

A Chain Rule for the Randomness Deficiency Function

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Abstract

This paper applies the addition equality theorem for algorithmic entropy in [G01] to the Cantor space and generalizes it to arbitrary joint distributions. This application implies that randomness deficiency of infinite sequences obeys the chain rule, analogous to the finite Kolmogorov complexity case. This is a generalization of van Lambalgen's Theorem. In addition, a dual integration trick shortens the proof of the original theorem.

1 Introduction

Prefix free Kolmogorov complexity, \mathbf{K} , obeys the chain rule, with for $x, y \in \{0, 1\}^*$,

$$\mathbf{K}(x, y) =^+ \mathbf{K}(x) + \mathbf{K}(y|x, \mathbf{K}(x)).$$

In this paper, we apply the addition equality theorem for algorithmic entropy in [G01] to the specific case of infinite sequences and generalize to arbitrary joint distributions. We also shorten the original proof using an integration trick. The consequence to this is a chain rule for the randomness deficiency function \mathbf{D} for computable probabilities μ over infinite sequences, $\mathbf{D}(\alpha|\mu, x) = \sup_n -\log \mu(\alpha[0..n] - \mathbf{K}(\alpha[0..n]|x))$. A measure $\mu\nu$ over the space $\{0, 1\}^\infty \times \{0, 1\}^\infty$ is computable if the marginal μ is computable and the conditional $\nu_x = \nu(\cdot|x)$ is computable. Conditional computability means $\nu(\beta[0..n]|\alpha[0..n])$ is computable. The randomness deficiency over the space $\{0, 1\}^\infty \times \{0, 1\}^\infty$, is $\mathbf{D}(\alpha, \beta|\mu\nu) = \sup_n -\log \mu\nu(\alpha[0..n]\beta[0..n]) - \mathbf{K}(\alpha[0..n]\beta[0..n])$. The discrete case for $\mathbf{d}(x|p) = -\log p(x) - \mathbf{K}(x)$ is trivial. The result detailed in this paper is as follows.

Theorem. *Relativized to probability $\mu\nu$ over $\{0, 1\}^\infty \times \{0, 1\}^\infty$,*

$$\mathbf{D}(\alpha, \beta|\mu\nu) =^+ \mathbf{D}(\alpha|\mu) + \mathbf{D}(\beta|\nu_\alpha, (\alpha, \lceil \mathbf{D}(\alpha|\mu) \rceil)).$$

This is a generalization of van Lambalgen's Theorem, which states (α, β) is ML random iff α is ML random and β is ML random with respect to α . An open question is whether \mathbf{D} follows the linear inequalities that parallel Shannon entropy \mathcal{H} , as Kolmogorov complexity was shown to do [HRSV00]:

Conjecture. *Given $\{\alpha_1 \dots \alpha_n\} \in \{0, 1\}^{\infty n}$, and random variables $\{\beta_1, \dots, \beta_n\}$ is it the case that over all $W \subseteq \{1, \dots, n\}$, and $\lambda_W \in \mathbb{R}$,*

$$\sum_W \lambda_W \mathbf{D}(\alpha^W) \leq 0 \Leftrightarrow \sum_W \lambda_W \mathcal{H}(\beta^W) \geq 0?$$

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2 Results

For the nonnegative real function f , we use $<^+ f$, $>^+ f$, and $=^+ f$ to denote $< f + O(1)$, $> f - O(1)$, and $= f \pm O(1)$. The Kolmogorov complexity of a lower computable function f is $\mathbf{K}(f)$, the size of the shortest program that enumerates it. As shown in [G01], $2^{\mathbf{D}(\alpha|\mu)} \stackrel{*}{=} \mathbf{t}_\mu(\alpha)$ where \mathbf{t}_μ is a universal lower computable μ -test. Furthermore, a modification of the proof Theorem 2.3.4 in [G01] to the $\{0, 1\}^\infty \times \{0, 1\}^\infty$ space can be used to show that $2^{\mathbf{D}(\alpha, \beta|\mu\nu)} \stackrel{*}{=} \mathbf{t}_{\mu\nu}(\alpha, \beta)$, where $\mathbf{t}_{\mu\nu}$ is a universal lower computable test over $\{0, 1\}^\infty \times \{0, 1\}^\infty$. For measure μ and lower continuous function f over $\{0, 1\}^\infty$, we use the notation $\mu^x f(x) = \int_{x \in \{0, 1\}^\infty} f(x) d\mu(x)$. Throughout this section, the universal Turing machine is assumed to be relativized to probabilities μ and ν_x over $\{0, 1\}^\infty$. This means that there is an $O(1)$ sized program that can compute $\mu(x\{0, 1\}^\infty)$ uniformly in $x \in \{0, 1\}^*$, and $\nu(y\{0, 1\}^\infty|x)$, uniformly in $x, y \in \{0, 1\}^n$.

Proposition 1 $\log \nu_x^y 2^{\mathbf{D}(x, y|\mu\nu)} <^+ \mathbf{D}(x|\mu)$.

