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## Specializations of symmetric functions

So far we haven't done much counting w/ symmetric f.n.'s. One way to get interesting sequences of #'s from sym f.n.'s is by specializing them, i.e., plugging in values.

Prop: (a)  $e_k(\overbrace{1,1,\dots,1}^n, 0, 0, \dots) = \binom{n}{k}$

(b)  $h_k(\overbrace{1,1,\dots,1}^n) = \left(\binom{n}{k}\right)^k = \binom{n+k-1}{k}$  "multi-choose"

Pf: Recall  $e_k(x_1, \dots, x_n) = \sum_{\substack{I \subseteq [n] \\ \#I = k}} \prod_{i \in I} x_i$  product of k distinct variables, so clearly setting  $x_i = 1 \ \forall 1 \leq i \leq n$  gives  $e_k(1, 1, \dots, 1) = \binom{n}{k}$ .

Similarly,  $h_k(\overbrace{1,1,\dots,1}^n) = \left(\binom{n}{k}\right)^k = \# \text{ } k\text{-multisets of } [n] = \{1, 2, \dots, n\}$

Recall that from "Stars and bars" we showed that  $\left(\binom{n}{k}\right)^k = \binom{n+k-1}{k}$  e.g.  $\{1, 1, 3, 4, 4\} \subseteq [5]$  multi-subset

$$\Rightarrow \begin{array}{ccccccccc} * & * & | & 1 & * & | & * & * & | \\ 1's & 2's & 3's & & 4's & & 5's & & \end{array}$$

In fact, we can similarly get the q-binomials;

DEF'N The q-binomial coefficient  $\begin{bmatrix} a+b \\ b \end{bmatrix}_q$  is the g.f.

$$\begin{bmatrix} a+b \\ b \end{bmatrix}_q := \sum_{\lambda \subseteq a \times b} q^{|\lambda|} \text{ of partitions in an } a \times b \text{ rectangle.}$$

e.g.  $\begin{bmatrix} 2+2 \\ 2 \end{bmatrix}_q = q^4 + q^3 + 2q^2 + q + 1$  since  $\emptyset, \begin{array}{|c|} \hline \square \\ \hline \end{array}, \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array}, \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}, \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array} \subseteq \begin{array}{|c|c|} \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \square & \square \\ \hline \end{array}$

We showed last semester that

$$\begin{bmatrix} a+b \\ b \end{bmatrix}_q = \frac{\begin{bmatrix} a+b \end{bmatrix}_q!}{\begin{bmatrix} a \end{bmatrix}_q! \begin{bmatrix} b \end{bmatrix}_q!}, \text{ where}$$

"q-number"

"q-factorial"

$$[n]_q = 1 + q + \dots + q^{n-1} = \frac{(1-q^n)}{(1-q)} \text{ and } [n]_q! = [n]_q \cdot [n-1]_q \dots [1]_q$$

e.g.  $\left[ \begin{smallmatrix} 2+2 \\ 2 \end{smallmatrix} \right]_q = \frac{[4]_q [3]_q [2]_q [1]_q}{[2]_q [1]_q [2]_q [1]_q} = (1+q+q^2) \cdot \frac{(1-q^4)}{(1-q^2)} = (1+q+q^2)(1+q^2)$

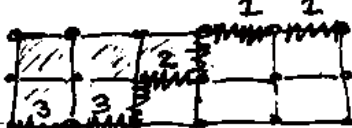
Thm (a)  $e_k(1, q, \dots, q^{n-1}) = q^{\binom{k}{2}} \cdot \left[ \begin{smallmatrix} n \\ k \end{smallmatrix} \right]_q$

(b)  $h_k(1, q, \dots, q^{n-1}) = \left[ \begin{smallmatrix} n+k-1 \\ k \end{smallmatrix} \right]_q$

Pf: We do (b) first. Observe that

$$h_k(1, q, \dots, q^{n-1}) = \sum_{1 \leq i_1 \leq i_2 \leq \dots \leq i_k \leq n} q^{\sum_{j=1}^k (i_j - 1)}$$

To any such k-multiset  $1 \leq i_1 \leq \dots \leq i_k \leq n$ , let's associate a partition  $\lambda$  inside the  $(n-1) \times k$  rectangle, as follows:

e.g.  $n=3, k=5, \lambda =$    $\leftrightarrow$  size 5 multiset  $\{1, 1, 2, 3, 3\} \subseteq [3]$

i.e., the values of the multiset tell us heights of horizontal steps on SE border of partition  $\lambda$  (where  $ht 1 = \text{top}$  and  $ht n = \text{bottom}$ )

Under this correspondence,  $|\lambda| = \sum_{j=1}^k (i_j - 1)$  (number of boxes).

