

# Scroll & zkEVM & Proving system

Ye Zhang



Scroll

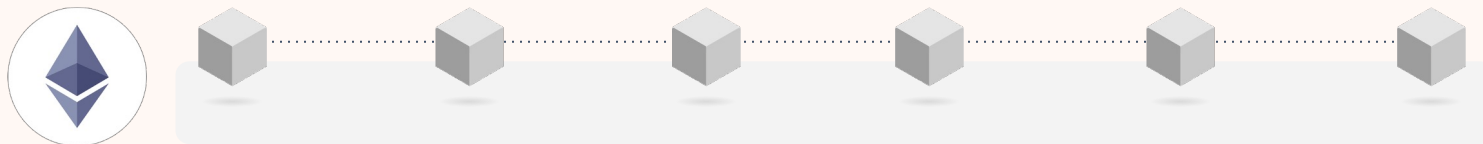
# **What is Scroll?**

**A scaling solution for Ethereum**

# What is Scroll?

An **EVM-equivalent** zk-Rollup

# The technical workflow



Pre-committed (1-2s)

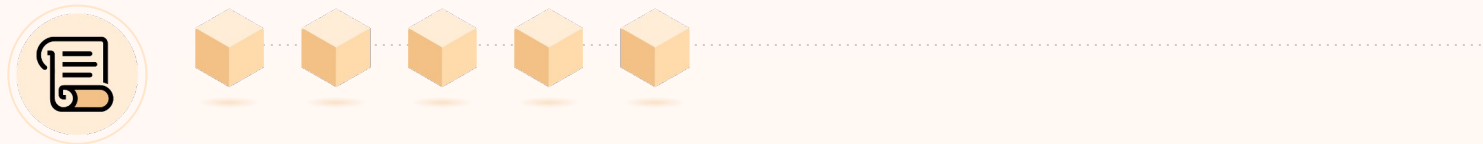
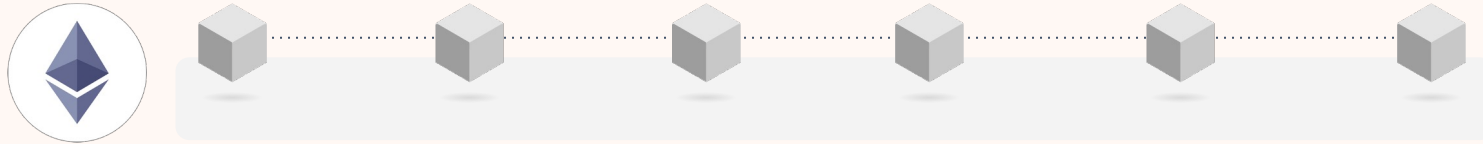


Committed (minutes)




Finalized (10+minutes)

# The technical workflow

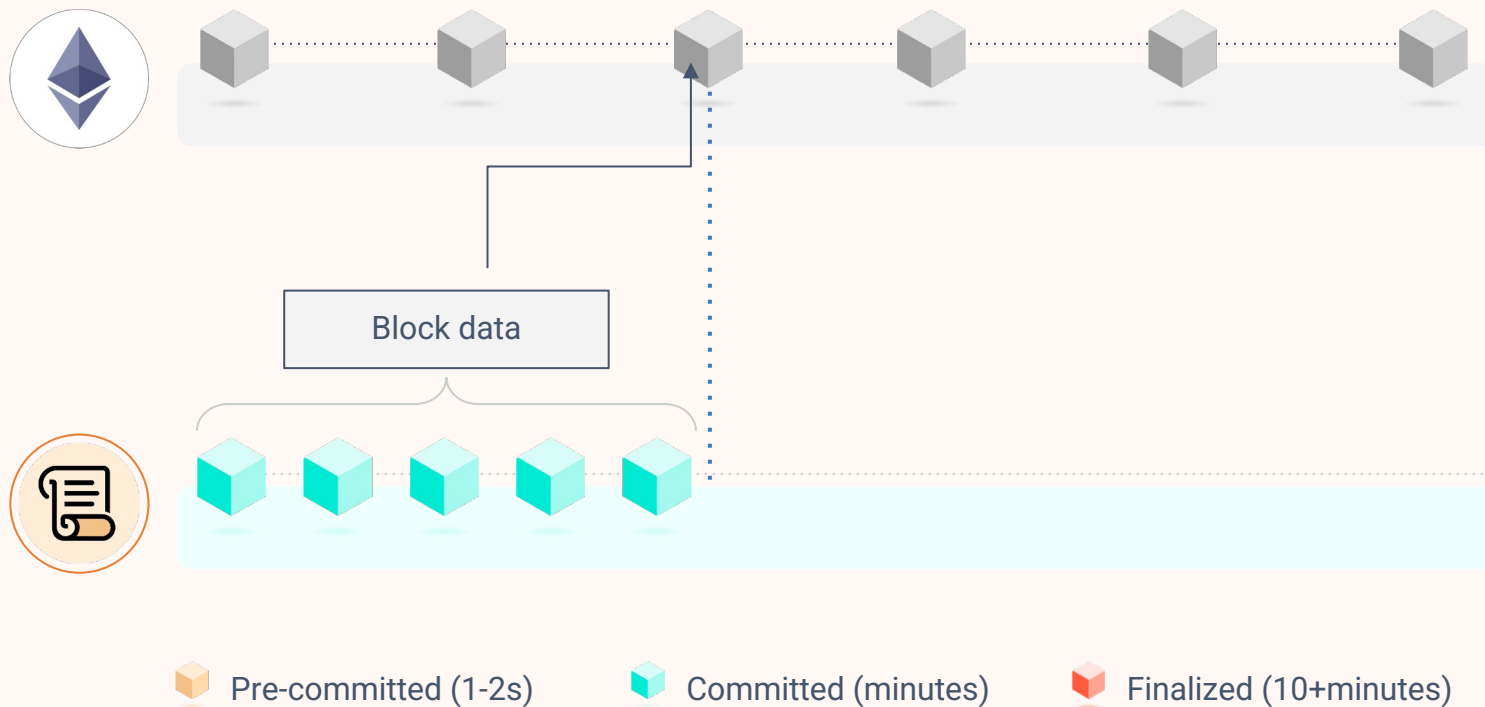


 Pre-committed (1-2s)

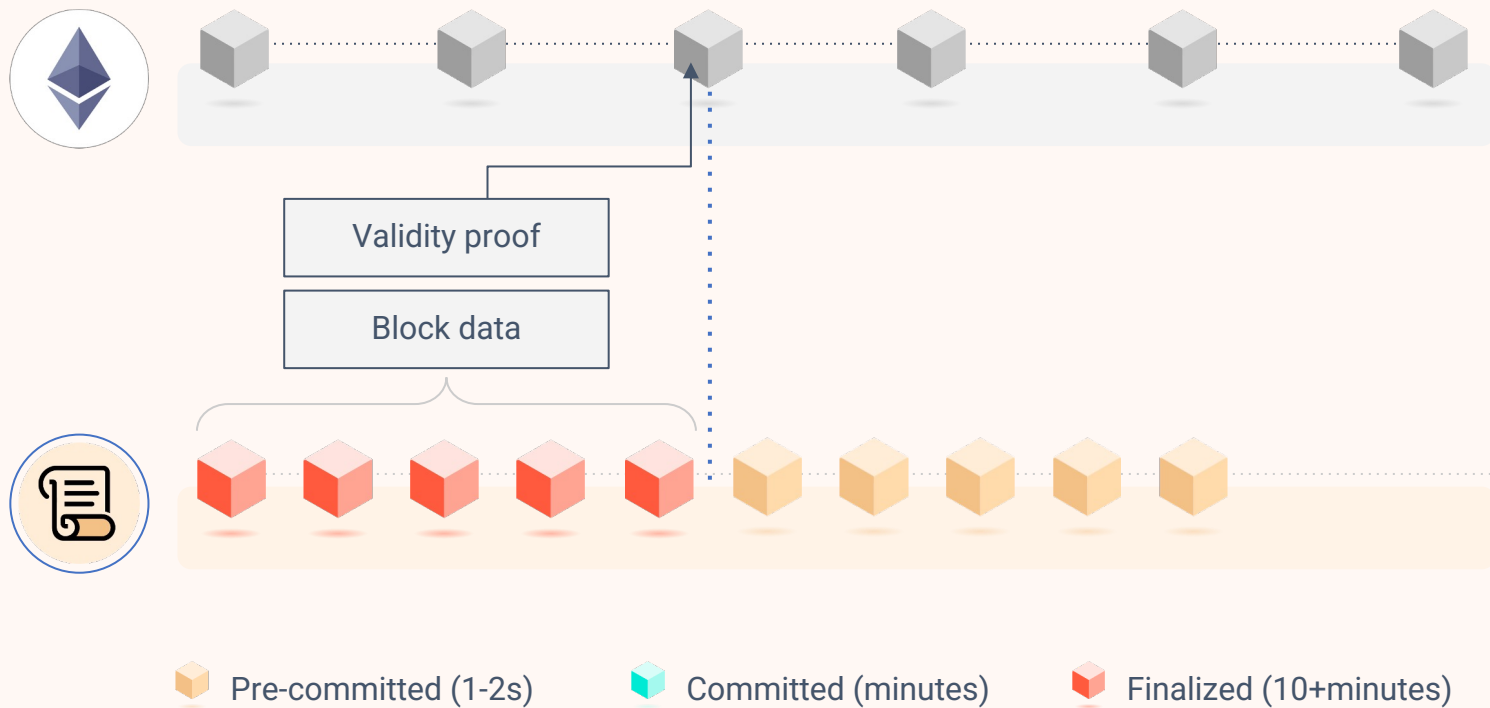
 Committed (minutes)

 Finalized (10+minutes)

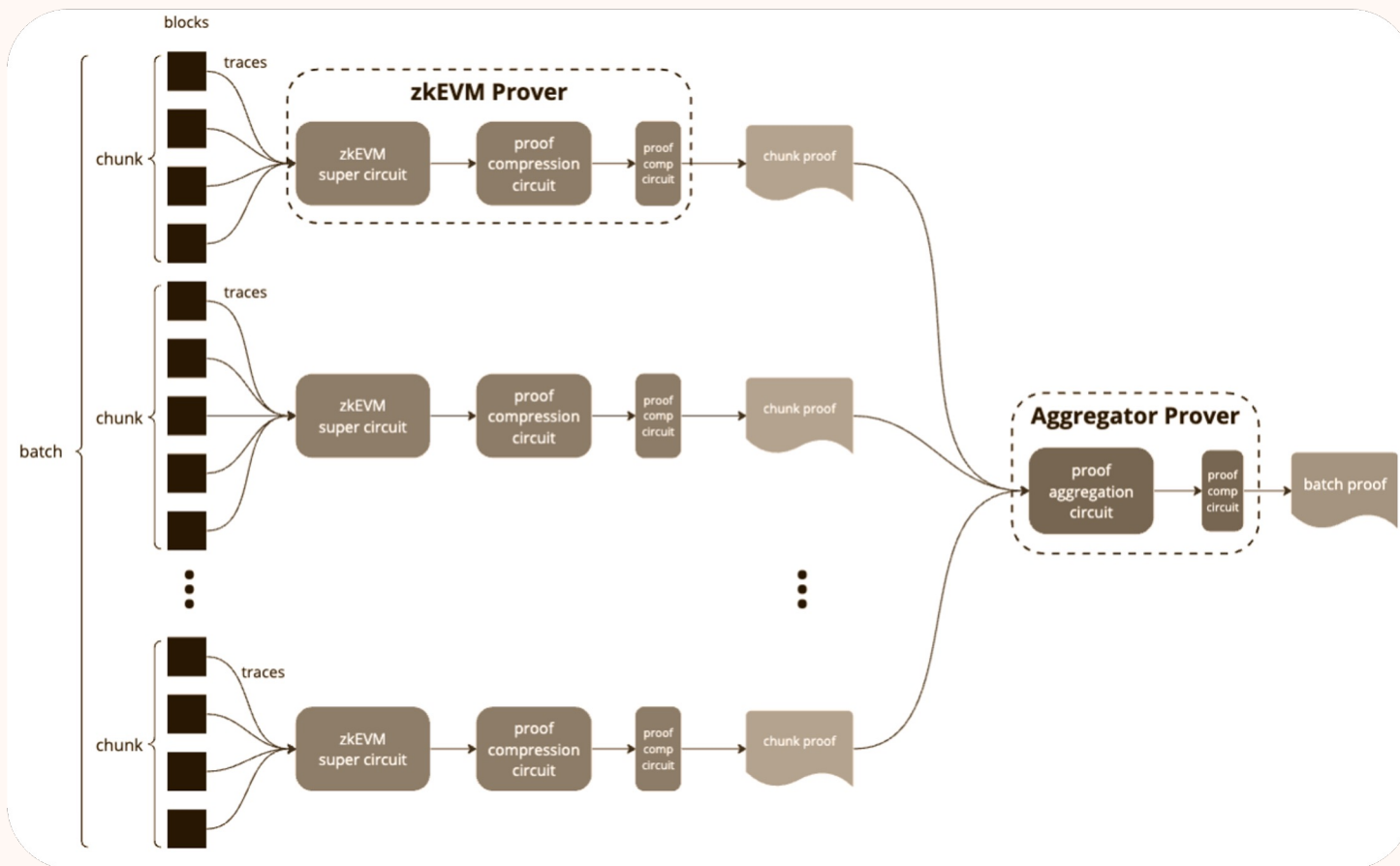
# The technical workflow



# The technical workflow



# The technical workflow






For users and devs:

**Scroll = Ethereum**


But cheaper and faster  
with a higher throughput

 Copy solidity code from dApp



 Paste code into Scroll interface



 Unlock all the benefits of  
Scroll's EVM-equivalent zk-  
rollup

For users and devs:

**Scroll = Ethereum**

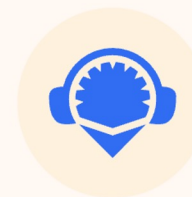
But cheaper and faster  
with a higher throughput



Hardhat



Foundry



Remix



EthersJS



Brownie



Truffle

Privacy & Scaling Explorations (PSE)

A zk working group within the Ethereum Foundation

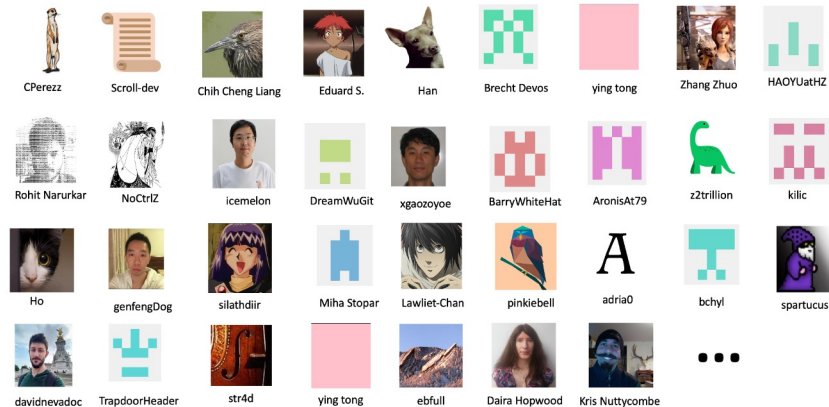


# The journey with community

Privacy & Scaling Explorations (PSE)  
A zk working group within the Ethereum Foundation



## Credit to all community members!

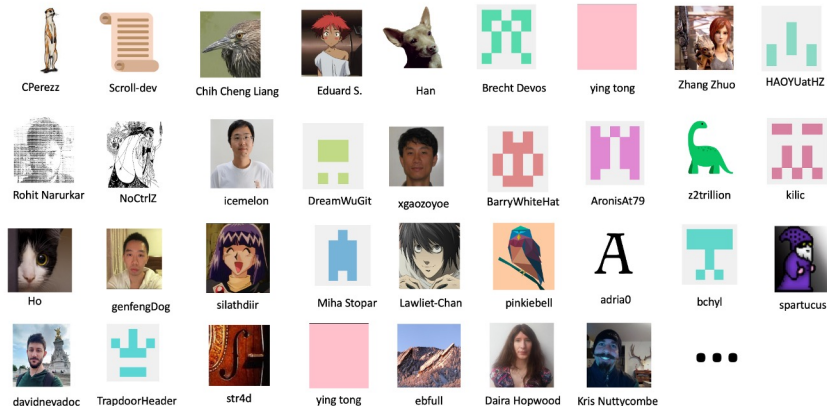


# The journey with community

Privacy & Scaling Explorations (PSE)  
A zk working group within the Ethereum Foundation



## Credit to all community members!



### Building and Testing Circuits with halo2-ce: An Introductory Workshop

March 20, 2023 3PM ET

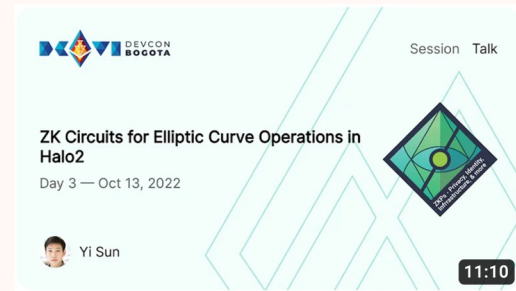
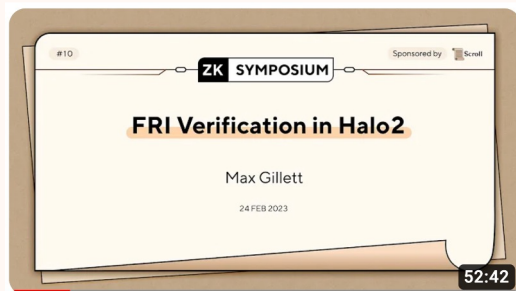


Scroll

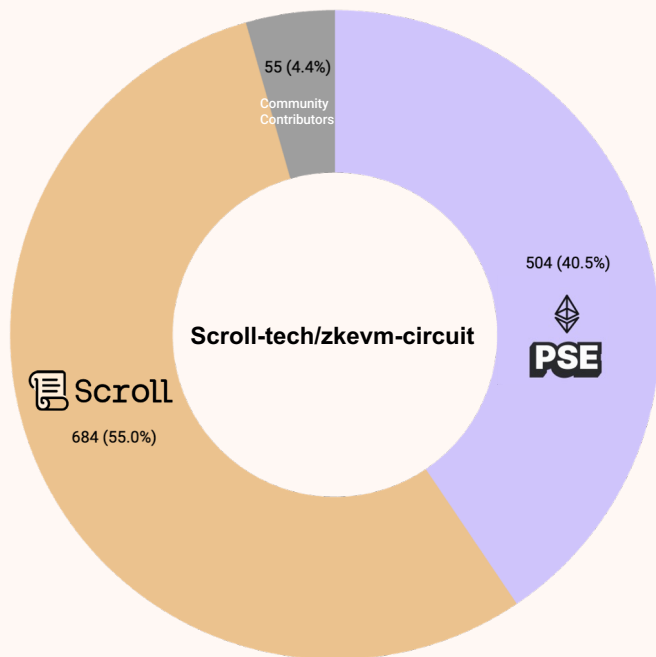


Mason Liang  
ZK Circuit Engineer @ Scroll

1:03:38



# Value: Openness and Community-driven Scroll



- Full transparency – anyone can review code
- The most open zkEVM from day 1
- The most community-driven zkEVM
- The only zkEVM that supports ecPairing

