

Problem Set 2

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Due: 10/8, 11:59pm.

Please typeset your solutions in LaTeX.

Problem 1 (On \models , borrowed from Aaron Potechin). Consider the following polynomial equation in 3 variables, x, y, z .

$$(x^2 + 1)y = z^2.$$

Because it implies $y = \frac{z^2}{x^2+1}$, any solution (x, y, z) to the above must have $y \geq 0$. We will see if sum-of-squares can capture this reasoning.

1. Construct a degree 4 pseudoexpectation $\tilde{\mathbb{E}}$ in variables x, y, z such that $\tilde{\mathbb{E}} \models (x^2 + 1)y = z^2$ but $\tilde{\mathbb{E}} y < 0$. (Computer-aided proofs are allowed.)

By $\tilde{\mathbb{E}} \models (x^2 + 1)y = z^2$, we mean that for any polynomial p of degree at most 1 in x, y, z , $\tilde{\mathbb{E}} p(x, y, z)(x^2 + 1)y = \tilde{\mathbb{E}} p(x, y, z)z^2$.

2. Despite the above, show that there exists a sum-of-squares refutation to the following system of polynomial inequalities, for any $c > 0$: $\{(x^2 + 1)y = z^2, y \leq -c\}$.

Problem 2. Suppose $\tilde{\mathbb{E}}$ is a pseudoexpectation of degree d , with d even, and $\tilde{\mathbb{E}} \models p \leq 0, p \geq 0$ for some polynomial p . (Informally, we have been writing $\tilde{\mathbb{E}} \models p = 0$.) Show that if p has even degree, for every q such that the degree of pq is at most d , we have $\tilde{\mathbb{E}} pq = 0$. Similarly, show that if p has odd degree, for every q such that the degree of pq is at most $d - 1$, we have $\tilde{\mathbb{E}} pq = 0$.

Problem 3. Unreleased.

Bonus Problem 4 (Integrality gaps for max-cut, borrowed from Pravesh Kothari). Let C_n be the cycle graph on vertex set $[n]$ with edge set E . Further suppose that n is odd. The size of the max-cut in C_n is $n - 1$. Recall from your solution to Problem 2 of the first problem set that this implies that for any degree 2 pseudoexpectation $\tilde{\mathbb{E}}$ on $\{\pm 1\}^n$, $\tilde{\mathbb{E}} \left[\frac{1}{4} \sum_{ij \in E} (x_i - x_j)^2 \right] \leq \left(1 - O\left(\frac{1}{n^2}\right) \right) n$. We will start by seeing that this is tight for degree 2 pseudoexpectations.

Let L the Laplacian of C_n defined by $L_{ii} = 2$ for each i , and L_{ij} is -1 if ij is an edge and 0 otherwise. Observe that for $x \in \{\pm 1\}^n$, the size of the cut associated to x is equal to $\frac{1}{4} \cdot x^\top L x$.

For each $0 \leq k \leq n/2$, let x_k, y_k be vectors with coordinates $(x_k)_i = \cos(2\pi i k/n)$ and $(y_k)_i = \sin(2\pi i k/n)$.

1. Prove that x_k and y_k are eigenvectors of L with eigenvalues $2 - 2 \cos(2\pi k/n)$.
2. Prove that the diagonal entries of the matrix $M_k = x_k x_k^\top + y_k y_k^\top$ are 1.

3. Prove that there is a degree 2 pseudoexpectation $\widetilde{\mathbb{E}}_k$ on $\{\pm 1\}^n$ with $\widetilde{\mathbb{E}} x = 0$ and $\widetilde{\mathbb{E}} x x^\top = M_k$. Using this, prove that for $k = \frac{n-1}{2}$, $\widetilde{\mathbb{E}} \left[\frac{1}{4} \sum_{ij \in E} (x_i - x_j)^2 \right] \geq \left(1 - O\left(\frac{1}{n^2}\right) \right) n$.

Next, we will see that degree 6 pseudoexpectations do not face such barriers (for the cycle graph).

4. Prove that for degree 6 pseudoexpectations $\widetilde{\mathbb{E}}$ over $\{\pm 1\}^n$, the squared triangle inequality holds: $\widetilde{\mathbb{E}}(x_i - x_j)^2 \leq \widetilde{\mathbb{E}}(x_i - x_k)^2 + \widetilde{\mathbb{E}}(x_k - x_j)^2$. For a harder exercise, prove this for degree 4 pseudoexpectations.
5. Prove that for any degree 6 pseudoexpectation $\widetilde{\mathbb{E}}$, $\widetilde{\mathbb{E}} \left[\frac{1}{4} \sum_{ij \in E} (x_i - x_j)^2 \right] \leq n - 1$.