# CSCE 689: Special Topics in Modern Algorithms for Data Science

Lecture 2

Samson Zhou

#### Last Time: Class Logistics

• Course materials: https://samsonzhou.github.io/csce689-2023

- LaTeX summary of lectures 20%
- Midterm presentation 35%
- Final project 45%

### Last Time: Probability Basics

• Conditional distribution:  $\Pr[X = x | Y = y]$  is the probability that X achieves the value x when Y achieves the value y

$$\Pr[X = x | Y = y] = \frac{\Pr[X = x, Y = y]}{\Pr[Y = y]}$$

• Implies Bayes' theorem

• Random variables X and Y are independent if  $\Pr[X = x] = \Pr[X = x | Y = y]$  for all possible outcomes  $x \in \Omega_X$ ,  $y \in \Omega_Y$ 

#### Warm-Up Question

- Suppose  $S_1$  is a "bad" event that occurs with probability  $\frac{0}{n}$
- Suppose  $S_2$  is a "bad" event that occurs with probability  $\frac{1}{n}$
- Suppose  $S_3$  is a "bad" event that occurs with probability  $\frac{2}{n}$

What is the probability that none of the bad events occurs?

#### Warm-Up Question

- Suppose  $S_1$  is a "bad" event that occurs with probability  $\frac{0}{n}$
- Suppose  $S_2$  is a "bad" event that occurs with probability  $\frac{1}{n}$
- Suppose  $S_3$  is a "bad" event that occurs with probability  $\frac{2}{n}$

 What is a lower bound on the probability that none of the bad events occur?

#### Warm-Up Question

- Suppose  $S_1$  is a "bad" event that occurs with probability  $\frac{0}{n}$
- Suppose  $S_2$  is a "bad" event that occurs with probability  $\frac{1}{n}$
- Suppose  $S_3$  is a "bad" event that occurs with probability  $\frac{2}{n}$

• What is *a lower bound* on the probability that none of the bad events occur?  $1 - \frac{3}{n}$ 

# Last Time: Union Bound (Boole's Inequality)

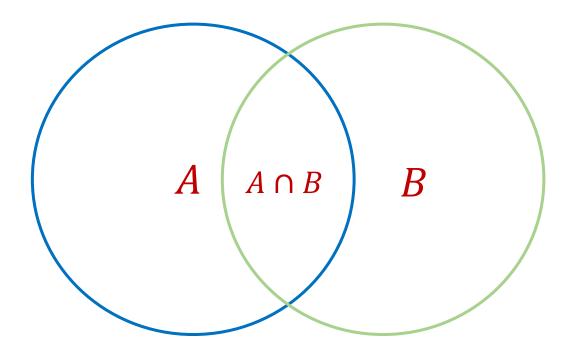
• Let  $S_1,...,S_k$  be a set of events that occur with probability  $p_1,...,p_k$ 

• The probability that at least one of the events  $S_1,...,S_k$  occurs is at most  $p_1+\cdots+p_k$ 

• Implication: the probability that NONE of the events  $S_1,...,S_k$  occur is at least  $1-(p_1+\cdots+p_k)$ 

#### Last Time: Union Bound

•  $Pr[A \cup B] = Pr[A] + Pr[B] - Pr[A \cap B]$ 



Proof by induction

# Today

- Hashing
- Abstraction: balls-in-bins
- Birthday paradox

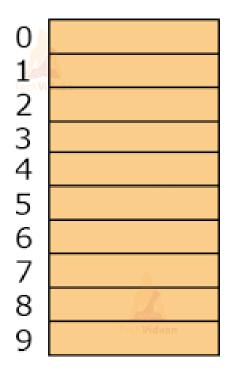
### Trivia Question #1 (Birthday Paradox)

• Suppose we have a fair n-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls? Example: 1, 5, 2, 4, 5

- $\Theta(1)$
- $\Theta(\log n)$
- $\Theta(\sqrt{n})$
- $\Theta(n)$

## Trivia Question #2 (Limits)

- Let c > 0 be a constant. What is  $\lim_{n \to \infty} \left(1 \frac{c}{n}\right)^n$ ?
- 0
- $\cdot \frac{1}{c}$
- $\bullet$   $\frac{1}{2c}$
- $\frac{1}{e^c}$
- 1





