

# 'votess: A multi-target, GPU-capable, parallel Voronoi tessellator'

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## Software

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## Statement of need

A Voronoi tessellation is a spatial decomposition that partitions space into a set of convex hulls based on proximity to a seed points with interesting to the applications in biology, data science, geography, and physics. One compute intensive application is its use in astrophysics, such as the analysis of matter distribution (Weygaert, 1994), optimal transport theory for early-universe reconstruction (Levy et al., 2021), and in observational data analysis and numerical simulations of cosmic structure formation (Springel, 2010).

The increasing size of datasets produced today have underscored the need for more efficient algorithms to both generate and analyse these datasets, and the rise of heterogenous computing facilities would enable such new algorithms to be run. There do exist several sequential and parallel implementations of the Voronoi diagram problem (Marot et al., 2019) (Wu et al., 2023) (The CGAL Project, 2018) (Inria, 2018), however, they are mostly restricted to CPU or specific GPU architectures, thus limiting their potential as a portable multi-architecture algorithm.

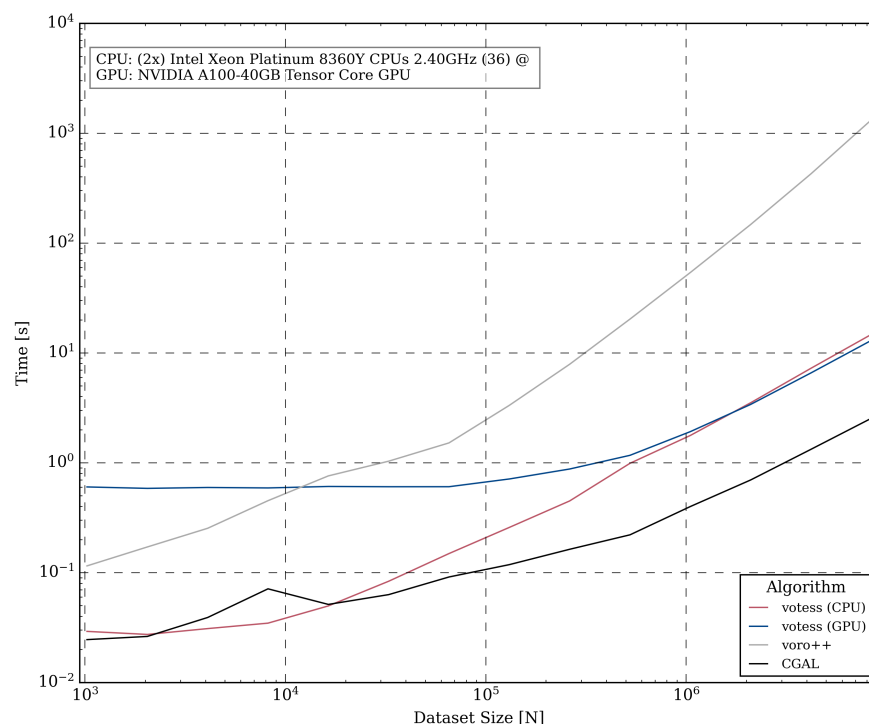
## Summary

votess is a library, implementing Ray's meshless algorithm (Ray et al., 2018), for computing parallel three dimensional Voronoi tessellations using the C++/SYCL framework for CPU, GPUs and other future architectures.

One advantage of this algorithm is the ability for each cell to be computed independently (Ray et al., 2018), making it suitable for parallel execution. It also produces the geometry of the Voronoi cells via their neighbor connectivity information, rather than a full combinatorial mesh data structure, thus making it more ammenable to data parallel architectures than alternatives such as sequential insertion or the Bowyer-Watson algorithm (Bowyer, 1981) (Watson, 1981).

The core method of votess consists of two main steps. First, the input set of points is sorted into a grid, and a k-nearest neighbors search is performed. Once the k nearest neighbors are identified for each point, the Voronoi cell is computed by iteratively clipping a bounding box using the perpendicular bisectors between the point and the identified neighbors. A *security radius* condition (Lévy & Bonneel, 2013) ensures that the resulting Voronoi cell is valid, and if the cell cannot be validated, an CPU fallback mechanism is used.

## 35 Performance



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37 In Figure 1, we benchmark votess against two well-established libraries: CGAL, which builds  
38 parallel CPU-based Delaunay meshes (The CGAL Project, 2018), and Voro++, which performs  
39 single-core, cell-based 3D tessellations (Rycroft, 2009). All tests use a float32 white-noise  
40 dataset, which provide a worst case scenario for uniform datasets (Ray et al., 2018). The use  
41 of clustered datasets lies outside of both votess and the reference implementation (Ray et  
42 al., 2018), but remains in mind for future iterations of votess. Other Multithreaded Voronoi  
43 tessellation codes exist, including ParVoro++ (Wu et al., 2023), and GEOGRAM (Inria, 2018).  
44 However, they do not natively support GPU architectures, but are acknowledged as established  
45 alternatives.

46 From the graph above, votess outperforms Voro++, however, when compared CGAL, both the  
47 CPU and GPU version falls short by around a factor of 6. Ongoing optimizations for both  
48 backends are underway to close this gap.

## 49 Features

50 votess is designed to be versatile. It supports various outputs, including the natural neighbor  
51 information for each Voronoi cell. This is a 2D jagged array of neighbor indices of the sorted  
52 input dataset.

53 Users can invoke votess in three ways: through the C++ library, a command-line interface  
54 clvotess, and a Python wrapper interface pyvotess. The C++ library offers a simple  
55 interface with a tessellate function that computes the mesh. The Python wrapper, mirrors  
56 the functionality of the C++ version, with native numpy array support, providing ease of use  
57 for Python-based workflows.

58 The behavior of votess can be fine-tuned with run time parameters in order to (optionally)

59 optimize runtime performance.

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