## Security Audit



Trail of Bits



OpenZeppelin



Zellic



Kalos



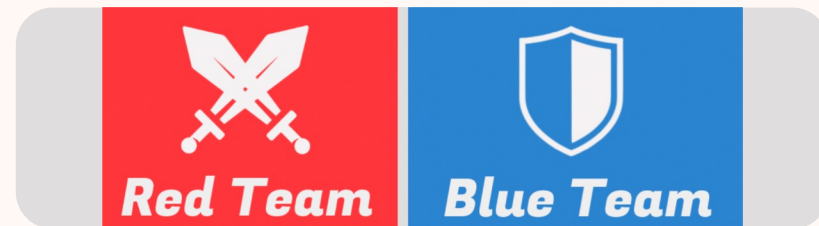
ABDK

Replay Tx's from different networks

Standard EVM test vectors

Fuzzing for zkEVM opcodes

## In-house Security Team



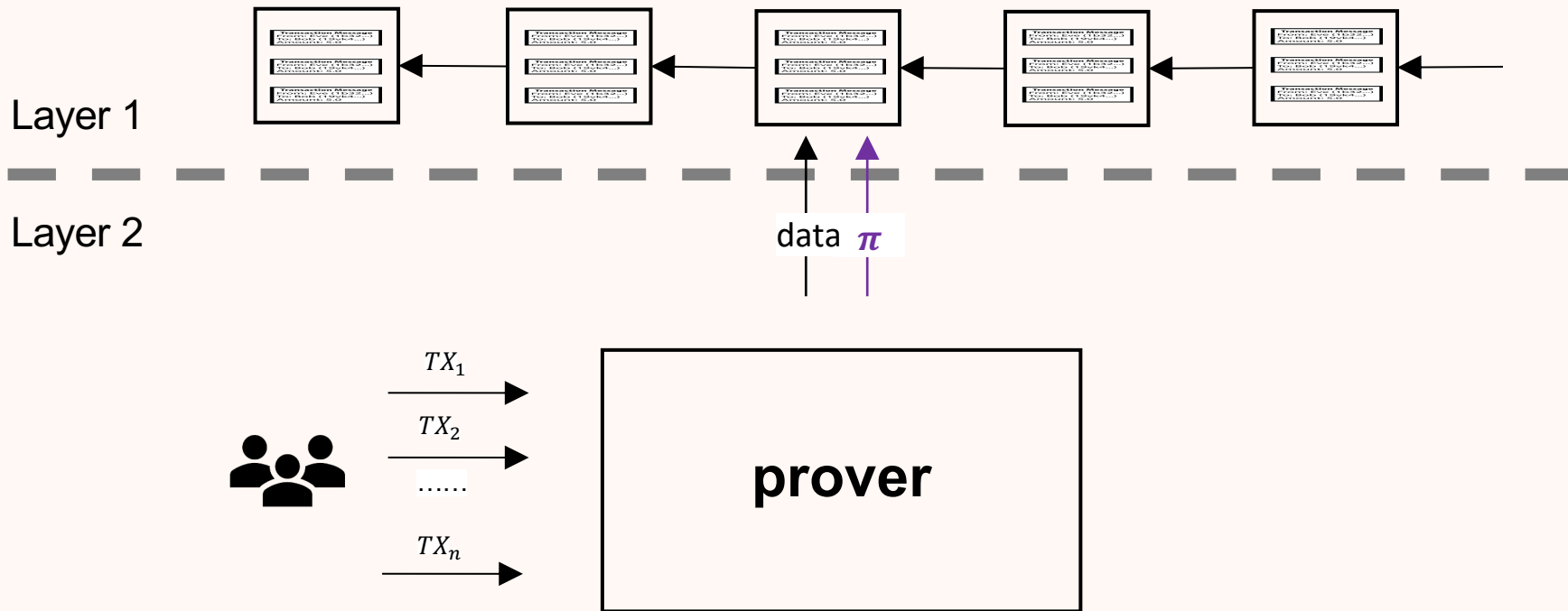
Reuse Geth

Fully open-source

Many more...

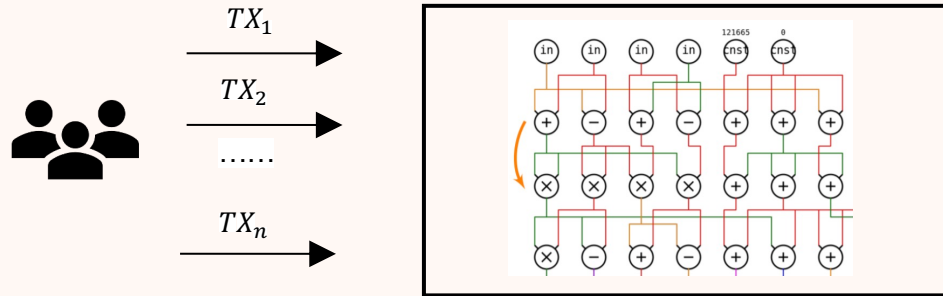
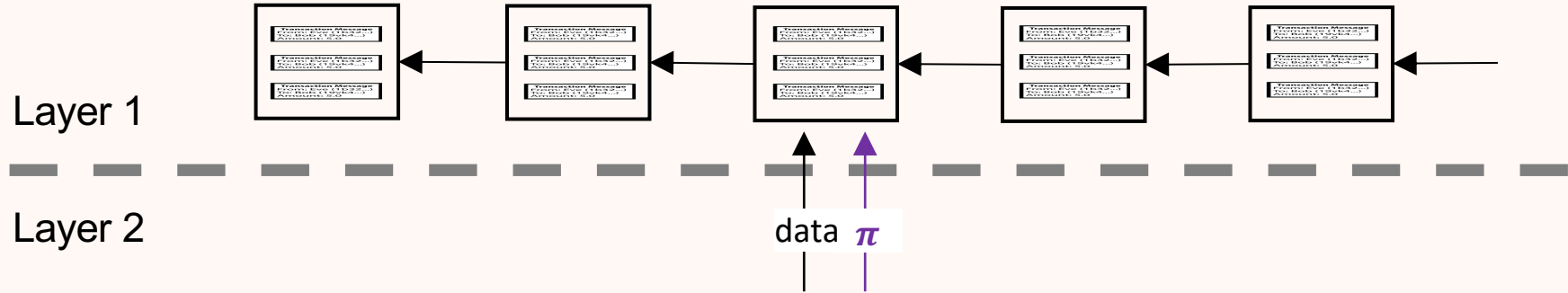
- Background & motivation
- zkEVM circuit arithmetization
- zkEVM prover optimization
- Something interesting to share





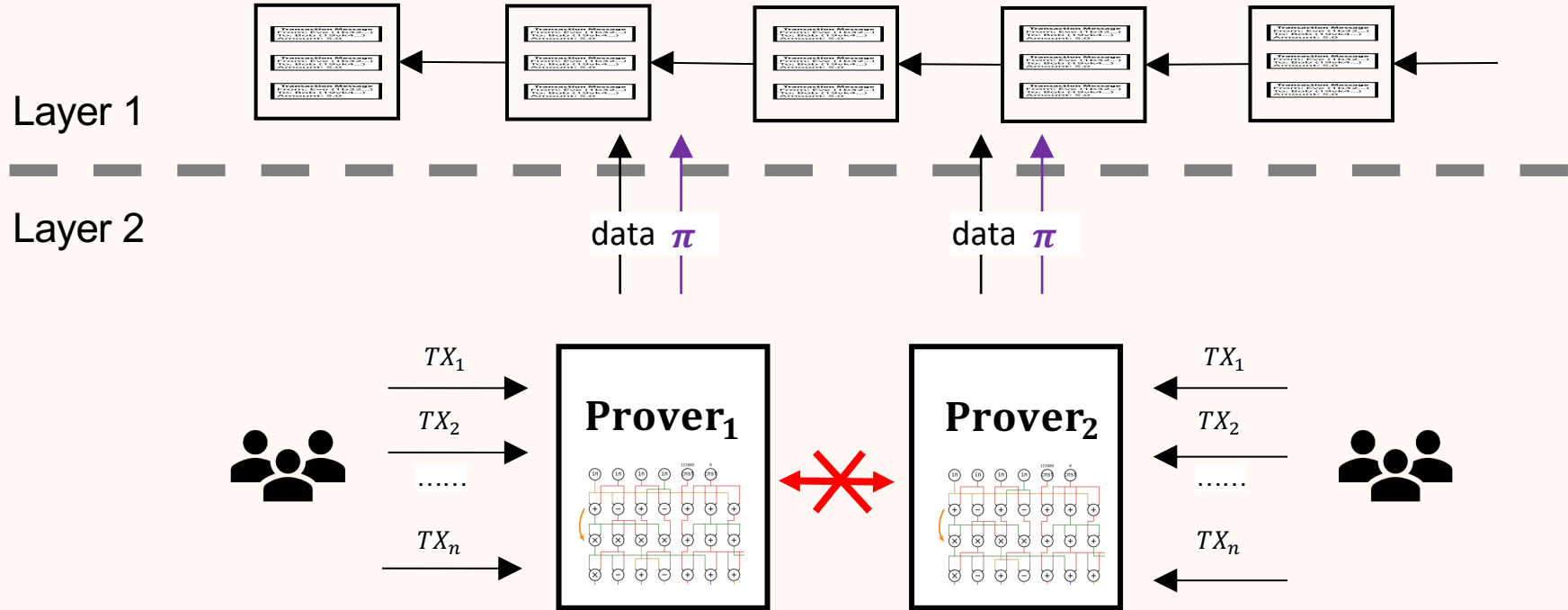
However, ...

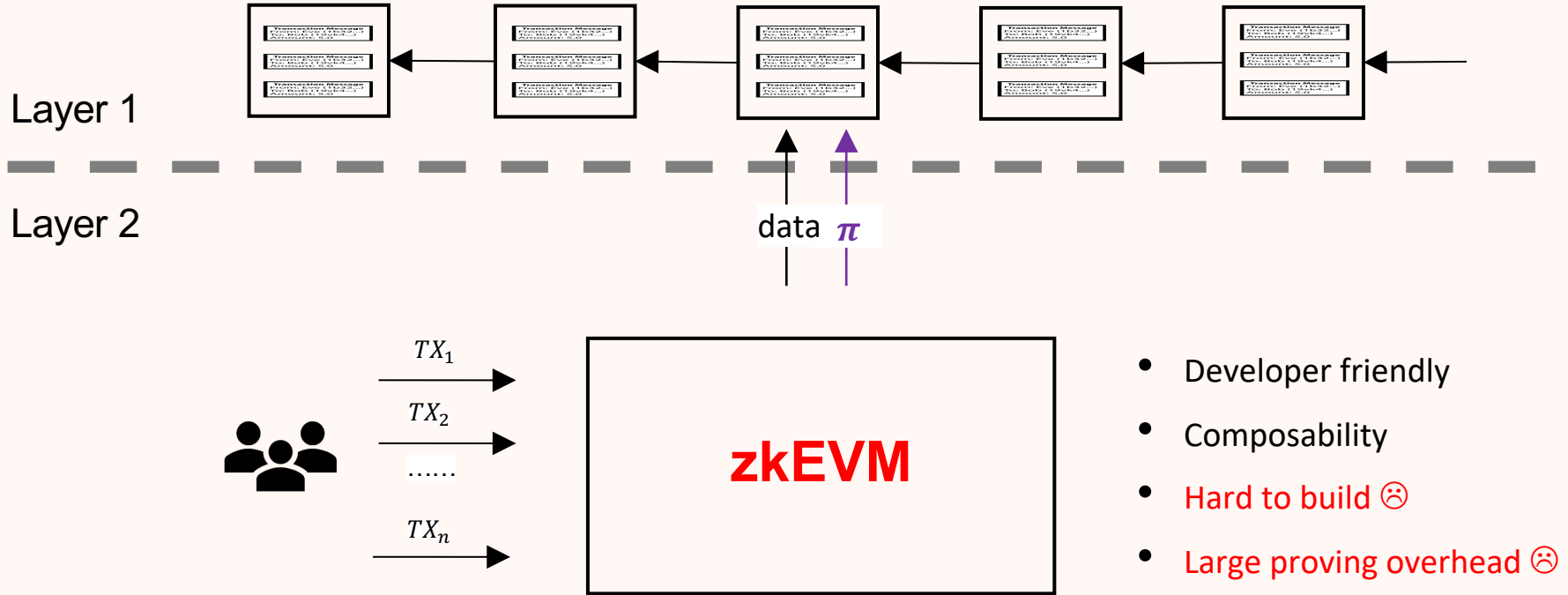
 Scroll

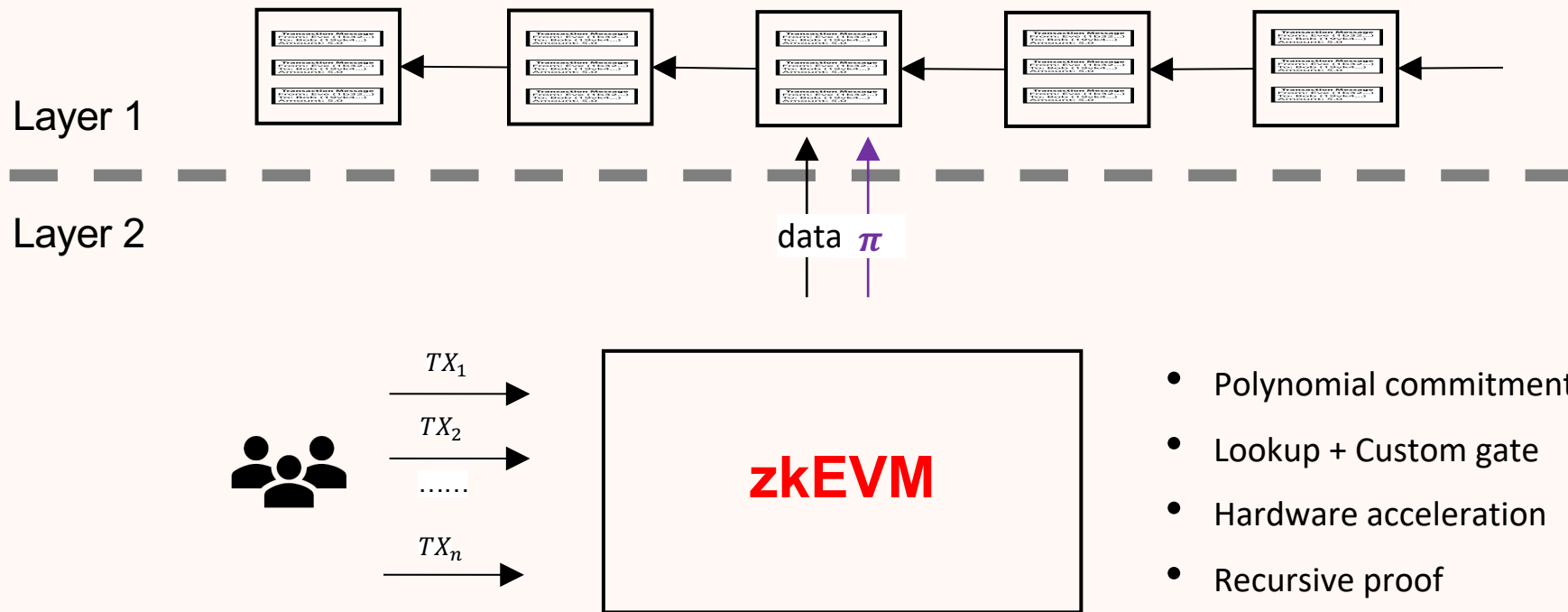


However, ...

 Scroll







- **Language level**  
Transpile an EVM-friendly language (Solidity or Yul) to a SNARK-friendly VM which differs from the EVM. This is the approach of Matter Labs and Starkware.
- **Bytecode level**  
Interpret EVM bytecode directly, though potentially producing different state roots than the EVM, e.g. if certain implementation-level data structures are replaced with SNARK-friendly alternatives. This is the approach taken by Scroll, Hermez, and Consensys.
- **Consensus level**  
Target full equivalence with EVM as used by Ethereum L1 consensus. That is, it proves validity of L1 Ethereum state roots. This is part of the "zk-SNARK everything" roadmap for Ethereum.

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- **Consensus level**

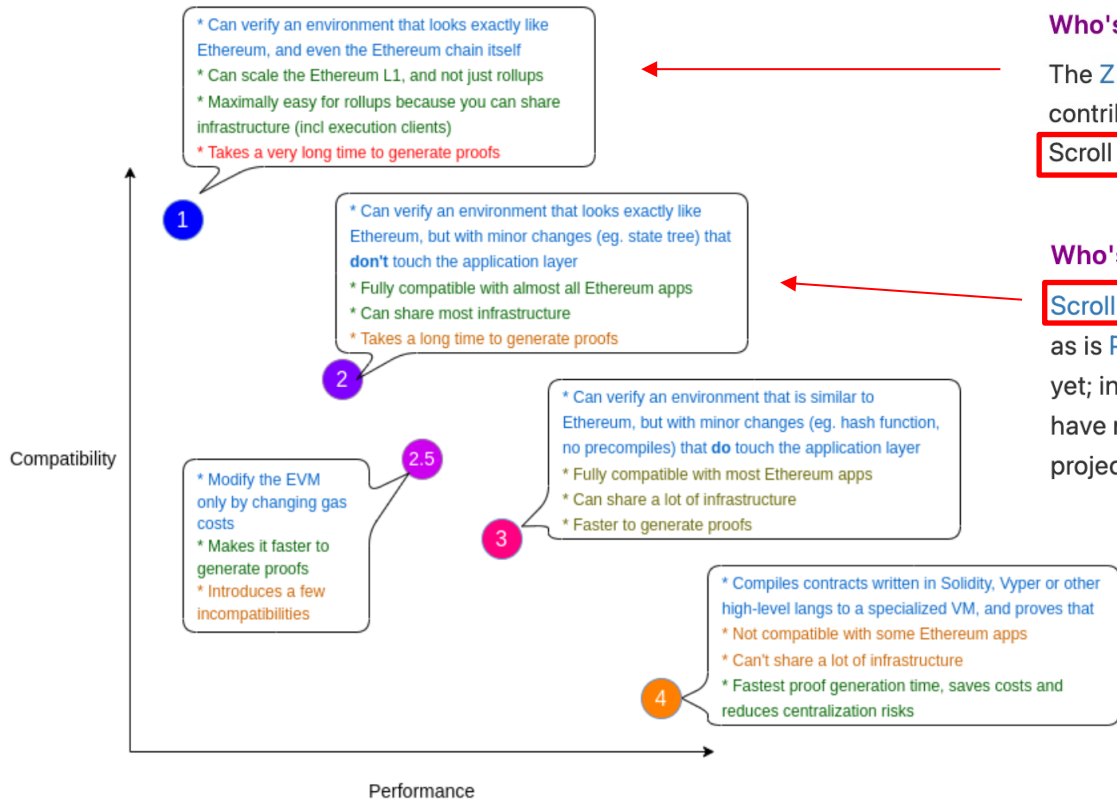
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# zkEVM flavors (by Vitalik)





# zkEVM flavors (by Vitalik)



## Who's building it?

The **ZK-EVM Community Edition** (bootstrapped by community contributors including **Privacy and Scaling Explorations**, the **Scroll team**, **Taiko** and others) is a Tier 1 ZK-EVM.

## Who's building it?

**Scroll's ZK-EVM project** is building toward a Type 2 ZK-EVM, as is **Polygon Hermez**. That said, neither project is quite there yet; in particular, a lot of the more complicated precompiles have not yet been implemented. Hence, at the moment both projects are better considered **Type 3**.

