

# Data Structures and Algorithms

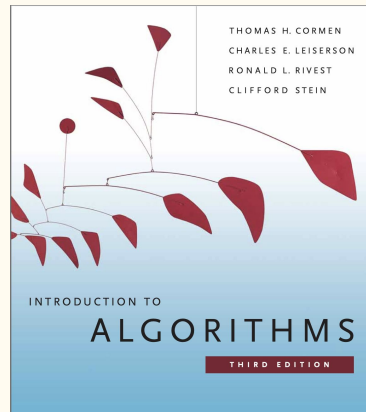
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Tutorial 7. Random BST. Red-Black Trees

# Today's topic is covered in detail in

T. H. Cormen, C. E. Leiserson, R. L. Rivest, and C. Stein.  
**Introduction to Algorithms.** The MIT Press 2009.

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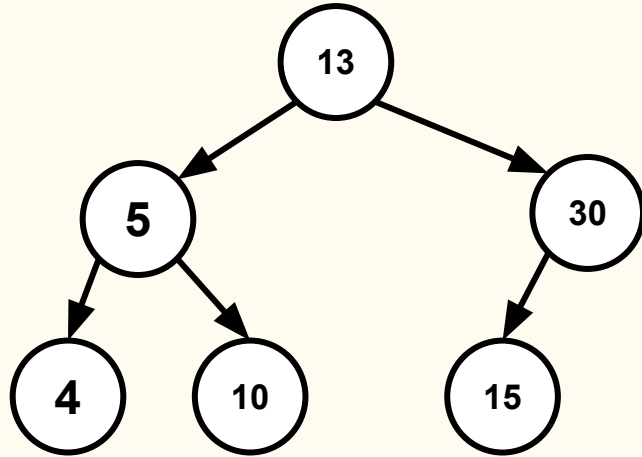


# Objectives

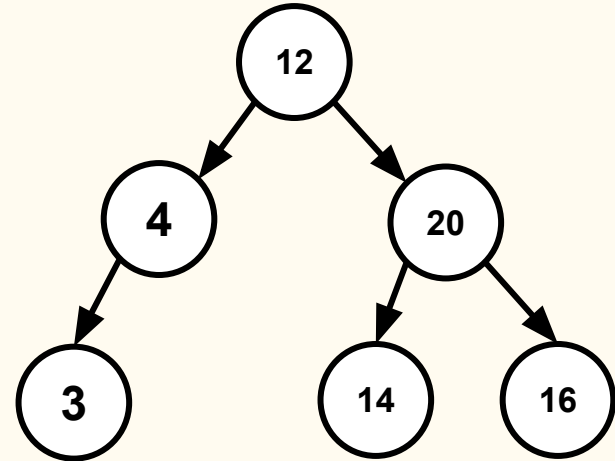
- Recap: binary search tree
- Height of a randomly build BST
- Red-Black Trees: invariant, insertion, deletion

## Red-Black Trees: exercise

**Exercise 7.0.** Which of the following are valid BSTs?



A



B

# Height of a randomly built BST

**Theorem.** The expected height of a randomly built binary search tree on  $n$  distinct keys is  $O(\log n)$ .

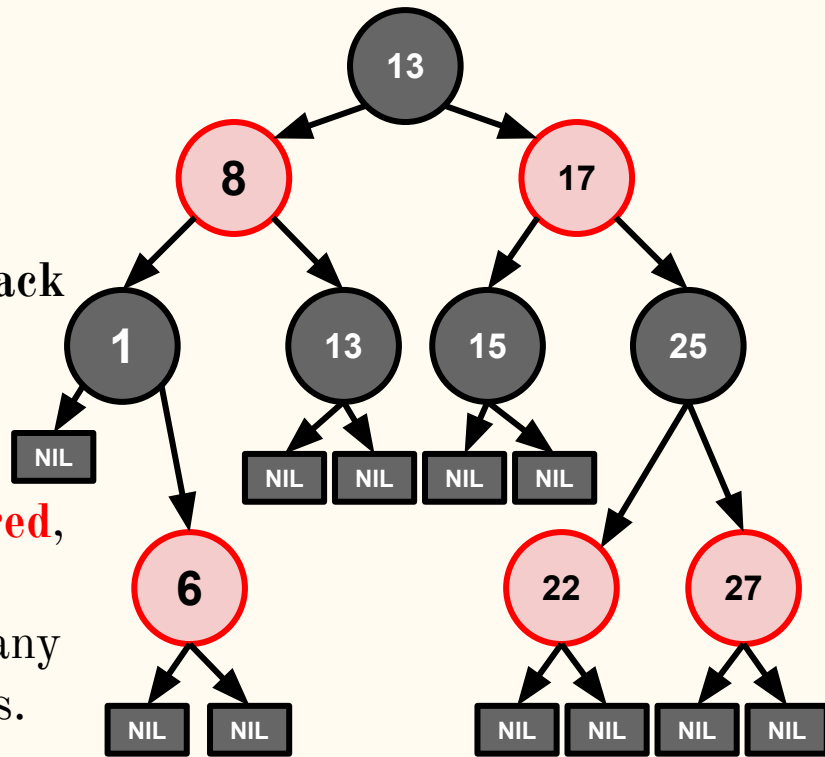
## Randomly build BST: exercise

**Exercise 7.1.** Show that **randomly build** BST on  $n$  keys is not the same as **randomly chosen** BST on  $n$  key.  
Hint: consider  $n = 3$ .

# Red-Black Trees: invariant

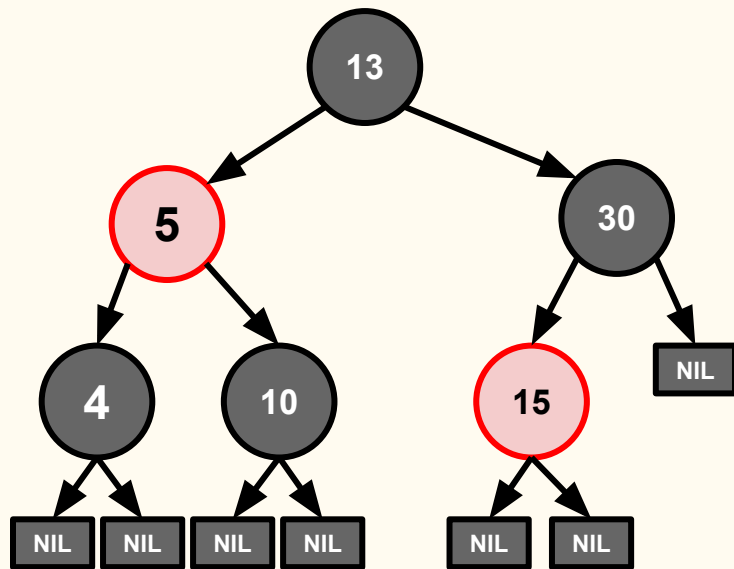
Red-Black Tree is a type of self-balancing BST:

1. Each node is either **red** or **black** (this information is stored in each node).
2. The root is always **black**.
3. Every leaf (NIL) is **black**.
4. If a node is **red**, then both its children are **black**.
5. For each node, all paths from this node to any leaf contain the same number of **black** nodes.

