As you can see induction can be a very powerful tool for proving statements about positive integer n, especially formulas involving n. But often there formulas have to be "guessed" by finding patterns, so induction goes hand-in-hand with "guess and Jack."

Sometimes when proving things by induction it can be helpful to know that PCKI is true for all k< n to show P(n) is true (and not just Pen-11).

The Strong Principle of Induction says that if:

- · P(no) is true for some no (base case)

 · P(n) is true whenever P(k) is true
 for an no≤ k < n, for all n ≥ no,

 then P(n) is true for all n ≥ no.

 (Notice also now we allow no to be different from I).
- E.g. Thm Using 2¢ and 5¢ stamps, we can make any amount n¢ for all n = 4.

Here's another example at strong induction: The Fibonacci numbers For for n21 one defined by F, = 1 and Fz=1 and Fn = Fn-, + Fn-z for n>2. $E_{.7}$. $F_3 = F_{1+}F_{2} = 1+1=2$ $F_4 = F_{2+}F_{3} = 1+2=3$ $F_5 = F_{3+}F_{4} = 2+3=5$... Inm Fn ≤ 2" for all n≥1. Pf: We use strong induction. We have two base Cases: n=1 no $F_1 = 1 \le 2^\circ = 2^{i-1}$ / n=2 $-> F_2 = 1 \le 2^i = 2^{2^{-i}}$ Now, for 1/2, assume that Fin \$2 n-2 and Fn-2 \le 2n-3 (using strong induction). Thus, $F_n = F_{n-2} + F_{n-1}$ $\leq 2^{n-3} + 2^{n-2}$ (by induction) $\leq 2^{n-2} + 2^{n-2} = 2 \cdot (2^{n-2}) = 2^{n-1}$ and so by induction we are done! The strong torm of induction is closely related to the well-ordering property for the nannagrative Integers, which say, that every nonempty set of ronnegative numbers has a min; mum. You can see in the book how the W.O. A. can be used to show we alway get a well-defined quotrent and remainder whom doing long division.

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Basic Mathematical Structures: Functions \$3.

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Having concluded our study of proofs (Chapter 2), where starting a new chapter, Chapter 3, which discusses basic mathematical structures. The most basic mathematical structures are sets, which we have already discussed in Chapter 1. The next most basic structures in math are functions which are procedures for going from one set to another,

There are many ways to think about functions.

One is that a function of from a set X to a set Y is a machine or a rule that takes something in X and Spits out Something in Y:

For example, consider the following procedure:

given a 10-digit number x like x = 1043213598

we sum together all the digits: 1+0+4+3+2+1+3+5+9+8=36

and then spit out" the ones digit of the resulting sum as our y = f(x) (here y = 6 in the example).

This describes a function of whose domain X is the set of 10-digit numbers and codomain Y is the set of one digit numbers.

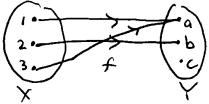
(This is a simplification of the "check sum" procedure for Credit courd numbers; the book describes the real procedure, which is more complicated.),

Wotrze that the domain of function f is the name we give to the input set X and codomain is the name we give to the output set Y.]

That was an intuitive definition of function as madrine. The formal definition of function user ordered pairs.

Defin A function of from set X (called the domain) to set Y (called the codomain) is a subset of X x Y (set of ordered poins (x,y) w x \in X, y \in Y), Such that: for every x \in X, there is a unique y \in Y with (x,y) in our sexbet.

E.g. We often represent functions by arrow dagrams.



This corresponds to the subset $\{(1,a), (2,b), (3,c)\}$ of $X \times Y$ with $X = \{1,2,3\}$ and $Y = \{4,b,c\}$.

Notice how for every $x \notin X$ there is a unique $y \notin Y$ with (x,y) in our subset; we write $(x,y) \notin Y$ thus $(x,y) \notin Y$ thus $(x,y) \notin Y$ the function $(x,y) \notin Y$ is named $(x,y) \notin Y$.

The function $(x,y) \notin Y$ is named $(x,y) \notin Y$.

In our credit card checksum example function we had f(1043213598) = 6, and more generally the ordered pairs in our subset will always be $\{X(X, f(X))\}$ so we use "f" as shorthard for the function.

