

# Data Structures and Algorithms

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Lab 10

Topological sorting

# Agenda

- Recall DFS, BFS
- Topological sorting
- Coding exercise

# DFS

Which data structure do we use?

What kind of tasks can it solve?

# DFS

Which data structure do we use?

- Stack

What kind of tasks can it solve?

- Path from a to b (lexicographical)
- Topological sorting
- Find a cycle
- Find connected components
- ...

# DFS pseudocode

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**Algorithm 1:** Recursive DFS

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**Data:**  $G$ : The graph stored in an adjacency list

$root$ : The starting node

**Result:** Prints all nodes inside the graph in the *DFS* order

$visited \leftarrow \{false\};$

$DFS(root);$

**Function**  $DFS(u)$ :

**if**  $visited[u] = true$  **then**

**return**;

**end**

$print(u);$

$visited[u] \leftarrow true;$

**for**  $v \in G[u].neighbors()$  **do**

$DFS(v);$

**end**

**end**

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# BFS

Which data structure do we use?

What kind of tasks can it solve?

# BFS

Which data structure do we use?

- Queue

What kind of tasks can it solve?

- Shortest path from a to b
- Find shortest cycle
- Find connected components
- ...

# BFS pseudocode

```
procedure BFS(G,s)

  for each vertex  $v \in V[G]$  do
    explored[ $v$ ]  $\leftarrow$  false
     $d[v] \leftarrow \infty$ 
  end for
  explored[ $s$ ]  $\leftarrow$  true
   $d[s] \leftarrow 0$ 
   $Q :=$  a queue data structure, initialized with  $s$ 
  while  $Q \neq \emptyset$  do
     $u \leftarrow$  remove vertex from the front of  $Q$ 
    for each  $v$  adjacent to  $u$  do
      if not explored[ $v$ ] then
        explored[ $v$ ]  $\leftarrow$  true
         $d[v] \leftarrow d[u] + 1$ 
        insert  $v$  to the end of  $Q$ 
      end if
    end for
  end while

end procedure
```



# Cycle in graph

How to find a cycle in given graph?

During DFS let's “color” vertices:

1. “Unvisited” vertices are **white**
2. “visited” are **gray**

**If during DFS we meet gray vertice – means we found a cycle.**

# Topological sorting

A linear ordering of vertices such that

- for every directed edge  $U \rightarrow V$
- vertex  $U$  comes before  $V$  in the ordering

We can apply topological sorting Directed Acyclic Graphs(DAG)

# Topological sorting: Example

Given a list of tasks:

- Task 1: Find a client
- Task 2: List requirements
- Task 3: Start working on the project
- Task 4: Build a team
- etc.

Suggest a possible order of execution considering that some tasks depend on result of others.

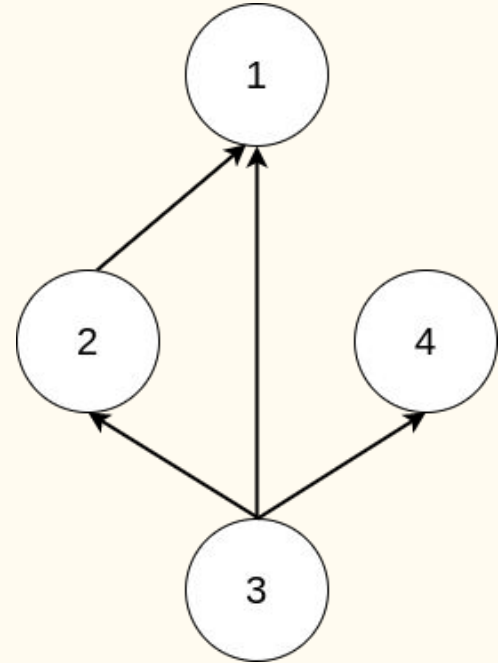
- Task 2 depends on task 1;
- Task 3 depends on task 1, 2 and 4;

# Topological sorting: Example

Given a list of tasks, some tasks depend on result of others.

Suggest a possible order of execution.

- Task 2 depends on task 1;
- Task 3 depends on task 1, 2 and 4;



How many ways to execute all of them?

