CSCE 411: Design and Analysis of Algorithms

Lecture 1: Intro, Asymptotic Runtimes, Divide and Conquer

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Course Logistics

- Read section 2.3, and chapter 4 for first week of classes.
- Read (or skim) chapters 1-3 to ensure familiarity with prerequisites
- Syllabus quiz is due Sat, Jan 18. HW 1 and intro video due Fri, Jan 24

1 Computational Problems and Algorithms

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Input: A sequence of n numbers: a_1, a_2, \ldots, a_n Output: A permutation σ of the input sequence so that $a_{\sigma(1)} \leq a_{\sigma(2)} \leq \cdots \leq a_{\sigma(n)}$ Instance of this problem is the sequence Input: An array of n numbers: $[a_1, a_2, \ldots, a_n]$	
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mple 2. Min element Input: An array of n numbers: $[a_1, a_2, \dots, a_n]$	
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Output: The smallest element in the array and its index.	
2. An algorithm is a computational procedure that	

An algorithm is said to be correct if it

2 Asymptotic Runtime Analysis (Chapter 3)

2.1 Rules for runtime analysis

- \bullet *n* denotes the size of the input
- Each basic operation takes constant time
- We focus on the _____ runtime
- We only care about the _____ of the runtime

2.2 Some initial examples

Question 1. Given an array of n items, find whether the array contains a negative number using the following steps:

```
for i = 1 to n do

if a_i < 0 then

Return (true, i)

end if

end for
```

What is the runtime of this method?

- O(1)
- O(n)
- $O(n^2)$
- D It depends
- E Other
- Don't know, I need a reminder for how this works.

Question 2. Given an $n \times n$ matrix A, what is the runtime of summing the upper triangular portion using the following algorithm? (same answers).

```
sum = 0

for i = 1 to n do

for j = i to n do

sum = sum + a_{ij}

end for

end for

Return sum
```

2.3 Formal Definitions

Let n be input size, and let f and g be functions over \mathbb{N} .

Definition 3. Big \mathcal{O} notation.

A function g(n) = O(f(n)) (we say, "g is big-O of f(n)") means:

Definition 4. Big Ω notation.

$$g(n) = \Omega(f(n))$$
 means:

Definition 5. Θ notation.

$$g(n) = \Theta(f(n))$$
 means:

Equivalently, this means

Additional runtime examples

1.
$$4n^4 + n^3 \log n + 100n \log n$$

2.
$$n + 2(\log n)^2$$

3.
$$2^n + 10^{100}n^45$$

Logarithms in Runtimes Which of the following runtimes are the same asymptotically? Which are not?

- $O(n \log n)$ and $O(n \lg n)$
- $O(\log n)$ and $O(\log^2 n)$
- $O(\log n)$ and $O(\log(n^2))$
- $O(n^{\log 100})$ and $O(n^{\lg 100})$

3 How to Present an Algorithm

Presenting and analyzing can be broken up into four steps.

1. Explain: the approach in basic English

2. **Pseudocode**: for formally presenting the algorithmic steps

3. Prove: the **correctness**

4. Analyze: the runtime complexity

As a rule it's a good idea to go through all steps when presenting an algorithm. Sometimes we will focus more on just a subset of these (e.g., you may be asked to prove a runtime complexity of an algorithm on a homework but not a correctness proof).

We will go through all four steps when we present the *merge sort* algorithm.

4 The Divide and Conquer Paradigm

The divide and conquer paradigm has three components:

• Divide:

• Conquer:

• Combine:

Example: Mergesort (Textbook, Chapter 2.3, 4) Given n numbers to sort, apply the following steps:

- ullet Divide the sequence of length n into
- Recursively
- Combine

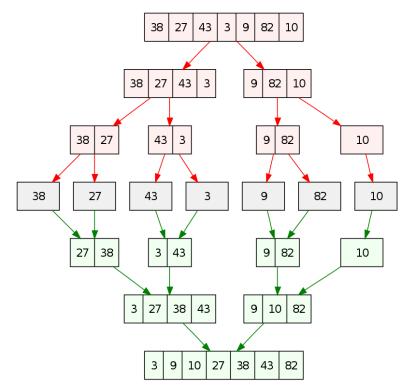


Image courtesy of Wikipedia: https://en.wikipedia.org/wiki/Merge_sort.

 $\overline{\text{MERGESORT}(A)}$

 $n = \mathbf{length}(A)$

if n == 1 then

else

 $m = \lfloor n/2 \rfloor$

end if

4.1 Analyzing The Merge Procedure

Correctness: To merge two sorted subarrays into a master array	
• Maintain a pointer to the	
• At each step,	
• Since subarrays are sorted, one of these numbers is	
• so place	
• At each step, we guarantee that:	
• So continuing until both subarrays are empty,	