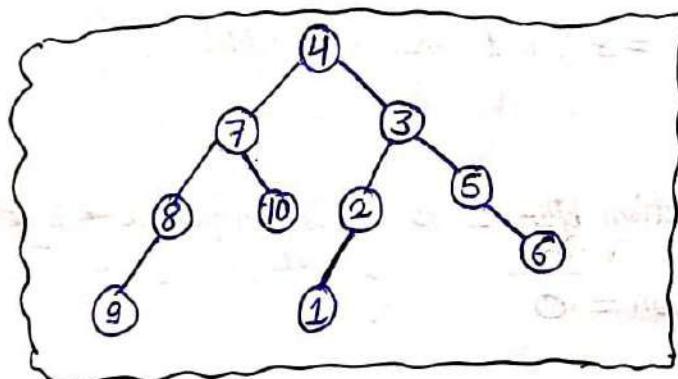
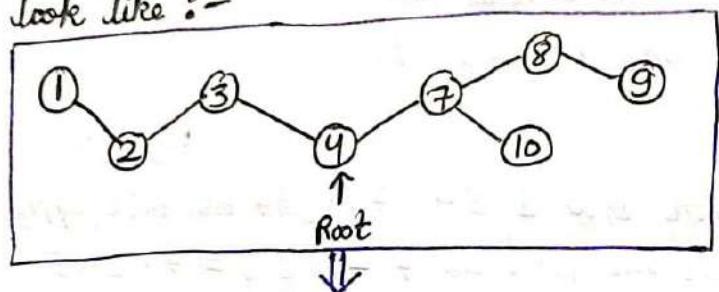


## Tree Foundations & Framework

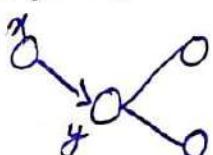
\* Trees → a general modification of graph which has no cycle and it is connected.

How trees look like :-



All concepts like BFS/DFS is applicable to trees. We will mostly use DFS instead of BFS because here in trees, there is no concept of shortest path but there is a concept of unique path.

Given nodes  $n$  &  $y$ , exploring  $y$ 's neighbours :



```
visited [y] = 1  
for (auto u : g[y])  
    if (visited [u] != 1)  
        DFS (u)
```

When performing DFS on tree you don't need a visited array. Because there are no cycles, you can't accidentally revisit a node.

∴  $\text{visited}[y] = 1$

```
for (auto u : g[y])
    if ( $u \neq \text{parent}[y]$ ) dfs(u)
```

DFS in Tree code :-

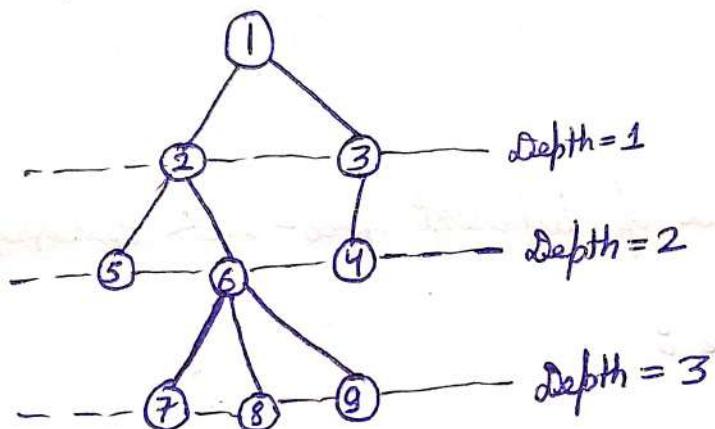
```
#include<bits/stdc++.h>
using namespace std;

int n;
vector<vector<int>> g;

void dfs(int nn, int pp) {
    cout << nn << endl;
    for (auto v : g[nn]) {
        if (v != pp) {
            dfs(v, nn);
        }
    }
}

void solve() {
    cin >> n;
    g.resize(n + 1);
    for (int i = 0; i < n - 1; i++) {
        int a, b;
        cin >> a >> b;
        g[a].push_back(b);
        g[b].push_back(a);
    }
    dfs(1, 0);
}
```

\* Find distance to all nodes from 1.