# Topological sorting: Example

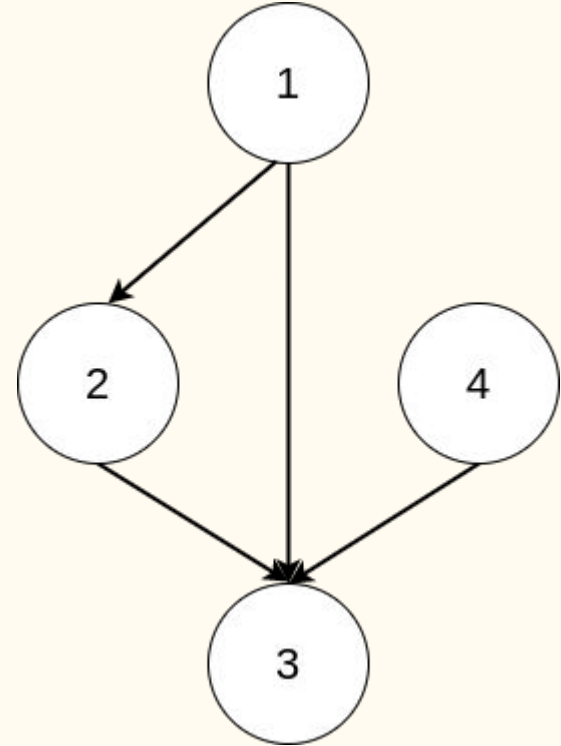
How should we modify the graph?

From:

- Task 2 depends on task 1;
- Task 3 depends on task 1, 2 and 4;

Transpose to:

- Task 1 executes before task 2 and 3
- Task 2 executes before task 3
- Task 4 executes before task 3



# Topological sorting

How to order vertices?

- Record when we leave a node
- The node we leave first has the most dependencies.
  - The vertex with the maximum in-degree
- The node we leave last has the least dependencies.
  - The vertex with the minimum in-degree

# Topological sorting: implementation

- In DFS, we print a vertex and then recursively call DFS for its adjacent vertices.
- In topological sorting, we need to print a vertex before its adjacent vertices.
- In DFS, we start from a vertex, we first print it and then recursively call DFS for its adjacent vertices.
- In topological sorting, we use a temporary stack. We don't print the vertex immediately, we first recursively call topological sorting for all its adjacent vertices, then push it to a stack.
- Finally, print contents of the stack.
- Note that a vertex is pushed to stack only when all of its adjacent vertices (and their adjacent vertices and so on) are already in the stack.

# Topological sorting: implementation

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**Algorithm 3:** DFS Algorithm for Topological Sort

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**Data:** A DAG  $G$

**Result:** A topological sort of all vertices in  $G$

Create an empty vertex list  $L$ ;

Create a *visited* array to indicate whether a vertex has been visited;

Initialize all elements in *visited* with *false*;

**foreach** vertex  $u \in G$  **do**

**if** *visited*[ $u$ ] *is false* **then**

        | `topologicalSortRecursive( $u$ );`

**end**

**end**

**return**  $L$ ;

**Function** `topologicalSortRecursive( $u$ ):`

    | *visited*[ $u$ ] = *true*;

