

# CSC 658: Randomized Algorithms

## Lecture 4

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# Last Time: Expected Value

- The expected value of a random variable  $X$  over  $\Omega$  is:

$$E[X] = \sum_{x \in \Omega} \Pr[X = x] \cdot x$$

- The “average value of the random variable”
- Linearity of expectation:  $E[X + Y] = E[X] + E[Y]$

# Last Time: Markov's Inequality

- Let  $X \geq 0$  be a non-negative random variable. Then for any  $t > 0$ :

$$\Pr[X \geq t \cdot \mathbb{E}[X]] \leq \frac{1}{t}$$

- Can rewrite as  $\Pr[X \geq t] \leq \frac{\mathbb{E}[X]}{t}$
- “Bounding the deviation of a random variable in terms of its average”

# Limitations of Markov's Inequality

- Let  $X$  be the outcome of a roll of a die. Then  $E[X] = 3.5 = \frac{7}{2}$

$$\Pr[X \geq 6] = \Pr\left[X \geq \frac{12}{7} \cdot \frac{7}{2}\right] \leq \frac{7}{12} \approx 0.5833$$

- We know  $\Pr[X \geq 6] = \frac{1}{6} \approx 0.167$

# Moments

- For  $p > 0$ , the  $p$ -th moment of a random variable  $X$  over  $\Omega$  is:

$$E[X^p] = \sum_{x \in \Omega} \Pr[X = x] \cdot x^p$$

# Variance

- The variance of a random variable  $X$  over  $\Omega$  is:

$$\text{Var}[X] = E[(X - E[X])^2]$$

- Can rewrite  $\text{Var}[X] = E[X^2] - (E[X])^2$  since  $E[E[X]] = E[X]$
- “On average, how far numbers are from the average”

# Variance

- Can rewrite  $\text{Var}[X] = E[(X - E[X])^2]$  since  $E[E[X]] = E[X]$

$$\begin{aligned} E[(X - E[X])^2] &= E[X^2 - 2X \cdot E[X] + (E[X])^2] \\ &= E[X^2] - 2E[X] \cdot E[E[X]] + (E[X])^2 \\ &= E[X^2] - 2E[X] \cdot E[X] + (E[X])^2 \\ &= E[X^2] - 2(E[X])^2 + (E[X])^2 \\ &= E[X^2] - (E[X])^2 = \text{Var}[X] \end{aligned}$$

# Variance

- The variance of a random variable  $X$  over  $\Omega$  is:

$$\text{Var}[X] = E[X^2] - (E[X])^2$$

- Linearity of variance for *independent* random variables:  $\text{Var}[X + Y] = \text{Var}[X] + \text{Var}[Y]$

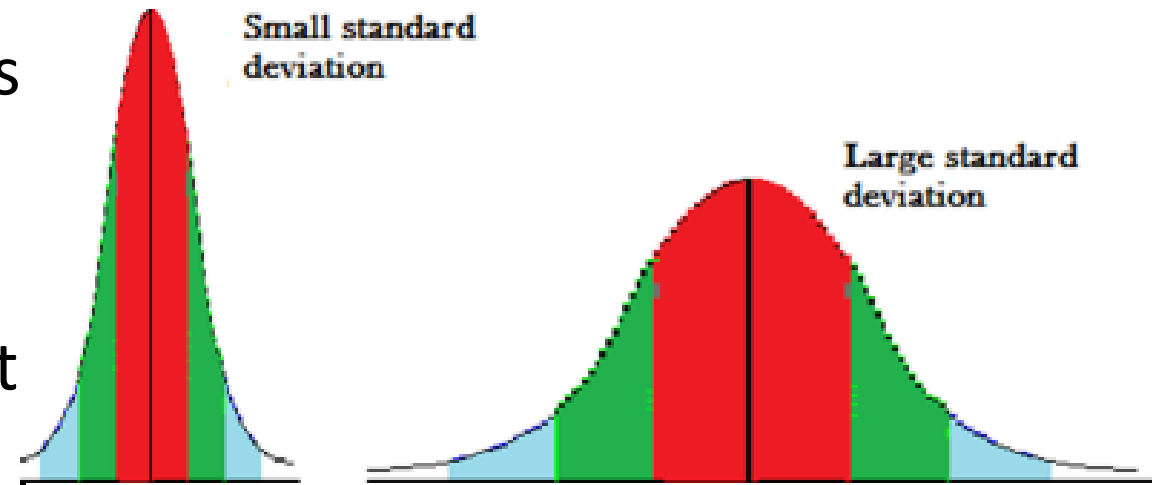


# Variance and Standard Deviation

- The variance of a random variable  $X$  over  $\Omega$  is:

$$\sigma^2 = \text{Var}[X] = E[X^2] - (E[X])^2$$

- The standard deviation  $\text{std}(X)$  of a random variable  $X$  is  $\sigma$ , and measures how far apart the outcomes are
- Standard deviation is in the same unit as the data set



# Variance

- Suppose  $X$  takes the value  $1$  with probability  $\frac{1}{2}$  and takes the value  $-1$  with probability  $\frac{1}{2}$
- What is  $E[X]$ ?
- What is  $\text{Var}[X]$ ? What is  $\text{std}(X)$ ?

# Variance

- Suppose  $Y$  takes the value  $100$  with probability  $\frac{1}{2}$  and takes the value  $-100$  with probability  $\frac{1}{2}$
- What is  $E[Y]$ ?
- What is  $\text{Var}[Y]$ ? What is  $\text{std}(Y)$ ?

# Markov's Inequality

- Let  $X \geq 0$  be a non-negative random variable. Then for any  $t > 0$ :

$$\Pr[X \geq t \cdot \mathbb{E}[X]] \leq \frac{1}{t}$$

- Can rewrite as  $\Pr[X \geq t] \leq \frac{\mathbb{E}[X]}{t}$

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- Let  $X \geq 0$  be a non-negative random variable. Then for any  $t > 0$ :

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- Can rewrite as  $\Pr[X \geq t] \leq \frac{\mathbb{E}[X]}{t}$
- We have  $\Pr[|X| \geq t] = \Pr[X^2 \geq t^2]$

# Using Markov's Inequality

- We have  $\Pr[|X| \geq t] = \Pr[X^2 \geq t^2]$

$$\Pr[|X| \geq t] = \Pr[X^2 \geq t^2] \leq \frac{E[X^2]}{t^2}$$

- Plug in  $X - E[X]$  for  $X$

$$\Pr[|X - E[X]| \geq t] \leq \frac{E[(X - E[X])^2]}{t^2}$$

# Toward Chebyshev's Inequality

$$\Pr[|X - E[X]| \geq t] \leq \frac{E[(X - E[X])^2]}{t^2}$$

# Chebyshev's Inequality

$$\Pr[|X - E[X]| \geq t] \leq \frac{E[(X - E[X])^2]}{t^2}$$

- Recall that  $\text{Var}[X] = E[X^2] - (E[X])^2 = E[(X - E[X])^2]$
- $\Pr[|X - E[X]| \geq t] \leq \frac{\text{Var}[X]}{t^2}$



# Chebyshev's Inequality

- Let  $X$  be a random variable with expected value  $\mu := E[X]$  and variance  $\sigma^2 := \text{Var}[X]$

- $\Pr[|X - E[X]| \geq t] \leq \frac{\text{Var}[X]}{t^2}$  becomes  $\Pr[|X - E[X]| \geq t] \leq \frac{\sigma^2}{t^2}$

