

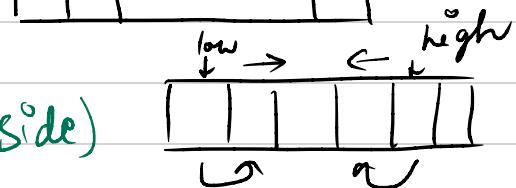
August 10, 2025

## TWO POINTERS FORM 1

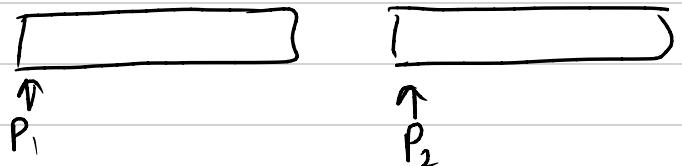
Form 0 : Sliding Window (Ex: Monotone Deque ques in STL)



Form 1 : Variable Window



Form 2 : 3Sum Type (Opposite side)



Form 3 : Multi-Sequence

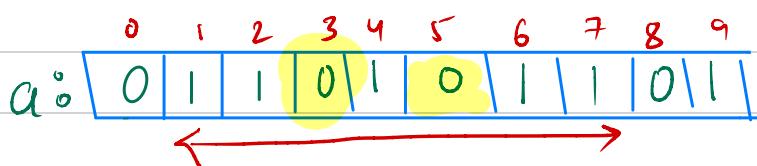
\* Form 1, what kind of problems can we solve using form 1?

For subarrays : Finding max length  
Finding min length  
Count of subarray  
Sum of length

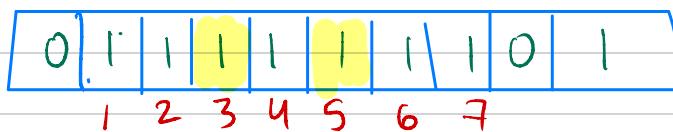
→ (given some condition  
that needs to be  
satisfied)

Ques 1. Given an array of 0's and 1's. Find the max length of 1's that can be created by flipping at most K positions with 0's.

N=10, k=2



For  $k=2$ ,  
that means we  
can flip at most  
2 zeroes in our  
array.



Answer = 7

We will be using form 1 to solve this question. For that, first we have to find the condition that should be satisfied at all times:

Find max length of subarray with  $\leq K$  0's → condition

Step 2:

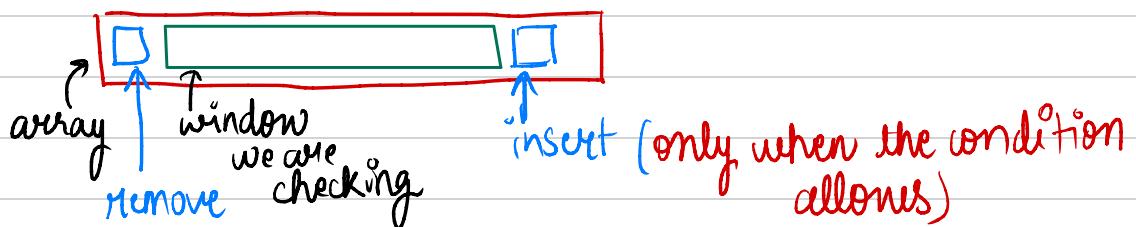
If a condition is true for  $l$  to  $r$ , then it must hold true for  $l$  to  $r-1$  as well.

Here,

$[L, R]$  satisfy  
 $[L, R-1]$  Satisfy

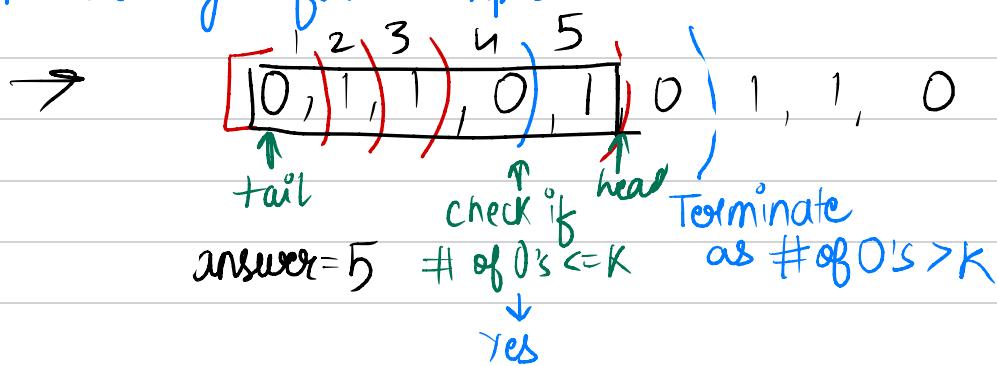
Step 3:

Design a data structure which can check for the identifiable condition. In this case, our check fn will check that if we take an element in our window, then will the condition still remain true?