**foreach**  $u$ 's neighboring vertex  $v$  **do**

        | **if** *visited*[ $v$ ] *is false* **then**

            | `topologicalSortRecursive( $v$ );`

        | **end**

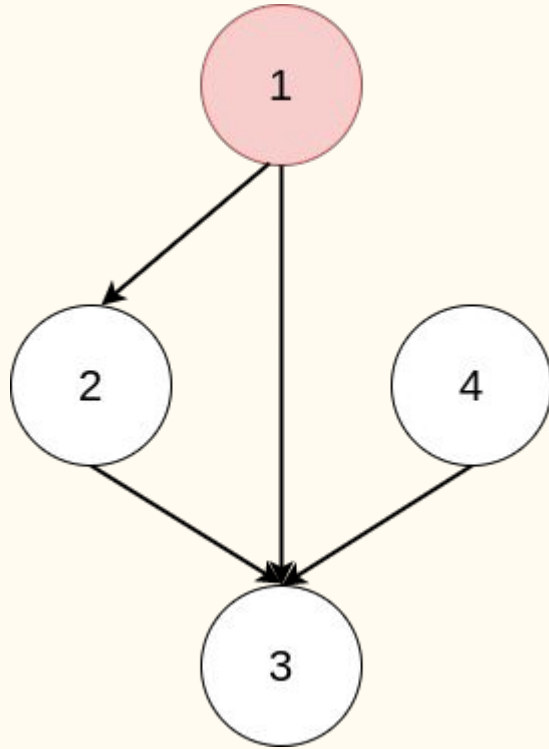
**end**

    | *Add  $u$  to the front of list  $L$ ;*

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# Sorting



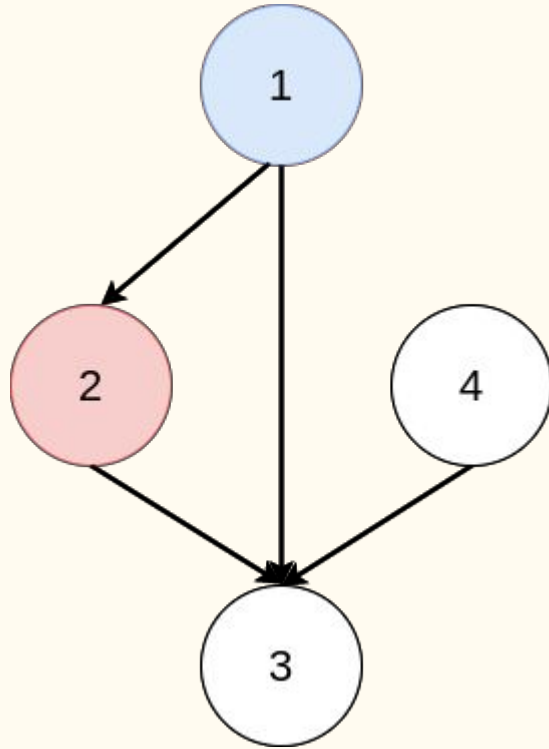
DFS Stack

1			
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T.Sort

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# Sorting



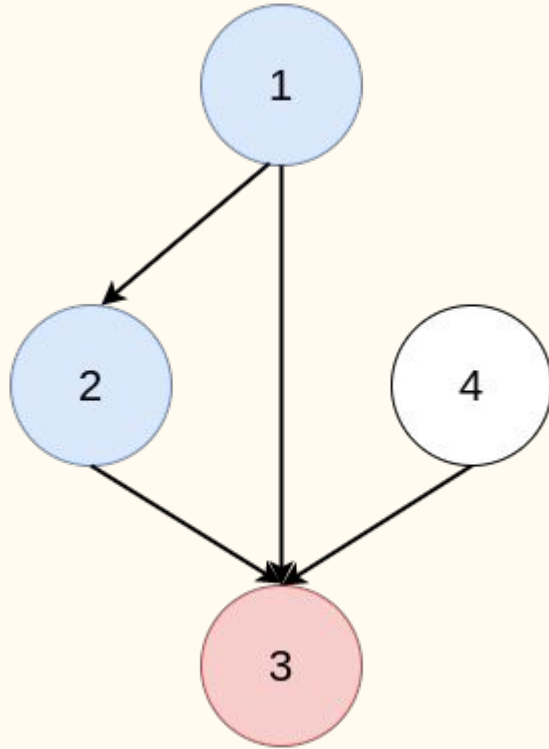
DFS Stack

1	2		
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T.Sort

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# Sorting



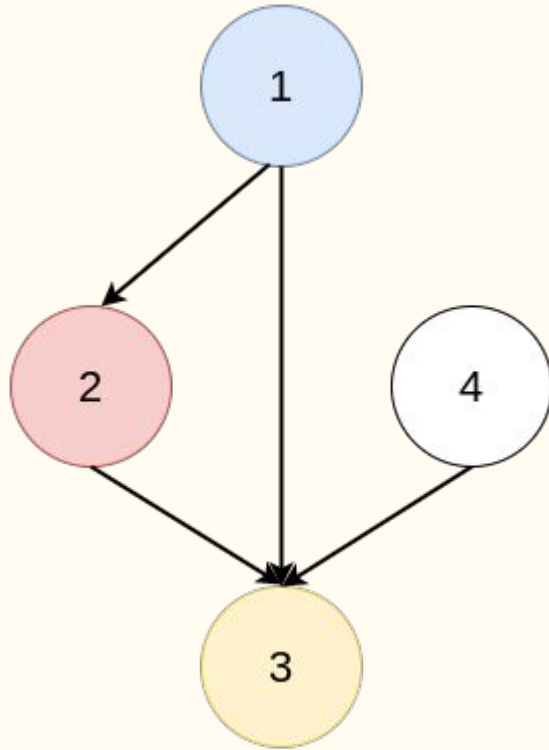
DFS Stack

<b>1</b>	<b>2</b>	<b>3</b>	
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T.Sort