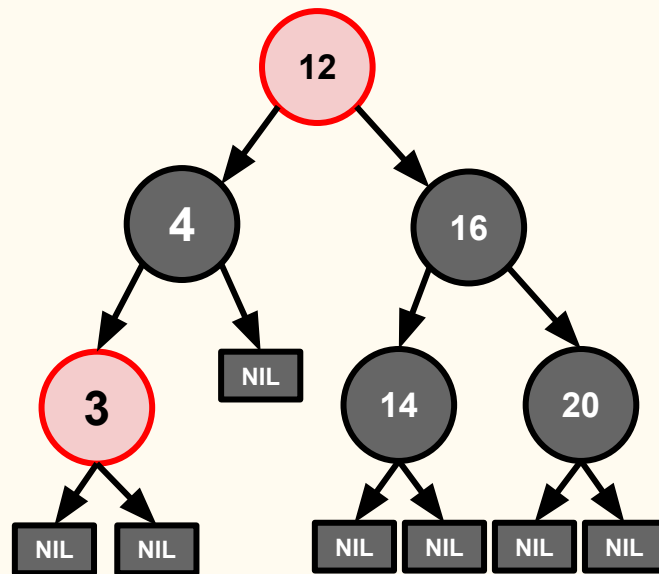


# Red-Black Trees: exercise

**Exercise 7.1.** Which of the following are valid RBTs?



A

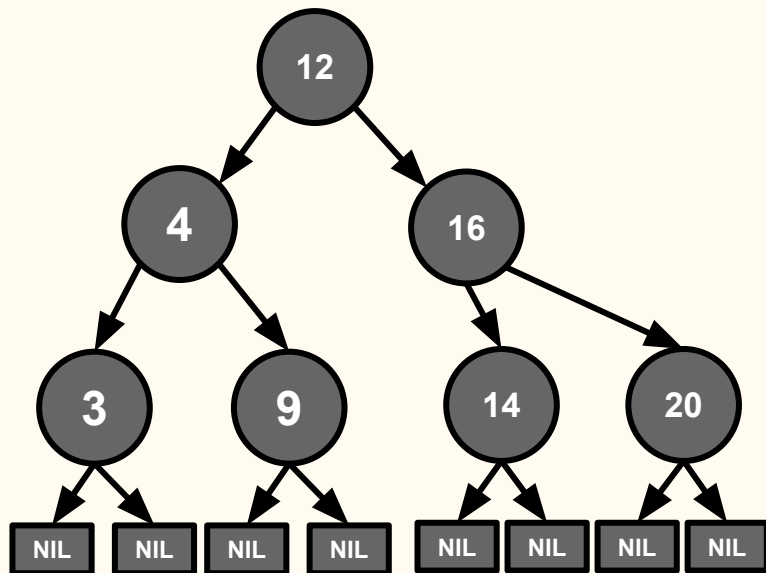


B

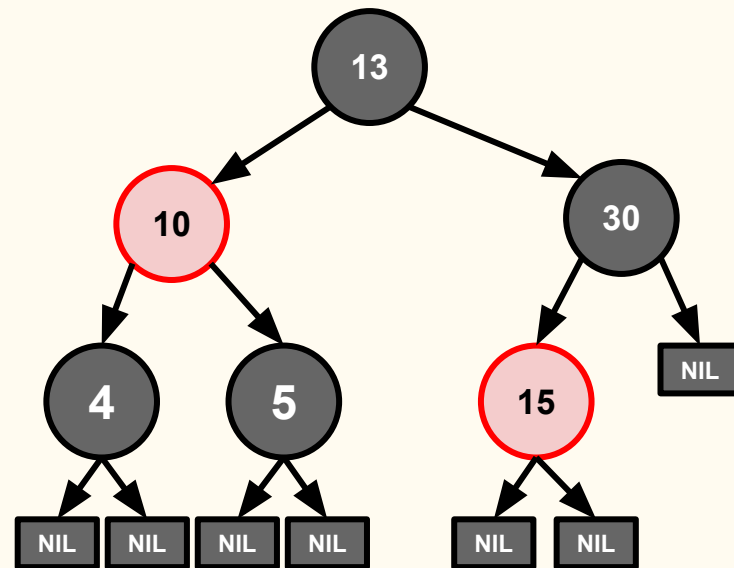


# Red-Black Trees: exercise

**Exercise 7.1.** Which of the following are valid RBTs?



C



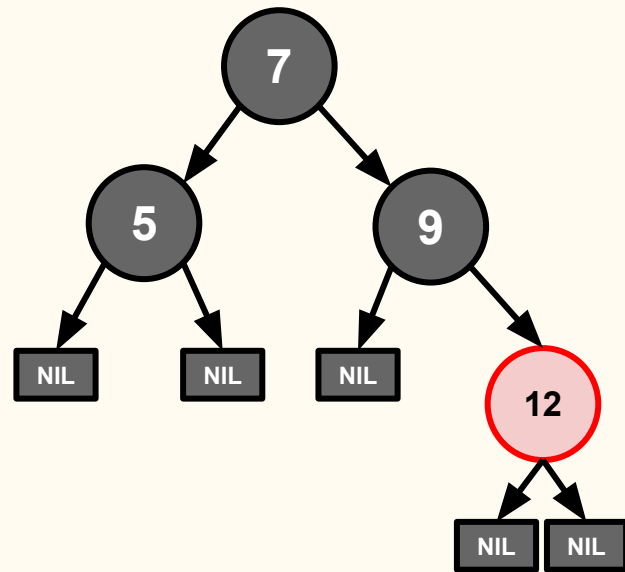
D

# Red-Black Trees: exercise

## Exercise

7.2.

Insert keys 8, 11, 10 in this RBT.  
What color should the new nodes have?



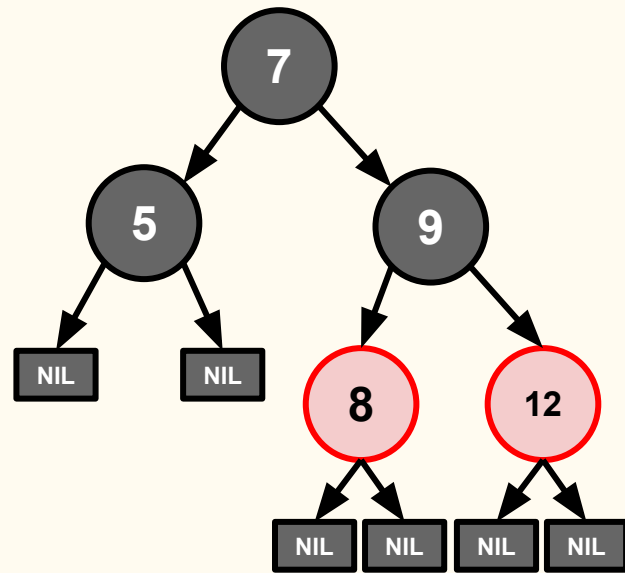
Valid RBT

# Red-Black Trees: exercise

## Exercise

Insert keys 8, 11, 10 in this RBT.  
What color should the new nodes have?

7.2.



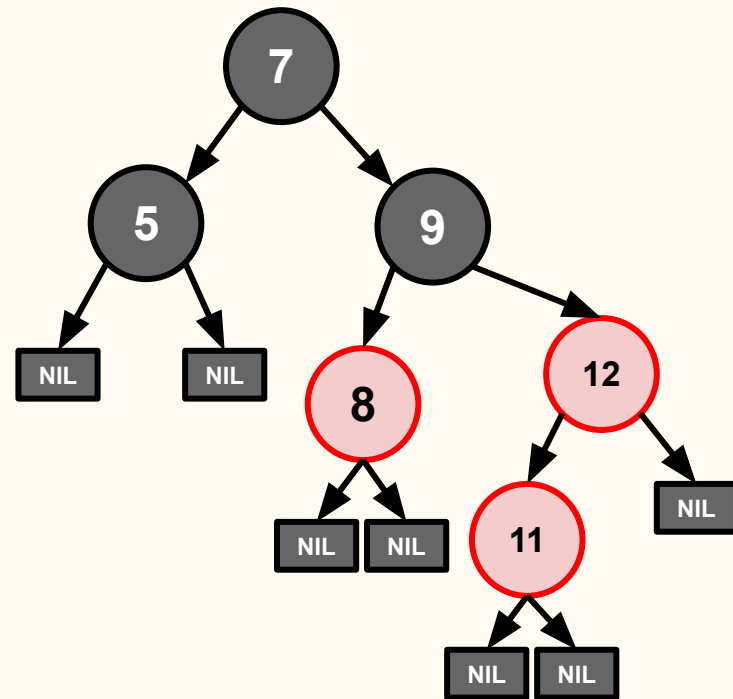
Valid RBT

# Red-Black Trees: exercise

## Exercise

Insert keys 8, 11, 10 in this RBT.  
What color should the new nodes have?

