

Cloud Computing

Hardware virtualization

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Introduction



Hardware-level Virtualization

➤ An abstract execution environment in terms of computer

hardware on top of which a guest operating system can be run.

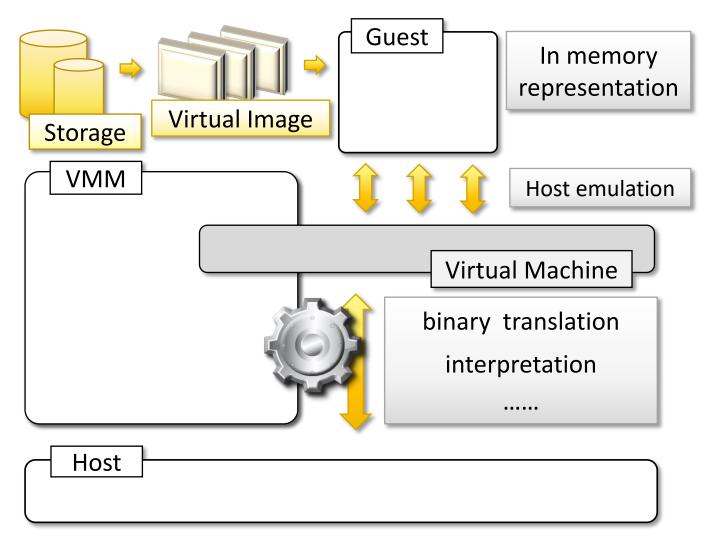
Concept	Represented by	
Guest	Operating system	
Host	Physical computer hardware	
Virtual machine	Its emulation	
Virtual machine manager	Hypervisor	

What is Hypervisor?

Hypervisor is a program enabling the abstraction of the underlying physical hardware.

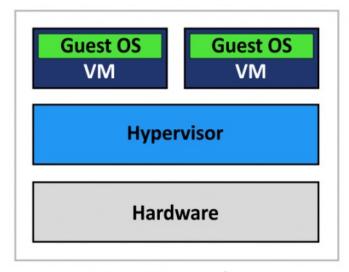
Hypervisor is also called Virtual Machine Manager (VMM)

Hardware-level Virtualization



Types of Hypervisor

- > Type I hypervisors (native VM)
 - Run directly on top of the hardware.
 - Take the place of the operating systems
 - Interact directly with the ISA interface



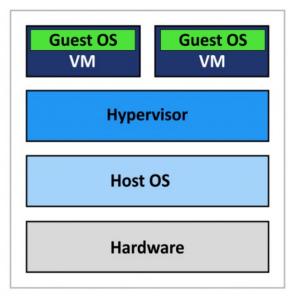
Type 1 Hypervisor (Bare-Metal Architecture)

Source: http://:

https://www.nakivo.com/blog/hyper-v-virtualbox-one-choose-infrastructure/

Types of Hypervisor (cont.)

- > Type II hypervisors (hosted VM)
 - Require the support of an operating system
 - Are programs managed by the operating system
 - Interact with operating system through the ABI.



Type 2 Hypervisor (Hosted Architecture)

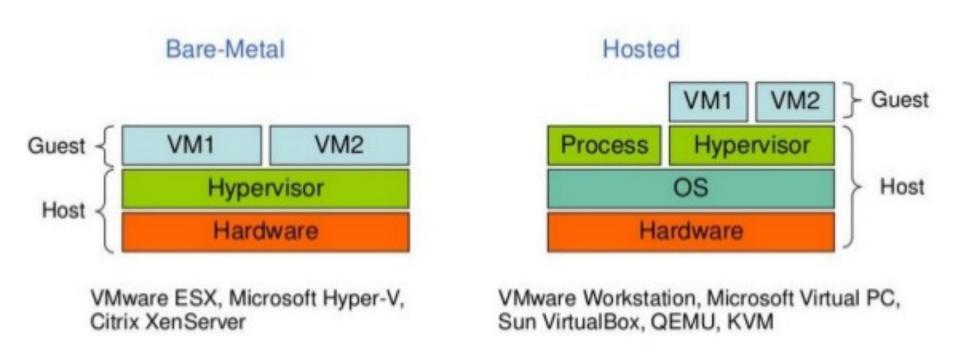
Source: http://:

https://www.nakivo.com/blog/hyp

er-v-virtualbox-one-choose-

infrastructure/

Type of Hypervisors (cont.)



