



# Data Structures

## Lecture 7: Graph

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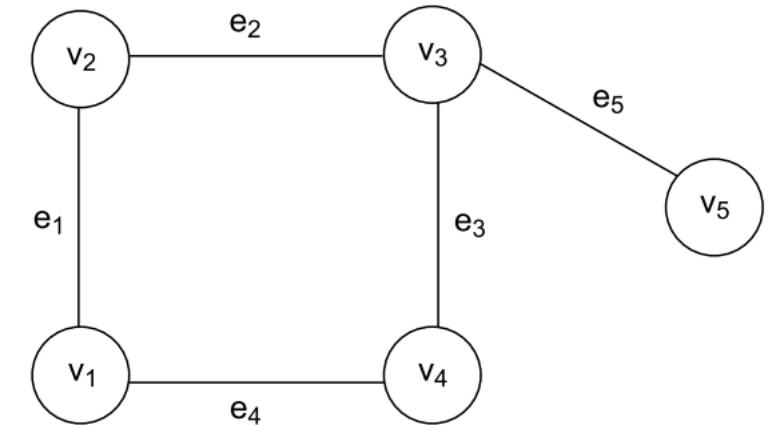
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- Concept of Graphs
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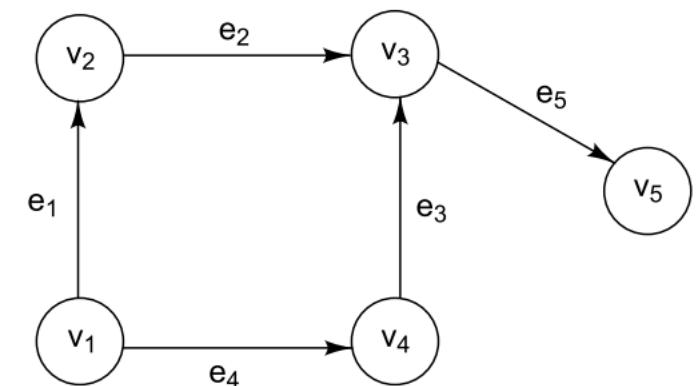
# Graph definition

- A Graph is a non-linear data structure that consists of vertices (nodes) and edges.
- A graph **G** consists of the following elements:
  - A set **V** of vertices or nodes, where  $V=\{v_1, v_2, v_3, \dots, v_n\}$
  - A set **E** of edges also called arcs where,  $E=\{e_1, e_2, e_3, \dots, e_n\}$
  - Hence,  $\mathbf{G}=(V,E)$
- **If  $e=(u,v)$  and  $e=(v, u)$  means same then the graph is undirected.**

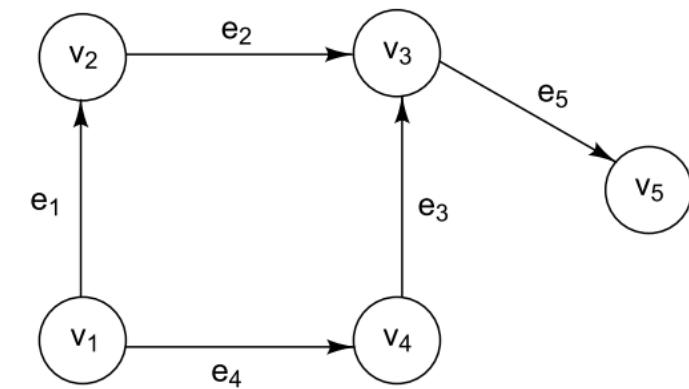
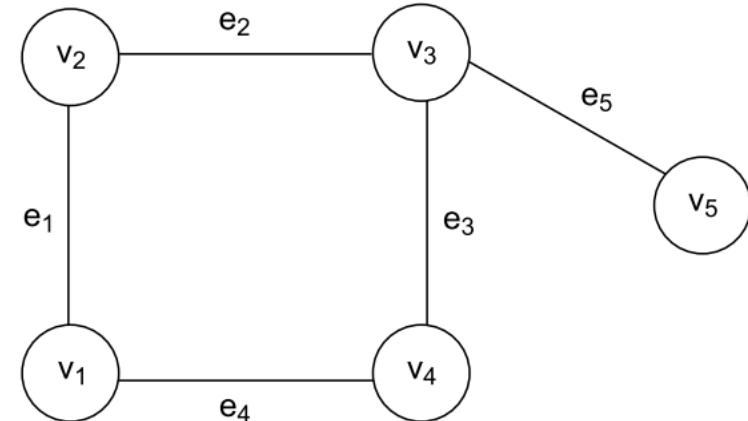