7.2.



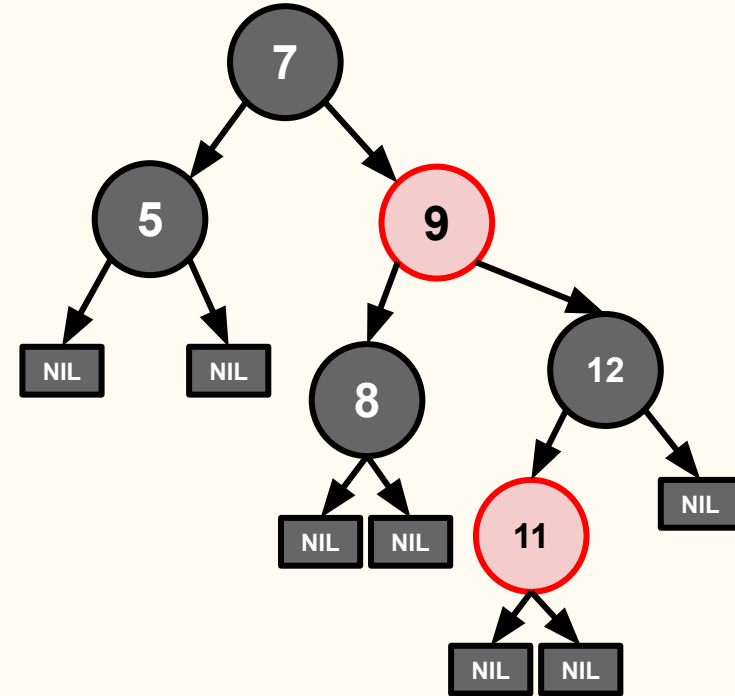
Invalid RBT

# Red-Black Trees: exercise

## Exercise

7.2.

Insert keys 8, 11, 10 in this RBT.  
What color should the new nodes have?



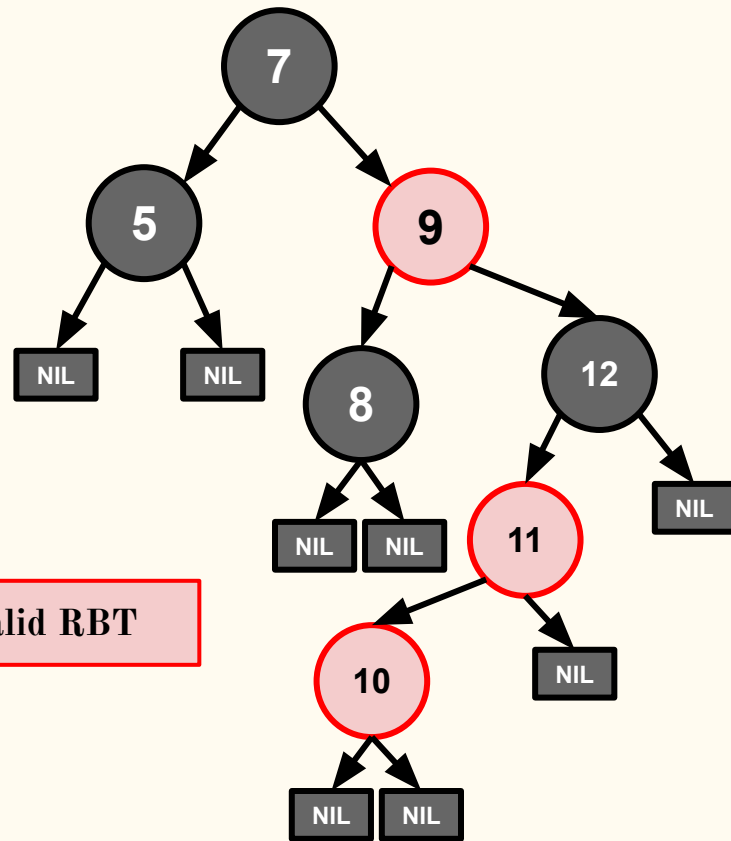
Valid RBT

# Red-Black Trees: exercise

## Exercise

Insert keys 8, 11, 10 in this RBT.  
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7.2.