Source: https://www.slideshare.net/PraveenHanchinal/virtualizationthe-cloud-enabler-by-inspiregroups/18-Types of hypervisors VMM

Approaches of Executing

Guest Instructions



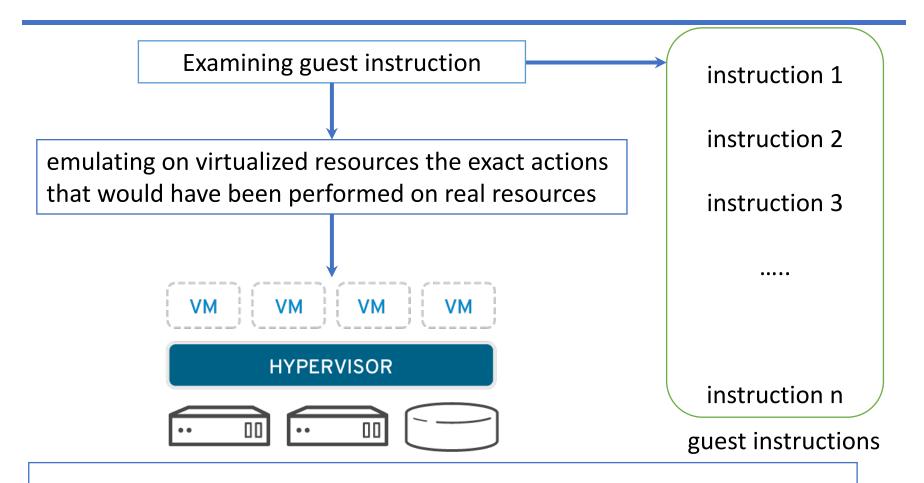
Executing Guest Instructions

- **Emulation**
- **▶** Direct native execution

Emulation

"the process of implementing the interface and functionality of one system or subsystem on a system or subsystem having a different interface and functionality..."

Emulation (cont.)



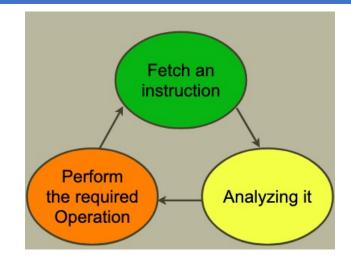
Only available mechanism when the ISA of the guest is different from the ISA of the host.



Emulation Approaches

≻Interpretation

Done in software,one instruction at a time

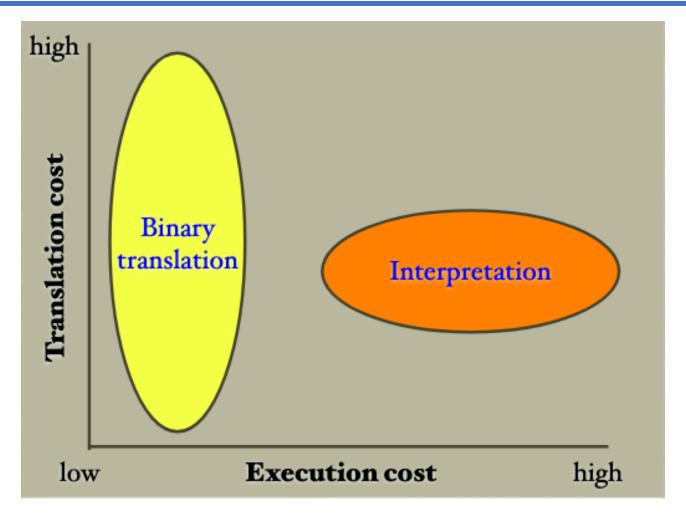


http://cse.unl.edu/~witty/class/emb
edded/material/note/emulation.pdf

▶ Binary translation

- Translating a block of source instructions to target instructions.
- Saving the translated code for repeated use

