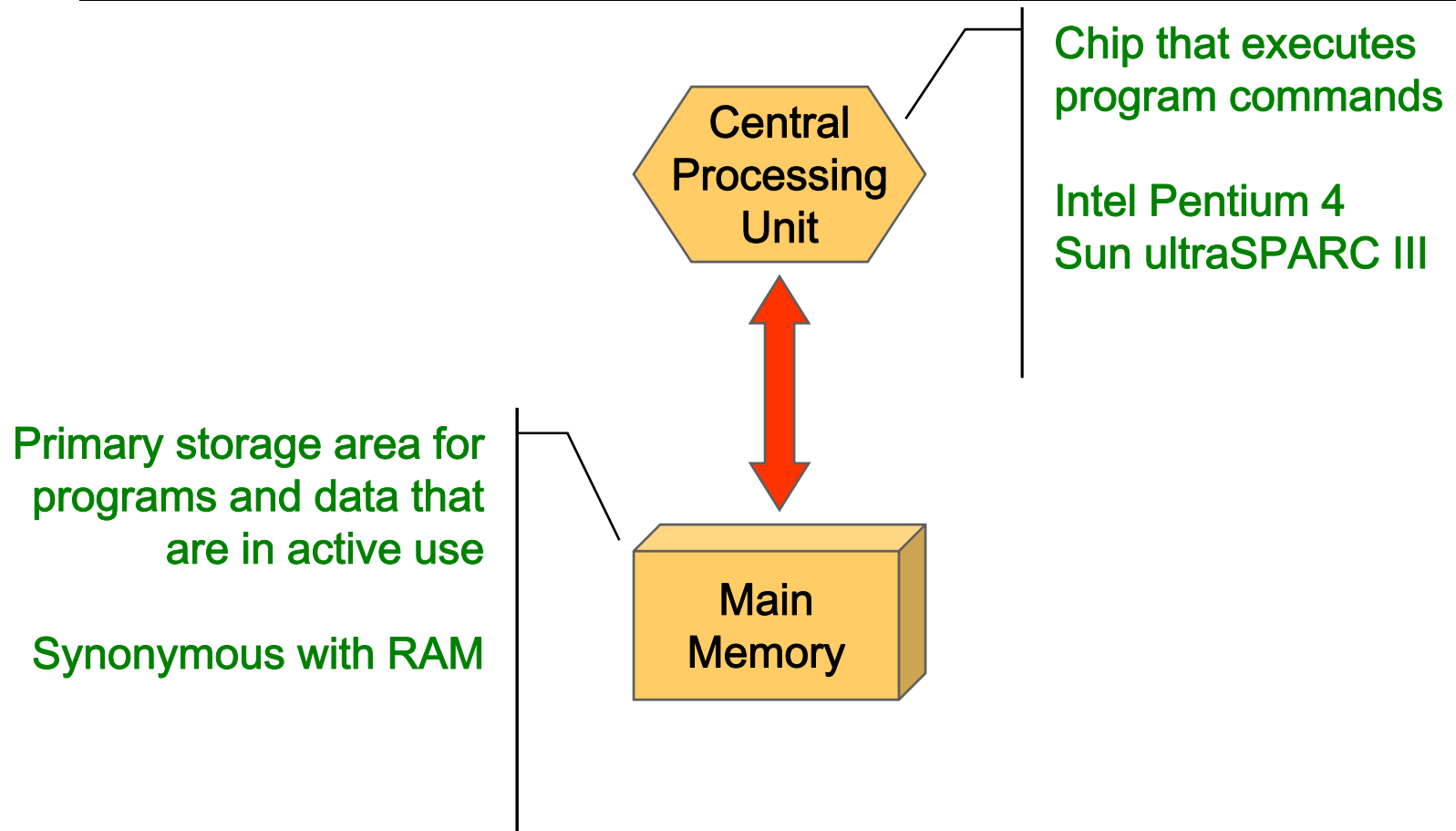




Introduction to Computer System and Architecture

CPU and Main Memory

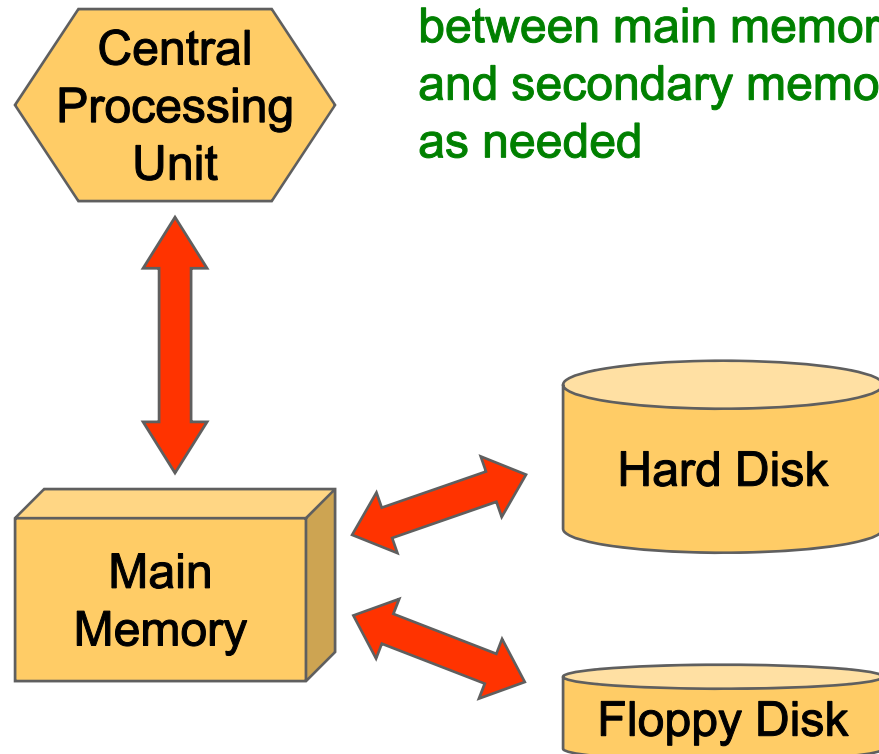


