

SOLUTIONS FOR CONTEST 0

SOLUTION 1

LANGUAGE: C++

```
/*
```

```
    Time complexity:  $O(N)$ 
```

```
    Space complexity:  $O(N)$ 
```

```
    where 'N' is the length of the array A.
```

```
*/
```

```
int minCostFlip (int n, vector <int> &a, vector <int> &b) {
```

```
    // Find prefix suffix '0's.
```

```
    vector <int> prefix(n), suffix(n);
```

```
    for (int i = 0; i < n; i++) {
```

```
        if (b[i] == 0) {
```

```
            prefix[i] = 1 + (i - 1 >= 0 ? prefix[i - 1] : 0);
```

```
        }
```

```
    }
```

```
    for (int i = n - 1; i >= 0; i--) {
```

```
        if (b[i] == 0) {
```

```
            suffix[i] = 1 + (i + 1 < n ? suffix[i + 1] : 0);
```

```
        }
```

```
    }
```

```
    int mxLen = 0, cost = 0;
```

```
    for (int i = 0; i < n; i++) {
```

```

// Find the max length you can make by flipping the current '1' if available.
int curLen = (i - 1 >= 0 ? prefix[i - 1] : 0) + (i + 1 < n ? suffix[i + 1] : 0);
if (b[i] == 1) {
    curLen += 1;
    if (mxLen < curLen) {
        mxLen = curLen;
        cost = a[i];
    }
    else if (curLen == mxLen) {
        cost = min(cost, a[i]);
    }
}
}
return cost;
}

```

LANGUAGE: JAVA

/*

Time complexity: O(N)

Space complexity: O(N)

where 'N' is the length of the array A.

*/

```
import java.util.ArrayList;
```

```
import java.util.Collections;
```

```
public class Solution {
```

```
    static int minCostFlip(int n, ArrayList<Integer> a, ArrayList<Integer> b) {
```

```
        // Find prefix and suffix '0's.
```

```
        ArrayList<Integer> prefix = new ArrayList<>(Collections.nCopies(n, 0));
```

```
        ArrayList<Integer> suffix = new ArrayList<>(Collections.nCopies(n, 0));
```

```
        for (int i = 0; i < n; i++) {
```

```

        if (b.get(i) == 0) {
            prefix.set(i, 1 + (i - 1 >= 0 ? prefix.get(i - 1) : 0));
        }
    }
    for (int i = n - 1; i >= 0; i--) {
        if (b.get(i) == 0) {
            suffix.set(i, 1 + (i + 1 < n ? suffix.get(i + 1) : 0));
        }
    }

    int mxLen = 0, cost = 0;
    for (int i = 0; i < n; i++) {
        // Find the max length you can make by flipping the current '1' if available.
        int curLen = (i - 1 >= 0 ? prefix.get(i - 1) : 0) + (i + 1 < n ? suffix.get(i + 1) : 0);
        if (b.get(i) == 1) {
            curLen += 1;
            if (curLen > mxLen) {
                mxLen = curLen;
                cost = a.get(i);
            } else if (curLen == mxLen) {
                cost = Math.min(cost, a.get(i));
            }
        }
    }
    return cost;
}
}

```

LANGUAGE: PYTHON

.....

Time complexity: $O(N)$

Space complexity: $O(N)$

where 'N' is the length of the array A.

.....

```

from typing import *

def min_cost_flip(n: int, a: List[int], b: List[int]):

    # Find prefix suffix '0's.
    prefix = [0] * n
    suffix = [0] * n
    for i in range(n):
        if b[i] == 0:
            prefix[i] = 1 + (prefix[i - 1] if i - 1 >= 0 else 0)
    for i in range(n - 1, -1, -1):
        if b[i] == 0:
            suffix[i] = 1 + (suffix[i + 1] if i + 1 < n else 0)

    mx_len = 0
    cost = 0
    for i in range(n):
        # Find the max length you can make by flipping the current '1' if available.
        cur_len = (prefix[i - 1] if i - 1 >= 0 else 0) + (suffix[i + 1] if i + 1 < n else 0)
        if b[i] == 1:
            cur_len += 1
            if mx_len < cur_len:
                mx_len = cur_len
                cost = a[i]
            elif cur_len == mx_len:
                cost = min(cost, a[i])
    return cost

```

SOLUTION 2:

LANGUAGE : C++

/*

Time Complexity: $O(N \log N)$

Space Complexity: $O(1)$

```

    where 'N' is the number of elements in the array.
*/

int minPrice(int n, int k, vector<int> &a){

    sort(a.begin(), a.end());

    // 'l' denotes the leftover elements after selecting 'k' elements.
    int l=n-k;
    int ans=a[l-1]-a[0];

    // Selecting 'k' elements from the array and
    // finding the minimum difference between the maximum and minimum element.
    for(int i=l;i<n;i++){
        ans=min(ans,a[i]-a[i-l+1]);
    }
    return ans;
}

```

LANGUAGE : JAVA

```

/*
    Time Complexity: O(N log N)
    Space Complexity: O(1)

    where 'N' is the number of elements in the array.
*/

import java.util.Arrays;

public class Solution {

    static int minPrice(int n, int k, int[] a) {
        Arrays.sort(a);

        // 'l' denotes the leftover elements after selecting 'k' elements.
        int l = n - k;
        int ans = a[l - 1] - a[0];

        // Selecting 'k' elements from the array and
        // finding the minimum difference between the maximum and minimum element.