Interpretation versus Binary Translation



http://se.unl.edu/~witty/class/embedded/material/note/emulation.pdf

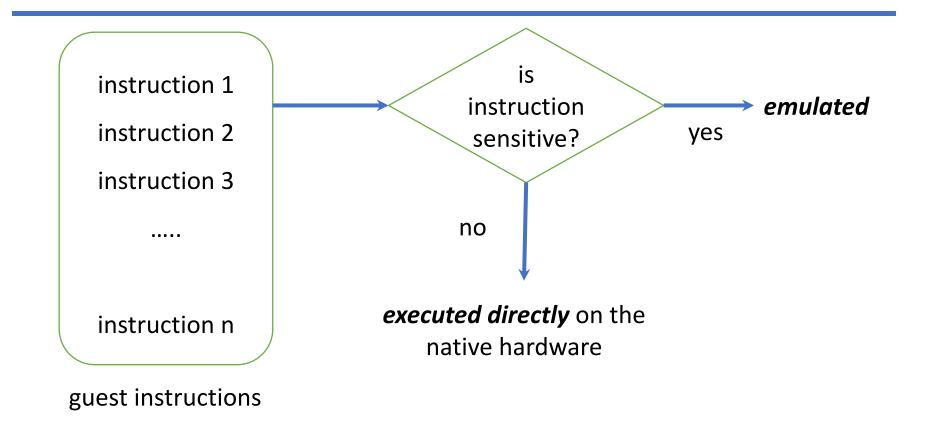


Interpretation versus Binary Translation (cont.)

	Implementation	Performance
Interpretation	simple and easy	low
Binary Translation	complex	high initial translation cost, small execution cost

http://www.ittc.ku.edu/~kulkarni/teaching/EECS768/slides/chapter2.pdf

Direct Native Execution



Only if the ISA of the host is identical to the ISA of the guest.

Hardware Virtualization Methods



Hardware Virtualization Methods

> Full Virtualization

- Binary Translation
- Hardware-assisted virtualization

→ Paravirtualization

Full Virtualization

- > Run a program directly on top of a VM and without any modification
 - The program thought it were run on the raw hardware.

- ➤ The principal advantage of full virtualization
 - Complete isolation → enhanced security
 - Ease of emulation of different architectures
 - Coexistence of different systems on the same platform.

Full Virtualization (cont.)

- ➤ In some architectures, some sensitive instructions are not privileged
 - They cannot be virtualized in the classic way.
 - Like the non-hardware-assisted x86

> Two technologies:

- Binary translation
- Hardware-assisted virtualization

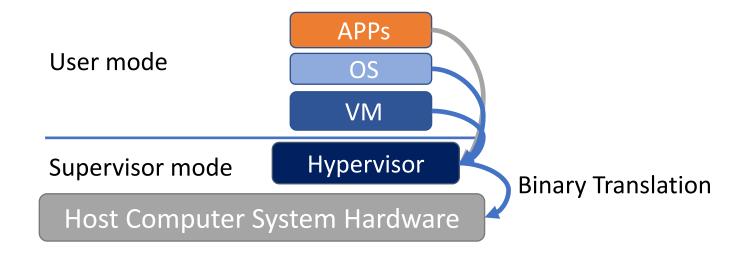
First scenario: Host ISA != Guest ISA



Main Challenge

- Guest Instruction can not directly executed on the host
- ➤ Feasible virtualization technologies
 - Interpretation
 - Binary Translation

Binary Translation



Binary Translation (cont.)

- **►** Static Binary Translation
 - On a full program
- Dynamic Binary Translation
 - Introduces an additional overhead.