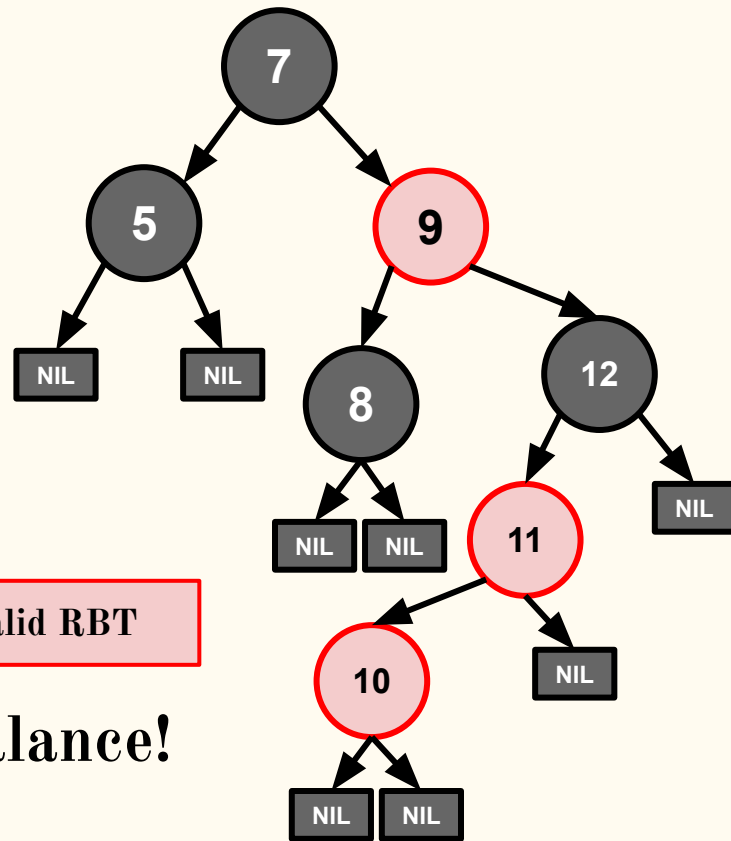


# Red-Black Trees: exercise

## Exercise

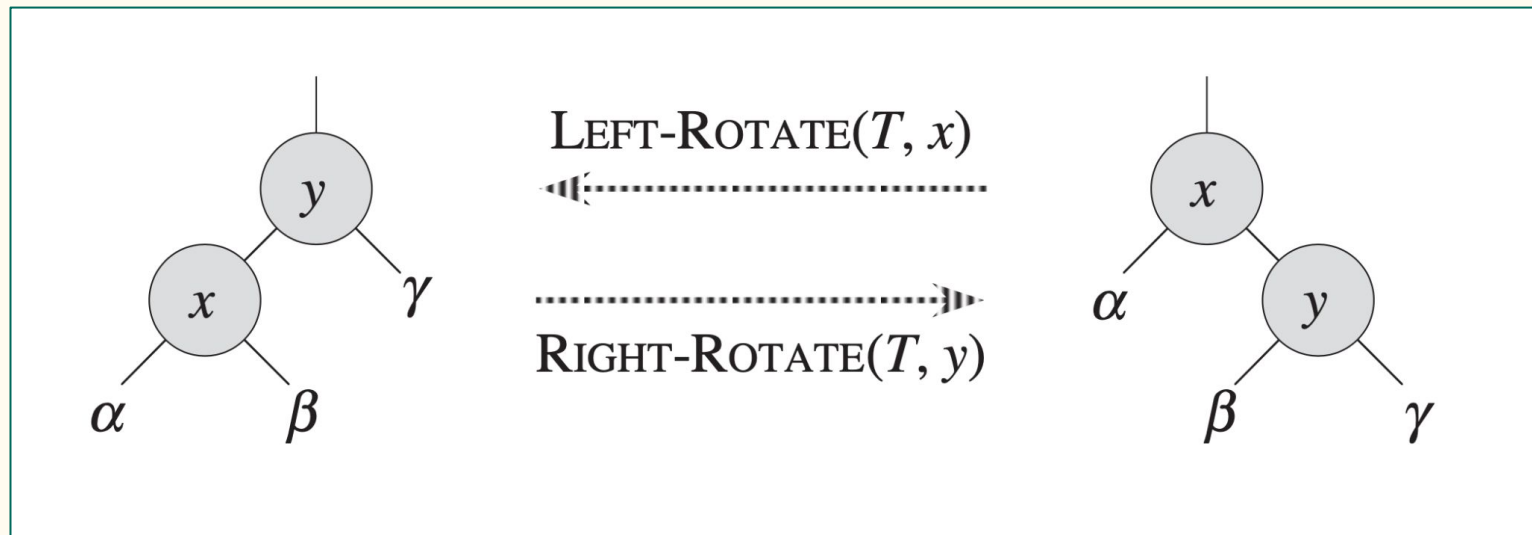
7.2.

Insert keys 8, 11, 10 in this RBT.  
What color should the new nodes have?



**Cannot fix by recoloring. Have to rebalance!**

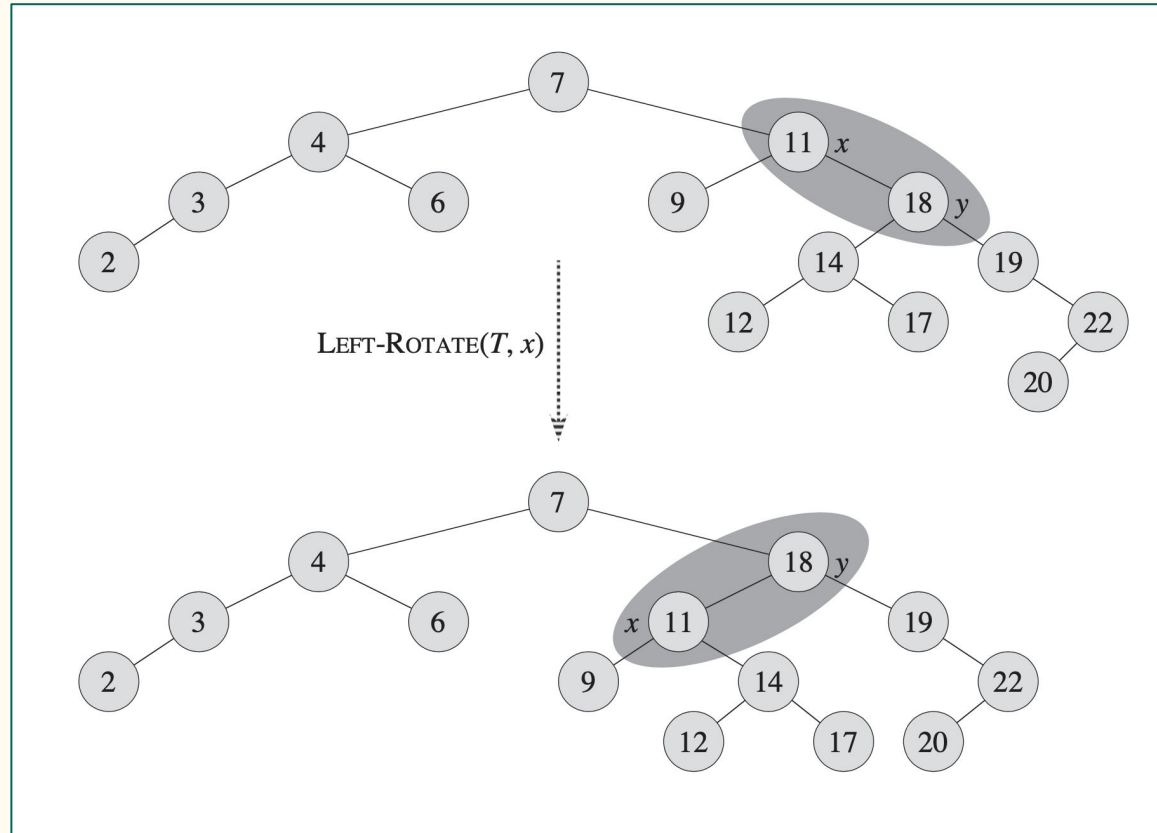
# BST rotations



Idea: change the shape of the tree, preserving BST property.



