

## Weighted Polya counting, a.k.a. Polya-Redfield enumeration

Notice that the example of

$G = \langle \sigma = (1, 2, 3, 4) \rangle \curvearrowright$  colorings of  $[4]$   
w/ 2 colors (e.g. black+white)

is the same as

$G \curvearrowright$  subsets of  $[4]$   
(Just ~~determine~~ use subset of black vertices.)

But we also know that an action on subsets like this preserves size of subset; e.g. we looked at

$G \curvearrowright$  size 2-subsets of  $[4]$

which in the language of colorings would be same as

$G \curvearrowright$  colorings of  $[4]$  w/ 2 vertices white  
+ 2 vertices black.

In general we might want to keep track of precise number of each color used, as in:

Q: Up to rotation, how many colorings of vertices of square ~~have~~ exactly 2 red, 1 blue, 1 green vertex?

2/1 To answer this question, we need more notation.

for a coloring  $f: X \rightarrow Y = \{1, 2, \dots, k\}$  define

monomial  $\vec{y}^f := \prod_{x \in X} y_{f(x)} \in \mathbb{C}[y_1, y_2, \dots, y_k]$

e.g. ~~coloring~~  $R - R$  if we decide  
Coloring  $f = \begin{cases} 1 & R \\ 1 & G \\ 2 & B \end{cases} \rightsquigarrow y_1^2 y_2 y_3 \quad \begin{cases} R=1 \\ B=2 \\ G=3 \end{cases}$

Notice: If  $G \curvearrowright X$  then  $\vec{y}_f = \vec{y}_{g \cdot f} \quad \forall g \in G$ .

DEF'N Let  $G \curvearrowright X$  and hence on colorings  $Y^X$ .  
 The pattern inventory polynomial of  $G \curvearrowright Y^X$  is

$$P(y_1, y_2, \dots, y_n) = \sum_{\mathcal{O}} \bar{y}^{\mathcal{O}} \in \mathbb{C}[y_1, \dots, y_n]$$

where the sum is over all orbits  $\mathcal{O}$  of  $G \curvearrowright Y^X$   
 and  $\bar{y}^{\mathcal{O}} := \bar{y}^f$  for any coloring  $f \in \mathcal{O}$ .

e.g. For  $G = \langle \sigma = (1, 2, 3, 4) \rangle \curvearrowright X = [4]$  and  $Y = \{0, 1, 2\} = \{1, 2\}$ ,

$$P = 1y_1^4 + 1y_1^3y_2 + 2y_1^2y_2^2 + 1y_1y_2^3 + 1y_2^4$$

Given the pattern inventory poly.  $P$ , we can then answer  
 questions like: "how many (symmetry classes of) colorings  
 use 2 white + 2 black vertices?" by extracting coefficients.

So our goal will now be to give a formula for  $P(y_1, \dots, y_n)$ .

To do that, we need to keep track of more refined  
cycle information of elements  $g: X \rightarrow X, g \in G$ .

Set  $c_i(g) := \# \text{ } i\text{-cycles of permutation } g: X \rightarrow X$ .

DEF'N The cycle index polynomial of  $G \curvearrowright X$  is

$$\mathcal{Z}_G(t_1, t_2, \dots, t_n) = \frac{1}{\#G} \sum_{g \in G} \prod_{i=1}^n t_i^{c_i(g)} \in \mathbb{C}[t_1, \dots, t_n]$$

This is the key to Poly-a counting!

e.g. With  $G = \langle \sigma = (1, 2, 3, 4) \rangle \curvearrowright X = [4]$ , have

$$Z_G = \frac{1}{4} \left( \underbrace{t_1^4}_{\sigma = (1)(2)(3)(4)} + \underbrace{2t_4}_{\sigma = (1, 2, 3, 4)} + \underbrace{t_2^2}_{\sigma^2 = (1, 3)(2, 4)} \right)$$