Secondary Memory Devices

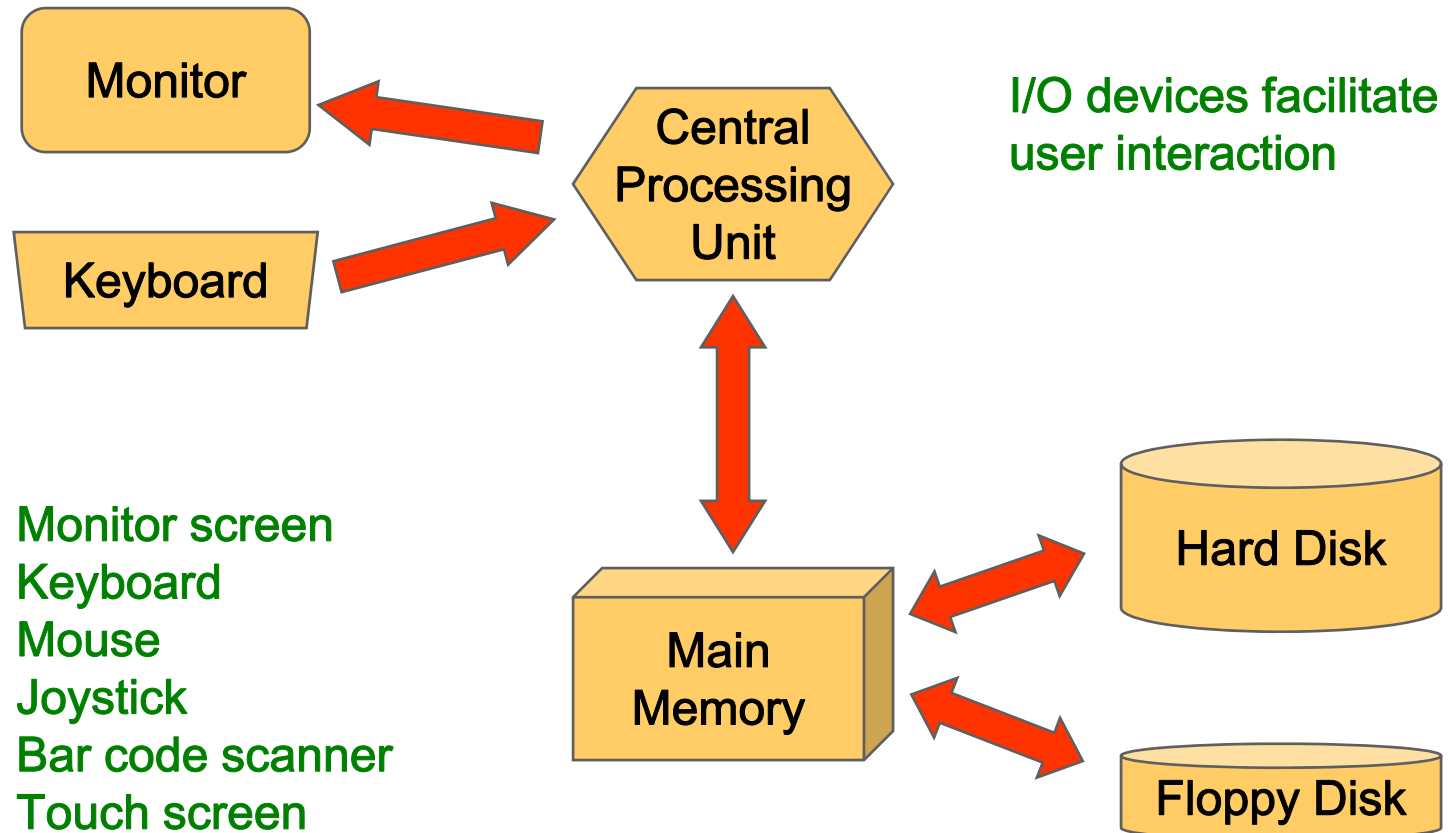
Secondary memory devices provide long-term storage

