

ASSIGNMENT

Course Code: CIS 122

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Submitted To:

Name: Md. Faruk hosen

Designation:

Department: CIS

Daffodil International University

Submitted By:

Name: MD. Samiul islam

ID: 242-16-009

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Daffodil International University

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Introduction

lata structures are important for managing and processing data in computer science. This project looks at how data structures like queues, graphs, binary search trees, and linked lists are used in real situations, showing how they help solve problems and improve efficiency. for example, an online ticket booking system Uses a queue to hundle nequests one at a time, making sume evenyone is freated fairly. Graph theory helps find the best travel routes between historical sites by analysing various factors. The binary search free (BST) makes it easy to seanth for student ID. while a linked list shows how to insent, delete, and merense items.

By using and studying these data structures, this project helps us understand how to manage data well, preparing us for more complexe problem - solving in computer science.

Task-1

Ans. to the question no -1

In the scenario, Mr. Sisir is developing an online ticket booking system for buses and trains, while Mr. Tusher is creating a music play list application. Each of these projects Utilizes different data structures that are well - suited to their specific requirements. Here are mainly used two data structure.

1 Queue

1 Singly linked list / linked list

Mr. Sisiris online Ticket Booking system;

For the online ticket booking system, Mr Sisir would likely use a queue data Structure. Queues are particularly effective in scenarios where requests need to be processed in a first-come, first - served manner. Basically it is a linear data structure that follows the FIFO (First In First Out) principle. In a ticket booking system, users typically wait in line to book their tickets, making queues an ideal choice.

operations: The queue allows for efficient operations such as enqueue (adding a new booking request) and dequeue (processing the next booking request). This ensures that the system can handle multiple users simultaneously while maintaining the order of requests.

Mr. Tusheris Music play hist Application.

In contrast, Mr. Tusher's music playlist application would benefit from using a linked list. Linked lists are advantageous for applications that require frequent insentions and deletions, such as adding or removing songs from a playlist.

promine size: Unlike armays, linked lists can grow and shrink dynamically, which is essential for a playlist where users may frequently modify their selections. Each song can be represented as a node in the linked list, allowing for easy traversal, insention, and deletion of songs without the need for contiguous memory allocation.

In summary, Mr. Sisir's Project utilizes a queue to manage ticket bookings efficiently, while Mr. Tusher's application employs a linked hist to facilitate dynamic playlist management. Understanding these data structures enables both developers to optimize their applications for penformance and user experience.

Ans. to the question no-2

To prove that the adjacency matrix of the directed graph in Figure of is asymmetric, we need to understand the properties of directed graphs and their adjancency matrices.

Direct Graph: In a directed graph, edges have a direction, meaning that an edge from ventex A to ventex B does not imply an edge from B to A.

Adjancency Matrix! The adjacency matrix for a directed graph is a square matrix used to represent the graph. The element at now i and column is indicates the presence (and possibly the weight) of an edge from ventex i to ventex it.

An advancency matrix A of a directed weighted graph is asymmetric if:

 $A[i][j] \neq A[j][i]$ for at least one pair (i, j)

Example from the Graph:

- · Consider the vertices coxis Bazar (6) and Rangamati (5);
- (i) The edge from Coxis Baran (6) to Rangamati (5) has a weight of -11
- (ii) The edge from Rangamati (5) to cox's Bazar (6) does not exist (or has a weight of o if not explicitly stated)

This means:

A[6][5] = -11 and A[5][6] = 15 $A[6][5] \neq A[5][6]$

In summary, The adjancency matrix of the directed graph shown in Figure as is asymmetric because the travel cost from one location to another is not necessarily equal to the cost in the reverse direction. This is a fundamental characteristic of directed weighted graphs.

Ans. to the question no- 3

Mor. Amif Successfully solved the Challenge by following a structured approach, combining BST construction with binary search principles to derive his student ID. Hereis a detailed breakdown of his solution:

We need to distinct so integers between so and 200.

Fixed position:

1st number = 26

4th number = XX

7th number = 86

9th number = 52

xx calculation:

Let's assume Mr. Aris's student ID ends with 34.

XX = (Last two digits of student ID % 10) + lowest number selected.

= (34 % 10) + lowest number

= 4+26

= 30

Here,

only two numbers (35: 26 and XX(30)

No more than four numbers > 110

To satisfy all constraints, Mr. Anif selected the following to numbers:

26, 28, 40, 30, 50, 90, 86, 100, 52, 120

50, two numbers (35: 26 and 30.

Four numbers > 100 : 120 (only one in this example; adjust if needed).

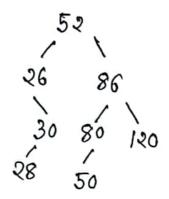
construct the Binary search Tree (BST)

For any node, BST order follows, '
lef subtree values < node value < right subtree
values

Insention process:

- 1. Start with first number: 20 (which is root)
- 2. Insent 28 (right of 26)
- 3. Insent 40 (night of 28).
- 4. Insert 30 (left of 40).
- 5. Continue untill all numbers are insented.

BST Structure:



Link BSTto to student ID via Binary search

Travse the BST to locate XX = 30

Start af most 52 -> move left 26 -> right 30

Given XX = (last two digits of ID% 10) t

lowest number (26)

 $\Rightarrow 30 = (JD\% 10) + 26 \Rightarrow (JD\% 10) = 4 \Rightarrow ID$ end with "4"

If xx were 20 -> ID ends with "3".

So, Mr. Arif by this way solved his problem. And this success came from carefully analysing each requirement and verifying the solution at each step.