Thm (Pólya-Redfield enumeration theorem)

The pattern inventory polynomial of  $G \curvearrowright Y^X$  is

$$P = Z_G \left( \sum_{i \in Y} y_i, \underbrace{\sum_{i \in Y} y_i^2}_{t_1}, \underbrace{\sum_{i \in Y} y_i^3}_{t_2}, \dots, \underbrace{\sum_{i \in Y} y_i^n}_{t_n} \right).$$

e.g. Let  $G = \langle \sigma = (1, 2, 3, 4) \rangle \curvearrowright X = [4]$  and consider set of colors  $\mathcal{Y} = \{R, G, B\} = \{1, 2, 3\}$ .

$$\text{Then } P = \frac{1}{4} (1y_1 + y_2 + y_3)^4 + 2(y_1^4 + y_2^4 + y_3^4) + (y_1^2 + y_2^2 + y_3^2)$$

$$= \dots = y_1^4 + y_2^4 + y_3^4 + y_1^3 y_2 + y_1^3 y_3 + y_2^3 y_1 + y_2^3 y_3 \\ \xrightarrow{\text{lots of algebra!}} + y_3^3 y_1 + y_3^3 y_2 + 2(y_1^2 y_2^2 + 2y_1^2 y_3^2 + 2y_2^2 y_3^2 + 3y_1^2 y_2 y_3 + 3y_2^2 y_1 y_3 + 3y_3^2 y_1 y_2).$$

To figure out how many colorings have 2 R, 1 B, 1 G, we extract coeff. of  $y_1^2 y_2 y_3$  from  $P$ :

$$\text{A: } [y_1^2 y_2 y_3] P = 3 \text{ colorings w/ 2R, 1B, 1G.}$$

Note: Setting  $y_i = 1$  for all  $i \in Y$ , we recover the unweighted Pólya counting formula for total number of colorings (ignoring patterns).

2/9

### Pf of Polya-Redfield Thm:

The proof is very similar to unweighted result; we just need to make sure we keep track of weights.

- First observe that for any orbit  $O$  of  $G \curvearrowright Y^X$ :

$$\vec{y}^O = \sum_{f \in O} \vec{y}^f / \#O_f, \text{ so}$$

$$P = \sum_O \vec{y}^O = \sum_{f \in Y^X} \frac{\vec{y}^f}{\#O_f} = \frac{1}{\#G} \sum_{f \in Y^X} \#G_f \cdot \vec{y}^f,$$

whereas before we used the Orbit-Stabilizer Thm.

By the same "summing over rows" vs. "summing over columns" trick, applied to matrix

$$M_{(g, f)} = \begin{cases} \vec{y}^f & \text{if } g \cdot f = f \\ 0 & \text{otherwise} \end{cases}, \text{ get that}$$

$$P = \sum_O \vec{y}^O = \frac{1}{\#G} \sum_{g \in G} \sum_{f \in (Y^X)^g} \vec{y}^f$$

- So we again need to think about  $(Y^X)^g$  for  $g \in G$ .

Recall that  $f \in (Y^X)^g \iff f(x) = f(x')$  whenever

$x, x'$  belong to same cycle of  $g: X \rightarrow X$

$$\text{e.g. } g = (\underbrace{x_1, x_2, x_3}_{\substack{\text{color all red} \\ \text{or all blue} \\ \text{or all green}}}) \underbrace{(x_4)}_{\substack{\text{color red} \\ \text{or green}}} \underbrace{(x_5, x_6)}_{\substack{\text{color both red} \\ \text{or both blue} \\ \text{or both green}}} : X \rightarrow X$$

$$\sum_{f \in (Y^X)^g} \vec{y}^f = (y_1^3 + y_2^3 + y_3^3) \cdot (y_4 + y_5 + y_6) \cdot (y_1^2 + y_2^2 + y_3^2)$$

$$\text{So in general } \sum_{f \in (Y^X)^g} \vec{y}^f = \prod_{\substack{\text{cycles} \\ \text{of } g: X \rightarrow X}} \sum_{y \in Y} y^{\text{size of } c}$$

This precisely means  $P(y_1, \dots, y_K) = Z_G(\sum_i y_i, \sum_i y_i^2, \dots, \sum_i y_i^n)$ .

