

# Three Eggs from the Chicken

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In this note, we show an example of taking a theorem produced from a non-constructive probabilistic proof and produce a three derandomization theorems, one that involves Kolmogorov complexity, one that involves resource bounded Kolmogorov complexity, and one involving games.

Our note deals with hypergraphs. A *hypergraph* is a pair  $J = (V, E)$  of vertices  $V$  and edges  $E \subseteq \mathcal{P}(V)$ . Thus each edge can connect  $\geq 2$  vertices. A hypergraph is  $k$ -regular of the size  $|e| = k$  for all edges  $e \in E$ . A 2-regular hypergraph is just a simple graph. A valid  $C$ -coloring of a hypergraph  $(V, E)$  is a mapping  $f : V \rightarrow \{1, \dots, C\}$  where every edge  $e \in E$  is not *monochromatic*  $|\{f(v) : v \in e\}| > 1$ . The following classic result [EL] is the first proved consequence of Lovász Local Lemma.

**Theorem (Probabilistic Method)** *Let  $J = (V, E)$  be a  $k$ -regular hypergraph. If for each edge  $f$ , there are at most  $2^{k-1}/e - 1$  edges  $h \in E$  such that  $h \cap f \neq \emptyset$ , then there exists a valid 2-coloring of  $J$ .*

We can now use derandomization, from [Eps22a], to produce bounds on the Kolmogorov complexity of the simplest such 2-coloring of  $G$ .

**Theorem 1 (Derandomization)** *Let  $J = (V, E)$  be a  $k$ -regular hypergraph with  $|E| = m$ . If, for each edge  $f$ , there are at most  $2^{k-1}/e - 1$  edges  $h \in E$  such that  $h \cap f \neq \emptyset$ , then there exists a valid 2-coloring  $x$  of  $J$  with*

$$\mathbf{K}(x) <^{\log} \mathbf{K}(n) + 4me/2^k + \mathbf{I}(J; \mathcal{H}).$$

The function  $\mathbf{K}$  is the prefix free Kolmogorov complexity.  $\mathbf{I}(J; \mathcal{H}) = \mathbf{K}(J) - \mathbf{K}(J|\mathcal{H})$  is the amount of asymmetric information the halting sequence  $\mathcal{H} \in \{0, 1\}^\infty$  has about the graph  $J$ . We can now use resource derandomization, from [Eps22b], to achieve bounds for the smallest time-bounded Kolmogorov complexity  $\mathbf{K}^t(x) = \min\{p : U(p) = x \text{ in } t(\|x\|) \text{ steps}\}$  of a 2-coloring of  $J$ .

**Assumption 1 *Crypto*** *is the assumption that there exists a language in  $\mathbf{DTIME}(2^{O(n)})$  that does not have size  $2^{o(n)}$  circuits with  $\Sigma_2^P$  gates.*

**Theorem 2 (Resource Bounded Derandomization)** *Assume **Crypto**. Let  $J_n = (V, E)$  be a  $k(n)$ -regular hypergraph where  $|V| = n$  and  $|E| = m(n)$ , uniformly polynomial time computable in  $n$ . Furthermore, for each edge  $f$  in  $J_n$  there are at most  $2^{k(n)-1}/e - 1$  edges  $h \in E$  such that  $h \cap f \neq \emptyset$ . Then there is a polynomial  $p$ , and a valid 2-coloring  $x$  of  $J_n$  with*

$$\mathbf{K}^p(x) < 4m(n)e/2^{k(n)} + O(\log n).$$

We define the following game involving hypergraphs and is from [Eps23]. The player has access to a list of vertices and the goal of the player is to produce a valid 2-coloring of the hypergraph. We assume that for each edge  $f$  of the graph, there are at most  $2^{k-1}/e - 1$  edges  $h$  such that  $f \cap h \neq \emptyset$ .

The game proceeds as follows. For the first round, environment gives the number of vertices to the player. The player has  $n$  vertices, each with starting color 1. At each subsequent turn, the environment sends to the player the edges which are monochromatic. The player can change the color of up to  $k$  vertices and sends these changes to the environment. The game ends when the player has a valid 2-coloring of the graph.

**Theorem 3 (Game Derandomization)** *For  $k \geq 6$ , there exists a deterministic player  $\mathbf{p}$  that can beat the environment  $\mathbf{q}$  in  $(1 + \epsilon)(n/k)$  turns of complexity  $\mathbf{K}(\mathbf{p}) <^{\log} \mathbf{I}(\mathbf{q}; \mathcal{H}) - \log \epsilon$ .*

## References

- [EL] P. Erdos and L. Lovász. Problems and results on 3-chromatic hypergraphs and some related questions. *Infinite and finite sets*, 10:609–627.
- [Eps22a] S. Epstein. 22 examples of solution compression via derandomization. *CoRR*, abs/2208.11562, 2022.
- [Eps22b] S. Epstein. Derandomization under different resource constraints. *CoRR*, abs/2211.14640, 2022.
- [Eps23] Samuel Epstein. Game derandomization, 2023. <https://www.jptheorygroup.org/doc/Game.pdf>.