

Samuel Lee
ID# 20412795

Question #1

$$T(n) = T\left(\frac{7n}{10}\right) + T\left(\frac{1n}{4}\right) + n$$
$$\alpha = 1, \beta = 0$$
$$\left(\frac{7}{10}\right)^r + \left(\frac{1}{4}\right)^r = 1$$
$$\left(\frac{38}{40}\right)^r = 1 \therefore r < 1 \therefore r < \alpha \therefore \Theta(n^1 \log^0 n)$$
$$= \Theta(n)$$

Question #2

$$T(n) = 3T\left(\frac{n}{4}\right) + T\left(\frac{n}{4}\right) + n$$
$$\alpha = 1, \beta = 0$$
$$3\left(\frac{1}{4}\right)^r + \left(\frac{1}{4}\right)^r = 1$$
$$\left(\frac{1}{4}\right)^r = 1 \therefore \Theta(n^1 \log^{B+1} n) = \Theta(n \log n)$$

Question 3

$$\text{if } n=1, A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$

assume given eq. for $n=k$ is true

$$A^k = \begin{pmatrix} F_{k+1} & F_k \\ F_k & F_{k-1} \end{pmatrix}$$

$$\text{to prove } A^{k+1} = \begin{pmatrix} F_{k+2} & F_{k+1} \\ F_{k+1} & F_k \end{pmatrix}$$

$$\begin{aligned} A^{k+1} &= A \cdot A^k = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F_{k+1} & F_k \\ F_k & F_{k-1} \end{pmatrix} = \begin{pmatrix} F_{k+1} + F_k & F_k + F_{k-1} \\ F_{k+1} & F_k \end{pmatrix} \\ &= \begin{pmatrix} F_{k+2} & F_{k+1} \\ F_{k+1} & F_k \end{pmatrix} = A \cdot A^k = A^{k+1} \end{aligned}$$

$$\therefore \text{by strong induction if } A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, A^n = \begin{pmatrix} F_{n+1} & F_n \\ F_n & F_{n-1} \end{pmatrix}$$

Question 4:

Fibonacci(matrix An, int n)

int A[2][2] = {1,1,1,0};

if (n==0)

return An[0][1];

else

An *= A;

Fibonacci(An, n-1)

Question 5

input (int n , A)

priorityQueue q;

for ($i=1$ to n), $O(\sqrt{\log n})$] $n \cdot \sqrt{\log n}$

for ($i=1$ to n), $O(\sqrt{\log n})$] $n \cdot \sqrt{\log n}$

return A

fastest time is $O(n \sqrt{\log n})$ not $O(n \log n)$

\therefore lying because $O(n \log n)$ is faster