Information is moved between main memory and secondary memory as needed

Hard disks
Floppy disks
ZIP disks
Writable CDs
Writable DVDs
Tapes



Input / Output Devices



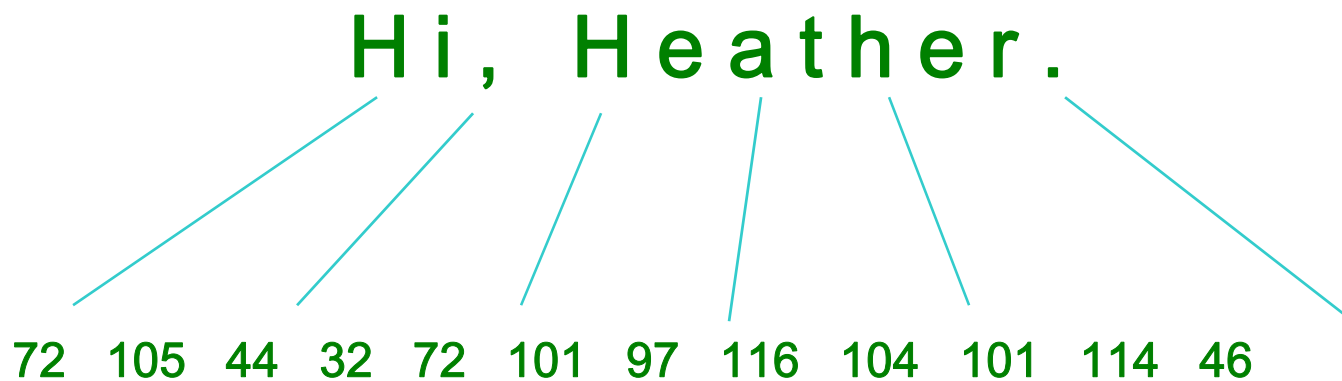


Software Categories

- Operating System
 - controls all machine activities
 - provides the user interface to the computer
 - manages resources such as the CPU and memory
 - Windows XP, Unix, Linux, Mac OS
- Application program
 - generic term for any other kind of software
 - word processors, missile control systems, games
- Most operating systems and application programs have a *graphical user interface* (GUI)

Representing Text Digitally

- Every character is stored as a number, including spaces, digits, and punctuation
- Corresponding upper and lower case letters are separate characters



- A common code is ASCII (American Standard Code for Information Interchange).



Binary Numbers

- Once information is digitized, it is represented and stored in memory using the *binary number system*
- A single binary digit (0 or 1) is called a *bit*
- Devices that store and move information are cheaper and more reliable if they have to represent only two states
- A single bit can represent two possible states, like a light bulb that is either on (1) or off (0)
- Permutations of bits are used to store values

Bit Permutations

- Each permutation can represent a particular item
- There are 2^N permutations of N bits
- Therefore, N bits are needed to represent 2^N unique items

**How many
items can be
represented by**

1 bit ?	$2^1 = 2$ items
2 bits ?	$2^2 = 4$ items
3 bits ?	$2^3 = 8$ items
4 bits ?	$2^4 = 16$ items
5 bits ?	$2^5 = 32$ items

