



Northeastern  
University

# Lecture 18: Stack Implementations - 1

Prof. Chen-Hsiang (Jones) Yu, Ph.D.  
College of Engineering

Materials are edited by Prof. Jones Yu from

Data Structures and Abstractions with Java, 5<sup>th</sup> edition. By Frank M. Carrano and Timothy M. Henry.  
ISBN-13 978-0-13-483169-5 © 2019 Pearson Education, Inc.

# Outline

---

- A Linked Implementation
- An Array-Based Implementation
- A Vector-Based Implementation

# Outline

---

- A Linked Implementation
- An Array-Based Implementation
- A Vector-Based Implementation

# A Linked Implementation

# Stack Implementation

---

- Each operation involves the top of stack
  - » push
  - » pop
  - » peek
- The head of linked list is easiest, fastest to access
  - » Let this be the top of the stack

# Linked Implementation

---

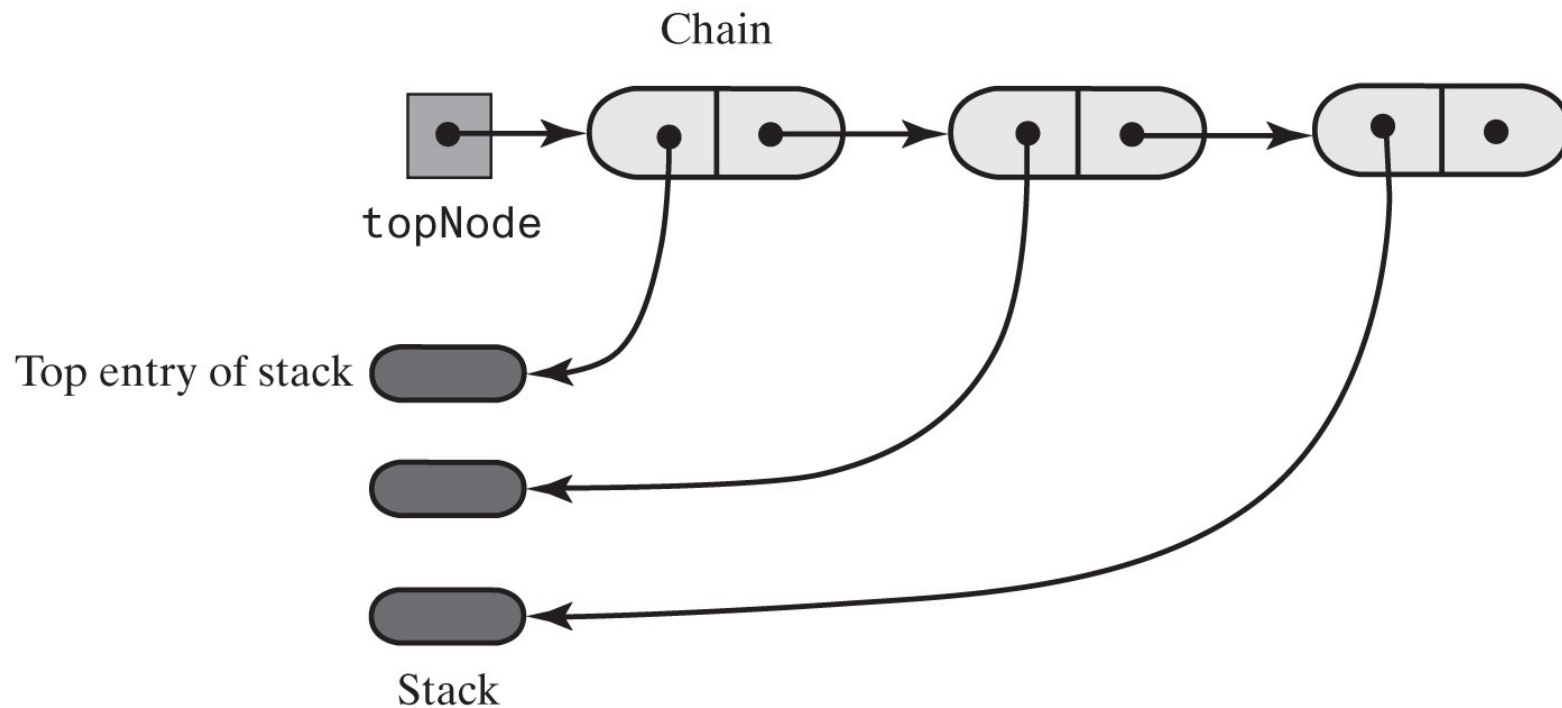


Figure 6-1: A chain of linked nodes that implements a stack

