

# CSCS 689: Special Topics in Modern Algorithms for Data Science

## Lecture 14

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# Presentation Schedule

- September 25: Team DAP, Team Bokun, Team Jason
- September 27: Galaxy AI, Team STMI
- September 29: Jung, Anmol, Chunkai

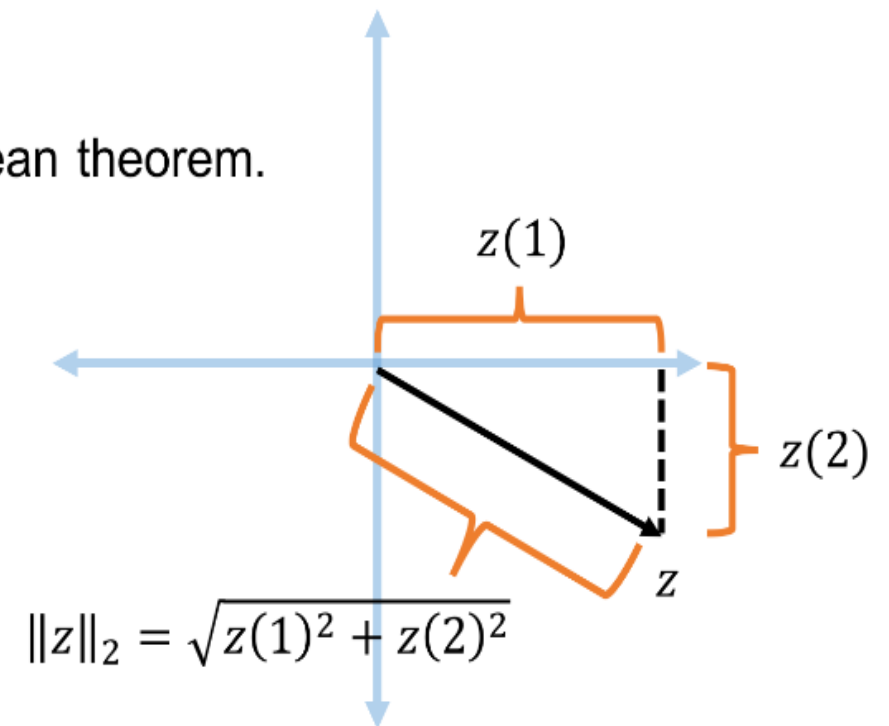
# Recall: Euclidean Space and $L_2$ Norm

- For  $z \in R^n$ , the  $L_2$  norm of  $z$  is denoted by  $\|z\|_2$  and defined as:

$$\|z\|_2 = \sqrt{z_1^2 + z_2^2 + \cdots + z_n^2}$$

- For  $x, y \in R^n$ , the distance function  $D$  is denoted by  $\|\cdot\|_2$  and defined as  $\|x - y\|_2$

Pythagorean theorem.



# Recall: CountSketch Summary

- CountSketch solves the  $L_2$  heavy-hitters problem: Given a set  $S$  of  $m$  elements from  $[n]$  that induces a frequency vector  $f \in R^n$  and a threshold parameter  $\varepsilon \in (0, 1)$ , output a list that includes:
  - The items from  $[n]$  that have frequency at least  $\varepsilon \cdot \|f\|_2$
  - No items with frequency less than  $\frac{\varepsilon}{2} \cdot \|f\|_2$
- Space usage:  $O\left(\frac{1}{\varepsilon^2} \log^2 n\right)$  bits of space

# $L_2$ Estimation

- **Goal:** Given a set  $S$  of  $m$  elements from  $[n]$  that induces a frequency vector  $f \in \mathbb{R}^n$  and an accuracy parameter  $\varepsilon \in (0, 1)$ , output a  $(1 + \varepsilon)$ -approximation to  $\|f\|_2$
- Find  $Z$  such that  $(1 - \varepsilon) \cdot \|f\|_2 \leq Z \leq (1 + \varepsilon) \cdot \|f\|_2$
- Find  $Z'$  such that  $(1 - \varepsilon) \cdot \|f\|_2^2 \leq Z' \leq (1 + \varepsilon) \cdot \|f\|_2^2$

# $F_2$ Moment Estimation

- **Goal:** Find  $Z'$  such that  $(1 - \varepsilon) \cdot \|f\|_2^2 \leq Z' \leq (1 + \varepsilon) \cdot \|f\|_2^2$

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1 7 7 7 3 7 7 1 4 1 1 1 1 5 1 1 7 1 7 5 1 7 7

# $F_2$ Moment Estimation

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1 7 7 7 3 7 7 1 4 1 1 1 1 5 1 1 7 1 7 5 1 7 7

