## A Short Proof on the Existence of Anomalies

Samuel Epstein\*

July 14, 2022

## Abstract

The Independence Postulate (IP) is a finitary Church-Turing Thesis, postulating that mathematical sequences are independent from physical ones. IP implies that anomalies are found in the physical world.

## Anomalies

 $\mathbf{K}(x|y)$  is the conditional prefix Kolmogorov complexity. For probability p over  $\mathbb{N}$ , the deficiency of randomness is  $\mathbf{d}(a|p) = \lfloor -\log p(a) \rfloor - \mathbf{K}(a)$ .  $\mathbf{I}(a;\mathcal{H}) = \mathbf{K}(a) - \mathbf{K}(a|\mathcal{H})$ , where  $\mathcal{H}$  is the halting sequence. An elementary probability measure over  $\mathbb{N}$  has finite support and a range in  $\mathbb{Q}$ .  $<^+f$  is < f + O(1) and  $<^{\log}f$  is  $< f + O(\log(f+1))$ . Stochasticity is

$$\Lambda(a|b) = \min\{\mathbf{K}(Q|b) + 3\log\max\{\mathbf{d}(a|Q,b), 1\}$$
  
: Q is an elementary probability measure\.

$$\Lambda(a|b) < \Lambda(a) + O(\log \mathbf{K}(b)).$$

The following definition is from [Lev74].

Definition 1 (Information) 
$$I(\alpha:\beta) = \log \sum_{x,y} 2^{K(x)+K(y)-K(x,y)-K(x|\alpha)-K(y|\beta)}$$
.

The Independence Postulate [Lev84, Lev13] statement is:

**IP**: Let  $\alpha$  be a sequence defined with an n-bit mathematical statement, and a sequence  $\beta$  can be located in the physical world with a k-bit instruction set. Then  $\mathbf{I}(\alpha:\beta) < k+n+c$  for some small absolute constant c.

There are many proofs in the literature that stochastic numbers have high mutual information with the halting sequence. One detailed proof is in [Eps21].

Lemma 1  $\Lambda(x) < \log \mathbf{I}(x; \mathcal{H})$ .

**Lemma 2** For computable probability p over  $\mathbb{N}$  and for  $D \subset \mathbb{N}$ ,  $|D| = 2^s$ ,  $s < \max_{a \in D} \mathbf{d}(a|p) + \Lambda(D) + \mathbf{K}(s) + \mathbf{K}(p) + O(\log \mathbf{K}(s)\mathbf{K}(p))$ .

**Proof.** We relativize the universal Turing machine to p and s. Let Q be an elementary probability measure that realizes  $\Lambda(D)$ . Let  $d = \max\{\mathbf{d}(D|Q), 1\}$ . Let  $F \subseteq \mathbb{N}$  be a random set where each element  $a \in \mathbb{N}$  is selected independently with probability  $cd2^{-s}$ , where  $c \in \mathbb{N}$  is chosen later.  $\mathbf{E}[p(F)] \leq cd2^{-s}$ . Furthermore

$$\mathbf{E}[Q(\{G : |G| = 2^s, G \cap F = \emptyset\})] \le \sum_{G} Q(G)(1 - cd2^{-s})^{2^s} < e^{-cd}.$$

Thus finite  $W \subset \mathbb{N}$  can be chosen such that  $p(W) \leq 2cd2^{-s}$  and  $Q(\{G: |G| = 2^s, G \cap W = \emptyset\}) \leq e^{1-cd}$ .  $D \cap W \neq \emptyset$ , otherwise, using the Q-test,  $t(G) = e^{cd-1}$  if  $(|G| = 2^s, G \cap W = \emptyset)$  and t(G) = 0 otherwise, we have

$$\begin{split} \mathbf{K}(D|Q,d,c) <^+ &- \log Q(D) - (\log e)cd \\ &(\log e)cd <^+ - \log Q(D) - \mathbf{K}(D|Q) + \mathbf{K}(d,c) \\ &(\log e)cd <^+ d + \mathbf{K}(d,c), \end{split}$$

which is a contradiction for large c. Thus there is an  $a \in D \cap W$ , where

$$\mathbf{K}(a) <^{+} -\log p(a) + \log d - s + \mathbf{K}(d) + \mathbf{K}(Q)$$
$$s <^{+} \mathbf{d}(a|p) + \Lambda(D).$$