# Linked Implementation

---

```
/** A class of stacks whose entries are stored in a chain of nodes. */
public final class LinkedStack<T> implements StackInterface<T>
{
    private Node topNode; // References the first node in the chain

    public LinkedStack()
    {
        topNode = null;
    } // end default constructor

    // < Implementations of the stack operations go here. >
    // . . .

    private class Node
    {
        private T    data; // Entry in stack
        private Node next; // Link to next node
        // < Implementations of the node operations go here. >
    } // end Node
} // end LinkedStack
```

Listing 6-1: An outline of a linked implementation of the ADT stack

# Linked Implementation - Push Operation

---

(a) A new node that references the node at the top of the stack

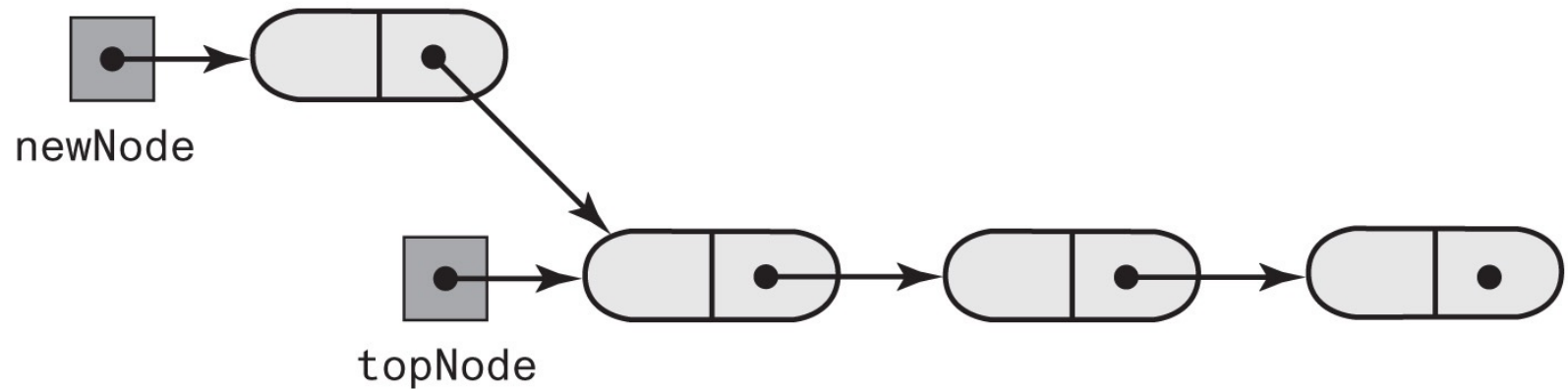


Figure 6-2: (a) A new node that references the node at the top of the stack;