$f_1$	$f_2$	$f_3$	$f_4$	$f_5$	$f_6$	$f_7$
10	0	1	1	2	0	9



# Johnson-Lindenstrauss Lemma

- **Distributional Johnson-Lindenstrauss Lemma:** Given  $\Pi \in R^{m \times n}$  with  $m = O\left(\frac{\log 1/\delta}{\varepsilon^2}\right)$  and each entry drawn from  $\frac{1}{\sqrt{m}}N(0,1)$ , then for any  $x \in R^n$  and setting  $y = \Pi x$ , then with probability at least  $1 - \delta$

$$(1 - \varepsilon)\|x\|_2 \leq \|y\|_2 \leq (1 + \varepsilon)\|x\|_2$$

# $F_2$ Moment Estimation

- **Algorithm**: Generate  $\Pi \in R^{m \times n}$  with  $m = O\left(\frac{\log 1/\delta}{\varepsilon^2}\right)$  and each entry drawn from  $\frac{1}{\sqrt{m}}N(0,1)$ . Set  $g = \Pi \cdot f$
- Whenever there is an update to a coordinate of  $f$ , update  $g$

# $F_2$ Moment Estimation

- **Algorithm**: Generate  $\Pi \in R^{m \times n}$  with  $m = O\left(\frac{\log 1/\delta}{\varepsilon^2}\right)$  and each entry drawn from  $\frac{1}{\sqrt{m}}N(0,1)$ . Set  $g = \Pi \cdot f$

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1 7 7 7 3 7 7 1 4 1 1 1 1 5 1 1 7 1 7 5 1 7 7

- Whenever there is an update to a coordinate of  $f$ , update  $g$
- $f = f + e_1$
- $f = f + e_7$
- $f = f + e_7$

# $F_2$ Moment Estimation

- **Algorithm:** Generate  $\Pi \in R^{m \times n}$  with  $m = O\left(\frac{\log 1/\delta}{\varepsilon^2}\right)$  and each entry drawn from  $\frac{1}{\sqrt{m}}N(0,1)$ . Set  $g = \Pi \cdot f$

1 7 7 7 3 7 7 1 4 1 1 1 1 5 1 1 7 1 7 5 1 7 7

- Whenever there is an update to a coordinate of  $f$ , update  $g$
- $f = f + e_1, g = g + \Pi e_1$
- $f = f + e_7, g = g + \Pi e_7$
- $f = f + e_7, g = g + \Pi e_7$

# AMS Algorithm

- Generate a random sign vector  $s \in \{-1, +1\}^n$
- Maintain  $Z = \langle s, f \rangle$
- Output  $W := Z^2$

1

1



2

1

2

1

1

1

2

1



1

2

2

2

1

# AMS Algorithm

- What values of  $Z$  did you get?
- $Z = \langle s, f \rangle = s_1 f_1 + s_2 f_2 + \cdots + s_n f_n$
- What values of  $W$  did you get?
- $W = Z^2 = \sum_{i,j} s_i s_j f_i f_j$

# AMS Algorithm

- What values of  $W$  did you get?

- $W = Z^2 = \sum_{i,j} s_i s_j f_i f_j$

$f_1$	$f_2$	$f_3$	$f_4$	$f_5$	$f_6$	$f_7$
9	6	0	0	0	0	0

# AMS Algorithm

- What is  $E[W]$ ?
- $Z = \langle s, f \rangle = s_1 f_1 + s_2 f_2 + \cdots + s_n f_n$
- $W = Z^2 = \sum_{i,j} s_i s_j f_i f_j$
- $E[W] = \sum_{i,j} E[s_i s_j f_i f_j] = \sum_i E[f_i^2] = \|f\|_2^2$



# AMS Algorithm

- What is  $\text{Var}[W]$ ?
- $Z = \langle s, f \rangle = s_1 f_1 + s_2 f_2 + \cdots + s_n f_n$
- $W^2 = Z^2 = \sum_{a,b,c,d} s_a s_b s_c s_d f_a f_b f_c f_d$
- $E[W^2] = \sum_{a,b,c,d} E[s_a s_b s_c s_d f_a f_b f_c f_d] = \sum_i E[f_i^4] + 6 \sum_{i \neq j} E[f_i^2 f_j^2] \leq 6 \|f\|_2^4$

# AMS Algorithm

- By Chebyshev's inequality,  $W$  will be a 9-approximation to  $\|f\|_2^2$  with probability  $2/3$

# AMS Algorithm

- How to get  $(1 + \varepsilon)$ -approximation?
- Repeat  $O\left(\frac{1}{\varepsilon^2}\right)$  times and take the average

# AMS Algorithm

- Space of algorithm:  $O\left(\frac{1}{\varepsilon^2}\right)$  words of space or  $O\left(\frac{1}{\varepsilon^2} \log m\right)$  bits of space