

Prim's Algorithm

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Course Title: Design and Analysis of Algorithm Lab



Basic Concepts

Spanning Tree:

- A subgraph T of an undirected graph G (V,E) is a spanning tree of G, if it is a tree and contains every vertex of G.
- There is a fixed number of edges in the spanning tree which is equal to one less than the total number of vertices (E = V 1).
- There can be many possible spanning trees for a graph.
- To find the minimum spanning tree two algorithms can be used.
 - 1. Prim's Algorithm.
 - 2. Kruskal's Algorithm.

Introduction

To Prim's Algorithm

- ➤ Prim's algorithm is used to find the Minimum Spanning Tree for a given graph.
- The algorithm finds the subset of edges for graph G (V, E) such that the summation of edge weights is minimum.
- > Prim's algorithm is a Greedy Algorithm.

Prim's Algorithm





Problem:

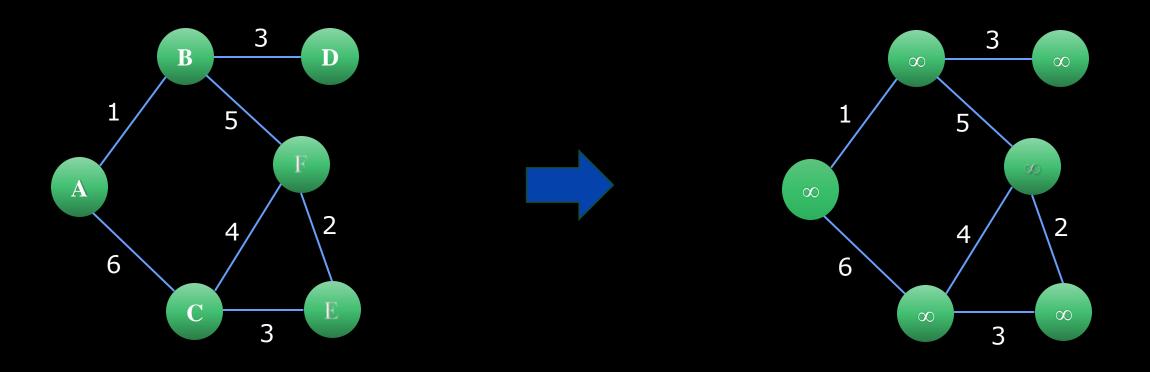
- In computer networks, routing data efficiently between different points is important for fast and smooth communication.
- Prim's algorithm helps by figuring out the best way to connect these points, like finding the shortest path between them.
- It's like planning the fastest root on a map to deliver something from one place to another.

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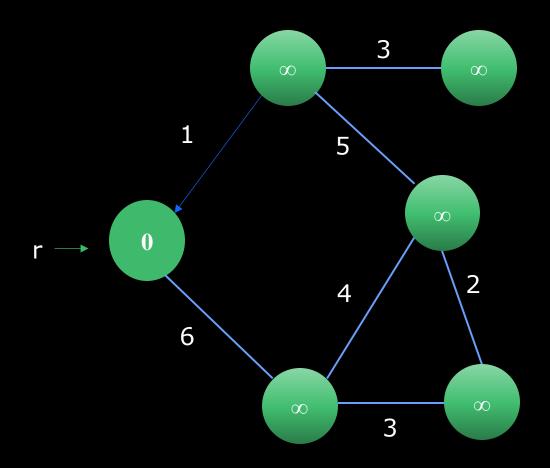
Pseudo Code:

```
MST-Prim (G, w, r) 
Q = V[G];
for each u \in Q
\text{key}[\mathbf{u}] = \infty;
\text{key}[\mathbf{r}] = 0;
p[r] = NULL;
while (Q not empty)
u = ExtractMin(Q);
for each v \in Adj[u]
if (v \in Q \text{ and } w(u,v) < \text{key}[v])
p[v] = u;
\text{key}[v] = w(u,v);
```

