

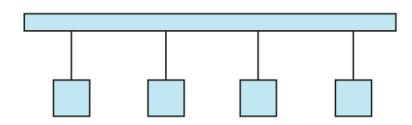


CSE3103 : Database FALL 2020

Nazmus Sakib
Assistant Professor
Department of Computer Science and Engineering
Ahsanullah University of Science and Technology

Interconnection Network Architectures

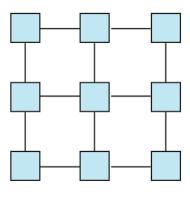
- **Bus**. System components send data on and receive data from a single communication bus;
 - Does not scale well with increasing parallelism.



(a) bus

Interconnection Network Architectures

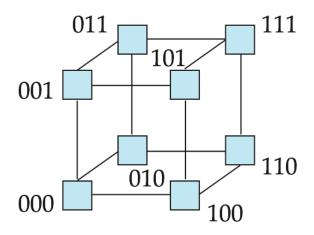
- **Mesh**. Components are arranged as nodes in a grid, and each component is connected to all adjacent components
 - Communication links grow with growing number of components, and so scales better.
 - But may require $2\sqrt{n}$ hops to send message to a node (or \sqrt{n} with wraparound connections at edge of grid).



(b) mesh

Interconnection Network Architectures

- **Hypercube**. Components are numbered in binary; components are connected to one another if their binary representations differ in exactly one bit.
 - *n* components are connected to *log(n)* other components and can reach each other via at most *log(n)* links; reduces communication delays.



(c) hypercube

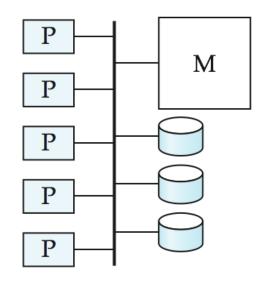
Parallel Database Architectures

- Shared memory -- processors share a common memory
- Shared disk -- processors share a common disk
- **Shared nothing** -- processors share neither a common memory nor common disk
- **Hierarchical** -- hybrid of the above architectures

Shared Memory

- Processors and disks have access to a common memory, typically via a bus or through an interconnection network.
- Extremely efficient communication between processors

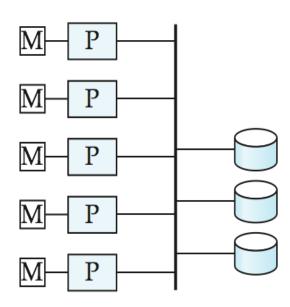
 data in shared memory can be accessed by any processor without having to move it using software.
- Downside architecture is not scalable beyond 32 or 64 processors since the bus or the interconnection network becomes a bottleneck
- Widely used for lower degrees of parallelism (4 to 8).



(a) shared memory

Shared Disk

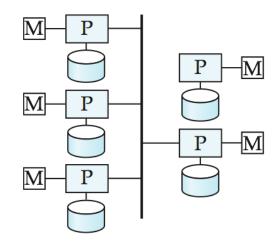
- All processors can directly access all disks via an interconnection network, but the processors have private memories.
 - The memory bus is not a bottleneck
 - Architecture provides a degree of fault-tolerance if a processor fails, the other processors can take over its tasks since the database is resident on disks that are accessible from all processors.
- Examples: IBM Sysplex and DEC clusters (now part of Compaq) running Rdb (now Oracle Rdb) were early commercial users
- Downside: bottleneck now occurs at interconnection to the disk subsystem.
- Shared-disk systems can scale to a somewhat larger number of processors, but communication between processors is slower.



(b) shared disk

Shared Nothing

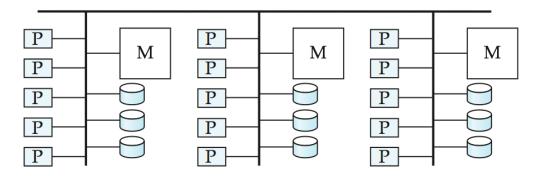
- Node consists of a processor, memory, and one or more disks. Processors at one node communicate with another processor at another node using an interconnection network. A node functions as the server for the data on the disk or disks the node owns.
- Examples: Teradata, Tandem, Oracle-n CUBE
- Data accessed from local disks (and local memory accesses) do not pass through interconnection network, thereby minimizing the interference of resource sharing.
- Shared-nothing multiprocessors can be scaled up to thousands of processors without interference.
- Main drawback: cost of communication and non-local disk access; sending data involves software interaction at both ends.



(c) shared nothing

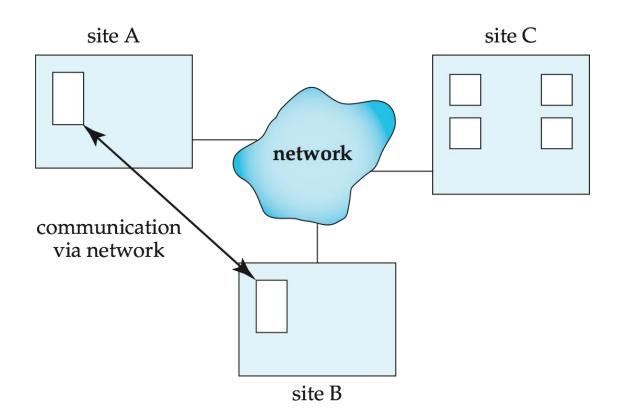
Hierarchical

- Combines characteristics of shared-memory, shared-disk, and shared-nothing architectures.
- Top level is a shared-nothing architecture nodes connected by an interconnection network, and do not share disks or memory with each other.
- Each node of the system could be a shared-memory system with a few processors.
- Alternatively, each node could be a shared-disk system, and each of the systems sharing a set of disks could be a shared-memory system.
- Reduce the complexity of programming such systems by **distributed virtual-memory** architectures
 - Also called non-uniform memory architecture (NUMA)



Distributed Systems

- Data spread over multiple machines (also referred to as **sites** or **nodes**).
- Network interconnects the machines
- Data shared by users on multiple machines



Distributed Databases

- Homogeneous distributed databases
 - Same software/schema on all sites, data may be partitioned among sites
 - Goal: provide a view of a single database, hiding details of distribution
- Heterogeneous distributed databases
 - Different software/schema on different sites
 - Goal: integrate existing databases to provide useful functionality
- Differentiate between *local* and *global* transactions
 - A **local transaction** accesses data in the *single* site at which the transaction was initiated.
 - A **global transaction** either accesses data in a site different from the one at which the transaction was initiated or accesses data in several different sites.

