

Homework #9

Name _____

Sec _____

Questions:	Answers:																
<p>1. Translate the following into predicate calculus. For each answer, also state your assumed universe of discourse.</p> <p>a) “Anyone who was an ancient Roman and tried to kill Caesar was not loyal to Caesar.”</p> <p>b) “All cats which are calico, are female.”</p> <p>c) “Some Texans have never left the state of Texas.”</p>	<p>a) Let $A(x)$ represent "x was an ancient Roman," and $K(x)$ represent "x tried to kill Caesar." Let $L(x)$ represent "x is loyal to Caesar."</p> $\forall x(A(x) \wedge K(x) \rightarrow \neg L(x))$ <p>Assumed universe of discourse: All individuals.</p> <p>b) Let $C(x)$ represent "x is a cat" and $F(x)$ represent "x is female." Let $H(x)$ represent "x is calico."</p> $\forall x((C(x) \wedge H(x)) \rightarrow F(x))$ <p>Assumed universe of discourse: All animals.</p> <p>c) Let $T(x)$ represent "x is a Texan" and $L(x)$ represent "x left the state of Texas."</p> $\exists x(T(x) \wedge \neg L(x))$ <p>Assumed universe of discourse: All individuals.</p>																
<p>2. A universe contains the three individuals a, b, and c. For these individuals, a predicate $Q(x, y)$ is defined, and its truth values are given by the following table:</p> <table><tr><th>$x \backslash y$</th><th>a</th><th>b</th><th>c</th></tr><tr><th>a</th><td>T</td><td>F</td><td>T</td></tr><tr><th>b</th><td>F</td><td>T</td><td>F</td></tr><tr><th>c</th><td>F</td><td>T</td><td>T</td></tr></table> <p>Write each of the following expressions without quantifiers (i.e. convert them to expressions with ANDs and ORs or both) and then evaluate the expression.</p> <p>a) $\forall x \exists y Q(x, y)$</p> <p>b) $\forall y Q(y, b)$</p> <p>c) $\forall y Q(y, y)$</p>	$x \backslash y$	a	b	c	a	T	F	T	b	F	T	F	c	F	T	T	<p>a) $\forall x \exists y Q(x, y)$</p> $Q(a, a) \vee Q(a, b) \vee Q(a, c)$ $Q(b, a) \vee Q(b, b) \vee Q(b, c)$ $Q(c, a) \vee Q(c, b) \vee Q(c, c)$ $Q(a, a) \vee Q(a, b) \vee Q(a, c)$ $T \vee F \vee T = T$ $Q(b, a) \vee Q(b, b) \vee Q(b, c)$ $F \vee T \vee F = T$ $Q(c, a) \vee Q(c, b) \vee Q(c, c)$ $F \vee T \vee T = T$ <p>$\forall x \exists y Q(x, y)$ evaluates to T</p> <p>b) $\forall y Q(y, b)$</p> $Q(a, b) \wedge Q(b, b) \wedge Q(c, b)$ $Q(a, b) \wedge Q(b, b) \wedge Q(c, b)$ $F \wedge T \wedge T = F$ <p>$\forall y Q(y, b)$ evaluates to F</p> <p>c) $\forall y Q(y, y)$</p> $Q(a, a) \wedge Q(b, b) \wedge Q(c, c)$ $Q(a, a) \wedge Q(b, b) \wedge Q(c, c)$ $T \wedge T \wedge T = T$ <p>$\forall y Q(y, y)$ evaluates to T</p>
$x \backslash y$	a	b	c														
a	T	F	T														
b	F	T	F														
c	F	T	T														

<p>5. Using the facts in the class university database in Discussion #15, write predicate logic statements to answer the following questions.</p> <p>a) What are the names of students who live at 12 Apple St.?</p> <p>b) What are the names of students who are getting an A in CS101?</p> <p>Your predicate logic statements must answer these questions for any state of the database, not just the one in the discussion slides.</p>	
<p>6. Algebraically transform:</p> $\neg \forall x(P(x) \wedge Q(y) \Rightarrow \exists zR(z)) \text{ to } \exists x \forall z(P(x) \wedge Q(y) \wedge \neg R(z))$ <p>Justify each step with one or more laws.</p>	<p>De Morgan's Law: $\neg \forall x(P(x) \wedge Q(y) \Rightarrow \exists zR(z))$ becomes $\exists x \neg(P(x) \wedge Q(y) \Rightarrow \exists zR(z))$</p> <p>Implication Law: $\neg(P \Rightarrow Q)$ is equivalent to $Q \wedge \neg P$</p> <p>Distributive Law: $\exists x(P(x) \wedge Q(y) \wedge \neg \exists zR(z))$ is equivalent to $\exists x(P(x) \wedge Q(y) \wedge \forall z \neg R(z))$</p> <p>Putting these together:</p> $\neg \forall x(P(x) \wedge Q(y) \Rightarrow \exists zR(z)) \equiv \exists x \neg(P(x) \wedge Q(y) \Rightarrow \exists zR(z)) \text{ (De Morgan's Law)}$ $\exists x \neg(P(x) \wedge Q(y) \Rightarrow \exists zR(z)) \equiv \exists x(P(x) \wedge Q(y) \wedge \neg \exists zR(z)) \text{ (Implication Law)}$ $\exists x(P(x) \wedge Q(y) \wedge \neg \exists zR(z)) \equiv \exists x(P(x) \wedge Q(y) \wedge \forall z \neg R(z)) \text{ (Distributive Law)}$ <p>$\neg \forall x(P(x) \wedge Q(y) \Rightarrow \exists zR(z))$ algebraically transforms to $\exists x \forall z(P(x) \wedge Q(y) \wedge \neg R(z))$</p>
<p>7. Consider the following expression:</p> $\forall x \exists y(P(y, x) \wedge \exists x Q(y, x)) \vee R(x) \wedge \exists y R(y).$ <p>a) Identify the subexpression in the scope of $\forall x$.</p> <p>b) Identify the free variables.</p> <p>c) Identify each bound variable and the quantifier to which it is bound.</p> <p>d) Rectify the expression. (“Rectification is also called “standardizing variables apart.”)</p>	