Šifra predmeta: R265 7.03.2023.

## Uvod u interaktivno dokazivanje teorema Vežbe 3

Zadatak 1 Intuicionistička pravila prirodne dedukcije u iskaznoj logici

Diskutovati o pravilima uvođenja i pravilima eliminacije prirodne dedukcije iskazne logike. Pomoću ključne reči *thm* ispitati svako pravilo prirodne dedukcije. Primeniti odgovarajuće pravilo prirodne dedukcije na jednostavnim formulama i diskutovati o cilju koga treba dokazati pre i posle primene tog pravila.

```
Uvodjenje konjukcije: conjI
thm conjI
lemma A \wedge B
 apply (rule conjI)
Uvodjenje disjunkcije: disjI1/disjI2
thm disjI1
thm disjI2
lemma A \vee B
 apply (rule disjI2)
Uvodjenje implikacije: impI
thm impI
lemma A \longrightarrow B
 apply (rule \ impI)
Uvodjenje ekvivalencije: iffI
thm iffI
lemma A \longleftrightarrow B
 apply (rule iffI)
Uvodjenje negacije: notI
thm notI
lemma \neg A
 apply (rule notI)
```

Eliminacija konjukcije. conjE

```
thm conjE

lemma A \wedge B \Longrightarrow C
apply (erule\ conjE)

Eliminacija disjunkcije. disjE

thm disjE

lemma A \vee B \Longrightarrow C
apply (erule\ disjE)

Eliminacija implikacije. impE

thm impE

lemma A \longrightarrow B \Longrightarrow C
apply (erule\ impE)

Eliminacija ekvivalencije. iffE

thm iffE

lemma A \longleftrightarrow B \Longrightarrow C
apply (erule\ iffE)
```