--	--	--	--

# Sorting



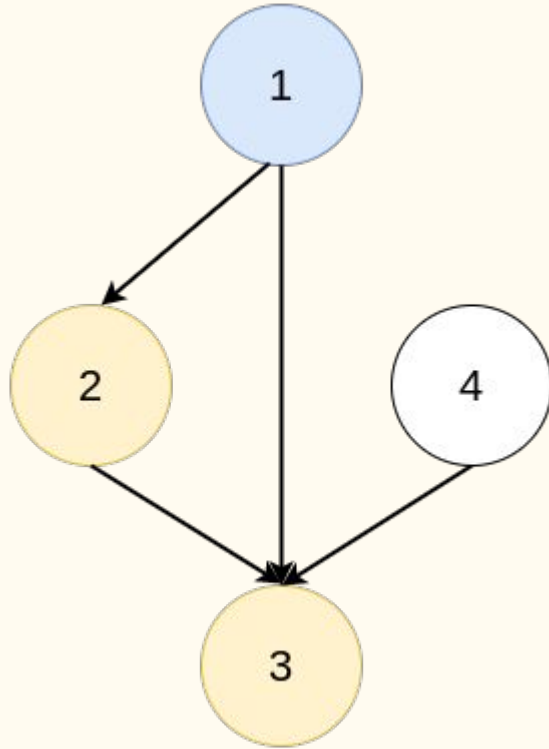
DFS Stack

<b>1</b>	<b>2</b>		
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T.Sort

3			
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# Sorting



DFS Stack

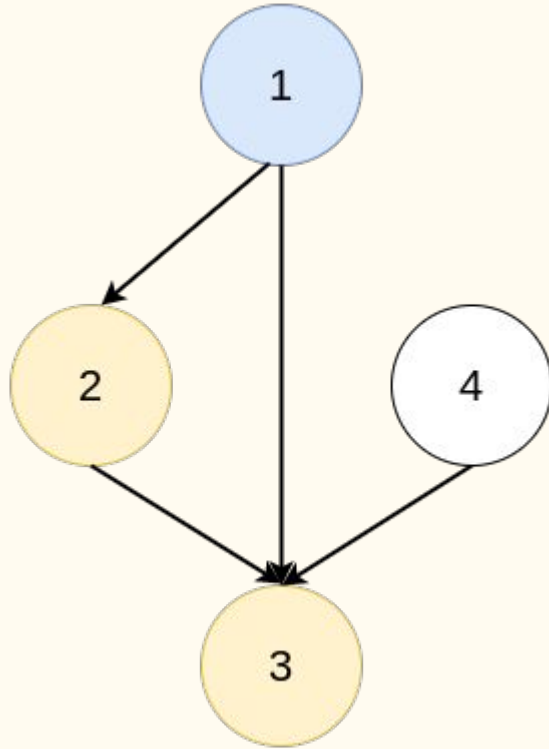


T.Sort



Where to put 2 ?

# Sorting



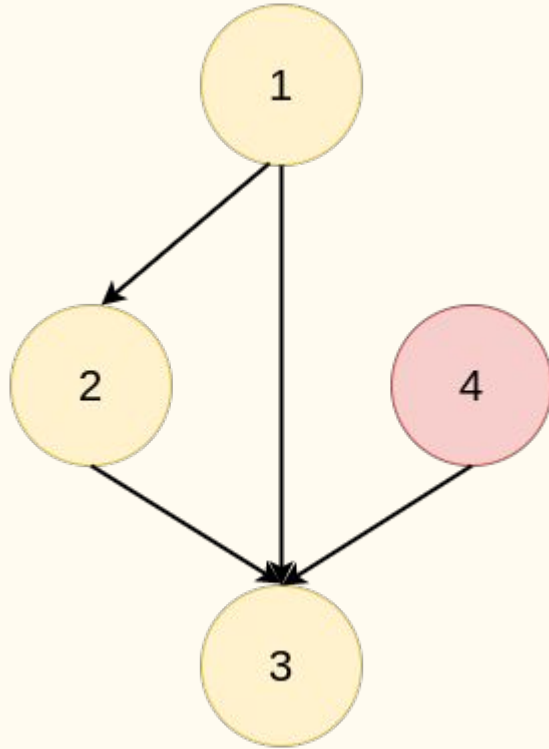
DFS Stack

<b>1</b>			
----------	--	--	--

T.Sort

<b>3</b>	<b>2</b>		
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# Sorting