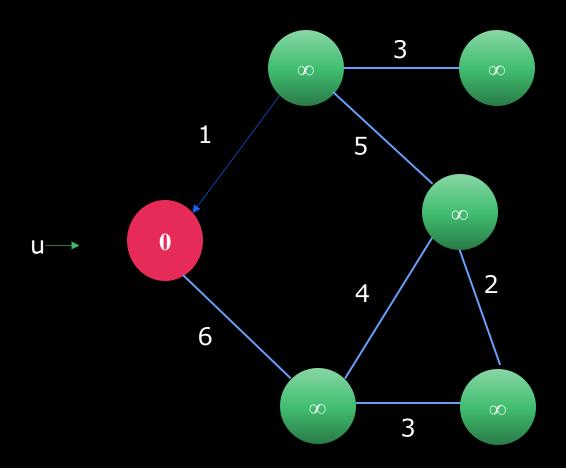
Step 1: Initially set the value of all vertices to infinity.



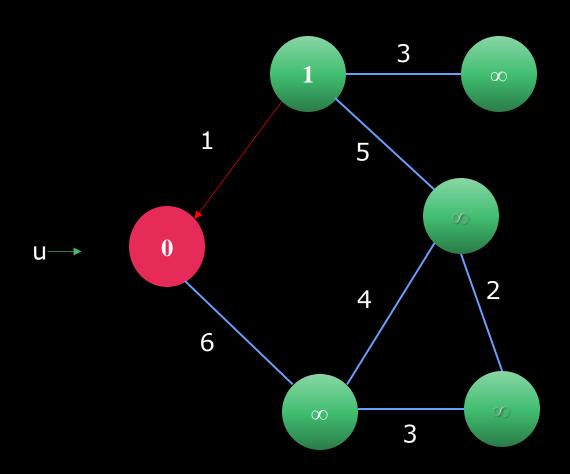
Step 2: Choose an arbitrary starting vertex r.

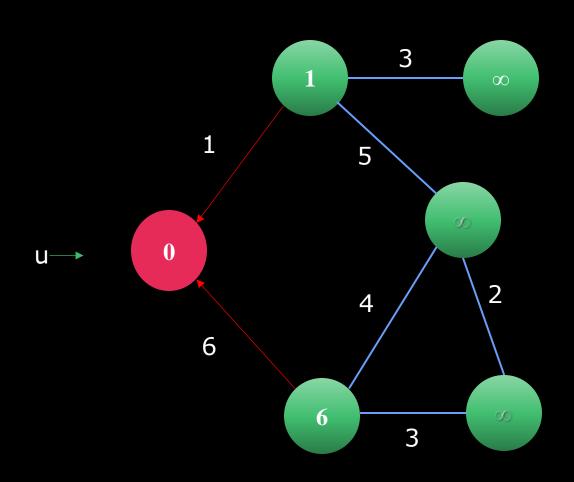


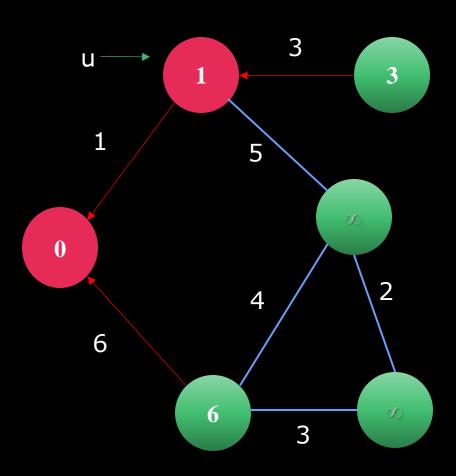
Step 3: Red vertices have been removed from Q.