- Background & motivation
- zkEVM circuit arithmetization
- zkEVM prover optimization
- Something interesting to share

# The workflow of zero-knowledge proof Scroll

**Program**

**Constraints**

**Proof**

```
def hcf(x, y):  
    if x > y:  
        smaller = y  
    else:  
        smaller = x  
  
    for i in range(1, smaller + 1):  
        if (x % i == 0) and (y % i == 0):  
            hcf = i  
  
    return hcf
```



R1CS  
Plonkish  
AIR

```
x * x == var1  
var1 * x == y  
(y+x) * 1 == var2  
(var2+5) * 1 == out
```



Polynomial IOP  
+  
PCS

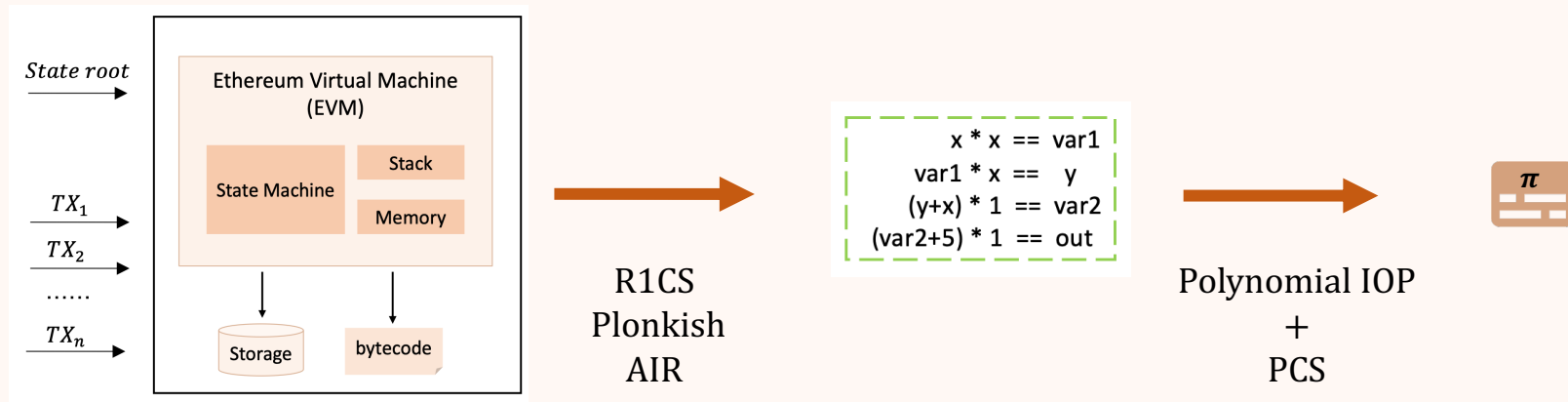


# The workflow of zero-knowledge proof Scroll

**Program**

**Constraints**

**Proof**

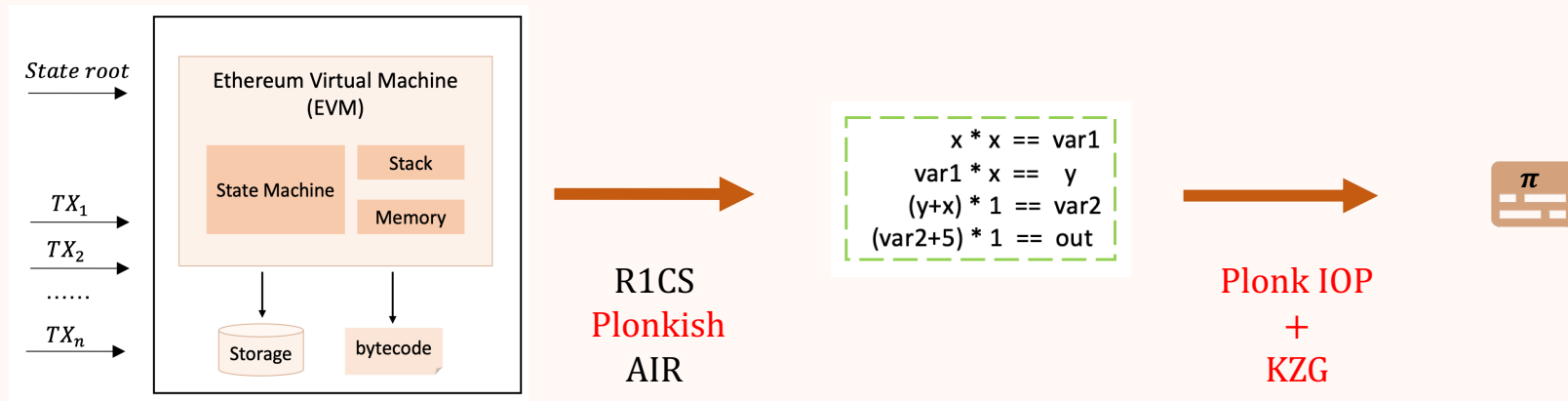


# The workflow of zero-knowledge proof Scroll

**Program**


**Constraints**

**Proof**




# Plonkish Arithmetization

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$



# Plonkish Arithmetization

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					



witness                      Table 1                      Table 2

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

witness

Table 1

Table 2

$$va_3 * vb_3 * vc_3 - vb_4 = 0$$



$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

witness
Table 1
Table 2

$$va_3 * vb_3 * vc_3 - vb_4 = 0$$

- High degree
- More customized

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
<i>input<sub>0</sub></i>	<i>input<sub>1</sub></i>	<i>input<sub>2</sub></i>		<i>output</i>					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

witness
Table 1
Table 2

$$vb_1 * vc_1 + vc_2 - vc_3 = 0$$

- High degree
- More customized

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

witness
Table 1
Table 2

$$vc_1 + va_2 * vb_4 - vc_4 = 0$$

- High degree
- More customized

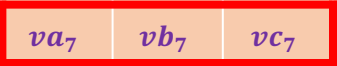

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

Diagram illustrating the permutation of variables in the Plonkish Arithmetization. The table is divided into three sections: **witness** (columns  $a_0$  to  $a_4$ ), **Table 1** (columns  $T_0$  to  $T_1$ ), and **Table 2** (columns  $T_2$  to  $T_4$ ). Red boxes highlight the variables  $vb_4$ ,  $vc_6$ ,  $va_6$ , and  $vb_6$ . Red arrows indicate the permutation:  $vb_4 \rightarrow vc_6 \rightarrow vb_6 \rightarrow va_6$ .

$$vb_4 = vc_6 = vb_6 = va_6$$

# Plonkish Arithmetization – Lookup argument Scroll

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					





witness
Table 1
Table 2

$$(va_7, vb_7, vc_7) \in (T_0, T_1, T_2)$$

# Plonkish Arithmetization – Lookup argument Scroll

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$	0000				
$va_1$	$vb_1$	$vc_1$		$vd_1$	0001				
$va_2$	$vb_2$	$vc_2$		$vd_2$	0010				
$va_3$	$vb_3$	$vc_3$		$vd_3$	0011				
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$	0100				
$va_5$	$vb_5$	$vc_5$		$vd_5$	0101				
$va_6$	$vb_6$	$vc_6$		$vd_6$	.....				
$va_7$	$vb_7$	$vc_7$			1101				
$va_6$	$vb_6$	$vc_6$		$vd_6$	1110				
$va_7$	$vb_7$	$vc_7$		$vd_7$	1111				


 Lookup

witness                      Table 1                      Table 2

$vc_7 \in [0, 15]$

# Plonkish Arithmetization – Lookup argument Scroll

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$	0000	0000	0000		
$va_1$	$vb_1$	$vc_1$		$vd_1$	0000	0001	0001		
$va_2$	$vb_2$	$vc_2$		$vd_2$	0000	0010	0010		
$va_3$	$vb_3$	$vc_3$		$vd_3$	0000	0011	0011		
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$	0000	0100	0100		
$va_5$	$vb_5$	$vc_5$		$vd_5$	0000	0101	0101		
$va_6$	$vb_6$	$vc_6$		$vd_6$	.....	.....	.....		
$va_7$	$vb_7$	$vc_7$			1111	1101	0010		
$va_6$	$vb_6$	$vc_6$		$vd_6$	1111	1110	0001		
$va_7$	$vb_7$	$vc_7$		$vd_7$	1111	1111	0000		




witness
Table 1
Table 2

$$vc_7 \in [0, 15]$$

$$va_7 \oplus vb_7 = vc_7$$

# Plonkish Arithmetization – Lookup argument Scroll

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$	0000	0000	0000		
$va_1$	$vb_1$	$vc_1$		$vd_1$	0000	0001	0001		
$va_2$	$vb_2$	$vc_2$		$vd_2$	0000	0010	0010		
$va_3$	$vb_3$	$vc_3$		$vd_3$	0000	0011	0011		
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$	0000	0100	0100		
$va_5$	$vb_5$	$vc_5$		$vd_5$	0000	0101	0101		
$va_6$	$vb_6$	$vc_6$		$vd_6$	.....	.....	.....		
$va_7$	$vb_7$	$vc_7$			1111	1101	0010		
$va_6$	$vb_6$	$vc_6$		$vd_6$	1111	1110	0001		
$va_7$	$vb_7$	$vc_7$		$vd_7$	1111	1111	0000		



witness
Table 1
Table 2


$$vc_7 \in [0, 15]$$


$$va_7 \oplus vb_7 = vc_7$$


**RAM operation**



$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	$T_0$	$T_1$	$T_2$	$T_3$	$T_4$
$input_0$	$input_1$	$input_2$		$output$					
$va_1$	$vb_1$	$vc_1$		$vd_1$					
$va_2$	$vb_2$	$vc_2$		$vd_2$					
$va_3$	$vb_3$	$vc_3$		$vd_3$					
$va_4$	$vb_4$	$vc_4$	.....	$vd_4$					
$va_5$	$vb_5$	$vc_5$		$vd_5$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					
$va_6$	$vb_6$	$vc_6$		$vd_6$					
$va_7$	$vb_7$	$vc_7$		$vd_7$					

  
**witness**

  
**Table 1**

  
**Table 2**

$$vb_1 * vc_1 + vc_2 - vc_3 = 0$$

$$va_3 * vb_3 * vc_3 - vb_4 = 0$$

$$vb_4 + vc_6 * vb_6 - va_6 = 0$$

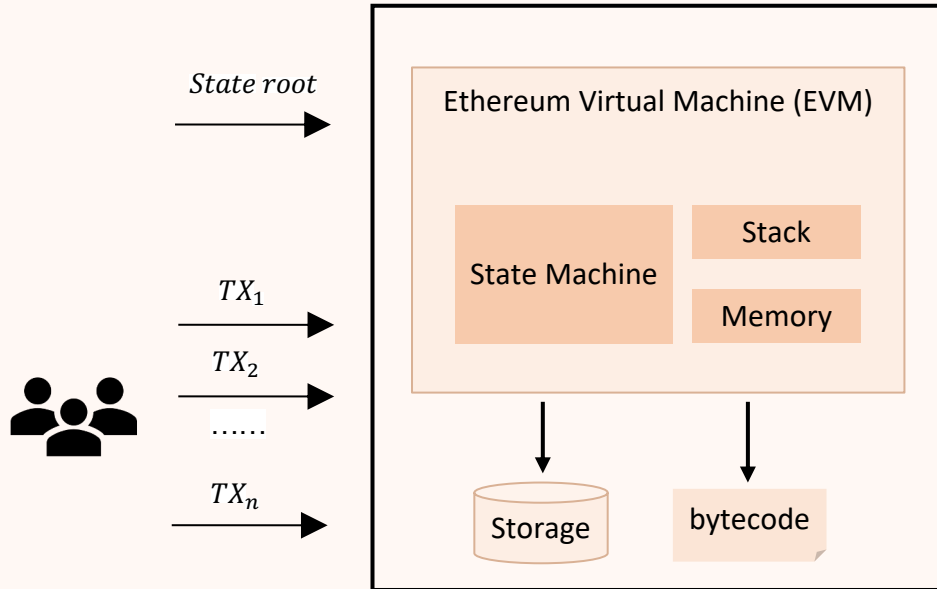
.....

$$vb_4 = vc_6 = vb_6 = va_6$$

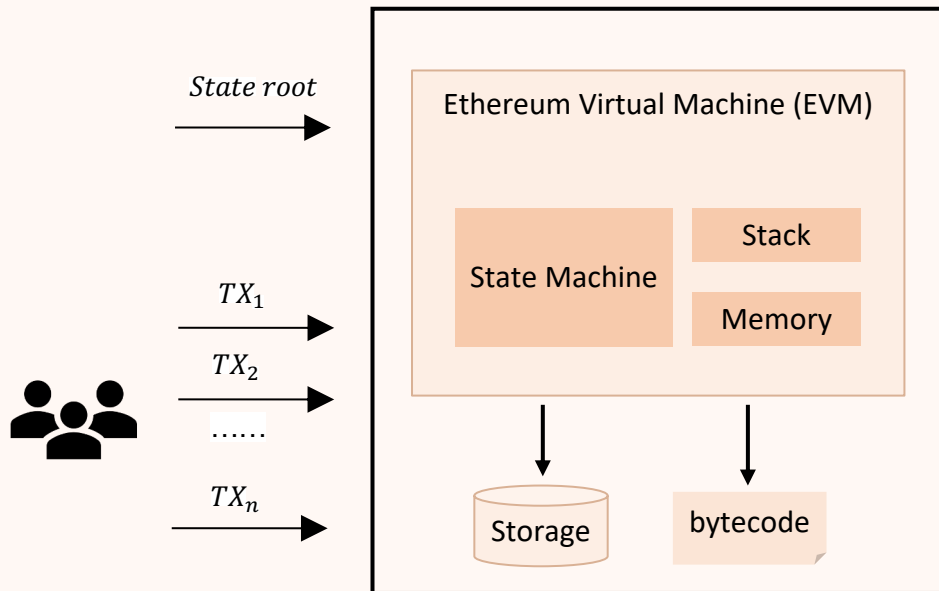
.....

$$(va_7, vb_7, vc_7) \in (T_0, T_1, T_2)$$

## Computation

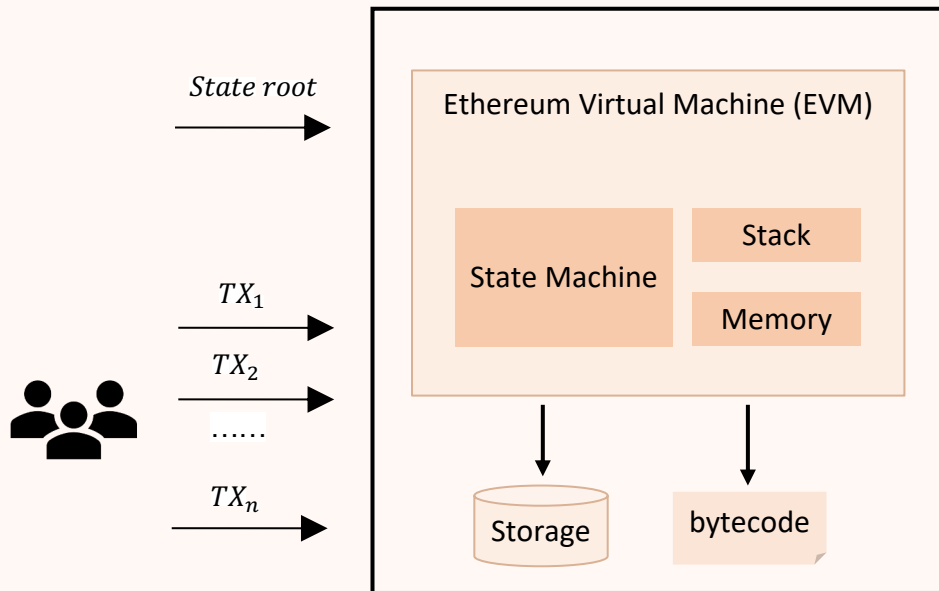


## Computation



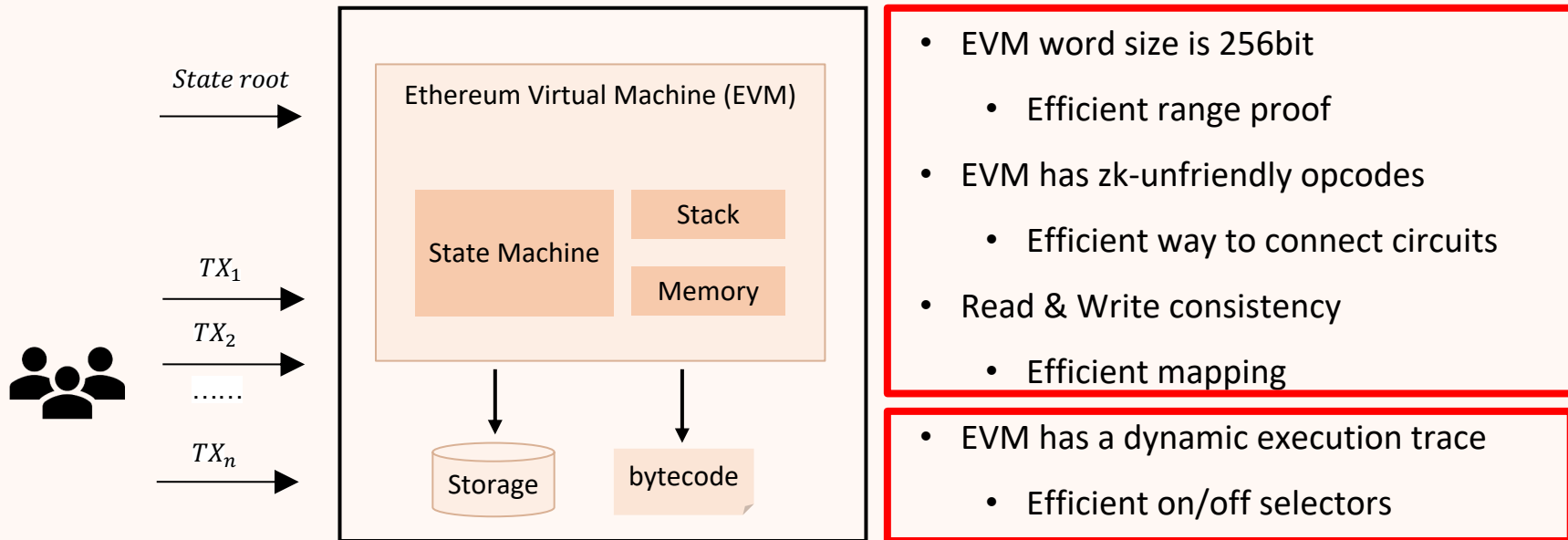
- EVM word size is 256bit
  - Efficient range proof
- EVM has zk-unfriendly opcodes
  - Efficient way to connect circuits
- Read & Write consistency
  - Efficient mapping
- EVM has a dynamic execution trace
  - Efficient on/off selectors

