

Math 4707: The Tutte polynomial

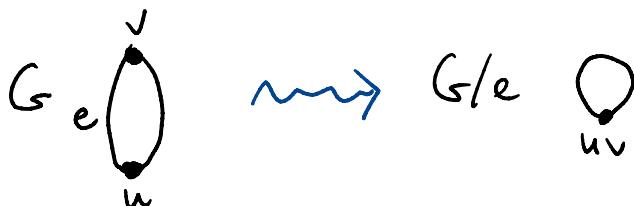
Reminder: • Final due in 1 week on Wed., 5/5.

In the past two classes we've seen **deletion-contraction** is a very powerful tool for understanding graphs. Today we'll see "how far" you can go w/ deletion-contraction. It turns out there is a universal deletion-contraction invariant, called the **Tutte polynomial**.

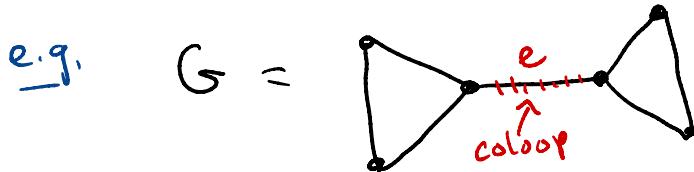
Before we define Tutte polynomial, need to slightly modify our definition of **contraction**. Previously we were always working w/ **simple graphs**, but now it's important to allow multiple edges and loops. So now when we contract an edge e , we do not remove multiple edges:



And if we contract a multiple edge, we create loops:



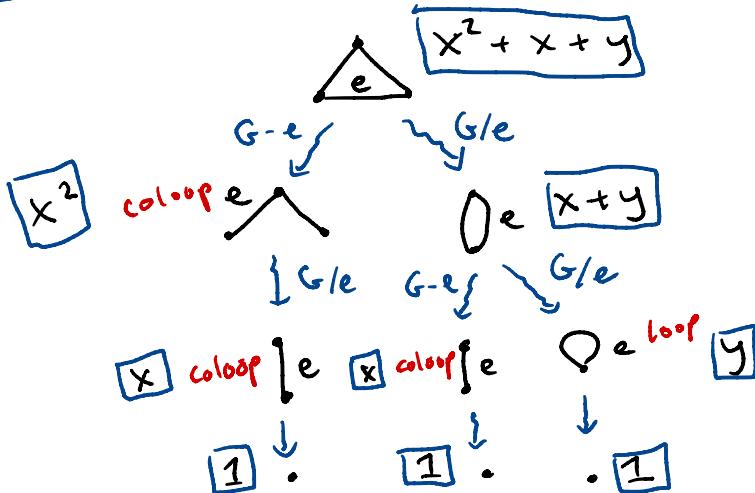
We need one more notion to define Tutte poly.:
 an edge e of G is a **coloop** a.k.a. **isthmus** if
deleting e increases the # of connected components:



Def'n The **Tutte polynomial** of G , denoted $T_G(x, y)$, is the unique polynomial in two variables x and y s.t.:

- $T_G(x, y) = T_{G-e}(x, y) + T_{G/e}(x, y)$ for any edge e which is not a loop or coloop.
- $T_G(x, y) = x T_{G/e}(x, y)$ if e is a coloop.
- $T_G(x, y) = y T_{G-e}(x, y)$ if e is a loop.
- $T_G(x, y) = 1$ if G has no edges.

e.g. to compute $T_G(x, y)$ for $G=K_3$ triangle:



Here we write

$$T_G(x, y)$$

next to
graph G

Rmk: Implicit in def'n of $T_G(x, y)$ is the fact that it does not matter which order we delete/contract the edges in. Actually, this is a **Theorem**, which can be Proved using a different def'n of $T_G(x, y)$.

~ Essentially by definition, the Tutte polynomial is the **universal deletion-contraction invariants**: we get others via **specialization**. For instance ...

Prop. The **chromatic polynomial** of G is

$$\chi(G, k) = (-1)^{\#V - c(G)} k^{c(G)} T_G(1-k, 0)$$

where $\#V = \#\text{vert's of } G$ and $c(G) = \#\text{ components of } G$.

~ e.g. For $G = K_3$ we saw before that

$$\chi(G, k) = k(k-1)(k-2)$$

and $(-1)^{3-1} k^1 T_G(1-k, 0) = k((1-k)^2 + (1-k) + 0) = k(k-1)(k-2) \checkmark$

~ The proof of prop. is same as we have seen before: show they satisfy same recurrence.

Similarly, from last class we get ...

Prop. $T_G(2, \sigma) = \#$ acyclic orientations of G .

Another very important specialization of Tutte poly. is $x=y=1$.

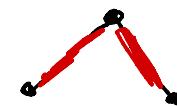
Prop. For a connected graph G ,

$T_G(1, 1) = \#$ spanning trees of G .

Pf. Suffices to show recurrence

$\#\text{span. tree}(G) = \#\text{span. tree}(G-e) + \#\text{span. tree}(G/e)$
for a non-loop/cloop e . We explained why this
is many classes ago... there are bijections:

Spanning trees
of G w/out e \iff Spanning trees
of $G-e$



Spanning trees
of G w/ e \iff Spanning trees
of G/e



E.g. $\#\text{span. tree}(K_3) = 3 = (x^2 + x + y)|_{x=y=1}$.



Let's finish w/ a cool symmetry of $T_G(x, y)$:

Thm Suppose G is planar graph, G^* its **dual**.
Then $T_G(x, y) = T_{G^*}(y, x)$.

In other words, planar duality swaps x and y !
Basic idea of pf:

deletion of $e \Leftrightarrow$ contraction of e^*
in G in G^*

e is loop in $G \Leftrightarrow e^*$ is coloop in G^* .

Rmk: We've already seen $\#\text{span.tree}(G) = \#\text{sp.tree}(G^*)$.

Since $T_G(1, 1) = T_{G^*}(1, 1)$, the tutte poly. gives another proof of this.

e.g. for $G = \Delta$, $G^* = \emptyset$ and $T_{G^*}(x, y)$ is ...

Here we write

$T_G(x, y)$

next to
graph G

$$\begin{array}{c} \emptyset e \quad y^2 + y + x = T_G(y, x) \\ \downarrow G-e \quad \downarrow G/e \\ x+y \quad e \emptyset \quad \infty e \quad y^2 \\ \downarrow G-e \quad \downarrow G/e \quad \downarrow G-e \quad \downarrow G-e \\ \times \quad \text{coloop} \quad \text{loop} \quad \text{loop} \quad \text{loop} \\ \downarrow G/e \quad \downarrow G/e \quad \downarrow G-e \quad \downarrow G-e \\ 1 \quad \cdot \quad 1 \quad \cdot \quad \cdot \quad 1 \end{array}$$

Now let's take a 5 min. break...

and when we come back we can do the final worksheet of the semester, on the Tutte poly., in breakout groups.