# Introduction To Relational Data Models Lecture 3, Chapter 2

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#### The goal of the next several lectures is to

- 1. develop a method of representing data as sets of tuples;
- 2. create a mathematical language which will allow us to query these sets;
- 3. translate this framework into languages, algorithms, and data-structures which can actually be implemented in software.

Recall from Lecture 1 that we wish to store and query "instances" of "structures". For example,

```
struct Movie {
    char title [128];
    int year;
    int length;
    char genre [128];
Movie m1 = \{
    "Gone With the Wind", 1939, 231, "drama"
};
Movie m2 = { "Star Wars", 1977, 124, "scifi" };
```

Recall from Lecture 1 that we wish to store and query "instances" of "structures". For example,

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struct Movie {
    char title [128];
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```

In C, we might store these instances in some data-structure which we could then query via loops and if statements.

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To do this we represent the struct Movie instances as tuples, and we "store" our tuples in sets. For example,

 $\{(\mathsf{Gone}\ \mathsf{With}\ \mathsf{the}\ \mathsf{Wind}, 1939, 231, \mathsf{drama}),\ (\mathsf{Star}\ \mathsf{Wars}, 1977, 124, \mathsf{scifi})\}$ 

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(("title", \textit{string}), ("year", \textit{int}), ("length", \textit{int}) ("genre", \textit{string}))
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```

A set of tuples along with their schema is called a **relation** (or sometimes relation instance.)

## The Table Representation

We will often make use of a notational convention where a relation is depicted as a table. For example,

Table: Example relation for the Movie schema.

title	year	length	genre
Gone With the Wind	1939	231	drama
Star Wars	1977	124	sciFi
Wayne's World	1992	95	comedy

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The header of the table is made from the names in the schema, and the rows of the table are just the tuples in our relation.

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... Admittedly, this seems very unlikely.

We can constrain the Movies schema so that no two movies have the same name and the same year attributes by asserting that {title, year} is a **key**.

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A set of attributes is called a key if no two tuples in the relation instance can have the same values for all elements of the key.

### Compact Schema Representation

Sometimes we will need to name and present a schema in a compact way.

In these cases we will just write the schema name as a prefix to a tuple containing the attribute names. If we need to emphasize a key, we will underline the names forming it. For example,

 ${\tt Movie}(\underline{{\tt title}},\ \underline{{\tt year}},\ {\tt length},\ {\tt genre})$ 

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We may take a look at JSON towards the end of the semester.