So indeed  $h_k(1, q, \dots, q^{n-1}) = \left[ \begin{smallmatrix} n+k-1 \\ k \end{smallmatrix} \right]_q$  ✓

For (a): Similar. Can use trick of changing k-subset  $i_1 < i_2 < \dots < i_k$  to k-multisubset  $i_1 < i_2 - 1 < \dots < i_k - (k-1)$ . The difference  $1 + 2 + \dots + (k-1) = \binom{k}{2}$  explains factor of  $q^{\binom{k}{2}}$ .

Cor ("Principal specialization" of  $e_k$  and  $h_k$ )

(a)  $e_k(1, q, q^2, \dots) = q^{\binom{k}{2}} \cdot \prod_{i=1}^k (1 - q^i)$

(b)  $h_k(1, q, q^2, \dots) = \prod_{i=1}^k (1 - q^i)$

Pf: Note  $\lim_{n \rightarrow \infty} \left[ \begin{smallmatrix} n \\ k \end{smallmatrix} \right]_q = \lim_{n \rightarrow \infty} \frac{[n]_q!}{[n-k]_q! [k]_q!} = \lim_{n \rightarrow \infty} \frac{(1-q^n)(1-q^{n-1}) \dots (1-q^{n-k+1})}{(1-q^k)(1-q^{k-1}) \dots (1-q)} = \prod_{i=1}^k (1 - q^i)$  ■

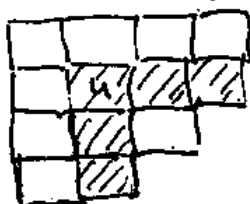
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## Principal specialization of Schur functions

We could ask about specialization of other sym-fn's, like  $p_k$ 's or  $m_\lambda$ 's. <sup>(HW #2)</sup> But now we discuss  $S_\lambda$ 's:

DEF'N Let  $\lambda$  be a partition, viewed as a Young diagram, and let  $u \in \lambda$  be a box of the Young diagram. The hook of  $u$  is all boxes below or to the right of  $u$ , together with  $u$  itself:

e.g.



boxes in  
hook = hook  $h(u) = 5$

The hook length  $h(u) := \# \text{ boxes in hook of } u$ .

Thm (Principal specialization of Schur function)


Let  $\lambda = (\lambda_1, \lambda_2, \dots, \lambda_k)$  be a partition. Then,

$$S_\lambda(1, q, q^2, \dots) = q^{b(\lambda)} \cdot \prod_{u \in \lambda} \frac{1}{1 - q^{h(u)}}$$

where  $b(\lambda) = 0 \cdot \lambda_1 + 1 \cdot \lambda_2 + 2 \cdot \lambda_3 + \dots = \sum_{i=1}^k (i-1) \cdot \lambda_i$ .

e.g.

$$e_k(1, q, q^2, \dots) = S_{(1^k)}(1, q, \dots) = q^{\binom{k}{2}} \cdot \frac{1}{(1-q^k)(1-q^{k-1}) \dots (1-q)}$$

since hook lengths are  for single column.

Similarly,  $h_k(1, q, \dots) = S_{(k)}(1, q, \dots) = \frac{1}{(1-q^k) \dots (1-q)}$

since  are hook lengths for single row.  
We saw these cases last class.

e.g. Let  $\lambda = (2, 1)$ . Then

$$S_\lambda(1, q, q^2, \dots) = \sum_{\text{SSYT } T, \text{sh}(T) = \lambda} q^{\text{sum of entries of } T - |\lambda|}$$

$$= q + 2q^2 + 3q^3 + 5q^4 + \dots$$

$$= q \cdot \frac{1}{(1-q)^2(1-q^2)} = q^{b(\lambda)} \cdot \prod_{u \in \lambda} \frac{1}{1-q^{h(u)}} \quad \text{since hook length}$$

In fact, even have a "finite version":

Thm (Stanley's hook-content formula)

$$S_\lambda(1, q, q^2, \dots, q^{n-1}) = q^{b(\lambda)} \cdot \prod_{u \in \lambda} \frac{1 - q^{c(u) + n}}{1 - q^{h(u)}}$$

where  $c(u) := j - i$  for box  $u = (i, j)$ .

