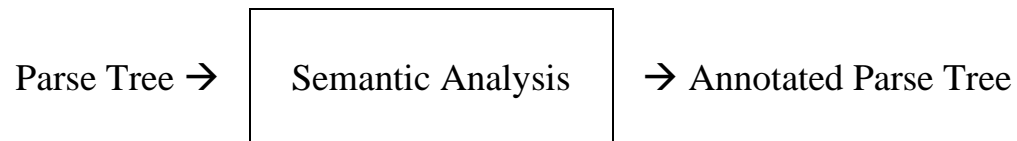


A Brief Overview of Semantic Analysis

- Semantic analysis helps evaluate language constructs, so that the source code is progressed toward executable action sequences or target code.
- Firstly, semantics help correct ‘simple errors’, that is, interpret tokens, their types, and their relations with each other.
- Secondly, semantics associate information with recognized constructs for finding the meaning of them.



➤ Simple semantic errors

- Type mismatch
- Undeclared variable
- Reserved word misuse
- Multiple declaration of variable in a scope
- Accessing an out-of-scope variable
- Actual and formal parameter mismatch
- and so on.

➤ Correction of simple semantic errors

- Most of the above semantic errors are detected and corrected during the earlier phases of compilation, that is, lexical analysis and syntactic analysis.
- Construction and use of symbol tables during lexical analysis and syntactic analysis are the right points where correction of simple semantic errors may be arranged with a very little effort.

➤ Syntax Directed Translation

- Parse trees generated by CFGs during syntax analysis lack information of how to evaluate a construct that has been recognized.
- To have interpreted a parse tree, that is, to find meaning of the construct, semantic rules need to be associated with the productions of the CFGs that generate it.
- CFG + semantic rules = Syntax Directed Definitions (SDD).
- Syntax Directed Definitions during parsing thus help do useful things toward code generation. The process is known as *Syntax Directed Translation (SDT)*.

- Two ways to do Syntax Directed Translation:

- Top Down
- Bottom UP

➤ Attribute Grammar

- A set of attributes is associated with each grammar symbol.
- Some of the attributes are called *synthesized*, and some are called *inherited*.
- Actions written corresponding to a production manipulate these attributes effectively to do the desired translation.
- The parse tree with attributes is called an *annotated parse tree*.

➤ Synthesized attributes

These attributes get values from the attribute values of their child nodes. In the production

$$S \rightarrow ABC$$

if S is taking values from its child nodes A, B, C, then it is said to be a synthesized attribute.

➤ Inherited attributes

In contrast to synthesized attributes, inherited attributes can take values from parent and/or siblings. In the production,

$$S \rightarrow ABC$$

A can get values from S, B and C. B can take values from S, A, and C. Likewise, C can take values from S, A, and B.

Example of annotating a parse tree

Evaluating arithmetic expression using SDT:

Consider the following grammar:

$$E \rightarrow E+T \mid E-T \mid T$$

$$T \rightarrow \text{num}$$

Let us assume an attribute, val, with E, T and F. Following are the grammar rules with semantic actions.

Production	Semantic rule
$E \rightarrow E_1 + T$	$\{E.\text{val} = E_1.\text{val} \parallel '+' \parallel T.\text{val} \}$
$E \rightarrow E_1 - T$	$\{E.\text{val} = E_1.\text{val} \parallel '-' \parallel T.\text{val} \}$
$E \rightarrow T$	$\{E.\text{val} = T.\text{val} \}$
$T \rightarrow \text{num}$	$\{T.\text{val} = \text{num} \}$

The operator ' \parallel ' represents concatenation, *num* represents any number constant and E_i , T_i are used to distinguish two Es and Ts.

Consider the problem of generating postfix equivalents of infix expressions with the following grammar.

$E \rightarrow E+T \mid E-T \mid T$

$T \rightarrow \text{num}$

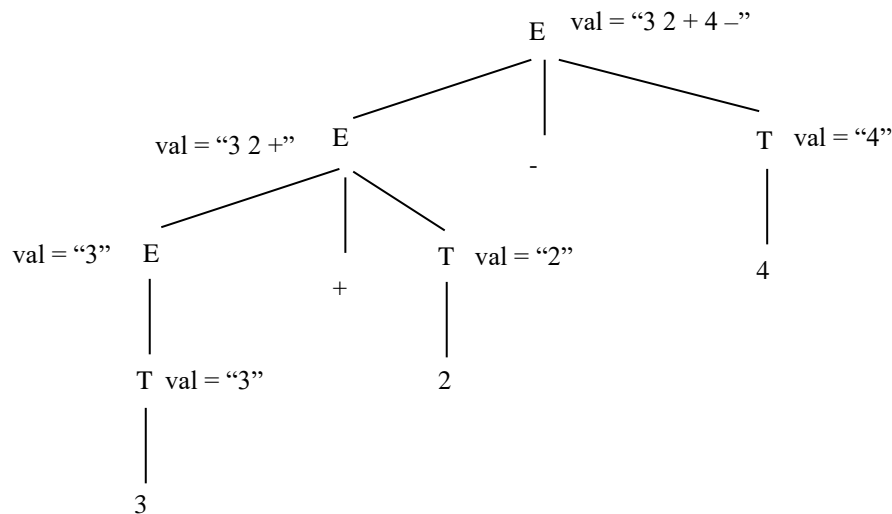
Let us assume an attribute, val, with E and T that holds the string corresponding to the postfix expression.

The semantic actions can be written as,

Production	Semantic rule
$E \rightarrow E_1+T$	$\{E.\text{val} = E_1.\text{val} \parallel T.\text{val} \parallel '+'\}$
$E \rightarrow E_1-T$	$\{E.\text{val} = E_1.\text{val} \parallel T.\text{val} \parallel '-'\}$
$E \rightarrow T$	$\{E.\text{val} = T.\text{val} \}$
$T \rightarrow \text{num}$	$\{T.\text{val} = \text{num} \}$

The operator ‘||’ represents concatenation, *num* represents any number constant and E_1 is used to distinguish two Es.

The **annotated parse tree** for the expression “3+2-4” is given below, which generates the postfix string “3 2 + 4 -”.



The parse trees produced by this stage will be used in later stages for code generation and other related activities.