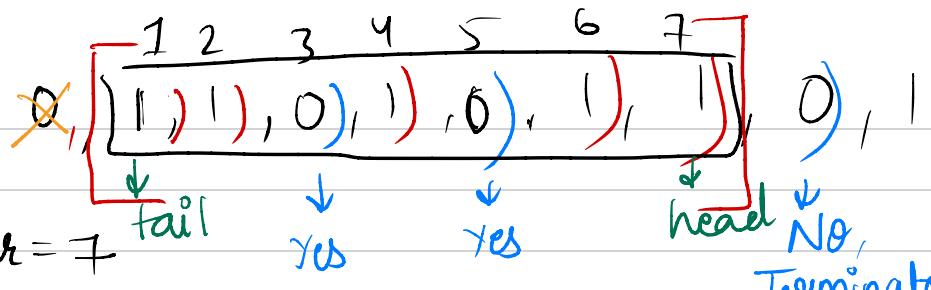
Can we take the next element or not? If it is 1, then we can continue to take it but if its 0 then we shall see if this incoming 0 is  $\leq K$ , if it is then we can continue otherwise return. Our check fn is essentially counting no. of 0's in a window.

• Runthrough for example



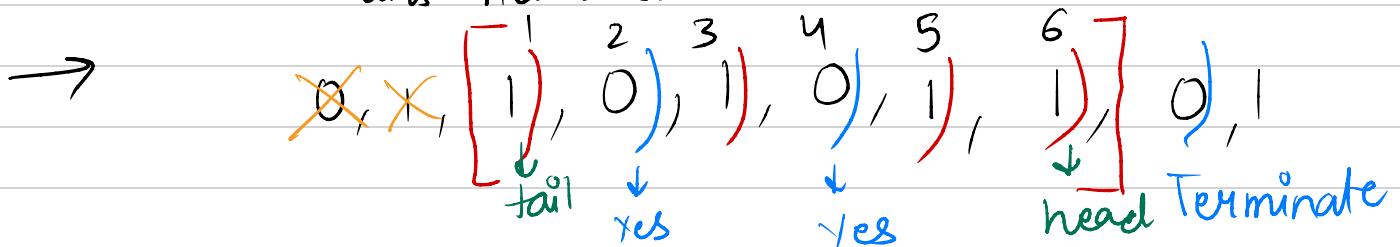
Red parentheses indicates that it can be taken.





Since new answer > answer

ans = new answer



New answer = 6

$6 < 7$ , So answer stays same.

Similarly the window slides through using the head & tail pointer and the best answer is returned

```
// tp
int ans = 0;
int head=-1,tail=0;
while(tail<n){ // O(N) Amortized
    while(head+1<n && (arr[head+1]==1 && cnt0<=k) ||
           (arr[head+1]==0 && cnt0<k)){
        head++;
        // insert head.
        if(arr[head]==0)cnt0++;
    }
    // update answer
    ans = max(ans, head-tail+1); O(1)

    // remove element from tail
    if(tail>=head){ // O(1)
        // remove from DS.
        if(arr[tail]==0)cnt0--;
        tail++;
    }else{
        tail++;
        head = tail-1;
    }
}
```



Ques 2. Find count of number of subarrays with # of distinct elements  $\leq K$ .

Just like last question, here our check would be:

If we take next element, will the # of distinct element will remain  $\leq K$ .

```
// ds  
int distinct = 0;  
int freq[1000100]; < used to keep count of distinct  
void insert(int x){  
    if(freq[x]==0)distinct++;  
    freq[x]++;  
}  
  
int check(int x){  
    int cnt = distinct;  
    if(freq[x]==0)cnt++;  
    return cnt;  
}  
  
int erase(int x){  
    freq[x]--;  
    if(freq[x]==0)distinct--;  
}
```

1/2

```
while(tail<n){  
    while(head+1<n && check(arr[head+1])<=k){  
        head++;  
        // insert head.  
        insert(arr[head]);  
    }  
    // update answer  
    ans += head-tail+1;  
  
    // remove element from tail  
    if(tail<=head){  
        // remove from DS.  
        erase(arr[tail]);  
        tail++;  
    }else{  
        tail++;  
        head = tail-1;  
    }  
  
cout<<ans<<endl;
```

2/2