Dynamic Binary Translation

- ➤ It is usually performed in small units called "basic blocks".
- ➤ A basic block is a set of instructions that ends with a branch instruction but does not have any branch instructions inside.
 - Be executed start to finish by a CPU
 - An ideal unit for translation
- > The translations of the basic blocks are cached
 - Overhead of translating only happens the first time a block is executed.

https://blogs.oracle.com/ravello/nested-virtualization-with-binary-translation

Static vs Binary Translation

	Input type	Granularity	Translation time
Static binary translation	Binary program	Full program	Before running program
Dynamic binary translation	Binary program	Basic block	At runtime

Demo

- ➤ Host: Mac M1 chip (ARM-based processor)
- ➤ Guest: Ubuntu ISO compiled for X86 processor (AMD)
- ➤ Start time of Ubuntu AMD ISO installation: 8:45
- Finish time: 10:05
- ➤ Why it takes so long?

Second scenario:
Host ISA = Guest ISA =
Original ISA Or
Non-hardware assisted X86

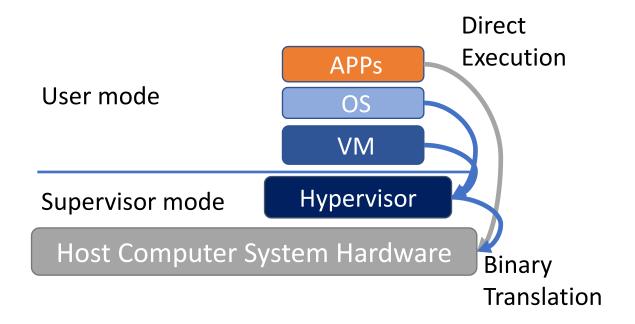


Main Challenge

- ➤ Some sensitive instructions are not privileged.
- They cannot be virtualized in the classic way.
- > Feasible virtualization technologies
 - Interpretation
 - Binary Translation and direct execution-based technique

Binary Translation and direct executionbased technique

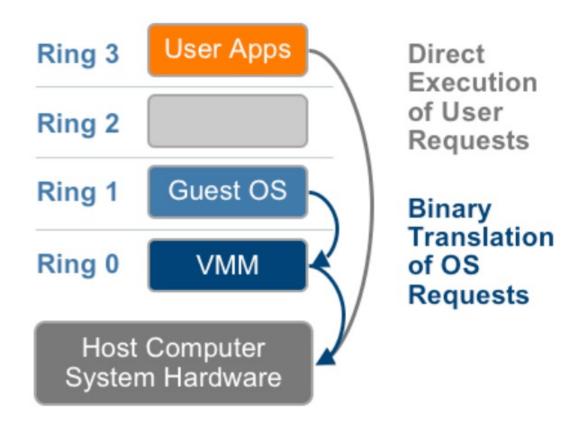
- The hypervisor translates all the OS instructions on the fly.
- The user code that typically runs on Ring 3 is directly executed.





Feasible Virtualization Technology

Binary Translation and direct execution-based technique



Third Scenario:
Host ISA = Guest ISA =
Hardware assisted X86
or
Hardware assisted ARM



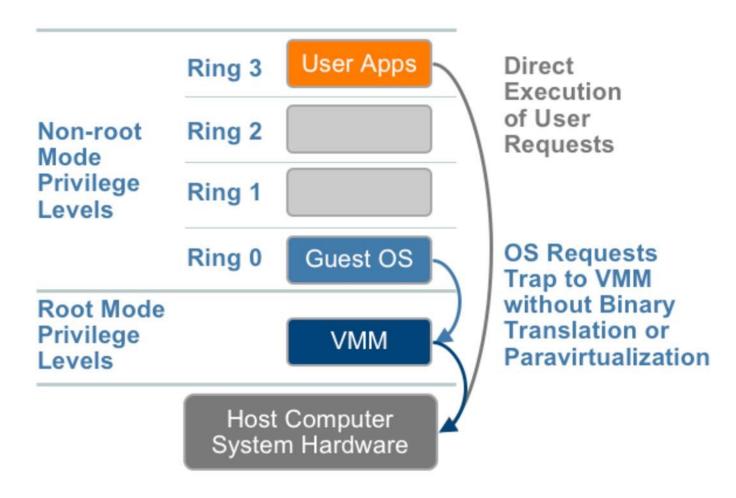
Hardware-assisted Virtualization

- ➤ Architectural support for building a VMM able to run a guest operating system in complete isolation.
- ➤ This technique was originally introduced in the IBM System/370.