As before can get principal specialization via limit  $n \rightarrow \infty$ .

Pf sketch: Starts w/ the "bialternant formula"

$$S_\lambda(x_1, x_2, \dots, x_n) = \frac{\det(x_j^{\lambda_i + n - i})}{\prod_{1 \leq i < j \leq n} (x_j - x_i)}, \text{ makes}$$

substitution  $x_i \rightarrow q^{i-1}$ , does some algebraic manipulations of the determinant. (See Stanley EC2).

3/21 We'd prefer a combinatorial proof, which we will give for the principal specialization. Starting point:

DEFN A reverse plane partition of shape  $\lambda$  is a filling of the boxes of  $\lambda$  with nonnegative integers that is weakly increasing in both rows and columns.

e.g. 

0	0	2
0	1	3
1	1	

 is an r.p.p. of  $\text{sh} = (3, 3, 2)$ .

Let  $RPP(\lambda) :=$  set of r.p.p.'s of  $sh = \lambda$ . There is a simple bijection  $\phi: SSYT(\lambda) \rightarrow RPP(\lambda)$  that subtracts  $i$  from all boxes in  $i$ th row;

e.g.  $\phi \left( \begin{array}{|c|c|c|} \hline 1 & 1 & 2 \\ \hline 2 & 3 & \\ \hline 3 & 5 & \\ \hline \end{array} \right) = \begin{array}{|c|c|c|} \hline 0 & 0 & 1 \\ \hline 0 & 1 & \\ \hline 0 & 2 & \\ \hline \end{array}$

Notice that via this bijection, sum of entries in  $T = b(\lambda) + |\lambda| + \text{sum of entries in } \phi(T)$ .  
So principal specialization of  $S_\lambda$  is equivalent to ...

Thm  $\sum_{\pi \in RPP(\lambda)} q^{|\pi|} = \prod_{u \in \lambda} \frac{1}{1 - q^{h(u)}}$

where  $|\pi| :=$  sum of entries of r.p.p.  $\pi$ .

We will explain a bijective proof of this thm.

To prove thm, it is enough to construct a bijection  $\phi: RPP(\lambda) \rightarrow \{ \text{arbitrary } \mathbb{N}\text{-fillings } A \text{ of boxes of } \lambda \}$

s.t.  $\text{sum of entries of } \pi = \sum_{u \in \lambda} A(u) \cdot h(u) \text{ for } A = \phi(\pi)$   
 $\Downarrow$   
 $\text{wt}(A)$

Why? Because then:

$$\sum_{\pi \in RPP(\lambda)} q^{|\pi|} = \sum_{A \text{ an } \mathbb{N}\text{-filling of } \lambda} q^{\text{wt}(A)} = \sum_A \prod_{u \in \lambda} q^{A(u) \cdot h(u)}$$

$$= \prod_{u \in \lambda} (1 + q^{h(u)} + q^{2 \cdot h(u)} + \dots) = \prod_{u \in \lambda} \frac{1}{1 - q^{h(u)}} \quad \checkmark$$

↑ ↑  
choose each value  $A(u)$  independently

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The bijection  $\phi$  is called the Hillman-Grassl algorithm. It is defined via a series of steps. We start off by writing our  $\Pi \in RPP(\lambda)$  next to the all 0's filling:

$RPP(\lambda) \ni \Pi = \begin{array}{|c|c|c|c|} \hline 1 & 2 & 2 & 2 \\ \hline 3 & 3 & 3 & \\ \hline 3 & & & \\ \hline \end{array}$   
 $\lambda = (4, 3, 1)$   
 $\Pi_0$

$A_0 = \begin{array}{|c|c|c|c|} \hline 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & \\ \hline 0 & & & \\ \hline \end{array}$

Then we find a path of boxes in  $\Pi (= \Pi_0)$  as follows:

- Start at northeastern most box  $(i, j)$  for which  $\Pi(i, j) \neq 0$ ,
- if we're at box  $(i, j)$ , then
  - move to  $(i, j-1)$  if  $\Pi(i, j) = \Pi(i, j-1)$
  - move to  $(i+1, j)$  otherwise
- repeat the previous until we exit  $\lambda$  (by leaving south of a column).

For example, w/ the above example we get this path

$\begin{array}{|c|c|c|c|} \hline 1 & 2 & 2 & 2 \\ \hline 3 & 3 & 3 & \\ \hline 3 & & & \\ \hline \end{array} \leftarrow \text{enter at row } 1 = i'$   
 $\downarrow \text{exit at column } 1 = j'$   
 We go left when value is same as where we are, otherwise go down.

