Floyd-Warshall

Also Bernard Roy

Floyd's allpairspath

https://open.kattis.com/problems/allpairspath

Also big truck

https://open.kattis.com/problems/bigtruck

And import spagetti

https://open.kattis.com/problems/importspaghetti

And some ICPC problem in 2008

https://www.urionlinejudge.com.br/repository/UOJ_2130_en.html

All Pairs Shortest Path (Kattis)



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Different algorithm

Floyd's

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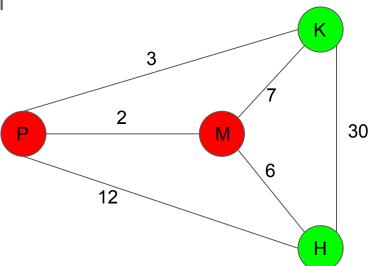
Traveling to the kindergarten was long to avoid some shady locations.

Suppose this is the graph

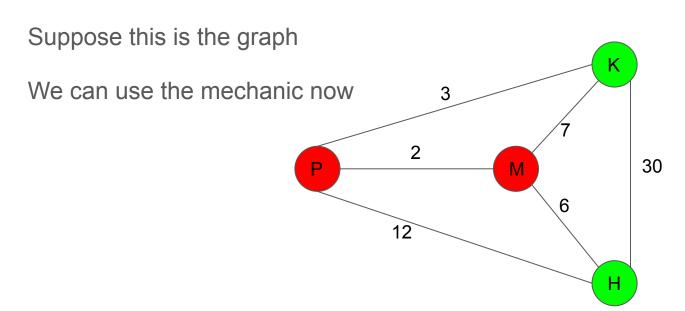
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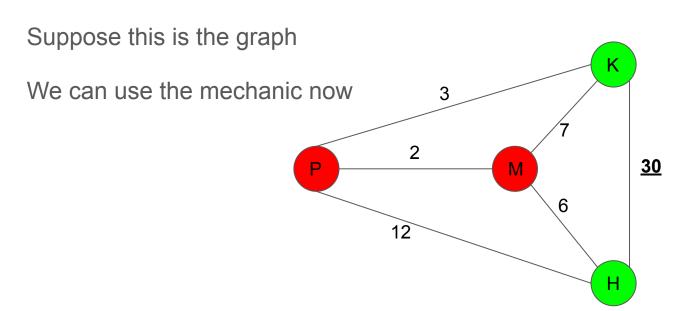
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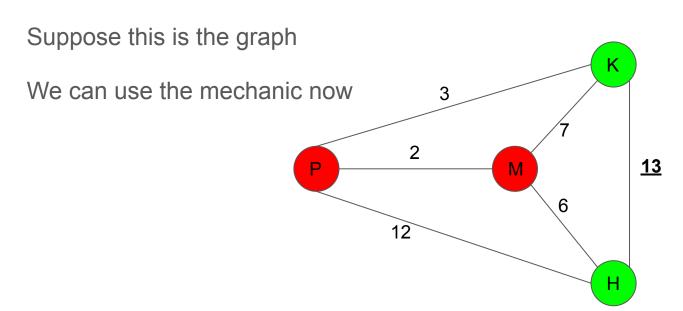
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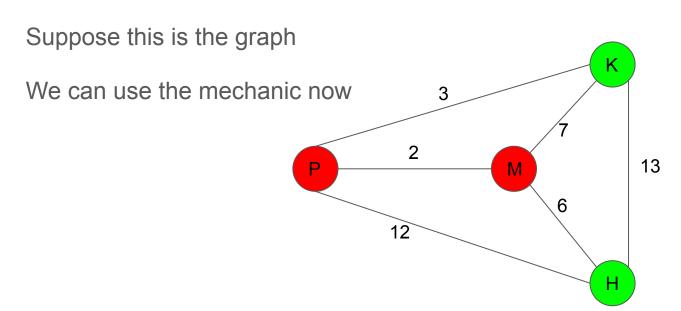
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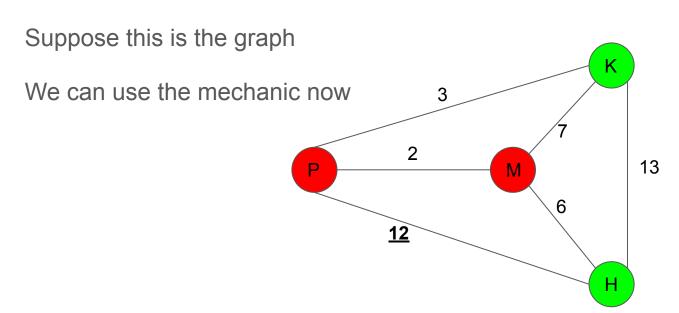
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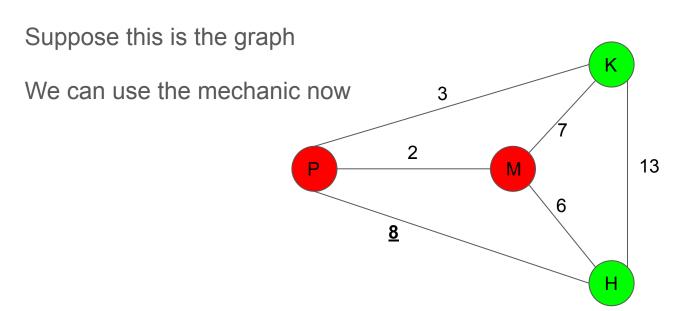
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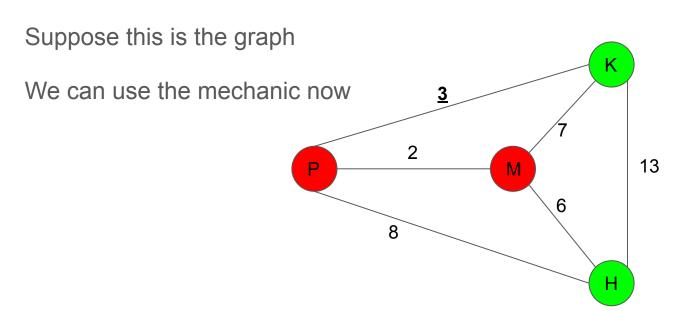
