# CVSD Final Project

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# 1 Algorithm

### 1.1 QR Decomposition Algorithm Introduction

There are many algorithms for QR decomposition, such as Householder transformation, Givens rotation, and modified Gram-Schmidt. Since modified Gram-Schmidt algorithm computes the results faster and it is more intuitive for implementation compared to other algorithms, we adopt modified Gram-Schmidt to implement the hardware for QR decomposition in this final project.

Modified Gram-Schmidt is an iterative algorithm that starts with the original matrix H and produces the orthogonal matrix Q and the upper triangular matrix R. It works by iteratively orthogonalizing the columns of H to create Q while simultaneously updating R. In each iteration, Modified Gram-Schmidt computes the orthogonal projection of the current column onto the previously processed columns, ensuring that Q remains orthogonal. This process continues until all columns have been orthogonalized, resulting in the QR decomposition of the original matrix H.

$$H = [h_1|h_2|h_3|h_4] = [e_1|e_2|e_3|e_4] \begin{bmatrix} h_1 \cdot e_1 & h_2 \cdot e_1 & h_3 \cdot e_1 & h_4 \cdot e_1 \\ 0 & h_2 \cdot e_2 & h_3 \cdot e_2 & h_4 \cdot e_2 \\ 0 & 0 & h_3 \cdot e_3 & h_4 \cdot e_3 \\ 0 & 0 & 0 & h_4 \cdot e_4 \end{bmatrix} = QR$$
(1)

# 2 Hardware Implementation

### 2.1 Hardware Scheduling

In our design, the result of one RE data will be calculated after twenty six cycles, and the calculation will be started in the  $17^{th}$  cycle (counter equals 16) of this RE. Figure(1) shows the flow for the  $1^{st}$  RE calculation and the details of each cycle. For the  $1^{st}$  RE, we start our calculations at the  $17^{th}$  cycle of the  $1^{st}$  RE input. After twenty six cycles of calculations, the result can be obtained at the  $3^{rd}$  cycle (counter equals 2) of the  $3^{rd}$  RE input. Since our calculation process takes more than twenty cycles, we parallel part of our design to allow all REs to start the process at the  $17^{th}$  cycle of each RE. Please note that we only show the  $1^{st}$  RE calculation in Figure(1). In our real design, the process for  $2^{nd}$  RE will be started at the  $17^{th}$  cycle of the  $2^{nd}$  RE while the process for the  $1^{st}$  RE have not been done.

The notations in Figure (1) are referred to the final introduction. The \* superscript means the multiplications needed for the intermediate value of the certain notation. Since the modified Gram-Schmidt needs many multiplications and the frequency requirement is above 200MHz, we have to pipeline the hardware into many stages. According to Figure (1), we can observe that we only do one stage calculations in most cycles in order to reach the higher frequency more than 200MHz. Because the multiplication delay is strict for high frequency design, we separate the multiplication from addition and subtraction to reduce the latency of each timing path.

input	RE1				RE2								
counter	16	17	18	19	0	1	2	3	4	5	6	7	8
	$R_{11}^{*}$	R <sub>11</sub>	$e_1$	$R_{12}^{*}$	R <sub>12</sub>	R <sub>13</sub>	R <sub>14</sub>	$R_{14}e_{1}$	h <sub>4</sub> <sup>(1)</sup>	$e_2$	R <sub>23</sub> *	R <sub>23</sub>	R <sub>24</sub>
					R <sub>13</sub> *	R <sub>14</sub> *	R <sub>13</sub> e <sub>1</sub>	h <sub>3</sub> <sup>(1)</sup>	R <sub>22</sub>	y_hat <sub>1</sub>	y_hat2*	y_hat2	R <sub>23</sub> e <sub>2</sub>
						$R_{12}e_1$	h <sub>2</sub> <sup>(1)</sup>	R <sub>22</sub> *	y_hat <sub>1</sub> *			R <sub>24</sub> *	
input	RE2											RE3	
counter	9	10	11	12	13	14	15	16	17	18	19	0	1
	R <sub>24</sub> e <sub>2</sub>	h <sub>4</sub> <sup>(2)</sup>	R <sub>33</sub>	$e_3$	R <sub>34</sub> *	R <sub>34</sub>	R <sub>34</sub> e <sub>3</sub>	h <sub>4</sub> <sup>(3)</sup>	R <sub>44</sub> *	R <sub>44</sub>	e <sub>4</sub>	y_hat <sub>4</sub> *	y_hat₄
	h <sub>3</sub> <sup>(2)</sup>	R <sub>33</sub> *			y_hat3*	y_hat <sub>3</sub>							

Figure 1: Flow Chart

## 2.2 Hardware Block Diagram

Figure (2) shows the block diagram of our design. In our design, we don't use SRAM since we want to fetch the data faster. There are other three modules designed by ourselves to simplify the code and improve the readability. There is a module functioning as a lookup table for square root value and reciprocal value. The other two modules functioning as complex multiplier with different fix-point format.

# QR\_Engine.v i\_clk o\_rd\_vld o\_rd\_vld o\_last\_data i\_trig i\_data complex\_mul\_type1 o\_y\_hat complex\_mul\_type2

Testbed.v

Figure 2: Block Diagram

# 2.3 Fixed Point Setting

The input data is in s1,22 fix point format, and the output data should be in s3,16 format. Since the precision for intermediate calculation directly affects the register needed and the output precision, selecting number of bits for each intermediate step should be considered carefully. Since the LUT for square root and reciprocal will be quite large if we keep all precision, we only consider ten bits in this LUT. According to the result by means of implementing the LUT, we can get about 1.23 performance gain while there are still margins for 10SNR and 15SNR cases. Thus, we reduce the width of some registers to minimize the area while maintain the same performance gain. In the end, we use about two thousand registers.