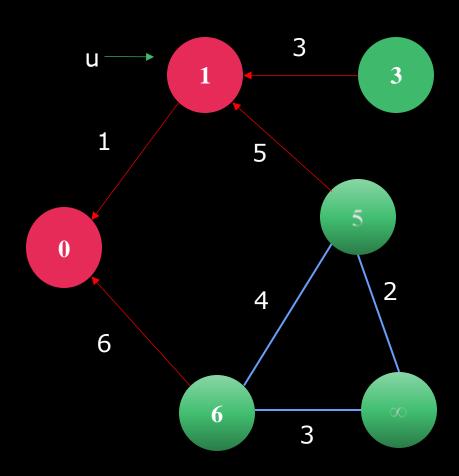


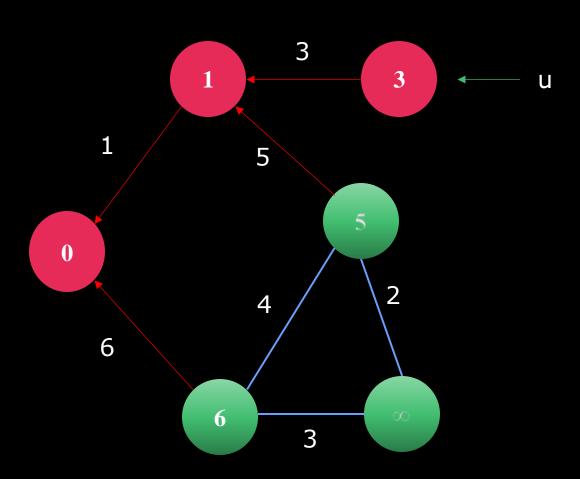
Step 4: Red arrows indicate parent vertices.