Eliminacija negacije. notE

 $\mathbf{lemma} \neg A \Longrightarrow B$   $\mathbf{apply} \ (erule \ not E)$ 

thm notE

## Zadatak 2 Dokazi u prirodnoj dedukciji

Pokazati da su sledeće formule tautologija u iskaznoj logici. Dozvoljeno je korišćenje samo intuicionističkih pravila prirodne dedukcije.

```
lemma A \wedge B \longrightarrow B \wedge A

apply (rule \ impI)

apply (erule \ conjE)

apply (rule \ conjI)

apply assumption +

done

lemma A \vee B \longrightarrow B \vee A

apply (rule \ disjE)

apply (erule \ disjI2)

apply (erule \ disjI2)

apply (erule \ disjI1)
```

```
apply assumption
  done
lemma A \wedge B \longrightarrow A \vee B
 apply (rule impI)
 apply (erule conjE)
 apply (rule disjI1)
 apply assumption
 done
lemma (A \wedge B \longrightarrow C) \longrightarrow (A \longrightarrow (B \longrightarrow C))
  apply (rule \ impI) +
 apply (erule impE)
  apply (rule conjI)
   apply assumption +
  done
lemma (A \longrightarrow (B \longrightarrow C)) \longrightarrow (A \land B \longrightarrow C)
 apply (rule impI) +
 apply (erule impE)
  apply (erule\ conjE)
  apply assumption
  apply (erule\ conjE)
 apply (erule impE)
  apply assumption +
  done
lemma \neg (A \lor B) \longrightarrow \neg A \land \neg B
  apply (rule impI)
 apply (rule conjI)
  apply (rule notI)
  apply (erule \ not E)
  apply (rule disjI1)
  apply assumption
 apply (rule notI)
 apply (erule notE)
 apply (rule disjI2)
 apply assumption
  done
lemma \neg A \land \neg B \longrightarrow \neg (A \lor B)
 apply (rule impI)
 apply (rule notI)
 apply (erule\ conjE)
 apply (erule \ disjE)
  apply (erule \ not E)
  apply assumption
 apply (erule notE) +
 apply assumption
  done
```

```
lemma \neg (A \longleftrightarrow \neg A)
 apply (rule notI)
 apply (erule iffE)
 apply (erule impE) back
  apply (rule notI)
  apply (erule impE)
   apply assumption
  apply (erule \ not E)
  apply assumption
 apply (erule impE)
  apply assumption
 apply (erule\ notE)
 apply assumption
 done
Dodatni primeri
lemma (Q \longrightarrow R) \land (R \longrightarrow P \land Q) \land (P \longrightarrow Q \lor R) \longrightarrow (P \longleftrightarrow Q)
 apply (rule \ impI)
 apply (rule iffI)
  apply (erule \ conjE) +
  apply (erule impE) back back
   apply assumption
  apply (erule \ disjE)
   apply assumption
  apply (erule impE) back
   apply assumption
  apply (erule conjE)
  apply assumption
 apply (erule \ conjE) +
 apply (erule \ impE)
  apply assumption
 apply (erule impE)
  {\bf apply} \ assumption
 apply (erule\ conjE)
 apply assumption
 done
lemma (P \longrightarrow Q) \land (Q \longrightarrow R) \longrightarrow (P \longrightarrow Q \land R)
 apply (rule impI) +
 apply (erule\ conjE)
 apply (erule impE)
  apply assumption
 apply (rule conjI)
  apply assumption
 apply (erule impE)
  apply assumption +
 done
lemma (P \longrightarrow Q) \land \neg Q \longrightarrow \neg P
 apply (rule impI)
 apply (rule notI)
```

```
apply (erule conjE)
  apply (erule impE)
  apply assumption
  apply (erule \ not E)
  apply assumption
  done
\mathbf{lemma}\ (P \longrightarrow (Q \longrightarrow R)) \longrightarrow (Q \longrightarrow (P \longrightarrow R))
  apply (rule impI) +
  apply (erule impE)
  {\bf apply} \ assumption
  apply (erule impE)
  apply assumption +
  done
lemma \neg (P \land \neg P)
  apply (rule notI)
  apply (erule conjE)
  apply (erule \ not E)
  apply assumption
  done
lemma A \wedge (B \vee C) \longrightarrow (A \wedge B) \vee (A \wedge C)
  apply (rule impI)
  apply (erule\ conjE)
  apply (erule disjE)
  apply (rule disjI1)
  apply (rule conjI)
   apply assumption +
  {f apply} \ ({\it rule \ disjI2})
  apply (rule\ conjI)
  apply assumption +
  done
lemma \neg (A \land B) \longrightarrow (A \longrightarrow \neg B)
  apply (rule impI) +
  apply (rule\ not I)
  apply (erule \ not E)
  apply (rule conjI)
  apply assumption +
  done
lemma (A \longrightarrow C) \land (B \longrightarrow \neg C) \longrightarrow \neg (A \land B)
  apply (rule impI)
  apply (rule notI)
  apply (erule conjE) +
  apply (erule impE)
  apply assumption
  apply (erule impE)
  apply assumption
```

```
apply (erule notE)
 apply assumption
  done
lemma (A \wedge B) \longrightarrow ((A \longrightarrow C) \longrightarrow \neg (B \longrightarrow \neg C))
 apply (rule \ impI) +
 apply (rule notI)
 apply (erule conjE)
 apply (erule impE)
  apply assumption
  apply (erule impE)
  apply assumption
  apply (erule \ not E)
  apply assumption
  done
lemma (A \longleftrightarrow B) \longrightarrow (\neg A \longleftrightarrow \neg B)
 apply (rule \ impI)
 apply (erule iffE)
 apply (rule iffI)
  apply (rule notI)
  apply (erule impE) back
   apply assumption
  apply (erule notE)
  apply assumption
  apply (rule notI)
  apply (erule impE)
  apply assumption
 apply (erule \ not E)
 apply assumption
  done
lemma A \longrightarrow \neg \neg A
 apply (rule impI)
 apply (rule notI)
 apply (erule \ not E)
 apply assumption
 done
lemma \neg (A \longleftrightarrow \neg A)
  apply (rule notI)
 apply (erule iffE)
 apply (erule impE) back
  apply (rule notI)
  apply (erule \ impE)
   apply assumption
  apply (erule notE)
  apply assumption
  apply (erule impE)
  apply assumption
 apply (erule \ not E)
```

```
{\bf apply} \ assumption
  done
lemma (A \longrightarrow B) \longrightarrow (\neg B \longrightarrow \neg A)
  apply (rule impI) +
  apply (rule notI)
  apply (erule impE)
  apply assumption
  apply (erule notE)
  {\bf apply} \ assumption
  done
\mathbf{lemma} \neg A \lor B \longrightarrow (A \longrightarrow B)
  \mathbf{apply} \ (\mathit{rule} \ \mathit{impI}) \ +
  apply (erule disjE)
  apply (erule notE)
  apply assumption +
  done
```