## Computation

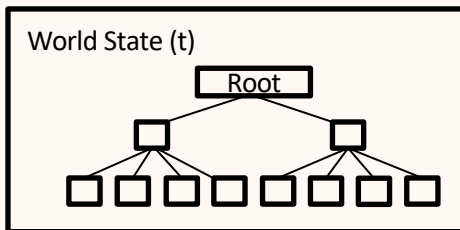


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## Computation



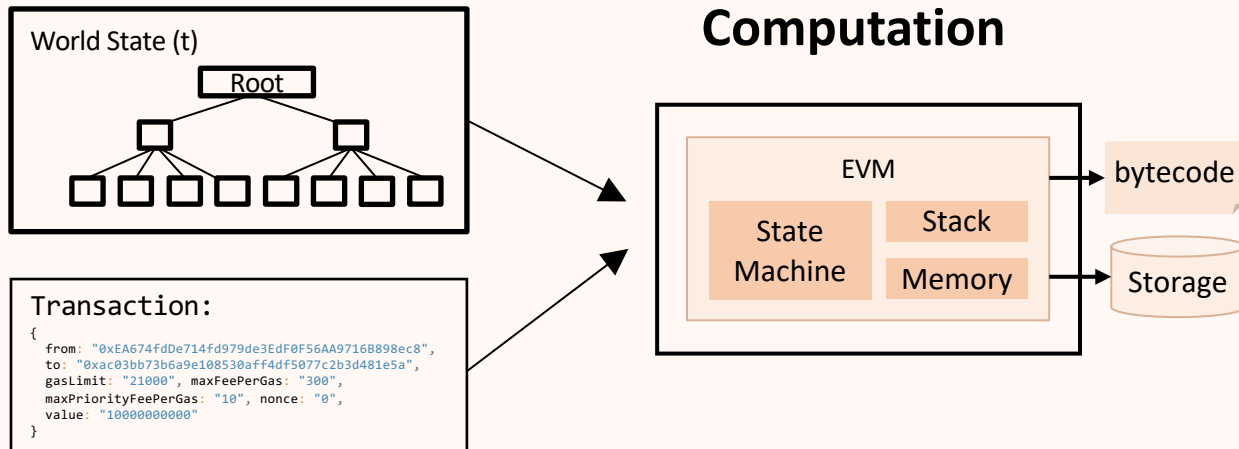
# What you need to prove



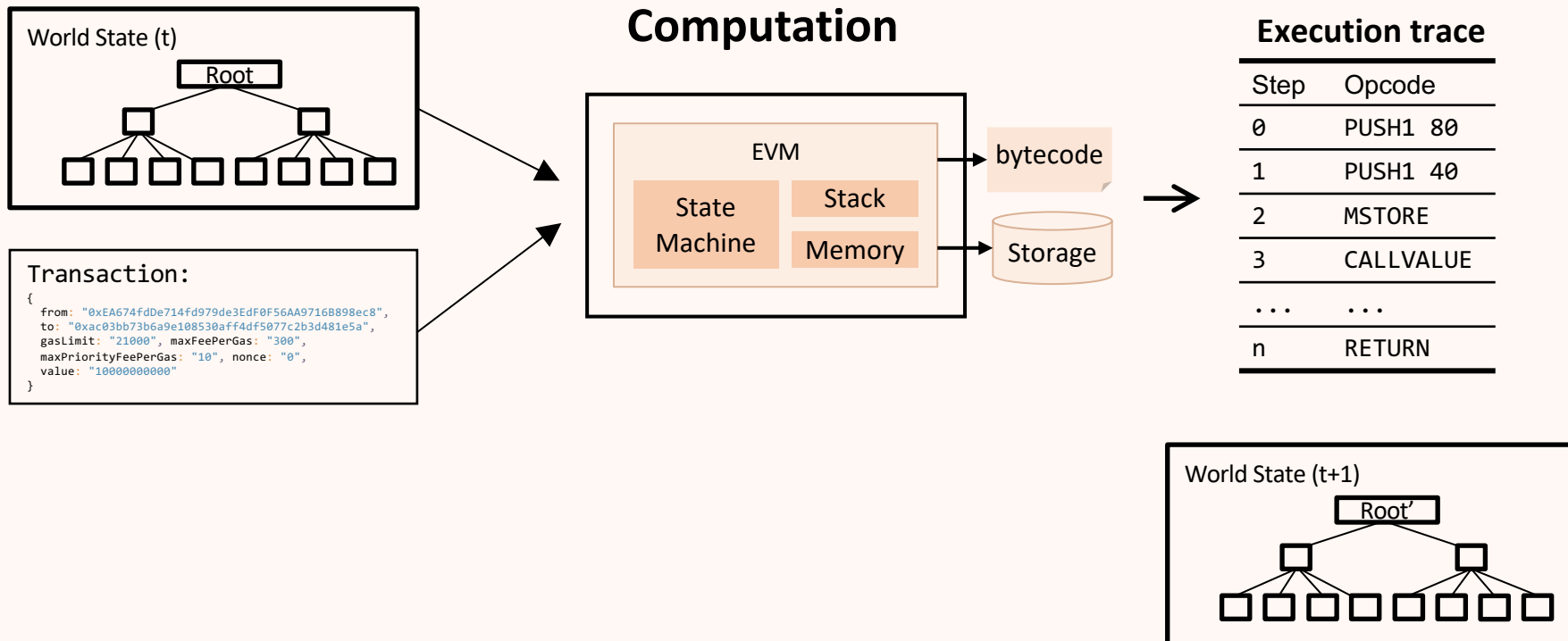
## Transaction:

```
{
  from: "0xEA674fdDe714fd979de3EdF0F56AA97168898ec8",
  to: "0xac03bb73b6a9e108530aff4df5077c2b3d481e5a",
  gasLimit: "21000", maxFeePerGas: "300",
  maxPriorityFeePerGas: "10", nonce: "0",
  value: "10000000000"
}
```