```

```

    for (int i = l; i < n; i++) {
        ans = Math.min(ans, a[i] - a[i - l + 1]);
    }
    return ans;
}
}

```

LANGUAGE: PYTHON

"""

Time Complexity: $O(N \log N)$

Space Complexity: $O(1)$

where 'N' is the number of elements in the array.

"""

from typing import *

def min_price(n: int, k: int, a: List[int]) -> int:

Sorting the array

a.sort()

'l' denotes the leftover elements after selecting 'k' elements.

l = n - k

Initialize the answer with the difference between the maximum and minimum element.

ans = a[l - 1] - a[0]

Selecting 'k' elements from the array and

finding the minimum difference between the maximum and minimum element.

for i in range(l, n):

ans = min(ans, a[i] - a[i - l + 1])

return ans

SOLUTION 3:

LANGUAGE: C++

/*

Time Complexity: $O(n)$

Space Complexity: $O(1)$

Where 'n' denotes the length of the vector 'v'.

*/

```
int maximumAlternateSum(int n, vector<int> &v) {
```

```
    // Initialize an integer 'suffix' equal to '0'.
```

```
    int suffix = 0;
```

```
    // Calculate the alternate sum of the vector 'v' and store into 'suffix'.
```

```
    for (int i = 0; i < n; i++) {
```

```
        if (i % 2 == 0) {
```

```
            suffix += v[i];
```

```
        }
```

```
        else {
```

```
            suffix -= v[i];
```

```
        }
```

```
    }
```

```
    // Initialize the integers 'prefix' and 'answer' equal to '0' and '-Infinity', respectively.
```

```
    int prefix = 0, answer = -1e9;
```

```
    // Iterate through every element of the vector 'v'.
```

```
    for (int i = 0; i < n; i++) {
```

```
        // Update the 'suffix' by removing the contribution of 'v[i]'.  
        if (i % 2 == 0) {
```

```

        suffix -= v[i];
    }
    else {
        suffix += v[i];
    }

    // Update the 'answer' based on the alternate sum obtained by removing the
    element 'v[i]'.
    int currentSum = prefix - suffix;
    answer = max(answer, currentSum);

    // Update 'prefix' by adding the contribution of 'v[i]'.
    if (i % 2 == 0) {
        prefix += v[i];
    }
    else {
        prefix -= v[i];
    }
}

// Return 'answer'.
return answer;
}

```

LANGUAGE: JAVA

/*

Time Complexity: $O(n)$

Space Complexity: $O(1)$

Where 'n' denotes the length of the array 'v'.

*/

```
public class Solution {
```

```
    public static int maximumAlternateSum(int n, int[] v) {
```

```
        // Initialize an integer 'suffix' equal to '0'.
```

```
        int suffix = 0;
```

```
        // Calculate the alternate sum of the array 'v' and store into 'suffix'.
```

```
        for (int i = 0; i < n; i++) {
```



```

        if (i % 2 == 0) {
            suffix += v[i];
        } else {
            suffix -= v[i];
        }
    }
}

// Initialize the integers 'prefix' and 'answer' equal to '0' and '-Infinity',
respectively.
int prefix = 0, answer = Integer.MIN_VALUE;

// Iterate through every element of the array 'v'.
for (int i = 0; i < n; i++) {

    // Update the 'suffix' by removing the contribution of 'v[i]'.
    if (i % 2 == 0) {
        suffix -= v[i];
    } else {
        suffix += v[i];
    }

    // Update the 'answer' based on the alternate sum obtained by removing the
    element 'v[i]'.
    int currentSum = prefix - suffix;
    answer = Math.max(answer, currentSum);

    // Update 'prefix' by adding the contribution of 'v[i]'.
    if (i % 2 == 0) {
        prefix += v[i];
    } else {
        prefix -= v[i];
    }
}

// Return 'answer'.
return answer;
}
}

```

LANGUAGE: PYTHON

.....

Time Complexity: $O(n)$

Space Complexity: $O(1)$

Where 'n' denotes the length of the vector 'v'.