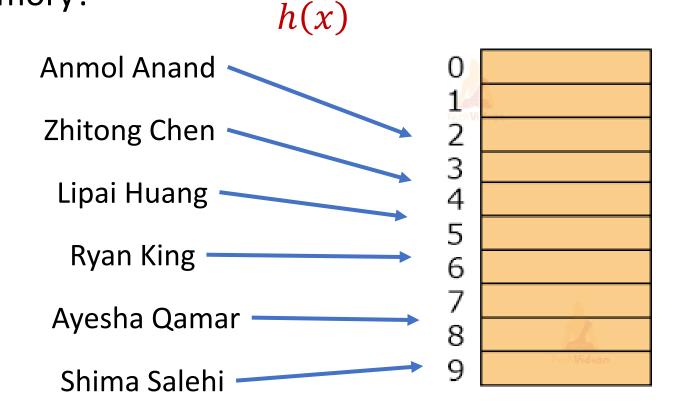
0	Anmol Anand
1	Zhitong Chen
2	Lipai Huang
3	Ryan King
4	Ayesha Qamar
5	Shima Salehi
6	
7	
8	
9	Fer NVidvan





	Anmol Anand
4	Zhitong Chen
	Lipai Huang
	Ryan King
4	Ayesha Qamar
	Shima Salehi
	J.K.
	Fer Nidvan





#### Hash Tables

• We have a set of m items from some large universe that we want to store into a database (images, text documents, IP addresses) with n locations

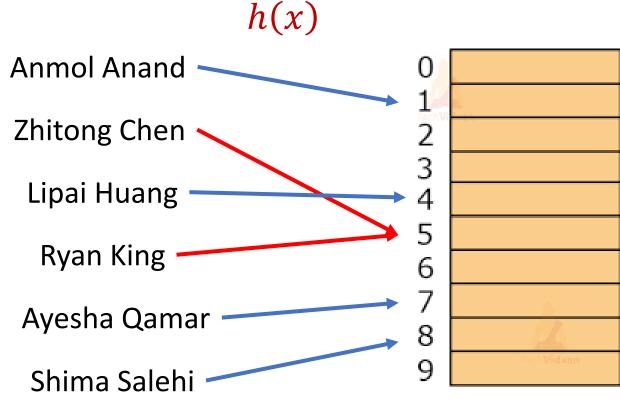
• Goal: query(x) to check if the database contains x in O(1) time

• Hash function  $h: U \to [n]$  maps items from the universe to a location in the database

#### Collisions

• Hash function  $h: U \to [n]$  maps items from the universe to a location in the database

- For  $|U| \gg n$ , many items map to the same location
- Collision: when multiple items should be stored in the same location



#### Dealing with Collisions

- Many ways of dealing with collisions
  - Store multiple items in the same location as a linked list
  - Bump item to the next available spot
  - Bump item to the next available spot using another hash function
  - Power-of-two-choices

### Dealing with Collisions

• Suppose we store multiple items in the same location as a linked list



• If the maximum number of collisions in a location is *c*, then could traverse a linked list of size *c* for a query

• Query runtime: O(c)

## Dealing with Collisions

• Goal: minimize c, the maximum number of collisions in a location

• In the worst case, all items could hash to the same location, c=m

Assume the hash function h is chosen "randomly"

#### Random Hash Function

• Let  $h: U \to [n]$  be a random hash function, so that for each  $x \in U$ , we have that  $\Pr[h(x) = i] = \frac{1}{n}$ , for all  $i \in [n]$ 

• Assume independence, i.e., h(x) and h(y) are independent for any  $x, y \in U$ 

• Suppose we insert m elements into a hash table with n locations using a random hash function. How do we analyze the number of pairwise collisions?

• Suppose we have a room with 367 people. What is the probability that two people share the same birthday?

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• Suppose we have a room with 23 people. What is the probability that two people share the same birthday?

$$\left(1-\frac{0}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)$$

$$\left(1 - \frac{0}{n}\right) \left(1 - \frac{1}{n}\right) \left(1 - \frac{2}{n}\right)$$

$$\left(1-\frac{0}{n}\right)\left(1-\frac{1}{n}\right)...\left(1-\frac{k-1}{n}\right)$$

$$\left(1 - \frac{0}{n}\right) \left(1 - \frac{1}{n}\right) \dots \left(1 - \frac{k-1}{n}\right) < \frac{1}{2}$$
 for  $k = O(\sqrt{n})$ 

• Suppose we have a fair n-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls?