# Linked Implementation - Push Operation (cont.)

---

(b) The new node is now at the top of the stack

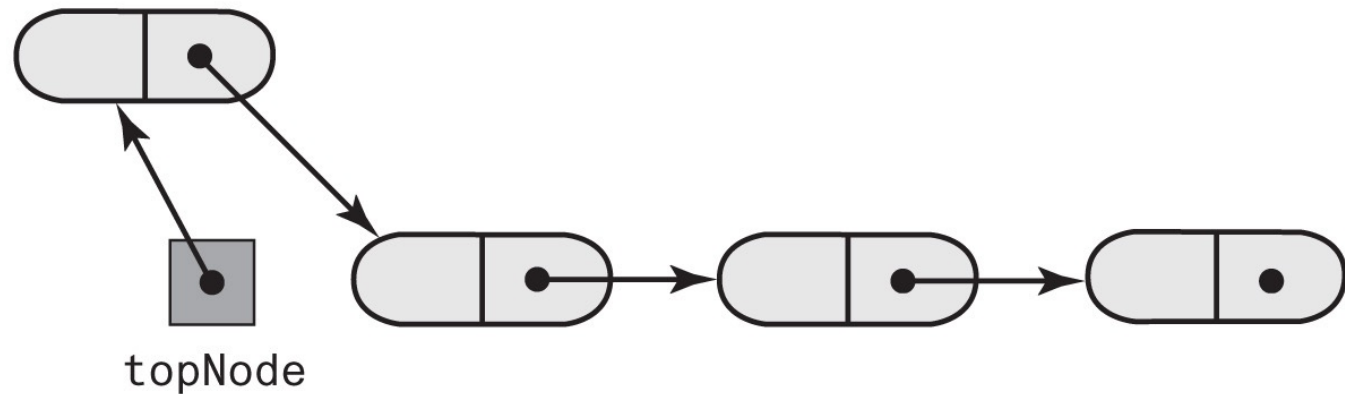
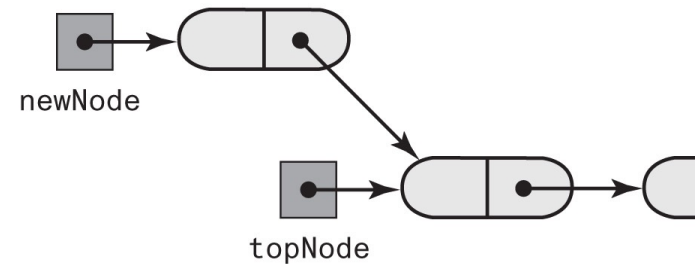


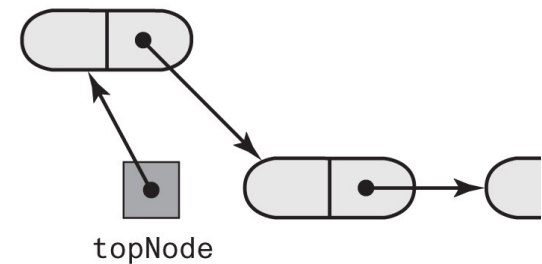
Figure 6-2: (b) the new node is now at the top of the stack

# Linked Implementation - Push Operation (cont.)



```
public void push(T newEntry)
{
    Node newNode = new Node(newEntry, topNode);
    topNode = newNode;
} // end push
```

