ how many moves we can search

(c) Max. value of $K = N * M$

Is $N * M$ possible? No as for even the farthest pt, we can cover $(N+M)$.

Use K constraint till $N+M$.

Code for (b) & (c) in live session chat