Ans. to the question no-4

Mr. Habib faced difficulties in finding his student ID using the Binary Tree (BST) he constructed. Mr. Anif, having successfully navigated the Challenge, provided assistance by explaining the fundamental principles of binary search and how they apply to BSTs.

Understanding Binary Search in BSTs:

A Binary Search Thee is a data structure where each node has at most two chidnen. The left child contains values less than the parent node, while the right child contains values greater than parent node. This property allows for efficient Seanching.

Binary Search process:

- (i) starting at the root node.
- (ii) Companing the tanget value with the current noders value.
- (iii) If the tanget value is less than the current moders value, the search continues in the left subtree.
- (iv) If the target value is greater, the search continues in the right subtree.

(v) This process repeats until the value is found on a leaf node is reached.

Steps Taken by Mr. Arif:

- clarifying the search steps: Mr. Anif likely walked Mr. Habib through the search process step-by-step, ensuring he understood how to navigate the tree based on comparisons.
- · Visualizing the BST: By drawing the BST and labeling the modes, Mr. Anif helped Mr. Habib visualize the structure making it easier to follow the search path.
- Debugging the Search Logic: If Mr. Habib was implementing the Search algorithm program-madically, Mr. Anif might have reviewed his code to identify any logical errors or incorrect companisons that could lead to unsuccess ful searches.

In summary, Mr. Arif explained things clearly and showed Mr. Habib how to use the features of the binary search tree (BST) and the binary Search method. This helped Mr. Habib find the his student ID in the tree. Working together not only solved Mr. Habibis problem but also helped both of them understand data structures and algorithm better.

Task-B - Ans. to the question no-5

To find the minimum travel cost from X, where X = (last two alignsts of your student ID % 6) th, we first need to calculate X using the provided student ID of 09.

Cakalation of x:

Last two digits of student ID=09 Now, (09%6)+1

= 341

Destination.

The destination cornesponding to X = 4 is Saint Martin.

from the & graph, we can see the following direct paths from DIU (ventex 7) to saint Mandin (ventex 4);

- (i) DIU to Saint Mantin: cost = 8
- (ii) DIU to Rangam ati (5): cost = 50, then
 Rangamati to Saint Mantin Cost = 45 (Total =
 50 + 25 = 75)
- (iii) DJU to Sajek valley (2): cost 294, then Sajeck valley to Saint Mantin: cost = 15 (Total = 94 + 13 = 109)

DIU to Ahsan Munzill (3): cost = 13, then Ahsan Manzill to Saint Muntin: cust = 8 (Potal = 13+8 -21)

Here, Directly to saint Mantin cast = 8
Via Rangamati's 75
Via Sajek valley: 109
Via Ahsan Manzil: 21

The minimum travel cost from DIU to Saint Mantin is 8.

Ans. to the question no - 6

The Chromatic number of a graph is defined as the smallest number of colors needed to color the vertices such that no two adjacent vertices share the same color. Heners a detailed explanation of how to defermine the Chromatic number for the graph shown in figure of.

Here, the graph consists of vertices representing historical sites: Coxis Bazar (1), Rangamati (2), Saint Mantin (3), Kuthibani (4), Satek Valley (5), Ahsan Manzill (6), and 0.70 (7).

The edges represent the connections between these ventices, with weights indicating thavel costs.

Graph Representation!

visualize the graph and note the connections. For example, if two ventices are directly connected by an edge they cannot shape the Same color,

Coloring:

Hene, (1) ventex 7 (DIU): Color 1

(ii) ventex 5 (Ranga mati): Color 2 (adian cent to

DIU)

(iii) Ventex 2 (Saiseak Valley): Color 2 (adiacent

to DIU)

(iv) Ventex q (saint Mantin): color 3 (adiacent to Ranga mati)

(v) Ventex 3 (Ahsan Mansill): Colon & Cadianal to OIU)

(vi) vertex 1 (cox's Basar); color 3 (addacent to Sadeek valley)

(vii) venter & (kuthibari): color 1 (addacent to Ahsan Mangill)

After applying the greedy coloning algorithm, if you find that you used three we found three different colons, then the chromatic number of the graph is 3. This means that at least three colons are required to colon the ventices of graph such that no two adjacent ventices share the same colon.

So, Chromatic Number of Figure 01 = 3

Ans. to the question no- \$

My SAPA is -> 4.00

2x => (4.00) + 3.21

-> 9.21

Figure look like

Head = 5x

Pse udo code:

1. Stant at the head of the linked list.

- 2. Traverse the list while keeping track of previous node.
- 3. If the cumment node contain \$ 7.21

 a. Change the pnerious modes next pointen

 to point to the next mode of 7.21.
- b. Delete the node containing 7.21 '
- 9. End travensal.

SB: Yes, it is possible to display the elements of a singly linked list in revens order, but we cannot thaverse backward directly, because a singly linked list only has pointers to next node.

However, we can achieve this in two days:

1) Using Recursion.

(1) Using stack.

a stack recursion allows us to neach the last node first and print values in reverse order.

Void print Reverse (Struct Mode x head) of

if (head = = NULL)

return;

print ("% 2f -> " head -> data");

Using a stack since a stack follows LIFO (Last In Fast Out) we can push elements into a stack and pop them to print in neverse order.

I. Therense the list and push each modes value and into a stack

s. pop elements from the stack to print them in verense onder.

3. Since a stack follows LIFO, the last node gets printed first.

Yes it is possible to print linked list in reverse order using this method.