10/28 Relations § 3.3

You can think of a relation from one set X to another set Y as a chart that 11sts now elements from X are "related" to elements from Y. For example, we can imagine a chart that 11sts for each student in a school the classes they're taking

Student | Class

Bill Economics

Bill English

Alexis English

Tordan Chemistry

Notice that unlike a function, each student can take multiple classes. Also, may be a student is faking no classes at all (e.g. the yire on a leave of absence).

Defin Formally, a relation R from set X to set Y is any subset of X x Y, i.e., any set of ordered pairs (xiy) wxxe X and y & Y.

If (x,y) & R then we write x Ry and say that "X is related to y."

E.g. with the student/class example the relation is $R = \{(Bill, Econ.), (Bill, Eng.), (Alexis, Eng.), (Jordan, Cham)...}$ and since Alexis is taking English we could also write Alexis R English.

Notice: A function f: X-> Y is a very special kind of relation from X to Y.

The most important relations are when X = Y'.

Defin It R is a relation from X to X, we say it is a relation on the Set X.

E.g. If $X = \{1,2,3\}$ then \leq defines a relation on X (i.e. "a is related to b" if and only if "a \leq b"). The set of creelened pairs for this relation is; $R = \{(1,1), (1,2), (1,3), (2,2), (2,3), (3,3)\}$. We can represent this same information as a digraph

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there we draw a viertex" (dot.) for each element of x; and draw an arrow a b whenever a R b.

Notice that if a Ra then we have a loop: ?

Defin The relation R on X is reflexive.

if x R x for all x E X.

Eig. The \leq relation on $\geq 1,2,33$ is reflexive (means we have a loop at every vertex). But if we did the relation given by \leq instead;



this is not reflexive (no loops at all have)

Defn the helation Ron X is symmetric if whenever x Ry for x,y \in X we also have y Rx.

Eig. The relation \(\lambda \text{ on } \gamma_{1,2,3} \) is not symmetric.

Since 1 = 2 but 2 \$1. For a symmetric relation we need arrows to look like:

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نهد نیک a. (or a no arraws)

E.g. An example of a symmetric relation Ris

X = Estudents at Howard? and

x Ry if "x has a class with y" for x,y \in X.

This is because if Person & his a class

with Person y, well then also Person y has a class with Person X (the same class that they're in!).

Relations & are "opposite" from symmetric, so:

Det'n The relation R on X is anti-symmetric if whenever x R y and y R x for x, y \in X then x = y.

E.g. The relation & lon X or or any set of numbers is antisymmetric since if

XSY and YSX then we must have X=Y.

(The relation < is also auti-symmetric since

there are no x,y at all with x < y and y < x).

Antisymmetric: No at but loops a OKV

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There is one more important property of relation \leq :

Defin A relation R on X is called transitive

if whenever x R y and y R z for $x, y, z \in X$,

we must have that x R z:

Eig. The relation E(ar2) is transitive because if $a \leq b$ and $b \leq c$ then certainly $a \leq c$.

Q' 15 the relation "has a class with" on Students transitue?

A: No! Manybe Bill has a class with Alexis (like English) and Alexis has a class with Cole (like Bis (ogy), but Bill has no class with Cole (ne's not taking Bislogy).

Defin A relation R on X that is;

• restexive

• and transitive

is calced a partial order on X.

E.g. & is a partial order on X= \(\xi\). \(\xi\) or any set of numbers)

Partial orders behave like \(\xi\): they

let us "compare" things in \(\xi\).

F. 9 : Consider a list of tasks you have to do to complete Some project. Maybe the project is "make a PB& I sandwich" and so the tasks one: 1. Toast two slives of bread. 2. Spread peanut butter on one slice. 3. Sprend jelly on the other slice. 4. Put the two slices together. Some of these tasks have to accome before others (e.g. have to do I before 2). So define relation K on the set of tasks by: i Rj if i=j or task i must bedore before task j. The digraph of this relation is: restexive anti-sy metre Notice how no arrows between 2 and 3 since spreading PB and I can be done in eather order. Also: notice we get a partial order on the tasks! If R is a partral order on X and x, y & X WC Say x and y are comparable if x Ry or y RX and say they are in comparable otherwise. E.g. In pB& J example, the tasks of spreadily PB and spreadily J are incomparable (can be done in any order). The partial order R on X is called a total order if every pair x, y EX is comparable. E.y. Relation & Con any set of H's) is a total order but "do before" relation on tasks net a total order

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Compositions of relations and inverse relations Now let's return to discussing relations R from X to Y. Recall that a function fix-yis a special such relation, and we can generalize to relations the important functional notions of composition and inversion.