# Directed Graph

- If we replace each edge of the Graph **G** with **arrows**, then it will become a ***directed graph or diagraph***.
- In this graph, the set of vertices and edges are:
  - $V(G)=\{v_1, v_2, v_3, v_4, v_5\}$
  - $V(E)=\{(v_1, v_2), (v_2, v_3), (v_1, v_4), (v_4, v_3), (v_3, v_5)\}$



Key Terms	Description
Adjacent node	If $e(u, v)$ represents an edge between $u$ and $v$ vertices then both $u$ and $v$ are called adjacent to each other. That means, $u$ is adjacent to $v$ and $v$ is adjacent to $u$ .
Predecessor node	If $e(u, v)$ represents a directed edge from $u$ to $v$ then $u$ is a predecessor node of $v$ .
Successor node	If $e(u, v)$ represents a directed edge from $u$ to $v$ then $v$ is a successor node of $u$ .
Degree	Degree of a vertex is the number of edges connected to a vertex. For example, in the graph shown in Fig. 9.1, the degree of vertex $v_3$ is 3.
Indegree	In a directed graph, indegree of a vertex is the number of edges ending at the vertex.
Outdegree	In a directed graph, outdegree of a vertex is the number of edges beginning at the vertex.
Path	A path is a sequence of vertices each adjacent to the next. For example, in the graph shown in Fig. 9.2, the path between the vertices $v_1$ and $v_5$ is $v_1-v_2-v_3-v_5$ .
Cycle	It is a path that starts and ends at the same vertex.
Loop	It is an edge whose endpoints are same that is, $e = (u, u)$ .
Weight	It is a non-negative number assigned to an edge. It is also called length.
Order	Order of a graph is the number of the vertices contained in the graph.
Labeled Graph	It is a graph that has labeled edges.
Weighted Graph	It is a graph that has weights assigned to each of its edges.
Connected Graph	It is an undirected graph in which there is a path between each pair of nodes.
Strongly Connected Graph	It is a directed graph in which there is a route between each pair of nodes.
Complete Graph	It is an undirected graph in which there is a direct edge between each pair of nodes.
Tree	It is a connected graph with no cycles.



# Types of Graphs

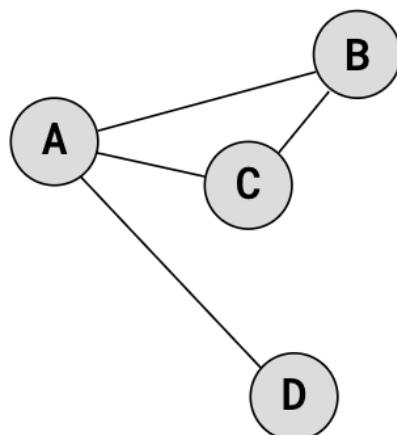
- Based on Edge Direction
  - Undirected Graph
  - Directed Graph
- Based on Weight of Edges
  - Weighted Graph
  - Unweighted Graph

# Graph Representation

- Adjacency Matrix (2D Array)
- Adjacency List (Linked List)

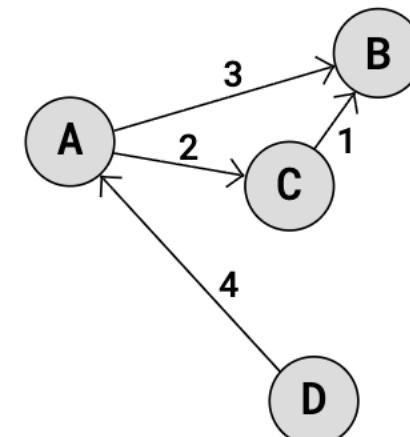
# Adjacency Matrix

- The Adjacency Matrix is a **2D array (matrix)** where each cell on index **( $i,j$ )** stores information about the edge from vertex  **$i$**  to vertex  **$j$** .
- 2D array  $\rightarrow A[V][V]$
- $A[i][j] = 1$ , if edge exists between vertex  $i$  and  $j$ , else **0**.