Definition of **push**



# Linked Implementation - Pop Operation

---

(a) Before pop

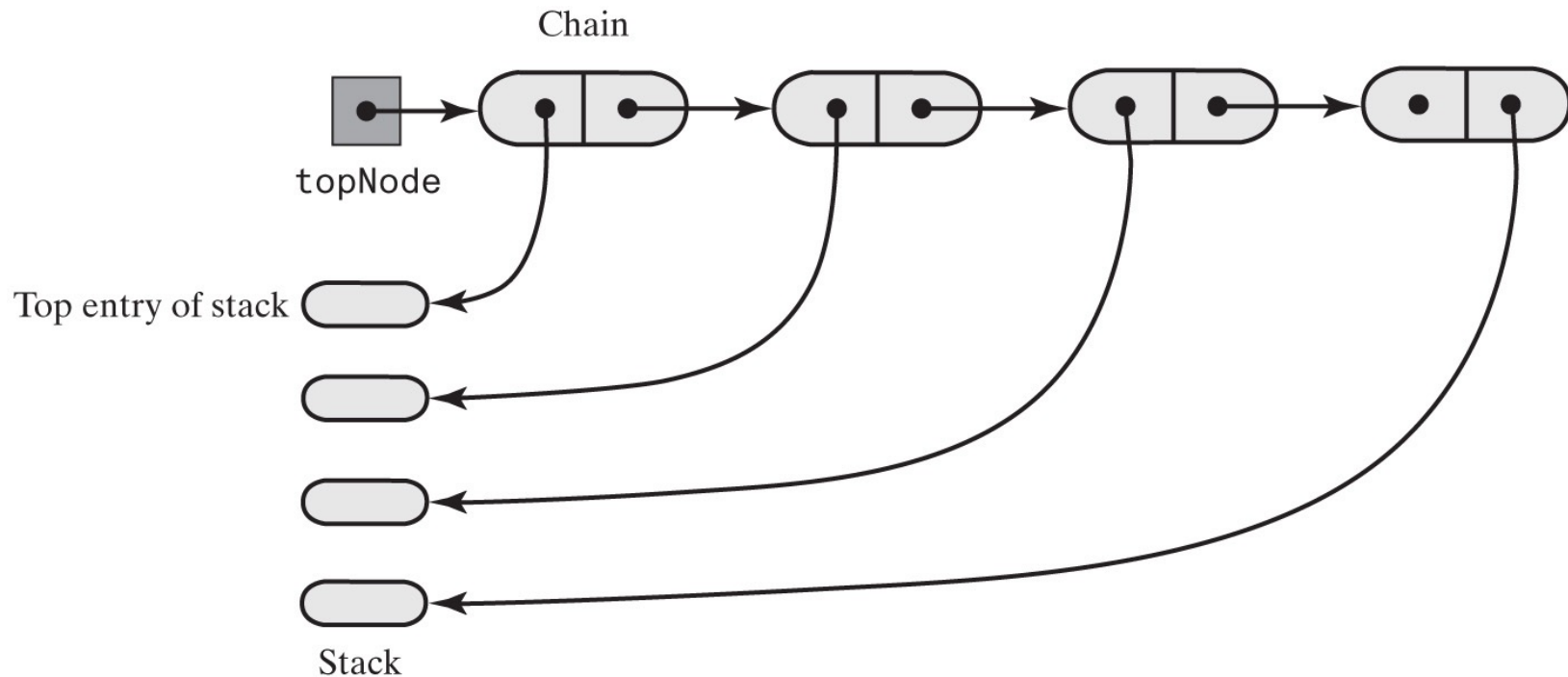


Figure 6-3: The stack (a) before the first node in the chain is deleted

# Linked Implementation - Pop Operation (cont.)

(b) After pop

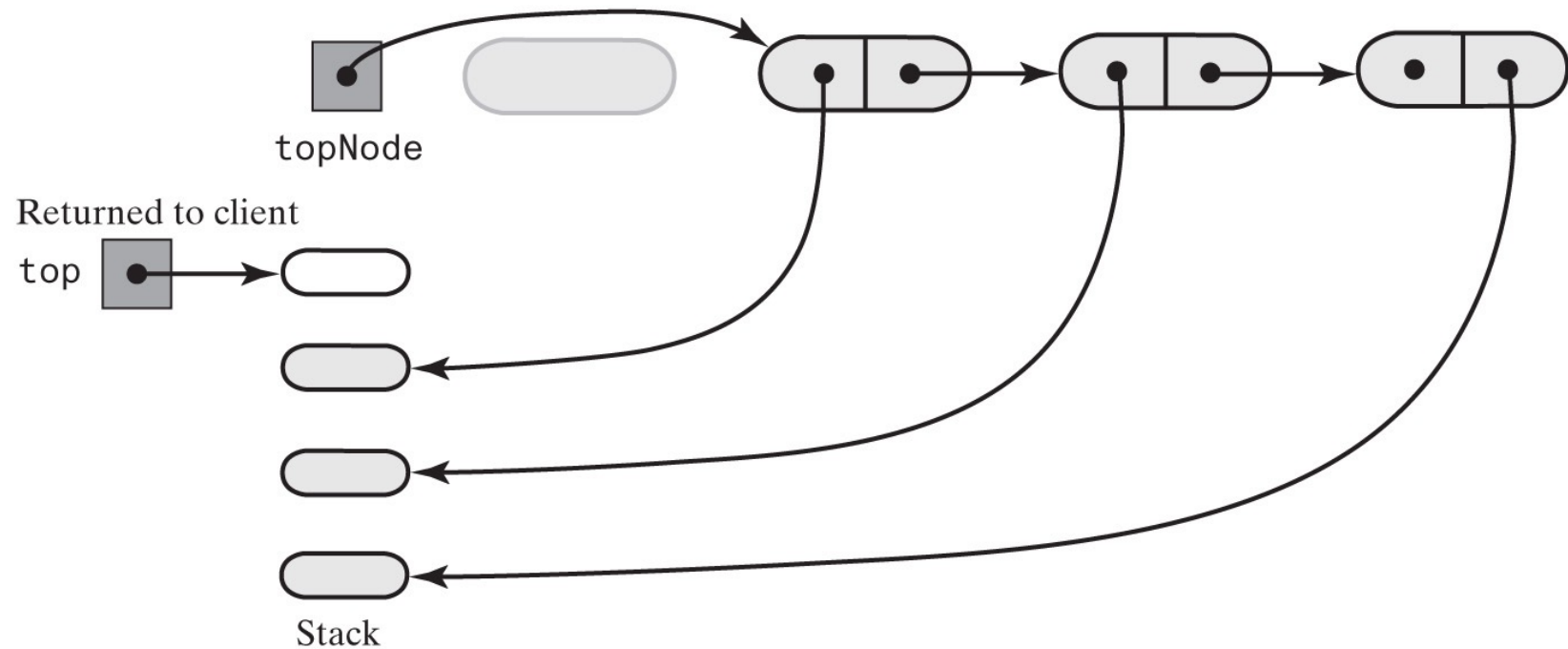


Figure 6-3: The stack (b) after the first node in the chain is deleted

# Linked Implementation - Peek & Pop Operations

---

```
public T peek()
{
    if (isEmpty())
        throw new EmptyStackException();
    else
        return topNode.getData();
} // end peek

public T pop()
{
    T top = peek(); // Might throw EmptyStackException

    // Assertion: topNode != null
    topNode = topNode.getNextNode();

    return top;
} // end pop
```

Definition of **peek** and **pop**

# Linked Implementation - isEmpty() & clear()

---

```
public boolean isEmpty()  
{  
    return topNode == null;  
} // end isEmpty
```

```
public void clear()  
{  
    topNode = null;  
} // end clear
```

Definition of **isEmpty** and **clear**

## Exercise

---

- Download “[L18\\_E1.zip](#)” from the Canvas
- Revise the implementation of `pop()` so that it does not call `peek()`

```
public T pop()
{
    T top = peek(); // Might throw EmptyStackException
    assert topNode != null;
    topNode = topNode.getNextNode();
    return top;
} // end pop
```

## Note: Assertions in Eclipse

---

- How to enable assertions in Eclipse?
  - » Open the Run Dialog ([Run > Run Configurations](#))
  - » Click on the tab, "(x)= Arguments."
  - » Under the field for "VM arguments," type [-ea](#) to enable assertions.
  - » Click on the "Apply" and "Run" button



# Answer

---

```
public T pop() {  
  
    if (topNode != null) {  
        T top = topNode.getData();  
        topNode = topNode.getNextNode();  
        numberOfEntries--;  
        return top;  
    } else {  
        throw new EmptyStackException();  
    }  
  
}
```

## Exercise

---

- Is an implementation of the ADT stack reasonable if the top of the stack is at the end of a chain of linked nodes instead of its beginning? Explain.

# Answer

---

- No
- Although adding a reference to the chain's last node enable you to access the stack's top entry or push a new entry onto the stack efficiently, it is not enough for pop the stack.
- You also need a reference to next-to-last node to remove the chain's last node.
- Therefore, placing the stack's top entry at the end of the chain is not as efficient or easy to implement as placing it at the beginning.