Defn let R, be a relation from X to Y and Rz a relation from Y to Z. The composition Rz = R, is a relation from X to Z where we have for xex, x (Rz = R) Z if and only if there is y & Y with x R, y and y Rz Z.

(x) R_1 (y) R_2 (z) => (x) $R_2 \circ R_1$ (z)

Def'n Let R be a relation from X to Y. The inverse relation R' is a relation from Y to X when R' = {Ly.x): (x,y) ∈ R}.

"reverse" every onelesed par.

Note: For a function $f: X \to Y$, the inverse $f^{-1}: Y \to X$ is defined only when f is a bijection (1-to-1 and onto). But inverse relation R^{-1} is a luxurys defined

Note: If Ris relation on X live. from X to X), then digraph of R' obtained from digraph of R by reversing direction of all arrows.

e.g. 12 304 (12 12 3)

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Equivalence Relations § 3.4

Let X be a set. Recall that a partition of X is a collection S of Subsets of X such that every element of X belongs to exactly one of the subsets in S.

E-g. If $\chi = \{1,2,3,4,5\}$ then one performance $\chi = \{1,3,4\}, \{2,5\}$

A partition of X is a way of "breaking X into groups" and we can use a partitition of to dofte a relation ont; we have x Ry if and only if x and y are in same subset in S.

- Eig. with previous set partition, the digraph of Ris:

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Theorem The relation R on set X defined from a partition Sof X is: • reflexive

· Symmetric

· and transitive.

PS: These are all easy to check directly. Reflexive: Ix is insome subset as x. Symmetric: if x is in Some subset asy, then y is sine as x. Transitive: if x and y in some subset, and y and Z, then some for x and Z. B

Defin A relation Ron X that is:

· reflexive

· symnetric

(compare to let of partial order)

, and transitive

is called an equivalence relation on X

An equivalence relation on X is a way elements of X can be "the same."

Eg. Relation R on IR where xRy if x=y2 is an equiv. relation.

E.g. Let n be any positive integer. We define relation R

on Z by xRy if X-y is a multiple of n.

Exercise: This is an equivalence relation on Z.

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We've seen that partitions give equiv. relations. Converse is also true.

Theorem Let R be an equiv. relation on X, Let $a \in X$ be any element and define $[a] := \{x \in X : x \text{ Ra}\}$ (things that Then $S = \{\{a\}: a \in X\}$ is a partition of X.

Pf: We need to show that every x EX belongs to exactly one subset in S. It clearly belongs to Ex3 (by reflex.) f So suppose it also belonged to EyJ. We want to show that ExJ=EyJ. So let ZE[x]. Then ZRx, and since xRy, we have ZRy, i.e. ZE EyJ (using trans.). By symmetry, also have yRx, so if ZE[y] then by same argument ZE[x]. Thus [x]-EyJ, as chimed. By

Des'n The sets [a] for a 6 x from the previous theorem are called the equivalence classes of the equiv. relation R.

E.g. with R being equiv. relation on IR w/ xRy if x2=y2, equivalence classes are {a;-a} for a ∈ IR, i.e. each number is grouped with its negative.

For the "X-y is a multiple of n" equil relations on the integers Z? 711/4

Combinatorics: Basic Counting Principles & G.A.

We are now starting a new chapter (our last of the senester): Chapter 60 on combinatorics which is just a fancy word for "counting!" We will tearn many techniques for country the elements of a finite set. We stort with some basic counting principles:

E.g. Suppose for a meal you get to choose

- · an appetiter: either soup or salad,
- · a main course: chicken, fish, or pasta,
- · a dessert: either ice cream or pie. How manytotal meals are possible with these choices?

A: We can represent all the droises by drawing a "decision tree":

soup Soul Soul of Soul

ne turbod about these a while ago

(soup, usour, director, cic.) pier

(Salad,
Parta,
Pie)

We see that there are 12 = 2 × 3 × 2 total meds: We multiply the choices at each step to get total.

Thru (Multiplication Principle for Country)

Suppose we make an object via a series of steps, where we have ki choices for step 1, kz choices for step 2, etc. down to km choices for step m. Then the total # of objects we can make is k, x kz x ··· x km

Remark: We saw before that for Cartesian product X, x X2x...x Xm of Sets we have # (X, x...x Xm)= # X, x * Xz x...x # Xm.
This is basically the same as multiplication principles

Let's see Some more examples of multiplication principle,

Q: A US telephone # is 10 dig. 75 long, and the

Ist digit cannot be a zero. How many telephone # fore

A: We have 9 digits to choose for the 1st digit and

To for each of the 9 other digits, 90 by mutiplication principle:

9 x 10 x 10 x ... x 10 = 9 x 109 = 9 billion.