Storing Information

9278

9279

9280

9281

9282

9283

9284

9285

9286



Each memory cell stores a set number of bits (usually 8 bits, or one *byte*)

Large values are stored in consecutive memory locations



Storage Capacity

- Every memory device has a *storage capacity*, indicating the number of bytes it can hold
- Capacities are expressed in various units:

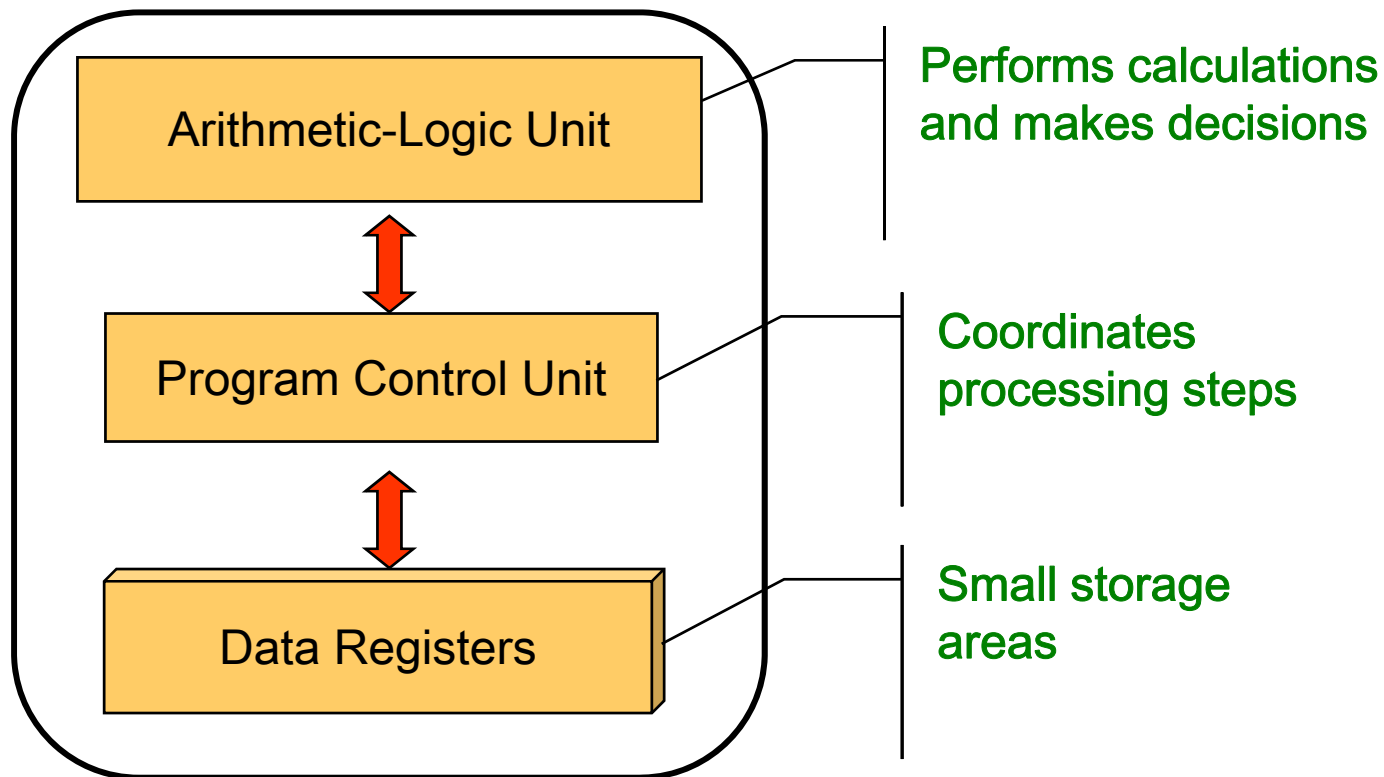
<u>Unit</u>	<u>Symbol</u>	<u>Number of Bytes</u>
kilobyte	KB	$2^{10} = 1024$
megabyte	MB	2^{20} (over 1 million)
gigabyte	GB	2^{30} (over 1 billion)
terabyte	TB	2^{40} (over 1 trillion)