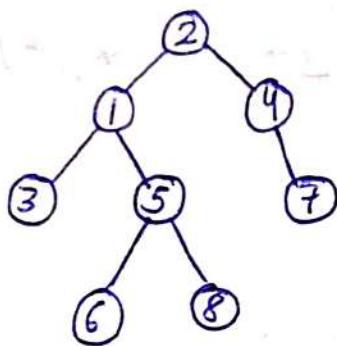
```
int n;
vector<vector<int>> g;
vector<int> depth;

void dfs(int nn, int pp, int dd) {
    depth[nn] = dd;
    for (auto v : g[nn]) {
        if (v != pp) {
            dfs(v, nn, dd + 1);
        }
    }
}

void solve() {
    cin >> n;
    g.resize(n + 1);
    depth.resize(n + 1);
    for (int i = 0; i < n - 1; i++) {
        int a, b;
        cin >> a >> b;
        g[a].push_back(b);
        g[b].push_back(a);
    }
    dfs(1, 0, 0);
}
```

• Subtree :-

For any particular node, it has child or its subchild.



Q: Given a tree. For every node : calculate no. of nodes in its subtree.

dfs( $x, pp$ )

$sub[x] = 1$

for ( $v : g[x]$ )

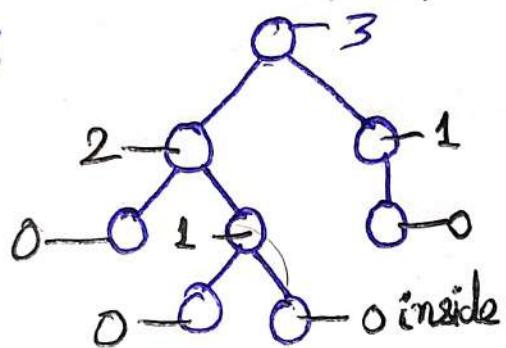
if ( $v \neq pp$ ) ~~dfs( $v, nn$ )~~

$dfs(v, nn)$

$sub[x] += sub[v]$

Q: Calculate how far, deep inside there is a node?

Ans:



```
#include<bits/stdc++.h>
using namespace std;

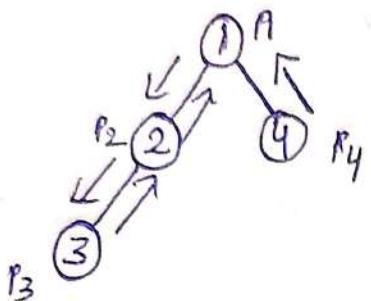
int n;
vector<vector<int>> g;
// ancestor
vector<int> depth,par;
// subtree
vector<int> subsz,subfar;

void dfs(int nn,int pp,int dd){
    depth[nn] = dd;
    par[nn] = pp;

    subsz[nn] = 1;
    subfar[nn] = 0;

    for(auto v:g[nn]){
        if(v!=pp){
            dfs(v,nn,dd+1);
            subsz[nn] += subsz[v];
            subfar[nn] = max(subfar[nn],1+subfar[v]);
        }
    }
}
```

- Pre Order traversal vs Post



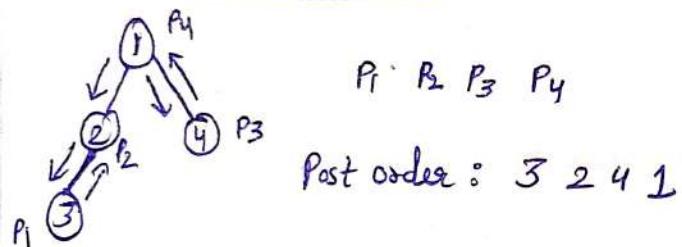
Pre Order : 1 2 3 4

OR

Pre Order : 1, 4, 2, 3

\* Preorder traversal of a tree is not unique.

