

Functions

Survey results

- ◆ 45 responded
- ◆ ~50% are saying the pace is fast
- ◆ ~85% are happy with content and the delivery
- ◆ Free form comments:
 - More examples, go slower
 - Material to study
 - More than 2 hours difficult to sit – can be split into two sessions on Saturday?
 - Compilers should not part of this course
 - Not audible – fix audio issues

Survey results...2

◆ Free form comments...contd....

- Call only last two digits during attendance 😊
- LHC lecture halls are better

◆ Course correction

- More examples will be given
- Material to study (given)
- Can we have two sessions on Saturday too?
1.5-2 hours * 3 per week?
- LHC can be considered
- Audio – will try and fix the issues

Syllabus

Lecture Series (hours)	Topics
1-4	Introduction and Motivation, Paradigms
5-10	Syntax and Semantics, BNF, Compilation
11-18	Data Types, Constructs, Functions, Activation Records , Names and Bindings
19-28	Functional PLs, Logical PLs, Lambda Calculus, Event driven programming, Concurrency
29-36	Virtual Machines, Managed Languages, JIT, Case study

Procedural Abstraction

Can be overloaded (e.g., binary +) Can pass arguments into the scope Contains local variable declarations and statements

- ◆ Procedure is a **named parameterized scope**
 - Allows programmer to focus on a function interface, ignoring the details of how it is computed
- ◆ Value-returning functions
 - Example: $x = (b*b - \text{sqrt}(4*a*c))/2*a$
- ◆ Non-value returning functions
 - Called “procedures” (Ada), “subroutines” (Fortran), “void functions/methods” (C, C++, Java)
 - Have a visible side effect: change the state of some data value not defined in the function definition
 - Example: `strcpy(s1,s2)`

System Calls

- ◆ OS procedures often return status codes
 - Not the result of computing some function, but an indicator of whether the procedure succeeded or failed to cause a certain side effect

```
int open(const char* file, int mode)
{
    if (file == NULL) {
        return -1; // invalid file name

    if (open(file, mode) < 0)
        return -2; // system open failed
    ...
}
```

Arguments and Parameters

- ◆ **Argument:** expression that appears in a function call
- ◆ **Parameter:** identifier that appears in function declaration
- ◆ Parameter-argument matching by number and position
 - Exception: Perl. Instead of being declared in a function header, parameters are available as elements of special array @_

```
int h, i;  
void B(int w) {  
    int j, k;  
    i = 2*w;  
    w = w+1;  
}  
void A(int x, int y) {  
    bool i, j;  
    B(h);  
}  
int main() {  
    int a, b;  
    h = 5; a = 3; b =  
    2;  
    A(a, b);  
}
```

The diagram illustrates parameter-argument matching by number and position. Arrows point from the arguments in the function calls to the parameters in the function declarations. Specifically, an arrow points from 'h' in 'B(h)' to 'h' in 'void B(int w)'. Another arrow points from 'a' in 'A(a, b)' to 'x' in 'void A(int x, int y)'. A third arrow points from 'b' in 'A(a, b)' to 'y' in 'void A(int x, int y)'. A fourth arrow points from 'h' in 'h = 5' to 'h' in 'int h, i;'.

Parameter Passing Mechanisms

- ◆ By value
- ◆ By reference
- ◆ By value-result
- ◆ By result
- ◆ By name

Pass by Value

- ◆ Caller passes r-value of the argument to function
 - Compute the value of the argument at the time of the call and assign that value to the parameter
 - Reduces “aliasing”
 - Aliasing: two names refer to the same memory location
- ◆ Function cannot change value of caller’s variable

```
int increment (int age)  
    { age = age + 1; return age; }  
main () {  
    int myAge=45;  
    int newAge increment(myAge);  
}
```

Pass by Value

- ◆ Caller passes r-value of the argument to function
 - Compute the value of the argument at the time of the call and assign that value to the parameter
 - Reduces “aliasing”
 - Aliasing: two names refer to the same memory location
 - ◆ Function cannot change value of caller’s variable
 - ◆ All arguments in C and Java are passed by value
 - To allow caller’s variables to be modified, pointers can be passed as arguments
 - Example: `void swap(int *a, int *b) { ... }`
- Is there a contradiction here?

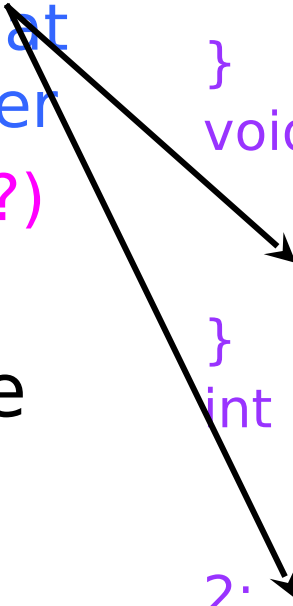
Pass by Reference

◆ Caller passes l-value of the argument to function

- Compute the address of the argument and assign that address to the parameter
- Increases aliasing (why?)

