

## Doubt Session

1. Maximize the fraction : Binary search

Approach :-

1. Sort : Order all pairs by their individual A/B ratios, from high to low.
2. Fill a priority queue of size k with the top k pairs to get an initial solution.
3. For each remaining pair, test if swapping it with any of the pair currently in the queue yields a better total ratio. If it does, perform the swap and update your best one.

Ex :-      10    9    7  
              3    5    4

① Pair by Ratio :  $(10, 3) \rightarrow 3.33$   
 $(9, 5) \rightarrow 1.8$   
 $(7, 4) \rightarrow 1.75$

② A priority quo (pq) of size k=2 is filled with first 2 pairs from the sorted list.

Add  $(10, 3)$  :  
 $\text{num\_sum} = 10, \text{den\_sum} = 3$

10,3
9,5

Add  $(9, 5)$  :  
 $\text{num\_sum} = 10 + 9 = 19, \text{den\_sum} = 3 + 5 = 8$

Priosity queue is now filled with  $[(10, 3), (9, 5)]$

③ Current Ratio :  $19/8 = 2.375$

④ Now consider the next pair  $(7, 4)$  and check if swapping it with a pair already in the queue improves the ratio. Replace with lowest value in the current set i.e.,  $(9, 5)$ .

→ Remove  $(9, 5)$  and add  $(7, 4)$

→ New sum : new num\\_sum =  $(19 - 9) + 7 = 17$   
new den\\_sum =  $(8 - 5) + 4 = 7$

→ New Ratio :  $\frac{17}{7} = 2.428$

So Compare and :  $2.428$  (New Ratio)  $> 2.375$  (Current ratio).

So, new pairs are  $(10, 3)$  and  $(7, 4)$  and max achievable ratio is  $2.428$ .

## 2. Codeforces [768B]

middle element =  $n \% 2$

$$\frac{n}{2} \quad \frac{n \% 2}{\downarrow} \quad \left[ \frac{n}{2} \right]$$

Ex:-  $10 \Rightarrow 10 \% 2 = 0, 10 / 2 = 5 \Rightarrow \boxed{5} 0 \boxed{5}^x$  not examined  
 $5 \% 2 = 1, 5 / 2 = 2 \Rightarrow \boxed{2} 1 \boxed{2}^x$   
 $2 \% 2 = 0, 2 / 2 = 1 \Rightarrow \boxed{1} 0 \boxed{1} \Rightarrow [101]$

Now,  $5 \rightarrow \boxed{101} 1 \boxed{101}$

Now  $10 \rightarrow \boxed{1011101} 0 \boxed{1011101}$

We can query L to R, to find overall 1's present there.

Don't form vectors, use strings

## 3. Cellular Network - codeforces [702C]

↳ Binary search on Ans

~~Instead of checking radius one by one find the best radius  $\sigma$  by guessing.~~

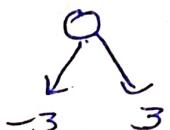
Max. range of tower =  $|c[0] - c[n-1]|$   
 $= \text{abs}(c[0] - c[n-1]).$

Ex:- - 2 2 4

Min = 0

Mid = 3

Max = 7



in this range

-6 --- 3 [From here, any tower + can provide service]

4 is outside range, so if any of the city left we have to inc. the range.

Minimise the range : store mid  $\rightarrow \sigma$  (it can be a probable ans).

Search in 0 --- mid-1 for better ans. Update if found better.