Proof. Let $f(x) = \log \nu_x^y 2^{\mathbf{D}(x, y|\mu\nu)}$. The function f is lower computable and has $\mu^x 2^{f(x)} \leq 1$. The proposition follows from the universal properties of \mathbf{t}_μ , where $2^f \stackrel{*}{<} \mathbf{t}_\mu$. \square

Proposition 2 If $i < j$, then

$$i + \mathbf{D}(x|\mu, j) <^+ j + \mathbf{D}(x|\mu, i).$$

Proof. By the properties of \mathbf{D} , we have

$$\mathbf{D}(x|\mu, j) <^+ \mathbf{D}(x|\mu, i) + \mathbf{K}(j - i) <^+ \mathbf{D}(x|\mu, i) + j - i.$$

\square

Definition 1 Let $F : \{0, 1\}^\infty \rightarrow \mathbb{Z} \cup \{-\infty, \infty\}$ be a lower semicomputable function. An (ν, F) -test is a function $t : \{0, 1\}^\infty \times \{0, 1\}^\infty \rightarrow \mathbb{Z} \cup \{-\infty, \infty\}$ that is lower semicomputable and $\nu_y^x t(x, y) \leq 2^{F(y)}$. There exists a maximal (ν, F) test, $\mathbf{t}_{(\nu, F)}$, such that $t \stackrel{*}{<} \mathbf{t}_{(\nu, F)}$.

Proposition 3 Let $F : \{0, 1\}^\infty \rightarrow \mathbb{Z} \cup \{-\infty, \infty\}$ be an upper semicomputable function of Kolmogorov complexity $O(1)$. For all x and with $\mathbf{t}_{(\nu, F)}(y) > -\infty$,

$$\mathbf{t}_{(\nu, F)}(x, y) \stackrel{*}{=} 2^{F(y)} \mathbf{t}_{\nu_y}(x|y, -F(y)).$$

Proof. To prove the inequality $\stackrel{*}{>}$, let $g(x, y, m) = \max_{i \geq m} 2^{-i} \mathbf{t}_{\nu_y}(x|y, i)$. This function is lower computable, and decreasing in m . The function $g(x, y) = g_{\nu_y}(x, y, -F(y))$ is lower semicomputable since $-F$ is upper semi-computable. The multiplicative form of Proposition 2 implies

$$\begin{aligned} g(x, y, m) &\stackrel{*}{=} 2^{-m} \mathbf{t}_{\nu_y}(x|y, m) \\ g(x, y) &\stackrel{*}{=} 2^{F(y)} \mathbf{t}_{\nu_y}(x|y, -F(y)). \end{aligned}$$

Since \mathbf{t}_{ν_y} is a test:

$$\begin{aligned} \nu_y^x 2^{-m} \mathbf{t}_{\nu_y}(x|y, m) &\leq 2^{-m} \\ \nu_y^x g(x, y) &\stackrel{*}{<} 2^{F(y)}, \end{aligned}$$

which implies $g(x, y) \stackrel{*}{<} \mathbf{t}_{(\nu, F)}(x, y)$ by the optimality of $\mathbf{t}_{(\nu, F)}$. We now consider the upper bound. Since, given fixed y , $2^{-F(y)}\mathbf{t}_{(\nu, F)}(x, y)$ is an x -test conditional on y and $-F(y)$, we have

$$2^{-F(y)}\mathbf{t}_{(\nu, F)}(x, y) \stackrel{*}{<} \mathbf{t}_{\nu_y}(x|y, -F(y))/\mathbf{m}(\mathbf{t}_{(\nu, F)}) \stackrel{*}{<} \mathbf{t}_{\nu_y}(x|y, -F(y))2^{\mathbf{K}(F)} \stackrel{*}{<} \mathbf{t}_{\nu_y}(x|y, -F(y)).$$

□

Theorem 1 *Relativized to probabilities μ and ν over $\{0, 1\}^\infty$,*

$$\mathbf{D}(x, y|\mu\nu) = {}^+ \mathbf{D}(x|\mu) + \mathbf{D}(y|\nu_x, (x, \lceil \mathbf{D}(x|\mu) \rceil)).$$

Proof. Let $f(x, y) = 2^{\mathbf{D}(x, y|\mu\nu)}$. Proposition 1 implies there exists $c \in \mathbb{N}$ with $\nu_x^y f(x, y) \leq 2^{\mathbf{D}(x|\mu)+c}$. Let $F(x) = \lceil \mathbf{D}(x|\mu) \rceil$. Note that if h is a lower computable function such that $\nu_x^y h(x, y) \stackrel{*}{<} 2^{\mathbf{D}(x|\mu)}$, then $\mu^x \nu_x^y h(x, y) \stackrel{*}{<} \mu^x \mathbf{t}_\mu(x) \stackrel{*}{<} 1$, so $h \stackrel{*}{<} f$, so f is a universal F -test. Proposition 3 (swapping x and y and noting that $\mathbf{K}(\lceil \mathbf{D}(\cdot|\mu) \rceil) = O(1)$) gives

$$\begin{aligned} \mathbf{D}(x, y|\mu\nu) &= \log f(x, y) = {}^+ F(x) + \mathbf{D}(y|\nu_x, (x, -F(x))) \\ \mathbf{D}(x, y|\mu\nu) &= {}^+ \mathbf{D}(x|\mu) + \mathbf{D}(y|\nu_x, (x, \lceil \mathbf{D}(x|\mu) \rceil)). \end{aligned}$$

□

References

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- [HRSV00] D. Hammer, A. Romashchenko, A. Shen, and N. Vereshchagin. Inequalities for shannon entropy and kolmogorov complexity. *Journal of Computer and System Sciences*, 60(2):442–464, 2000.