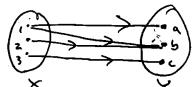
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The set $\{f(x): x \in X\}$ of values our function f actually takes on is called the range of f, and it is a subset of the codomain: e.g. in the arrow diagram example f above the codomain was $Y = \{A, b, c\}$, but the range is $\{a,b\}$ since there is no $x \notin X \bowtie f(x) = C$.

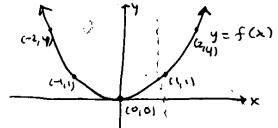
We also write $f: X \rightarrow Y$ to mean f is a function from X to Y. The arrow helps you remainly what it does: it takes something in X to something in Y. A diagram like:



is not the arrow diagram of a function because the key property of a function is that for every x EX there is a unique f(x) = y EY it is "sent to" and here I is "sent" to both a and b! From calculus you are probably used to function like $f(x) = x^2$

Whose domain and rodomain are the real numbers \mathbb{R} . Notice how ' $f(x) = \chi^2$ " is the "rule/machine" description of the function: it tells us for a given input χ how to produce the output of the function; $f(3) = 3^2 = 3 \times 3 = 9$.

But we can also represent a function f: R-> R by its graph like we are used to doing!

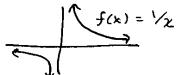


The graph of a function fire>IR is just a drawing of all the points {(x,f(x)): x f IR3, i.e., it is another virual representation of the ordered pairs definition of function. (Recall: "Vertical line test")

Some functions defined algebraically like

f(x) = 1/2.

have domains that are strict subsets of TR:



Here the domain (atnot range) of f(x) = 1/2 is {x ETR: x 703 since we are not allowed to divide by zero.

More about functions: Let f: X -> Y be a function. Defin The function fis called one-to-one if there are not two different $\chi_{i}, \chi_{i} \in X$ with $f(\chi_{i}) = f(\chi_{i})$. "Everything in X is sent to a different thing in Y" Eig. Let $f: \mathbb{Z} \to \mathbb{Z}$ be given by f(n) = 2n+1. [so f(0) = 1, f(1) = 3, f(-1) = -3.] This + is one-to-one since if Zuiti=Zhztl Defin The function f is called onto (or surjective) if for every y & Y, there is some x Ex with f(x) = y. "Everything in Y is mapped to by something in X' Onto Same as: range = codomain. not on to Eig. 1f f. 2-72 is given by f(n) = 2n+1 1 then f is not onto since there is no n = Z $\omega/2n+1=0$ for any even number).

(or a bijection) Defin The function f is called bijective if it is both one-to-one and onto. (a) is bijective. E.9. tig. f(n) = n+1: 2 -> Z is a bijection. (why?) Exercise If fix-> Y is a bijection between finite sets X and Y, then # X = #Y (the sets-have same # of elements). Defin If $f: X \to Y$ is a bijection, then we define its inverse function $f': Y \to X$ by f''(y) = X if and only if f(x) = yE.g. To check that a function fiR-DR is one-to-one we have the "nor. Funtal line test" $\int_{0}^{\infty} f(x) = x^{2}$ not one—to— The inverse of $f(x) = x^3$ is $f^{-1}(x) = \sqrt[3]{x}$ 17500=x3 (0

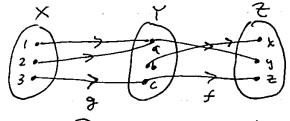
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The inverse function f' "undoes" whatever f does. Let's make this precise by talking about composition

Def'n Let $g: X \to Y$ and $f: Y \to Z$ be two functions. Their composition (fog): $X \to Z$ is defined by $(f \circ g)(x) = f(g(x))$ for all $x \in X$.

"Dog first and then do f to what you get."

E.p.



"Combine arrow diagrams"

to form

arrow dayson for compository

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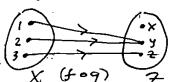


Fig. If $f(x) = 2^x : \mathbb{R} \rightarrow \mathbb{R}$ and $g(x) = x^3 : \mathbb{R} \rightarrow \mathbb{R}$ Then $(f \circ g)(x) = 2^{x^3}$ and $(g \circ f)(x) = (2^x)^3 = 2^{3x}$. Notice how $(f \circ g) \neq (g \circ f)!$ order matters.

Defin The identity function $Fd: X \to X$ on a set X is function with $Fd_X(X) = X$ $\forall X \in X$.

"The identity function 'does nothing': gives input "

If $f: X \to Y$ is a bijection. Then

($f' \circ f$) = Id_X identity.

Since of (f(x)) = x Vx fx. This is sense in which inverse function "undoes" original.