# BST rotations: example

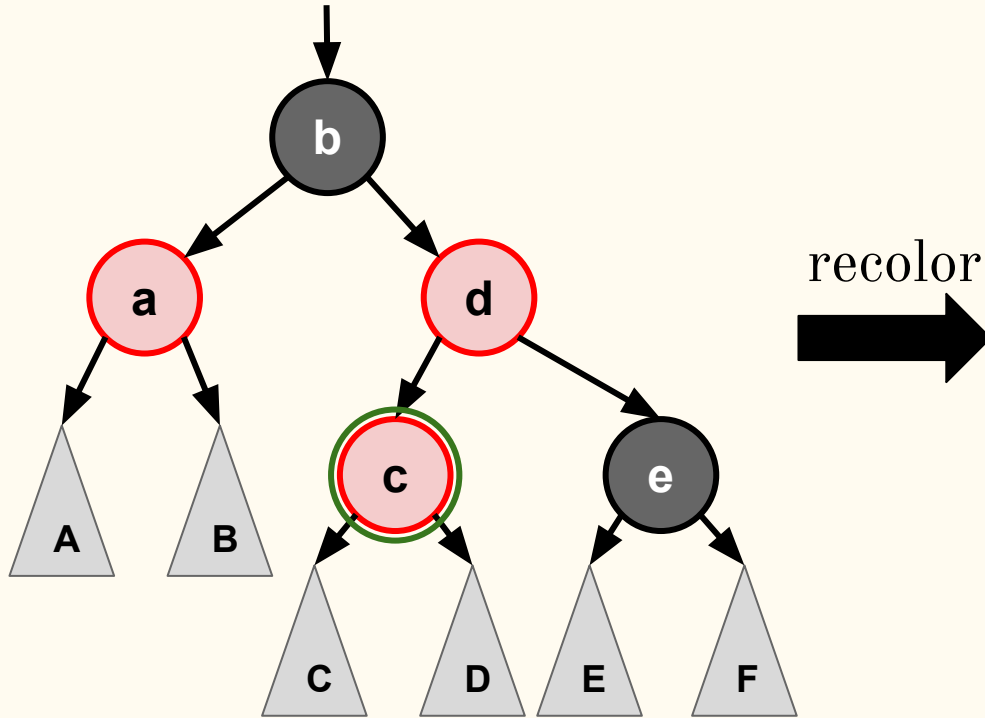


# Red-Black Trees: insertion

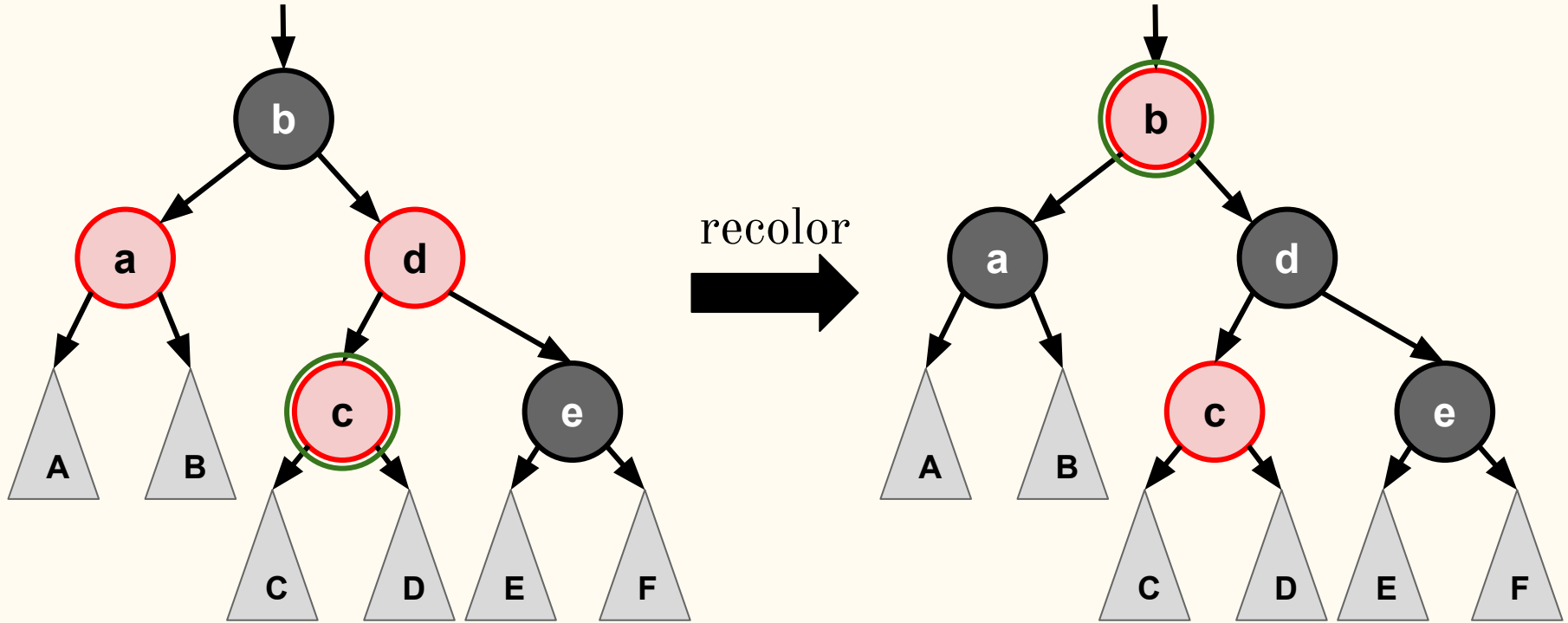
The idea for insertion is simple:

1. Insert a new **red** node using the default BST insertion
2. Fix the tree, recursively raising from the new node:
  - a. If current node's parent is **black** — stop
  - b. If current node's parent and uncle are **red** — recolor and go up
  - c. If current node's parent is **red** and uncle is **black**, then rotate, recolor, and go up

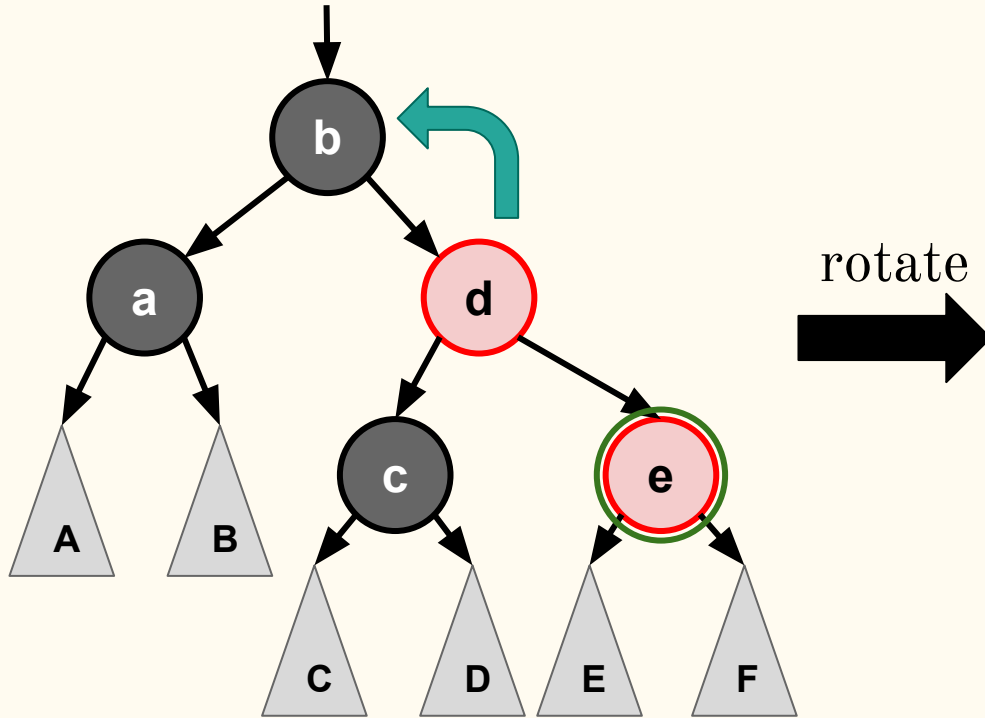
# Red-Black Trees: insertion (case 1)



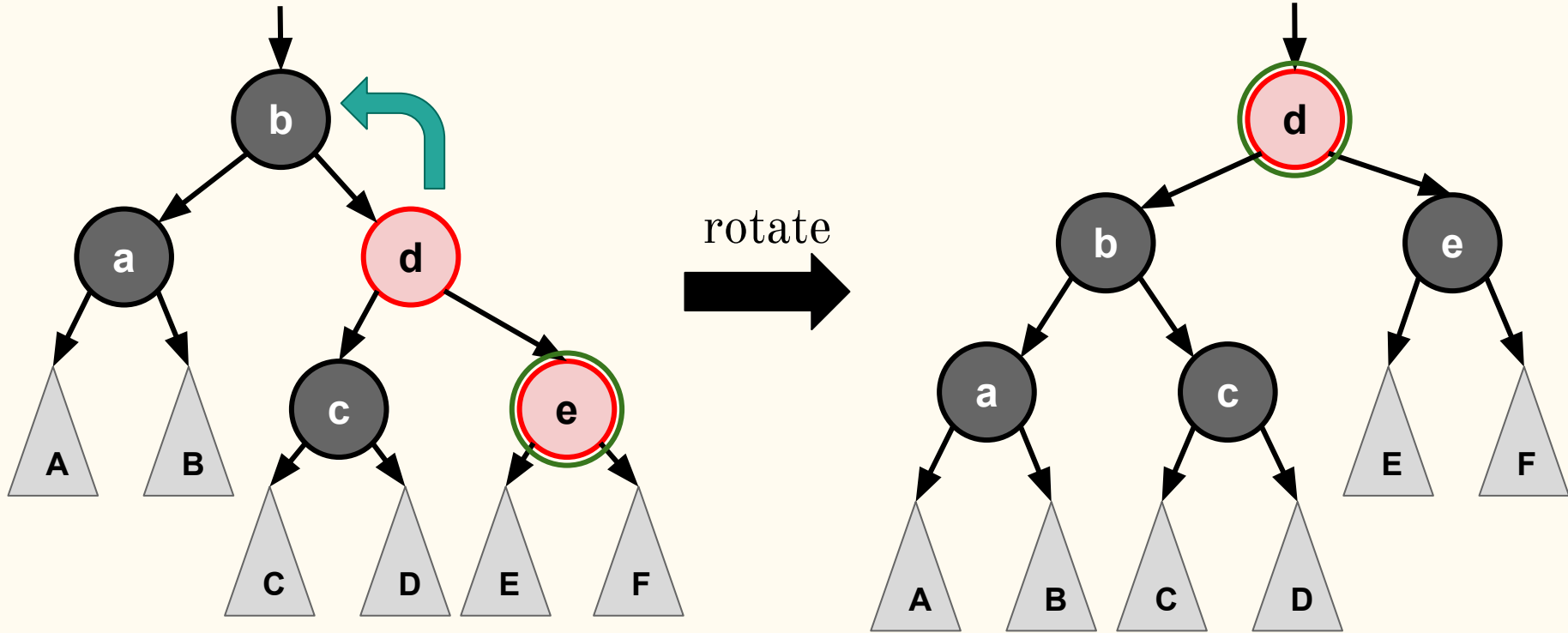
# Red-Black Trees: insertion (case 1)



