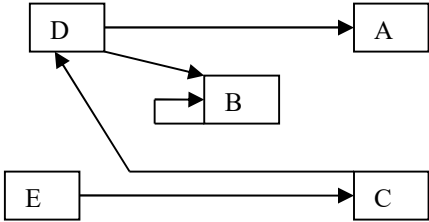
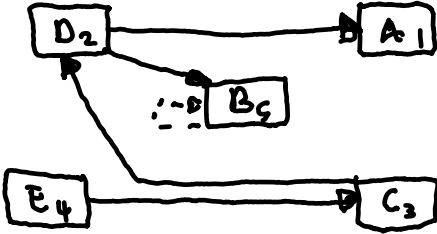


Questions:	Answers:
<p>1. Consider graphs with n nodes and m edges which are stored as adjacency lists. Think of various algorithms to solve the following graph problems. For each, give pseudo code and the worst-case big-Oh running time for the best algorithm you considered.</p> <p>a) For a given a node x, find all nodes in a directed graph that are a distance of two from node x.</p> <p>b) Find all edges with weight 2 in an undirected, weighted graph.</p>	
<p>2. For the graph G below, create the depth first search forest that starts with node A. Whenever there is a choice of which node to visit next, choose the node that comes alphabetically first. Use solid arrows for tree edges; add all other edges as dashed edges. Label all edges as tree, forward, backward, or cross edges. Add postorder numbers as subscripts to each node name.</p> 	<p>1. AD</p> <p>2. ADB</p> <p>3. ADBC</p> <p>4. ADBCE</p> 

<p>3. The appearance of a backward edge in a depth-first-search forest of a directed graph tells us that the graph is cyclic. Graph G in Problem #1 is cyclic. Identify the backward edge in G by giving its tail node x and head node y ($x \rightarrow y$). How do we recognize backward edges by testing postorder numbers? Explain, for the backward edge in G, how the edge $x \rightarrow y$ satisfies this test.</p>	
<p>4. A topological sort of a strict partial ordering R is a total ordering such that x precedes y if xRy.</p> <p>a) Do a DFS on graph G (from #2), choosing nodes alphabetically first, to generate a postorder of the vertices, then reverse the postorder. Your answer should be the reversed postorder. Note that the reverse of the postorder is <i>not</i> always the same as the preorder.</p> <p>b) Give the total ordering produced if we do a DFS choosing nodes alphabetically last.</p> <p>c) Let $R = \{(a, b), (a, c), (b, d), (b, e), (c, e), (d, f)\}$ and let the strict partial ordering be R^+. List all topological sorts of R^+. (You need not use any particular algorithm)</p>	

5. a) Give all depth-first-search forests for G in Problem #2 starting with node E .

b) Expressed in terms of reachability from node E , why do all of the “forests” only have one tree?

6. Let F be a depth-first-search forest for an undirected graph G . (Recall that when we create a depth-first-search forest for an undirected graph, we replace each edge with two directed edges, one in each direction, and then run the ordinary algorithm to create a depth-first-search forest.)

Prove: Two nodes, x and y , are in the same tree of F if and only if x and y are in the same connected component of G .

7. Create the adjacency matrix A for the graph below. Compute A^2 , A^3 , and A^4 and then the union of A , A^2 , A^3 , and A^4 , yielding the reachability matrix for A .