•  $O(\sqrt{n})$ 

• But is it  $\Theta(\sqrt{n})$ ?

• Suppose we have a fair n-sided die that we roll k=1,2,3,4,... times. What is the probability we see a repeated outcome among the rolls?

• Let  $S_i$  be the event that the i-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome

• 
$$\Pr[S_i] = \frac{i-1}{n}$$

•  $\Pr[S_1 \cup \cdots \cup S_k] \leq ???$ 

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• Let  $S_i$  be the event that the i-th roll is a repeated outcome, conditioned on the previous rolls not being a repeated outcome

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$$\Pr[S_i] = \frac{i-1}{n}$$

• 
$$\Pr[S_1 \cup \dots \cup S_k] \le \frac{0}{n} + \dots + \frac{k-1}{n} \le \frac{k^2}{n}$$

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**Union Bound** 

• Suppose we have a fair n-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls?

•  $\Theta(\sqrt{n})$ 

## Trivia Question #1 (Birthday Paradox)

• Suppose we have a fair n-sided die. "On average", how many times should we roll the die before we see a repeated outcome among the rolls? Example: 1, 5, 2, 4, 5

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- Let c > 0 be a constant. What is  $\lim_{n \to \infty} \left(1 \frac{c}{n}\right)^n$ ?
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# CSCE 689: Special Topics in Modern Algorithms for Data Science

Lecture 3

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## Trivia Question #3 (Coupon Collector)

• Suppose we have a fair n-sided die. "On average", how many times should we roll the die before we all possible outcomes among the rolls? Example: 1, 5, 2, 4, 1, 3, 1, 6 for n = 6

- $\bullet \Theta(n)$
- $\Theta(n \log n)$
- $\Theta(n\sqrt{n})$
- $\Theta(n^2)$

#### Trivia Question #4 (Max Load)

• Suppose we have a fair n-sided die that we roll n times. "On average", what is the largest number of times any outcome is rolled? Example: 1, 5, 2, 4, 1, 3, 1 for n = 7

- $\Theta(1)$
- $\widetilde{\Theta}(\log n)$
- $\widetilde{\Theta}(\sqrt{n})$
- $\widetilde{\Theta}(n)$

#### Expected Value

• The expected value of a random variable X over  $\Omega$  is:

$$E[X] = \sum_{x \in \Omega} \Pr[X = x] \cdot x$$

The "average value of the random variable"

• Linearity of expectation: E[X + Y] = E[X] + E[Y]

## **Expected Value**

• Suppose we roll a 6-sided die

• Let X be the outcome of the roll

• What is E[X]?

#### Moments

• For p > 0, the p-th moment of a random variable X over  $\Omega$  is:

$$E[X^p] = \sum_{x \in \Omega} \Pr[X = x] \cdot x^p$$

#### Variance

• The variance of a random variable X over  $\Omega$  is:

$$Var[X] = E[X^2] - (E[X])^2$$

• Linearity of variance for *independent* random variables: Var[X + Y] = Var[X] + Var[Y]

• "How far numbers are from the average"

#### Variance

• Suppose X takes the value 1 with probability  $\frac{1}{2}$  and takes the value -1 with probability  $\frac{1}{2}$ 

• What is **E**[*X*]?

What is Var[X]?

#### Variance

- Suppose Y takes the value  $\frac{100}{2}$  with probability  $\frac{1}{2}$  and takes the value
  - -100 with probability  $\frac{1}{2}$
- What is E[Y]?

What is Var[Y]?

## Chebyshev's Inequality

• Let X be a random variable with expected value  $\mu \coloneqq \mathrm{E}[X]$  and variance  $\sigma^2 \coloneqq \mathrm{Var}[X]$   $\mathrm{Pr}[|X - \mu| \ge k\sigma] \le \frac{1}{k^2}$ 

 "What is the probability a random variable is far away from its average?"