

Contribution Technique

Given an array of size N . Find the sum of all the subarrays.

① $N \leq 10^3$

$A: [1 \quad 3 \quad 2]$

② $N \leq 10^5$

Approach 1:

```
int ans = 0
```

```
for (L=0; L<N; L++) {
```

```
    for (R=L; R<N; R++) {
```

```
        sum = 0;
```

```
        for (i=L; i<=R; i++) {
```

```
            sum += arr[i];
```

```
            ans += sum;
```

```
        }
```

```
    }
```

$T.C: O(n^3)$

* Not valid for $N \leq 10^3$ or $N \leq 10^5$.

1			→	1
	3		→	3
		2	→	2
1	3		→	4
	3	2	→	5
1	3	2	→	6
				21

Approach 2:

```
int ans = 0;
```

```
for (L=0; L<N; L++) {
```

```
    sum = 0;
```

```
    for (R=L; R<N; R++) {
```

```
        sum += arr[R];
```

```
        ans += sum;
```

```
    }
```

```
}
```

$T.C: O(n^2)$

* valid for $N \leq 10^3$ but not $N \leq 10^5$

$L=0$ $R=0$	$R=1$	$R=2$
[1	3	2]

sum = ~~(1+3)~~ 1+3+2

ans = 1 + (1+3) + (1+3+2)

* An array has $\frac{N(N+1)}{2}$ ^{no. of} subarrays.

For $N \leq 10^5$, add column wise \Rightarrow check the no. of times i appears and add all.

$[1 \quad 3 \quad 2]$

1			→	1
	3		→	3
		2	→	2
1	3		→	4
	3	2	→	5
1	3	2	→	6
				21

$3 \times 1 + 4 \times 3 + 3 \times 2 = 21$

To find how many times i^{th} element appear, use this formula :



No. of times ^{ith element} will appear = ~~(N+1) × i~~ (i+1) × (N-i)

Approach 3:

```
int ans = 0;
```

```
for (int i=0; i<N; i++) {
```

$$T.C : O(n)$$
$$ans += (arr[i] * (i+1) * (N-i));$$

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Contribution

- Atomic Item contribution
- Extending End contribution
- Every [start/End] contribution

Q.

4	1	3	2
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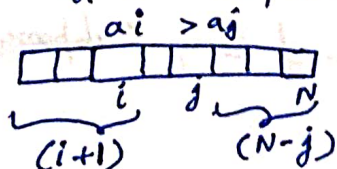
Find the sum of no. of (#) inversion of each subarray ? Constraint : $N \leq 10^3$

For exa: $\begin{array}{ccc} 1 & 3 & 2 \\ \hline & \underline{\quad} & \rightarrow +1 \\ \hline & \underline{\quad} & \rightarrow +1 \end{array} > = 2 \text{ inversions}$

What is inversion?

Then (i, j) exist such that
 $i < j$ and $\text{arr}[i] > \text{arr}[j]$.

* Kiti ne subareage me vol particulare inversion occur hoga.



Total subarrays contains (i, j) pair
 $= (i+1) * (N-j)$ [If $a_i > a_j$]

Approach:

```
void solve() {
```

```
int m;
```

$$c_j n \gg n_j$$

```
int arr[n];
```

```
for (int i=0; i<n; i++) {
```

```
cin >> arr[i];
```



```
int ans = 0;
```

```
for (int i = 0; i < n; i++) {
    for (int j = i + 1; j < n; j++) {
        if (arr[i] > arr[j]) {
            ans += (i + 1) * (n - j);
        }
    }
}
cout << ans << endl;
```

Ex: Q: For every subarray find its product and sum them.

$$\begin{array}{cccc}
 a & b & c & d \\
 a & b & \boxed{\begin{array}{c} c \\ bc + \\ abc + \end{array}} & d \\
 \hline
 a + b + ab + c + bc + abc & & &
 \end{array}
 \rightarrow (c + bc + abc) \times d = cd + bcd + abcd + d$$

$$\begin{array}{cccc}
 a & b & c & d \\
 a \times b & + b \times c & + c \times d & \\
 \hline
 a \times b & + a \times b \times c & + a \times b \times c \times d &
 \end{array}$$

Code :-

```
void solve() {
    int n;
    cin >> n;
    int arr[n];
    for (int i = 0; i < n; i++) {
        cin >> arr[i];
    }
    int ans = 0;
    int lastans = 0;
    for (int i = 0; i < n; i++) {
        lastans = lastans * arr[i] + arr[i];
        ans += lastans;
    }
    cout << ans << endl;
```

Q. Max. Subarray sum

code:-

```
int ans = 0;
```

```
int lastans = 0;
```

```
for (int i=0; i<n; i++) {
```

```
    lastans = max (lastans + arr[i], arr[i]);
```

```
    ans = max (ans, lastans);
```

```
}
```

```
cout << ans << endl;
```