# What you need to prove

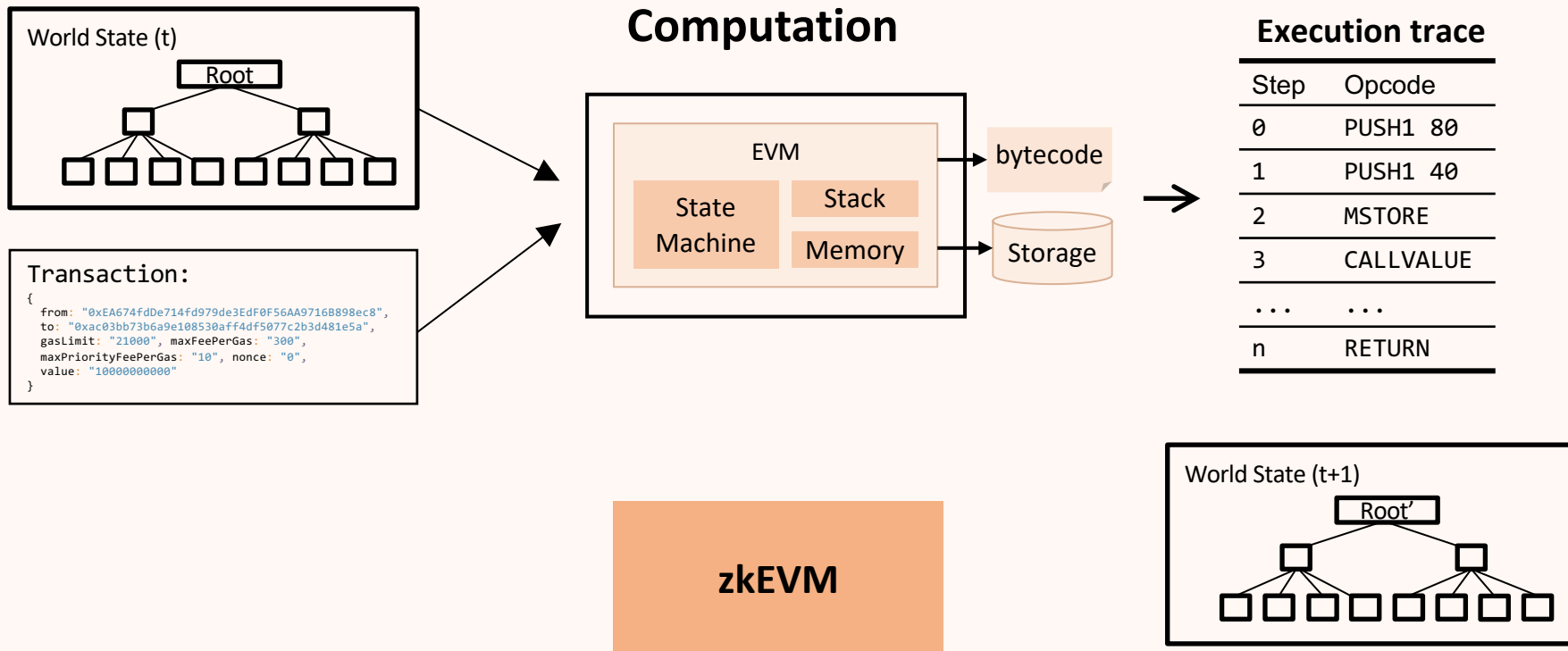


# What you need to prove

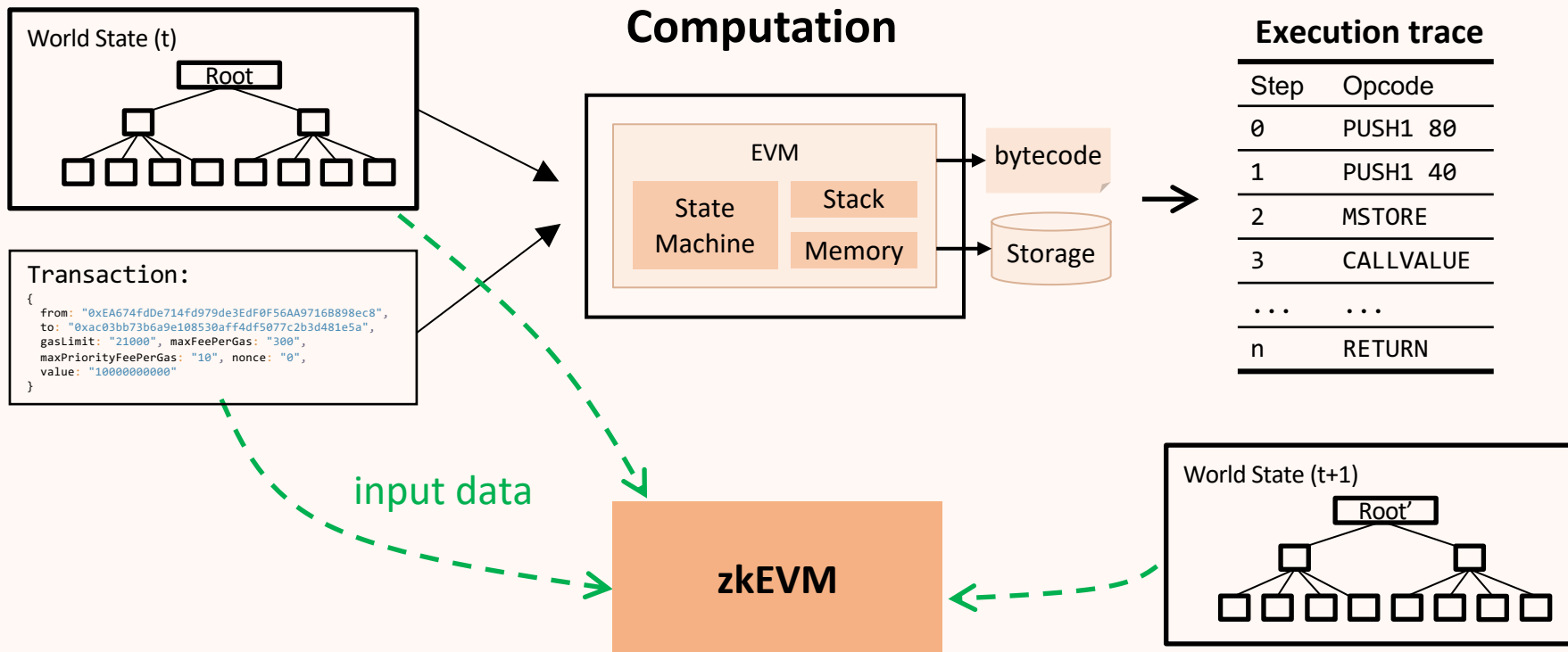




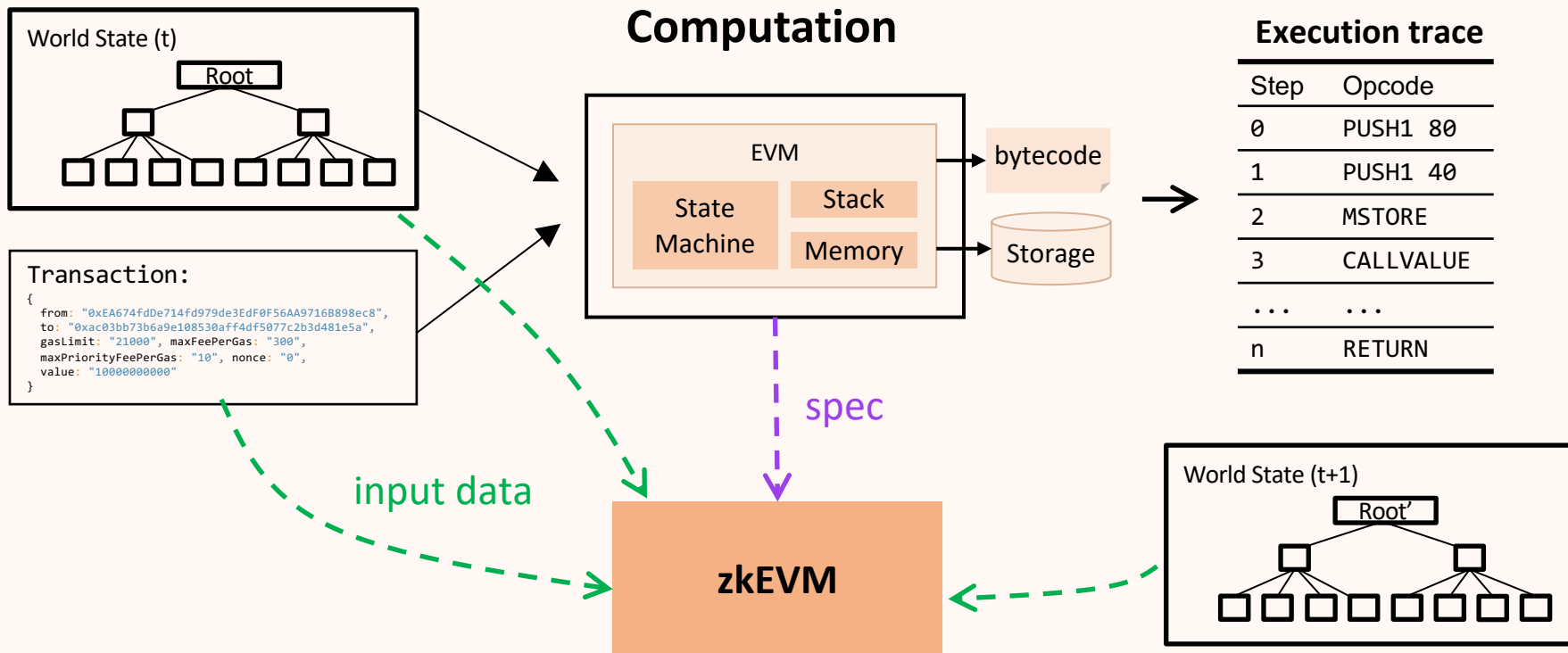
# What you need to prove



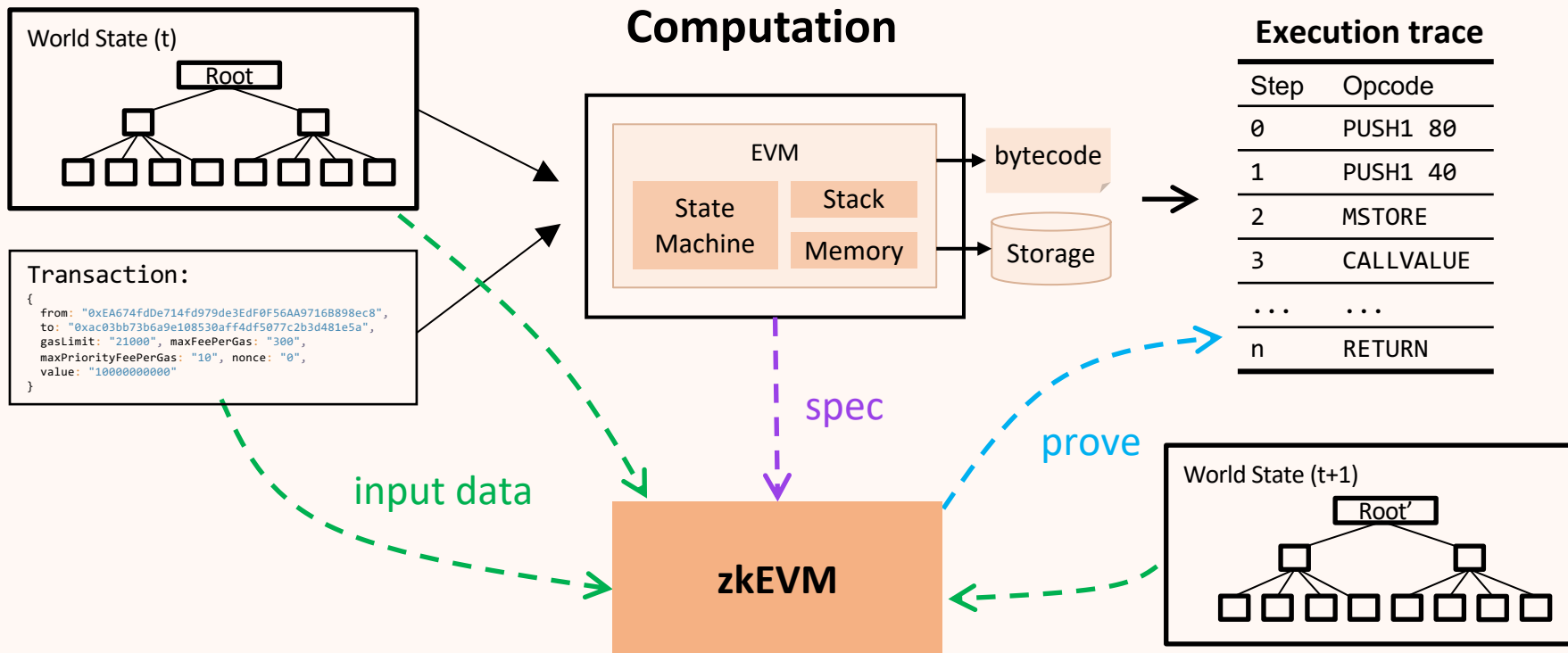
# What you need to prove

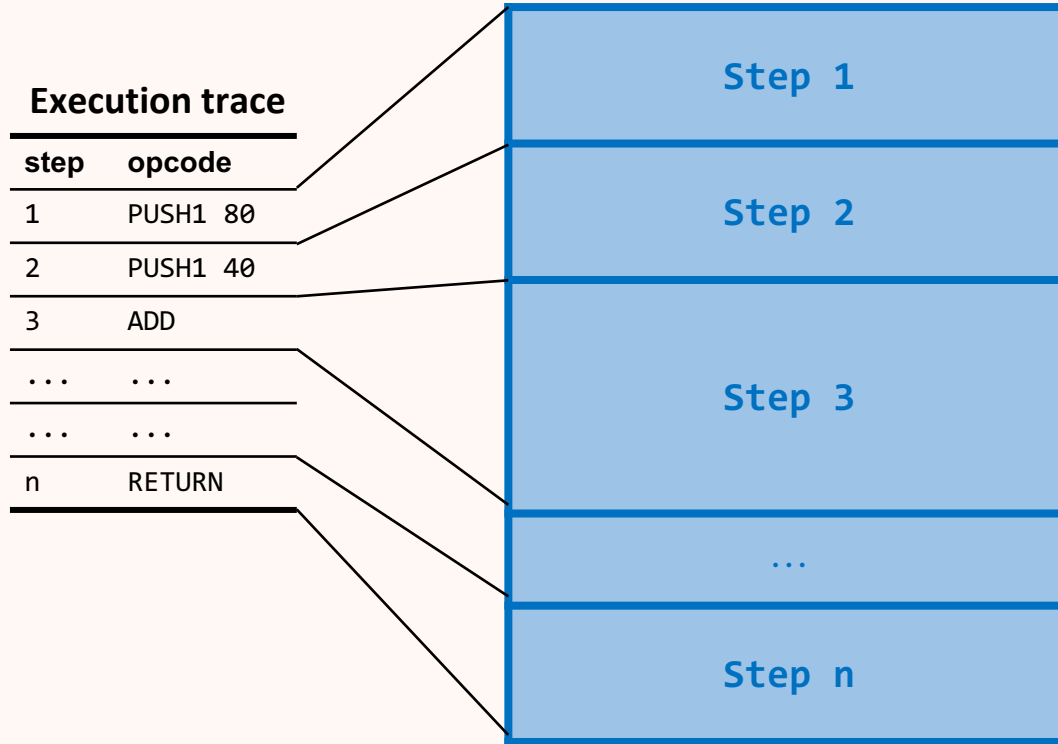


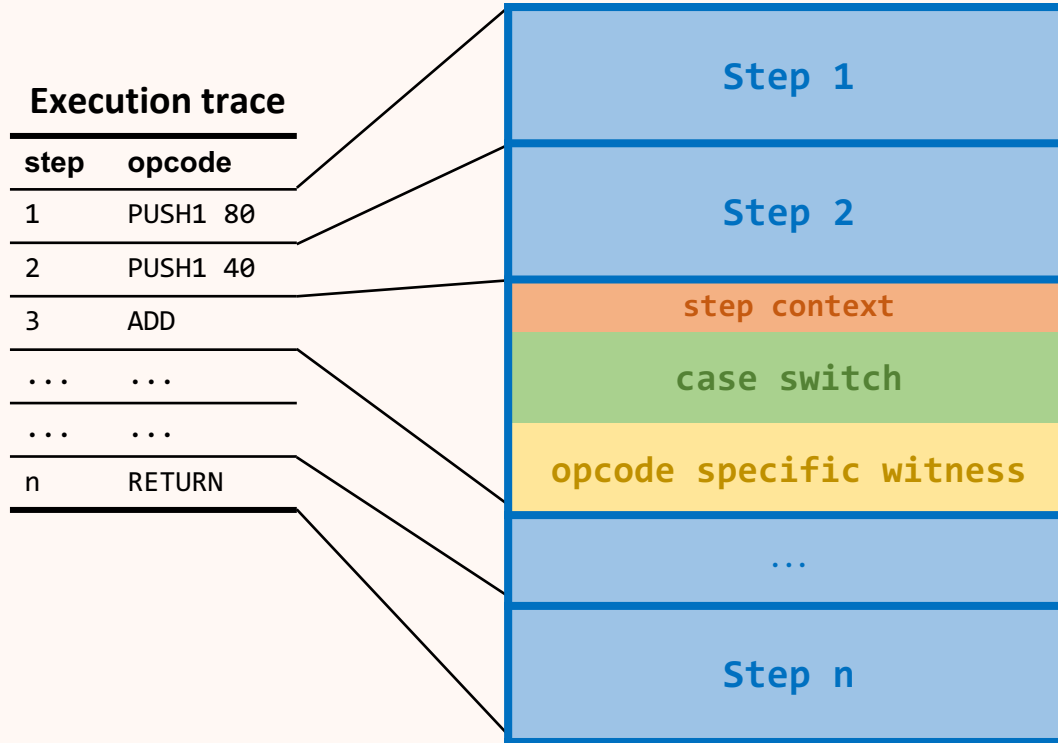
# What you need to prove

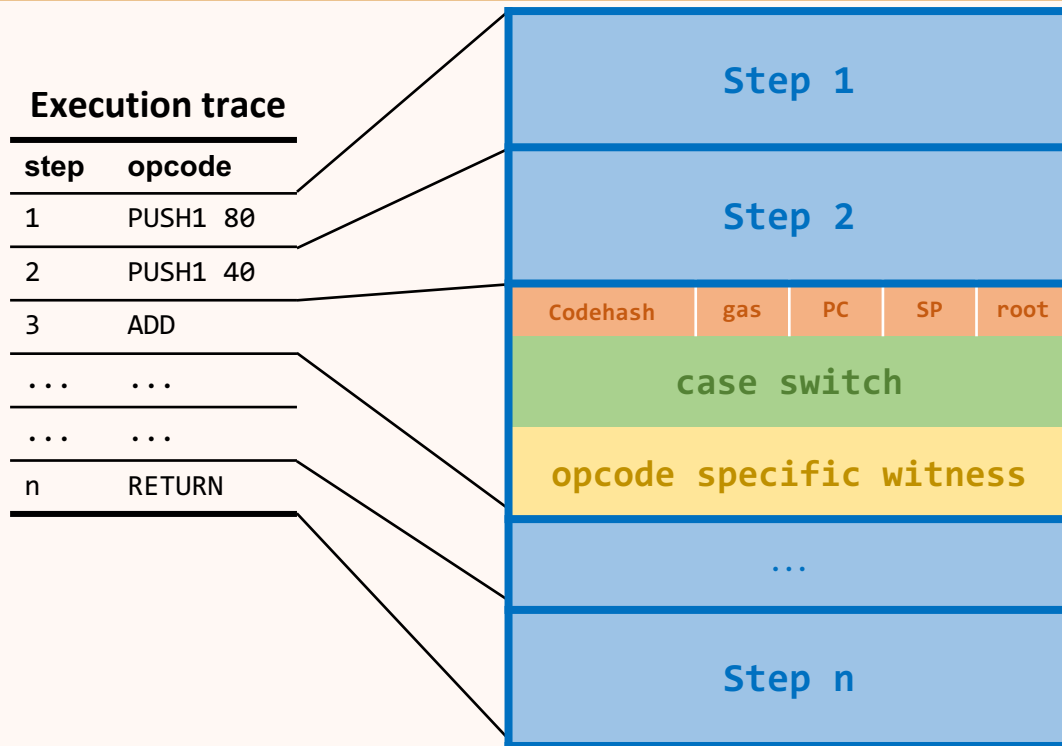


# What you need to prove



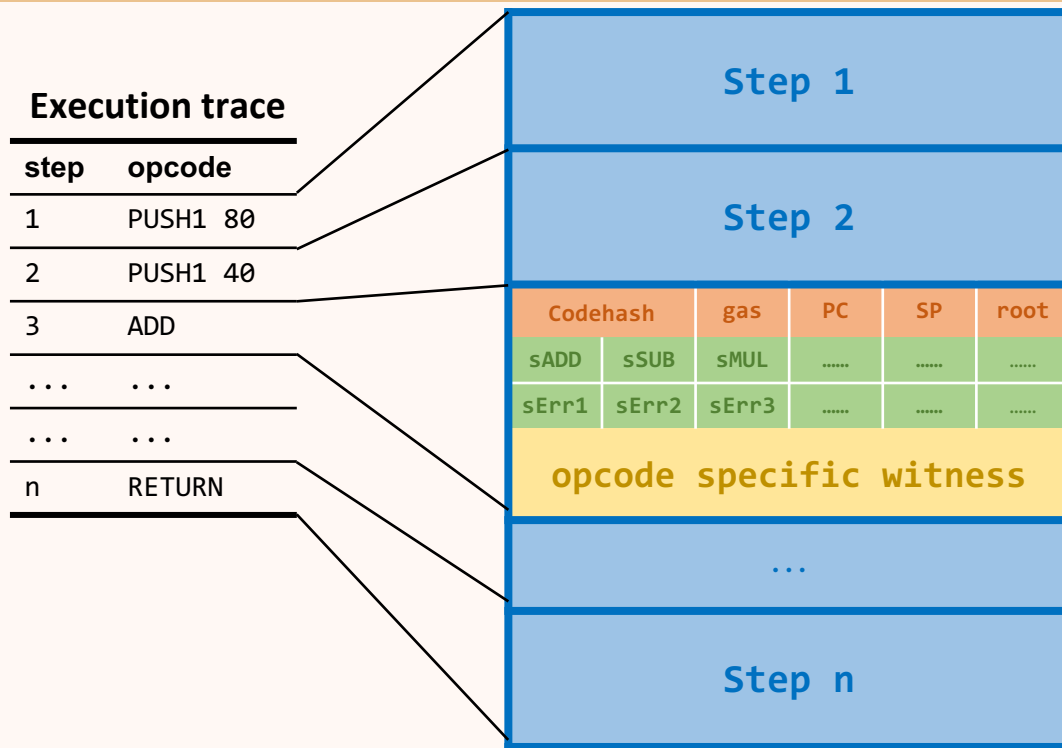






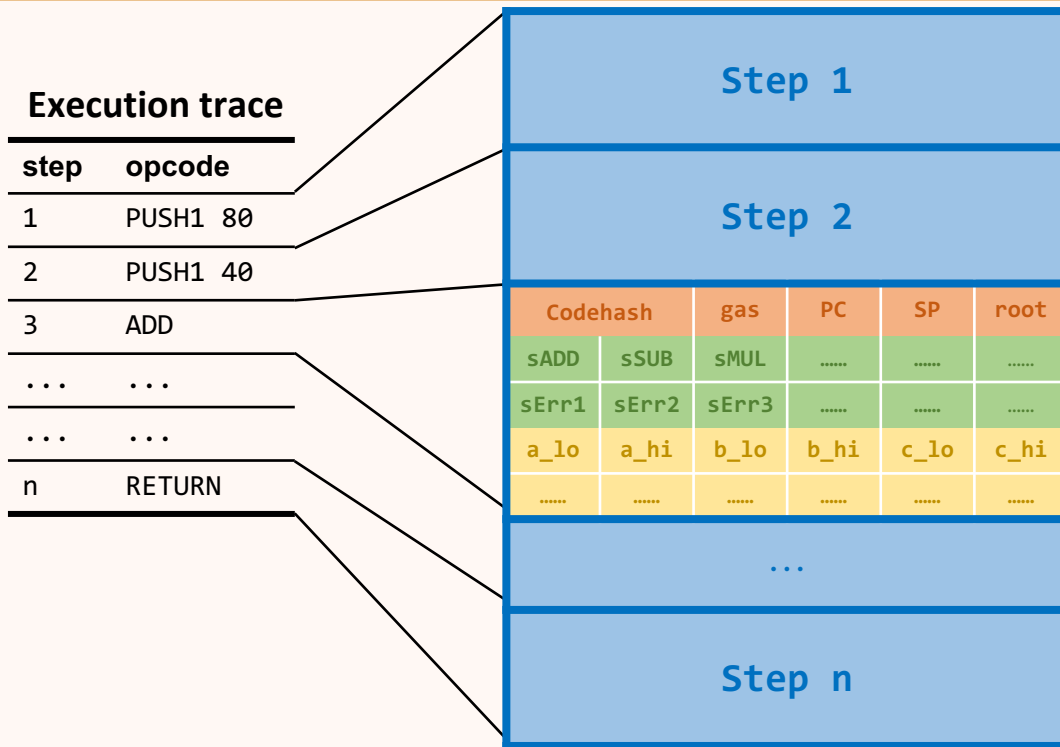
- Step context**

- Codehash
- Gas left
- Program counter, Stack pointer

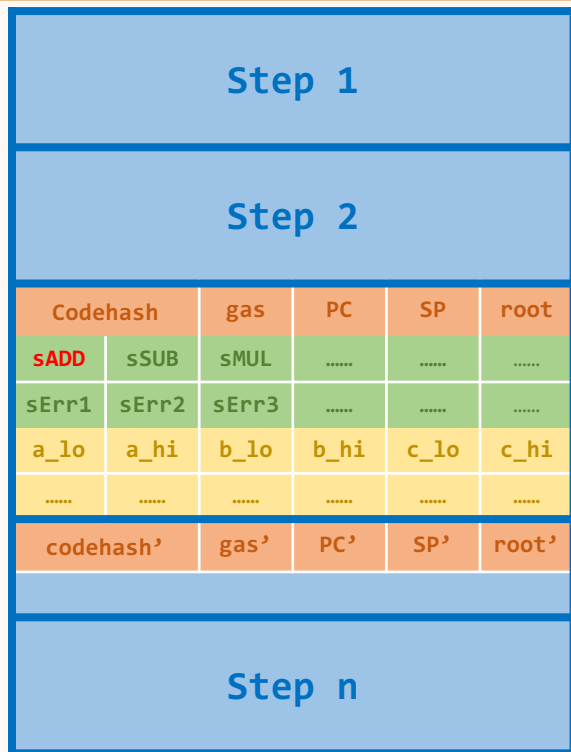


- **Step context**
  - Codehash
  - Gas left
  - Program counter, Stack pointer
- **Case switch**
  - Select opcodes & error cases
  - Exactly one is switched on





- **Step context**
  - Codehash
  - Gas left
  - Program counter, Stack pointer
- **Case switch**
  - Select opcodes & error cases
  - Exactly one is switched on
- **Opcode specific witness**
  - Extra witness used for opcodes
  - i.e. operands, carry, limbs, ...



- Step context

$$sADD * (pc' - pc - 1) == 0$$

$$sADD * (sp' - sp - 1) == 0$$

$$sADD * (gas' - gas - 3) == 0$$

- Case switch

$$sADD * (1 - sADD) == 0$$

$$sMUL * (1 - sMUL) == 0$$

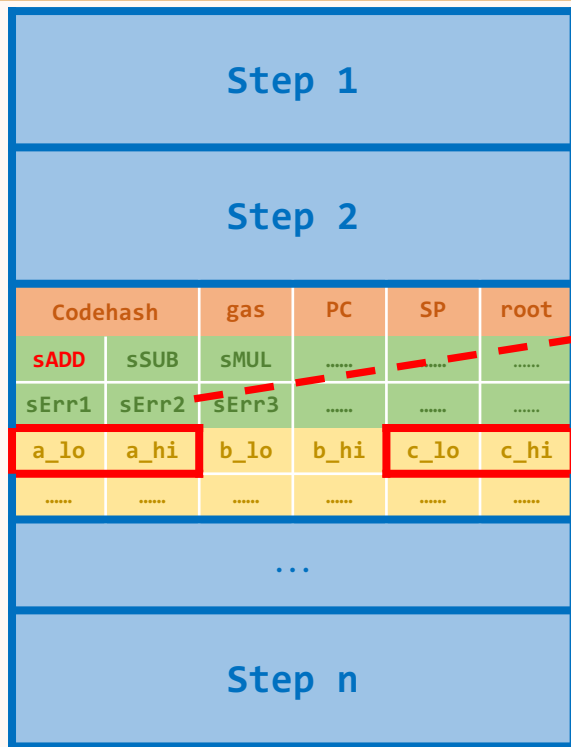
...

$$sADD + sMUL + \dots + sERRk == 1$$

- Opcode specific witness

$$sADD * (a\_lo + b\_lo - c\_lo - carry0 * 2^{128}) == 0$$

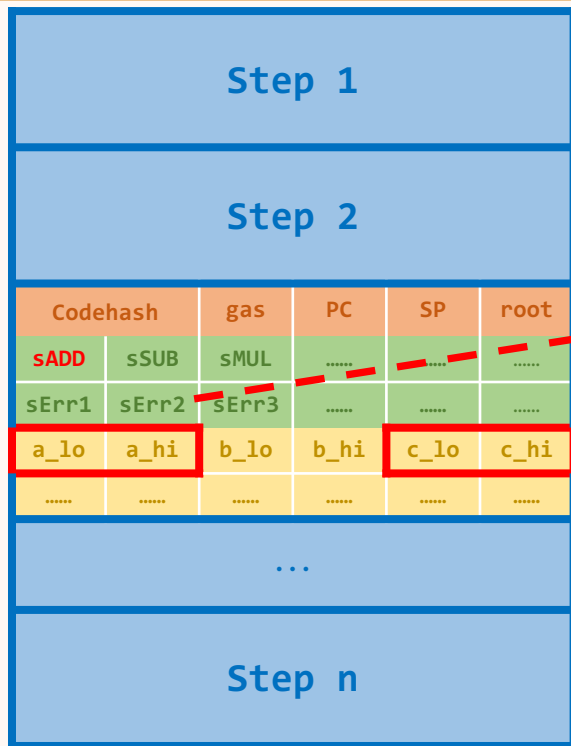
$$sADD * (a\_hi + b\_hi + carry0 - c\_hi - carry1 * 2^{128}) == 0$$



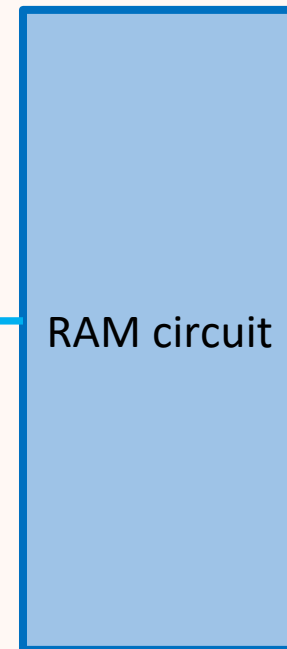
- Opcode specific witness

idx	tag	addr	R/W	value
1	STACK	1023	1	...
5	STACK	1022	0	word_a
6	STACK	1023	0	word_b
7	STACK	1023	1	word_c
...	STACK	...	...	...
...	MEMORY	0x40	1	...
...	MEMORY	...	...	...
...	STORAGE	...	...	...

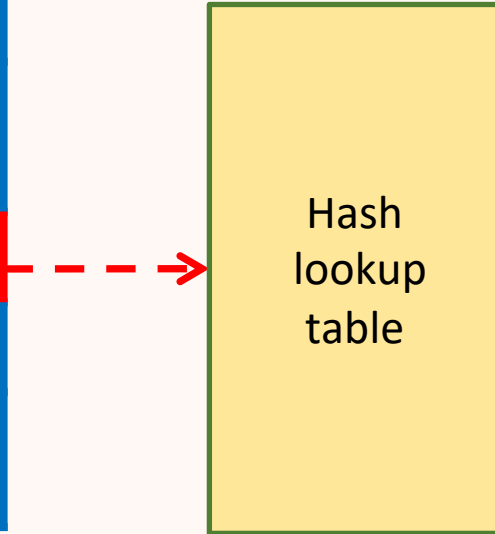
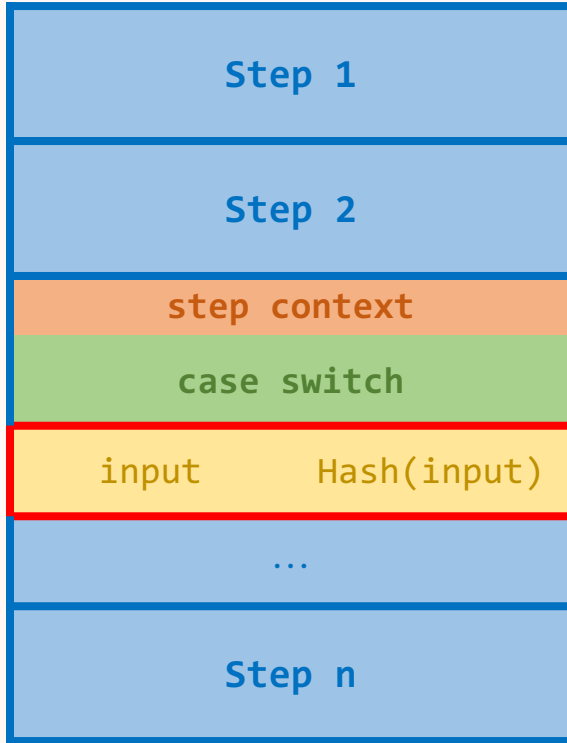
- Opcode specific witness