## 2.4 Area / Power / Latency Reports

```
*************
Report : area
Design : QR_Engine
Version: U-2022.12
Date : Sun Dec 17 00:37:51 2023
   slow (File: /home/raid7_2/course/cvsd/CBDK_IC_Contest/CIC/SynopsysDC/db/slow.db)
Number of ports:
Number of nets:
Number of cells:
Number of combinational cells:
Number of sequential cells:
Number of macros/black boxes:
Number of buf/inv:
Number of references:
Combinational area:
                             466793.486211
Buf/Inv area:
                            45848.471327
Noncombinational area:
                             80139.344679
Macro/Black Box area:
                                 0.000000
                            3304949.922974
Net Interconnect area:
Total cell area:
                             546932.830889
Total area:
                            3851882.753863
                      (a) Synthesis Area Report
           : 607217.50
: 575885
              Die Area(um^2)
                                    575885.39
              Chip Density (Counting Std Cells and MACROs and IOs): 94.840%
Core Density (Counting Std Cells and MACROs): 100.000%
                                  : 100.000%
              Average utilization
              Number of instance(s)
              Number of Macro(s)
              Number of IO Pin(s)
                                  : 533
           Number of Power Domain(s): 0
```

(b) APR Area Report

Figure 3: Area Report

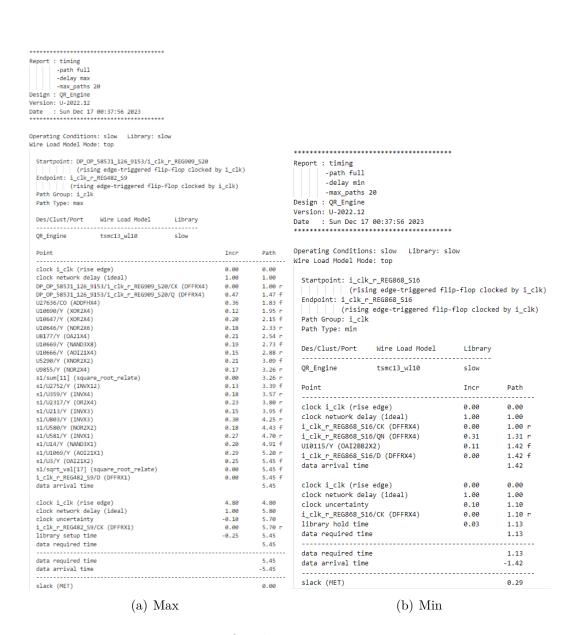


Figure 4: Synthesis Timing Report

```
report_power
Report : Time Based Power
Design : QR_Engine
Version: T-2022.03
Date : Mon Dec 18 09:42:37 2023
     i - Including register clock pin internal power
       u - User defined power group
        Internal Switching Leakage
                                                                     Total
Power Group Power Power Power ( %) Attrs

        clock_network
        2.814e-03 5.159e-04 8.190e-06 3.338e-03 (12.87%)

        register
        3.502e-03 1.290e-04 6.391e-05 3.695e-03 (14.24%) i

        combinational
        0.0116 6.956e-03 3.234e-04 0.0189 (72.89%)

        sequential
        0.0000 0.0000 0.0000 0.0000 0.0000 (0.00%)

sequential
                               0.0000 0.0000 0.0000 0.0000 (0.00%)
0.0000 0.0000 0.0000 0.0000 (0.00%)
0.0000 0.0000 0.0000 0.0000 (0.00%)
memory
io_pad
black_box
  Net Switching Power = 7.601e-03 (29.30%)
  Cell Internal Power = 0.0179 (69.17%)
Cell Leakage Power = 3.955e-04 (1.52%)
 Total Power
                            = 0.0259 (100.00%)
X Transition Power
                            = 3.595e-05
Glitching Power
                            = 1.169e-04
Peak Power
                            = 0.5933
Peak Time
                            = 524.761
```

Figure 5: Power Report

```
Pattern 1# ~ 100# are written
Pattern 101# ~ 200# are written
Pattern 201# ~ 300# are written
Pattern 301# ~ 400# are written
Pattern 301# ~ 400# are written
Pattern 501# ~ 500# are written
Pattern 601# ~ 700# are written
Pattern 601# ~ 800# are written
Pattern 701# ~ 800# are written
Pattern 801# ~ 900# are written
Pattern 901# ~ 1000# are written
Pattern 901#
```

Figure 6: Latency

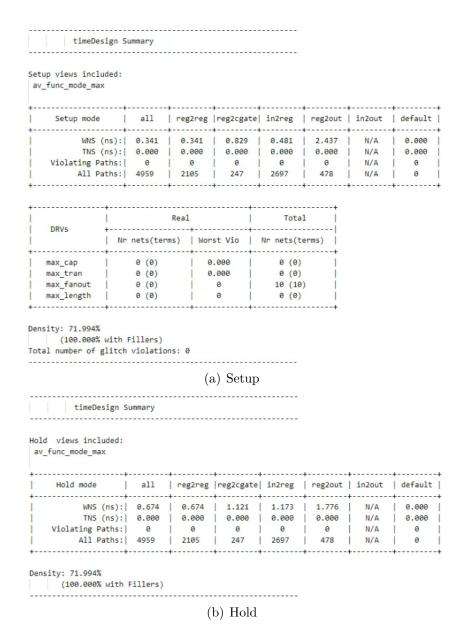


Figure 7: Route Timing Report