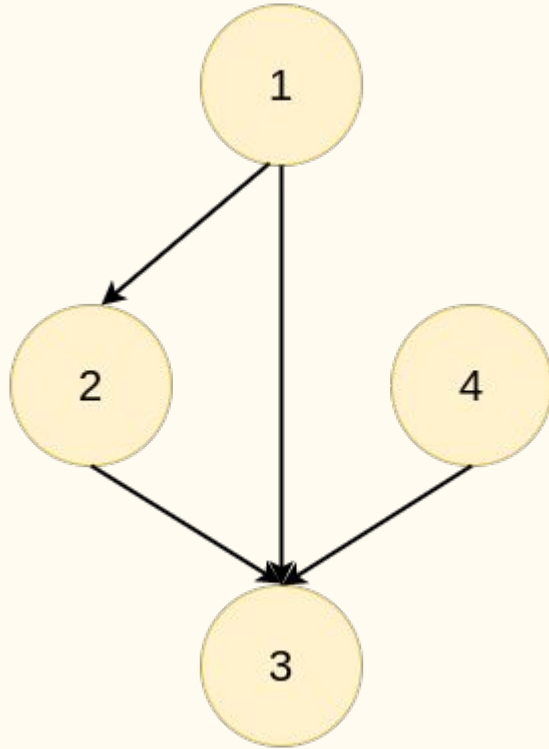
DFS Stack



T.Sort



# Sorting



DFS Stack

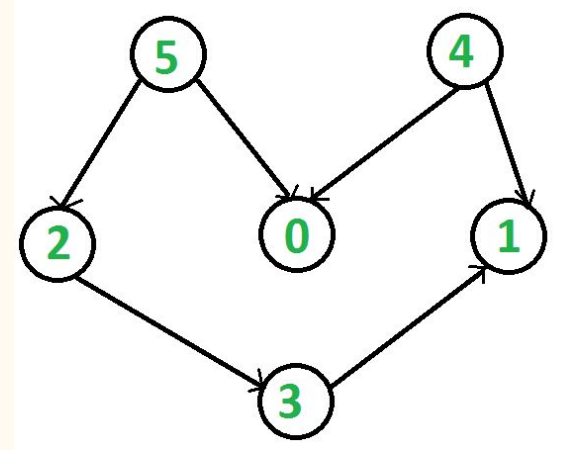


T.Sort

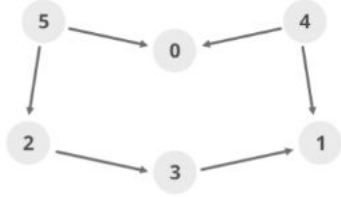




## Another example



- A topological sorting of the following graph is “5 4 2 3 1 0”.
- There can be more than one topological sorting for a graph.
- For example, another topological sorting of the following graph is “4 5 2 3 1 0”. The first vertex in topological sorting is always a vertex with in-degree as 0 (a vertex with no incoming edges).



Adja cent list (G)

- 0 →
- 1 →
- 2 → 3
- 3 → 1
- 4 → 0, 1
- 5 → 2, 0

	0	1	2	3	4	5
visited	false	false	false	false	false	false

Stack( empty )

**Step 1:**

Topological Sort( 0 ), visited[ 0 ] = true

List is empty. No more recursion call.

Stack 

0			
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**Step 2:**

Topological Sort( 1 ), visited[ 1 ] = true

List is empty. No more recursion call.

Stack 

0	1		
---	---	--	--

**Step 3:**

Topological Sort( 2 ), visited[ 2 ] = true

Topological Sort( 3 ), visited[ 3 ] = true

'1' is already visited. No more recursion call

Stack 

0	1	3	2
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**Step 4:**

Topological Sort( 4 ), visited[ 4 ] = true

'0' , '1' are already visited. No more recursion call

Stack 

0	1	3	2	4
---	---	---	---	---

**Step 5:**

Topological Sort( 5 ), visited[ 5 ] = true

'2' , '0' are already visited. No more recursion call

Stack 

0	1	3	2	4	5
---	---	---	---	---	---

**Step 6:**

Print all elements of stack from top to bottom



# Topological Sorting: Implementation

Implement topological sorting

See You next week!