Memory

- Main memory is *volatile* - stored information is lost if the electric power is removed
- Secondary memory devices are *nonvolatile*
- Main memory and disks are *direct access* devices - information can be reached directly
- The terms *direct access* and *random access* often are used interchangeably
- A magnetic tape is a *sequential access* device since its data is arranged in a linear order - you must get by the intervening data in order to access other information

The Central Processing Unit (CPU)



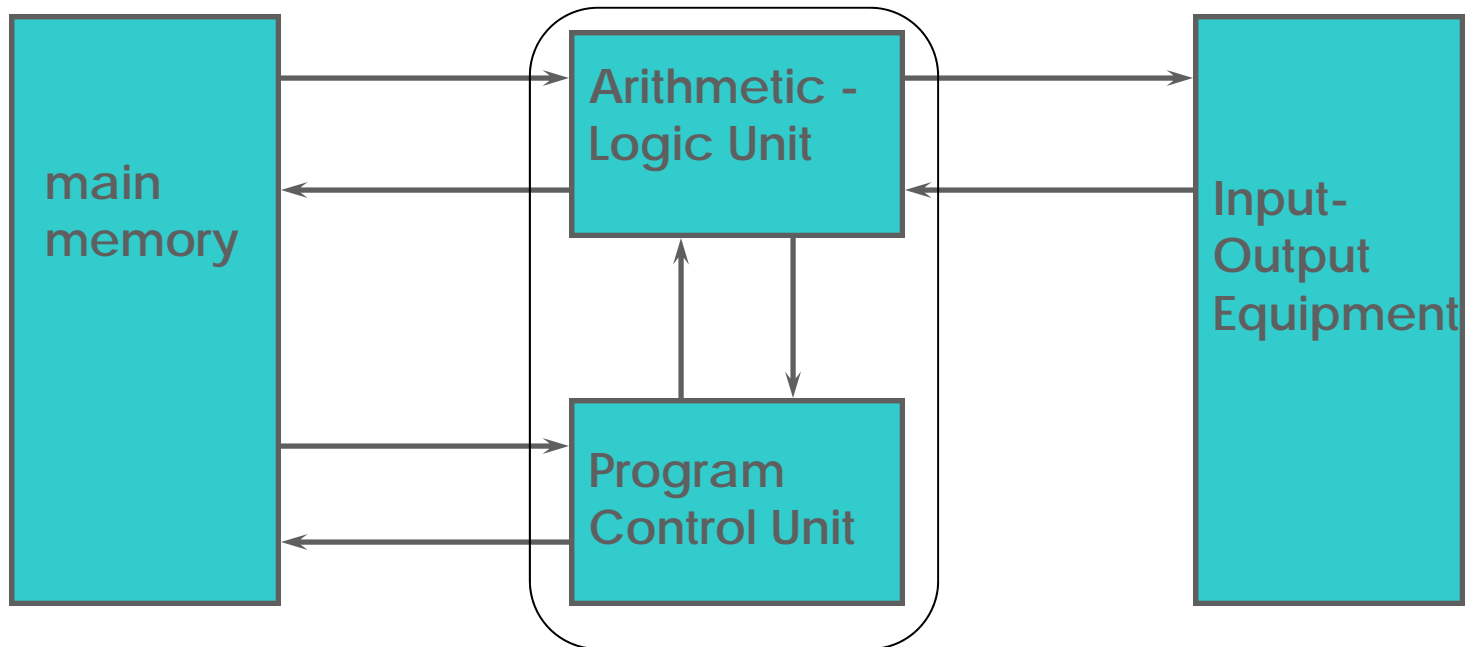
CPU - Microprocessor

- A CPU is on a chip called a *microprocessor*



Simple Machine Organization

- von Neumann machine
 - ALU performs transfers between memory and I/O devices



Simple CPU (no data registers)



CPU

- A CPU is a sequential circuit
- It repeatedly reads and executes an **instruction** from its memory
- Machine language program is a set of instructions drawn from CPU instruction set.
- An instruction is composed of one instruction **operation**, and a number of **operands**
 - Operand means the data needed for the operation
 - For example:
 - 1-operand format: use the accumulator (AC) register
 - e.g. Load A (It means $AC \leftarrow A$)