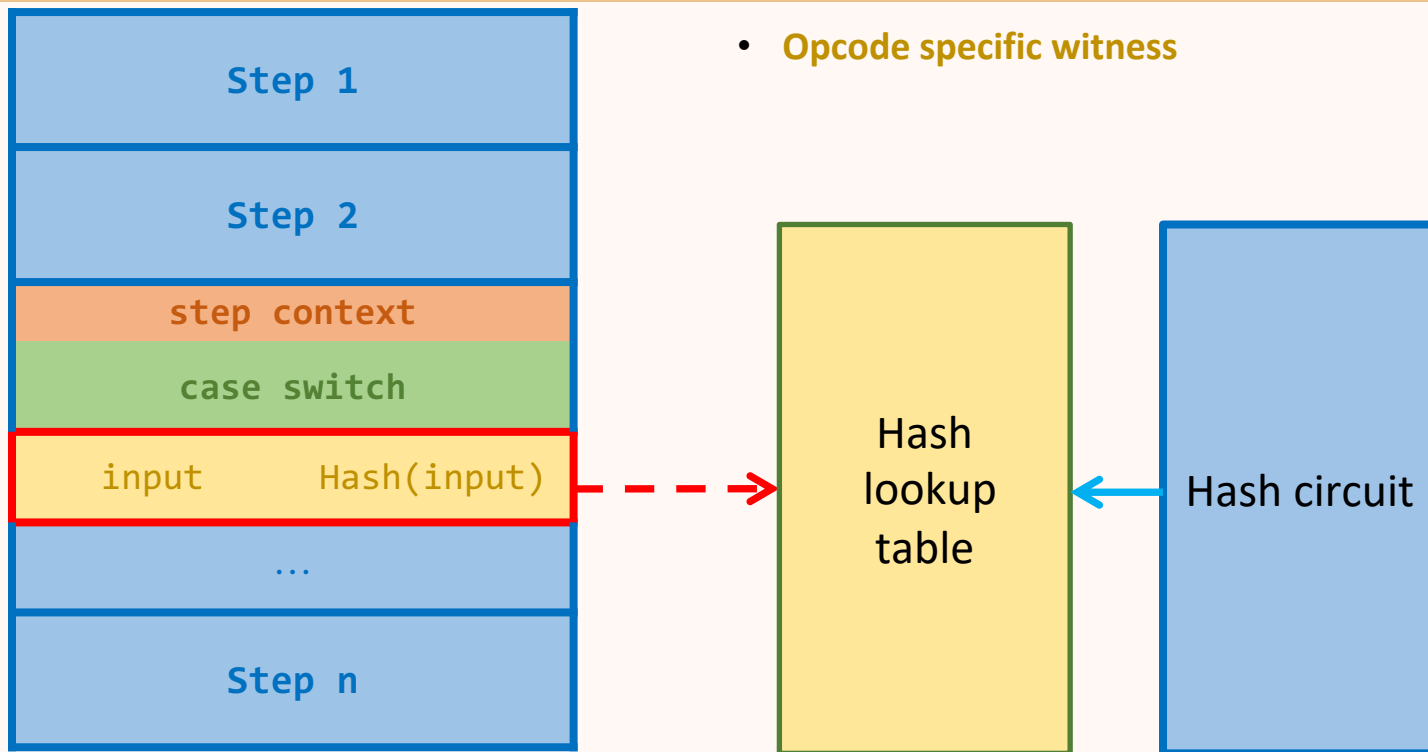
idx	tag	addr	R/W	value
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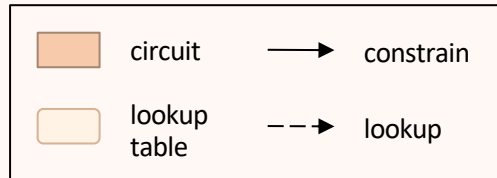
- Opcode specific witness

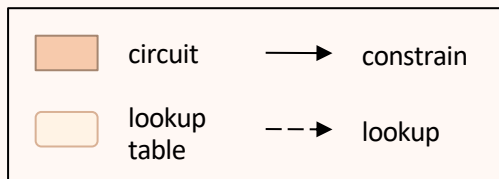
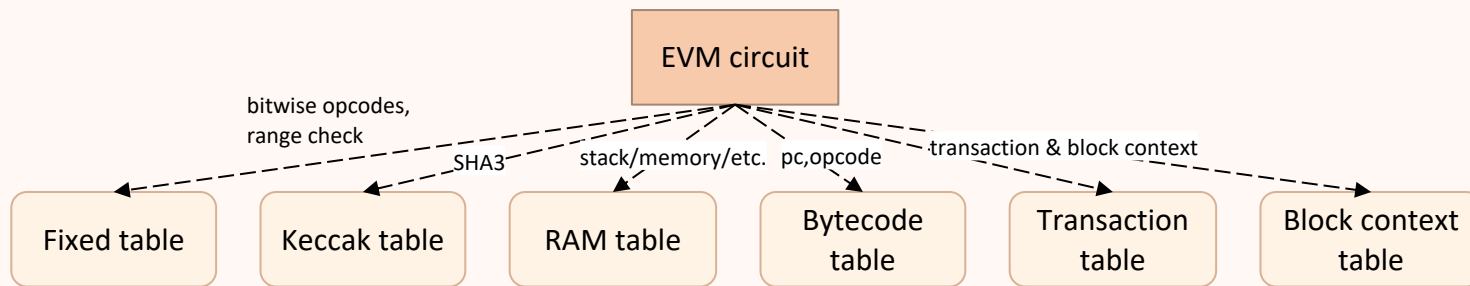


- Opcode specific witness



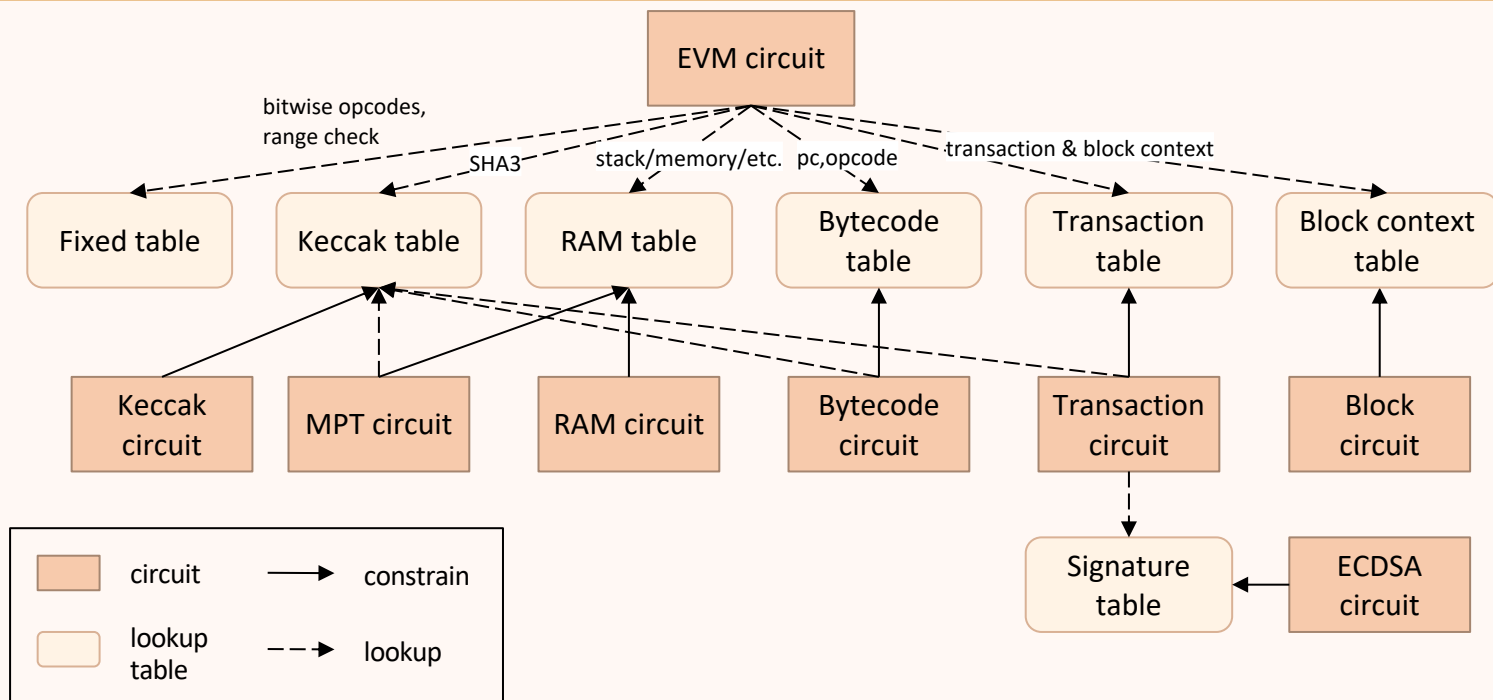
EVM circuit

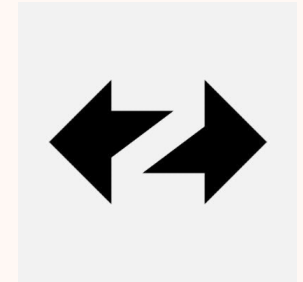
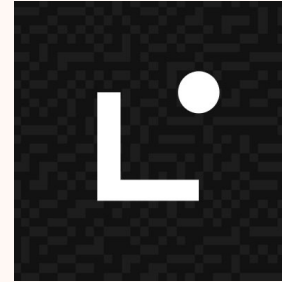




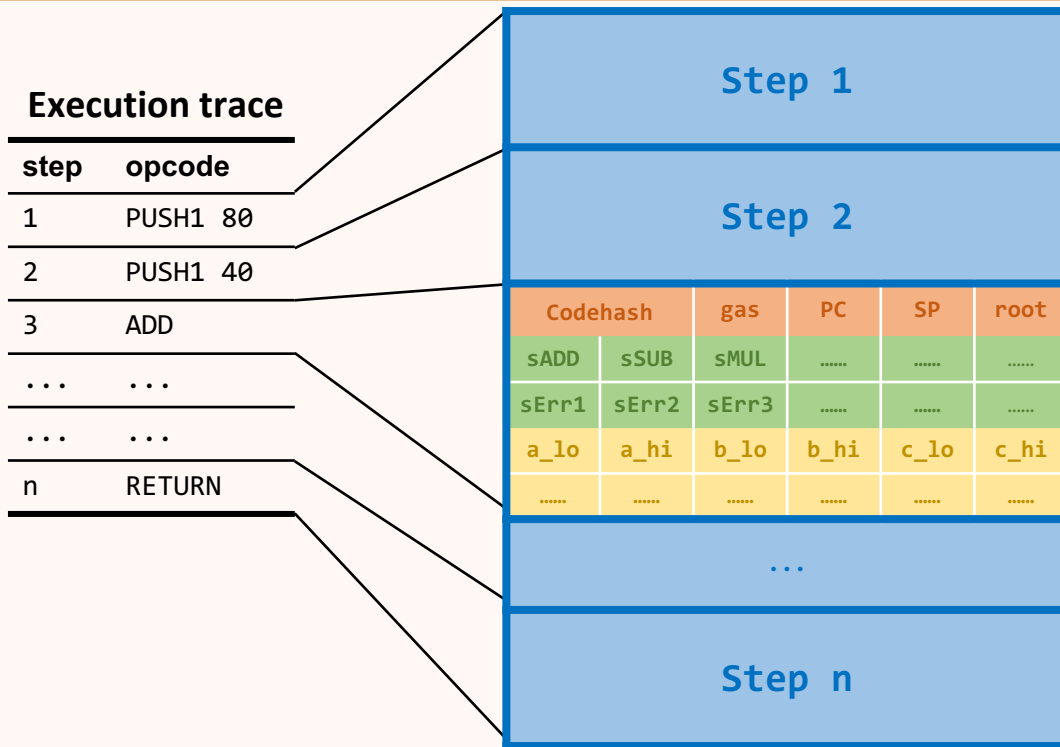


# The architecture of zkEVM circuits





Will walk through without slides :(



- Step context**

- Codehash
- Gas left
- Program counter, Stack pointer

- Case switch**

- Select opcodes & error cases
- Exactly one is switched on

- Opcode specific witness**

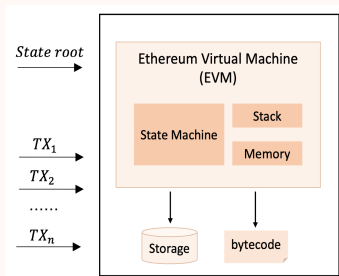
- Extra witness used for opcodes
- i.e. operands, carry, limbs, ...

- Background & motivation
- zkEVM circuit arithmetization
- zkEVM prover optimization
- Something interesting to share

# The workflow of zero-knowledge proof

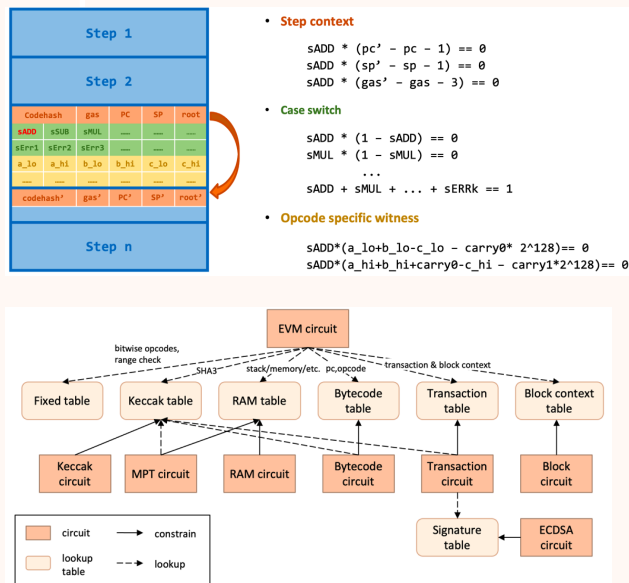


## Program



➔  
R1CS  
Plonkish  
AIR

## Constraints

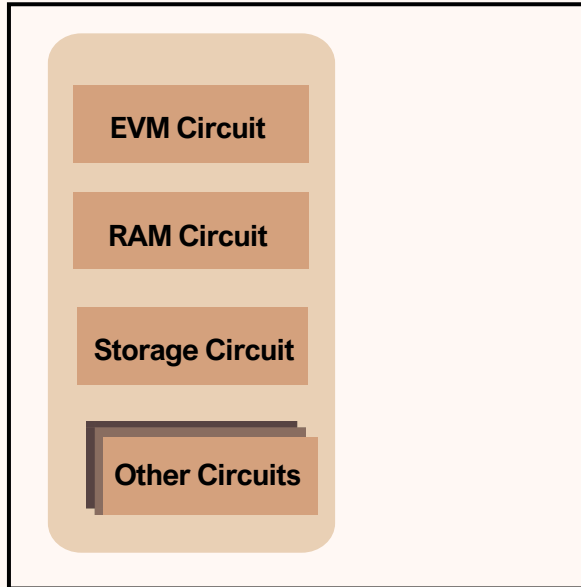


## Proof

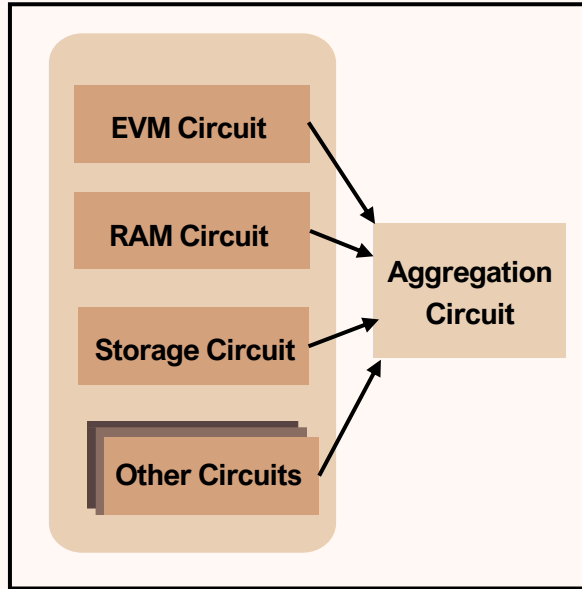
➔  
Plonk IOP  
+  
KZG



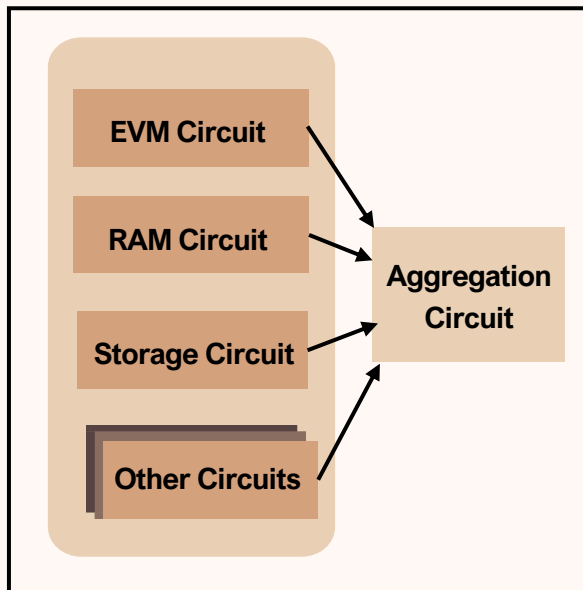
## zkEVM



## zkEVM



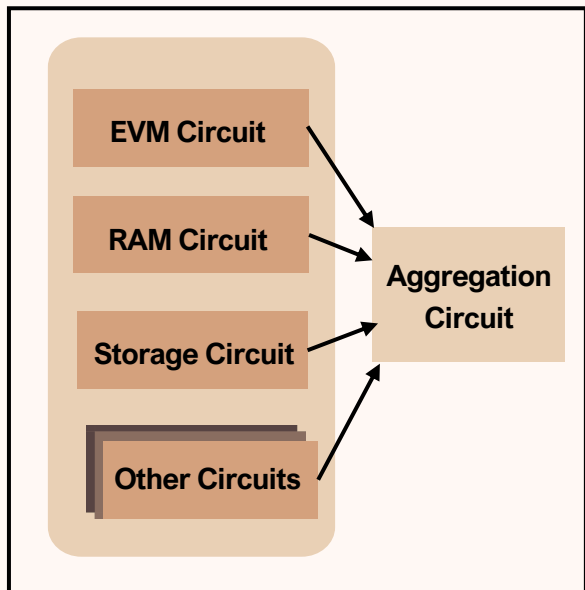
## zkEVM



- **The first layer needs to handle large computation**
  - Custom gate, Lookup support (“expressive”, customized)
  - Hardware friendly prover (parallelizable, low peak memory)
  - The verification circuit is small
  - Transparent or Universal trusted setup
- Some promising candidates
  - Plonky2/Starky /eSTARK
  - Halo2/Halo2-KZG
  - New IOP without FFTs (i.e. HyperPlonk, Plonk without FFT)
  - If Spartan/Virgo/... (sumcheck based) or Nova can support Plonkish

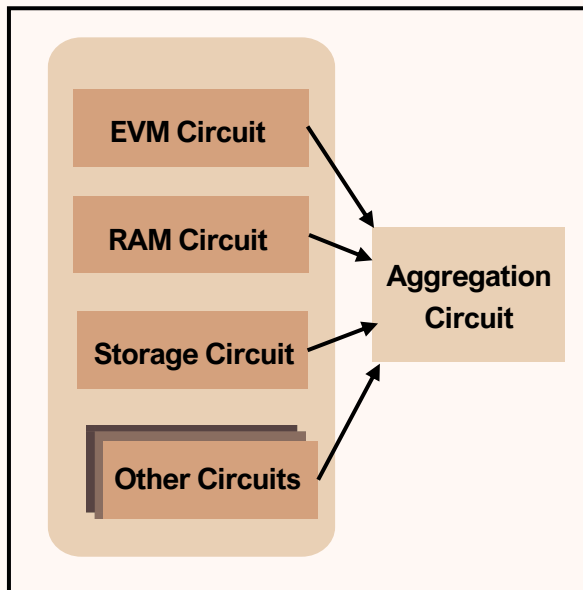


## zkEVM



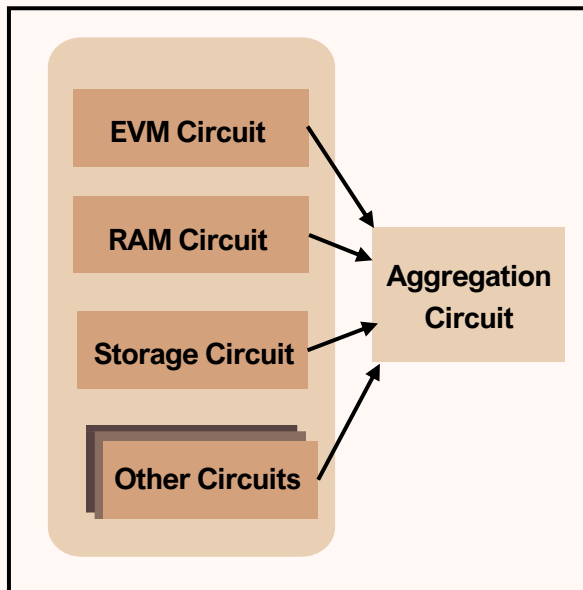
- **The second layer needs to be verifier efficient (in EVM)**
  - Proof is efficiently verifiable on EVM (small proof, low gas cost)
  - Prove the verification circuit of the former layer efficiently
  - Ideally, hardware friendly prover
  - Ideally, transparent or universal trusted setup
- Some promising candidates
  - Groth16
  - Plonk with very few columns
    - KZG/Fflonk/Keccak FRI (larger code rate)

