# CSE 4125: Distributed Database Systems Chapter – 4 (Part – B)

Distributed Database Design

# The Design of Horizontal Fragmentation

## Horizontal Fragments

A horizontal fragment  $R_i$  of a relation R consists of all the tuples of R that satisfy a minterm predicate  $m_i$ .

#### Which minterm predicate should we use?

-We have to decide on the set of simple predicates that are the basis for the minterm predicates.

# Desirable properties of the set of simple predicates

The set should be complete and minimal.

Informally, the set should include only predicates with attributes and conditions that are used in the applications

### Example

Consider the following global relation, fragmentation schema and applications.

J	<u>JNO</u>	JNAME	BUDGET	LOC
	J1	Instrumental	150,000	Montreal
	J2	GUI .	135,000	New York
	J3	CAD/CAM	250,000	New York
	J <b>4</b>	Database Dev	310,000	Orlando

**Fragmentation schema**: Three horizontal fragments JP1, JP2, JP3 based on locations respectively Montreal, New York, Orlando.

**Applications**: **APP1** = SELECT \* FROM J;

APP2 = SELECT \* FROM J WHERE BUDGET < 200,000;

Determine the set of simple predicates Pr to obtain J from its fragments.

If APP1 and APP2 are issued, do you think Pr will satisfy completeness property? If not, then make necessary changes to make the set complete. Justify your answers.

## Example of Completeness

Pr= { LOC="Montreal", LOC="New York", LOC="Orlando"

$$JP_1 = SL_{LOC = MONTREAL}$$
 J

JP1 JNO JNAME

JNO JNAME BUDGET LOC

J1 Instrumental 150,000 Montreal

 $JP_2 = SL_{LOC = NEW YORK}$  J

JP2

		2112 255	
JNO	<u> JNAME</u>	BUDGET	LOC
J2	GUI	135,000	New York
J3	CAD/CAM	250,000	New York

 $JP_3 = SL_{LOC = ORLANDO}$  J

JP3

JNO	JNAME	BUDGET	LOC
J4	Database Dev.	310,000	Orlando

Case 1: An application (app1) is issued in three sites. It wants to access the tuples according to the location (any location).

 $\mathbf{APP1} = SELECT * FROM \mathbf{J};$ 

In this case,  $P_r$  is complete because each tuple of each fragment has the same probability of being accessed.



JP1 (app1)	JNO	JNAME_	BUDGET	LOC
	<del>&gt;</del> J1	Instrumental	150,000	Montreal

JP2	JNO	JNAMF	BUDGET	LOC
	<u>→</u> J2	GUI	135,000	New York
(app1) ———————————————————————————————————	<b>&gt; J3</b>	CAD/CAM	250,000	New York

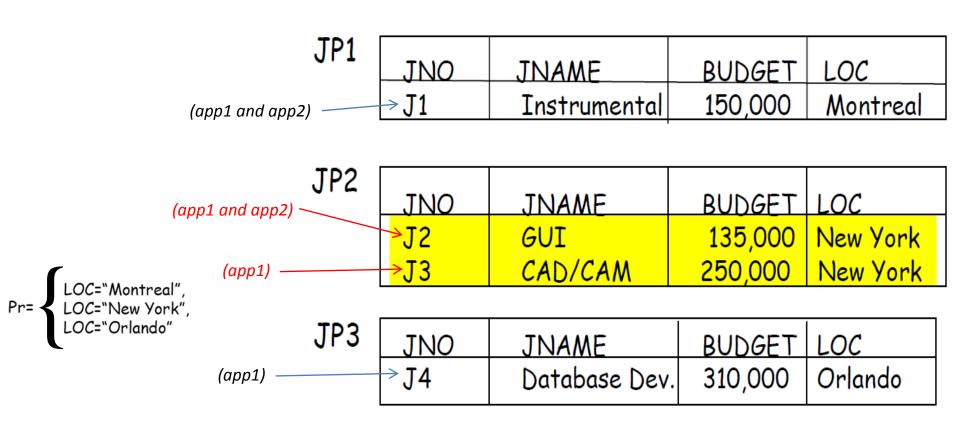
Is Loc="Orlando" relevant? Yes
Is Loc="New York" relevant? Yes
Is Loc="Montreal" relevant? Yes

JP3	JNO	JNAME	BUDGET	LOC
(app1)	<i>→</i> J4	Database Dev.	310,000	Orlando

Case 2: There is a second application (app2) which is also issued in three sites. It accesses only those tuples where budget is less than \$200,000.

APP2 = SELECT \* FROM J WHERE BUDGET < 200,000;

Tuple  $J_2$  has higher access probability than tuple  $J_3$  in  $JP_2$ . In this case,  $P_r$  is not complete since some tuples  $(J_i)$  in  $JP_i$  has higher access probability.