CPU speed

- The speed of a CPU is controlled by the *system clock*
- The system clock generates an electronic pulse at regular intervals
- The pulses coordinate the activities of the CPU
- The speed is usually measured in *gigahertz* (GHz)

e.g. 2.8 GHz means 2.8 billion pulses per second

Von Neumann Architecture

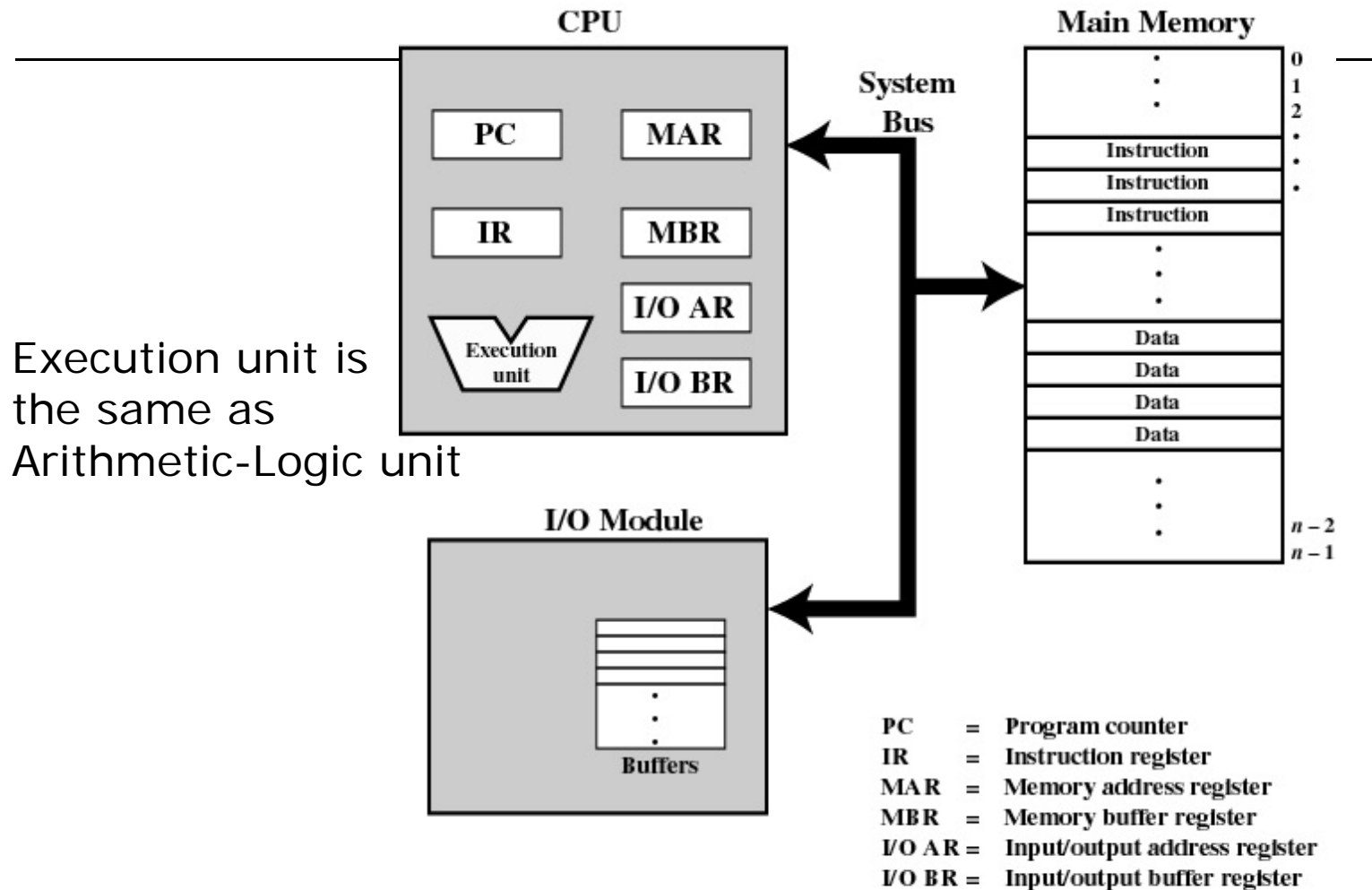


Figure 1.1 Computer Components: Top-Level View
SEEM 3460



Type of Instructions

- Processor-memory
 - transfer data between processor (register) and memory
- Processor-I/O
 - data transferred to or from a peripheral device
- Data processing
 - arithmetic or logic operation on data
- Control
 - alter sequence of execution

Instruction Set Principles

- Multiple data registers in CPU
- Data registers can be used to hold intermediate results during execution of instructions
- They are instruction-visible
e.g. R1, R2, R3
- 2-operand format: one of the operands is both a source and result
e.g. Add R1,B (It means $R1 \leftarrow R1 + B$)
- 3-operand format: result and two source operands
e.g. Add R3,R1,R2 (It means $R3 \leftarrow R1 + R2$)