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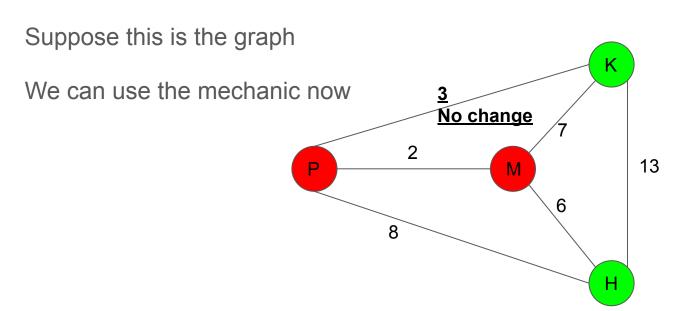
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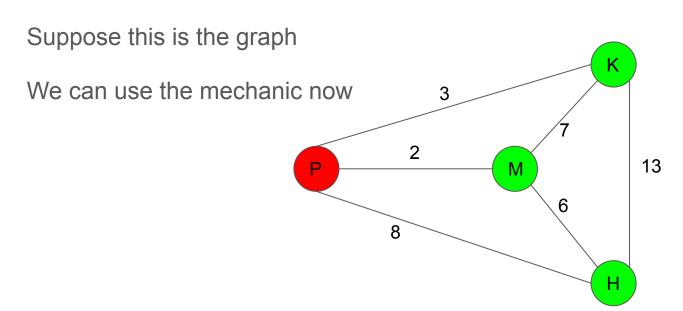
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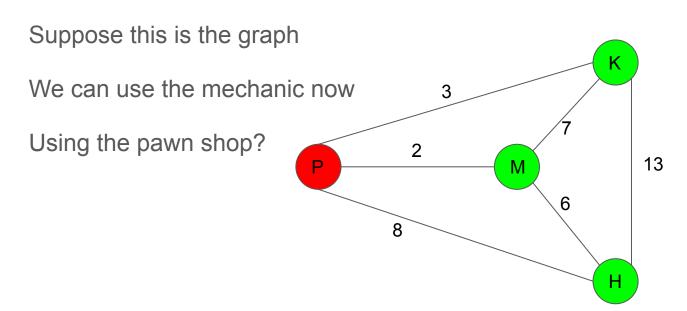
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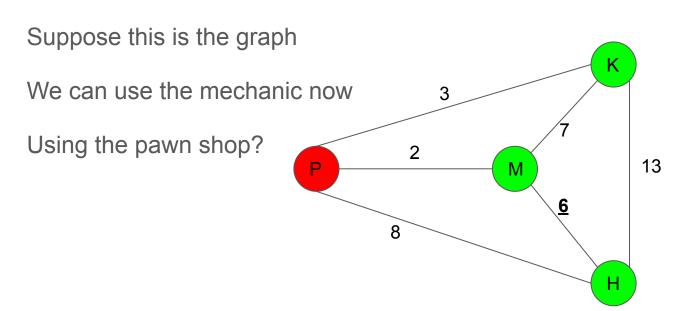
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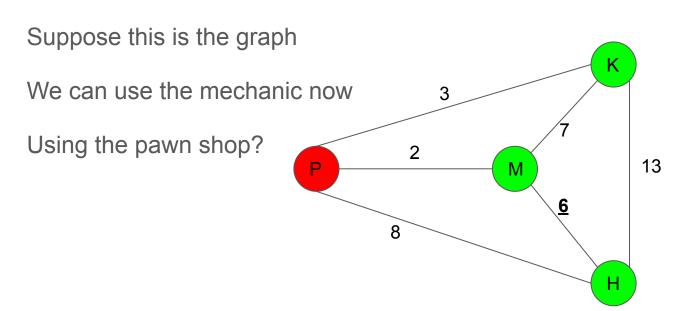
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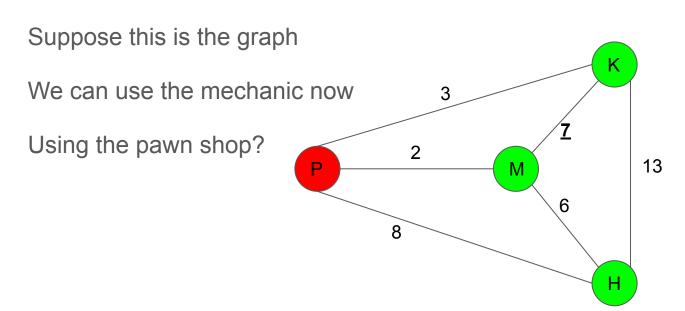
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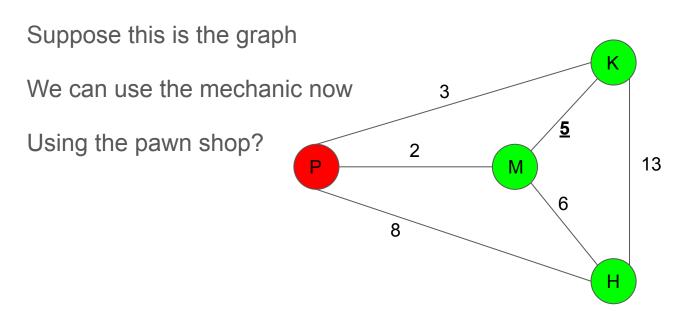
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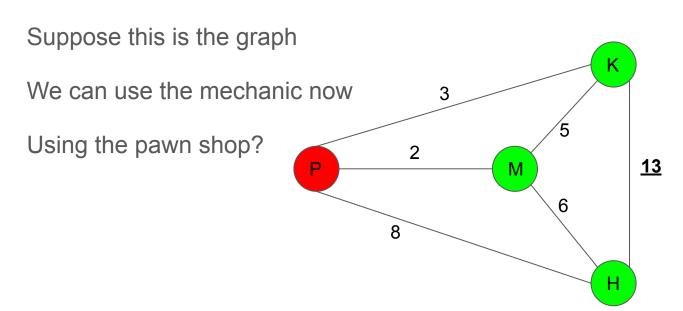
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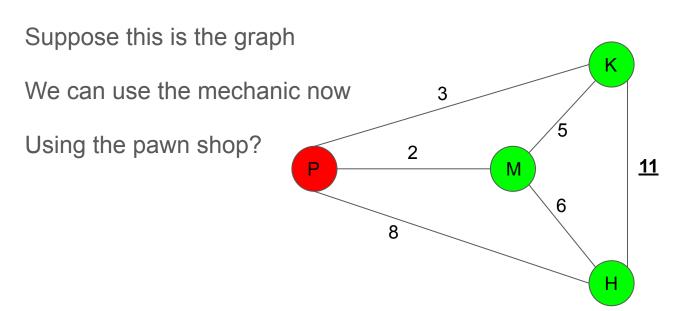
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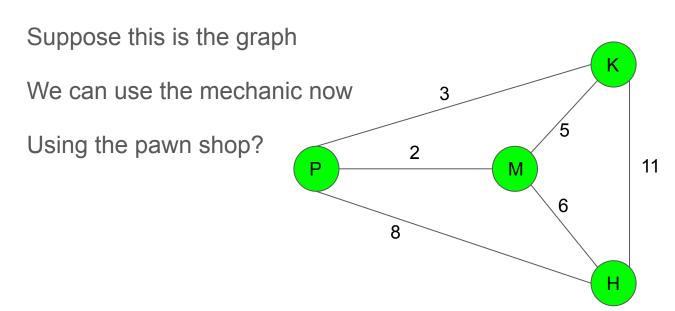
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Floyd's Implementation

The high level

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Look a every node as an intermediate point and adjust all pairs of points

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```
for (iter = 0; iter < n; iter++)
for (st = 0; st < n; st++)
  for (en = 0; en < n; en++)
     if (path from st to en is worse than
          path from st to iter combined with
          Path from iter to en)
     Update the path from st to en</pre>
```

Floyd's Usefulness

Can be used to find ALL pair of shortest paths on graphs with negative edges

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Can be used (in one pass) to find some nodes on negative cycles which can extend to finding all nodes that can

- Don't run it multiple times
- Check nodes that are on a negative cycle
- Check nodes that can reach and be reached by these negative cycles

Runtime

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 $O(N^3)$

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Don't use over 500 Nodes in competitive programming contests.

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Up to 300 Nodes is safe.

Other Applications

Can also be used to find the transitive closure of a graph (SCCs).

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Most people will use Tarjan's Low Link method or Kosaraju's for SCC.