## zkEVM



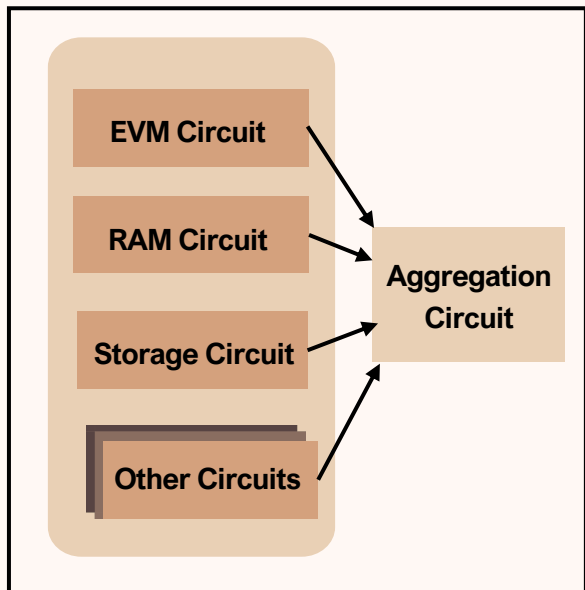
- **The first layer is Halo2-KZG** (Poseidon hash transcript)
  - Custom gate, Lookup support
  - Good enough prover performance (GPU prover)
  - The verification circuit is “small”
  - Universal trusted setup
- **The second layer is Halo2-KZG** (Keccak hash transcript)
  - Custom gate, Lookup support (express non-native efficiently)
  - Good enough prover performance (GPU prover)
  - The final verification cost can be configured to be really small

## zkEVM



- The first layer needs to be “expressive”
  - EVM circuit has **143 columns, 5700 custom gates, 58 lookups**
  - Highest custom gate degree: 9
  - For 1M gas, EVM circuit needs  **$2^{18}$  rows** (more gas, more rows)
- The second layer needs to aggregate proofs into one proof
  - Aggregation circuit has **23 columns, 1 custom gate, 7 lookups**
  - Highest custom gate degree: 5
  - For aggregating EVM, RAM, Storage circuits, it needs  **$2^{25}$  rows**

## zkEVM



- Our GPU prover optimization
  - MSM, NTT and quotient kernel
  - Pipeline and overlap CPU and GPU computation
  - Multi-card implementation, memory optimization
- The Performance
  - For EVM circuit
    - CPU prover takes 270.5s, GPU prover takes **30s (9x speedup!)**
  - For Aggregation circuit
    - CPU prover takes 2265s, GPU prover takes **149s (15x speedup!)**
  - For 5M gas, first layer takes 2 minutes, second layer takes 3 minutes

- Background & motivation
- zkEVM circuit arithmetization
- zkEVM prover optimization
- Something interesting to share

# An overview of modularized proof system



Scroll

## Frontend

R1CS / Plonkish/ AIR / CCS

## Backend

Polynomial IOP + PCS

## Frontend

R1CS / Plonkish/ AIR / CCS

## Backend

Polynomial IOP + PCS

- **Different front-end is suitable for different applications**
  - R1CS is good for linear combination and general
  - Plonkish/AIR is more uniform and customized
  - CCS has smaller witness, but sparse matrix opening is expensive
- **The PCS influences concrete properties a lot**
  - Trusted setup, Security assumption
  - Prover efficiency, Proof size, Verifier efficiency

## Frontend

R1CS / Plonkish/ AIR / CCS

## Backend

Polynomial IOP + PCS

- **KZG**
  - Universal trusted setup, DLOG
  - Prover is doing MSM, Verifier is doing pairing, small proof
- **FRI-based**
  - No trusted setup, hash collision
  - Prover is doing hashes & FFTs, Verifier is doing hashes, large proof
- **IPA**
  - No trusted setup, DLOG
  - Prover is doing MSM, Verifier is doing MSM, medium proof
- **Multilinear PCS** (for sum-check based constructions)



## Frontend

R1CS / Plonkish/ AIR / CCS

## Backend

Polynomial IOP + PCS

- **Commonly used**

- Halo2: Plonkish, Plonk IOP, IPA/KZG
- Plonky2: Plonkish, Plonk IOP, FRI
- STARK: AIR, STARK IOP, FRI
- (Groth16 is based on linear PCP, not fall into this category)

- **New protocols based on multilinear PCS**

- Spartan: R1CS, Spartan IOP, IPA derived
- HyperPlonk: Plonkish, HyperPlonk IOP, KZG/FRI derived
- Nova, SuperNova, HyperNova, Protostar

- **Ecosystem**
  - Compatibility with existing libraries
  - Existing projects and gadgets
- **Implementation**
  - Industrial vs academic implementation
  - Consider the best practice (GPUs)
  - License
  - Audit
- **Standardize**
  - Community
  - ZK ASICs

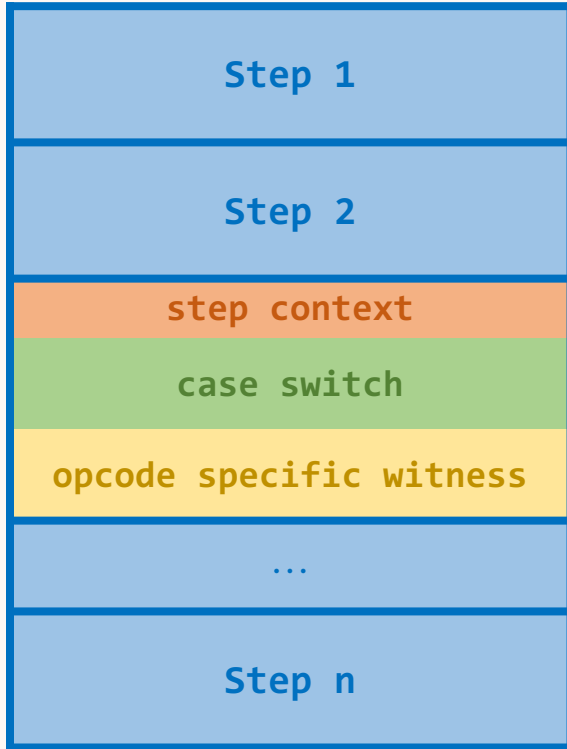
- Next generation zkEVM (faster, higher gas limit, etc)
- On-chain data compression, 4844 support
- Decentralized prover
- Decentralized sequencer
- Interoperability between zkRollups

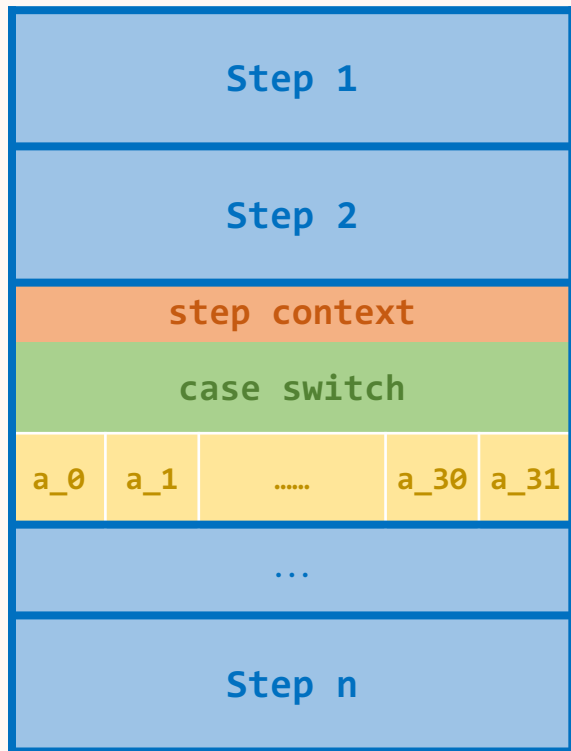
- **We are building cool things at Scroll!**
  - Scroll is a general purpose zkRollup scaling solution for Ethereum
  - Building a native zkEVM using very advanced circuit arithmetization + proof system
  - Building fast prover through hardware acceleration (GPU in production) + proof recursion
  - We are live on the mainnet with a production-level robust infrastructure
- **There are a bunch of interesting problems to be solved and we are hiring!**
  - Protocol design and mechanism design
  - Zk engineer & research for practical efficiency

# Thank you!

X / Twitter: @yezhang1998

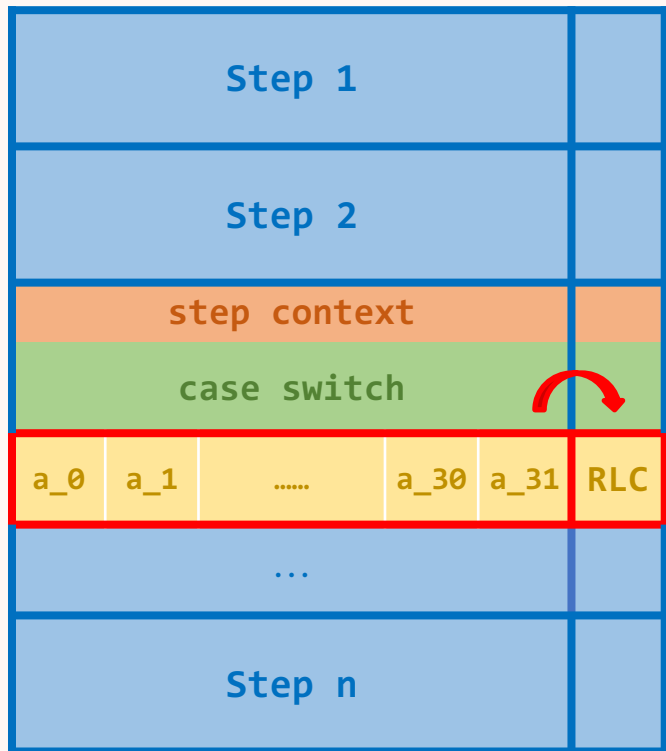






- Break down 256-bit word into 32 8-bit limbs.

$$A = a_0 + a_1 * 256 + a_2 * 256^2 + \dots + a_{31} * 256^{31}$$



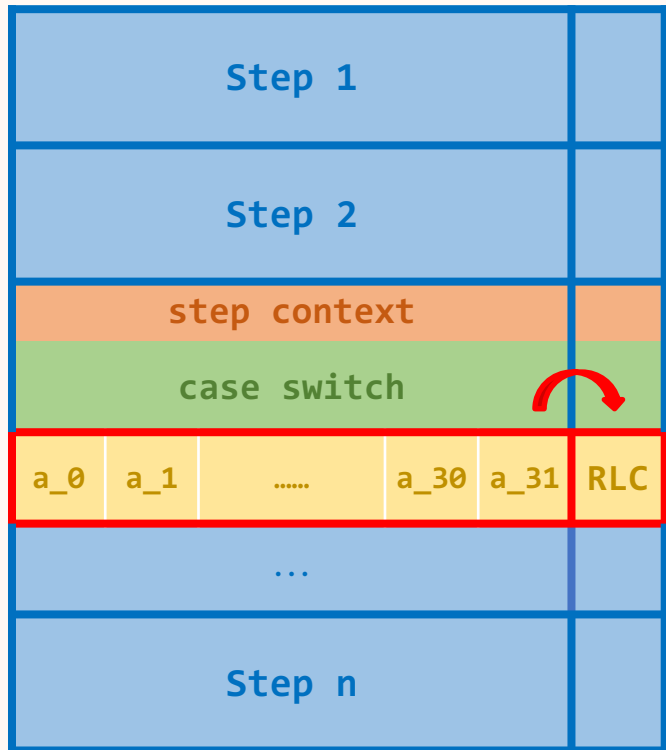
- Break down 256-bit word into 32 8-bit limbs.

$$A = a_0 + a_1 * 256 + a_2 * 256^2 + \dots + a_{31} * 256^{31}$$

- Encode EVM word using RLC (Random Linear Combination)

$$A_{RLC} \equiv a_0 + a_1 * \theta + a_2 * \theta^2 + \dots + a_{31} * \theta^{31} \pmod{F_p}$$





- Break down 256-bit word into 32 8-bit limbs.

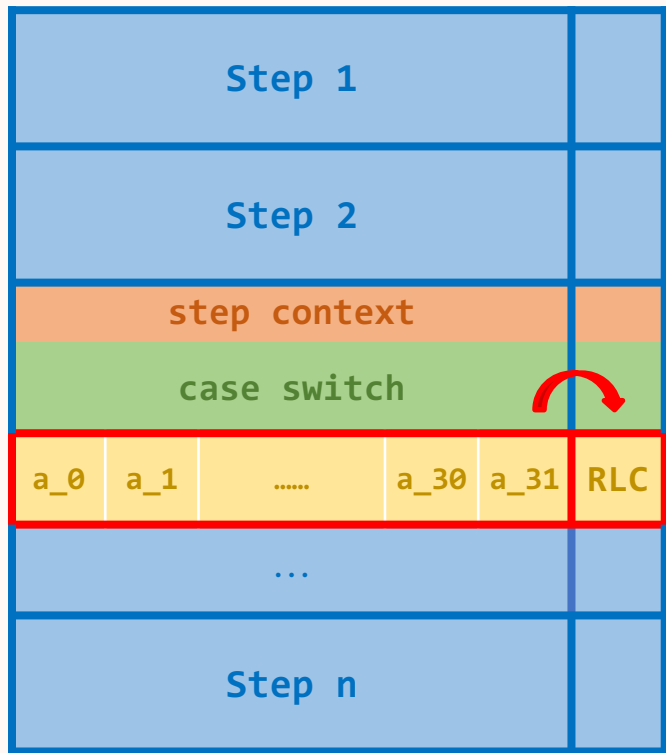
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- $\theta$  should be computed after  $a_0, \dots, a_{31}$  are fixed

- Multi-phase prover: synthesis part of witness, derive witness



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$$A = a_0 + a_1 * 256 + a_2 * 256^2 + \dots + a_{31} * 256^{31}$$

- Encode EVM word using RLC (Random Linear Combination)

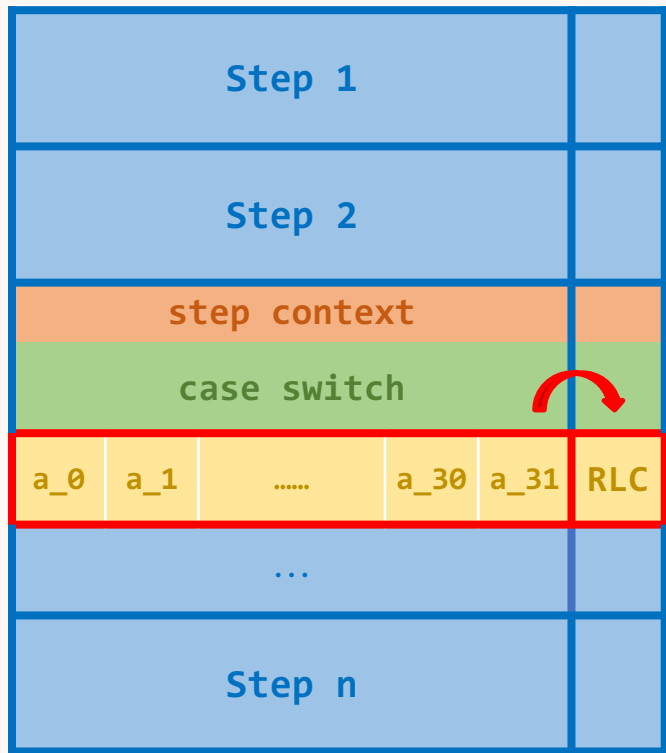
$$A_{RLC} \equiv a_0 + a_1 * \theta + a_2 * \theta^2 + \dots + a_{31} * \theta^{31} \pmod{F_p}$$

- $\theta$  should be computed after  $a_0, \dots, a_{31}$  are fixed

- Multi-phase prover: synthesis part of witness, derive witness

- RLC is useful in many places

- Compress EVM word into one value
- Encode dynamic length input
- Lookup layout optimization



- Break down 256-bit word into 32 8-bit limbs.

$$A = a_0 + a_1 * 256 + a_2 * 256^2 + \dots + a_{31} * 256^{31}$$

- Encode EVM word using RLC (Random Linear Combination)

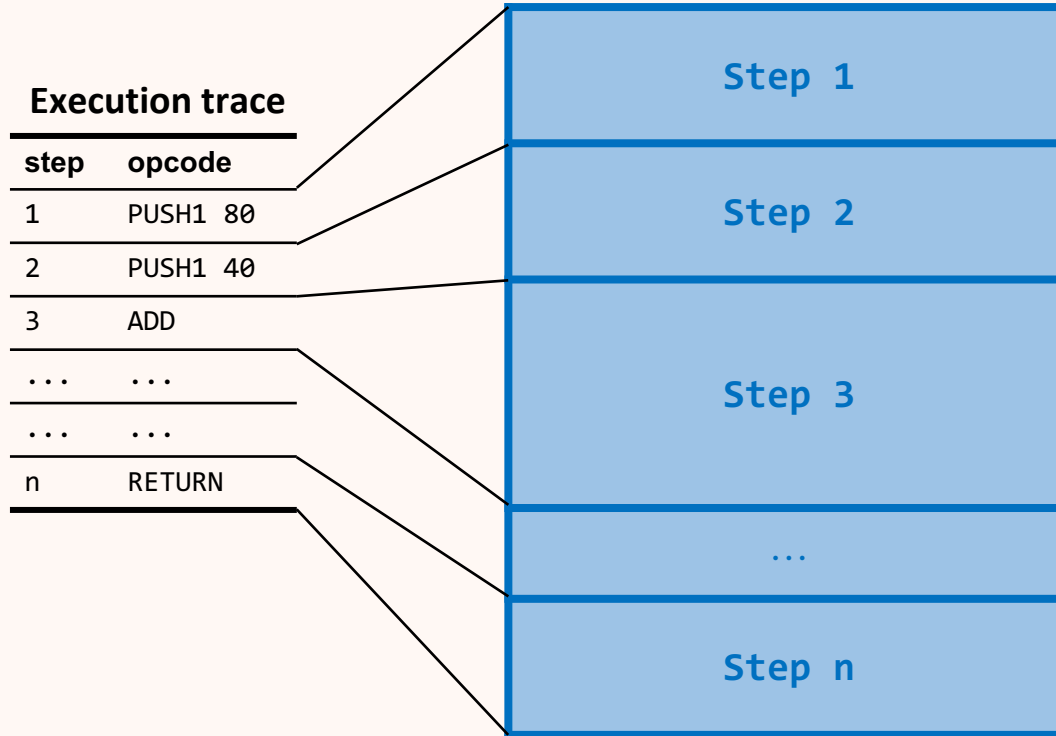
$$A_{RLC} \equiv a_0 + a_1 * \theta + a_2 * \theta^2 + \dots + a_{31} * \theta^{31} \pmod{F_p}$$