Recall: Weive see in that the total # of subsets of \$1,2,..., n3 is 2".

This is easy to see with the multiplication principle; to make a subset we decide: include 1 or not? (2 chairs)
include 2 or not? (2 chairs)

This is n steps, with 2 chures at each step, so in total, IF possibilities = $2 \times 2 \times 12^2 = 2^n$.

Sometimes it takes a little more thought to see how to apply the multiplication principle to a country problem; Q: How many reflexive relations on \$1,2,..., no are there?

A: Remember that a relation is a subset of X x X.

We know that (x,x) must be in our relation R for each x \in X so that it will be reflexive.

How many choices do we have for other elements of R?

Well, for each (x,y) \in X \times X with x \times y, we can include that (x,y) or not (= 2 choices). The # of (x,y) w/ x \times y is n x (n-1) since we can choose any x \in X for 11t coordinate, and then ore of the

(n-1) other elements of X for the Indicordante y. Thus, we have nx(n-1) binary choice, for make our neflexive nebotron R, so # of such nelations is: $2 \times 2 \times \cdots \times 2 = 2^{n(n-1)} - 2^{n^2-n}$ (A similar, but simpler, argument should that there are 2^{n^2} total relations on $X = \xi 1, 2, ..., n_3$.) Q: Let X = {1,2,..., n? How many ordered pairs (A,B) of subsets of X satisfying A C B C X are there? A: It is helpful to draw a Vern diagram of ar stration; B x extract there are 3 regions in this Venn dorgram · things in A, · things in B-A, · things in X-B. To make an ordered pair (A,B) of this form, we can therefre choose for each i=1,2,..., where to place i. . place 1 in A, B-A, or X-B? (3 choice) · Place 2 in A, B-A, or X -B? (3 chilles) ·Place un A, B-A, or X-B? (3 chara), Thus, we have a steps with 3 choicer at each step, So totalt of possibilities = 3x3x...x3 = [3] Exercise; What about (A, B, C) with A = B = C = \(\xi \), 2, -1, n ? ?

1/7 Addition Principle + Principle of Inclusion - Exclusion \$6.1 Sometimes we are try by to count objects that have multiple "kinds": E.g. Q: Let X = {a,b}. How many Strings in X* are How which have length 3 or length 4? H: The # of strings of length 3 in X*= 2x2x2=23 by multiple.
of strings of conth 4 = 2x2x2x2= 24 # of strings of length 3014 = 23+ 24 = 8+16 = 24. We see another counting principle maction here: Theorem (Addition Principle for Country) 1 If X, Xz,..., Xm are disjoint sets (mouning Xin Xj = & for all i = j, i.e. set have no common elevats) -then #(X,UXU...UXm) = +x, + +x2+...+ #Xm, We see that, as long as the set rare disjoint, we can count any grouping of sets just by adding together: Eg. Q: # of strings in Eq. 13th of length 3 or 4 or 5? A: 23+24+25, by the aldidron principle.

Eg. Alexis, Ben, Cole, David, and Erica area 5 person group.
They have + choose a: President, Vice Aeridant, & Treasurer.

a) flow may ways are there to do this?

b) How many ways are there if we require that either Alexic or Ben is the Prosident?

a): We can choose any of the 5 people for prez. Then for VP we can choose any of the remaining 4 and similarly for the owner we can choose my of the remaining 3.

By mult. principle: $5 \times 4 \times 3 = 60$.

b) It Alexis is president, we have $4 \times 3 = 12$ choses for treasure. It Ben is prez, we also have $4 \times B = 12$ disibilifier uptensmer. By add from principle # of choices is 12 + 12 = 24.

But what if the sets one not digistent? Then we use.

Theorem (Principle of Inclusion - Exclusion)

#(XUY) = #X + #Y - #(X NY) = xny= #

cre disjoint

To see why P.I. E. werks, look at venn diagram

when we all #x to #y, we

count things in Xny & double,

have to subtract -#(xny) to correct.

E.J. c) How many ways one there to pick prez, vi, treusous, where either Alexic is Prez. or Ben is VP?

c): Let X = assignments where Alexii is Proz.

Then #X = 4x3, # of choices of VP + treasurer.

Let Y = assignments where Ben is VP

Then #Y = 4x3, # of choices of Prez + treasurer.

We want to compute #(XUY). By P.I.E.,

We also need to know the Size of XNY:

#X NY = 3, since if A is Prez and Bis VP we

(an choice any of remaining 3 to be treasurer.

SO. #(XUY) = 12 + 12 - 3 = [21] ways to make Alexi, Aez.

dx #Y a(xny) or Beyn UP.