Different Kinds of Instruction Sets

- Different kinds of computer architecture and the corresponding instruction sets

	1-operand format	2-operand format	3-operand format
	(AC) Accumulator	Register (register-memory)	Register (load-store)
Push A	Load A	Load R1,A	Load R1,A
Push B	Add B	Add R1,B	Load R2,B
Add	Store C	Store C,R1	Add R3,R1,R2
Pop C			Store C,R3

FIGURE 2.1 The code sequence for $C = A + B$ for four instruction sets.

- Examples of instruction operation: Load, Add, Store
- Examples of memory address: A, B, C
- Examples of registers: Accumulator (AC), R1, R2, R3

Example of Program Compilation

- Suppose there is a C statement in the source program:

$A = A + B$

- Suppose the CPU supports single-operand accumulator-kind of instructions.
- After compilation, this statement is translated into 3 instructions:

LOAD B	($AC \leftarrow \text{memory address B}$)
ADD A	($AC \leftarrow AC + \text{memory address A}$)
STORE A	($\text{memory address A} \leftarrow AC$)

- LOAD B means that it loads the content referenced by address B to the internal accumulator register AC
- ADD A means that it adds the content referenced by address A to the internal accumulator register AC
- STORE A means that it stores the content of the internal accumulator register AC into the memory location referenced by address A



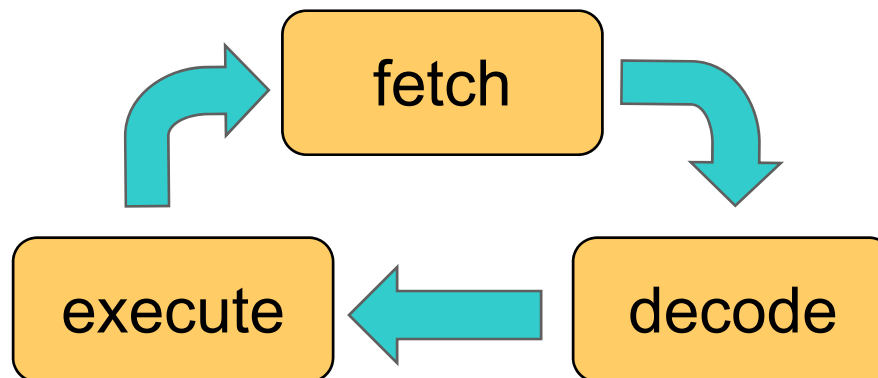
Instruction Set Principles – Control Registers

- Used by processor to control the operation of the processor
- Program Counter (PC)
 - Contains the address of an instruction to be fetched
- Instruction Register (IR)
 - Contains the instruction most recently fetched

Instruction Cycle

- The CPU continuously follows the *fetch-decode-execute cycle*:

Retrieve an instruction from main memory



Carry out the instruction

Determine what the instruction is



Instruction Fetch and Execute

- Program counter (PC) holds the address of the instruction to be fetched next.
- The processor fetches the instruction from memory.
- Program counter is incremented by one unit after each fetch so that the processor knows where the next instruction is.
- Fetched instruction is placed in the instruction register (IR).



Example of Program Execution

- Recall that the C statement $A = A + B$ has been compiled into the following instructions:
 - LOAD B ($AC \leftarrow \text{memory address B}$)
 - ADD A ($AC \leftarrow AC + \text{memory address A}$)
 - STORE A ($\text{memory address A} \leftarrow AC$)
- When the user executes the executable file, the operating system will:
 - place those three instructions starting from a memory address such as 300
 - allocate a memory address such as 940 for the variable B
 - allocate a memory address such as 941 for the variable A