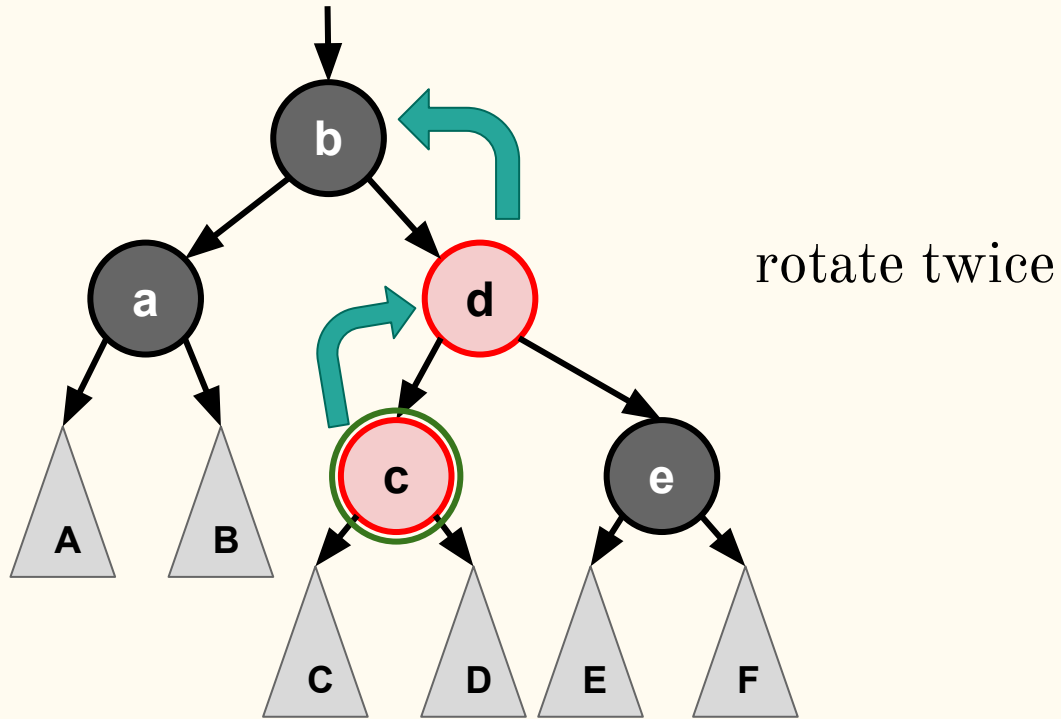
## Red-Black Trees: insertion (case 2)



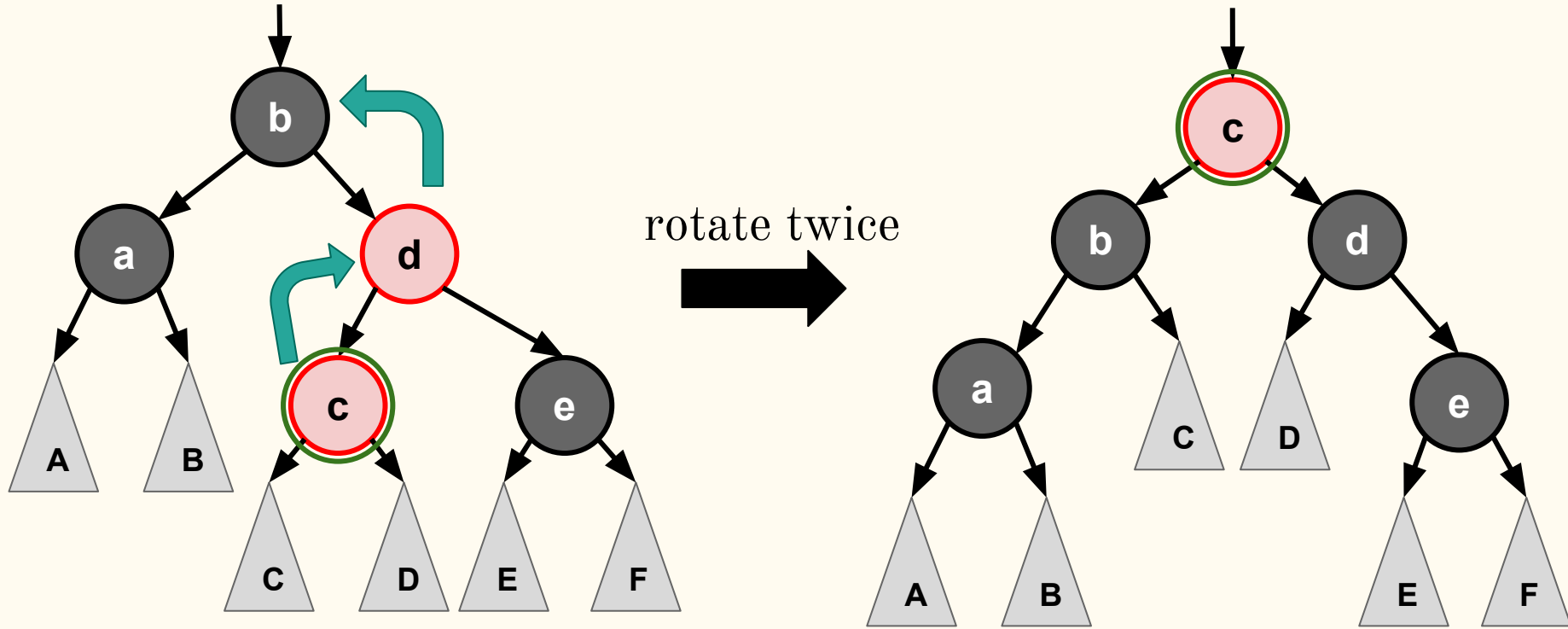
## Red-Black Trees: insertion (case 2)



# Red-Black Trees: insertion (case 3)



# Red-Black Trees: insertion (case 3)





# Red-Black Trees: insertion time complexity

The time complexity of insertion into a Red-Black Tree is

?