.....

```
from typing import List
```

```
def maximumAlternateSum(n: int, v: List[int]) -> int:
```

```
    # Initialize an integer 'suffix' equal to '0'.
```

```
    suffix = 0
```

```
    # Calculate the alternate sum of the vector 'v' and store into 'suffix'.
```

```
    for i in range(n):
```

```
        if i % 2 == 0:
```

```
            suffix += v[i]
```

```
        else:
```

```
            suffix -= v[i]
```

```
    # Initialize the integers 'prefix' and 'answer' equal to '0' and '-Infinity', respectively.
```

```
    prefix = 0
```

```
    answer = float('-inf')
```

```
    # Iterate through every element of the vector 'v'.
```

```
    for i in range(n):
```

```
        # Update the 'suffix' by removing the contribution of 'v[i]'.
```

```
        if i % 2 == 0:
```

```
            suffix -= v[i]
```

```
        else:
```

```
            suffix += v[i]
```

```
    # Update the 'answer' based on the alternate sum obtained by removing the  
    element 'v[i]'.
```

```
    currentSum = prefix - suffix
```

```
    answer = max(answer, currentSum)
```

```
    # Update 'prefix' by adding the contribution of 'v[i]'.
```

```
    if i % 2 == 0:
```

```
    prefix += v[i]
else:
    prefix -= v[i]

# Return 'answer'.
return answer
```

SOLUTION 4:

Will be discussed during the meet.

SOLUTION 5:

LANGUAGE: C++

/*

Time Complexity: $O(n^2)$

Space Complexity: $O(n)$

Where 'n' is the length of the array 'a'.

***/**

```
vector<int> nextGreaterElementII(vector<int>& a) {  
    int n = a.size();
```

```
    // Declare an array 'answer' of size 'n',  
    // to store the Next Greater Element for each element.  
    vector<int> answer(n);
```

```
    // Run a for loop from i = 0 to 'n' - 1.  
    for (int i = 0; i < n; i++) {
```

```
        // Initialise an integer variable 'currentAnswer' to -1.  
        int currentAnswer = -1;  
        for (int j = i + 1; j < n; j++) {
```

```

        // If a[(i + j) % n] > a[i] then update
        // 'currentAnswer' to a[(i + j) % n] and break.
        if (a[(i + j) % n] > a[i]) {
            currentAnswer = a[(i + j) % n];
            break;
        }
    }

    // Update answer[i] to 'currentAnswer'.
    answer[i] = currentAnswer;
}

// Return 'answer' as the answer to the problem.
return answer;
}

```

LANGUAGE: JAVA

```

/*
    Time Complexity: O(N^2)
    Space Complexity: O(N)

    Where 'N' is the length of the array 'A'.
*/

public class Solution {
    public static int[] nextGreaterElementII(int []a) {
        int n = a.length;

        // Initialize an array 'answer' of size 'N',
        // to store the Next Greater Element for each element.
        int []answer = new int[n];

        // Run a for loop from i=0 to N-1.
        for (int i = 0; i < n; i++) {

```

```

// Initialise an integer variable 'currentAnswer' to -1.
int currentAnswer = -1;
for (int j = 1; j < n; j++) {

    // If A[ (i+j)%N ] > A[ i ] then update
    // 'currentAnswer' to A[ (i+j)%N ] and break.
    if (a[(i + j) % n] > a[i]) {
        currentAnswer = a[(i + j) % n];
        break;
    }
}

// Update answer[ i ] to 'currentAnswer'.
answer[i] = currentAnswer;
}

// Return 'answer' as the answer to the problem.
return answer;
}
}

```

LANGUAGE: PYTHON

'''

Time Complexity: $O(N^2)$

Space Complexity: $O(N)$

Where 'N' is the length of the array 'A'.

'''

def nextGreaterElementII(a: int) -> int:

n = len(a)

```

# Initialize an array  $\tilde{answer}$  of size  $\tilde{N}$ ,
# to store the Next Greater Element for each element.

answer = [0 for i in range(n)]

# Run a for loop from i=0 to N-1.

for i in range(n):

    # Initialise an integer variable  $\tilde{currentAnswer}$  to -1.

    currentAnswer = -1

    for j in range(n):

        # If  $A[(i+j)\%N] > A[i]$  then update
        #  $\tilde{currentAnswer}$  to  $A[(i+j)\%N]$  and break.

        ind = int((i + j) % n)

        if (a[ind] > a[i]):

            currentAnswer = a[int((i + j) % n)]

            break

    # Update answer[ i ] to  $\tilde{currentAnswer}$ .

    answer[i] = currentAnswer

# Return  $\tilde{answer}$  as the answer to the problem.
return answer

```

SOLUTION 6:

LANGUAGE: C++

/*

Time Complexity: $O(n * m)$

Space Complexity: $O(n * m)$

Where 'n' and 'm' denote the number of vectors and the number of elements in each vector, respectively.

