

Basic Elements of Formal Languages

- What is Formal Language?
- What is Natural Language?
- How Formal Language is different from Natural Language?

Answers:

- A Natural Language or Ordinary Language is any language that has evolved naturally through use and repetition without conscious planning or premeditation.
- We discover the grammar of a natural language through empirical investigation.
- We don't discover the grammar of a formal language (artificial language), we stipulate it- we define it however we want.

BASIC ELEMENTS:

- **Symbol:** A *symbol* is an abstract entity that we shall not define formally, just as *point* and *line* are not defined in geometry.
- **Alphabet:** An *alphabet* is a finite, nonempty set of symbols. Conventionally, we use the symbol Σ (capital sigma) for an alphabet. Common alphabets are:
 1. $\Sigma = \{0, 1\}$, the binary alphabet.
 2. $\Sigma = \{a, b, \dots, z\}$, the set of all lower-case letters.
 3. The set of all ASCII characters, or the set of all printable ASCII characters.
- **String/Word:** A *string* (or sometimes *word*) is a finite sequence of symbols chosen from some alphabet. Example:
 - 01101 is a string from the binary alphabet $\Sigma = \{0, 1\}$.
 - 111 is another string from the same alphabet.

Empty String:

- The *empty string* is the string with zero occurrences of symbols.
- This string, denoted by ϵ (epsilon), is a string that may be chosen from any alphabet whatsoever.

Length of a String:

- The *length of a string* is the number of positions for symbols in the string. For example, 01101 has length 5.
- The number of symbols as the length: accepted but not strictly correct. The string 01101 has only two symbols, 0 and 1.

- The standard notation for the length of a string w is $|w|$.
- Example: $|011| = 3$ and $|\epsilon| = 0$.

Prefix of a String:

- A *prefix* of a string is any number of leading symbols of that string.
- Example: String abc has prefixes ϵ , a , ab , and abc .
- A prefix of a string, other than the string itself, is called a *proper prefix*.

Suffix of a String:

- A *suffix* of a string is any number of trailing symbols of that string.
- Example: String abc has suffixes ϵ , c , bc , and abc .
- A suffix of a string, other than the string itself, is called a *proper suffix*.

Concatenation of Strings:

- The *concatenation* of two strings is the string formed by writing the first, followed by the second, with no intervening space.
- Let x and y be strings of length i and j respectively. Then xy denotes the concatenation of x and y and the *length* of xy is $i+j$.
- Example: Let $x = 1101$ and $y = 0011$. Then $xy = 11010011$.
- The empty string is the *identity* for the concatenation operator. That is, $\epsilon w = w\epsilon = w$.

Powers of an Alphabet: Enigma (7c)

- If Σ is an alphabet, the set of all strings of a certain length from that alphabet can be expressed by using an exponential notation.
- Σ^k is defined as the set of strings of length k , each of whose symbols is in Σ .
- $\Sigma^0 = \epsilon$, no matter what the alphabet Σ is. In other words, ϵ is the only string of length 0.
- Example: If $\Sigma = \{a, b, c\}$ then $\Sigma^1 = \{a, b, c\}$, $\Sigma^2 = \{aa, ab, ac, ba, bb, bc, ca, cb, cc\}$, $\Sigma^3 = \{aaa, aab, aac, aba, abb, abc, aca, acb, acc, baa, bab, bac, bba, bbb, bbc, bca, bcb, bcc, caa, cab, cac, cba, cbb, cbc, cca, ccb, ccc\}$.
- The set of all possible strings of all possible lengths over an alphabet Σ is conventionally denoted by Σ^* (*Kleene Star*). For instance, $\{0, 1\}^* = \{\epsilon, 0, 1, 00, 01, 10, 11, 000, \dots\}$.
- The set of nonempty strings from alphabet Σ is denoted by Σ^+ (*Kleene Closure/Plus*).

** Confusion between Σ and Σ^1 ?

- We shall use the same notation for the two sets. Instead, it may be inferred from the context whether we are speaking about an alphabet or a set of strings.

- **Languages:**

- A set of strings all of which are chosen from some Σ^* , where Σ is a particular alphabet, is called a *(formal) language*.
- If Σ is an alphabet, and $L \subseteq \Sigma^*$, then L is a language over Σ . A language over Σ need not include strings with all the symbols of Σ .
- If L is a language over Σ , it is also a language over any alphabet that is a superset of Σ .
- Complement of a formal language, $\Sigma^* - L$.
- \emptyset , the empty language, is a language over any alphabet.
- $\{\epsilon\}$, the language consisting of only the empty string, is also a language over any alphabet. Notice that, $\emptyset \neq \{\epsilon\}$; the former has no string but the latter has one string.

Example:

- The language of all strings consisting of n 0's followed by n 1's, for some $n \geq 0$: $\{\epsilon, 01, 0011, 000111, \dots\}$;
- The set of strings of 0's and 1's with an equal number of each: $\{\epsilon, 01, 10, 0011, 0101, \dots\}$.
- Set-Formers to define Language: $\{w \mid w \text{ consists of an equal number of 0's and 1's}\}$.