## Red-Black Trees: insertion time complexity

The time complexity of insertion into a Red-Black Tree is

$$O(\log n)$$

## Red-Black Trees: insertion (exercise)

**Exercise 7.3.** Delete these keys in order from the given RBT:

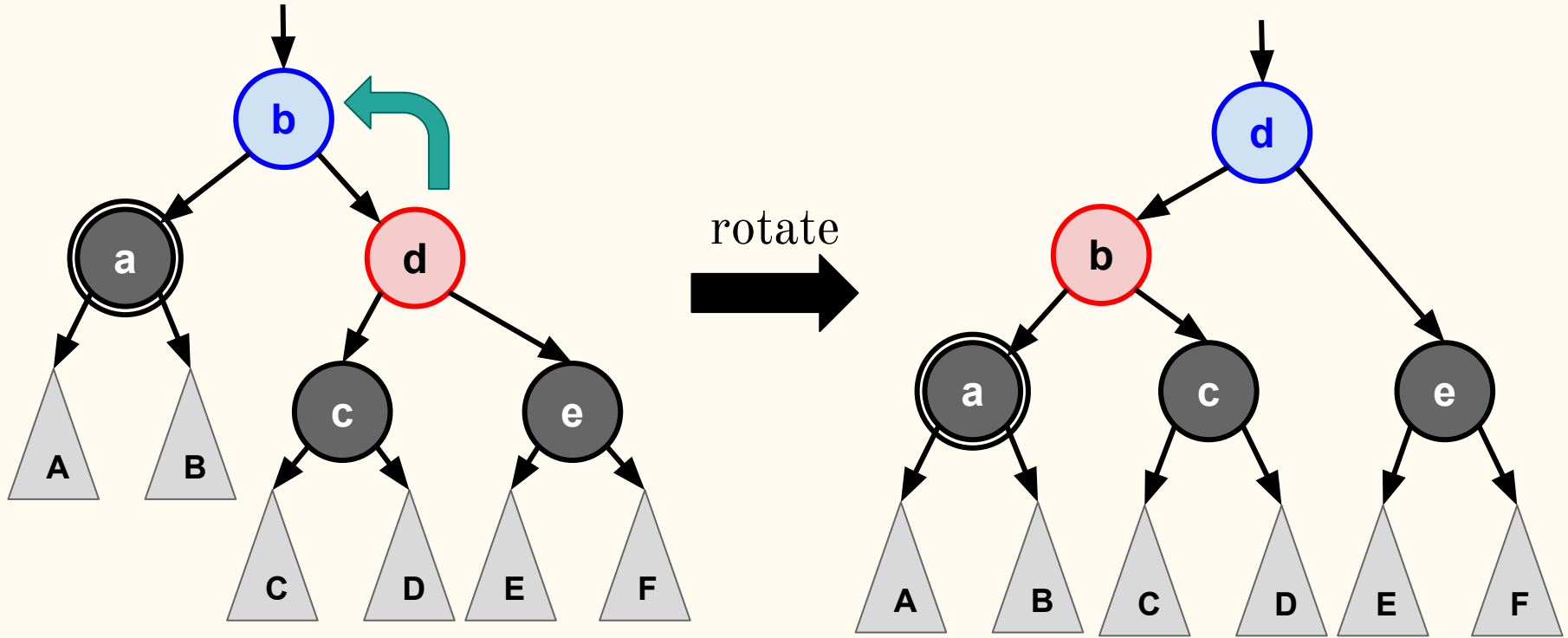
8, 12, 19, 31, 38, 41

# Red-Black Trees: deletion

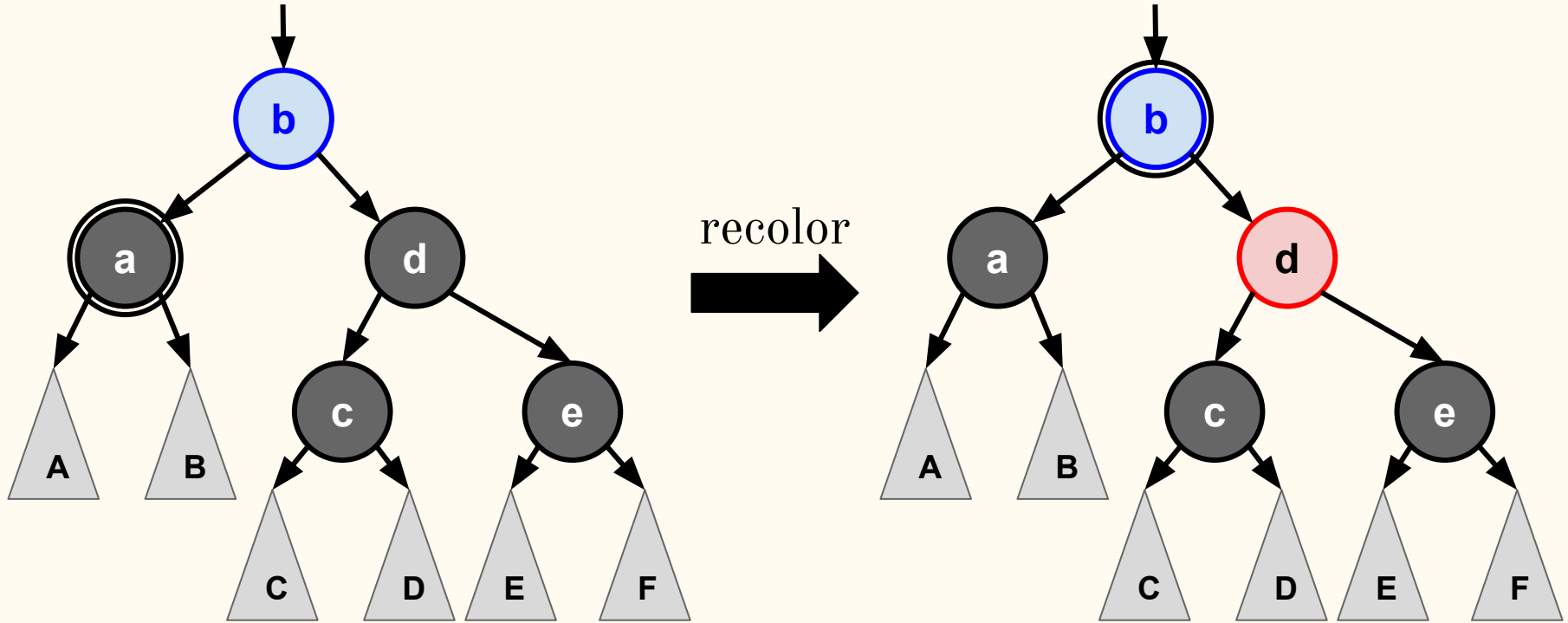
The idea for deletion is simple:

1. Delete a node using the default BST deletion
2. If deleted leaf node was **red**, nothing has to be adjusted.
3. Otherwise, fix the tree, recursively raising from the deleted node:
  - a. If current node is **red** — stop
  - b. Examine current node's sibling and nephews to perform rotation and recoloring correspondingly