### Post order Traversal :-



P1 P2 P3 P4

Post order : 3 2 4 1

```
int n;
vector<vector<int>> g;
```

```
void dfs(int nn, int app) {
    cout << nn << endl;
    for (auto v : g[nn]) {
        if (v == app) {
            cout << v << endl;
        } else {
            dfs(v, nn);
        }
    }
}
```

```
using namespace std;
int n;
vector<vector<int>> g;
vector<int> ans;
int a, b;
vector<int> c;
int main() {
    cin >> n;
    for (int i = 0; i < n - 1; i++) {
        int u, v;
        cin >> u >> v;
        g[u].push_back(v);
        g[v].push_back(u);
    }
    cin >> a >> b;
    dfs(a, b);
    cout << ans[0] << endl;
    for (int i = 1; i < ans.size(); i++) {
        cout << " " << ans[i];
    }
}
```

```
int n;
vector<vector<int>> g;

void dfs(int nn,int pp=1)
{
    cout<<n<<endl;
    if (v[pp])
        dfs(v,nn);
}

cout<<n<<endl;
}

void solve(){
    cin>>n;
    g.resize(n+1);
    for (int i=0;i<n-1;i++){
        int a,b;
        cin>>a>>b;
        g[a].push_back(b);
        g[b].push_back(a);
    }
    dfs(1,0);
}
```

Q: Given a tree, find out the number of different pre-order traversals poss?

When we are deciding for a node:

a → look at levels

b → in which order you'll visit them

c → order inside subtree.

For c,

$$\# \text{ of ways } (\text{child 1}) = 1$$

$$\text{ " " " } (\text{child 2}) = 2$$

$$\text{ " " " } (\text{child 3}) = 3$$

⋮

Depending on b,

$$(\# \text{ child 1}) \times (1 * 2 * 3 * \dots * k^{\text{th}} \text{ child ways})$$

```
int fact[100101];
int n;
vector<vector<int>> g;
vector<int> ways;
void dfs(int nn, int pp) {
    ways[nn] = 1;
    int child = 0;

    for (auto v : g[nn]) {
        if (v != pp) {
            dfs(v, nn);
            ways[nn] *= ways[v];
            child++;
        }
    }
    ways[nn] *= fact[child_count];
}
```

## Frameworks :-

### 1. Foundation (Basic properties)

- DFS
- subtree
- Ancestors

### 2. Diameter

Nodes with max distance in a tree determines its diameter.

### 3. Center

Mid point of node of diameter's path. There can be more than one center.

### 4

#### • Centroid

A node whose all subtrees size is less than half of the whole tree's size.

### 5. Contribution Technique

#### Edge

How to find diameter?

Diameter is the longest possible path between any two nodes in the tree.

1. Pick any random node (say node A) and run a DFS to find the node farthest from it. (B).
2. Now, run a second DFS starting from node B to find the node farthest from it (C).
3. ~~A & B are di~~ The path between B and C is the diameter of the tree.

#### • More on Frameworks

1. Ancestral Maintenance — Data structures on Ancestors
2. Small to Large Merging / BSU on sack — Data structures on subtrees
3. DS on Paths — Binary lifting, LCA finding

```
#include<bits/stdc++.h>
using namespace std;

int n;
vector<vector<int>> g;
// ancestor
vector<int> depth;

void dfs(int nn,int pp,int dd){
    depth[nn] = dd;
    for(auto v:g(nn)){
        if(v!=pp){
            dfs(v,nn,dd+1);
        }
    }
}

void solve(){
    cin>>n;
    g.resize(n+1);
    depth.resize(n+1);

    for(int i=0;i<n-1;i++){
        int a,b;
        cin>>a>>b;
        g[a].push_back(b);
        g[b].push_back(a);
    }

    dfs(1,0,0);

    int x = 1;
    for(int i=1;i<=n;i++){
        if(depth[i]>depth[x])x=i;
    }

    dfs(x,0,0);

    int y = 1;
    for(int i=1;i<=n;i++){
        if(depth[i]>depth[y])y=i;
    }

    cout<<x<<y<<endl;
}
```