# To make the set complete, we need to add – $(BUDGET \le 200,000, BUDGET \ge 200,000)$ to $P_r$ .

$$Pr = \begin{cases} LOC="Montreal", \\ LOC="New York", \\ LOC="Orlando" \end{cases} \longrightarrow Pr = \begin{cases} LOC="Montreal", \\ LOC="New York", \\ LOC="Orlando" \\ BUDGET <= 200,000, \\ BUDGET > 200,000 \end{cases}$$

#### SL LOC = 'Montreal' J

JP1

JNO JNAME BUDGET LOC
J1 Instrumental 150,000 Montreal

LOC="Montreal", LOC="New York", LOC="Orlando" BUDGET<= 200,000, BUDGET>200,000 JP21 JNO JNAME BUDGET LOC
J2 GUI 135,000 New York

JP22 JNO JNAME BUDGET LOC
J3 CAD/CAM 250,000 New York

JP3 JNO JNAME BUDGET LOC

J4 Database Dev. 310,000 Orlando

TD1				
JFI	JNO	JNAME	BUDGET	LOC
$(app_1 and app_2)$ >	J1	Instrumental	150,000	Montreal

 $(app_1 and app_2)$  ---- JP21 JNO JNAME BUDGET LOC J2 GUI 135,000 New York

Is Loc="Orlando" relevant? Yes
Is Loc="New York" relevant? Yes
Is Loc="Montreal" relevant? Yes
Is Budget <= 200,000 relevant? Yes
Is Budget > 200,000 relevant? Yes

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$(app_1)$ $JPZZ$	JNO	JNAME	BUDGET	LOC
>	J3	CAD/CAM	250,000	New York

Considering the previous app1 and app2, and the set  $P_r$ .

- $\triangleright$  If we add the predicate JNAME = "Instrument" to  $P_r$ .
- $\triangleright$  Is resulting  $P_r$  complete?

```
Pr= LOC="Montreal",
LOC="New York",
LOC="Orlando"
BUDGET<= 200,000,
BUDGET>200,000

JNAME = "Instrument" ??
```

Is Loc="Orlando" relevant? Yes
Is Loc="New York" relevant? Yes
Is Loc="Montreal" relevant? Yes
Is Budget <= 200,000 relevant? Yes
Is Budget > 200,000 relevant? Yes
Is Jname = "Instrumental" relevant? No

# Completeness

A set of simple predicate  $P_r$  is said to be **complete** if and only if –

There is an equal probability of access by every application (i.e. query) to any tuple belonging to any minterm fragment that is defined by  $P_r$ 

# Minimality

The set of predicates  $P_r$  is minimal if and only if there is at least one application (i.e. query) that accesses the fragment.

If all the predicates of a set  $P_r$  is relevant, then  $P_r$  is Minimal.

```
Pr= LOC="Montreal",

LOC="New York",

LOC="Orlando"

BUDGET<= 200,000

BUDGET>200,000
```

```
LOC="Montreal",
LOC="New York",
LOC="Orlando"
BUDGET<= 200,000,
BUDGET>200,000

JNAME = "Instrument" | ??
```

## Let's See

Say, only APP3 application accesses J wants to access project located in New York.

APP3 = SELECT \* FROM J WHERE LOC = "NEW YORK";

Pr= { LOC="Montreal", LOC="New York", LOC="Orlando" JP1

JNO	JNAME	BUDGET	LOC
J1	Instrumental	150,000	Montreal

Is Loc="Orlando" relevant? No Is Loc="New York" relevant? Yes Is Loc="Montreal" relevant? No

Is the set Complete? Minimal?

JP2

JNO	JNAME	BUDGET	LOC
J2	GUI	135,000	New York
<b>J</b> 3	CAD/CAM	250,000	New York

JP3

JNO	JNAME	BUDGET	LOC
J4	Database Dev.	310,000	Orlando

#### Exercise

Consider the following global relation, fragmentation schema and applications.

#### STUDENT

SNUM	SNAME	DEPT	SEM
1	A	CSE	1
2	В	EEE	2
3	С	CSE	3
4	D	EEE	4
5	Е	CSE	6
6	F	CSE	7

#### Fragmentation Schema:

```
STUDENT_1 = SL_{DEPT = "CSE"} STUDENT

STUDENT_2 = SL_{DEPT = "EEE"} STUDENT
```

#### Applications:

```
APP<sub>1</sub> = SELECT * FROM STUDENT WHERE DEPT = "CSE" OR DEPT = "EEE";
APP<sub>2</sub> = SELECT * FROM STUDENT WHERE SEM = 6;
```

#### Now, answer the following questions:

- Define simple predicate and minterm predicate.
- ii. Determine the set of simple predicates  $P_r$  to obtain **STUDENT** from its fragments. If **APP**<sub>1</sub> and **APP**<sub>2</sub> are issued, do you think  $P_r$  will be completed? If not, then make necessary changes to make the set complete. Justify your answer.

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