◆ Function can modify caller's variable via the address it received as argument

```
int h, i;
void B(int* w) {
    int j, k;
    i = 2*(*w);
    *w = *w+1;
}
void A(int* x, int* y) {
    bool i, j;
    B(&h);
}
int main() {
    int a, b;
    h = 5; a = 3; b =
2;
    A(&a, &b);
}
```



Pass by Reference in C++

- ◆ Special “reference type” indicates that l-value is passed as argument

- Recall that in C, only r-values can be arguments

```
void swap ((int& a, int& b)
           int temp = a;
           a = b;
           b = temp;
           }
```

l-values for C++ reference types are completely determined at compile-time

(why is this important?)

- ◆ & operator is overloaded in C++

- When applied to a variable, gives its l-value
- When applied to type name in parameter list, means pass the argument by reference

Two Ways To Pass By Reference

C or C++

```
void swap (int *a, int *b)
{
    int temp = *a;
    *a = *b;
    *b = temp;
}
```

```
int x=3, y=4;
swap(&x, &y);
```

C++ only

```
void swap (int& a, int&
b) {
    int temp = a;
    a = b;
    b = temp;
}
```

```
int x=3, y=4;
swap(x, y);
```

Which one is better? Why?

Pass by Value-Result

- ◆ Pass by value at the time of the call and/or copy the result back to the argument at the end of the call (copy-in-copy-out)
 - Example: “in out” parameters in Ada
- ◆ Reference and value-result are the same, except when aliasing occurs
 - Same variable is passed for two different parameters
 - Same variable is both passed and globally referenced from the called function

Pass by Name

- ◆ Textually substitute the argument for every instance of its corresponding parameter in the function body
 - Originated with Algol 60 but dropped by Algol's successors -- Pascal, Ada, Modula
- ◆ Example of late binding
 - Evaluation of the argument is delayed until its occurrence in the function body is actually executed
 - Associated with lazy evaluation in functional languages (e.g., Haskell)

An example

◆ Computes $\sum_{i=1}^{100} \frac{1}{i}$ in Algol 60

```
begin
  integer i;
  real procedure sum (i, lo, hi,
term);
    value lo, hi;
    integer i, lo, hi;
    real term;
    begin
      real temp;
      temp := 0;
      for i := lo step 1 until hi do
        temp := temp + term;
        sum := temp
      end;
      print (sum (i, 1, 100, 1/i))
    end
```

passed by name

becomes 1/i when sum is executed

Macro

◆ Textual substitution

```
#define swap(a,b) temp=a; a=b; b=temp;
```


```
...
```

```
int x=3, y=4;
```

```
int temp;
```

```
swap(x,y);
```

Textually expands to
temp=x; x=y; y=temp;



◆ Looks like a function definition, but ...

- Does not obey the lexical scope rules (i.e., visibility of variable declarations)
- No type information for arguments or result

Problems with Macro Expansion

```
#define swap(a,b) temp=a; a=b; b=temp;
```

□ □ □

if ($x < y$)

```
swap(x,y);
```

Textually expands to

```
if (x < y)
```

```
temp=x;
```

```
x=y;
```

```
y=temp;
```

Why not `#define swap(a,b) { int temp=a; a=b; b=temp; }`?

```
instead #define swap(a,b) do {  
    int temp=a; a=b; b=temp;  
} while(false);
```

Fixes type of swapped variables

Variable Arguments

- ◆ In C, can define a function with a variable number of arguments

- Example: `void printf(const char* format, ...)` Part of syntax!

- ◆ Examples of usage:

```
printf("hello, world");  
printf("length of '%s' = %d\n", str, str.length());  
printf("unable to open file descriptor %d\n", fd);
```

Format specification encoded by special %-encoded characters

- %d,%i,%o,%u,%x,%X – integer argument
- %s – string argument
- %p – pointer argument (void *)
- Several others (see C Reference Manual!)

Implementation of Variable Args

◆ Special functions `va_start`, `va_arg`, `va_end` compute arguments at run-time

```
void printf(const char* format, ...)
{
    int i; char c; char* s; double d;
    va_list ap; /* declare an "argument pointer" to a variable arg list */
    va_start(ap, format); /* initialize arg pointer using last known arg */

    for (char* p = format; *p != '\\0'; p++) {
        if (*p == '%') {
            switch (*++p) {
                case 'd':
                    i = va_arg(ap, int); break;
                case 's':
                    s = va_arg(ap, char*); break;
                case 'c':
                    c = va_arg(ap, char); break;
            }
            ... /* etc. for each % specification */
        }
    }
    ...

    va_end(ap); /* restore any special stack manipulations */
}
```

Implementation of Variable Args

- ◆ Special functions `va_start`, `va_arg`, `va_end` compute arguments at run-time

```
#include <stdarg.h> #include <stdio.h>
double average ( int num, ... )
{
    va_list arguments;
    double sum = 0;
    va_start ( arguments, num );
    for ( int x = 0; x < num; x++ )
        { sum += va_arg ( arguments, double ); }
    va_end ( arguments );
    return sum / num;
}
int main() {
    printf( "%f\n", average ( 3, 12.2, 22.3, 4.5 ) );

    printf( "%f\n", average ( 5, 3.3, 2.2, 1.1, 5.5, 3.3 ) );
}
```