- Extensions to x86-64b architecture
 - Introduced with Intel-VT and AMD-V.



Hardware-assisted Virtualization (cont.)



https://thecustomizewindows.com/2014/09/hardware-assisted-virtualization/

Intel-VT and AMD-V

- New CPU execution mode feature
- This allows the VMM to run in a new root mode below ring 0
 - Ring OP: privileged root mode (VMM)
 - Ring 0D : de-privileged non-root mode (Guest OS)

- Sensitive calls are set **to automatically trap** to the hypervisor and handled by hardware
 - Removing the need for either binary translation or para-virtualization.

Intel-VT

➤ Main feature: inclusion of the new VMX mode of operation.

	all four IA-32 privilege levels (rings)	VMX instructions
VMX non-root operation		
VMX root operation		

VMX Instructions

"VMX" stands for Virtual Machine Extensions

13 new instructions

VMPTRLD	VMPTRST	VMCLEAR	VMREAD	VMWRITE
VMCALL	VMLAUNCH	VMRESUME	VMXOFF	VMXON
INVEPT	INVVPID	VMFUNC		

Permit entering and exiting a *virtual execution mode* where the *guest OS perceives* itself as running with full privilege (ring 0), but the *host OS remains protected*.

Hardware-assisted Virtualization

The behavior of the processor in *non-root operation is limited* in some respects from its behavior on a normal processor.

➤ Critical shared resources are kept under the control of a monitor running in VMX root operation.

- VMM is run in VMX root mode
- Virtual machine and the guest OS are run in non-root mode.

Examples of Hardware-assisted Virtualization

➤ VirtualBox

➤ VMware

➤ Microsoft Hyper-V



Demo

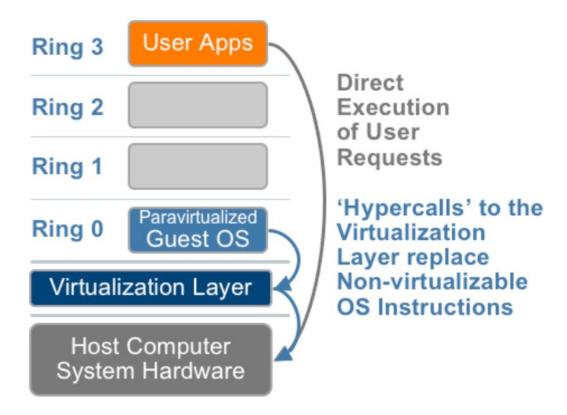
- ➤ Host: Mac M1 chip (ARM-based processor)
- ➤ Guest: Ubuntu ISO compiled for ARM processor
- Time taken to install Ubuntu ARM ISO: less than 10 mins
- ➤ Let's do a simple benchmark

Paravirtualization



Paravirtualization

➤ Paravirtualization refers to communication between the guest OS and the hypervisor to improve performance and efficiency.



Paravirtualization (cont.)

- ► It is not a transparent virtualization solution
 - Allows implementing thin virtual machine managers.
 - Remapping the performance-critical operations through the virtual machine software interface.

- ➤ Expose a software interface to the virtual machine that is slightly modified from the host
 - As consequence, guests need to be modified.

Paravirtualization (cont.)

- Provide the capability to demand the execution of performance-critical operations directly on the host
 - Preventing performance losses that would otherwise be experienced in managed execution.

- > Allows a simpler implementation of virtual machine managers
 - VMM have to simply transfer the execution of performance-critical operations directly to the host.
 - These instructions were hard to virtualize

Paravirtualization (Cont.)

>Xen is *the most popular implementation* of paravirtualization.