Then we define  $\Pi_1$  by subtracting 1 from all boxes on the path, and we define  $A_1$  by adding 1 to  $A_0$  in position  $(i', j')$ , where  $i'$  = row we entered at, and  $j'$  = column we exited at.

Thus,  $\Pi_1 = \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 3 & \\ \hline 2 & & & \\ \hline \end{array}$   $A_1 = \begin{array}{|c|c|c|c|} \hline 1 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 \\ \hline 0 & & & \\ \hline \end{array} \leftarrow \text{added 1 to } (1, 1)$

Notice that the # of boxes in the path must be the same as the number of boxes in the hook of  $(i', j')$ :

$\begin{array}{ccc} \bullet & \bullet & \bullet \\ \bullet & & \\ \bullet & & \end{array} = \begin{array}{ccc} \bullet & \bullet & \bullet \\ \bullet & & \\ \bullet & & \end{array}$  because the path is a "ribbon"

So we have that  $|\Pi_0| - |\Pi_1| = \text{wt}(A_1) - \text{wt}(A_0)$ .

Then we repeat: find a path in  $\Pi_1$  using the same rules,

and define  $\pi_2$  and  $A_2$  from  $\pi_1$  and  $A_1$  in same way:

$$\begin{aligned} \pi_1 &= \begin{array}{|c|c|c|c|} \hline 1 & 1 & 1 & 1 \\ \hline 2 & 2 & 3 & \\ \hline 2 & & & \\ \hline \end{array} \leftarrow A_1 = \begin{array}{cccc} 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & & & \end{array} \Rightarrow \pi_2 = \begin{array}{|c|c|c|c|} \hline 0 & 0 & 0 & 0 \\ \hline 1 & 2 & 3 & \\ \hline & & & \\ \hline \end{array} \leftarrow A_2 = \begin{array}{cccc} 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & & & \end{array} \\ \Rightarrow \pi_3 = \begin{array}{|c|c|c|c|} \hline 0 & 0 & 0 & 0 \\ \hline 1 & 2 & 2 & \\ \hline & & & \\ \hline \end{array} \leftarrow A_3 = \begin{array}{cccc} 2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & & & \end{array} \Rightarrow \pi_4 = \begin{array}{|c|c|c|c|} \hline 0 & 0 & 0 & 0 \\ \hline 1 & 1 & 1 & \\ \hline 1 & & & \\ \hline \end{array} \leftarrow A_4 = \begin{array}{cccc} 2 & 0 & 0 & 0 \\ 0 & 1 & 1 & \\ 0 & & & \end{array} \\ \Rightarrow \pi_5 = \begin{array}{|c|c|c|c|} \hline 0 & 0 & 0 & 0 \\ \hline 0 & 0 & 0 & 0 \\ \hline 0 & & & \\ \hline \end{array} \leftarrow A_5 = \begin{array}{cccc} 2 & 0 & 0 & 0 \\ 0 & 1 & 1 & \\ 0 & & & \end{array} \end{aligned}$$

We stop when we reach  $\pi_k$  being all 0's. Then we set  $\phi(\pi) = A_k$ .

Note that  $|\pi| = (|\pi_0| - |\pi_1|) + (|\pi_1| - |\pi_2|) + \dots + (|\pi_{k-1}| - |\pi_k|)$   
 $= (wt(A_k) - wt(A_{k-1})) + \dots + (wt(A_1) - wt(A_0)) = wt(A_k) = wt(\pi)$

So indeed we defined map  $\phi: RPP(X) \rightarrow \{\text{NN-fillings } A\}$  which has the correct behavior on the weights of the fillings.

Need to check that  $\phi$  is a bijection. To do that, we will show that it is invertible, i.e., that we can undo the steps.

To explain inverse procedure:

1) Note that if we increment  $(i'_1, j'_1)$  before  $(i'_2, j'_2)$ , then either  $i'_1 < i'_2$ , or  $i'_1 = i'_2$  and  $j'_1 > j'_2$ .

↳ This tells us the reverse order to decrement values of  $A$  in the reverse inverse procedure

2) Show that we can build reverse of any path by entering at bottom of column  $j'$ , and moving right when entry to the right is same, otherwise move up (stopping when we reach right of row  $i'$ ).

For the details of this proof of bijectivity, see Sagan.

Main takeaway: we can "locally" reverse the steps.  $\square$