Example of Program Execution

0001 1001 0100 0000
1 9 4 0
LOAD 940

0101 1001 0100 0001
5 9 4 1
ADD 941

0010 1001 0100 0001
2 9 4 1
STORE 941

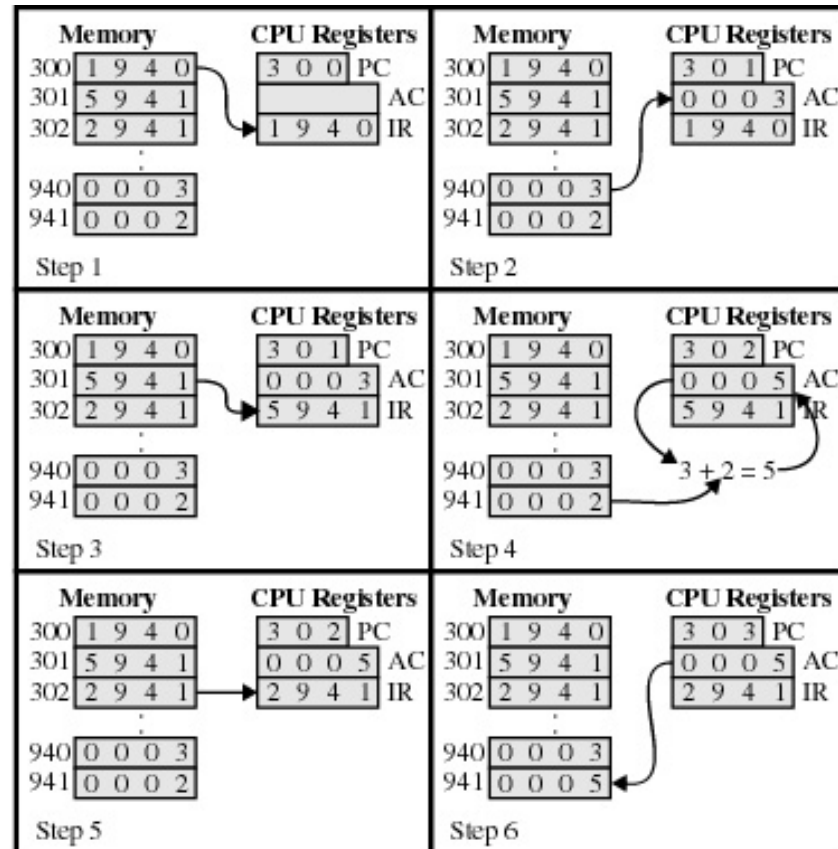


Figure 1.4 Example of Program Execution
(contents of memory and registers in hexadecimal)

SEEM 3460

Instruction Set Principles – Instruction Encoding Examples

Operation & no. of operands	Address specifier 1	Address field 1	...	Address specifier n	Address field n
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(a) Variable (e.g., VAX)

Operation	Address field 1	Address field 2	Address field 3
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(b) Fixed (e.g., DLX, MIPS, Power PC, Precision Architecture, SPARC)

Operation	Address specifier	Address field
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Operation	Address specifier 1	Address specifier 2	Address field
-----------	------------------------	------------------------	------------------

Operation	Address specifier	Address field 1	Address field 2
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(c) Hybrid (e.g., IBM 360/70, Intel 80x86)



Instruction Set Principles – Status Registers

- Used for controlling the execution of instructions
- Program Status Word (PSW)
 - Condition Codes or Flags
 - Bits set by the processor hardware as a result of operations
 - Can be accessed by a program but not altered
 - e.g. positive result, negative result, zero, overflow

Examples of Control Flow Instructions

Example instruction	Instruction name	Meaning
J name	Jump	$PC \leftarrow \text{name}; ((PC+4) - 2^{25}) \leq \text{name} < ((PC+4) + 2^{25})$
JAL name	Jump and link	$R31 \leftarrow PC+4; PC \leftarrow \text{name};$ $((PC+4) - 2^{25}) \leq \text{name} < ((PC+4) + 2^{25})$
JALR R2	Jump and link register	$\text{Regs}[R31] \leftarrow PC+4; PC \leftarrow \text{Regs}[R2]$
JR R3	Jump register	$PC \leftarrow \text{Regs}[R3]$
BEQZ R4, name	Branch equal zero	if $(\text{Regs}[R4] == 0)$ $PC \leftarrow \text{name};$ $((PC+4) - 2^{15}) \leq \text{name} < ((PC+4) + 2^{15})$
BNEZ R4, name	Branch not equal zero	if $(\text{Regs}[R4] \neq 0)$ $PC \leftarrow \text{name};$ $((PC+4) - 2^{15}) \leq \text{name} < ((PC+4) + 2^{15})$

FIGURE 2.24 Typical control-flow instructions in DLX.

ALU Components

- ALU does arithmetic and logical comparisons
- AC = accumulator holds results
 - MQ = memory-quotient holds second portion of long results
 - MBR = memory buffer register holds data while operation executes