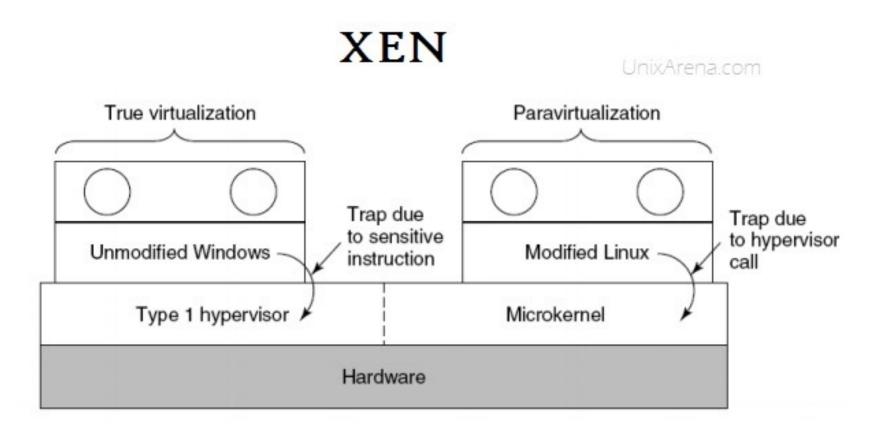
- The guest operating systems need to be changed
- The sensitive system calls need to be re-implemented with *hypercalls*
 - Are specific calls exposed by the virtual machine interface of Xen.

Paravirtualization (Cont.)

- ➤ With the use of *hypercalls*, the Xen hypervisor is able to
 - catch the execution of all the sensitive instructions
 - manage them,
 - and return the control

to the guest operating system by means of a supplied handler.

Xen Hypervisor



Xen supports both Full virtualization and Para-virtualization

source:https://www.unixarena.com/2017/12/para-virtualization-full-virtualization-hardware-assisted-virtualization.html/



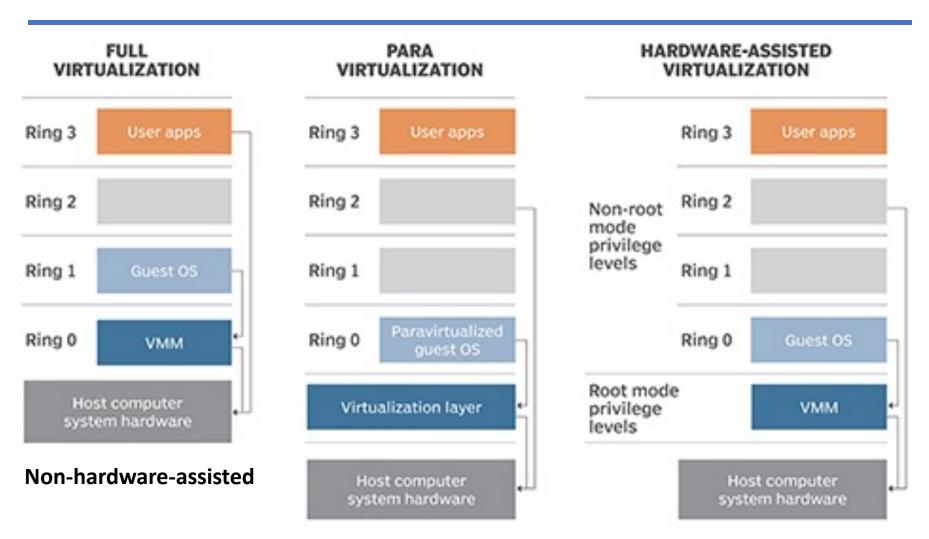
Paravirtualization (cont.)

- Open-source operating systems such as Linux can be easily modified
 - Their code is publicly available
 - Xen provides full support for their virtualization

Components of the Windows family are generally not supported by

Xen unless hardware-assisted virtualization is available.

System Virtualization Implementation



https://searchservervirtualization.techtarget.com/definition/hardware-assisted-virtualization