Modular arithmetic functions

Many of the functions you're familiar with from calculu, especially poincar functions like f(x) = 5x - 2 and polynomials like $f(x) = 3x^3 - 2x^2 + 4x - 1$ are important in discrete math too...

The "modulo n" function is another function that's very important in discrete math, and may be new to you.

Des'n For any integer $m \in \mathbb{Z}$, $m \mod n$ ("m modulo n") is the unique $r \in \{0,1,2,...,n-1\}$ such that r is the remainder when dividing m by n, i.e. $\exists k \in \mathbb{Z}$ such that $m = k \cdot n + r$.

E.g. 3 mod 5=3 and 8 mod 5=3 too since 8=5+3. 1247 mod 10 is 7 since we just look at 15 place. For any n, n mod n=0, and -1 mod n=n-1.

In this way, for every positive integer n we eget a function if: Z:>Z given by f(m) = m mod n.

(The range of fir £0,1,..., n-13 so we could take the codomain to be £0,1,..., n-13 instead...)

The mod n functions can be useful for clock or calendar problems, e.s.

Exercise If the first day of the year is a Tuesday, what day of the week is the 100th day of the year?

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1 10/26 Sequences § 3.2 A sequence is a list of things, like: 1,2,3,4,5,... 2, 4, 8, 16, 32, ... 1,2,3 etc. b, a, n,a, n, a It can be finitely long, or infinidely long. It can have repetations (like in the letters of banana) The important thing is that the order of a sequence matters, so 1,2,3 \$\frac{1}{2},3,1,2. Formally, we can a sequence by a finction s whose domain is a subset of the positive integers, (which we denote by Zzo or Zt) E13. S: 21,07 Z with s(n) = n gives the sequence 1,2,3,4, ---B.7.8:2/2042 with scn) = 2 gives sequence 2, 4, 8, 16, ... E.g. S. &1,2,3,4,5,6}-> {a,b,n} with S(1)=b, s(2)=a, s(3)=n, s(4)=a, s(5)=n, s(6)=a gives the sequence b, a, na, n, a Usually the domain is either all of Zoo (for an infinite Sequence) or [1,2,3,4,..., n] (for a finite sequence) We worde the sequence as S1, S2, S3, ... where S: = S(i) is 'sequence notation' We also sometimes write it as I sn 3 n=1

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If the codomain of the sequences is a set of numbers, we says is increasing if siks; when ikj and says is decreasing if Siss, when icj, Eig 2,4,8,16,32,... is increasing We defre nonincrersing sequences (w/ S; ≥ Sj) and nondereasing seq. (w/ SiESj) Similarly. If we have a finite sequence Esn3n=1, we define its sum \sum sn = S1 + S2 + ... + SK. E.g. We already saw (using induction) that $2^{\circ}+2^{\prime}+2^{2}+\cdots+2^{k}=\sum_{k=1}^{k+1}2^{k+1}=2^{k+1}-1$ Can define product It sn = s, x szx. xsk as well. If sis a sequence, a sinto sequence of sis a Sequence we get by selecting some of the items of the 15st s (not recessor by consecutive) in the same order) Fig. b, a is a subsequence of b, a, n, a, n,a, as is b, n and a, a, a and n, n,

but a, b is not a subsequence of b, a, n, a, n, a

If the sequence is $\xi S_n \xi$ then the subsequence usin be $S_{n_1}, S_{n_2}, S_{n_3}, \ldots$ where $\{n_1 < n_2 < \ldots \}$ is a subset of the domain of S.

E.g. 2, 4, 6, 8, ... is a subsequence of 1,7,3, 4,5,...

Strings § 3.2 If X is a finishe set, then a String over X is any finishe sequence of elements from X. We use X* to denote all Strings over X.

Eig. If X= {a,b} then some elements of X* are a,b, aa, ab, bba, baba, etc.

Another string that's always in Xthis the null Strong denoted &, that doesn't have any letters.

If X, B + X* are two strings, their concatenation XB is what we get by patting X right before B:

Eg. X = aba, B=bba, then &B = ababba.

Notice that the length (# of letters in) & B 13 the sum of length of & and length of B.

Finally, a substring of dEX* is a string of consecutive letters from a.

E.g. d=aba then ab and ba are substrings, but ba is not.

Exercise Show B is a substitute of X if and only if x = 8p% for some strings 8 and 8.