$$\Pr[|X - \mu| \geq k\sigma] \leq \frac{1}{k^2}$$

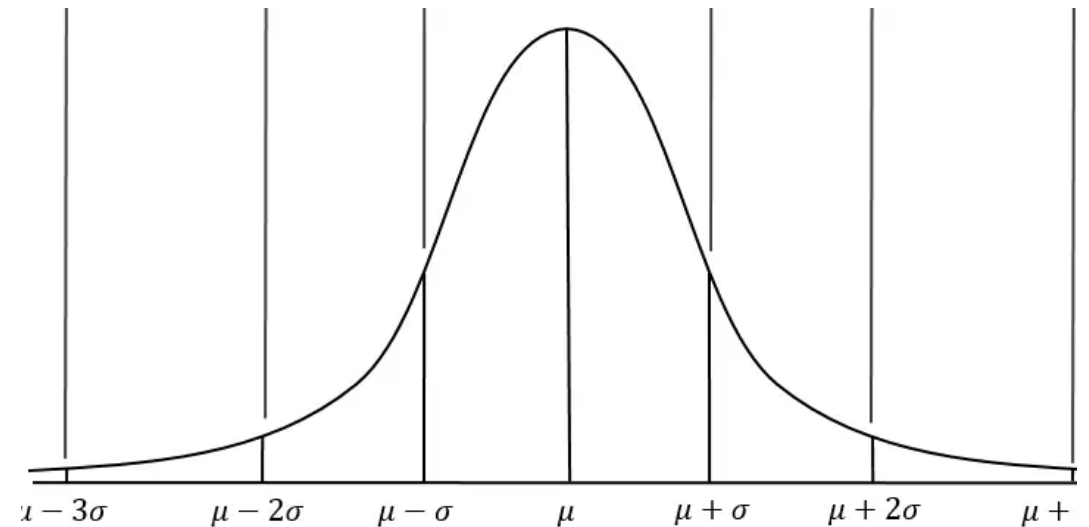
- “Bounding the deviation of a random variable in terms of its standard deviation / variance”

# Chebyshev's Inequality

- Let  $X$  be a random variable with expected value  $\mu := E[X]$  and variance  $\sigma^2 := \text{Var}[X]$

$$\Pr[|X - \mu| \geq k\sigma] \leq \frac{1}{k^2}$$

- Do not require assumptions about  $X$



# Chebyshev's Inequality

- Let  $X$  be the outcome of a roll of a die. Then  $E[X] = 3.5 = \frac{7}{2}$  and  $\text{Var}[X] = \frac{35}{12} \approx 2.92$  so  $\text{std}(X) \approx 1.71$

$$\begin{aligned}\Pr[X \geq 6] &= \Pr[X - 3.5 \geq 2.5] \\ &= \Pr[X - 3.5 \geq 1.41 \cdot 1.71] \\ &\leq \frac{1}{1.41^2} \approx 0.4667\end{aligned}$$

- Recall that Markov's inequality bounded this by **0.5833**

# Law of Large Numbers

- Let  $X_1, \dots, X_n$  be random variables that are independent identically distributed (i.i.d.) with mean  $\mu$  and variance  $\sigma^2$
- Consider the sample average  $X = \frac{1}{n} \sum_i X_i$ . How does it compare to  $\mu$ ?
- $\text{Var}[X] = \frac{1}{n^2} \sum_i \text{Var}[X_i] = \frac{\sigma^2}{n}$
- By Chebyshev's inequality,  $\Pr[|S - \mu| \geq t] \leq \frac{\sigma^2}{nt}$

# Law of Large Numbers

- By Chebyshev's inequality,  $\Pr[|S - \mu| \geq t] \leq \frac{\sigma^2}{nt}$
- **Law of Large Numbers:** The sample average will always concentrate to the mean, given enough samples

# Use Case

- Suppose we design a randomized algorithm  $A$  to estimate a hidden statistic  $\Theta$  of a dataset and we know  $0 < \Theta \leq 1000$
- Suppose each time we use the algorithm  $A$ , it outputs a number  $X$  such that  $E[X] = \Theta$  and  $\text{Var}[X] = 100\Theta^2$
- What can we say about  $A$ ?
- $\Pr[|X - \Theta| \geq 30\Theta] \leq \frac{1}{9}$  and  $Z \leq 1000$  so  $\Pr[|X - \Theta| < 30,000] > \frac{8}{9}$

# Accuracy Boosting

- How can we use  $A$  to get additive error  $\varepsilon$ ?

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- How can we use  $A$  to get additive error  $\varepsilon$ ?
- Repeat  $A$  a total of  $\frac{10^{12}}{\varepsilon^2}$  times and take the average
- The variance of the average is  $\frac{\varepsilon^2}{10^{10}} \Theta$  and  $\Pr[|X - \mu| \geq k] \leq \frac{\sigma^2}{k^2}$
- $\Pr[|X - \Theta| \geq \varepsilon] \leq \frac{\Theta}{10^{10}}$  and  $\Theta \leq 1000$  so  $\Pr[|X - \Theta| < \varepsilon] > 0.999$



# Accuracy Boosting

- Algorithmic consequence of Law of Large Numbers
- To improve the accuracy of your algorithm, run it many times independently and take the average