Making the relativization of p and s explicit,

$$s < -\log p(a) - \mathbf{K}(a|p,s) + \Lambda(D|p,s)$$
  

$$s < \max_{a \in D} \mathbf{d}(a|p) + \Lambda(D) + \mathbf{K}(s) + \mathbf{K}(p)$$
  

$$+ O(\log \mathbf{K}(s)\mathbf{K}(p)).$$

<sup>\*</sup>JP Theory Group. samepst@jptheorygroup.org

For  $\tau \in \mathbb{N}^{\mathbb{N}}$ ,  $\langle \tau \rangle = \langle \tau[1] \rangle \langle \tau[2] \rangle \langle \tau[3] \rangle \dots$  Let  $\tau(n)$  be the first  $2^n$  unique numbers found in  $\tau$ . The sequence  $\tau$  is assumed to have an infinite amount of unique numbers, and represents a series of observations.

**Theorem 1** For probability p over  $\mathbb{N}$  and  $\tau \in \mathbb{N}^{\mathbb{N}}$ , let  $s_{\tau,p}$  be the smallest number such that  $\max_{a \in \tau(n)} \mathbf{d}(a|p) > n - 4\mathbf{K}(n) - s_{\tau,p}$ . Then  $s_{\tau,p} <^{\log} \mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p)$ .

**Proof.** By Lemmas 1 and 2, and the fact that  $I(x; \mathcal{H}) <^+ I(\alpha : \mathcal{H}) + K(x|\alpha)$ ,

$$\begin{split} n &< \max_{a \in \tau(n)} \mathbf{d}(a|p) + \mathbf{I}(\tau(n);\mathcal{H}) + \mathbf{K}(p) + 2\mathbf{K}(n) \\ &+ O(\log(\mathbf{I}(\tau(n);\mathcal{H}) + \mathbf{K}(p))), \\ n &< \max_{a \in \tau(n)} \mathbf{d}(a|p) + 4\mathbf{K}(n) + \mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p) \\ &+ O(\log(\mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p))), \\ n &- 4\mathbf{K}(n) - (\mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p) \\ &+ O(\log(\mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p)))) < \max_{a \in \tau(n)} \mathbf{d}(a|p). \end{split}$$

Let k be the physical address of an infinite sequence of numbers  $\tau \in \mathbb{N}^{\mathbb{N}}$ . The halting sequence can be described by a small mathematical statement. By Theorem 1 and IP,

$$s_{\tau,p} <^{\log} \mathbf{I}(\langle \tau \rangle : \mathcal{H}) + \mathbf{K}(p) <^{\log} k + c + \mathbf{K}(p).$$

Thus sequences  $\tau$  with large  $s_{\tau,p}$  will have large physical addresses. So, assuming IP, it's hard to find observations which do not have large anomalies, and impossible to find observations with no anomalies.

## References

- [Eps21] Samuel Epstein. All sampling methods produce outliers. *IEEE Transactions on Information Theory*, 67(11):7568–7578, 2021.
- [Lev74] L. A. Levin. Laws of Information Conservation (Non-growth) and Aspects of the Foundations of Probability Theory. *Problemy Peredachi Informatsii*, 10(3):206–210, 1974.

- [Lev84] L. A. Levin. Randomness conservation inequalities; information and independence in mathematical theories. *Information and Control*, 61(1):15–37, 1984.
- [Lev13] L. A. Levin. Forbidden information. J. ACM, 60(2), 2013.