Cor Let  $G \curvearrowright X$ . Then  $\sum_{k=0}^n \#(\text{orbits of } G \curvearrowright \text{size } k \text{ subsets of } X) \cdot t^k = Z_G(1+t, 1+t^2, \dots, 1+t^n)$ .

Pf: Use 2 colors in weighted polygon counting.  $\blacksquare$

2/9 e.g. Recall that a graph consists of a vertex set  $V$  and a set of edges  $E$ , unordered pairs of vertices.

$$G = \begin{array}{c} 2 \\ \bullet \\ \swarrow \quad \searrow \\ 1 \quad 3 \\ \downarrow \quad \downarrow \\ 4 \quad 5 \end{array} \Rightarrow V = [5], E = \{\{1, 2\}, \{2, 3\}, \{4, 5\}\}.$$

Q: How many graphs  $G$  with vertex set  $V = [n]$ ?

A:  $2^{\binom{n}{2}}$  since there are  $\binom{n}{2}$  possible edges, and we can choose any subset of edges.

But... what if we want to count unlabeled graphs, i.e., graphs up to isomorphism?

DEF'N An isomorphism between graphs  $G = (V, E)$  and  $G' = (V', E')$  is a bijection  $\phi: V \rightarrow V'$  on vertices s.t.  $\{i, j\} \in E \Leftrightarrow \{\phi(i), \phi(j)\} \in E'$ .

$$\begin{array}{c} 2 \\ \bullet \\ \swarrow \quad \searrow \\ 1 \quad 3 \end{array} = G \sim \begin{array}{c} 2 \\ \bullet \\ \swarrow \quad \searrow \\ 1 \quad 3 \end{array} = G'$$

Q: How many isomorphism classes of graphs w/  $n$  vertices are there?

Even better, what is  $\sum_{G, \text{ graph on } n \text{ vertices}} t^{\#\text{edges}(G)}$ ?

A: By weighted Polya counting, answer is

$$Z_G(1+t, 1+t^2, \dots, \cancel{t^{n-1}}, 1+t^{\binom{n}{2}})$$

where  $G = S_n \curvearrowright X = \{\text{size 2 subsets of } [n]\}$

$\leftarrow$  WARNING! Not  $X = [n]$

Let's first consider case  $n=3$ :

cycle type	$\sigma \in S_3$ w/ this type	cycle structure of $\sigma: X \rightarrow X$	monomial $\sum \sigma$	# $\sigma \in S_3$ w/ type $\lambda$
$(1, 1, 1)$	$e = (1)(2)(3)$	$(\{1, 2\}, \{1, 3\}, \{2, 3\})$	$t_1^3$	1
$(2, 1)$	$(1, 2)(3)$	$(\{1, 3\}, \{2, 3\}, \{1, 2\})$	$t_2 t_1$	3
$(3)$	$(1, 2, 3)$	$(\{1, 2\}, \{1, 3\}, \{2, 3\})$	$t_3$	2

$$\text{So } Z_G(t_1, t_2, t_3) = \frac{1}{3!} (t_1^3 + 3t_2 t_1 + 2t_3)$$

$$\begin{aligned} \text{and } Z_G(1+t, 1+t^2, 1+t^3) &= \frac{1}{6} ((1+t)^3 + 3(1+t^2)(1+t) + 2(1+t^3)) \\ &= t^3 + t^2 + t + 1 \end{aligned} \quad \leftarrow \text{g.f. of graphs on } n=3 \text{ vertices, by #edges.}$$