	A	B	C	D
A	0	1	1	1
B	1	0	1	0
C	1	1	0	0
D	1	0	0	0

An undirected Graph  
and the adjacency matrix

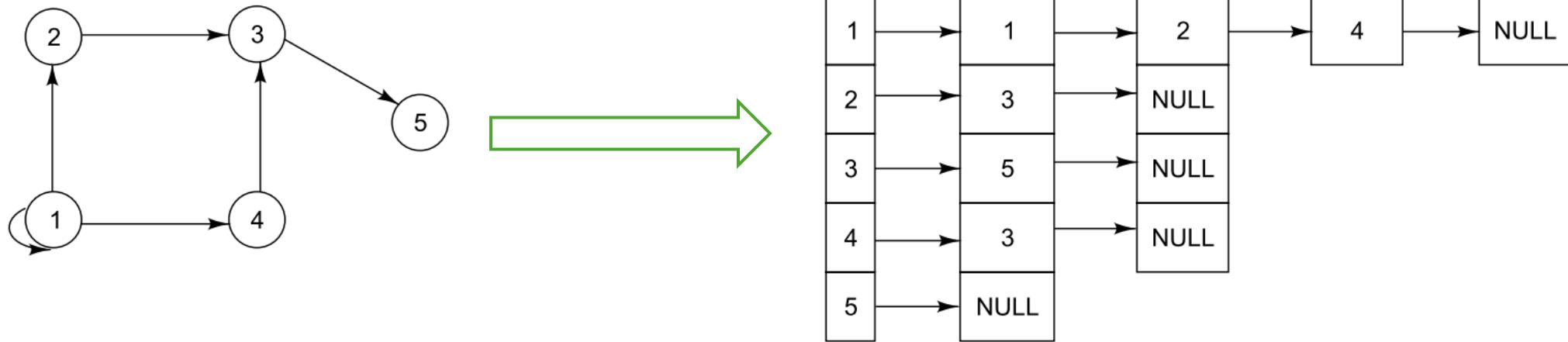


	A	B	C	D
A	0	3	2	0
B	0	0	0	0
C	0	1	0	0
D	4	0	0	0

A directed and weighted Graph,  
and its adjacency matrix.

# Adjacency List

- Array of ***linked lists (or dynamic arrays)***
- Each vertex stores a list of its adjacent vertices.
- Suitable for ***sparse*** graphs.

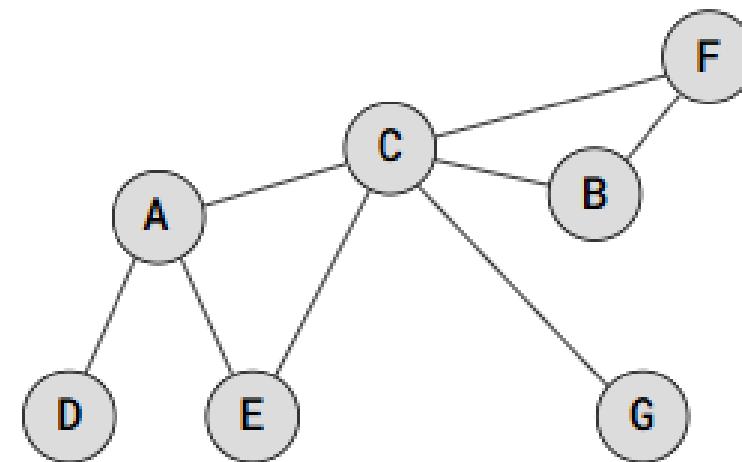


# Graph Traversal Techniques

- **Depth First Search (DFS):**
  - Explores as far as possible along each branch before backtracking.
  - Uses a stack (or recursion).
- **Breadth First Search (BFS)**
  - Explores all neighbors of a vertex before moving to the next level.
  - Uses a queue.