***/**

vector<int> lexicographicallyMinimum(int n, int m, vector<vector<int>> &v) {

// Initialize a two-dimensional vector 'valid' to store the dp states.

vector<vector<bool>> valid(n, vector<bool>(m, false));

// Initialize an integer 'nextMax' to infinity.

int nextMax = 1e9;

for (int j = m - 1; j >= 0; j--) {

// Initialize an integer 'currentMax' equal to '0'.

int currentMax = 0;

for (int i = 0; i < n; i++) {

// Calculate the value of 'valid[i][j]'.

if (v[i][j] <= nextMax) {

valid[i][j] = true;

currentMax = max(currentMax, v[i][j]);

}

}

// Update 'nextMax' to 'currentMax'.

nextMax = currentMax;

}

// Initialize a vector 'answer'.

vector<int> answer(m);

// Initialize an integer 'previous' equal to '0'.

int previous = 0;

for (int j = 0; j < m; j++) {

for (int i = 0; i < n; i++) {

// Find the minimum value of 'j' to perform the replacement operation.

if (valid[i][j] && v[i][j] >= previous) {


```

        previous = v[i][j];
        answer[j] = i;
        break;
    }

    // Return '-1' if there is no such 'j'.
    if (i == n - 1) {
        return {-1};
    }
}

// Return the 'answer'.
return answer;
}

```

LANGUAGE: JAVA

/*

Time Complexity: $O(n * m)$

Space Complexity: $O(n * m)$

Where 'n' and 'm' denote the number of vectors and the number of elements in each vector, respectively.

*/

```
import java.util.*;
```

```

public class Solution {
    public static int[] lexicographicallyMinimum(int n, int m, int[][] v) {
        // Initialize a two-dimensional array 'valid' to store the dp states.
        boolean[][] valid = new boolean[n][m];

        // Initialize an integer 'nextMax' to infinity.
        int nextMax = (int) 1e9;
        for (int j = m - 1; j >= 0; j--) {
            // Initialize an integer 'currentMax' equal to '0'.
            int currentMax = 0;
            for (int i = 0; i < n; i++) {

```

```

        // Calculate the value of 'valid[i][j]'.
        if (v[i][j] <= nextMax) {
            valid[i][j] = true;
            currentMax = Math.max(currentMax, v[i][j]);
        }
    }
    // Update 'nextMax' to 'currentMax'.
    nextMax = currentMax;
}

// Initialize an array 'answer'.
int[] answer = new int[m];

// Initialize an integer 'previous' equal to '0'.
int previous = 0;
for (int j = 0; j < m; j++) {
    for (int i = 0; i < n; i++) {
        // Find the minimum value of 'j' to perform the replacement operation.
        if (valid[i][j] && v[i][j] >= previous) {
            previous = v[i][j];
            answer[j] = i;
            break;
        }
        // Return '-1' if there is no such 'j'.
        if (i == n - 1) {
            return new int[]{-1};
        }
    }
}
// Return the 'answer'.
return answer;
}
}

```

LANGUAGE: PYTHON

.....

Time Complexity: $O(n * m)$

Space Complexity: $O(n * m)$

Where 'n' and 'm' denote the number of vectors and the number of elements in each vector, respectively.

.....

```
from typing import List
```

```
def lexicographicallyMinimum(n: int, m: int, v: List[List[int]]) -> List[int]:
```

```
    # Initialize a two-dimensional list 'valid' to store the dp states.
```

```
    valid = [[False] * m for _ in range(n)]
```

```
    # Initialize an integer 'nextMax' to infinity.
```

```
    nextMax = float('inf')
```

```
    for j in range(m - 1, -1, -1):
```

```
        # Initialize an integer 'currentMax' equal to '0'.
```

```
        currentMax = 0
```

```
        for i in range(n):
```

```
            # Calculate the value of 'valid[i][j]'.
```

```
            if v[i][j] <= nextMax:
```

```
                valid[i][j] = True
```

```
                currentMax = max(currentMax, v[i][j])
```

```
    # Update 'nextMax' to 'currentMax'.
```

```
    nextMax = currentMax
```

```
    # Initialize a list 'answer'.
```

```
    answer = [0] * m
```

```
    # Initialize an integer 'previous' equal to '0'.
```

```
    previous = 0
```

```
    for j in range(m):
```

```
        for i in range(n):
```

```
            # Find the minimum value of 'j' to perform the replacement operation.
```

```
            if valid[i][j] and v[i][j] >= previous:
```

```
                previous = v[i][j]
```

```
                answer[j] = i
```

```
                break
```

```
    # Return '-1' if there is no such 'j'.
```

```
    if i == n - 1:
```

```
        return [-1]
```

```
# Return the 'answer'.  
return answer
```