

Segment Tree Applications

1. Given an array of n elements. You are given a query from L to R of value K . Find # of elements $\geq K \rightarrow [L, R]$

$$N \leq 10^5$$

$$Q \leq 10^5$$

$$arr[i], K \leq 10^5$$

Approach :-

1. Sort array elements in descending order of value.
2. Sort queries in descending order of K .
3. Tree stores 1 at index i if $arr[i] > \text{current } K$, otherwise 0.
4. So, range sum give count of elements $> K$.

```
int n;  
int arr[100101];  
  
int t[400400];  
void build(int id, int l, int r){  
    if(l==r){  
        t[id]=1;  
        return;  
    }  
    int mid = (l+r)/2;  
    build(id<<1, l, mid);  
    build(id<<1|1, mid+1, r);  
    t[id] = t[id<<1] + t[id<<1|1];  
}
```

```
void update(int id, int l, int r, int pos, int val) {  
    if (pos > r || pos < l) return;  
    if (l == r) {  
        t[id] = val;  
        return;  
    }  
    int mid = (l + r) / 2;  
    update(id << 1, l, mid, pos, val);  
    update(id << 1 | 1, mid + 1, r, pos, val);  
    t[id] = t[id << 1] + t[id << 1 | 1];  
}
```

```
int range_sum(int id,int l,int r,int lq,int rq){  
    if(l>rq||lq>r)return 0;  
    if(lq<=l&&r<=rq){  
        return t[id];  
    }  
    int mid = (l+r)/2;  
    return range_sum(id<<1,l,mid,lq,rq) + range_sum(id<<1|1,mid+1,r,lq,rq);  
}
```

```
using query = pair<pair<int,int>,pair<int,int>>;
#define F first
#define S second;
```

```
void solve(){
    cin>>n;
    for(int i=0;i<n;i++){
        cin>>arr[i];
    }

    int q;
    cin>>q;
    vector<query> allq;
    for(int i=0;i<q;i++){
        int l,r,k;
        cin>>l>>r>>k;
        l--;r--;
        allq.push_back({{k,i},{l,r}});
    }
```

```
cin>>n;
vector<pair<int,int>> vals;
for(int i=0;i<n;i++){
    cin>>arr[i];
    vals.push_back({arr[i],i});
}
build(1,0,n-1);
int cur = 0;
sort(vals.begin(),vals.end());

int q;
cin>>q;
vector<query> allq;
for(int i=0;i<q;i++){
    int l,r,k;
    cin>>l>>r>>k;
    l--;r--;
    allq.push_back({{k,i},{l,r}});
}
sort(allq.begin(),allq.end());
```

```
sort(allq.begin(),allq.end());
int ans[q];
for(int i=0;i<q;i++){
    while(cur<n && vals[cur].F <= allq[i].F.F){
        update(1,0,n-1,vals[cur].S,0);
        cur++;
    }
    ans[allq[i].F.S] = range_sum(1,0,n-1,allq[i].S.F,
    allq[i].S.S);
}
for(int i=0;i<q;i++){
    cout<<ans[i]<<endl;
}
}
```

Q. K-Query online

Build a segment Tree where each node stores a sorted list of its range, and answer each query by combining binary search with segment tree traversal to count elements greater than K


```

vector<int> t[400400];
void build(int id,int l,int r){
    if(l==r){
        t[id]={arr[l]};
        return;
    }
    int mid = (l+r)/2;
    build(id<<1,l,mid);
    build(id<<1|1,mid+1,r);
    t[id].resize(r-l+1);
    merge(t[id<<1].begin(),t[id<<1].end(),
        t[id<<1|1].begin(),t[id<<1|1].end(),
        t[id].begin());
}

int query(int id,int l,int r,int lq,int rq,int k){
    if(l>rq||lq>r)return 0;
    if(lq<=l&&r<=rq){

```

```

int query(int id,int l,int r,int lq,int rq,int k){
    if(l>rq||lq>r)return 0;
    if(lq<=l&&r<=rq){
        return t[id].end()- upper_bound(t[id].begin(),t[id].
            end(),k);
    }
    int mid = (l+r)/2;
    return query(id<<1,l,mid,lq,rq,k) + query(id<<1|1,mid+1,
        r,lq,rq,k);
}

```

```

void solve()
{
    cin >> n;
    for (int i = 0; i < n; i++) {
        cin >> arr[i];
    }
    build(1, 0, n - 1);

    int q;
    cin >> q;
    int last_ans = 0;
    for (int i = 0; i < q; i++) {
        int l, r, k;
        cin >> l >> r >> k;
        l ^= last_ans;
        r ^= last_ans;
        k ^= last_ans;
        l--; r--;
        last_ans = query(1, 0, n - 1, l, r, k);
        cout << last_ans << endl;
    }
}

```