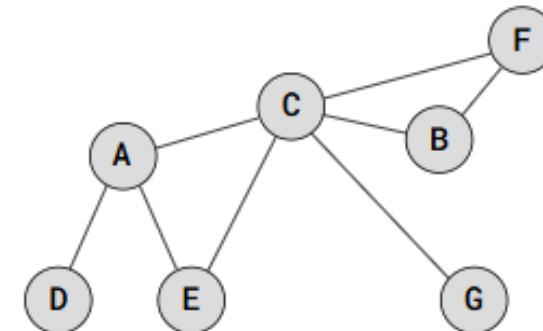
# Depth First Search (DFS)

- Start DFS traversal on a vertex.
- Uses a stack (or recursion).
- Do a recursive DFS traversal on each of the adjacent vertices as long as they are not already visited.
- Path: **D,A,C,B,F,E,G**



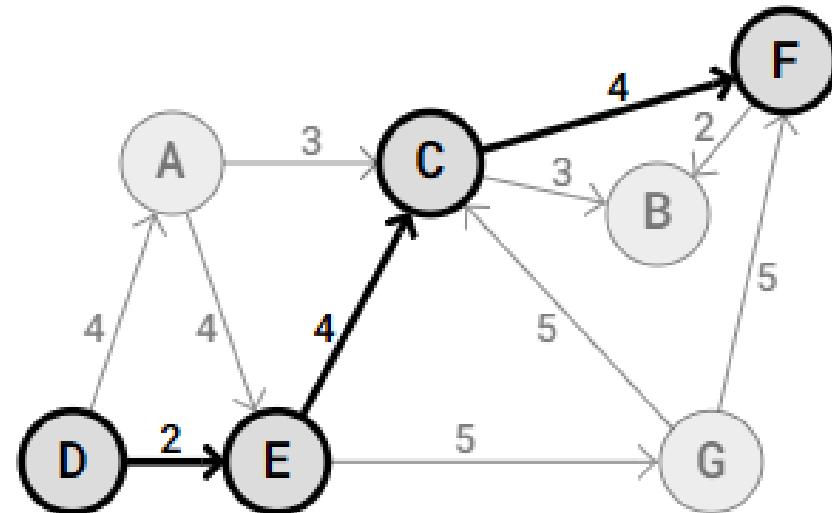
# Breath First Search (BFS):

- Put the starting vertex into the **queue**.
- For each vertex taken from the queue, visit the vertex, then put all unvisited adjacent vertices into the queue.
- Continue as long as there are vertices in the queue.
- Path: **D,A,C,E,B,F,G**



# Shortest Path

- To solve the shortest path problem means to find the **shortest possible route or path** between two vertices (or nodes) in a Graph.
- In the shortest path problem, a Graph can represent anything from a **road network** to a communication network, where the vertices can be intersections, cities, or routers, and the edges can be **roads, flight paths, or data links**.
- **Solutions (SPP): Dijkstra's algorithm and the Bellman-Ford algorithm** find the shortest path from one start vertex, to all other vertices.



# Graph Algorithms

Algorithm	Purpose	Concept
Dijkstra's Algorithm	Shortest path (single source)	Greedy algorithm
Floyd-Warshall	All-pairs shortest paths	Dynamic programming
Prim's Algorithm	Minimum Spanning Tree	Greedy (build tree step-by-step)
Kruskal's Algorithm	Minimum Spanning Tree	Greedy (sort edges, union-find)
Topological Sort	Ordering in a DAG	DFS-based sorting
Bellman-Ford	Shortest path (with negative weights)	Dynamic programming

# Real life applications of Graphs

- **Social Networks:** Each person is a vertex, and relationships (like friendships) are the edges. Algorithms can suggest potential friends.
- **Maps and Navigation:** Locations, like a town or bus stops, are stored as vertices, and roads are stored as edges. Algorithms can find the shortest route between two locations when stored as a Graph.
- **Internet:** Can be represented as a Graph, with web pages as vertices and hyperlinks as edges.
- **Biology:** Graphs can model systems like neural networks or the spread of diseases.

# References

- **Chapter 9:**
  - **Data Structures using C** by E. Balagurusamy

# Thank You