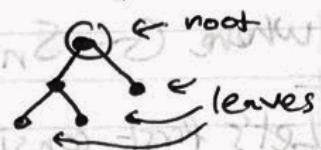
$n=4: \lambda$	$\sigma \in S_4$	cycle structure of $\sigma: X \rightarrow X$	$\sum \sigma$	# $\sigma \in S_4$
$(1, 1, 1, 1)$	$e = (1)(2)(3)(4)$	$(\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\})$	$t_1^6$	1
$(2, 1, 1)$	$(1, 2)(3)(4)$	$(\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\})$	$t_2^2 t_1^2$	$\binom{4}{2} = 6$
$(2, 2)$	$(1, 2)(3, 4)$	$(\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\})$	$t_2^2 t_1^2$	$\binom{4}{2}/2 = 3$
$(3, 1)$	$(1, 2, 3)(4)$	$(\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\})$	$t_3^2$	$4 \cdot 2 = 8$
$(4)$	$(1, 2, 3, 4)$	$(\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3\}, \{2, 4\}, \{3, 4\})$	$t_4 t_2$	$3! = 6$

$$\text{So } Z_G(t_1, t_2, t_3, t_4) = \frac{1}{4!} (t_1^6 + 6t_2^2 t_1^2 + 8t_3^2 + 6t_4 t_2)$$

$$\begin{aligned} \text{and } Z_G(1+t, 1+t^2, 1+t^3, 1+t^4) &= \frac{1}{24} ((1+t)^6 + 6(1+t^2)^2 (1+t)^2 \\ &\quad + 8(1+t^3)^2 + 6(1+t^4)(1+t^2), \\ &= t^6 + t^5 + 2t^4 + 3t^3 + 2t^2 + t + 1 \end{aligned} \quad \leftarrow \begin{array}{l} \text{g.f. for} \\ \text{graphs} \\ \text{on 4 vertices} \end{array}$$

Cultural aside on trees: Polya developed Polya counting to enumerate trees, motivated by problems in molecular chemistry!

A rooted binary tree looks like:  
each node has 2 or 0 children

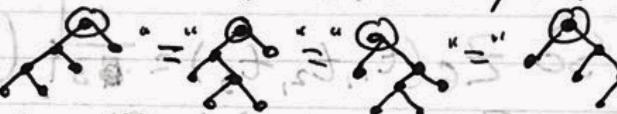


# rooted binary trees w/  $n+1$  leaves = Catalan number  $C_n = \frac{1}{n+1} \binom{2n}{n}$  & last sentence

$$\text{e.g. } C_3 = 5 \text{ and}$$



But ... what is we wanted to count "structurally different" binary trees, i.e.)



Let  $a_n := \#$  structurally different rooted binary trees w/  $n+1$  leaves

$$\begin{aligned} n &= 0, 1, 2, 3, 4, 5, \dots \\ C_n &= 1, 1, 2, 5, 14, 42, \dots \\ a_n &= 1, 1, 1, 2, 3, 6, \dots \end{aligned}$$

$$\text{Set } C(x) := \sum_{n \geq 0} C_n x^n \text{ and } A(x) := \sum_{n \geq 0} a_n x^n$$

$$\text{We saw } C(x) = 1 + x \cdot C(x)^2 \text{ account for symmetry!}$$

Polya counting  $\Rightarrow A(x) = 1 + \frac{x}{2} (A(x)^2 + A(x^3))$

Even leads to asymptotics for all trees!

Thm (Otter, 1948) Let  $t_n := \#$  unlabeled, unrooted trees on  $n$  vertices

$$\text{Then } t_n \sim C x^{n-5/2} \text{ w/ } \alpha \approx 2.955.$$

$$C \approx 0.5349 \dots$$

$$\text{Compare: } n^{n-2} / n! \sim \frac{1}{\sqrt{2\pi}} e^{-n} n^{-5/2} \approx 2.71 \dots$$

Cayley's formula for ~~labeled~~ labeled trees!