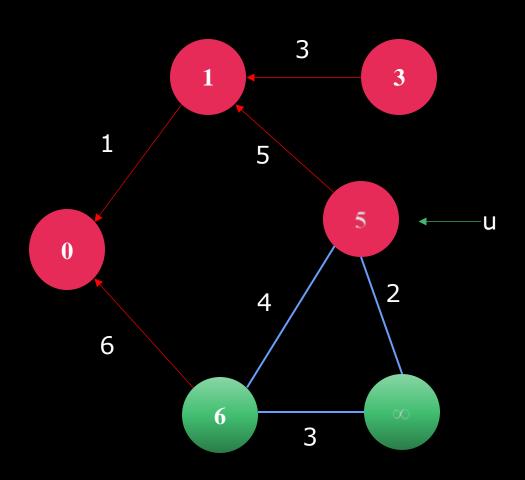


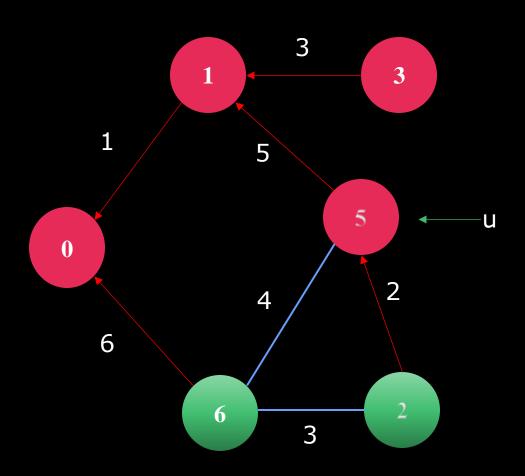


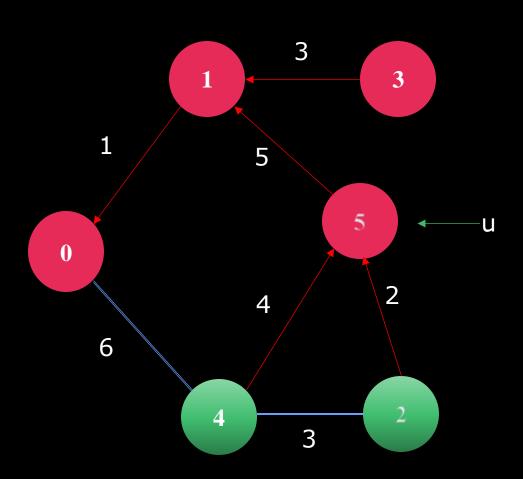


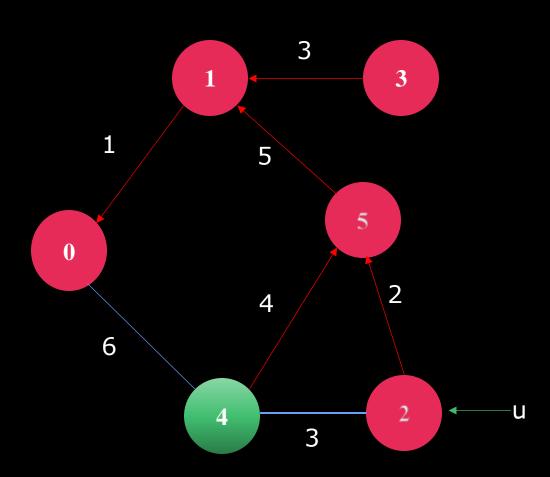


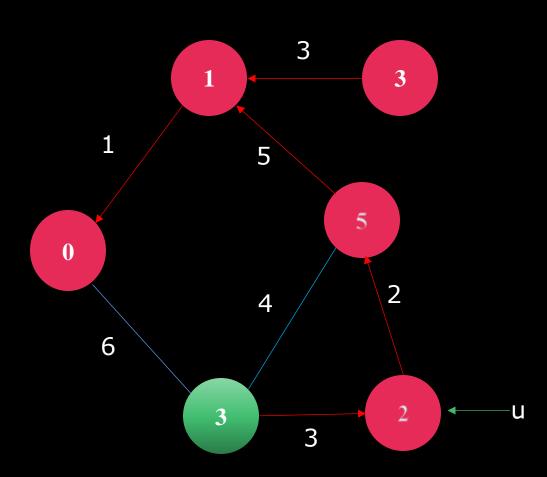


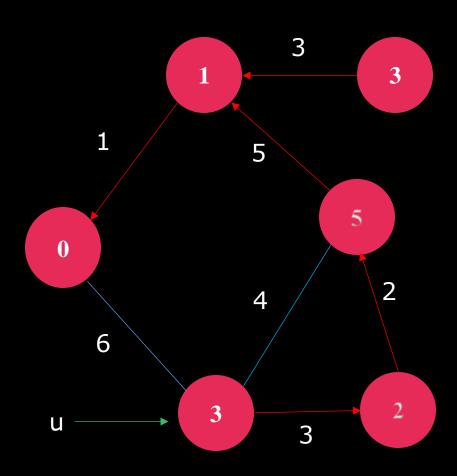


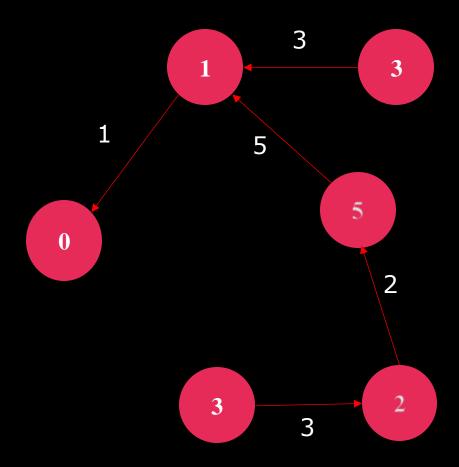






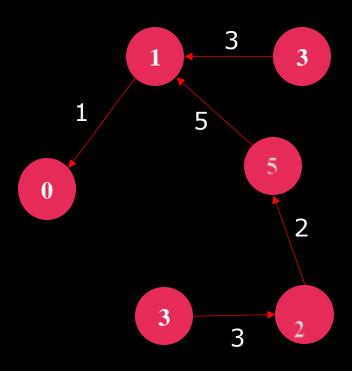






Weight: 1+3+5+2+3=14

Running Time



The Final Minimum Spanning Tree

What will be the running time?

A: Depends on queue

Binary heap: O(E lg V)

Fibonacci heap: O(V lg V + E)

