

Doubt Session

1. Maximize the fraction : Binary Search

Approach :-

1. Sort : Order all pairs by their individual A/B ratio, from high to low.
2. Fill a priority queue of size k with the top k pairs to get an initial solution.
3. For each remaining pair, test if swapping it with any of the pairs currently in the queue yields a better total ratio. If it does, perform the swap and update your best ans.

Exa :

10	9	7
3	5	4

① Pair by Ratio : $(10, 3) \rightarrow 3.33$
 $(9, 5) \rightarrow 1.8$
 $(7, 4) \rightarrow 1.75$

② A priority queue (pq) of size $k=2$ is filled with first 2 pairs from the sorted list.

Add $(10, 3)$:
num-sum = 10, den-sum = 3

Add $(9, 5)$:

num-sum = $10 + 9 = 19$, den-sum = $3 + 5 = 8$

10, 3
9, 5

Priority queue is now filled with $[(10, 3), (9, 5)]$

③ Current Ratio : $19/8 = 2.375$

④ Now consider the next pair $(7, 4)$ and check if swapping it with a pair already in the queue improves the ratio. Replace with lowest value in the current set i.e., $(9, 5)$.

→ Remove $(9, 5)$ and add $(7, 4)$

→ New sum : new num-sum = $(19 - 9) + 7 = 17$

new den-sum = $(8 - 5) + 4 = 7$

New
→ Ratio : $\frac{17}{7} = 2.428$

5. Compare and : 2.428 (New Ratio) $>$ 2.375 (Current ratio).

So, new pairs are $(10, 3)$ and $(7, 4)$ and max. achievable ratio is 2.428 .

2. Codeforces [768B]
 middle element = $n/2$

$$\left\lfloor \frac{n}{2} \right\rfloor \quad \frac{n/2}{\downarrow} \quad \left\lceil \frac{n}{2} \right\rceil$$

0/1

Ex:- 10
 $10 \% 2 = 0, 10/2 = 5 \Rightarrow \boxed{5} 0 \boxed{5}^x$ not required
 \downarrow
 Break down this

$5 \% 2 = 1, 5/2 = 2 \Rightarrow \boxed{2} 1 \boxed{2}^x$

$2 \% 2 = 0, 2/2 = 1 \Rightarrow \boxed{1} 0 \boxed{1} \Rightarrow [1 0 1]$

Now, 5 $\rightarrow \boxed{101} 1 \boxed{101}$

Now 10 $\rightarrow \boxed{1011101} 0 \boxed{1011101}$ $\begin{matrix} L & R \end{matrix}$

We can query L to R, to find overall 1's present there.

Don't form vectors, use strings.

3. Cellular Network - codeforces [702C]

\hookrightarrow Binary search on Ans

~~Instead of checking radii one by one find the best radius by guessing.~~

Max. range of tower = $|c[0] - c[n-1]|$
 $= \text{abs}(c[0] - c[n-1])$

Ex:- -2 2 4

Min = 0 Mid = 3

Max = 7



In this range

-6 --- 3 [From here, any tower can provide service]

4 is outside range, so if any of the city left we have to inc. the range.

Minimise the range: store mid \rightarrow (it can be a probable ans).

search in 0 --- mid-1 for better ans. Update if found better.