

Context Discussion

1. String should have \rightarrow atleast 1 vowel & 1 consonant

$S = "BCW"$ X If length = 1 \rightarrow 0 is any

'A' X

'ABCE' ✓

* Print all valid strings of length $n \geq 2$

Use Recursion

1. We will go to ^{each} index and by brute force will make all strings and in last we will check if it is valid or not (1 vowel + 1 consonant). If yes \rightarrow Print string.

Optimal :-

All consonant $\rightarrow 21^N$ } Not all consonants or & not all vowels
 All vowels $\rightarrow 5^N$ } needed.

$$\therefore \text{ans} = 26^N - (21^N + 5^N)$$

Whenever constraints is greater than $10^9 \rightarrow$ use smathy

2:

$$\begin{matrix} a \\ b & c \\ d & e & f \end{matrix} \quad a < b \& a < c$$

$$d < e < f$$

1) Sort the array given.

2) Now place the smallest element at starting and then keep the rest at b, c, d, e and f posⁿ.

3) Now check the condⁿ.

3:

$$[1, 2, 3, 2, 1, 4]$$

$$\text{Size} = 3$$

$$\text{Operation, } M = 5$$

$$[2 \underline{3} \ 4 \ \underline{2 \ 1 \ 4}]$$

$$[\underline{2 \ 2 \ 3} \ 3 \ 2 \ 5]$$

$$\text{Max - Min value in array} = ?$$

$$[3 \ 4 \ 5 \ \underline{3 \ 2 \ 5}]$$

$$[\underline{3 \ 4 \ 5} \ 4 \ 3 \ 6]$$

$$[4 \ 5 \ 6 \ \underline{4 \ 3 \ 6}]$$

If not 4 we can make it (4)

if greater than we can make in right side

$$-\ - \underline{1 \ 2 \ 3} \ 4 - -$$

$\therefore \text{high} = \text{mid} + 1 \quad // \text{go left}$ $\text{low} = \text{mid} + 1 \quad // \text{go to right}$

check if can bring ③

$$[1 \ 2 \ 3 \ 2 \ 1 \ 4] \Rightarrow [3 \ 4 \ 5 \ 2 \ 1 \ 4]$$

$$[3 \ 4 \ \cancel{5} \ 2 \ 1 \ 4] \xrightarrow{(5>3)} \Rightarrow [3 \ 4 \ 5 \ 3.2 \ 5]$$

$$A [3 \ 4 \ 5 \ 3 \ \frac{+1}{2} \ \frac{+1}{5}] \Rightarrow [3 \ 4 \ 5 \ 3 \ 3 \ 6] \cancel{x_{(6>3)}}$$

Pseudocode :-

```

check(x) {
    int step = 0;
    for(i=0; i<n; i++) {
        if(a[i] < x) {
            int add = x - a[i], step += add;
            for(j=i; j<i+step; j++) a[j] += add;
        }
        if(step <= M) return 1;
    }
    return 0;
}
  
```

$$T.C : O(n^2)$$

Optimal Approach :- Using Prefix sum

Instead of updating the array a directly at every step, it updates a range indirectly by making increments at start and decrements off at the end.

Code :-

```

int n;
vector<int> s;
int check(int n) {
    vector<int> partial(n, 0);
    for(int i=0; i<n; i++) {
        if(i>0) partial[i] += partial[i-1];
        a[i] += partial[i];
        if(a[i] < n) {
            int add = n - a[i];
            partial[i] += add;
            if(i+m < n) partial[i+m] -= add;
            a[i] += add;
        }
    }
}
  
```

$\Rightarrow [1 \ 2 \ 3 \ 3 \ 2 \ 5]$

Subarray \rightarrow Two pointers, sliding window

subset \Rightarrow If i will sort \rightarrow then would i will get ans

$[1 \ 2 \ 2 \ 3 \ 3 \ 5]$
↑ ↓
 $i+s$

$\Rightarrow [1 - 4]$

$[6 - 12]$

$[19 - 20]$

$P_1 \ P_2 \ P_3 \ P_4 \ P_5$
1, 4, 2, 6, 10
 3 2 4 9

$\boxed{Z = 2}$ (min closest pair value)

We have to maximize it

If smaller, then left If bigger, go right

Binary search:
— — — X — —

Check (5)

$P_1 \ P_2 \ P_3 \ P_4 \ P_5$
1 6 11 19 X

] If min dist (Z) = 5 \rightarrow Not psbl.