First AC: Costin-Andrei Oncescu, Romania (31:07) #AC = 16

problem author: Karol Pokorski

Easiest possible version

$$F_i = 1, f_i = 1$$
  
 $C_i = 1, c_i = 1$   
 $n = 1$  (one machine)

(consider the most profitable order)

#### Standard version

$$F_i = 1, f_i = 1$$
  
 $C_i = 1, c_i = 1$   
 $C_i = 1, c_i = 1$   
 $C_i = 1, c_i = 1$ 

$$O(m \times c_1)$$

dp[cores] - the largest profit to have so many cores

Double version

$$F_i = 1, f_i = 1$$
  
 $C_i = 1, c_i = 1$   
 $n = 1$  (one machine)

 $O(n \times (n \times C) + m \times (m \times C))$ two knapsacks

Double version

$$F_{i} \le f_{i} \leftarrow \text{works too}$$

$$C_{i} = 1, C_{i} = 1$$

$$n = 1 \text{ (one machine)}$$

 $O(n \times (n \times C) + m \times (m \times C))$ two knapsacks

One knapsack with modified items, e.g.:

a machine with weight -7 and value -15 - a task with weight 5 and value 20

We must end with total weight 0 or smaller.

$$O((n + m) \times (n \times C))$$

Sort by f, F, decreasingly.

is 0 or smaller at every moment of time. Then just guarantee that the total weight

$$O((n + m) \times (n \times C))$$

The alternative knapsack

$$V_i = 1, \ V_i = 1$$

dp[cores] → dp[money]

$$(u \times (m + u))O$$