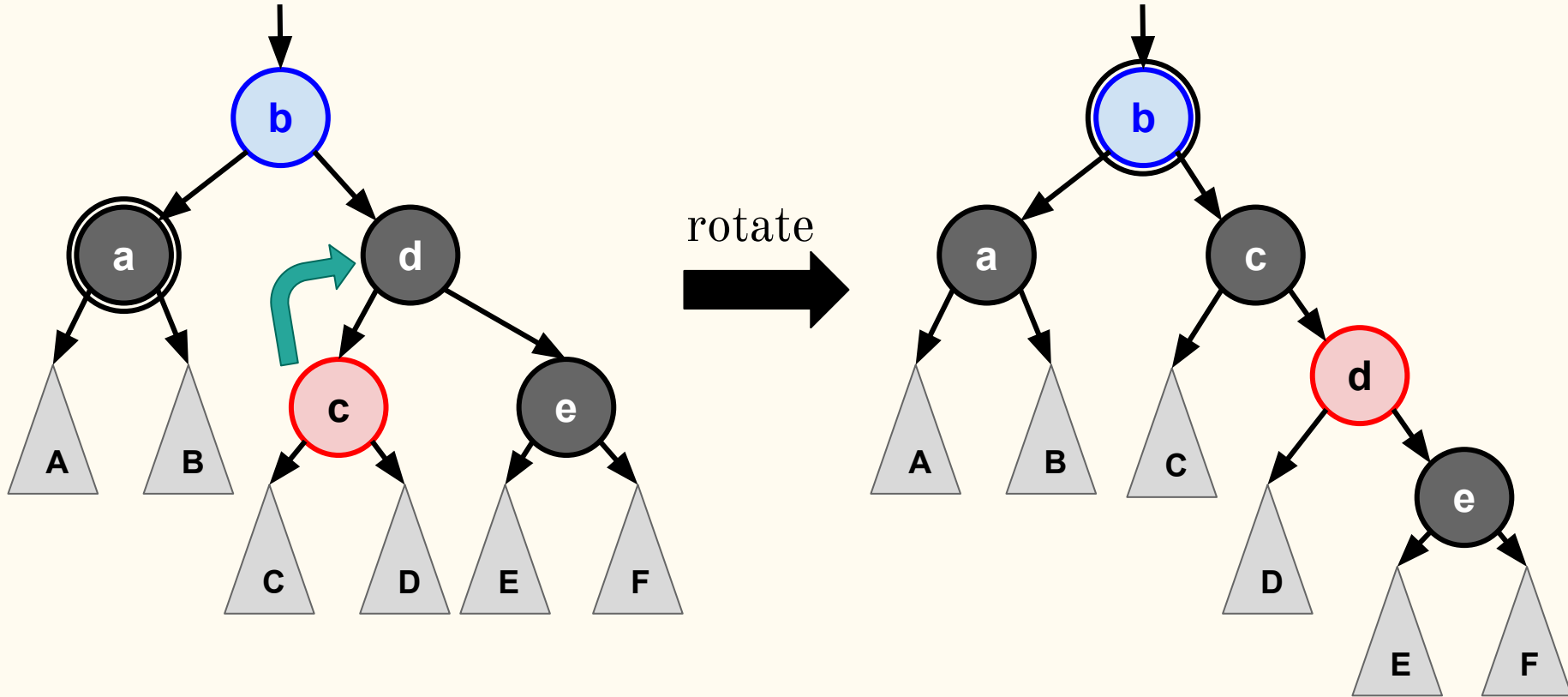
# Red-Black Trees: deletion (case 1)



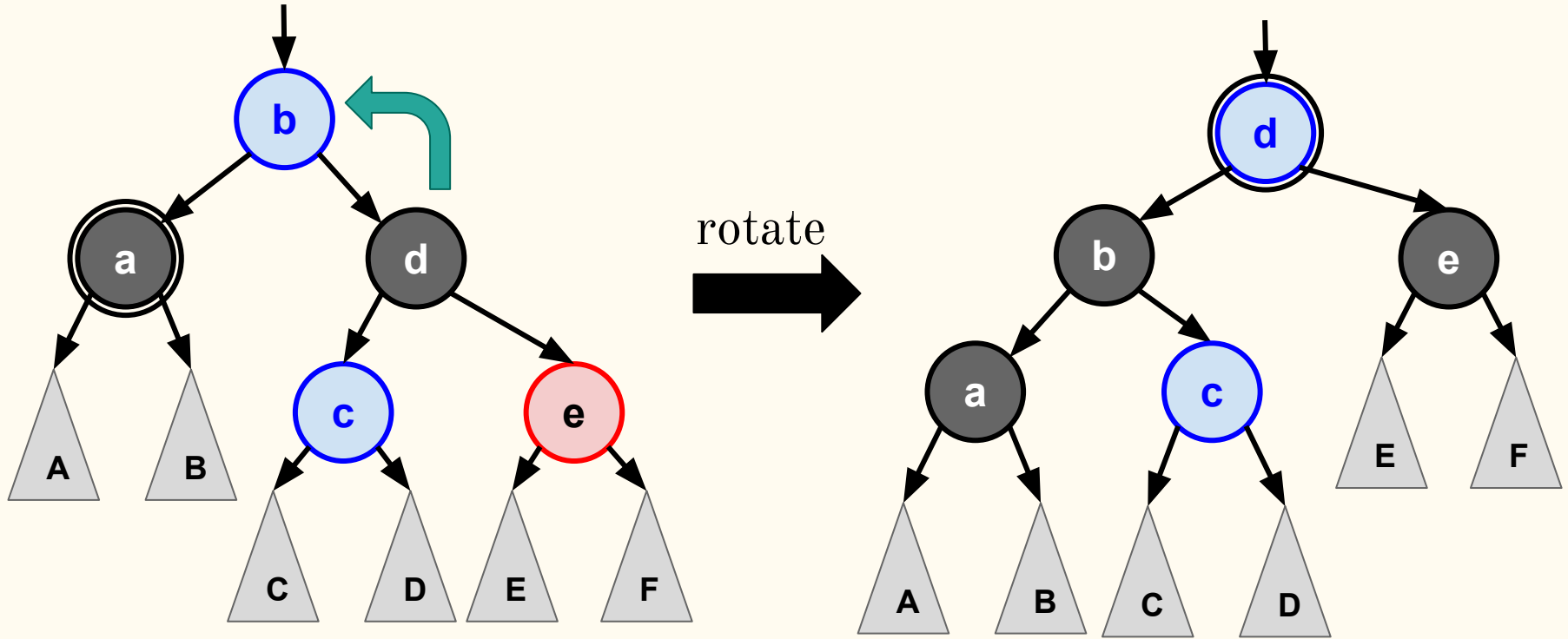
## Red-Black Trees: deletion (case 2)



## Red-Black Trees: deletion (case 3)



## Red-Black Trees: deletion (case 4)





# Red-Black Trees: deletion time complexity

The time complexity of deletion from a Red-Black Tree is

?

## Red-Black Trees: deletion time complexity

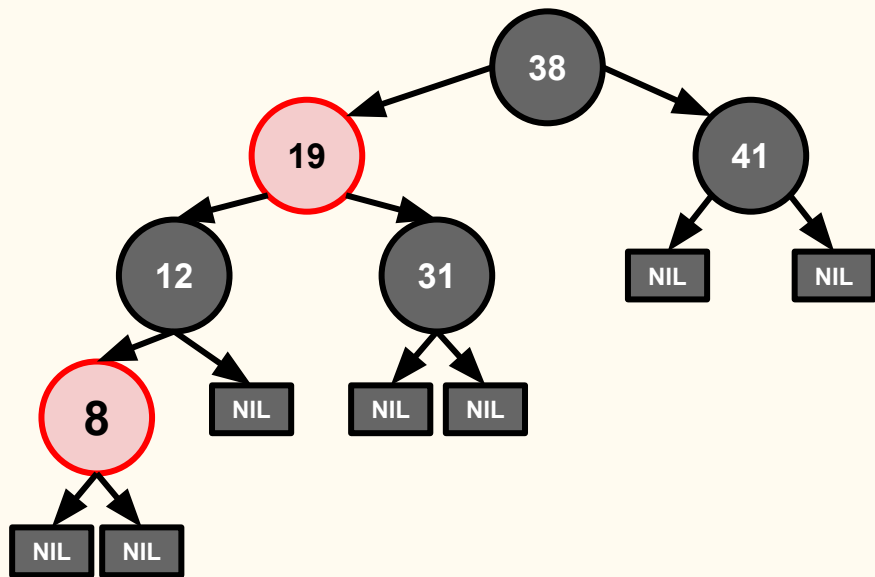
The time complexity of deletion from a Red-Black Tree is

$$O(\log n)$$

## Red-Black Trees: deletion (exercise)

**Exercise 7.4.** Build a RBT by inserting these keys in order:

11, 19, 8, 16, 17, 31, 26, 41, 61



# Summary

- Randomly build BST has expected height of  $\Theta(\log n)$
- Red-Black tree is a kind of self-balancing BST
  - Height of an RBT is  $O(\log n)$
  - Insertion into an RBT takes  $O(\log n)$
  - Deletion from an RBT takes  $O(\log n)$

# Summary

- Randomly build BST has expected height of  $\Theta(\log n)$
- Red-Black tree is a kind of self-balancing BST
  - Height of an RBT is  $O(\log n)$
  - Insertion into an RBT takes  $O(\log n)$
  - Deletion from an RBT takes  $O(\log n)$

See you next week!