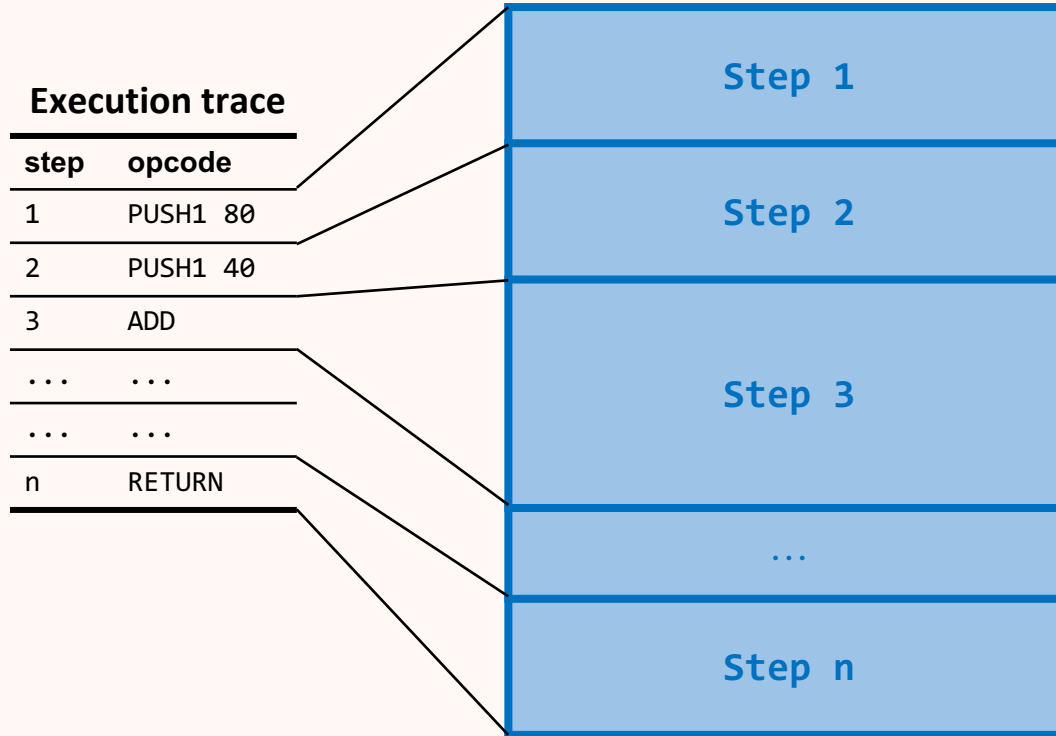
- $\theta$  should be computed after  $a_0, \dots, a_{31}$  are fixed

- Multi-phase prover: synthesis part of witness, derive witness

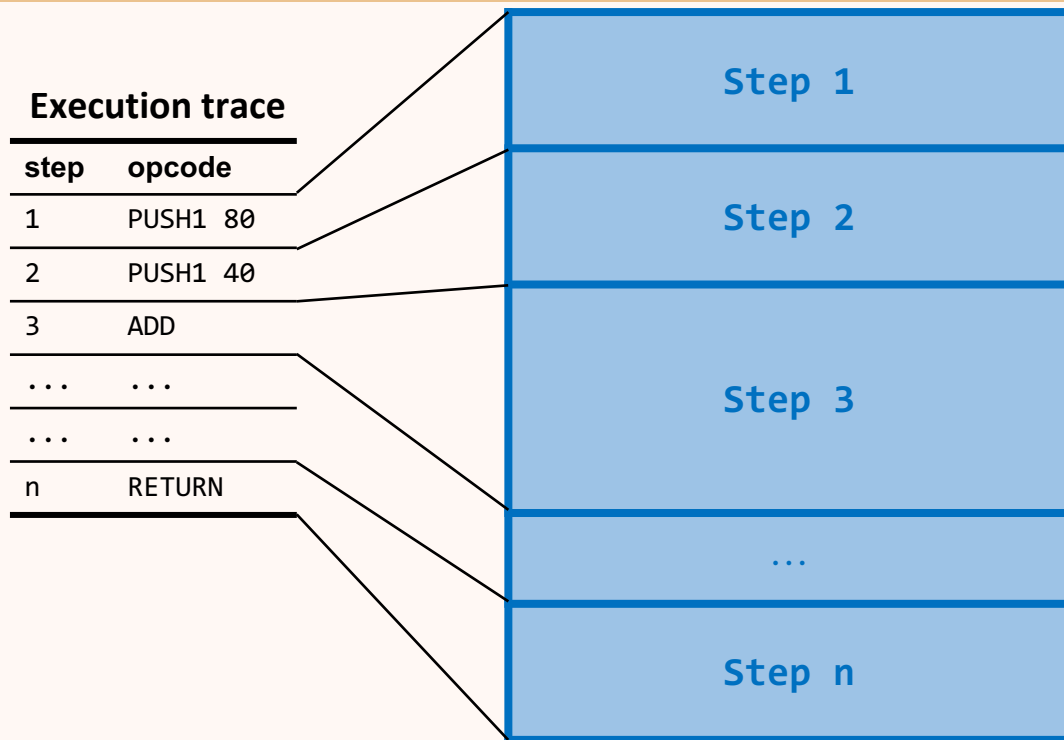
- **RLC is useful in many places, remove it?**

- Compress EVM word into one value  $\rightarrow$  high, low for EVM word
- Encode dynamic length input  $\rightarrow$  fixed chunk, dynamic times
- Lookup layout optimization



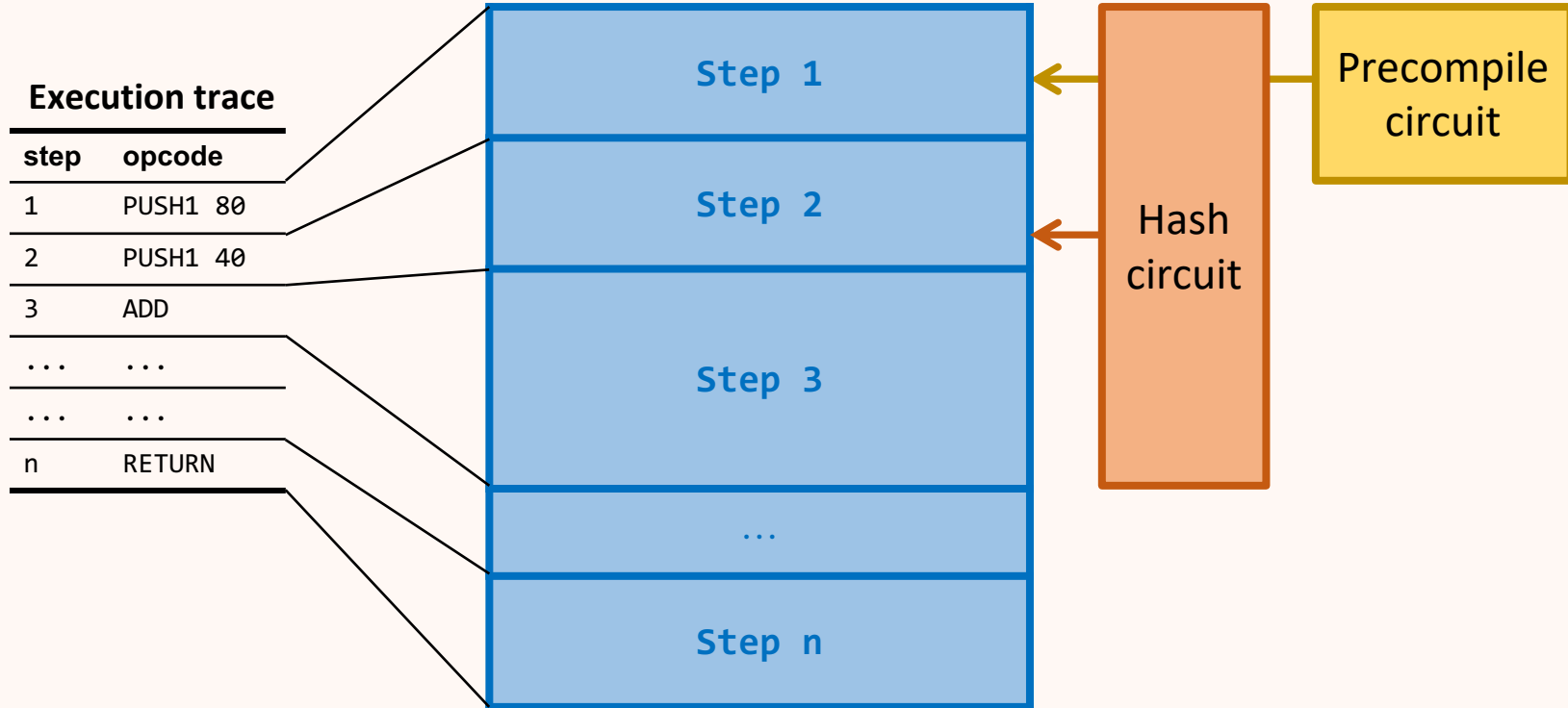


- The execution trace is dynamic
  - enable different constraints
  - permutation is not fixed
  - hard to use standard gates

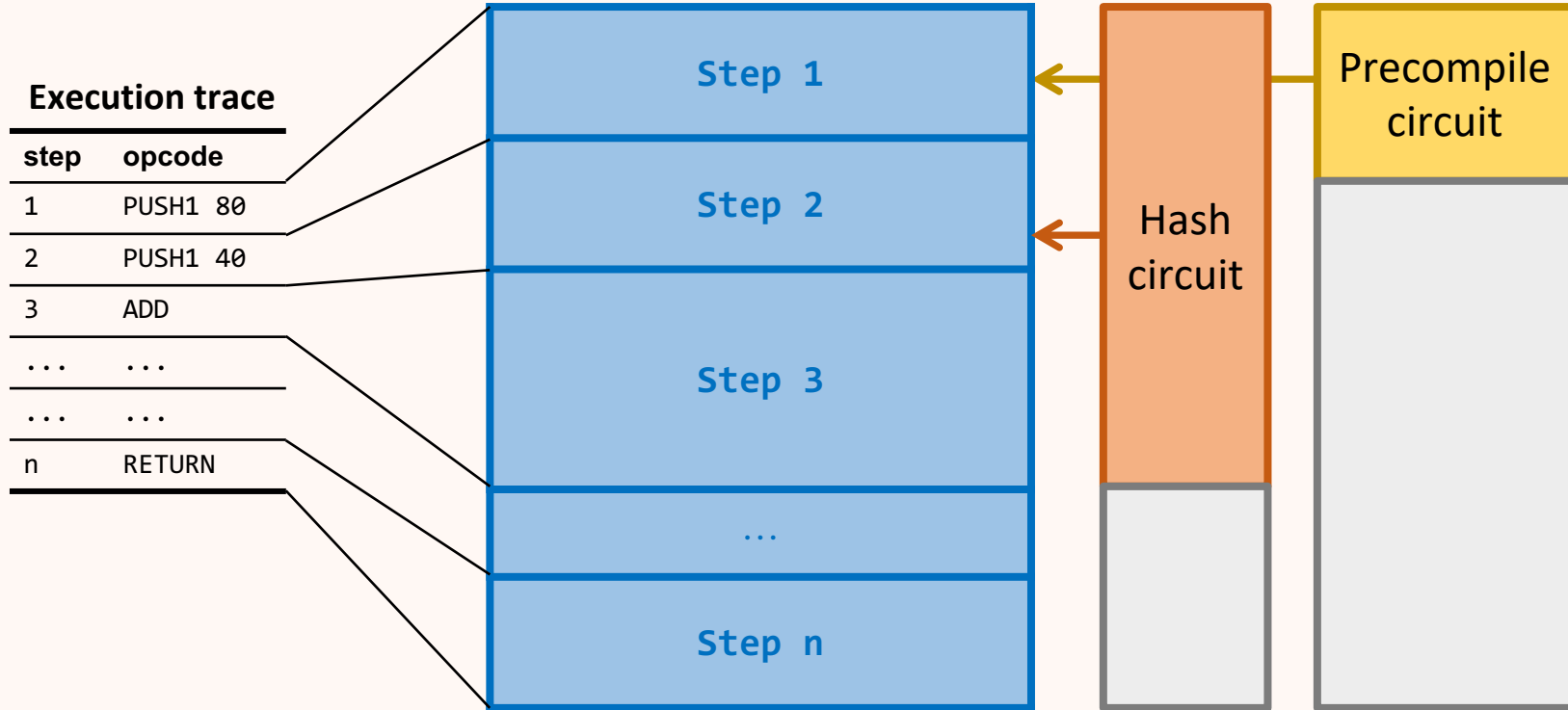


- The execution trace is dynamic
  - enable different constraints
  - permutation is not fixed
  - hard to use standard gates
- **Better way to layout?**
  - We have 2000+ custom gates
  - Different rotation to access cells

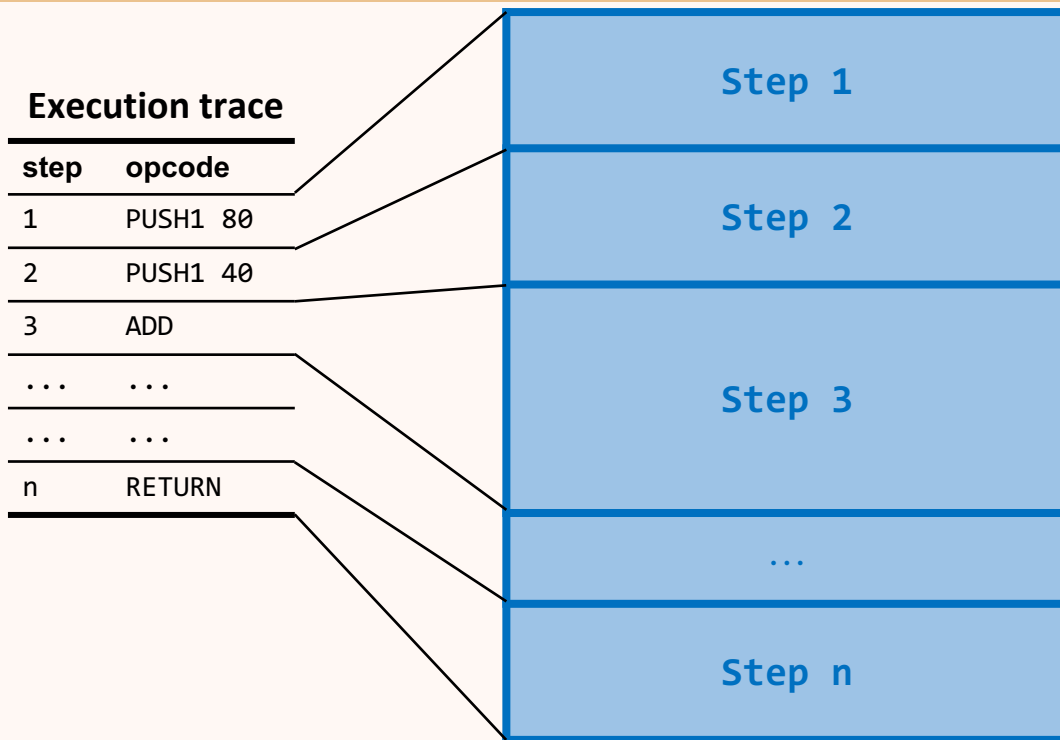
# Circuit - Dynamic size



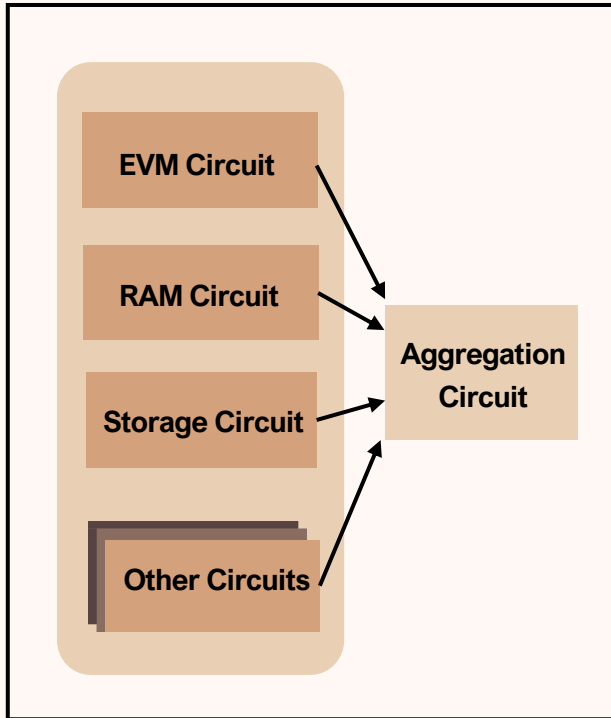
# Circuit - Dynamic size

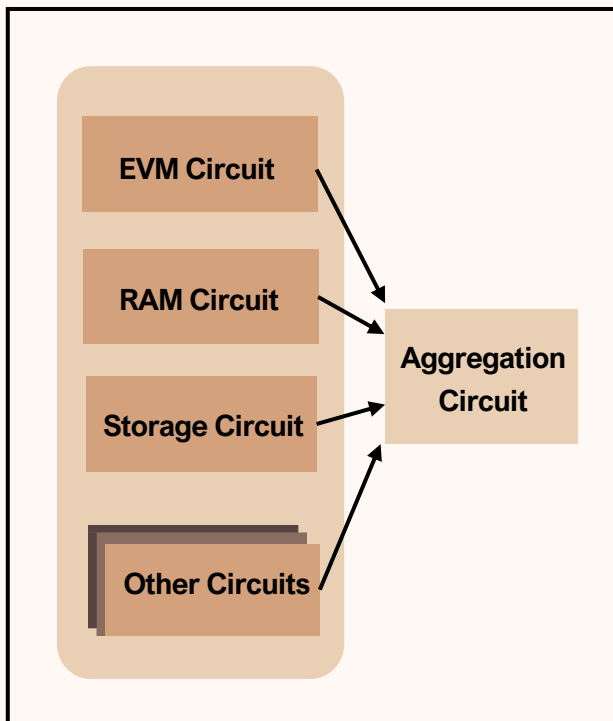




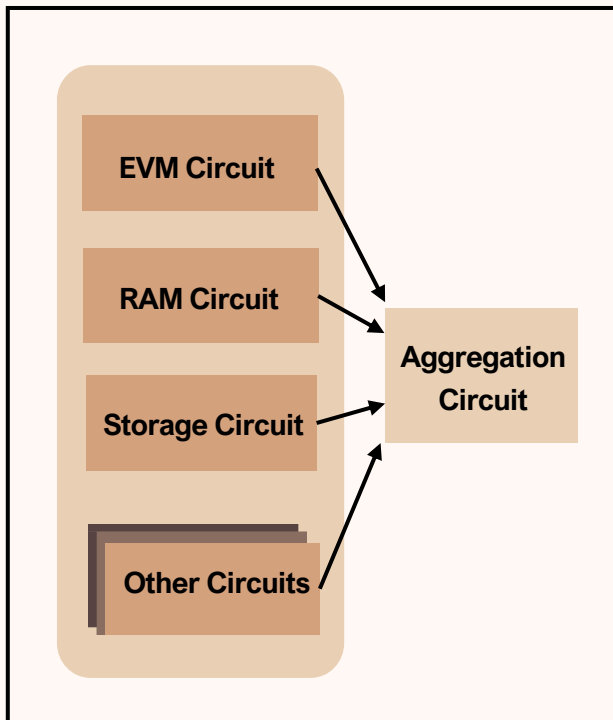


- Some bad influences
  - i.e. Mload is more costly (more rows)
  - i.e. Maximum number of Keccaks
  - i.e. Pay larger proving cost for padding
- **Can we make zkEVM dynamic?**





- Our prover
  - GPU can make MSM & NTT really fast  
Bottleneck moves to witness generation & data copy
  - Need large CPU memory (1TB -> 300GB+)
- **Hardware friendly prover?**
  - Parallelizable & Low peak memory
  - Don't ignore the witness generation
  - Run on cheap machines, more decentralized



- **Best way to compose different proof system?**
  - The first layer needs to be “expressive”
  - The second layer needs to be verifier efficient (in EVM)
  - **Should we move to smaller field?**  
(Breakdown/FRI with Goldilocks, Mersenne prime)
  - **Should we stick to EC-based constructions?**  
(SuperNova, Cyclic elliptic curve with fast MSM)
  - More options waiting for you → Reach out to us!